

Discovering Maximal Motif Cliques in Large Heterogeneous Information Networks

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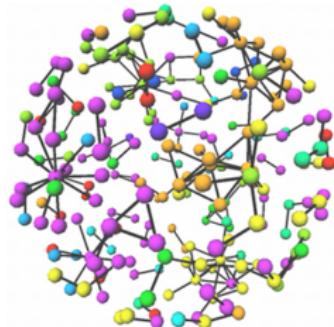
Graphs are Everywhere



Social Network

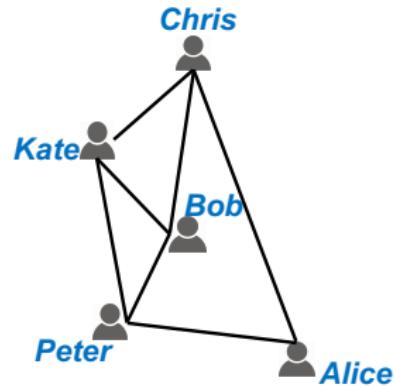


E-Commerce



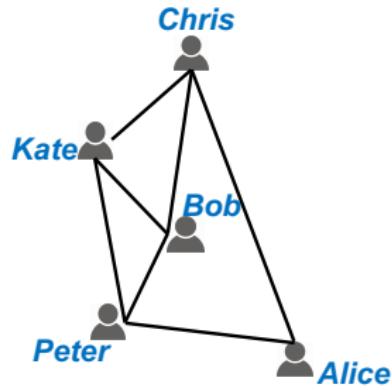
Protein-Protein
interaction Network

Homogeneous VS Heterogeneous Graphs

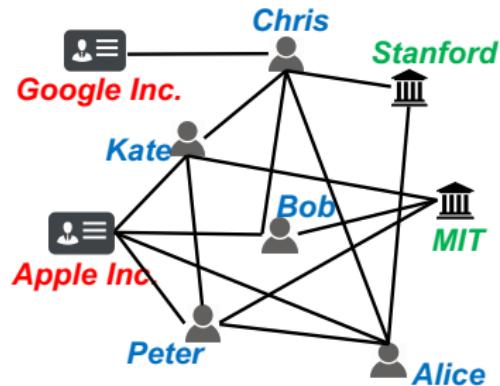


(a) Homogeneous

Homogeneous VS Heterogeneous Graphs

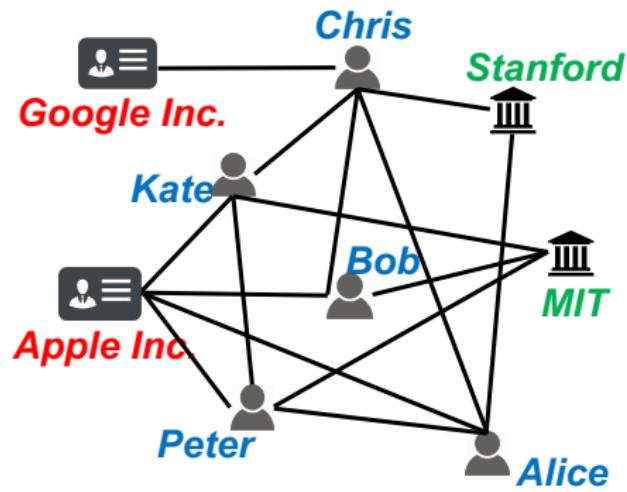


(a) Homogeneous



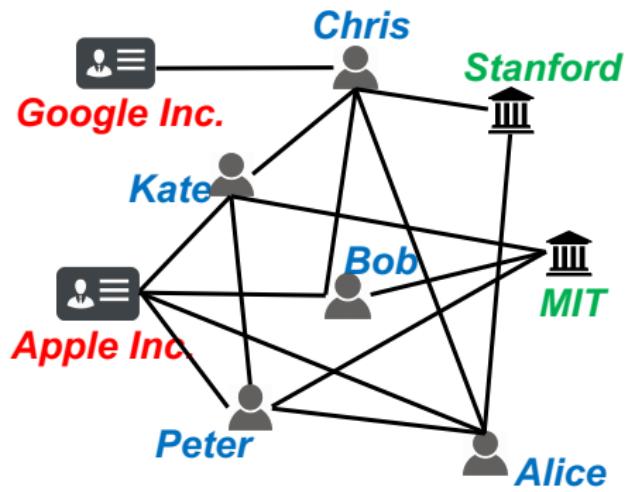
(b) Heterogeneous

Heterogeneous Information Networks (HINs)



