

početní funkce

1) nulová funkce $\xi : \mathbb{N}^0 \rightarrow \mathbb{N}$
 $\xi : () \mapsto 0$

2) následující $\sigma : \mathbb{N} \rightarrow \mathbb{N}$
 $\sigma(x) = x + 1$

3) projekce $\pi_k^n : \mathbb{N}_*^n \rightarrow \mathbb{N}$
 $\pi_k^n (x_1, \dots, x_{k-1}, x_k, x_{k+1}, \dots, x_n) = x_k$

primitivní rekursivní funkce

4) kombinace $F : \mathbb{N}^2 \rightarrow \mathbb{N}^m$ $g : \mathbb{N}^2 \rightarrow \mathbb{N}^n$

$F \times g : \mathbb{N}^2 \rightarrow \mathbb{N}^{m+n}$

$F \times g(\bar{x}) = (F(\bar{x}), g(\bar{x})) \quad \bar{x} \in \mathbb{N}^2$

5) composition $F: \mathbb{N}^2 \rightarrow \mathbb{N}^m$ $g: \mathbb{N}^m \rightarrow \mathbb{N}^n$

$$g \circ F: \mathbb{N}^2 \rightarrow \mathbb{N}^n$$

$$g \circ F(\bar{x}) = g(F(\bar{x})) \quad \bar{x} \in \mathbb{N}^2$$

6) primitive recursion

$$F: \mathbb{N}^{2+1} \rightarrow \mathbb{N}^m$$

$$g: \mathbb{N}^2 \rightarrow \mathbb{N}^m \quad h: \mathbb{N}^{2+m+1} \rightarrow \mathbb{N}^m$$

$$F(\bar{x}, 0) = g(\bar{x})$$

$$F(\bar{x}, y+1) = h(\bar{x}, y, F(\bar{x}, y)) \quad \bar{x} \in \mathbb{N}^2$$

$$\begin{array}{l} \cancel{F(x, 0) = x} \\ F(x, 0) = \pi_1^1(x) \end{array} \quad \left| \quad \begin{array}{l} \cancel{F(x, y+1) = x} \\ F(x, y+1) = \pi_1^3(x, y, F(x, y)) \end{array} \right.$$

7) minimization

$$g: \mathbb{N}^{n+1} \rightarrow \mathbb{N}$$

$$F: \mathbb{N}^n \rightarrow \mathbb{N}$$

$$F(\bar{x}) = \mu y [g(\bar{x}, y) = 0]$$

$g(\bar{x}, y)$ je def. $\forall z < y$

undef: $\mathbb{N}^0 \rightarrow \mathbb{N}$

$$\text{undef}() = \mu y [1 = 0]$$

complete induction Funder:
 $\mathbb{N}^0 \rightarrow \mathbb{N}^2$

$$a) ((\sigma \circ \xi) \times \xi)(\cdot) =$$

$$(\underbrace{\sigma \circ \xi(\cdot)}_{\sigma(\xi(\cdot))}, \xi(\cdot)) = (\sigma(0), 0) = (1, 0)$$

$$b) (\pi_2^3 \times \pi_3^3 \times \pi_2^3)(5, 6, 7) = (\pi_2^3(5, \underline{6}, 7), \pi_3^3(5, \underline{6}, 7), \pi_2^3(5, \underline{6}, 7)) =$$

$$= (6, 7, 6) \quad \pi_0^3 = ()$$

$$c) ((\sigma \times \sigma) \circ \pi_2^2)(4, 7) = ((\sigma \times \sigma) \circ \pi_2^2(4, 7))$$

$$= ((\sigma \times \sigma)(7)) = (\sigma(7), \sigma(7)) = (8, 8)$$

d)

$$f(x, 0) = \sigma(x)$$

$$f(x, y+1) = \pi_3^3(x, y, f(x, y))$$

$$f(5, 3) = \pi_3^3(5, 2, f(5, 2)) = \pi_3^3(5, 2, 6) = \underline{\underline{6}}$$

