

CS2105

An $\text{\AA}\omega\epsilon\sigma\omicron\mu\epsilon$ Introduction to Computer Networks

Lectures 6&7: The Network Layer



Department of Computer Science
School of Computing

Lectures 6&7: The Network Layer

After this class, you are expected to understand:

- ❖ the basic services network layer provides.
- ❖ the purpose of DHCP and how it works.
- ❖ IP address, subnet, subnet mask and address allocation.
- ❖ how longest prefix forwarding in a router works.
- ❖ the purpose of routing protocols on the Internet.
- ❖ the principle of Bellman-Ford equation.
- ❖ the workings of distance vector algorithm.
- ❖ the purpose of NAT and how it works.
- ❖ the Internet Protocol (IP) and how datagram fragmentation works.

Lectures 6&7: Roadmap

4.1 Overview of Network Layer

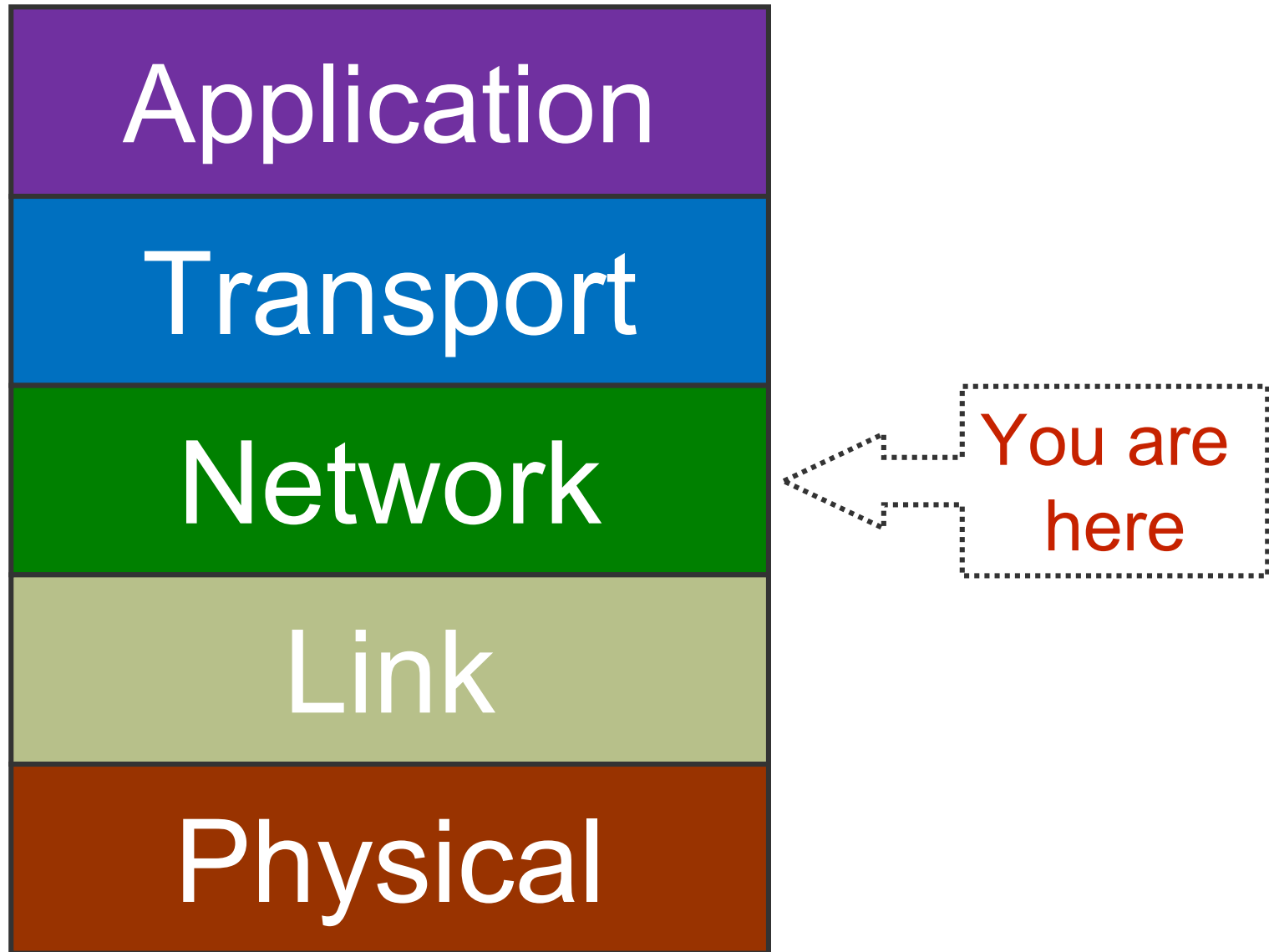
4.2 What's Inside a Router

4.3 The Internet Protocol (IP)

5.2 Routing Algorithms

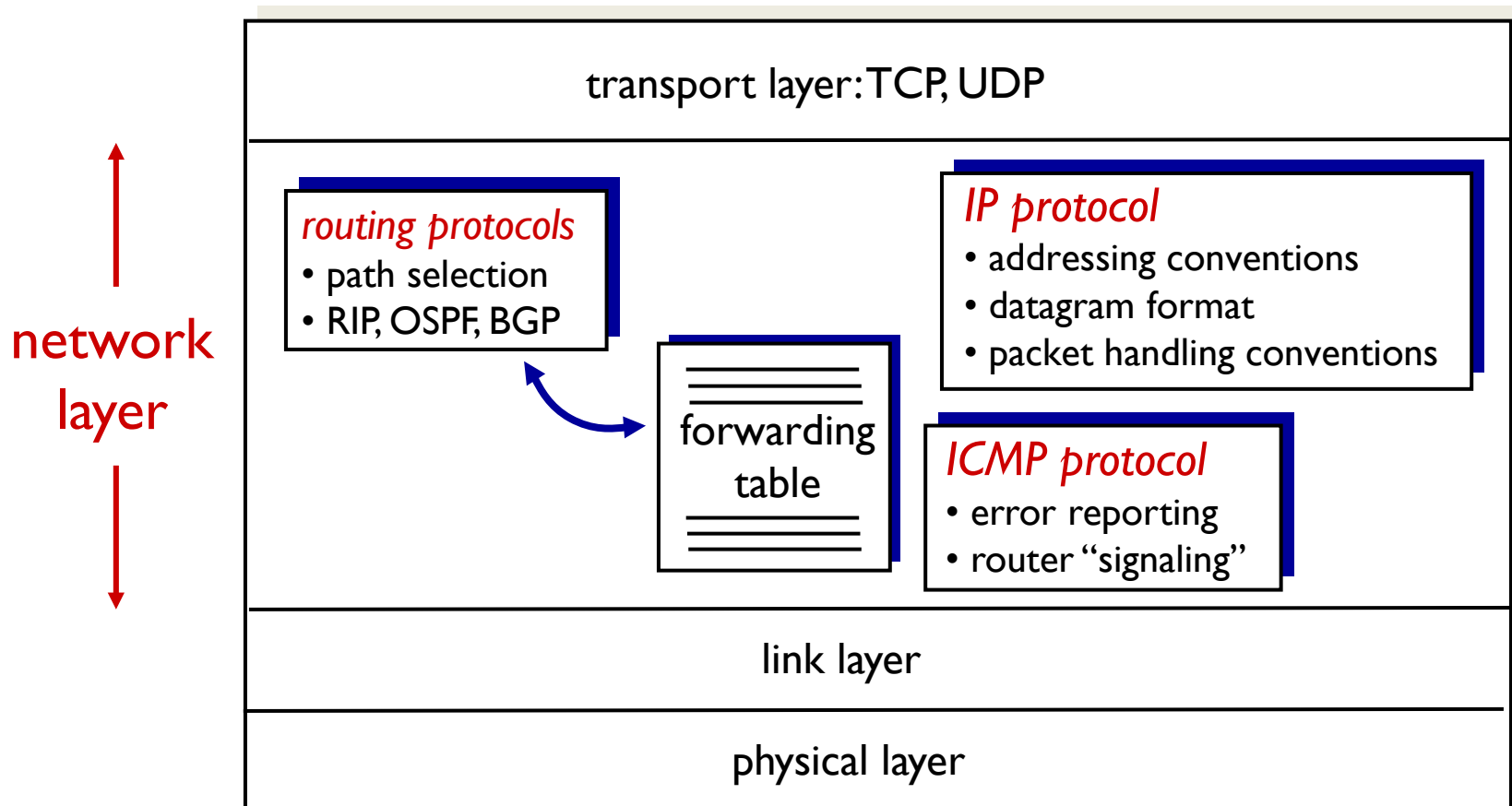
5.6 ICMP

Kurose Textbook, Chapters 4&5
(Some slides are taken from the book)



Network Layer Services

- ❖ Network layer delivers packets to receiving hosts.
 - Routers examine header fields of IP datagrams passing it.