(a) An HIN

Heterogeneous Information Networks (HINs)



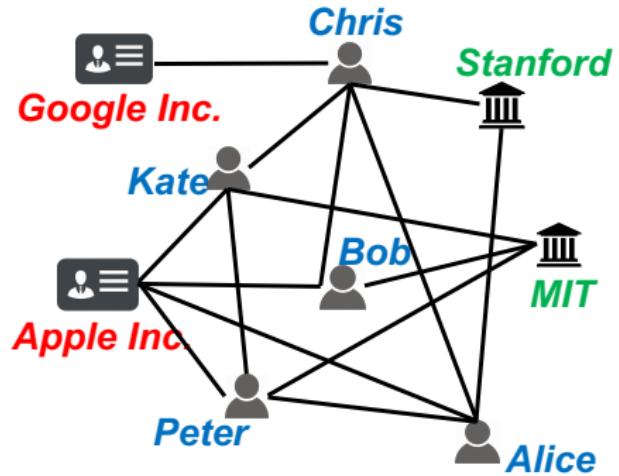
(a) An HIN



(b) Schema

How to Find Cliques for HINs?

- Clique: complete subgraph [E. Akkoyunlu, 1973]



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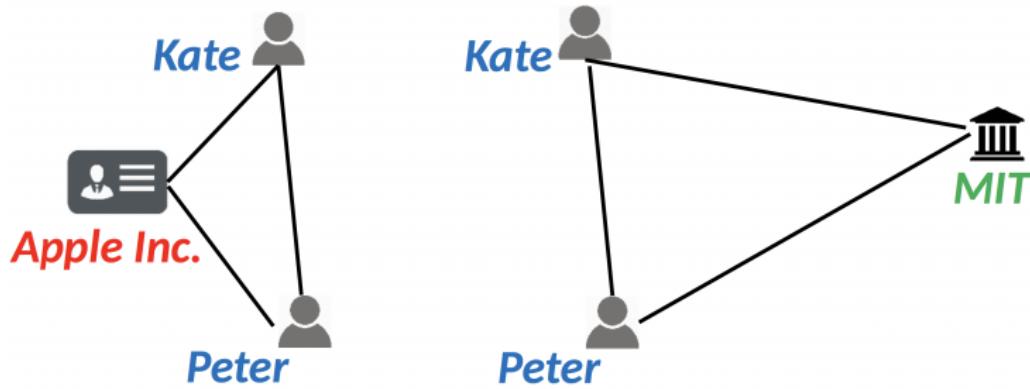
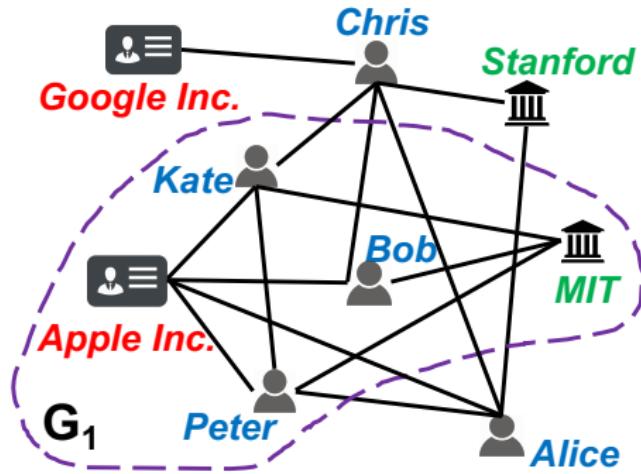


Figure: Traditional cliques

How to Find Cliques for HINs?

