

Advanced Topics in Malware Analysis

High Level Language Constructs
In Assembly

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Function calls, If/then/else, Switch



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Learning Objectives

- Interpret high level structures in assembler
- Identify function calls
- Reconstruct switch expressions to index table of targets
- Employ loops and compiler-induced elimination of branching
- Discover arrays
- Identify structures in assembly code



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Understanding HLL Structures in ASM

Need to understand representation of compiled high-level structures in assembler

- Functions
 - Calls and Returns discussed earlier
- Control structures
 - Loops
 - Conditional branches
 - if/then/else
 - switch statements
- Data structures
 - structs/unions
 - Arrays
 - Linked structures



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IDA Annotated Function Calls

Full annotation
(PUSHed args)

.text:004010F3	push	10h	; namelen
.text:004010F5	lea	ecx, 0[ebp+name]	
.text:004010F8	push	ecx	; name
.text:004010F9	mov	edx, 0[ebp+s]	
.text:004010FF	push	edx	; s
.text:00401100	call	connect	

Partial annotation
(MOVs)

.text:004011A5	mov	[esp+244h+var_23C], 10h
.text:004011AD	lea	eax, 0[ebp+var_28]
.text:004011B0	mov	[esp+244h+var_240], eax
.text:004011B4	mov	eax, 0[ebp+var_C]
.text:004011B7	mov	[esp+244h+var_244], eax
.text:004011BA	call	connect



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if/then/else

Typically an assembly sequence that:

1. Computes the expression in the if portion
2. Performs a test or cmp to ensure that CPU flags are set
3. Performs a jnz/jge/etc. and a jmp

```

cmp [z], 5      ; if z < 5 then
jge Else
call func_if    ; func_if()
jmp After
Else:
call func_else ; func_else()
After:          ; end if

```



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Switch Case Statements

Several possible implementations:

- **Table Implementation**
 - Use a table of targets for an unconditional JMP
 - Use switch expression as index into the table of targets
 - This implementation isn't appropriate when there are numerous gaps in the case statements
 - Problem: Jump table may be embedded in the code section—problem for linear sweep disassembly
- **Tree Implementation**
 - Implements a binary search for the matching case

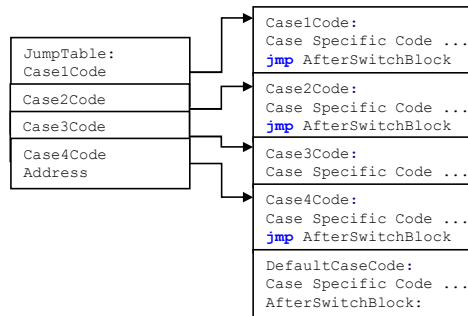


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Switch: Table Implementation “Idea”

```
switch (a) {
  case 1:
    Case Specific Code ...
    break;
  case 2:
    Case Specific Code ...
    break;
  case 3: // no break!
    Case Specific Code ...
  case 4:
    Case Specific Code ...
    break;
  default:
    Case Specific Code ...
}
```

```
movzx eax, BYTE PTR [a]
sub eax, 1
cmp eax, 3
ja DefaultCaseCode
jmp DWORD PTR [JumpTable + eax * 4]
```



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Switch: Table Implementation Real Code

- Compiled with gcc on 64-bit Mac with no optimizations (i.e., -O0)

```
#include <stdio.h>
void main(void) {
  int a;
  scanf("%d", &a);
  switch (a) {
    case 0:
      puts ("Zero");
      break;
    case 1:
      puts ("One");
      break;
    case 2:
      puts ("Two");
      break;
    case 3:
      puts ("Three");
      break;
    default:
      puts("No idea.");
  }
}
```

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Switch: Table Implementation Real Code

```
#include <stdio.h>
void main(void) {
    int a;
    scanf("%d", &a);
    switch (a) {
        case 0:
            puts ("Zero");
            break;
        case 1:
            puts ("One");
            break;
        case 2:
            puts ("Two");
            break;
        case 3:
            puts ("Three");
            break;
        default:
            puts("No idea.");
    }
}
```

```

cstring:00000000100000F46 ; =====
cstring:00000000100000F46
cstring:00000000100000F46 ; Segment type: Pure data
cstring:00000000100000F46 cstring segment byte public
cstring:00000000100000F46 assume cs: cstring
cstring:00000000100000F46 ;org 100000F46h
cstring:00000000100000F46 ; char aD[]
cstring:00000000100000F46 ad db 'd',0
cstring:00000000100000F49 ; char aZero[] db 'Zero',0
cstring:00000000100000F49 aZero
cstring:00000000100000F4E ; char aOne[] db 'One',0
cstring:00000000100000F4E aOne
cstring:00000000100000F52 ; char aTwo[] db 'Two',0
cstring:00000000100000F52 aTwo
cstring:00000000100000F56 ; char aThree[] db 'Three',0
cstring:00000000100000F56 aThree
cstring:00000000100000F5C ; char aNOIdea_[] db 'No idea.',0
cstring:00000000100000F5C aNOIdea_
cstring:00000000100000F5C cstring ends
cstring:00000000100000F5C

