

Stochastic Processes

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January 20, 2022

1 Moment Generating Function

- Moment Generating Function: $\mathbb{E}[e^{tX}]$
 - Property:
 - * $\mathbb{E}[e^{tX}] = \int_{-\infty}^{\infty} e^{tx} f_X(x) dx$
 - * $\mathbb{E}[e^{tX}] = \sum_{k=0}^{\infty} E[X^k] \frac{t^k}{k!}$
 - $e^{tx} = \sum_{k=0}^{\infty} \frac{(tx)^k}{k!}$
 - $E[e^{tX}] = E[\sum_{k=0}^{\infty} \frac{(tX)^k}{k!}] = \sum_{k=0}^{\infty} E[X^k] \frac{t^k}{k!}$
 - * $\frac{d\mathbb{E}[e^{tX}]}{dt} = \mathbb{E}[X]$
 - * $\mathbb{E}[e^{t(aX+b)}] = e^{tb} \mathbb{E}[e^{taX}]$
 - * Not all random variables have Moment generating function
 - Characteristic Function: $\mathbb{E}[e^{itX}]$
 - Property:
 - * All random variables have Moment generating function
 - Joint Moment Generating Function: $G(x, y) = \mathbb{E}[e^{xX} e^{yY}]$
 - Property:
 - (Joint) moment generating function uniquely determines the (joint) CDF
 - Example
 - Trapped miner's random walk
 - * Miner has probability of $\frac{1}{3}$ to waste 3 hours in vain, $\frac{1}{3}$ to waste 5 hours in vain, and $\frac{1}{3}$ to spend 2 hours to go out of the mine.
 - * X is the random variables of the hours to go out of the mine
 - * Y_i is the random variables of the hours for the i -th action.
 - * $\mathbb{E}[e^{tX}] = \mathbb{E}[e^{tX} | Y_1 = 2] + \mathbb{E}[e^{tX} | Y_1 = 3] + \mathbb{E}[e^{tX} | Y_1 = 5]$
 - $= \mathbb{E}[e^{2t}] + \mathbb{E}[e^{t(X+3)}] + \mathbb{E}[e^{t(X+5)}]$
 - * Find expectation and variance by joint moment generating function

2 Expectation

- Suppose N is a integer random variable
- Suppose $X_1, \dots, X_i, \dots, X_N$ are i.i.d random variables with mean μ and variance σ^2
- $Y = \sum_{i=1}^N X_i$
- $\mathbb{E}[Y] = \mathbb{E}[N]\mu$
 - $\mathbb{E}[Y] = \sum_{n=1}^{\infty} \mathbb{E}[\sum_{i=1}^N X_i | N = n] P[N = n]$
 - $= \mu \times \sum_{n=1}^{\infty} n P[N = n] = \mathbb{E}[N]\mu$

- $\mathbb{E}[Y^2] = \mathbb{E}[N]\mathbb{E}[X^2] + \mathbb{E}[N^2]\mu^2 - \mathbb{E}[N]\mu^2$
 - $\mathbb{E}[Y^2] = \sum_{n=1}^{\infty} \mathbb{E}[(\sum_{i=1}^N X_i)^2 | N = n] P[N = n] = \sum_{n=1}^{\infty} (n\mathbb{E}[X_i^2] + n(n-1)\mu^2) P[N = n]$
 $= \mathbb{E}[N]\mathbb{E}[X^2] + \mathbb{E}[N^2]\mu^2 - \mathbb{E}[N]\mu^2$
- $Var(Y) = \mathbb{E}[N]\sigma^2 + Var(N)\mu^2$

3 Inequality

- Markov Inequality

Definition:

- Suppose $X \geq 0$, then $P[X \geq \epsilon] \leq \frac{\mathbb{E}[X]}{\epsilon}$

Proof:

1. $\mathbb{E}[X] = \int_0^{\infty} x f_X(x) dx \geq \int_{\epsilon}^{\infty} x f_X(x) dx \geq \epsilon \int_{\epsilon}^{\infty} f_X(x) dx = \epsilon P[X \geq \epsilon]$
2. $X(\omega) \geq \epsilon \mathbb{1}_{X(\omega) \geq \epsilon}, \forall \omega \in S$
 - Calculate expectation on both side.
 - $\mathbb{E}[X] \geq \epsilon P[X \geq \epsilon]$

Property:

- The equality happens when $P[X = k] = 0, \forall k \notin \{0, \epsilon\}$.