- Clique: complete subgraph [E. Akkoyunlu, 1973]

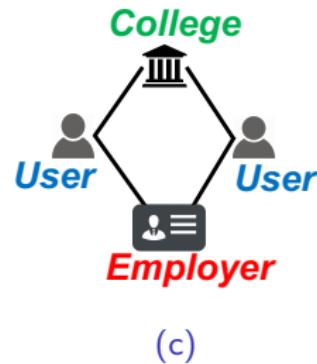


Motif

- Small pattern (higher-order structure)
- Building blocks of large and complex networks [Science'16]

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- Existing works:
 - ▶ motif discovery [SIGMOD'15, DMKD'18]
 - ▶ graph node clustering [Science'16, KDD'17]
 - ▶ motif frequency estimation [WWW'15, WSDM'17, TKDD'17]

Motif clique (or m-clique)

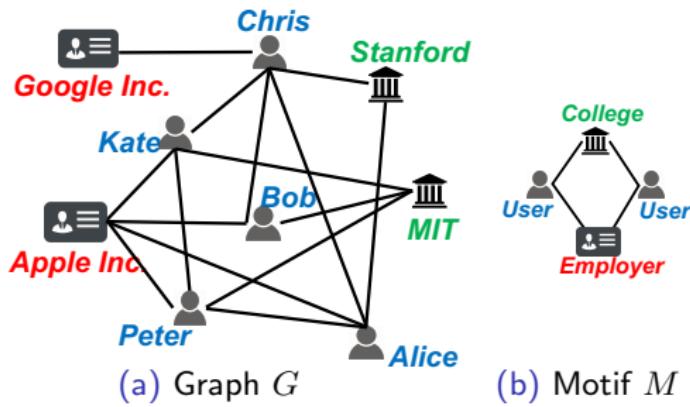
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M : a small connected HIN which **follows the schema** of G

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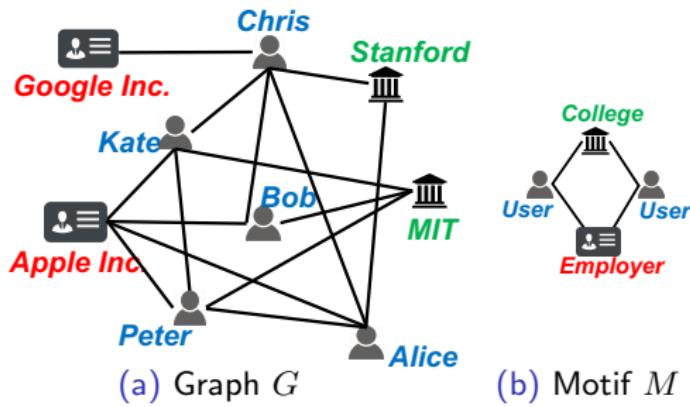
Label-matched sets:

- {MIT, Kate, Peter, Apple}
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- ...

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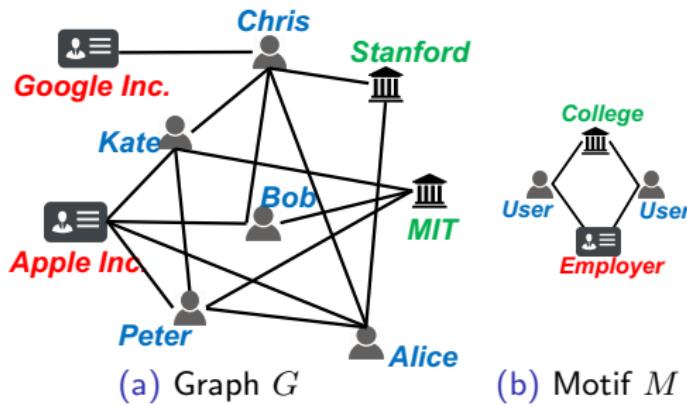
An induced subgraph G' of G is an **m-clique** of M :

- ① G' has the **same set of labels** with M ;
- ② \forall label-matched set H in G' , M is **subgraph isomorphic** to $G'[H]$.

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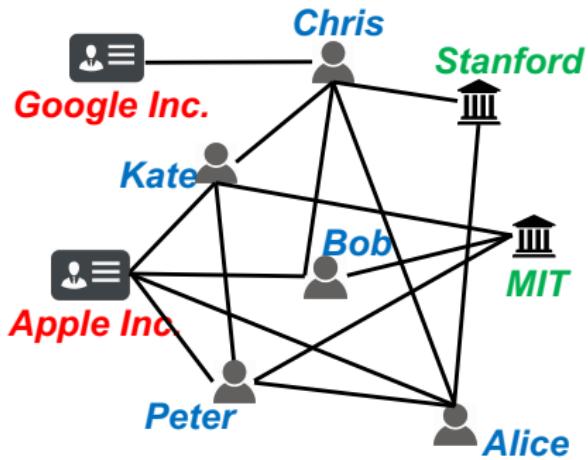
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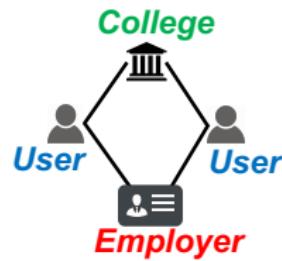
Maximal m-clique: m-clique & not contained in any other m-clique.