```



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Switch: Table Implementation Real Code

```
#include <stdio.h>
void main(void) {
    int a;
    scanf("%d", &a);
    switch (a) {
        case 0:
            puts ("Zero");
            break;
        case 1:
            puts ("One");
            break;
        case 2:
            puts ("Two");
            break;
        case 3:
            puts ("Three");
            break;
        default:
            puts("No idea.");
    }
}
```

```

_main                                     public _main
proc near                                ; CODE XREF: start

var_10                                  = qword ptr -10h
var_4                                   = dword ptr -4

push    rbp
mov     rbp, rsp
sub     rsp, 10h
lea     rax, [rbp+var_4]
xor     cl, cl
lea     rdx, aD                          ; "d"
mov     rdi, rdx                          ; char *
mov     rsi, rax
mov     al, cl
call    _scanf
eax, [rbp+var_4]
cmp     eax, 3                            ; switch 4 cases
mov     eax, eax
mov     [rbp+var_10], rax
short loc_100000EDC                       ; default case
lea     rax, off_100000EF4
mov     rcx, [rbp+var_10]
mov     rdx, rcx
movsxd rdx, dword ptr [rax+rdx*4]
lea     rax, [rdx+rax]
jmp     rax                                ; switch jump
;

```



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Switch: Table Implementation Real Code

```
#include <stdio.h>
void main(void) {
    int a;
    scanf("%d", &a);
    switch (a) {
        case 0:
            puts ("Zero");
            break;
        case 1:
            puts ("One");
            break;
        case 2:
            puts ("Two");
            break;
        case 3:
            puts ("Three");
            break;
        default:
            puts("No idea.");
    }
}
```

```
loc_100000EDC:  lea     rax, aNoIdea_      ; CODE XREF: _main+2E1j
                  mov     rdi, rax      ; jumtable 00000001000000E96 default
                  call    _puts        ; char *

loc_100000EEB:  add     rsp, 10h          ; CODE XREF: _main+571j
                  pop     rbp          ; _main+681j ...
_main:          retn
                endp
```



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Switch: Table Implementation Real Code

```
#include <stdio.h>
void main(void) {
    int a;
    scanf("%d", &a);
    switch (a) {
        case 0:
            puts ("Zero");
            break;
        case 1:
            puts ("One");
            break;
        case 2:
            puts ("Two");
            break;
        case 3:
            puts ("Three");
            break;
        default:
            puts("No idea.");
    }
}
```

```
_main          public _main
                proc near          ; CODE XREF: start

var_10          = qword ptr -10h
var_4           = dword ptr -4

                push    rbp
                mov     rbp, rsp
                sub     rsp, 10h
                lea     rax, [rbp+var_4]
                xor     cl, cl
                lea     rdx, aD      ; "%d"
                mov     rdi, rdx      ; char *
                mov     rsi, rax
                mov     al, cl
                call    _scanf
                mov     eax, [rbp+var_4]
                cmp     eax, 3        ; switch 4 cases
                mov     eax, eax
                mov     [rbp+var_10], rax
                ja      short loc_100000EDC ; default case
                lea     rax, off_100000EF4
                mov     rcx, [rbp+var_10]
                mov     rdx, rcx
                movsxd  rdx, dword ptr [rax+rdx*4]
                lea     rax, [rdx+rax]
                jmp     rax          ; switch jump
```



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Switch: Table Implementation Real Code

```
#include <stdio.h>
void main(void) {
    int a;
    scanf("%d", &a);
    switch (a) {
    case 0:
        puts ("Zero");
        break;
    case 1:
        puts ("One");
        break;
    case 2:
        puts ("Two");
        break;
    case 3:
        puts ("Three");
        break;
    default:
        puts("No idea.");
    }
}
```

```

; -----
; align 4
off_10000EF4 dd offset loc_10000E98 - 10000EF4h
; DATA XREF: main+307o
dd offset loc_10000EA9 - 10000EF4h ; jump table for switch
dd offset loc_10000EBA - 10000EF4h
dd offset loc_10000ECB - 10000EF4h
ends
__text

```



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Switch: Table Implementation Real Code

```
#include <stdio.h>
void main(void) {
    int a;
    scanf("%d", &a);
    switch (a) {
    case 0:
        puts ("Zero");
        break;
    case 1:
        puts ("One");
        break;
    case 2:
        puts ("Two");
        break;
    case 3:
        puts ("Three");
        break;
    default:
        puts("No idea.");
    }
}
```

```

__main
var_10
var_4

public __main
proc near
; CODE XREF: start

    = qword ptr -10h
    = dword ptr -4

    push rbp
    mov rbp, rsp
    sub rsp, 10h
    lea rax, [rbp+var_4]
    xor cl, cl
    lea rdx, ad ; "%d"
    mov rdi, rdx ; char *
    mov rsi, rax
    mov al, cl
    call _scanf
    mov eax, [rbp+var_4]
    cmp eax, 3 ; switch 4 cases
    mov eax, eax
    mov [rbp+var_10], rax
    ja short loc_10000EDC ; default case
    lea rax, off_10000EF4
    mov rcx, [rbp+var_10]
    mov rdx, rcx
    movsxd rdx, dword ptr [rax+rdx*4]
    lea rax, [rdx+rax]
    jmp rax ; switch jump

