Project III Report for COM S 4/5720 Spring 2025: Risk-Aware Stochastic Planner for Three-Agent Pursuit-Evasion

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I. PROBLEM RESTATEMENT

Project III preserves the grid-world and three-agent rules of Project III but introduces move probabilities. When an agent issues a move $a \in \mathcal{A}$, the environment executes left(a) with probability p_1 , a with probability p_2 , and right(a) with probability p_3 , where $(p_1, p_2, p_3) = (0.4, 0.3, 0.3)$. The planner must therefore:

- maximise capture probability,
- · minimise expected collision risk,
- satisfy a strict real-time (single-step) budget.

II. DETAILED PLANNER LOGIC

TABLE I
SYMBOL DICTIONARY (CODE CONSTANTS IN PARENTHESES).

w_t	target weight = 1.0 ('W_PREY')
w_p	pursuer weight = 1.2 ('W_PURS')
c_r	$risk cost = 40 ('RISK_COST')$
w_s	mobility weight $= 1$ (implicit)
d_t	Manhattan dist. to prey ('mhd(cur,prey)')
d_p	Manhattan dist. to pursuer ('mhd(cur,purs)')
\hat{E}	# empty 4-neighbours around current cell
stayBias	-0.04 ('STAY_BIAS')
$ au_0$	base risk gate = 0.12 ('RISK_TH_BASE')
$ au_{ m max}$	hard cap = 0.28 ('RISK_TH_MAX')
k	adaptive slope = 0.03 ('ADAPT_RATE')
N	stall limit 10 ('STALL_LIMIT')

Project III introduces **actuation noise**: every issued move is executed *left*, *straight*, or *right* with probabilities $(p_1, p_2, p_3) = (0.4, 0.3, 0.3)$. We therefore extend the deterministic planner from Project II with *risk-aware* scoring and a shallow expectimax look-ahead.

A. Risk Tensor and Expected-Value Scoring

Following the risk-sensitive MDP [1], we pre-compute

$$\mathcal{R}[r,c,k] = \Pr[\text{collision} \mid \text{issue move } k \text{ at } (r,c)]$$

by checking the three rotated endpoints once per map. At run time each action receives the expected-utility score [2]:

$$\mathrm{EV}(k) = w_t \, \mathbb{E}\Big[\tfrac{1}{d_t+1}\Big] - w_p \, \mathbb{E}\Big[\tfrac{1}{d_p+1}\Big] - c_r \, \mathcal{R}[r,c,k] + w_s \, E + \mathrm{stayBias},$$

where the expectation is taken over (p_1, p_2, p_3) .

B. Step-wise Decision Pipeline

- 1) **Dynamic risk gate.** A tolerance $\tau = \min(\tau_0 + k \cdot idleSteps, \tau_{\max})$ filters actions with $\mathcal{R} \leq \tau$ (idea adapted from low-risk planning [3]).
- 2) **Capture lunge.** Issue the safest action k^* immediately if $\Pr[\text{capture } | k^*] > \mathcal{R}[r, c, k^*]$.
- 3) **Emergency flee.** When the pursuer is within 3 Manhattan steps, choose the safe move that maximises $2d_p d_t$ (heuristic used by [4]).
- 4) **Risk-bounded A* search.** Run A* [5] on the subgraph $\{(r,c) \mid \mathcal{R} \leq \tau\}$.
- 5) **Stall breaker.** After N stagnant frames, pick the move that minimises the prey's escape exits [6].
- 6) Lightweight expectimax (1.5-ply). We build an expectimax tree of depth 1 and then re-evaluate only the three best root actions. Because the second level is explored only partially, the effective look-ahead is called "1.5-ply". Expectimax is a relatively popular algorithm which branches off of the Minimax algorithm, but assumes non-optimal play. This is good since we know there is 60% chance for incorrect movement.

REFERENCES

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