Maximal M-Clique **Enumeration (MMCE)**: extract all maximal m-cliques in G .

Example

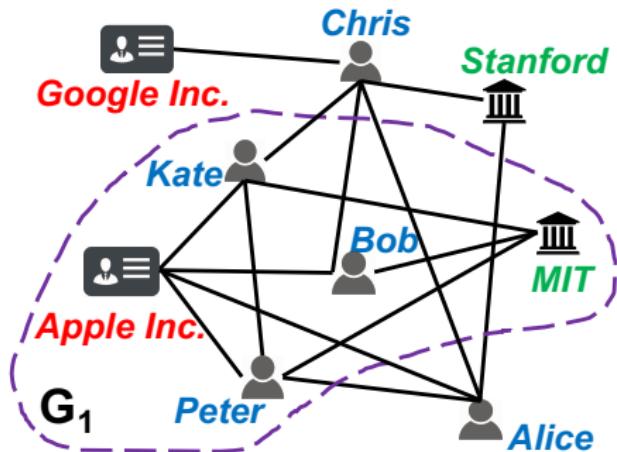


(a) Social network G

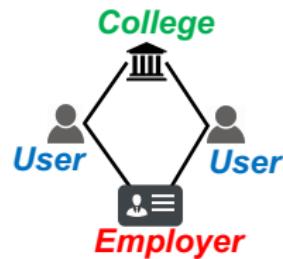


(b) Motif M

Example



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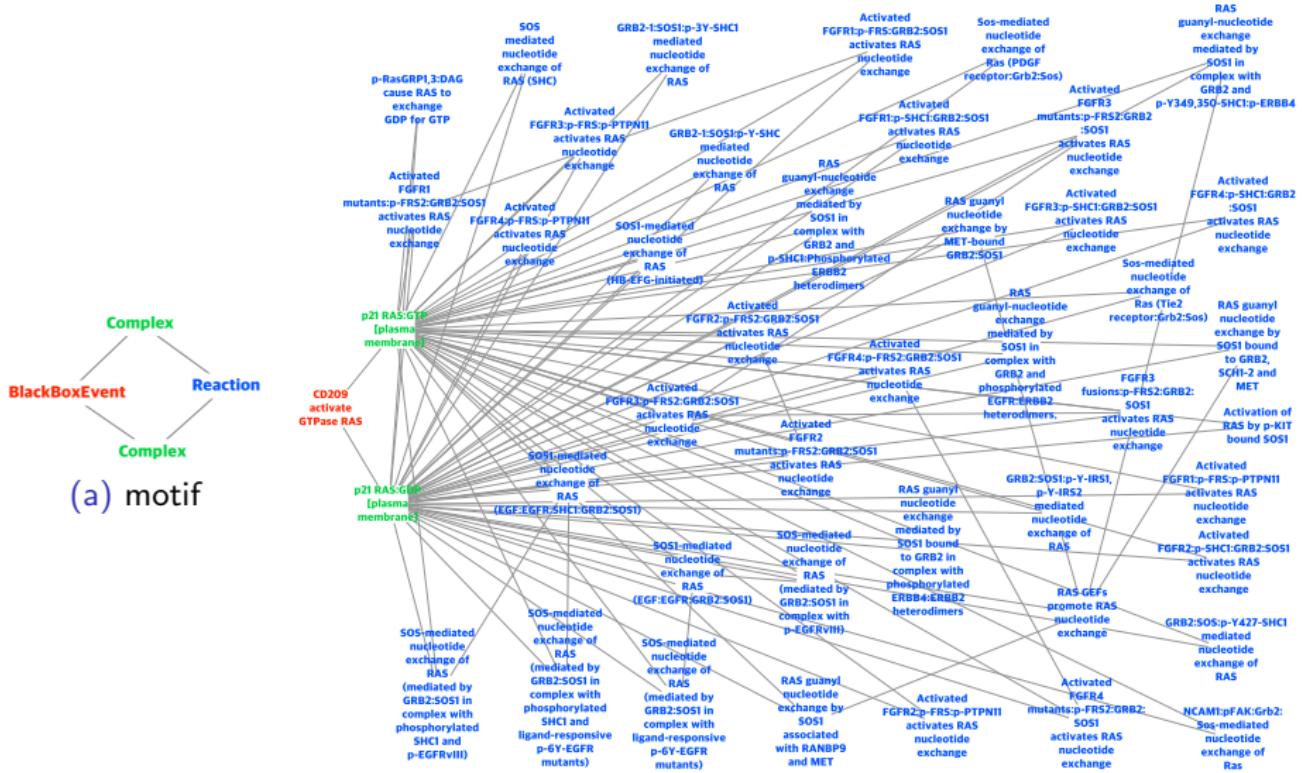