Lectures 6&7: Roadmap

4.1 Overview of Network Layer

4.2 What's Inside a Router

- 4.2.1 Destination-Based Forwarding

4.3 The Internet Protocol (IP)





- 4.3.3 IPv4 Addressing

5.2 Routing Algorithms

5.6 ICMP

IP Address

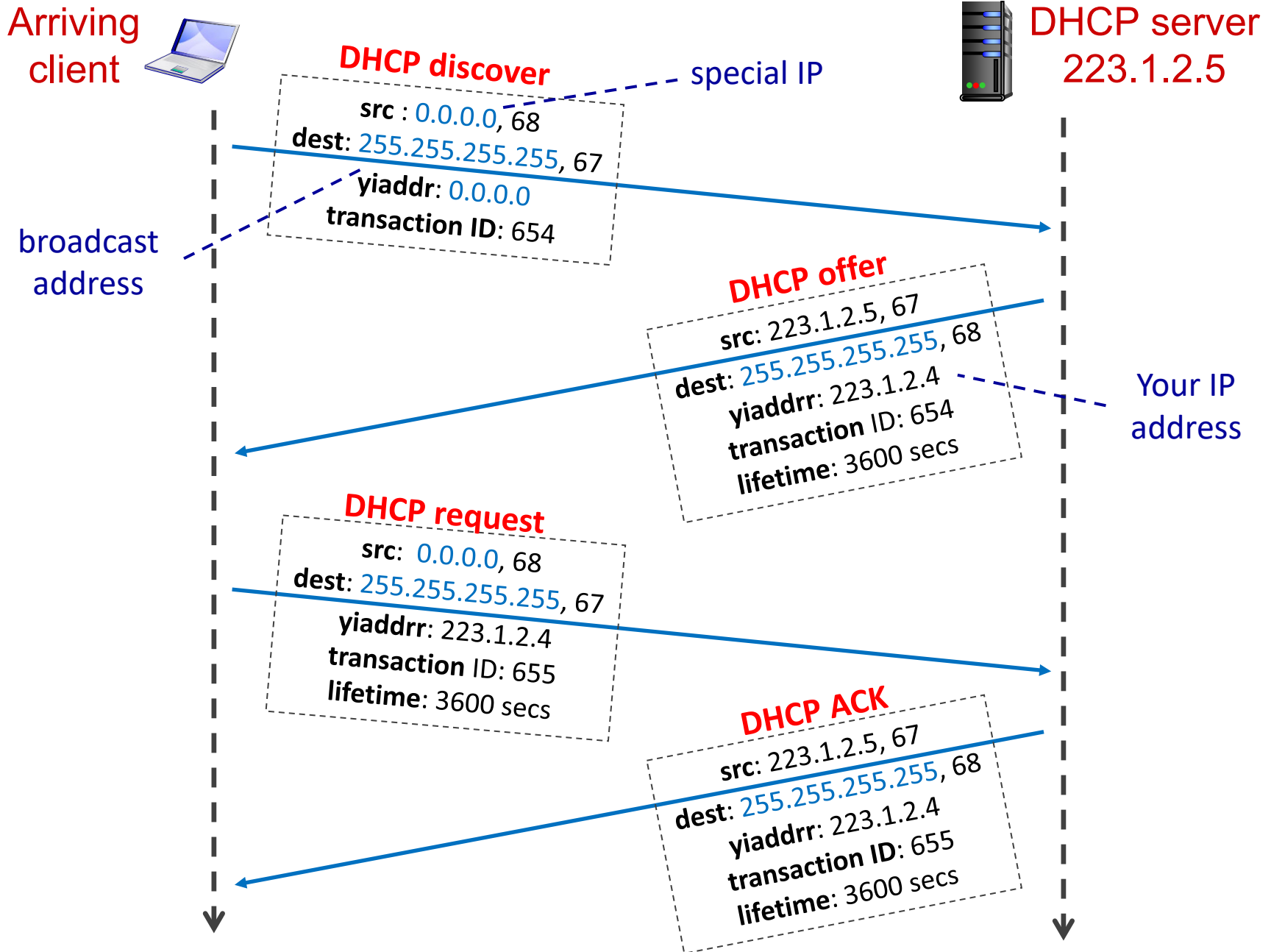
- ❖ **IP address** is used to identify a host (or a router).
 - A 32-bit integer expressed in either binary or decimal

Binary:	00000001	00000010	00000011	10000001
				
Decimal:	1	2	3	129

- ❖ How does a host get an IP address?
 - manually configured by system administrator, or
 - automatically assigned by a **DHCP** (**Dynamic Host Configuration Protocol**) server.

Dynamic Host Configuration Protocol

- ❖ **DHCP** allows a host to dynamically obtain its IP address from DHCP server when it joins network.
 - IP address is renewable
 - allow reuse of addresses (only hold address while connected)
 - support mobile users who want to join network.
- ❖ **DHCP**: 4-step process:
 - 1) Host broadcasts “DHCP discover” message
 - 2) DHCP server responds with “DHCP offer” message
 - 3) Host requests IP address: “DHCP request” message
 - 4) DHCP server sends address: “DHCP ACK” message



More on DHCP

- ❖ In addition to host IP address assignment, DHCP may also provide a host additional network information:
 - IP address of first-hop router
 - IP address of local DNS server
 - Network mask (indicating network prefix versus host ID of an IP address)

- ❖ DHCP runs over UDP
 - DHCP server port number: 67
 - DHCP client port number: 68

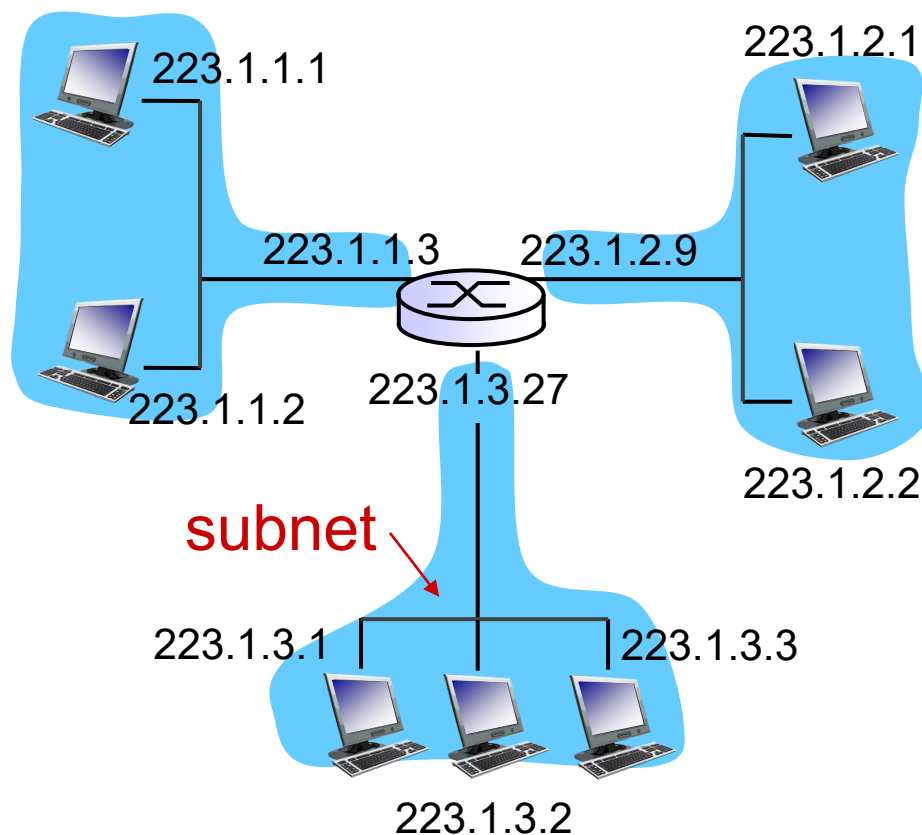
Some Special IP Addresses

Special Addresses	Present Use
0.0.0.0/8	Non-routable meta-address for special use
127.0.0.0/8	Loopback address. A datagram sent to an address within this block loops back inside the host. This is ordinarily implemented using only 127.0.0.1/32.
10.0.0.0/8 172.16.0.0/12 192.168.0.0/16	Private addresses, can be used without any coordination with IANA or an Internet registry.
255.255.255.255/32	Broadcast address. All hosts on the same subnet receive a datagram with such a destination address.

The full list of special IP addresses can be found in RFC5735:
<https://tools.ietf.org/rfc/rfc5735.txt>

IP Address and Network Interface

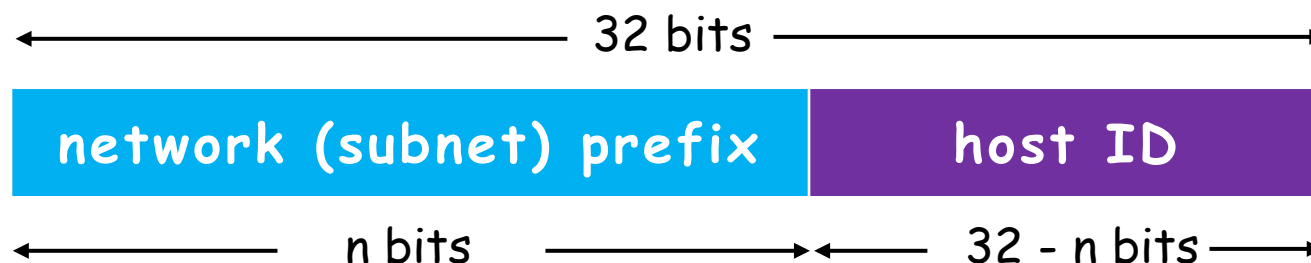
- ❖ An IP address is associated with a **network interface**.
 - A host usually has one or two network interfaces (e.g. wired Ethernet and WiFi).
 - A router typically has multiple interfaces.



A network consisting of 3 subnets
(first 24 bits of IP addr. are network prefix)

IP Address and Subnet

- ❖ An IP address logically comprises two parts:



- ❖ **Subnet** is a network formed by a group of “directly” interconnected hosts.
 - Hosts in the same subnet have the same network prefix of IP address.
 - Hosts in the same subnet can physically reach each other without intervening router.
 - They connect to the outside world through a router.

IP Address: CIDR

- ❖ The Internet's IP address assignment strategy is known as **Classless Inter-domain Routing (CIDR)**.
 - Subnet prefix of IP address is of arbitrary length.
 - Address format: **a.b.c.d/x**, where **x** is the number of bits in subnet prefix of IP address.

← subnet prefix → host ID →
 11001000 00010111 00010000 00101010

this subnet contains 2^9 IP addresses
 subnet prefix: 200.23.16.42/23

/23 indicates the no. of
 bits of subnet prefix

Subnet Mask

- ❖ **Subnet mask** is used to determine which subnet an IP address belongs to.
 - made by setting all subnet prefix bits to "1"s and host ID bits to "0"s.
- ❖ Example: for IP address 200.23.16.42/23:

	← subnet prefix →		← host ID →	
IP address in binary	11001000	00010111	00010000	00101010
Subnet mask	11111111	11111111	11111110	00000000
Subnet mask in decimal	255.255.254.0			

Quiz

- ❖ For the following 4 IP addresses, which one is in a different subnet from the rest 3?
- a. 172.26.185.128/26
 - b. 172.26.185.130/26
 - c. 172.26.185.160/26
 - d. 172.26.185.192/26

IP Address Allocation

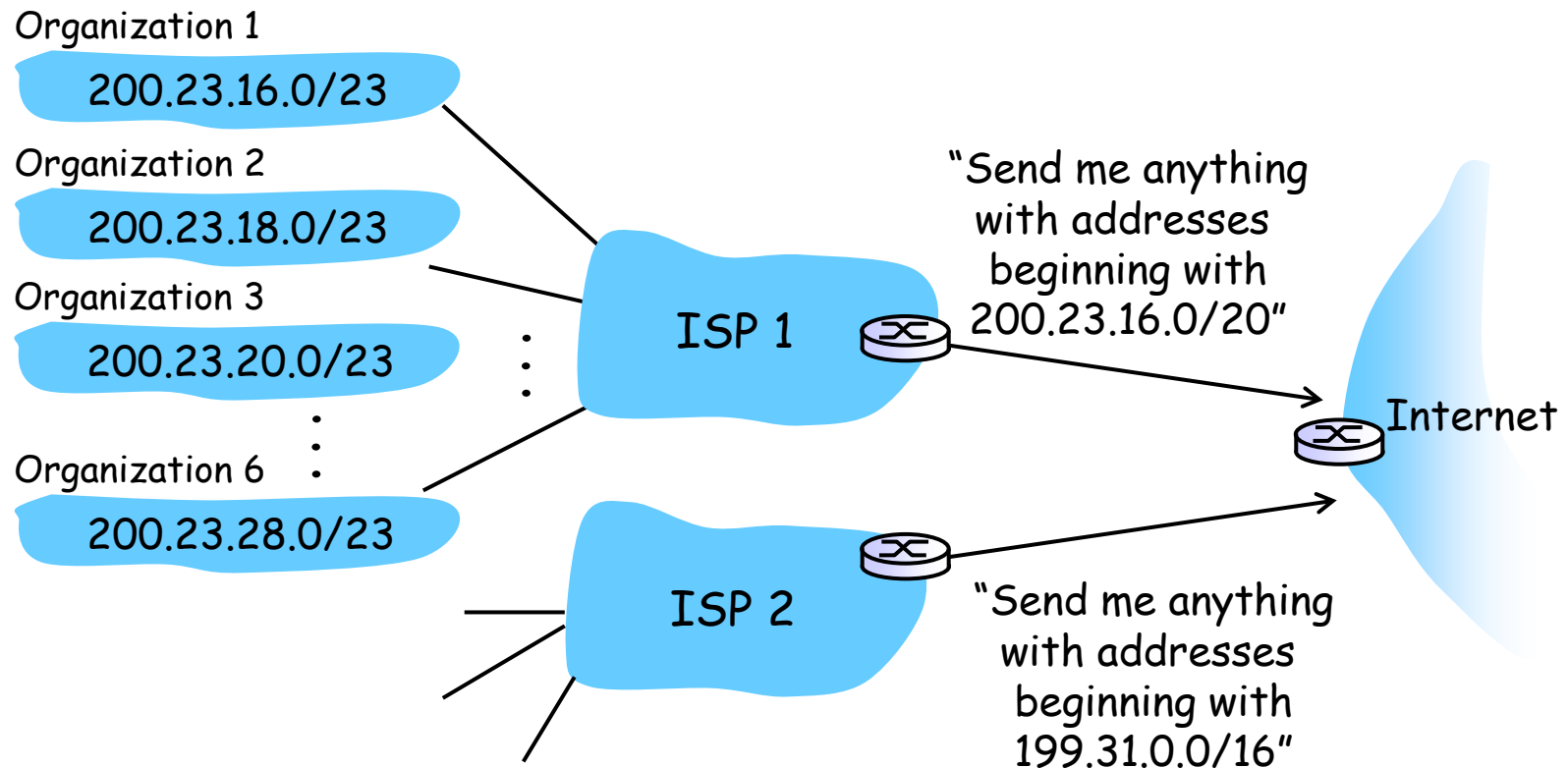
- ❖ **Q:** How does an organization obtain a block of IP addresses?
- ❖ **A:** Buy from registry or rent from ISP's address space.