```



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Switch: Table Implementation Real Code

```
#include <stdio.h>
void main(void) {
    int a;
    scanf("%d", &a);
    switch (a) {
        case 0:
            puts ("Zero");
            break;
        case 1:
            puts ("One");
            break;
        case 2:
            puts ("Two");
            break;
        case 3:
            puts ("Three");
            break;
        default:
            puts("No idea.");
    }
}
```

```
loc_10000E98:                ; CODE XREF: _main+467j
                                ; DATA XREF: _text:off 10000EF4j
                                ; jumtable 0000000100000E96 case 0
                                ; char "
lea     rax, aZero
mov     rdi, rax
call    _puts
jmp     short loc_10000EEB

loc_10000EA9:                ; CODE XREF: _main+467j
                                ; DATA XREF: _text:off 10000EF4j
                                ; jumtable 0000000100000E96 case 1
                                ; char "
lea     rax, aOne
mov     rdi, rax
call    _puts
jmp     short loc_10000EEB

loc_10000EBA:                ; CODE XREF: _main+467j
                                ; DATA XREF: _text:off 10000EF4j
                                ; jumtable 0000000100000E96 case 2
                                ; char "
lea     rax, aTwo
mov     rdi, rax
call    _puts
jmp     short loc_10000EEB

loc_10000ECB:                ; CODE XREF: _main+467j
                                ; DATA XREF: _text:off 10000EF4j
                                ; jumtable 0000000100000E96 case 3
                                ; char "
lea     rax, aThree
mov     rdi, rax
call    _puts
jmp     short loc_10000EEB

loc_10000EDC:                ; CODE XREF: _main+2E7j
                                ; jumtable 0000000100000E96 default
                                ; char "
lea     rax, aNoIdea
mov     rdi, rax
call    _puts

loc_10000EEB:                ; CODE XREF: _main+577j
                                ; _main+687j ...
add     rsp, 10h
pop     rbp
ret

_main
endp

off_10000EF4 align 4
dd offset loc_10000E98 - 10000EF4h ; DATA XREF: _main+307o
dd offset loc_10000EA9 - 10000EF4h ; jumtable 0000000100000E96 case 1
dd offset loc_10000EBA - 10000EF4h ; jumtable 0000000100000E96 case 2
dd offset loc_10000ECB - 10000EF4h ; jumtable 0000000100000E96 case 3
ends

_text
```



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Compiled with -O2

```
#include <stdio.h>
void main(void) {
    int a;
    scanf("%d", &a);
    switch (a) {
        case 0:
            puts ("Zero");
            break;
        case 1:
            puts ("One");
            break;
        case 2:
            puts ("Two");
            break;
        case 3:
            puts ("Three");
            break;
        default:
            puts("No idea.");
    }
}
```

```
lea     rdi, aD
xor     rax, rax
call    _scanf
mov     eax, [rbp+var_4]
cmp     rax, 3
ja      short loc_10000EE0 ; default case
lea     rcx, off_10000EEC
mov     rcx, dword ptr [rcx+rax*4]
add     rax, rcx
jmp     rax ; switch jump

loc_10000EB3:                ; CODE XREF: _main+317j
                                ; DATA XREF: _text:off 10000EECj
                                ; jumtable 0000000100000EB1 case 0
lea     rdi, aZero
call    _puts
add     rsp, 10h
pop     rbp
ret

loc_10000EBA:                ; CODE XREF: _main+4C7j
                                ; _main+557j ...
call    _puts

loc_10000EC5:                ; CODE XREF: _main+317j
                                ; DATA XREF: _text:off 10000EECj
                                ; jumtable 0000000100000EB1 case 1
lea     rdi, aOne
jmp     short loc_10000EBA

loc_10000ECE:                ; CODE XREF: _main+317j
                                ; DATA XREF: _text:off 10000EECj
                                ; jumtable 0000000100000EB1 case 2
lea     rdi, aTwo
jmp     short loc_10000EBA

loc_10000ED7:                ; CODE XREF: _main+317j
                                ; DATA XREF: _text:off 10000EECj
                                ; jumtable 0000000100000EB1 case 3
lea     rdi, aThree
jmp     short loc_10000EBA

loc_10000EE0:                ; CODE XREF: _main+217j
                                ; jumtable 0000000100000EB1 default
lea     rdi, aNoIdea
jmp     short loc_10000EBA

_main
endp

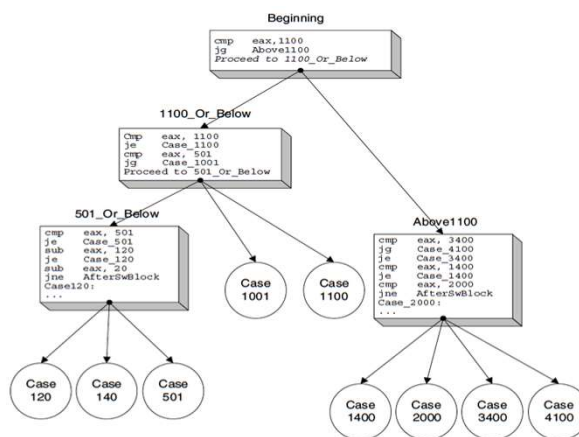
off_10000EEC align 4
dd offset loc_10000EB3 - 10000EECh ; DATA XREF: _main+237o
dd offset loc_10000EC5 - 10000EECh ; jumtable 0000000100000EB1 case 1
dd offset loc_10000ECE - 10000EECh ; jumtable 0000000100000EB1 case 2
dd offset loc_10000ED7 - 10000EECh ; jumtable 0000000100000EB1 case 3
```



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Switch: Binary Search

```
switch (value) {
  case 120: Code...
  break;
  case 140: Code...
  break;
  case 501: Code...
  break;
  case 1001: Code...
  break;
  case 1100: Code...
  break;
  case 1400: Code...
  break;
  case 2000: Code...
  break;
  case 3400: Code...
  break;
  case 4100: Code...
  break;
};
```



Source: Eilam, E. (2011). *Reversing: secrets of reverse engineering*. Hoboken, NJ: John Wiley & Sons.



