

NATIONAL UNIVERSITY OF SINGAPORE

ANSWERS**CS2100 – COMPUTER ORGANISATION**

(Semester 2: AY2017/18)

ANSWER BOOKLET

Time Allowed: 2 Hours

INSTRUCTIONS TO CANDIDATES

1. This answer booklet consists of **SIX (6)** printed pages.
2. Fill in your Student Number **with a pen clearly** below. Do **NOT** write your name.
3. You may write your answers in pencil (2B or above).

STUDENT NUMBER
(fill in with a **pen**):

A	0	1						
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For examiner's use only		
<i>Question</i>	<i>Total</i>	<i>Marks</i>
Q1	10	
Q2	15	
Q3	20	
Q4	12	
Q5	15	
Q6	14	
Q7	14	
<i>Total</i>	100	

Write your answers in the box/space provided.

1a.
[4]

11 A
33 f

1b.
[3]

123.125

42F64000 = 0 100 0010 1111 0110 0100 0000 0000 0000
10000101 = 133 - 127 = 6
 $1.111011001_2 \times 2^6 = 1111011.001_2 = 123.125$

1c.
[3]

F = 1

Circuit delay = **9t**

Q1:

/10

2a.
[10]

$$DA = \mathbf{A \cdot C' + B \cdot C' \cdot D'}$$

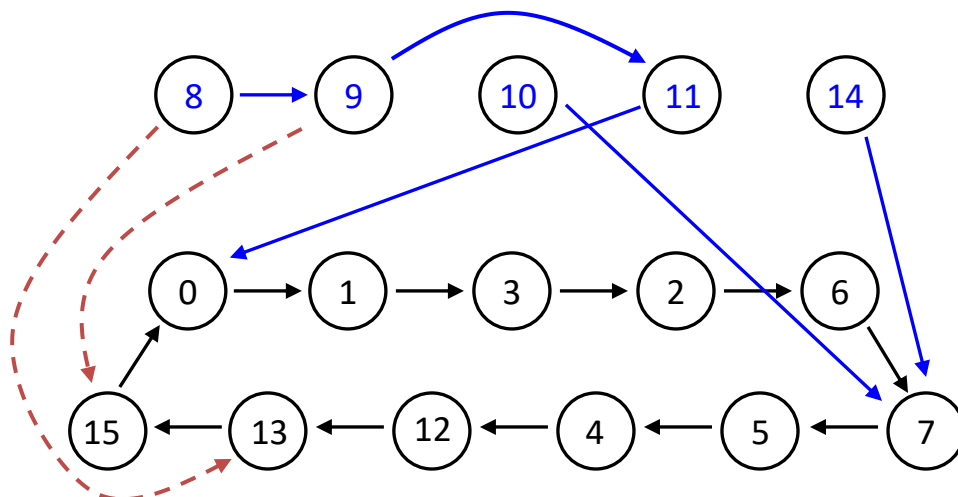
$$DB = \mathbf{A' \cdot B + C \cdot D' + B \cdot C'}$$
 or $\mathbf{A' \cdot B + C \cdot D' + A \cdot C'}$

$$TC = \mathbf{A \cdot D + B \cdot C \cdot D + B' \cdot C' \cdot D}$$

$$JD = \mathbf{A + B \cdot C + B' \cdot C'}$$

$$KD = \mathbf{A \cdot C + B' \cdot C + A' \cdot B \cdot C'}$$

2b.
[5]



There are two sets of answers, where states 8 and 9 may transit to different states.

Set 1: $8 \rightarrow 9$; $9 \rightarrow 11$ (Correspond to $DB = A' \cdot B + C \cdot D' + B \cdot C'$)

Set 2: $8 \rightarrow 13$; $9 \rightarrow 15$ (Correspond to $DB = A' \cdot B + C \cdot D' + A \cdot C'$)

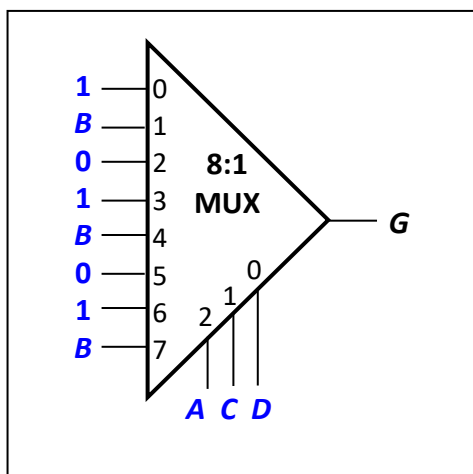
Q2:

/15

3a.
[4]

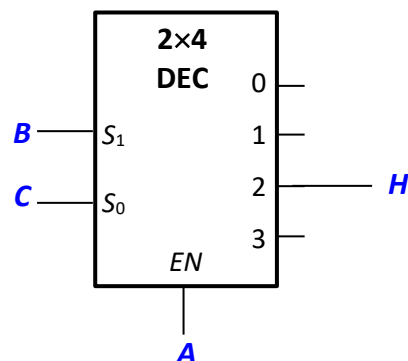
$$F = A$$

3b.
[4]



3c.
[4]

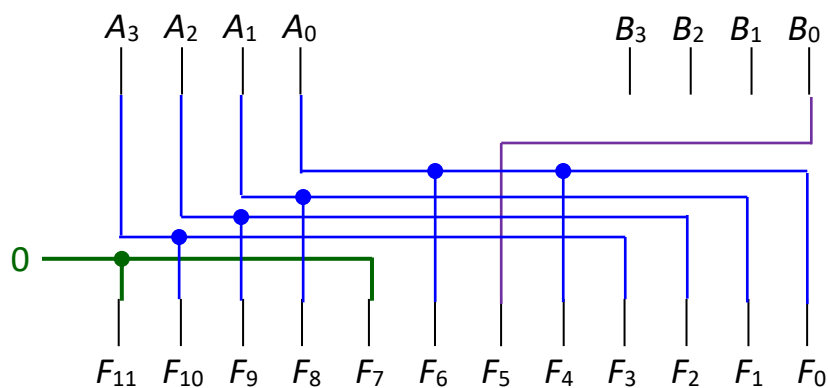
Many alternative answers.



3d.
[8]

A				5×A							
A ₃	A ₂	A ₁	A ₀								
0	0	0	0	0	0	0	0	0	0	0	0
0	0	0	1	0	0	0	0	0	1	0	1
0	0	1	0	0	0	0	1	0	0	0	0
0	0	1	1	0	0	0	1	0	1	0	1
0	1	0	0	0	0	1	0	0	0	0	0
0	1	0	1	0	0	1	0	0	1	0	1
0	1	1	0	0	0	1	1	0	0	0	0
0	1	1	1	0	0	1	1	0	1	0	1
1	0	0	0	0	1	0	0	0	0	0	0
1	0	0	1	0	1	0	0	0	1	0	1

Complete your circuit here.



Q3:

/20

4a.
[2]

Maximum total instructions = **379**

$$(2^6 - 5) + (5 \times 2^6) = 59 + (5 \times 64) = 59 + 320 = 379$$

4b.
[3]

Stuck-at-0 fault at bit 6 of the instruction

Eg: **addi \$t0, \$zero, 64**

This is supposed to put the value 64 into \$t0. However, due to the stuck-at-0 error at bit 6, the value in \$t0 will be 0.

Other similar examples using the same argument are acceptable, for example, some students use **sll \$t1, \$t1, 1** (pre-condition: \$t1 cannot be zero), but this would require comparing the old and new values of \$t1. The **addi** solution above is simpler.

4c.
[3]

Stuck-at-0 fault at ALUSrc

If the instruction is correctly carried out, **lw \$t1, 0(\$t0)** would have loaded the value at address **12** into \$t1.

With stuck-at-0 fault at ALUSrc, the instruction would have loaded the value at address **46** into \$t1 instead.

Hence, we could first load different values in addresses **12** and **46** before using the test.

4d.
[4]

Adding bne instruction

Since ALUop code 11 is not used, we may use it for bne.

Hence, branch taken = **ALUop1 AND ALUop2 AND !(isZero)**
where (isZero) is the output from the ALU.

or

Use ALUop code 01 (same as beq).

Hence, branch taken = **!ALUop1 AND ALUop2 AND !(isZero)**
where (isZero) is the output from the ALU, or
branch taken = **Branch AND !(isZero)**.

I tried to be more liberal in marking this part as most students didn't give a sufficiently rigorous answer.