	Binary Address	Decimal Address
ISP's block	11001000 00010111 0001 000 0 00000000	200.23.16.0/20
Organization 1	11001000 00010111 0001 000 0 00000000	200.23.16.0/23
Organization 2	11001000 00010111 0001 001 0 00000000	200.23.18.0/23
Organization 3	11001000 00010111 0001 010 0 00000000	200.23.20.0/23
...
Organization 6	11001000 00010111 0001 101 0 00000000	200.23.28.0/23

use 3 more bits to differentiate
6 organizations

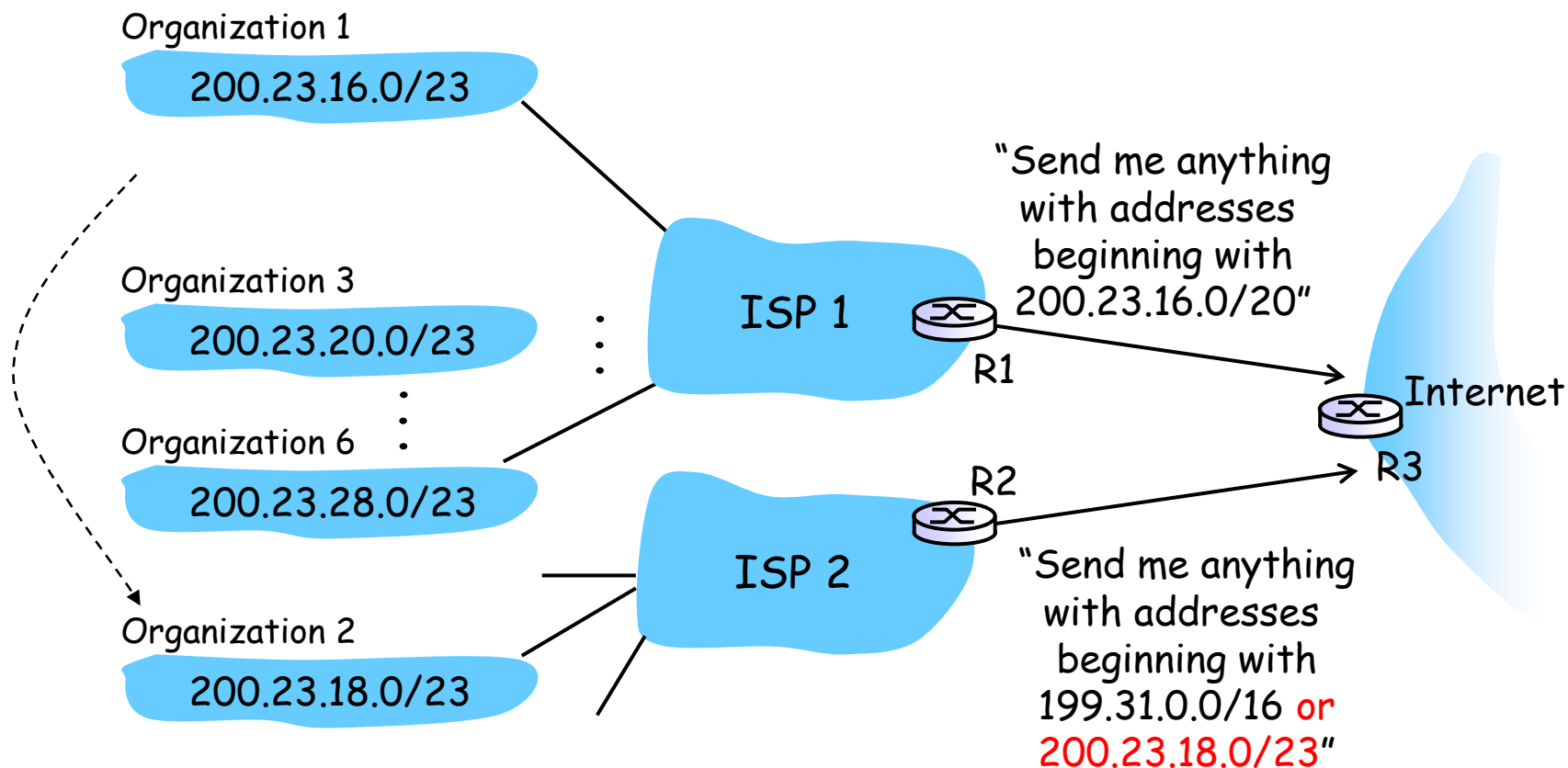
Hierarchical Addressing

Hierarchical addressing allows efficient advertisement of routing information:



Hierarchical Addressing

Suppose Organization 2 now switches to ISP 2, but doesn't want to renumber all of its routers and hosts.

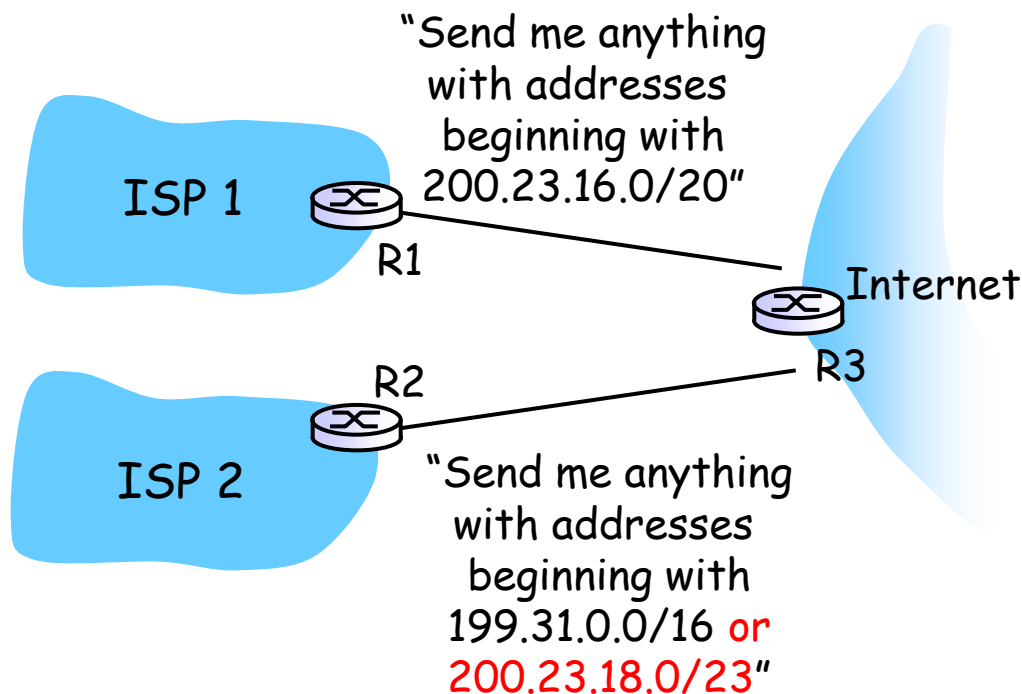


Longest Prefix Match (1/2)

- ❖ **Question:** which router to deliver to,
- if a packet has destination IP **200.23.20.2**?
 - if a packet has destination IP **200.23.19.3**?

Forwarding Table at R3

Net mask	Next hop
200.23.16.0/20	R1
200.23.18.0/23	R2
199.31.0.0/16	R2
...	...



Longest Prefix Match (2/2)

- ❖ Packet with destination IP **200.23.20.2** ⇒ R1
 - (Binary: **11001000** **00010111** **00010100** 00000010)
- ❖ Packet with destination IP **200.23.19.3** ⇒ R2
 - (Binary: **11001000** **00010111** **00010011** 00000011)

Forwarding Table at R3

Net mask	Net mask in binary	Next hop
200.23.16.0/20	11001000 00010111 00010000 00000000	R1
200.23.18.0/23	11001000 00010111 00010010 00000000	R2
199.31.0.0/16	11000111 00011111 00000000 00000000	R2
...		...

match the
longest prefix

More on IP Address Allocation

- ❖ **Q1:** How does an organization obtain a block of IP addresses?
- ❖ **A1:** Buy from registry or rent from ISP's address space.
- ❖ **Q2:** How does an ISP get a block of addresses?
- ❖ **A2:** ICANN: Internet Corporation for Assigned Names and Numbers
 - Allocates addresses
 - Manages DNS
 - Assigns domain names, resolves disputes