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High Level Language Constructs In Assembly

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Loops



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Loops

- Loops are much as you'd expect, except...
- Pretested loops (e.g., while/do) can end up being post-tested (do/while)

```
c = 0;
while (c < 1000)
{
    array[c] = c;
    c++;
}
```

<code>mov ebx, DWORD PTR [array]</code>	<code>; c = 0</code>
<code>xor ecx, ecx</code>	<code>; do {</code>
<code>Loop:</code>	<code>; array[c]=c</code>
<code>mov DWORD PTR [ebx+ecx*4], ecx</code>	<code>; c++</code>
<code>add ecx, 1</code>	<code>; }</code>
<code>cmp ecx, 1000</code>	<code>; while (c < 1000)</code>
<code>jle Loop</code>	



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Loops: Rep Instruction

- **rep** instruction prefix
 - Syntax: REP [INS] (INS operands)
- Repeats an instruction the number of times specified in the RCX/ECX/CX register or until the indicated condition of the ZF flag is no longer met

Repeat Prefix	Termination Condition 1	Termination Condition 2
REP	ECX=0	None
REPE/REPZ	ECX=0	ZF=0
REPNE/REPNZ	ECX=0	ZF=1

- Used to replace loops which consists of only one instruction (most often moving data)
 - To repeat a block of instructions, the LOOP instruction can be used (very rare) or a logical looping construct



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Rep Example: Libc memcpy

```

__memcpy_g_internal proc near                ; CODE XREF: __strcat_c+364p
                                              ; __memcpy_c+264p

arg_0      = dword ptr  4
arg_4      = dword ptr  8
arg_8      = dword ptr 0Ch

    push    ebp                            ; save clobbered regs
    push    edi
    push    esi
    mov     eax, [esp+0Ch+arg_0] ; eax = dest
    mov     ebp, [esp+0Ch+arg_4] ; ebp = source
    mov     ecx, [esp+0Ch+arg_8] ; ecx = size
    mov     edi, eax                ; edi = dest
    mov     esi, ebp                ; esi = source
    cld                             ; direction flag = 0, inc pointer after iteration
    shr     ecx, 1                  ; ecx = ecx / 2
    jnb     short loc_83879         ; jump if bit shifted off ecx was 0, otherwise ...
    movsb                                     ; move 1 byte from [esi] to [edi] & inc esi and edi

loc_83879:
    shr     ecx, 1                  ; CODE XREF: __memcpy_g_internal+164j
    jnb     short loc_8387F         ; ecx = ecx / 2
    movsw                                     ; jump if bit shifted off ecx was 0, otherwise ...
    movsw                                     ; move 1 word from [esi] to [edi] & add 2 to esi and edi

loc_8387F:
    rep movsd                         ; CODE XREF: __memcpy_g_internal+1B4j
    ; while(ecx > 0)
    ; move 1 DWORD from [esi] to [edi]
    ; add 4 to esi and edi
    ; decrement ecx

    pop     esi
    pop     edi
    pop     ebp
    retn                                     ; restore saved regs
                                              ; return to caller

```



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Loop Unrolling

- Duplicate loop body and reduces the number of iterations
- **Goal:** Eliminate as much branching as possible at the expense of increased code size

```
int main(int argc, char* argv[])
{
    char dest[64];
    for(int i = 0; i < 64; i++)
        dest[i] = argv[1][i];
    puts(dest);
}
```

- gcc -W -Wall -Wextra -Wpedantic -O2 -S -masm=intel -funroll-loops fill_array.c -o fill_array.s
- Compiled on 64-bit Windows

```
main:
    push rbx ;save rbx
    sub rsp, 96
    mov rbx, rdx ; rbx = argv
    mov rdx, QWORD PTR 8[rbx] ; rdx = argv[1]
    lea rcx, 32[rsp] ; rcx = dest
    xor eax, eax

.L2:
    movzx r8d, BYTE PTR [rdx+rax]
    mov BYTE PTR [rcx+rax], r8b
    movzx r9d, BYTE PTR 1[rdx+rax]
    mov BYTE PTR 1[rcx+rax], r9b
    movzx r10d, BYTE PTR 2[rdx+rax]
    mov BYTE PTR 2[rcx+rax], r10b
    movzx r11d, BYTE PTR 3[rdx+rax]
    mov BYTE PTR 3[rcx+rax], r11b
    movzx ebx, BYTE PTR 4[rdx+rax]
    mov BYTE PTR 4[rcx+rax], bl
    movzx r8d, BYTE PTR 5[rdx+rax]
    mov BYTE PTR 5[rcx+rax], r8b
    movzx r9d, BYTE PTR 6[rdx+rax]
    mov BYTE PTR 6[rcx+rax], r9b
    movzx r10d, BYTE PTR 7[rdx+rax]
    mov BYTE PTR 7[rcx+rax], r10b
    add rax, 8
    cmp rax, 64
    jne .L2
    call puts
    xor eax, eax
    add rsp, 96
    pop rbx
    ret
```



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Loop Unrolling

- **Extreme Case:** Complete unroll
 - No loop — just replicated loop body

```
int main(int argc, char* argv[])
{
    char dest[64];
    for(int i = 0; i < 64; i++)
        dest[i] = argv[1][i];
    puts(dest);
}
```

- gcc -W -Wall -Wextra -Wpedantic -O3 -S -masm=intel fill_array.c -o fill_array.s
- Compiled on 64-bit Windows

```
main:
    push rbx ;save rbx
    sub rsp, 96
    mov rbx, rdx ; rbx = argv
    mov r8, QWORD PTR 8[rbx] ; r8=argv[1]
    ; copy 64 bytes unrolled
    mov rax, QWORD PTR [r8]
    mov QWORD PTR 32[rsp], rax
    mov rax, QWORD PTR 8[r8]
    mov QWORD PTR 40[rsp], rax
    mov rax, QWORD PTR 16[r8]
    mov QWORD PTR 48[rsp], rax
    mov rax, QWORD PTR 24[r8]
    mov QWORD PTR 56[rsp], rax
    mov rax, QWORD PTR 32[r8]
    mov QWORD PTR 64[rsp], rax
    mov rax, QWORD PTR 40[r8]
    mov QWORD PTR 72[rsp], rax
    mov rax, QWORD PTR 48[r8]
    mov QWORD PTR 80[rsp], rax
    mov rax, QWORD PTR 56[r8]
    mov QWORD PTR 88[rsp], rax

    call puts
    xor eax, eax
    add rsp, 96
    pop rbx
    ret
```