(b) Motif M

Apply m-clique in Biological Analysis

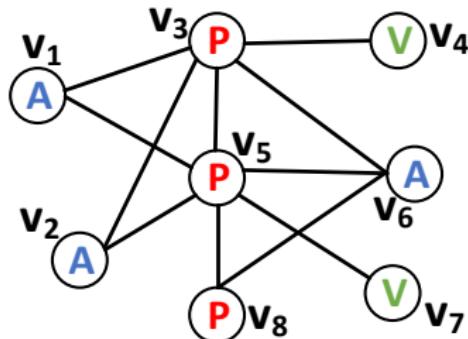
- Dataset: **Reactome**, a well known bioinformatic HIN [A. Fabregat et al., 2017]
- **Complex**: physical entities, e.g., proteins.
- **Reaction**: biochemical reactions which have balanced input and output entities.
- **BlackBoxEvent**: reactions or complex processes where details are not yet established.

Apply m-clique in Biological Analysis

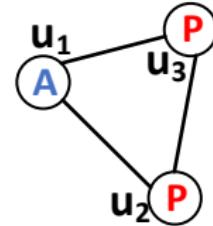


(b) maximal m-clique

Baseline Algorithm

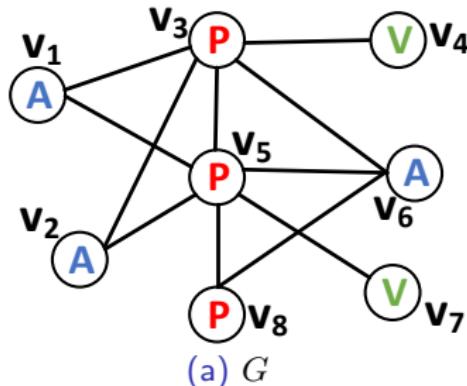


(a) G

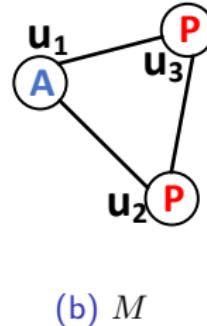


(b) M

Baseline Algorithm



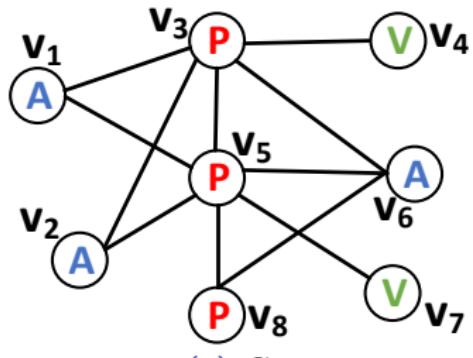
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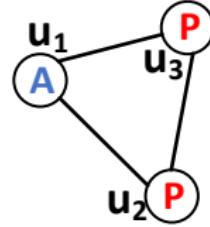
(b) M

- Find all matched subgraphs
 - ▶ $S_1 = \{v_1, v_3, v_5\}$; $S_2 = \{v_2, v_3, v_5\}$; $S_3 = \{v_6, v_3, v_5\}$; $S_4 = \{v_6, v_5, v_8\}$
- For each matching, find all maximal m-cliques containing it.
 - ▶ $S_1 \rightarrow \{v_1, v_2, v_3, v_5, v_6\}$
 - ▶ $S_2 \rightarrow \{v_1, v_2, v_3, v_5, v_6\}$
 - ▶ $S_3 \rightarrow \{v_1, v_2, v_3, v_5, v_6\}$
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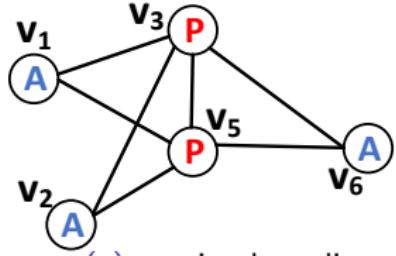
Baseline Algorithm