- Chebyshev Inequality

Definition:

- Suppose $m = \mathbb{E}[X], \sigma^2 = Var(X)$, then $P[|X - m| \geq \epsilon] \leq \frac{\sigma^2}{\epsilon^2}$

Proof:

- $P[|X - m| \geq \epsilon] = P[(X - m)^2 \geq \epsilon^2]$
- $P[(X - m)^2 \geq \epsilon^2] \leq \frac{\mathbb{E}[(X - m)^2]}{\epsilon^2}$ (by Markov Inequality)

Property:

- The equality happens when $P[X = k] = 0, \forall k \notin \{m - \epsilon, m, m + \epsilon\}$.
- Might be tighter than Markov Inequality since it requires m, σ

- Chernoff Inequality

Definition:

- Suppose X_1, \dots, X_n are independent identically distributed Bernoulli random variable with probability p and $X = \sum_{i=1}^n X_i$
- $P[X \geq \epsilon] \leq \frac{(pe^t + 1 - p)^n}{e^{t\epsilon}} \leq \frac{e^{np(e^t - 1)}}{e^{t\epsilon}}$
 - * $P[X \geq \epsilon] = P[e^{tX} \geq e^{t\epsilon}] \leq \frac{\mathbb{E}[e^{tX}]}{e^{t\epsilon}} = \frac{(E[e^{tX_i}])^n}{e^{t\epsilon}} = \frac{(pe^t + 1 - p)^n}{e^{t\epsilon}} \leq \frac{e^{np(e^t - 1)}}{e^{t\epsilon}}$
- $P[X \geq np(1 + \epsilon)] \leq (\frac{e^{\epsilon}}{(1 + \epsilon)^{1 + \epsilon}})^{np} \leq \begin{cases} e^{\frac{-\epsilon^2 np}{3}} & \text{if } 0 \leq \epsilon \leq 1 \\ e^{\frac{-\epsilon^2 np}{(2 + \epsilon)}} & \text{if } \epsilon \geq 1 \end{cases}$
 - * Substitute ϵ with $np(1 + \epsilon)$
 - * Substitute t with $\log(1 + \epsilon)$
 - * the last inequality is without proof
- $P[X \leq \epsilon] \leq \frac{(pe^{-t} + 1 - p)^n}{e^{-t\epsilon}} \leq \frac{e^{np(e^{-t} - 1)}}{e^{-t\epsilon}}$
 - * $P[X \leq \epsilon] = P[e^{-tX} \geq e^{-t\epsilon}] \leq \frac{\mathbb{E}[e^{-tX}]}{e^{-t\epsilon}} = \frac{(E[e^{-tX_i}])^n}{e^{-t\epsilon}} = \frac{(pe^{-t} + 1 - p)^n}{e^{-t\epsilon}} \leq \frac{e^{np(e^{-t} - 1)}}{e^{-t\epsilon}}$
- $P[X \leq np(1 - \epsilon)] \leq (\frac{e^{-\epsilon}}{(1 - \epsilon)^{1 - \epsilon}})^{np} \leq e^{\frac{-\epsilon^2 np}{2}}$

- * Substitute ϵ with $np(1 - \epsilon)$
- * Substitute t with $-\log(1 - \epsilon)$
- * the last inequality is without proof

- Chernoff/ Hoeffding Lemma

Definition:

- Suppose X_1, \dots, X_n are independent distributed random variable and $a_i \leq X_i \leq b_i$
- Suppose $X = \sum_{i=1}^n X_i$ and $\mu = \mathbb{E}[X]$
- $P[|X - \mu| \geq \epsilon] \leq 2e^{\frac{-2\epsilon^2}{\sum_{i=1}^n (b_i - a_i)^2}}$ without proof

- Application:

- Balls in Bins

Definition: Throw n balls into n bins, find bounds for the maximum number of balls in all bins

- * $P[\text{maximum number of balls in all bins} \geq \epsilon]$
 $= P[\cup_{i=1}^n \text{number of balls in } i\text{-th bin} \geq \epsilon]$
 $\leq n \times P[\text{number of balls in one bin} \geq \epsilon]$
- * By Markov inequality:
 $\cdot P[\text{number of balls in one bin} \geq \epsilon] \leq \frac{1}{\epsilon} \rightarrow \text{useless}$
- * By Chebyshev inequality:
 $\cdot P[\text{number of balls in one bin} \geq \epsilon] \leq \frac{(1 - \frac{1}{n})}{\epsilon^2}$
 $\cdot P[\text{maximum number of balls in all bins} \geq n^{\frac{1}{2} + \epsilon}] \leq \frac{(1 - \frac{1}{n})}{n^{2\epsilon}}$
 $\cdot \text{when } n \rightarrow \infty, \text{ the maximum number of balls should less than } n^{\frac{1}{2} + \epsilon}$
- * By Chernoff inequality:
 $\cdot P[\text{number of balls in one bin} \geq 2 \log n] \leq \frac{e^{np(\epsilon^t - 1)}}{n^{2t}}$
 $\cdot P[\text{maximum number of balls in all bins} \geq 2 \log n] \leq \frac{e^{np(\epsilon^t - 1)}}{n^{2t-1}}$
 $\cdot \text{when } t \text{ is a constant } \geq 0.5 \text{ and } n \rightarrow \infty, \text{ the maximum number of balls should less than } 2 \log n$

4 Law of Large Numbers

- $\{X_i\}_{i=1}^\infty$ is a sequence of pairwise uncorrelated random variable with
 $\mathbb{E}[X_i] = m, \text{Var}(X_i) = \sigma_i^2$.
- $M_n = \frac{1}{n} \sum_{i=1}^n X_i$
- $M_n \rightarrow m$ almost surely, in mean square and in probability.