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Arrays



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Arrays

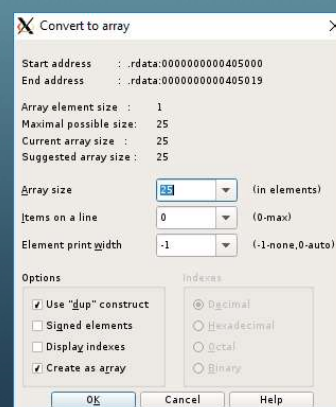
- Arrays indexed by a variable can generally be located easily in ASM
- Sometimes possible to determine length of array by closely analyzing code
- Arrays indexed by a constant are harder
 - Compiler will optimize
 - ASM looks like structure accesses or accesses to individual variables
- Accesses to dynamically allocated arrays are easier to deduce
 - There will generally be a call to a memory allocation function
 - Accesses to individual elements are offsets from the base address of the allocation



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Arrays In IDA

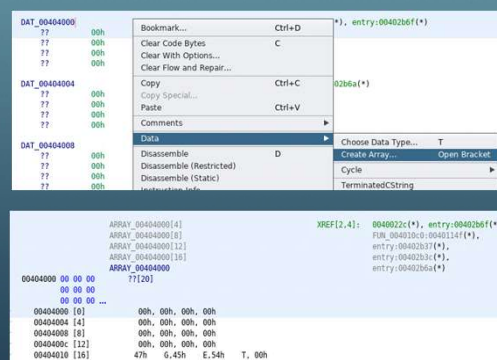
- IDA will try to be smart about array identification
- It can miss some
 - Especially arrays of union data types!
- IDA can be coaxed into displaying arrays properly via Right-Click->Array
 - see Eagle, Chapter 8 for more



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Arrays In Ghidra

- Ghidra has similar functionality
 - A little less fleshed out
- Right Click → Data → Create Array
- Then set the size in bytes



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Global Array Example

```
int global_array[3];

int main() {
    int idx = 2;
    global_array[0] = 10;
    global_array[1] = 20;
    global_array[2] = 30;
    global_array[idx] = 40;
}
```



```
.text:00401000 _main
.text:00401000
.text:00401000 idx
.text:00401000
.text:00401000
.text:00401001
.text:00401003
.text:00401004
.text:00401008
.text:00401015
.text:0040101F
.text:00401029
.text:0040102C
.text:00401037
.text:00401039
.text:0040103B
.text:0040103C
.text:0040103C _main

proc near
= dword ptr -4

push    ebp
mov     ebp, esp
push    ecx
mov     [ebp+idx], 2
①mov    dword_40B720, 10
②mov    dword_40B724, 20
③mov    dword_40B728, 30
mov     eax, [ebp+idx]
④mov    dword_40B720[edx*4], 40
xor     eax, eax
mov     esp, ebp
pop     ebp
retn
endp
```

Global/static array example

- (1), (2), (3) are constant accesses
- (4) is an indexed access
- Indexed access points out size of individual elements

Source: Eagle, C. (2011). *The Ida Pro Book: The unofficial guide to the worlds most popular disassembler*. San Francisco: No Starch Press.

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heap Array Example

```
int main() {
    int *heap_array = (int*)malloc(3 * sizeof(int));
    int idx = 2;
    heap_array[0] = 10;
    heap_array[1] = 20;
    heap_array[2] = 30;
    heap_array[idx] = 40;
}
```



- Heap-allocated array
- (5) reveals size
- (1) – (4) are indexed accesses

Source: Eagle, C. (2011). *The Ida Pro Book: The unofficial guide to the worlds most popular disassembler*. San Francisco: No Starch Press.

```
.text:00401000 _main      proc near
.text:00401000
.text:00401000 heap_array  = dword ptr -8
.text:00401000 idx          = dword ptr -4
.text:00401000
.text:00401000      push     ebp
.text:00401001      mov      ebp, esp
.text:00401003      sub      esp, 8
.text:00401006      push     0Ch          ; size_t
.text:00401008      call     _malloc
.text:0040100D      add      esp, 4
.text:00401010      mov      [ebp+heap_array], eax
.text:00401013      mov      [ebp+idx], 2
.text:0040101A      mov      eax, [ebp+heap_array]
.text:0040101D      mov      dword ptr [eax], 10
.text:00401023      mov      ecx, [ebp+heap_array]
.text:00401026      mov      dword ptr [ecx+4], 20
.text:0040102D      mov      edx, [ebp+heap_array]
.text:00401030      mov      dword ptr [edx+8], 30
.text:00401037      mov      eax, [ebp+idx]
.text:0040103A      mov      ecx, [ebp+heap_array]
.text:0040103D      mov      dword ptr [ecx+eax*4], 40
.text:00401044      xor      eax, eax
.text:00401046      mov      esp, ebp
.text:00401048      pop      ebp
.text:00401049      retn
.text:00401049 _main      endp
```



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Advanced Topics in Malware Analysis