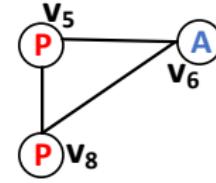
(a) G



(b) M

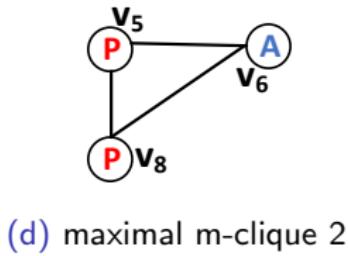
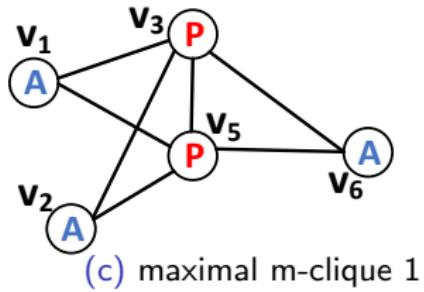
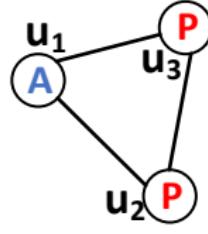
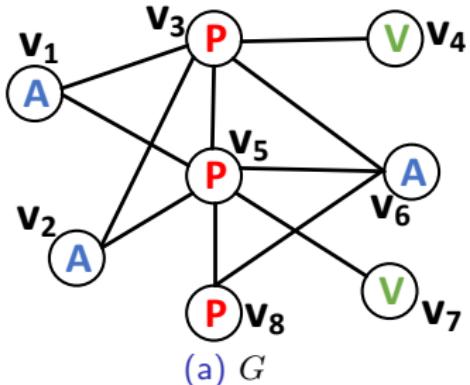


(c) maximal m-clique 1



(d) maximal m-clique 2

Baseline Algorithm



Time cost on Reactome: > 1000 seconds/query

Challenge 1: Node Expansion is NP-hard

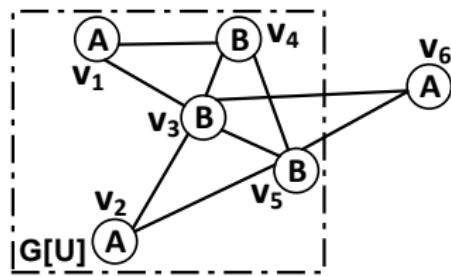


Figure: G

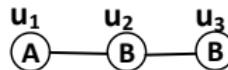


Figure: Motif

Possible label matched sets:

- $\{v_6, v_3, v_4\}$
- $\{v_6, v_3, v_5\}$
- $\{v_6, v_4, v_5\}$

Challenge 1: Node Expansion is NP-hard

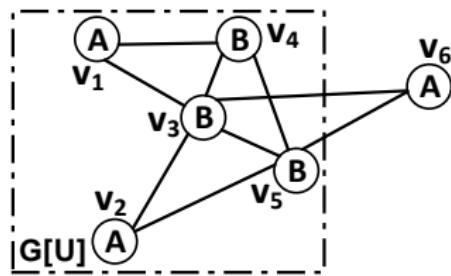


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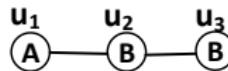
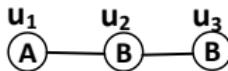
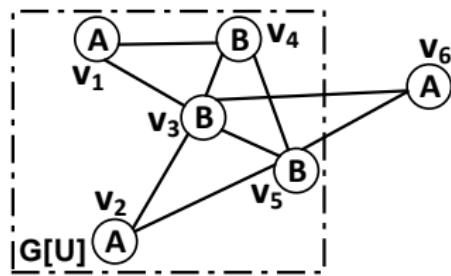


Figure: Motif

Dominance relationship:

- $v = v_6$; $u = v_2$
- $L(v_6) = L(v_2) = \text{"A"}$
- v_2 is dominated by v_6