Lectures 6&7: Roadmap

4.1 Overview of Network Layer

4.2 What's Inside a Router

4.3 The Internet Protocol (IP)

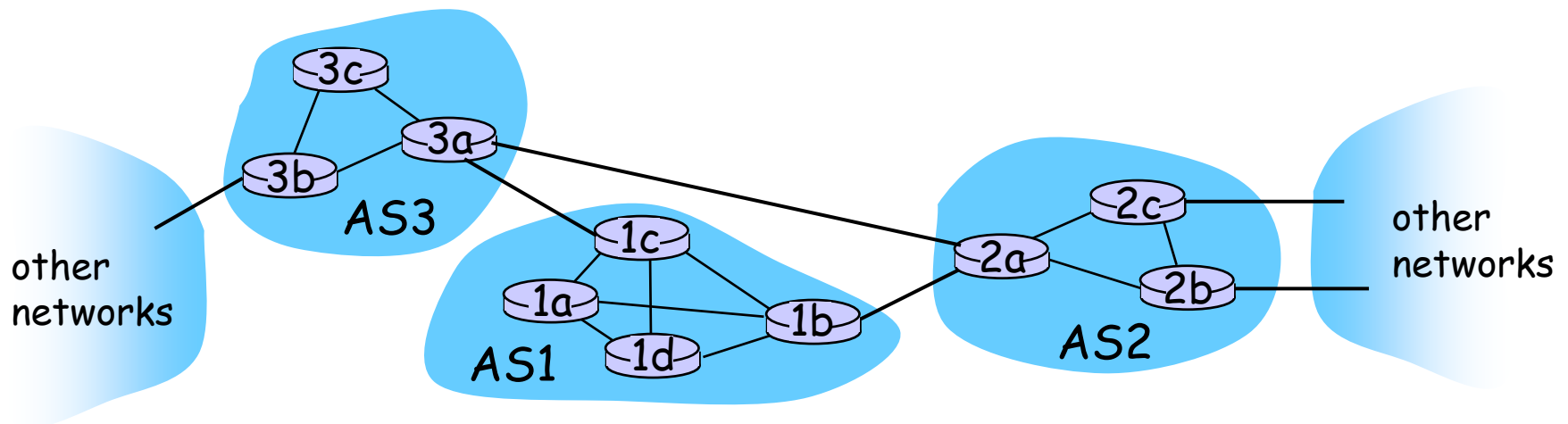
5.2 Routing Algorithms

- 5.2.2 The Distance Vector Routing Algorithm

5.6 ICMP

Internet: Network of Networks

- ❖ The Internet is a “network-of-networks”.
 - A hierarchy of Autonomous Systems (AS), e.g., ISPs, each owns routers and links.
- ❖ Due to the size of the Internet and the decentralized administration of the Internet, routing on the Internet is done hierarchically.



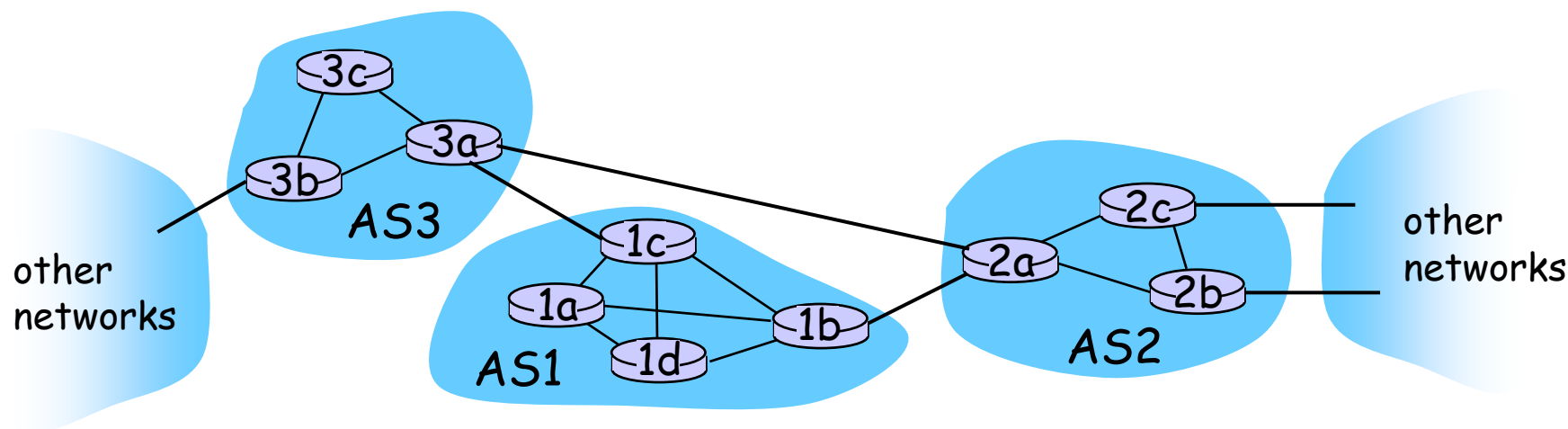
Routing in The Internet

❖ Intra-AS routing

- Finds a good path between two routers within an AS.
- Commonly used protocols: **RIP**, **OSPF**

❖ Inter-AS routing (not covered)

- Handles the interfaces between ASs.
- The de facto standard protocol: **BGP**



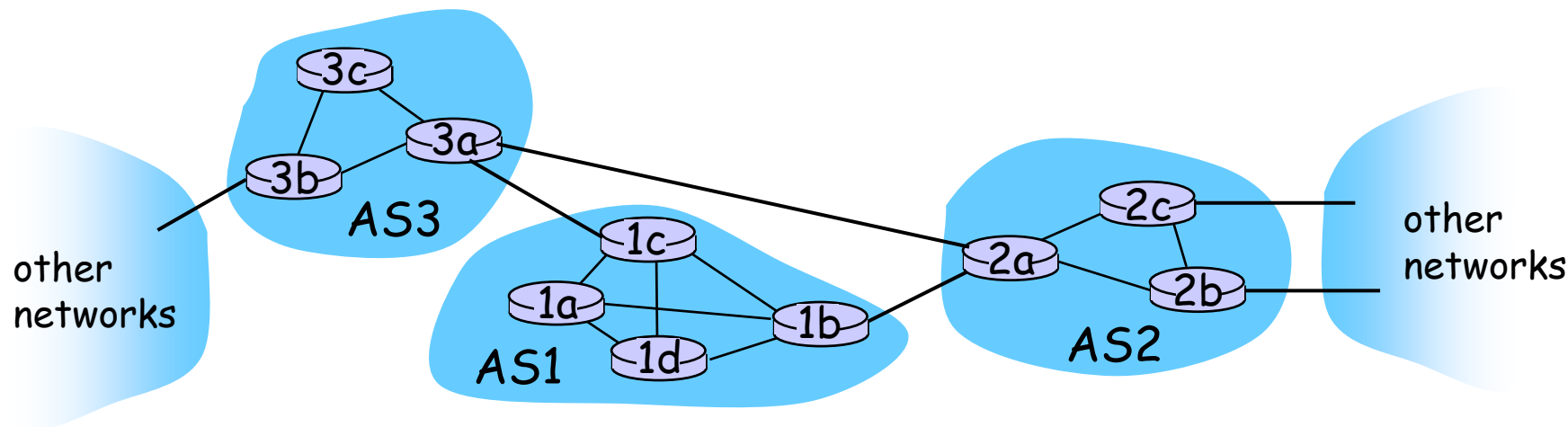
Routing in The Internet

❖ Intra-AS routing

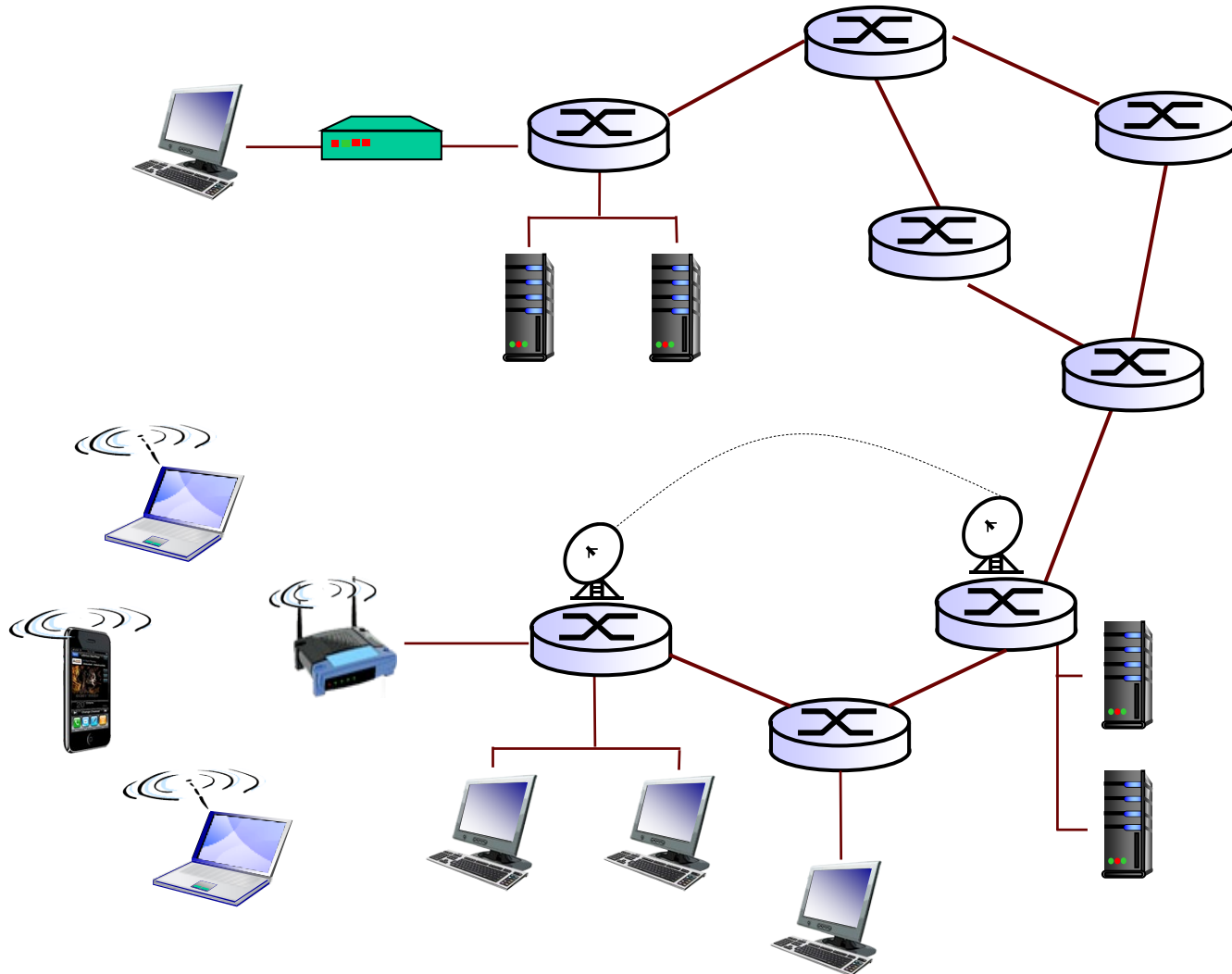
- Single admin, so no policy decisions are needed.
- Routing mostly focus on performance.

❖ Inter-AS routing (not covered)

- Admin often wants to control over how its traffic is routed, who routes through its net, etc.
- Policy may dominate over performance.