High Level Language Constructs
In Assembly

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Struct/Union Analysis



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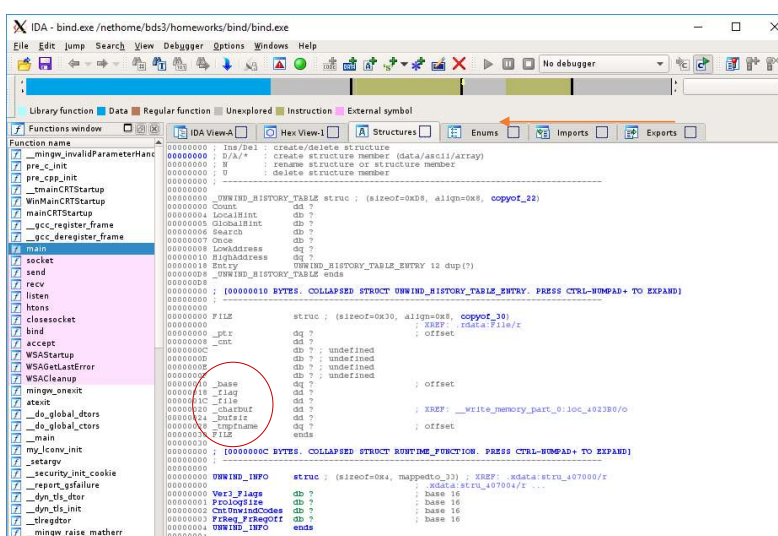
Struct/Union/etc.

- Field names are lost
 - Only remain with debugging symbols (which malware doesn't like)
 - Will be replaced with offsets by the compiler
- Global/static allocation:
 - Accesses to structure elements look like accesses to arrays via constant indices
- Dynamic allocation:
 - Similar to arrays, seeing memory allocation may help
- Struct / Union analysis is complicated by alignment issues!
- Good news! For common structures (e.g., struct sockaddr, etc.), IDA Pro may be able to automatically recognize struct accesses
- If recognition for particular struct types isn't available, you can add it manually in IDA



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Struct/Union in ida

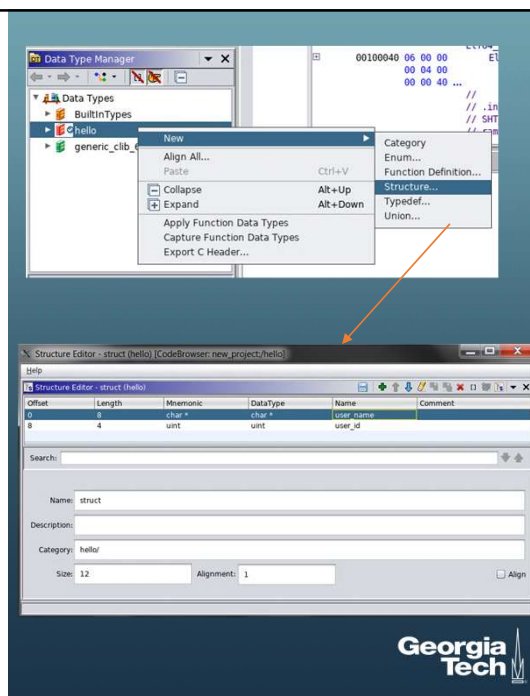


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Struct/Union in Ghidra

- Ghidra does structs/unions much cleaner than IDA!
- Plus this functionality is extremely useful when paired with Ghidra's decompiler
- E.g.: If Ghidra detects a class, you can define each of its members in memory
- Makes decompiled code *significantly* more readable



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Structure Alignment Issues

- Structures defined in source code only provide field names & types
- Compilers must lay those fields out in memory
 - Different compilers/flags will choose different field alignments
- For example: A memory-efficient compilation should choose a minimal size (packed)
- But: A processor-cache-efficient compilation may choose a power of 2 alignment
- This can confuse reverse engineers --- is the extra space actually another field?

```

struct ch8_struct { //Size  Minimum offset  Default offset
  int field1;      // 4      0              0
  short field2;    // 2      4              4
  char field3;     // 1      6              6
  int field4;      // 4      7              8
  double field5;   // 8      11             16
};                //Minimum total size: 19  Default size: 24

```

Source: Eagle, Chris. The IDA pro book.
No Starch Press, 2011



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Global struct Example

```

struct ch8_struct global_struct;

int main() {
  global_struct.field1 = 10;
  global_struct.field2 = 20;
  global_struct.field3 = 30;
  global_struct.field4 = 40;
  global_struct.field5 = 50.0;
}

```



```

.text:00401000 _main
.text:00401000
.text:00401001
.text:00401003
.text:0040100D
.text:00401016
.text:0040101D
.text:00401027
.text:0040102D
.text:00401033
.text:00401035
.text:00401036
.text:00401036 _main

```

```

proc near
push    ebp
mov     ebp, esp
mov     dword_40EA60, 10
mov     word_40EA64, 20
mov     byte_40EA66, 30
mov     dword_40EA68, 40
fld     ds:dbl_40B128
fstp    dbl_40EA70
xor     eax, eax
pop     ebp
retn
endp

```

Global struct

- All offsets computed at compile time
- No runtime computation
- Hard to distinguish between struct and array access, except that elements of C arrays have equal size

Source: Eagle, C. (2011). *The Ida Pro Book: The unofficial guide to the worlds most popular disassembler*. San Francisco: No Starch Press.