Q4:

/12

5a.
[1]

lw \$s2, 0(\$t0)

5b.
[4]

Array A:

11	10	31	14	9	42	6	11
----	----	----	----	---	----	---	----

5c.
[4]

```

int i;
for (i=n-1; i>=0; i-=2) {
    if (B[i]%4 == 3)
        A[i]++;
    else
        A[i] += B[i];
}

```

These elements are unchanged.

5d.
[2]

0018c080

5e.
[2]

0810001c

5f.
[2]

0300782a

Q5:

/15

6a.
[2]

Minimum = 1203 (3 + 100×12)

Maximum = 1303 (3 + 100×13)

6b.
[6]

(Due to control) Inst2, Inst4, Inst11, Inst13, Inst14
 (Due to data) Inst8, Inst10, Inst17

6c.
[3]

24 cycles

6d.
[3]

26 cycles

Q6:

/14

- 7a. [2] Index: 5 bits; Offset: 4 bits
- 7b. [4] $A[1023] \rightarrow$ Index 31; $B[1023] \rightarrow$ Index 15;
- 7c. [2] Array A: 1024 accesses; Array B: 512 accesses
- 7d. [2] Array A: 75 %; Array B: 50 %
- 7e. [4] Misses: 9; Hits: 6138

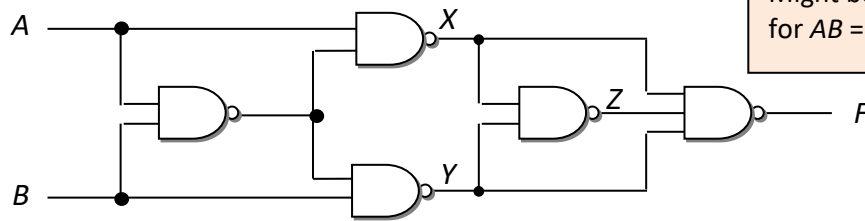
Q7:

/14

=== END OF PAPER ===

Workings

Q1c.



Might be easier just to trace for $AB = 00, 01, 10, 11$.

$$X = ((A \cdot B)' \cdot A)' = A \cdot B + A' = B + A'; \quad Y = ((A \cdot B)' \cdot B)' = A \cdot B + B' = A + B'$$

$$Z = (X \cdot Y)' = X' + Y' = A \cdot B' + A' \cdot B$$

$$F = (X \cdot Y \cdot Z)' = [(B + A') \cdot (A + B') \cdot (A \cdot B' + A' \cdot B)]' = [(A \cdot B + A' \cdot B') \cdot (A \cdot B' + A' \cdot B)]' = 0' = 1$$

Q3a.

A	B	F
0	0	0
0	1	0
1	0	1
1	1	1

Q3b.

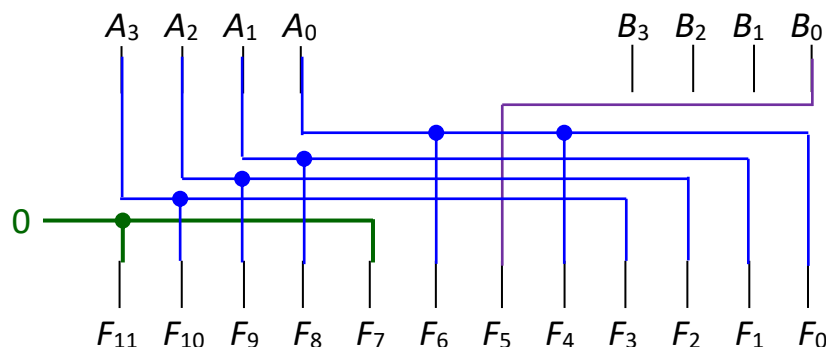
A	B	C	D	G
0	0	0	0	1
0	0	0	1	0
0	0	1	0	0
0	0	1	1	1
0	1	0	0	1
0	1	0	1	1
0	1	1	0	0
0	1	1	1	1
1	0	0	0	0
1	0	0	1	0
1	0	1	0	1
1	0	1	1	0
1	1	0	0	1
1	1	0	1	0
1	1	1	0	1
1	1	1	1	1

Q3c. $H = A \cdot B \cdot C' \cdot D' + A \cdot B \cdot C' \cdot D$
 $= A \cdot B \cdot C'$

Q3d. Note that $50 \times A$ on $A=0, 1, 2, 3, 4, \dots, 9$ gives us 000, 051, 102, 153, 204, ..., 459. It can be seen that the left-most digit 0, 0, 1, 1, 2, 2, 3, 3, 4, 4 is simply $A/2$. Hence the left-most 4 bits $F_{11}F_{10}F_9F_8$ are $0A_3A_2A_1$.