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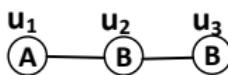
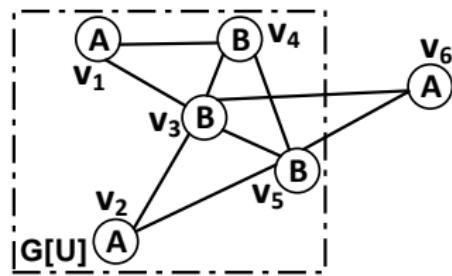
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Pruning strategies:

- Advanced node expansion
- Early stop pruning

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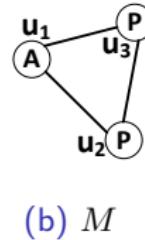
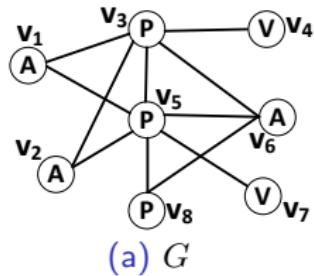
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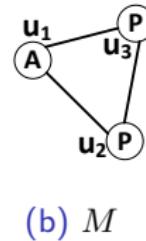
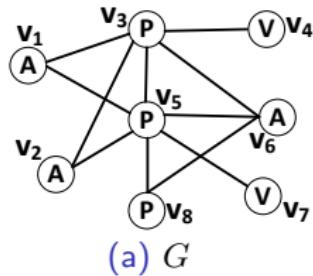
Time cost on Reactome: around **10 seconds/query**

Challenge 2: Duplication Avoidance



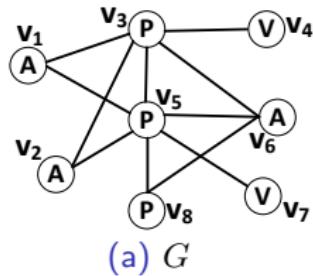
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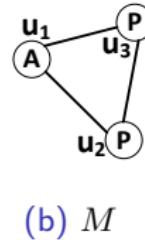


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- Pruning strategy: **set-trie tree**
 - ▶ Dynamically build the set-trie tree and check candidates.

Challenge 2: Duplication Avoidance



(a) G



(b) M

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- Pruning strategy: **set-trie tree**
 - ▶ Dynamically build the set-trie tree and check candidates.

Time cost on Reactome: around **0.1 seconds/query**

META: Maximal m-clique EnumeraTion Algorithm

Basic framework + the following pruning strategies:

- Dominance relationship between nodes
 - ▶ Advanced node expansion
 - ▶ Early stop pruning
- Duplication avoidance (set-trie tree)