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Heap struct Example

```

.text:00401004      ①push    24                ; size_t
.text:00401006      call     _malloc
.text:00401008      add     esp, 4
.text:0040100E      mov     [ebp+heap_struct], eax
.text:00401011      mov     eax, [ebp+heap_struct]
.text:00401014      ①mov     dword ptr [eax], 10
.text:0040101A      mov     ecx, [ebp+heap_struct]
.text:0040101D      ②mov     word ptr [ecx+4], 20
.text:00401023      mov     edx, [ebp+heap_struct]
.text:00401026      ③mov     byte ptr [edx+6], 30
.text:0040102A      mov     eax, [ebp+heap_struct]
.text:0040102D      ④mov     dword ptr [eax+8], 40
.text:00401034      mov     ecx, [ebp+heap_struct]

```

Heap-allocated struct

- As for dynamically allocated arrays, size is known from malloc() call
- Next slide: struct is packed (single byte alignment)

Source: Eagle, C. (2011). *The Ida Pro Book: The unofficial guide to the worlds most popular disassembler*. San Francisco: No Starch Press.



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Same struct, but compiler uses different alignment

```

.text:0040100E      mov     [ebp+heap_struct], eax
.text:00401011      mov     eax, [ebp+heap_struct]
.text:00401014      mov     dword ptr [eax], 10
.text:0040101A      mov     ecx, [ebp+heap_struct]
.text:0040101D      mov     word ptr [ecx+4], 20
.text:00401023      mov     edx, [ebp+heap_struct]
.text:00401026      mov     byte ptr [edx+6], 30
.text:0040102A      mov     eax, [ebp+heap_struct]
.text:0040102D      mov     dword ptr [eax+7], 40

```

Source: Eagle, C. (2011). *The Ida Pro Book: The unofficial guide to the worlds most popular disassembler*. San Francisco: No Starch Press.



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Array of Structs, Data Structures, Oh My!



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Recall The Structure Alignment

struct ch8_struct {	//Size	Minimum offset	Default offset
int field1;	// 4	0	0
short field2;	// 2	4	4
char field3;	// 1	6	6
int field4;	// 4	7	8
double field5;	// 8	11	16
}; //Minimum total size: 19 Default size: 24			

Source: Eagle, C. (2011). *The Ida Pro Book: The unofficial guide to the worlds most popular disassembler*. San Francisco: No Starch Press.



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Getting tougher: Arrays of structs

```

int main() {
    int idx = 1;
    struct ch8_struct *heap_struct;
    heap_struct = (struct ch8_struct*)malloc(sizeof(struct ch8_struct) * 5);
    heap_struct[idx].field1 = 10;
}

```

.text:00401000	_main	proc near
.text:00401000		
.text:00401000	idx	= dword ptr -8
.text:00401000	heap_struct	= dword ptr -4
.text:00401000		
.text:00401000		push ebp
.text:00401001		mov ebp, esp
.text:00401003		sub esp, 8
.text:00401006		mov [ebp+idx], 1
.text:0040100D		push 120 ; size_t
.text:0040100F		call _malloc
.text:00401014		add esp, 4
.text:00401017		mov [ebp+heap_struct], eax
.text:0040101A		mov eax, [ebp+idx]
.text:0040101D		imul eax, 24
.text:00401020		mov ecx, [ebp+heap_struct]
.text:00401023		mov dword ptr [ecx+eax], 10
.text:0040102A		xor eax, eax
.text:0040102C		mov esp, ebp
.text:0040102E		pop ebp
.text:0040102F		retn
.text:0040102F	_main	endp

Source: Eagle, C. (2011). *The Ida Pro Book: The unofficial guide to the worlds most popular disassembler*. San Francisco: No Starch Press.



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Reverse Engineering A Linked List

- Compiled with gcc on 64-bit Windows with light optimizations (i.e., -O1)
- Recall: Windows ABI says 1st arg in RCX!

```
struct node {
    struct node *next;
    int data;
};

void print(struct node *ptr) {
    while (ptr != NULL) {
        printf("%d ", ptr->data);
        ptr = ptr->next;
    }
    printf("\n");
}
```

```
print:
    push    rsi
    push    rbx
    sub     rsp, 40
    mov     rbx, rcx
    test    rcx, rcx
    je      .L7
    lea     rsi, .LC0[rip] ; "%d "
.L8:
    mov     edx, DWORD PTR 8[rbx]
    mov     rcx, rsi
    call    printf
    mov     rbx, QWORD PTR [rbx]
    test    rbx, rbx
    jne     .L8
.L7:
    mov     ecx, 10
    call    putchar
    add     rsp, 40
    pop     rbx
    pop     rsi
    ret
```



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Always More Complex Structures

- Trees
- Heaps
- Trees of Heaps of Stacks
- ...
- C++ / object oriented languages
- Borland Delphi (!)
- Objective C is a counterexample because of rich runtime system
- Much more difficult
- Eventually, you can't identify a structure... just a collection of data blobs
- Pray for malware to be written in ASM



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Importing Shared Function



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Importing Shared Functions

- External code (i.e., libraries) can be compiled statically or dynamically
- Static is easy, you have the entire library function baked into the single binary
- Dynamic is harder, the loader needs to patch the function addresses at runtime
- Different executable formats follow different conventions

```

mov     rax, [rbp+var_10]
add     rax, 8
mov     rax, [rax]
rsi,    rax
mov     edi, offset format ; "Hello %s\n"
mov     eax, 0
call    _printf

; Segment type: Pure code
; Segment permissions: Read/Execute
_plt
segment para public 'CODE' use64
assume cs: _plt
;org 4003F0b
assume es:nothing, ss:nothing, ds:_data, fs:no
dq 2 dup(?)

; int printf(const char *format, ...)
_printf
proc near
jmp     cs:off_601018 ; CODE XREF: main+2D4ip
endp

; Segment type: Pure data
; Segment permissions: Read/Write
; Segment alignment 'qword' can not be represented in assembly
_got_plt
segment para public 'DATA' use64
assume cs:_got_plt
;org 601000
GLOBAL OFFSET TABLE db ? ;
off_601018 dq offset _printf ; DATA XREF: _printf$r
off_601020 dq offset __libc_start_main ; DATA XREF: __libc_start_main$r

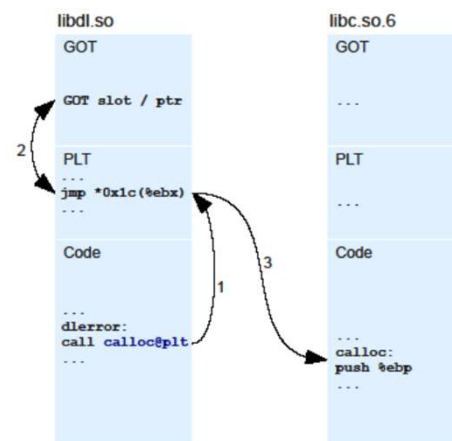
```