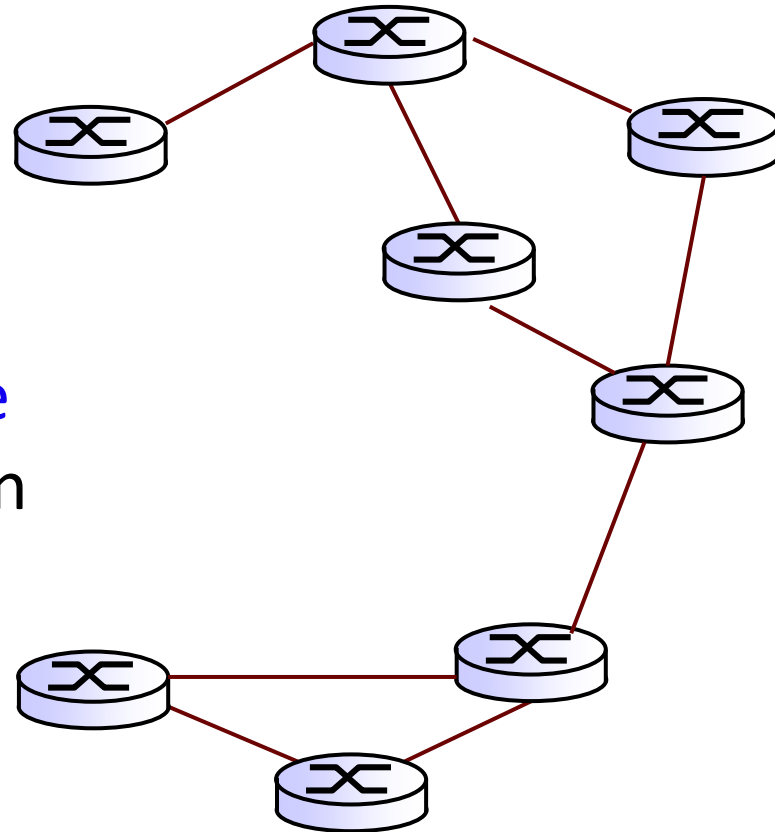


Abstract View of Intra-AS Routing



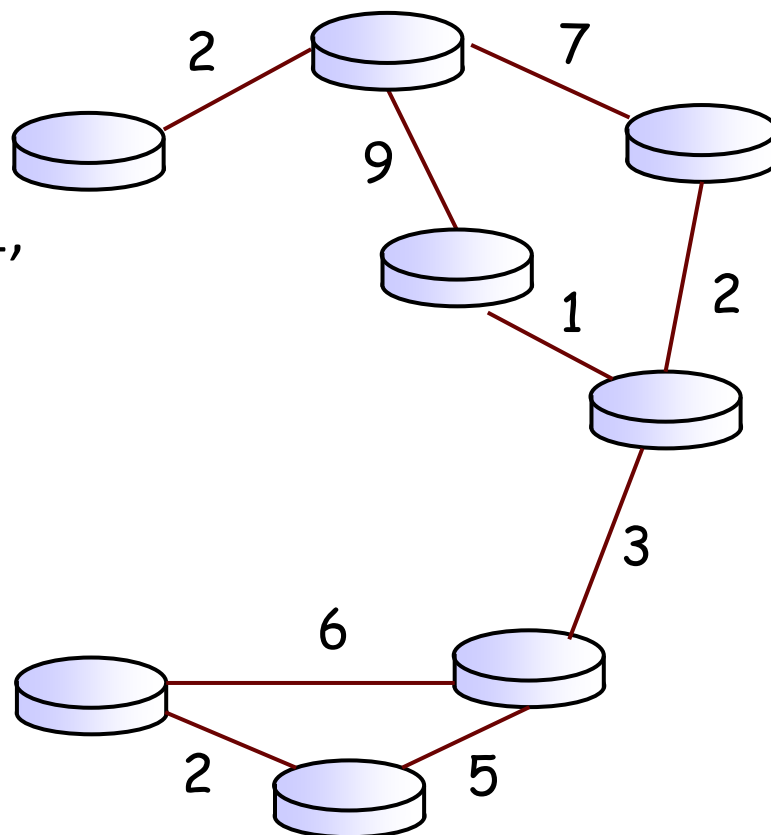
Abstract View of Intra-AS Routing

- ❖ We can abstractly view a network of routers as a **graph**, where **vertices are routers** and **edges are physical links** between routers.



Abstract View of Intra-AS Routing

- ❖ We can associate a **cost** to each link.
 - cost could always be 1, or inversely related to bandwidth, or related to congestion.

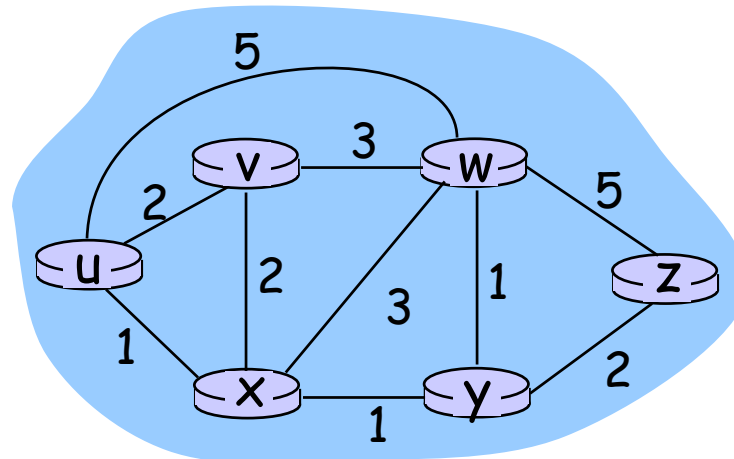


Routing: finding a
least cost path
between two
vertices in a graph

Routing Algorithms Classification

“link state” algorithms

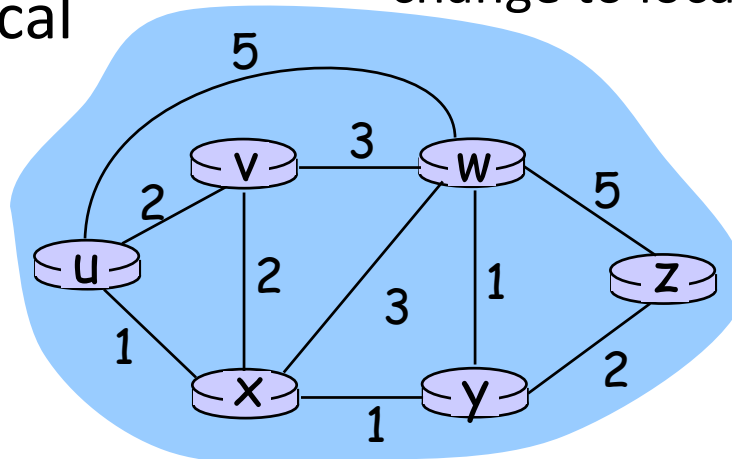
- ❖ All routers have the complete knowledge of network topology and link cost.
 - Routers periodically broadcast link costs to each other.
- ❖ Use Dijkstra algorithm to compute least cost path locally (using global map).
- ❖ Non-examinable 😊



Routing Algorithms Classification

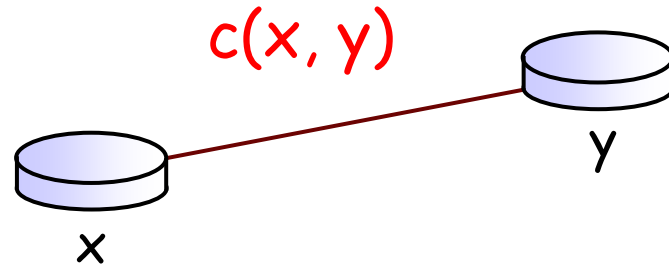
“distance vector” algorithms

- ❖ Routers know physically-connected neighbors and link costs to neighbors.
- ❖ Routers exchange “local views” with neighbors and update own “local views” (based on neighbors’ view).
- ❖ Iterative process of computation
 1. Swap local view with direct neighbours.
 2. Update own’s local view.
 3. Repeat 1 - 2 till no more change to local view.

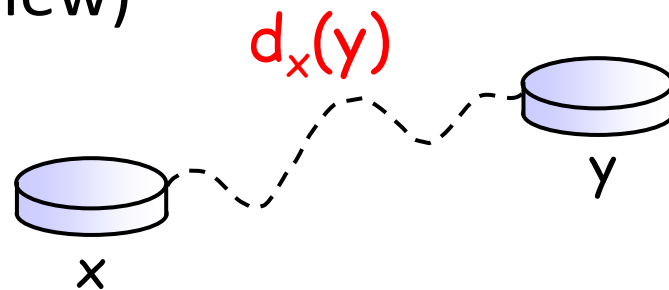


Some Graph Notations

- ❖ $c(x, y)$: the cost of link between routers x and y
 - $= \infty$ if x and y are not direct neighbours



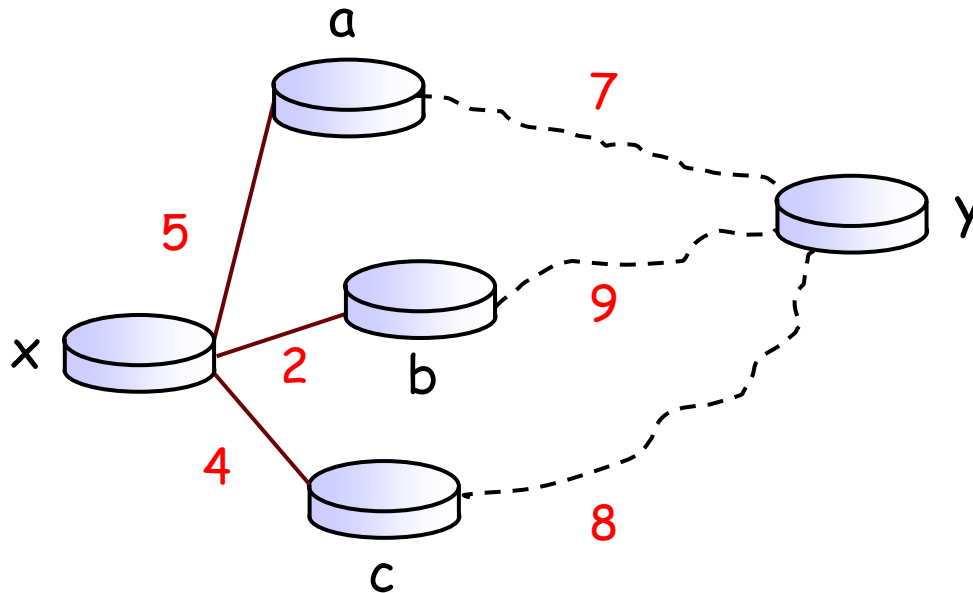
- ❖ $d_x(y)$: the cost of the least-cost path from x to y (from x 's view)