$$f(5, 2) = \pi_3^3(5, 1, f(5, 1)) = \pi_3^3(5, 1, 6) = 6$$

$$f(5, 1) = \pi_3^3(5, 0, f(5, 0)) = \pi_3^3(5, 0, 6) = 6$$

$$f(5, 0) = \sigma(5) = 6$$

Ukaže, že funkce swap patří mezi jednoduché trojici
 (x, y, z) dvojici (u, v) , která vznikne 2-násobnou záměnou
 x a y , je primitivně rekurzivní.

$$\begin{aligned} g: \mathbb{N}^2 &\rightarrow \mathbb{N}^2 \\ h: \mathbb{N}^5 &\rightarrow \mathbb{N}^2 \end{aligned}$$

$$\begin{aligned} \text{swap}(x, y, 0) &= \pi_1^2 \times \pi_2^2(x, y) \\ \text{swap}(x, y, z+1) &= \pi_5^5 \times \pi_4^5(x, y, \text{swap}(x, y, z)) \end{aligned}$$

Ukaže, že funkce plus je PR

$$\begin{aligned} \text{plus}(x, 0) &= \pi_1^1(x) \\ \text{plus}(x, y+1) &= \sigma \circ \pi_3^3(x, y, \text{plus}(x, y)) \end{aligned}$$

Ukaže, že funkce mult je PR

$$\begin{aligned} \text{mult}(x, 0) &= \xi \circ \pi_0^1(x) \\ \text{mult}(x, y+1) &= \text{plus}(\pi_1^3 \times \pi_2^3)(x, y, \text{mult}(x, y)) \end{aligned}$$

$$\xi \circ \pi_0^1(x)$$

$$x \cdot (y+1) = x + x \cdot y$$

$$\text{fact}(x) = x!$$

$$(x+1)! = (x+1) \cdot x!$$

$$\text{fact}(0) = \xi()$$

$$\text{fact}(x+1) = \pi_1(x, \text{fact}(x))$$

$$\text{mult}_0 \left(\begin{pmatrix} 0 & \pi_2^2 \\ 1 & \end{pmatrix} \times \pi_2^2 \right)$$

$(x+1) \cdot x!$

~~$$(x+1)!$$~~

$$\text{pred}(x) = \begin{cases} 0 & \text{pozd } x=0 \\ x-1 & \text{pozd } x>0 \end{cases}$$

$$\text{pred}(0) = \xi()$$

$$\text{pred}(x+1) = \pi_1^2(x, \text{pred}(x))$$

$$\text{minus}(x, y) = \begin{cases} x-y & \text{pozd } x \geq y \\ 0 & \text{jinak} \end{cases} \quad x \div (y+1)$$

$$\text{minus}(x, 0) = \pi_1^1(x)$$

$$\text{minus}(x, y+1) = \text{pred} \circ \pi_3^3(x, y, \text{minus}(x, y))$$

$$x \div y$$

$$\text{even}(x): \begin{cases} 1 & \text{pozdol } x \text{ je sude nebo } 0 \\ 0 & \text{jina2} \end{cases}$$

$$\text{even}: \mathbb{N}^1 \rightarrow \mathbb{N}^1$$

~~even =~~

$$\text{even} = \text{ ~~} \pi_0^1 \circ \pi_0^1 \text{ } \rangle~~$$

$$\text{swap}(0,1,3)$$

$$\begin{aligned} (0,1) &\xrightarrow{1} (1,0) \\ &\xrightarrow{2} (0,1) \\ &\xrightarrow{3} (1,0) \end{aligned}$$

$$\pi_2^2 \circ \text{swap} \circ \left(\underbrace{\left(\pi_0^1 \right)}_0 \times \underbrace{\left(\pi_0^1 \right)}_1 \times \underbrace{\pi_1^1}_x \right)$$