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ELF Shared libraries (PLT + GOT)

- Global Offset Table
 - GOT ... not Game of Thrones ☹
 - Contains pointers (**data**) to symbols imported from other shared objects (libraries)
- Procedure Linkage Table
 - PLT
 - Contains **code** that transfers control through the GOT to a symbol in the shared object
- The entries in the GOT are initialized to point to a function in the loader to resolve symbols on-the-fly
 - Before first use, all GOT entries point to loader
 - During first call, the loader executes, loads the symbol (or loads the whole library), and patches the GOT address
 - All later calls go directly to shared object



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Check it out in IDA

```

mov     rax, [rbp+var_10]
add     rax, 8
mov     rax, [rax]
mov     rsi, rax
mov     edi, offset format ; "Hello %s\n"
mov     eax, 0
call    _printf

; Segment type: Pure code
; Segment permissions: Read/Execute
; Segment: .plt
; Segment para public 'CODE' use64
; assume cs:_plt
; org 4003F0B
; assume es:nothing, ss:nothing, ds:_data, fs:nothing
; dq 2 dup(?)

; int printf(const char *format, ...)
; _printf
; _printf
; proc near
; jmp     cs:off_601018 ; CODE XREF: main+2Dip
; endp

; =====
; Segment type: Pure data
; Segment permissions: Read/Write
; Segment alignment 'qword' can not be represented in assembly
; Segment: .got_plt
; Segment para public 'DATA' use64
; assume cs:_got_plt
; org 60100Bh
; dq 60100Bh
; GLOBAL OFFSET TABLE db ?
; off_601018 dq offset printf ; DATA XREF: _printf
; off_601020 dq offset __libc_start_main ; DATA XREF: __libc_start_main

```



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PE Shared libraries (IAT + .idata)

- The Windows loader is part of the OS Kernel
- The .idata section of the PE file contains
 - Import Lookup Table
 - Array of structures, each contains an ordinal or RVA of a name for each imported function
 - The ordinal represents the function's position in the DLL's Export Address Table
 - Structure indices are parallel to those in the Import Address Table (IAT)

- **Import Address Table (IAT)**
 - Also an array of the same structures
 - Initially both the Import Lookup Table and the IAT contain similar entries
 - The loader fills in the addresses of each imported routine in this table
- The Windows loader will proactively load some external DLLs but lazily link symbols
- Depends on how the DLL dependency was declared in the source code & how they were compiled



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PE Shared library calls: Flavor 1 of 2

- Similar to ELF, all external symbol addresses are initialized to a loader function
 - In the case of Windows this is a helper function that hands control to the kernel
 - If the DLL has been proactively loaded, it just links
 - If not, it loads the DLL and then links
- If the original source code declared an included function as external
 - Then the compiler will optimize by just directly reading the address and calling it

```

.text:0000000000001655      mov     r8d, 0
.text:000000000000165B      mov     edx, 1          ; protocol
.text:0000000000001660      mov     ecx, 2
.text:0000000000001665      mov     rax, cs: _imp_socket
.text:000000000000166C      call    rax             ; _imp_socket

; SOCKET __stdcall socket(int af, int type, int protocol)
; extrn __imp_socket:qword ; DATA XREF: main+1051r
; socket ir

```



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PE Shared library calls: Flavor 2 of 2

- If the original source code **did not** declare an included function as external
 - Then the compiler will not know that it should optimize
 - Thus a stumpy function will be made in the .text section
 - No special section like ELF's PLT
- All roads lead to the .idata section eventually

```

.text:00000000000015DA      mov     rax, [rbp+640h+arg_8]
.text:00000000000015E1      mov     rax, [rax]
.text:00000000000015E4      mov     rdx, rax
.text:00000000000015E7      lea     rcx, aYouCanChangeTh ; "[*] You can change this,
.text:00000000000015EE      call    printf

; ===== SUBROUTINE =====
; Attributes: thunk
; int printf(const char *Format, ...)
; public printf
; proc near
; CODE XREF: main+6C1p
; main+8E1p ...
; jmp     cs:_imp_printf
; endp
; printf

```

```

; .idata:00000000000009A4 ; int printf(const char *Format, ...)
; .idata:00000000000009A4      extrn __imp_printf:qword ; DATA XREF: printf1r

```



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Excellent additional reading (Optional)

- Symantec blogger reverse engineered the loading and linking of external symbols on Linux and Windows!
- <https://www.symantec.com/connect/articles/dynamic-linking-linux-and-windows-part-one>
- <https://www.symantec.com/connect/articles/dynamic-linking-linux-and-windows-part-two>



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Lesson Summary

- Interpret high level structures in assembler
- Identify function calls
- Reconstruct switch expressions to index table of targets
- Employ loops and compiler-induced elimination of branching
- Discover arrays
- Identify structures in assembly code



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