Bellman-Ford Equation

$$d_x(y) = \min_v \{c(x, v) + d_v(y)\}$$

where min is taken over all direct neighbors v of x

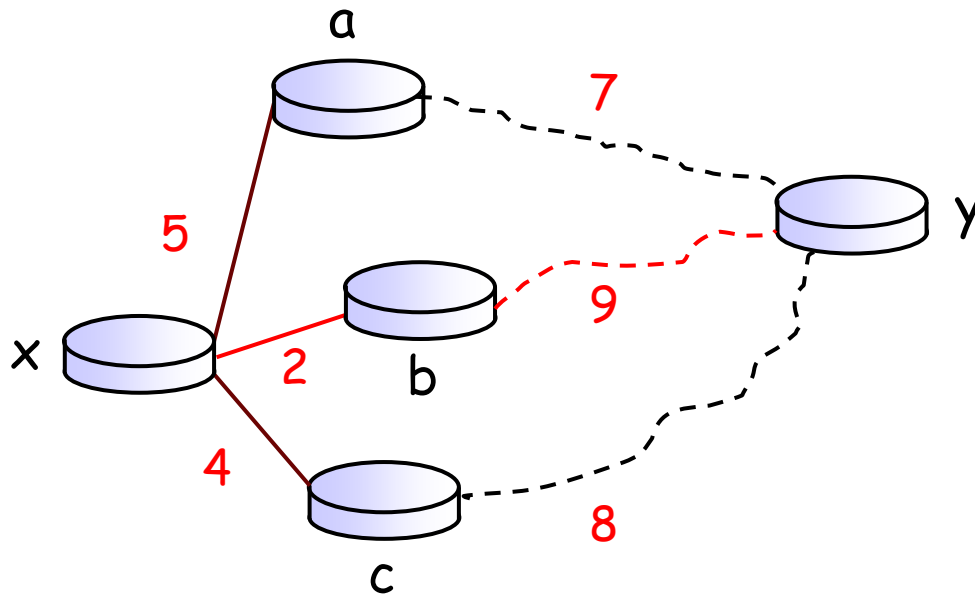


$$d_x(y) = \min_v \{ c(x, a) + d_a(y), \\ c(x, b) + d_b(y), \\ c(x, c) + d_c(y) \}$$

$$= \min \{12, 11, 12\} = 11$$

Bellman-Ford Equation

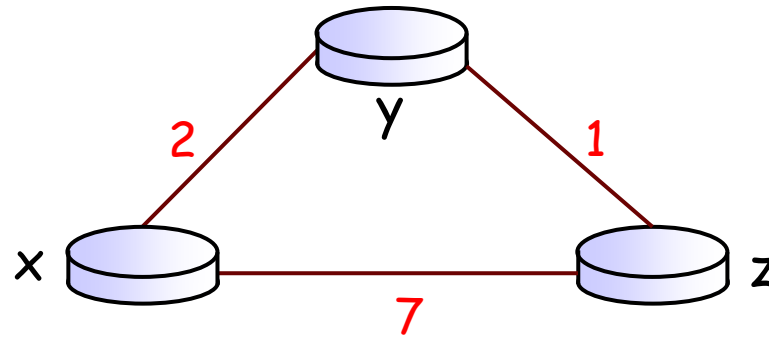
- ❖ To find the least cost path, x needs to know the cost from each of its direct neighbour to y .
- ❖ Each neighbour v sends its **distance vector** (y, k) to x , telling x that the cost from v to y is k .



Now x knows, to reach y , packet should be forward to b and the total cost would be 11.

Bellman-Ford Example

$$d_x(y) = \min_v \{c(x, v) + d_v(y)\}$$



cost to

	x	y	z
x			
y			
z			

x' view

cost to

	x	y	z
x			
y			
z			

y' view

cost to

	x	y	z
x			
y			
z			

z' view

Distance Vector Algorithm

- ❖ Every router, \mathbf{x} , \mathbf{y} , \mathbf{z} , sends its distance vectors to its directly connected neighbors.
- ❖ When \mathbf{x} finds out that \mathbf{y} is advertising a path to \mathbf{z} that is cheaper than \mathbf{x} currently knows,
 - \mathbf{x} will update its distance vector to \mathbf{z} accordingly.
 - In addition, \mathbf{x} will note down that all packets for \mathbf{z} should be sent to \mathbf{y} . This info will be used to create forwarding table of \mathbf{x} .
- ❖ After every router has exchanged several rounds of updates with its direct neighbors, all routers will know the least-cost paths to all the other routers.

RIP

- ❖ **RIP (Routing Information Protocol)** implements the DV algorithm. It uses **hop count** as the cost metric (i.e., insensitive to network congestion).
- ❖ Exchange routing table every 30 seconds over UDP port 520.
- ❖ “Self-repair”: if no update from a neighbour router for 3 minutes, assume neighbour has failed.

Lectures 6&7: Roadmap

4.1 Overview of Network Layer

4.2 What's Inside a Router

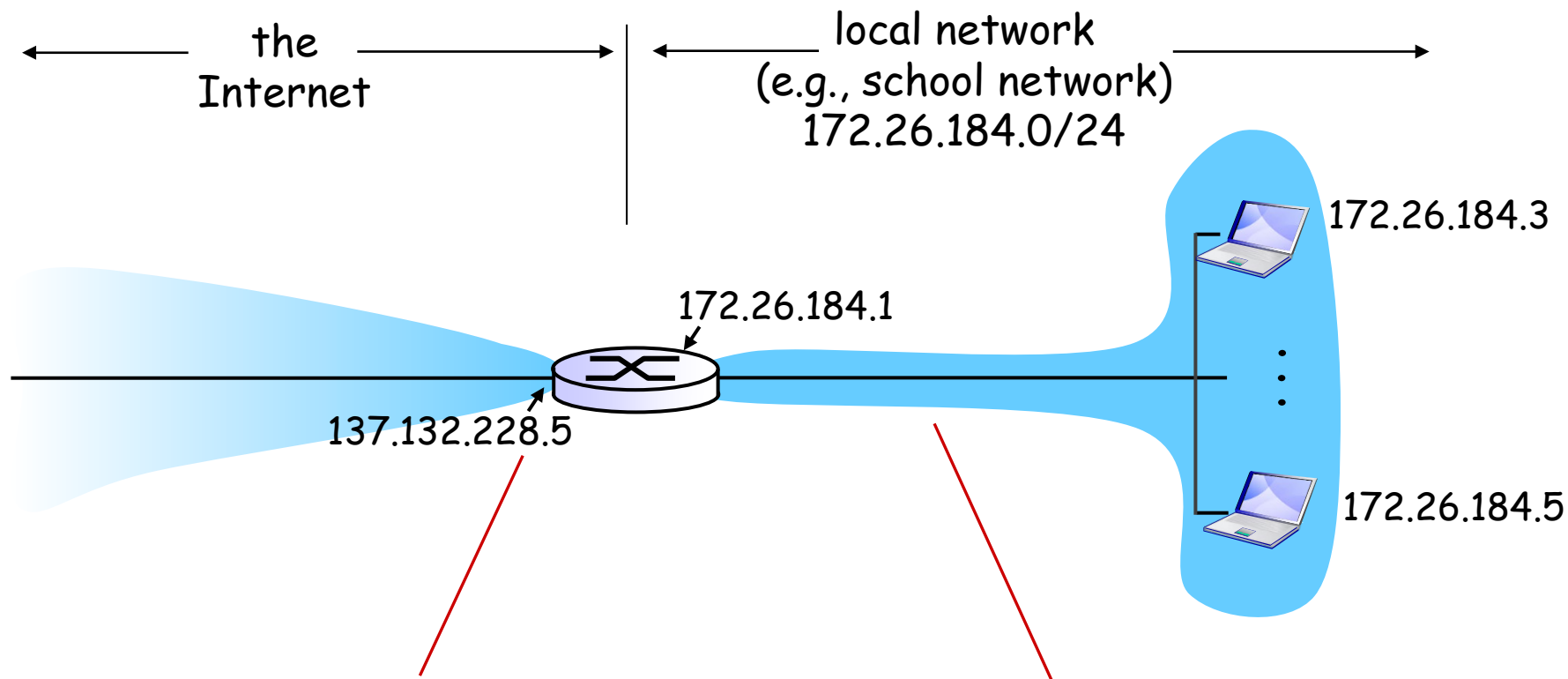
4.3 The Internet Protocol (IP)

- 4.3.4 Network Address Translation

5.2 Routing Algorithms

5.6 ICMP

NAT: Network Address Translation



all datagrams *leaving* local network have the *same* source NAT IP address: 137.132.228.5

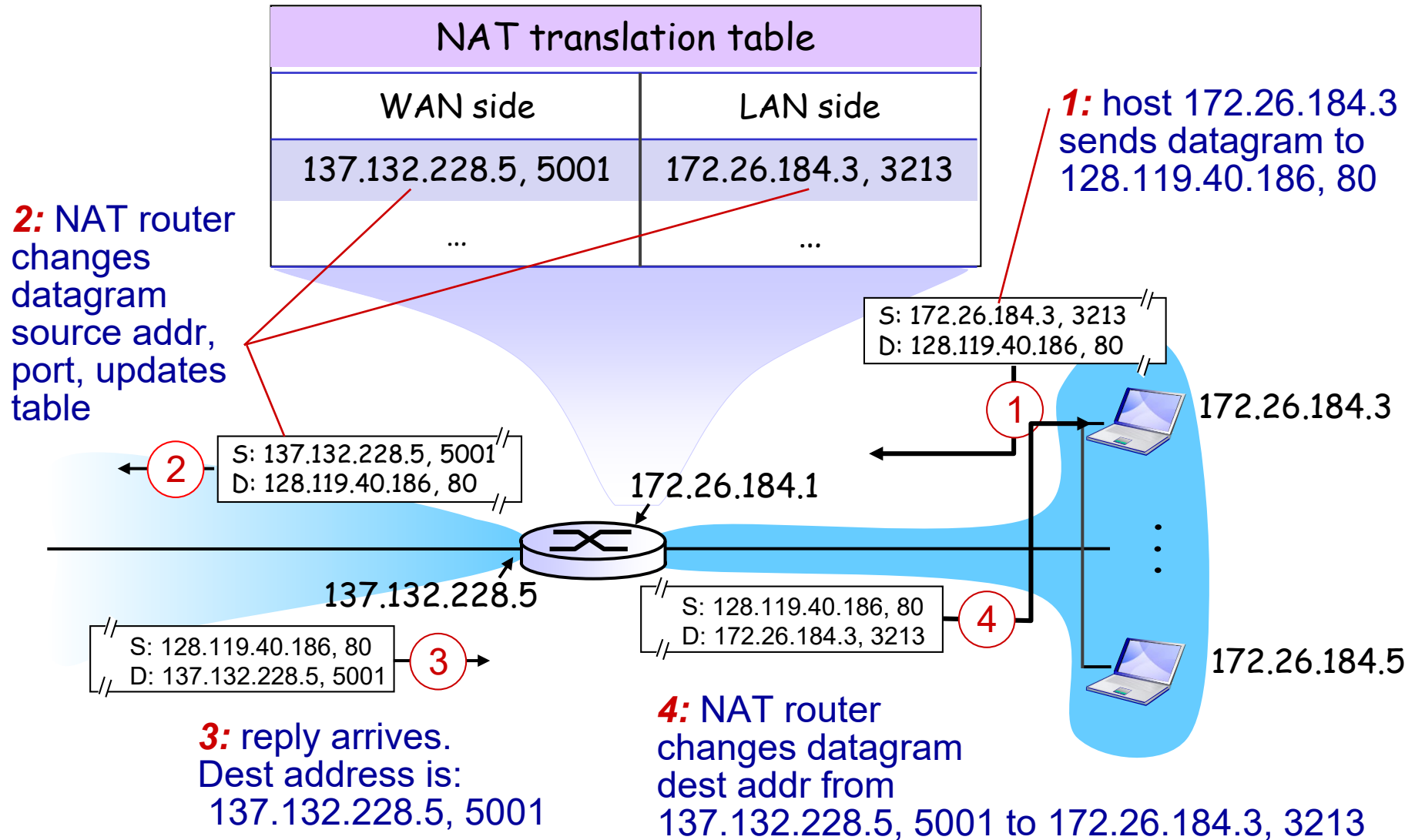
Within local network, hosts use private IP addresses 172.26.184.* for communication

NAT: Implementation

❖ NAT routers must:

- **Replace** (source IP address, port #) of every **outgoing datagram** to (NAT IP address, new port #).
- **Remember** (in NAT translation table) the mapping from (source IP address, port #) to (NAT IP address, new port #).
- **Replace** (NAT IP address, new port #) in destination fields of every **incoming datagram** with corresponding (source IP address, port #) stored in NAT translation table.