$$\text{even}(x) \quad \left\lfloor \frac{x}{2} \right\rfloor \cdot 2 = x$$

$$\text{neg}(x) \begin{cases} 0 & \text{pozdol } x > 0 \\ 1 & \text{pozdol } x = 0 \end{cases}$$

pozitive minus

$$\text{neg} = \text{minus}(1, x)$$

$$1 - x$$

$$F(x, y, z) = \begin{cases} x & \text{je-li } z \text{ sude-} \\ y & \text{je-li } z \text{ liche-} \end{cases}$$

$$F(x, y, z) = \pi_2^2(\text{swap}(y, x, z))$$

$$F(x, y, z) = x \cdot \text{even}(z) + y \cdot \text{neg}(\text{even}(z))$$

Dokážte, že pro parciálně rel. fce

$$f: \mathbb{N} \rightarrow \mathbb{N} \text{ a } g: \mathbb{N} \rightarrow \mathbb{N}$$

není funkce $h: \mathbb{N} \rightarrow \mathbb{N}$ parciálně rel.

$$\text{Zde } h(x) = \begin{cases} f(x) & \text{je-li } f(x) \text{ def.} \\ g(x) & \text{je-li } f(x) \text{ undef.} \end{cases}$$

Díky sporem

Předpoklad je: h je par. rel.

$$f(x) = 1 \quad \text{požad } x \text{ je číslo, jehož binární} \\ \text{zápis odpovídá } \langle \pi \rangle \langle \# \rangle \langle w \rangle \\ \text{Zde } w \in L(\pi) \text{ pro TS } \pi$$

$$g(x) = 0$$

$$h(x) = \begin{cases} 1 & \text{požad } x = \langle \pi \rangle \langle \# \rangle \langle w \rangle \text{ a } w \in L(\pi) \\ 0 & \text{požad } x = \text{---} \text{---} \text{---} \text{ a } w \notin L(\pi) \\ & \text{nebo } x \text{ je neplatný kód} \end{cases}$$

h rozhoduje čistě v RE jazyce \leadsto spor



Načte 2 dispozici funkce $F(x, y)$ a ~~neg~~ (x, y) .

Implementujte v jazyce C funkci

$g(x) = \text{my}[\text{neg}(F(x, y), 1) = 0]$

unsigned $g(\text{unsigned } x) \{$

unsigned $y = 0;$

$\text{while}(F(x, y) \neq 1) \{$

$\quad ++y;$

$\}$

$\text{return } y;$

$\}$

5 funkcij: for, minus, mult, div, eq, neg zapiske
 funkcije $\text{gcd}(x, y)$ - največji skupni delitelj x in y

$$\text{gcd}(x, y) = \max(\{d \in \mathbb{N} \mid \exists a \in \mathbb{N}. x = a \cdot d \wedge \exists b \in \mathbb{N}. y = b \cdot d\})$$

predpostavljamo $x \leq y$

$$1 \leq \text{gcd}(x, y) \leq x$$

pač se ničeno staviti najit največji-2 faktor, če

$$\text{gcd}(x, y) = x \div 2$$

$$\text{gcd}_{x \leq y}(x, y) = \text{minus}(x, \mu 2 [\dots = 0])$$

$$\text{divides}(x, y) = \begin{cases} 1 & \text{if } x \mid y \quad (y \% x = 0) \\ 0 & \text{if } \neg(x \mid y) \quad (y \% x > 0) \\ \text{undef} & \text{if } x = 0 \end{cases}$$

$$\text{divides}(x, y) = \text{eq}(\text{mult}(\text{div}(y, x), x), y)$$

$$\left\lfloor \frac{y}{x} \right\rfloor \cdot x = y$$

$$\text{gcd}_{x \leq y}(x, y) = \text{minus}(x, \mu 2 [\text{neg}(\text{divides}(x \div 2, x)) + \text{neg}(\text{divides}(x \div 2, y)) = 0])$$