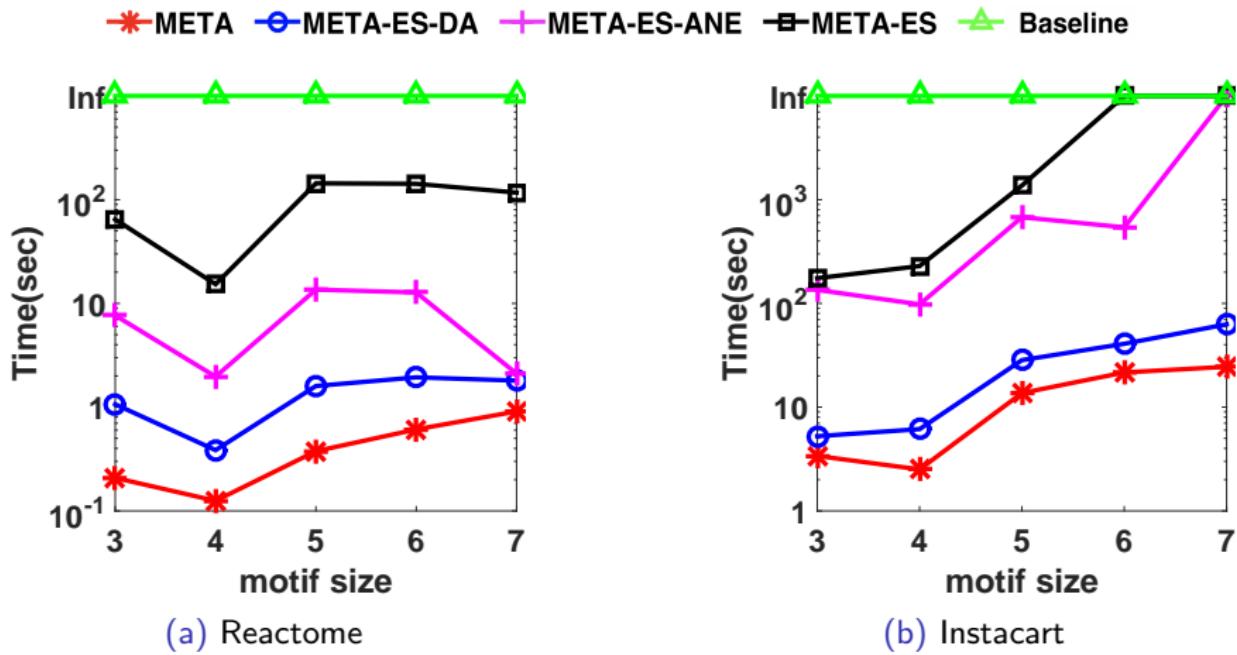
Experiments

Real graphs:

Dataset	#Nodes	#Edges	#Labels	Avg. Degree
DBLP	15K	51K	4	6.6
Amazon	548K	1.78M	4	6.5
Reactome	54K	98K	15	3.6
Yeast	3K	13K	71	8.1
Instacart	5K	13K	21	4.9

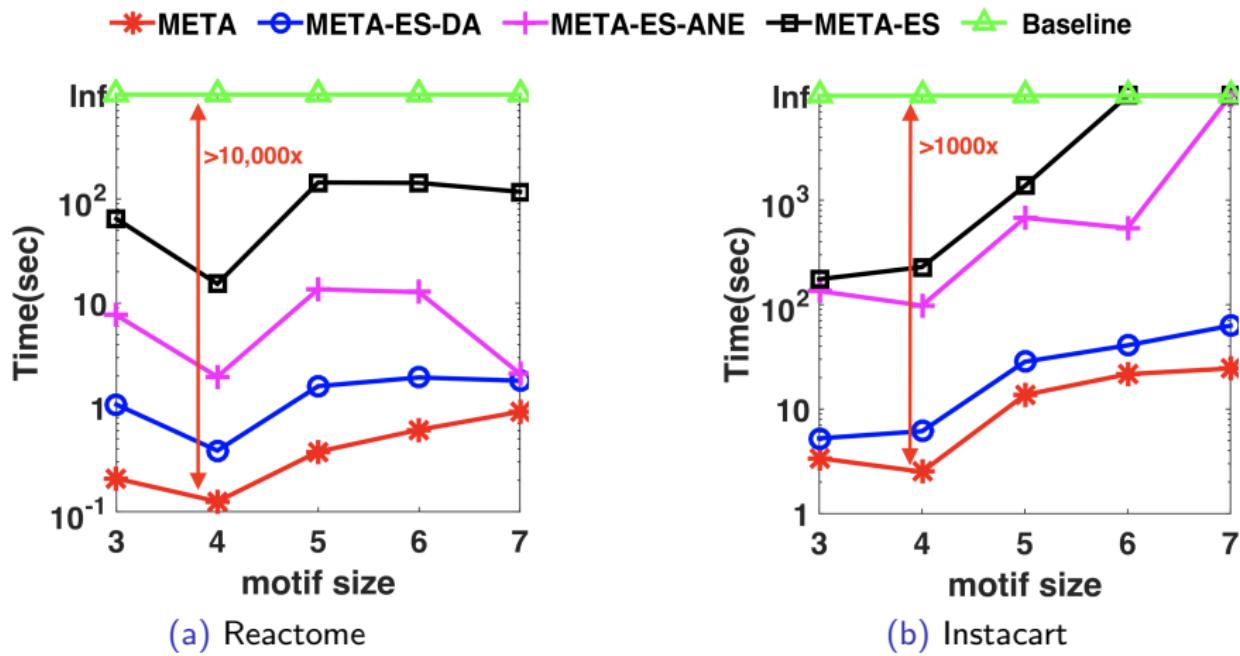
Time Cost by Varying Motif Size

The time of finding 10^3 maximal m-cliques.



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Conclusions

- m-clique for HINs
- The META algorithm
- Evaluation on both real and synthetic datasets

Future Work

- Extend META to handle more rich information on HINs
 - ▶ nodes with multiple labels
 - ▶ edges with directions and labels
 - ▶ ...
- Study other fundamental graph problems based on motif
 - ▶ motif-based path
 - ▶ motif-based connected component
 - ▶ ...

Thanks!

Q&A

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