NAT: Illustration



NAT: Motivation and Benefits

- ❖ No need to rent a range of public IP addresses from ISP: just one public IP for the NAT router.
- ❖ All hosts use private IP addresses. Can change addresses of hosts in local network without notifying the outside world.
- ❖ Can change ISP without changing addresses of hosts in local network.
- ❖ Hosts inside local network are not explicitly addressable and visible by outside world (a security plus).

Lectures 6&7: Roadmap

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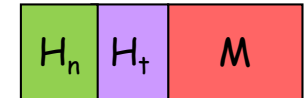
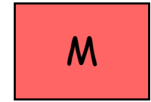
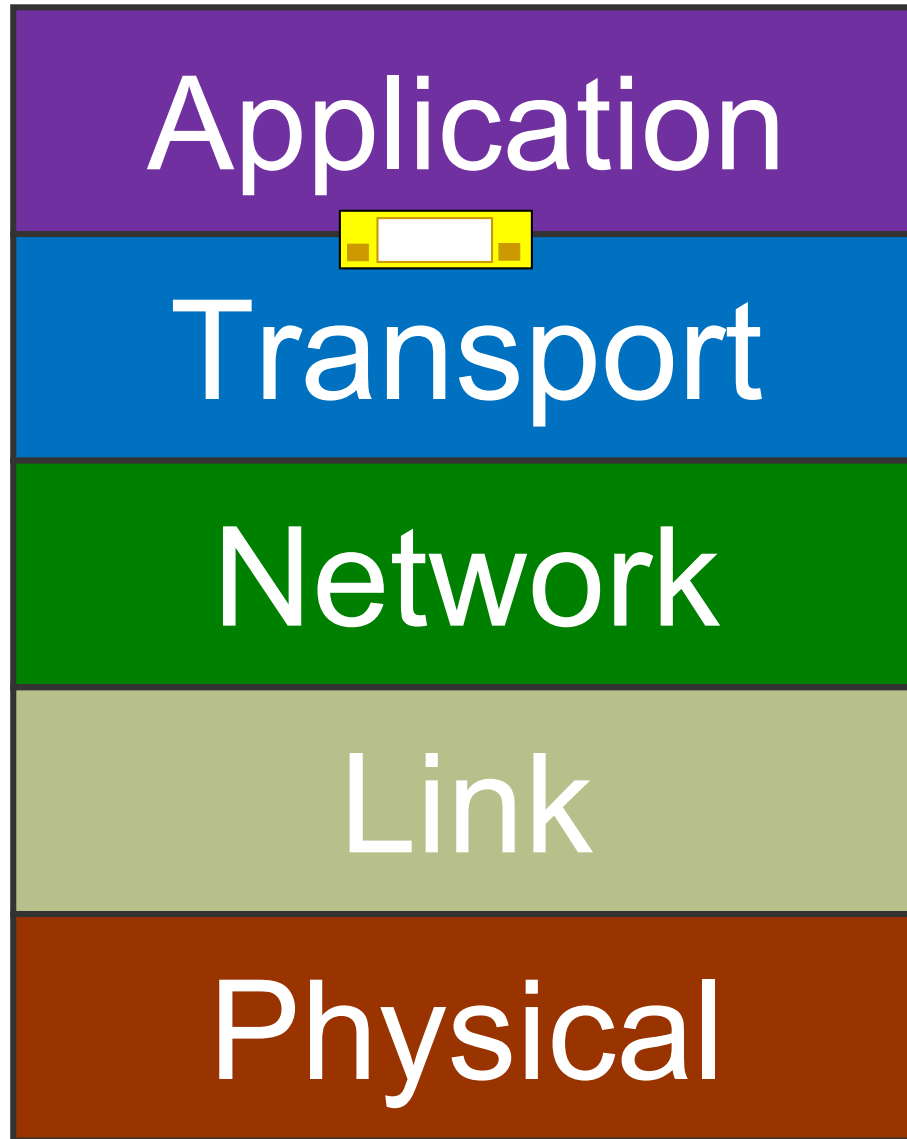
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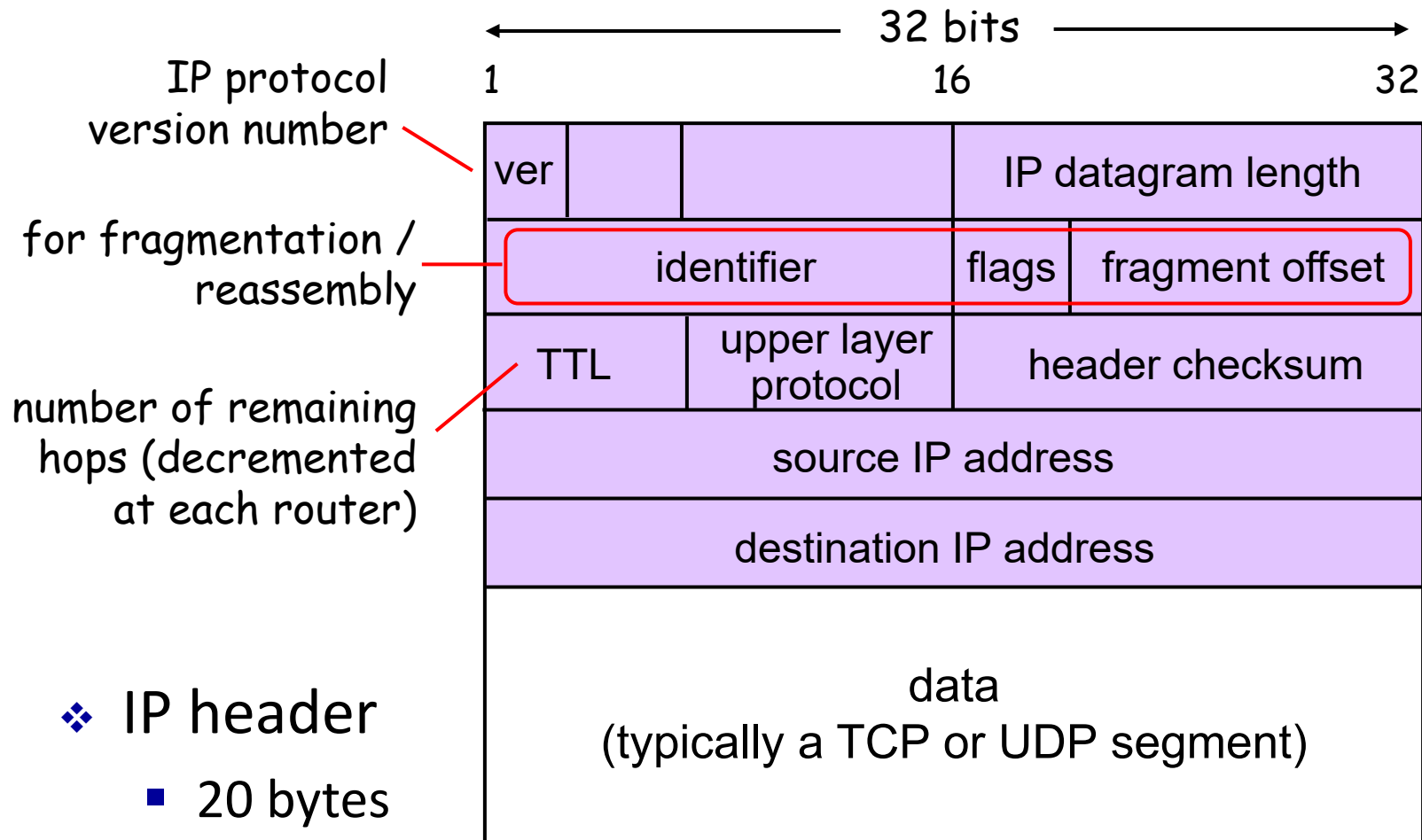
- 4.3.1 IPv4 Datagram Format
- 4.3.2 IPv4 Datagram Fragmentation
- 4.3.5 IPv6 (non-examinable)

5.2 Routing Algorithms

5.6 ICMP



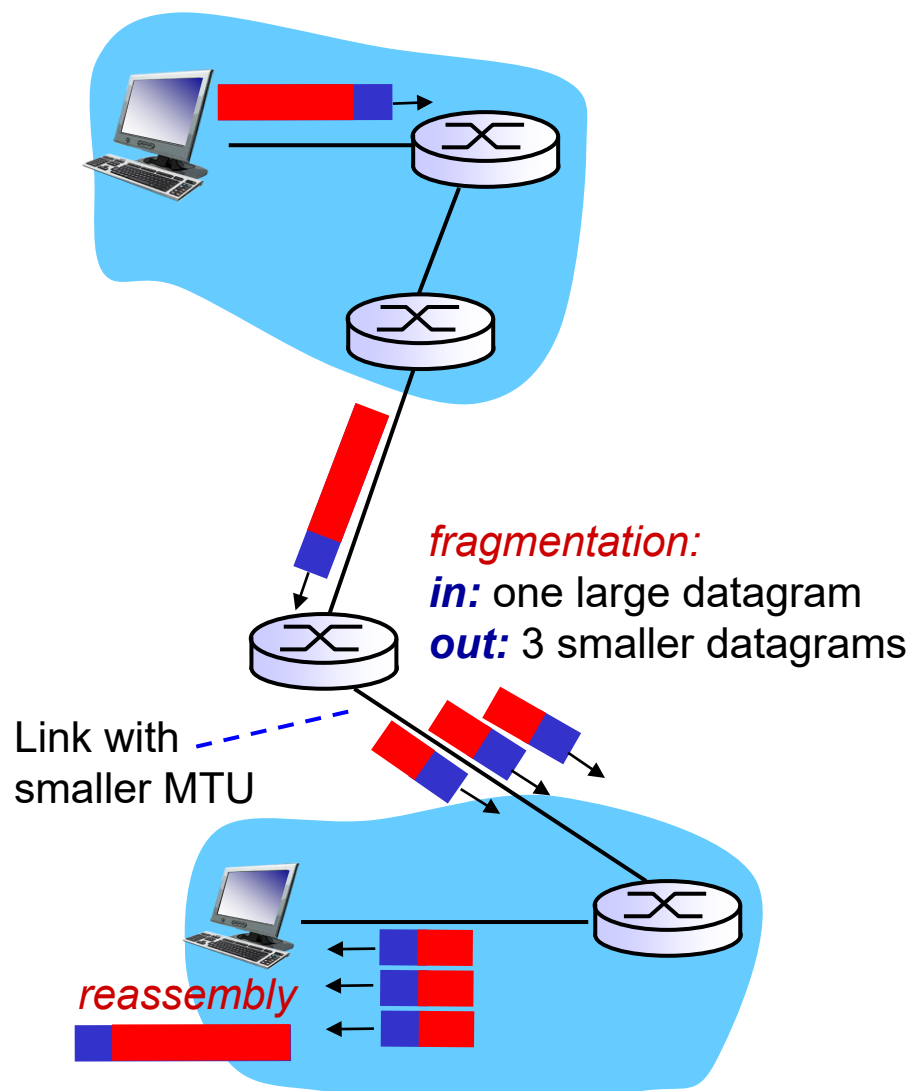
IPv4 Datagram Format



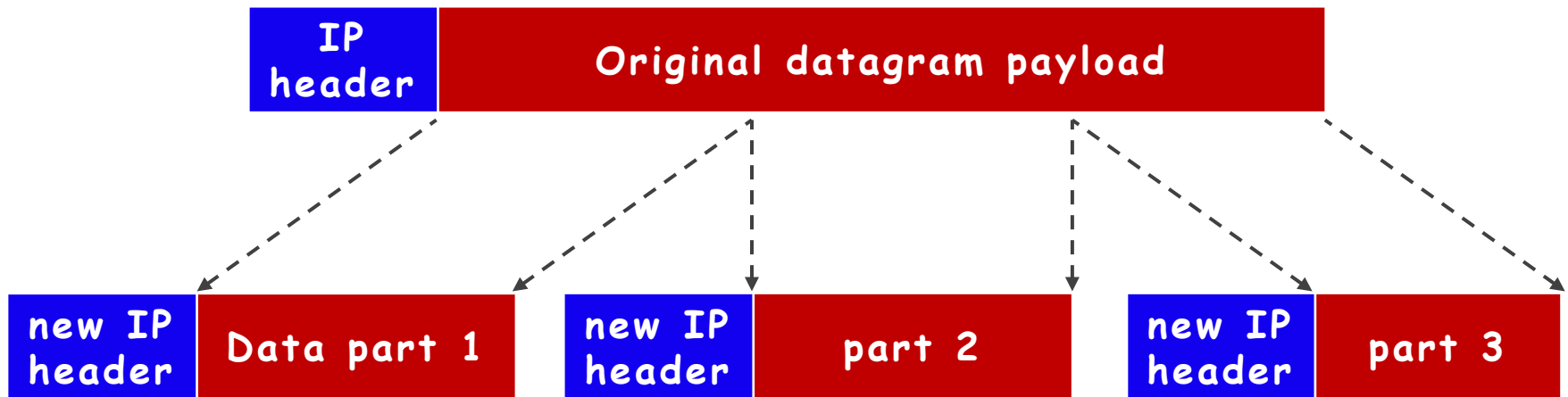
(some fields are not shown)

IP Fragmentation & Reassembly

- ❖ Different links may have different **MTU (Max Transfer Unit)** – the maximum amount of data a link-level frame can carry.
- ❖ “Too large” IP datagrams may be fragmented by routers.



IP Fragmentation Illustration



- ❖ Destination host will reassemble the packet.
- ❖ IP header fields are used to identify fragments and their relative order.

IP Fragmentation

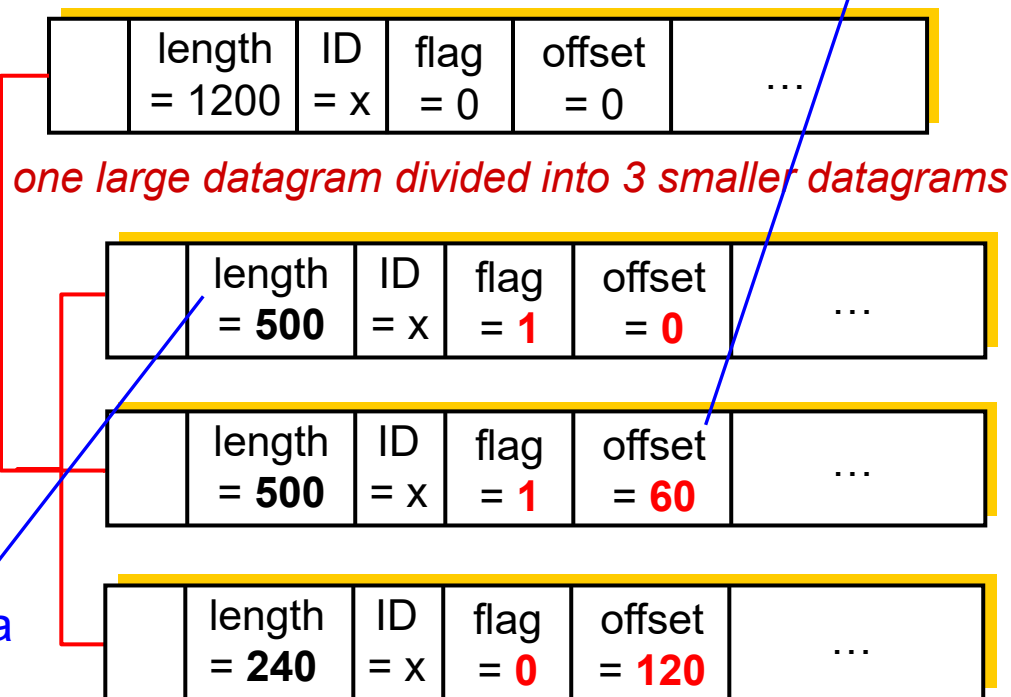
			length	
identifier			flags	offset
source IP address				
destination IP address				

- ❖ Flag (frag flag) is set to
 - **1** if there is next fragment from the same segment.
 - **0** if this is the last fragment.
- ❖ Offset is expressed in unit of 8-bytes.

❖ Example

- 20 bytes of IP header
- 1,200 byte IP datagram
- MTU = 500 bytes

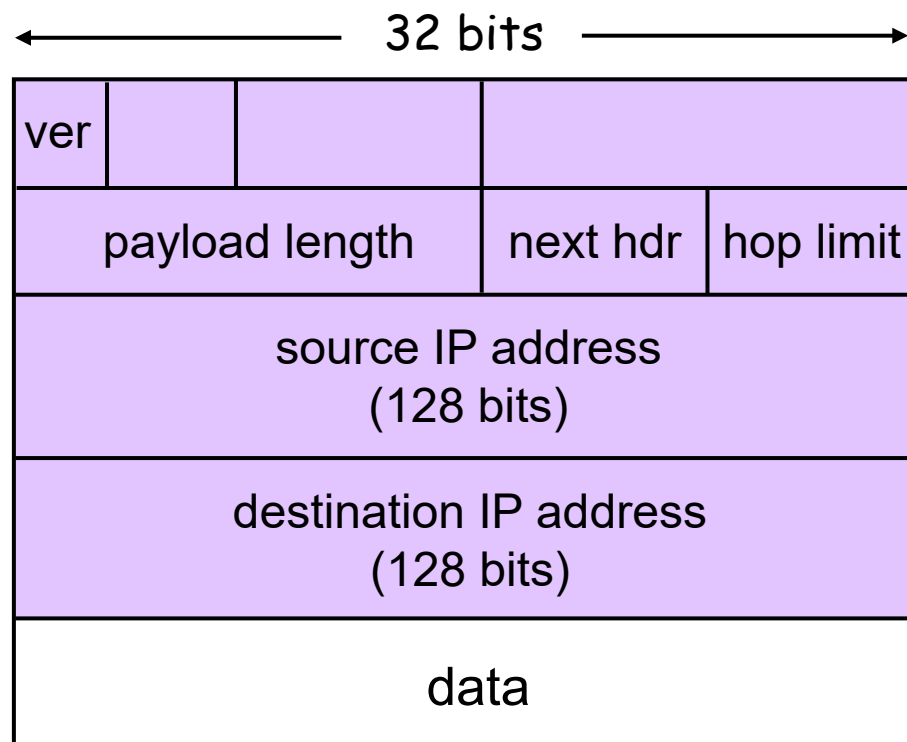
carry 480
bytes of data



IPv6

**Non-
examinable**

- ❖ IPv6 is designed to replace IPv4.
- ❖ Primary motivation: 32-bit IPv4 address space is soon to be completely allocated.
- ❖ IPv6 datagram:
 - 40 byte header



(some fields are not shown)

Example IPv6 address (in hexadecimal):
2001:0db8:85a3:0042:1000:8a2e:0370:7334

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ICMP

- ❖ **ICMP (Internet Control Message Protocol)** is used by hosts & routers to communicate network-level information.
 - Error reporting: unreachable host / network / port / protocol
 - Echo request/reply (used by ping)
- ❖ ICMP messages are carried in IP datagrams.
 - ICMP header starts after IP header.

ICMP Type and Code

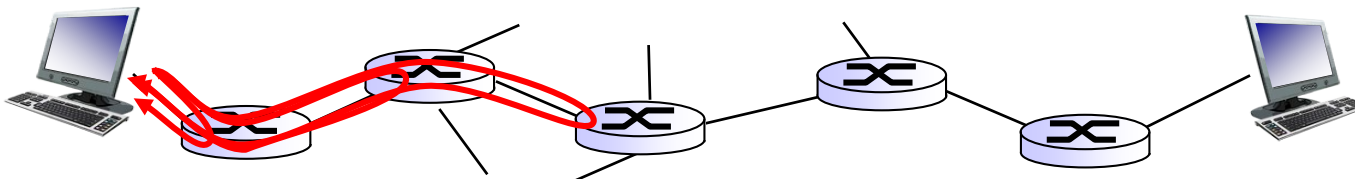
❖ ICMP header: Type + Code + Checksum + others.

Type	Code	Description
8	0	echo request (ping)
0	0	echo reply (ping)
3	1	dest host unreachable
3	3	dest port unreachable
11	0	TTL expired
12	0	bad IP header

Selected ICMP Type and subtype (Code)

Examples: *ping* and *traceroute*

- ❖ The command **ping** sees if a remote host will respond to us – do we have a connection?
- ❖ The command **traceroute** sends a series of small packets across a network, and attempts to display the route (or path) that the messages would take to get to a remote host.



Lectures 6&7: Summary

- ❖ An IP address is associated with a network interface. A device may have multiple network interfaces, thus multiple IP addresses.
- ❖ DHCP automates the assignment of IP addresses in an organization's network.
- ❖ On TCP/IP networks, subnets are defined as all devices whose IP addresses have the same network (subnet) prefix.
- ❖ Subnet mask is useful in checking if two hosts are on the same subnet.

Lectures 6&7: Summary

- ❖ Routing is the process of selecting best paths in a network.
- ❖ **NAT** maps one IP addresses space into another.
 - Commonly used to hide an entire private IP address space behind a single public IP address.
 - NAT router uses stateful translation tables to remember the mapping.
- ❖ **ICMP** is used by routers to send error messages.
 - E.g. when TTL is 0, a packet is discarded and an ICMP error message is sent to the datagram's source address.