The middle digit is 0 (or 0000 in binary) if A is even (i.e. $A_0=0$), or 5 (or 0101 in binary) if A is odd (i.e. $A_0=1$). The right-most digit is A itself.

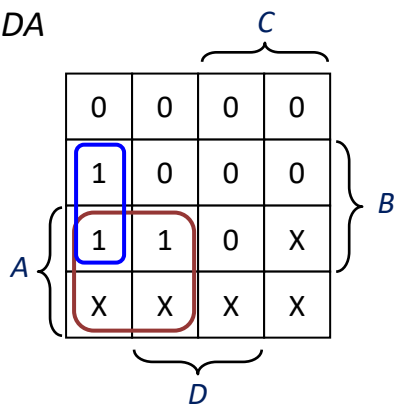
Finally, $20 \times (B\%2)$ is achieved by putting B_0 into F_5 .



Q2.

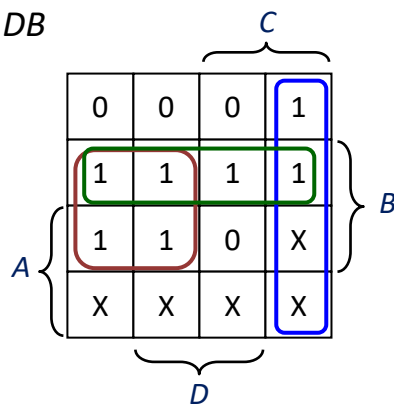
Current state				Nex state						
A	B	C	D	A ⁺	B ⁺	C ⁺	D ⁺	TC	JD	KD
0	0	0	0	0	0	0	1	0	1	X
0	0	0	1	0	0	1	1	1	X	0
0	0	1	0	0	1	1	0	0	0	X
0	0	1	1	0	0	1	0	0	X	1
0	1	0	0	1	1	0	0	0	0	X
0	1	0	1	0	1	0	0	0	X	1
0	1	1	0	0	1	1	1	0	1	X
0	1	1	1	0	1	0	1	1	X	0
1	0	0	0	X(1)	X(0,1)	X(0)	X(1)	X(0)	X(1)	X(0)
1	0	0	1	X(1)	X(0,1)	X(1)	X(1)	X(1)	X(1)	X(0)
1	0	1	0	X(0)	X(1)	X(1)	X(1)	X(0)	X(1)	X(1)
1	0	1	1	X(0)	X(0)	X(0)	X(0)	X(1)	X(1)	X(1)
1	1	0	0	1	1	0	1	0	1	X
1	1	0	1	1	1	1	1	1	X	0
1	1	1	0	X(0)	X(1)	X(1)	X(1)	X(0)	X(1)	X(1)
1	1	1	1	0	0	0	0	1	X	1

DA



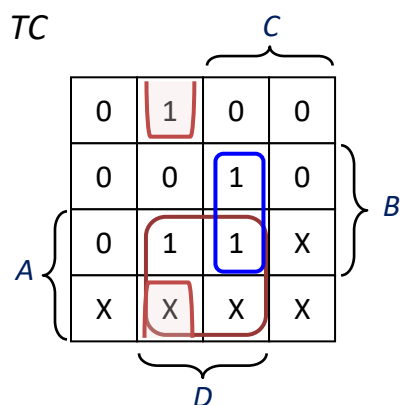
$$DA = A \cdot C' + B \cdot C' \cdot D'$$

DB



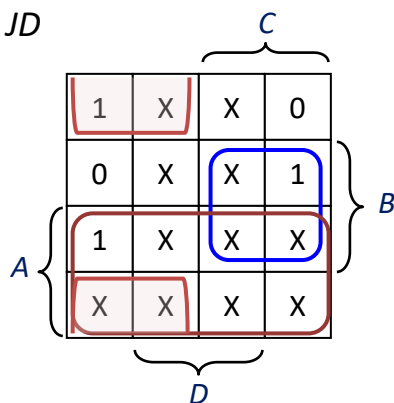
$$DB = A' \cdot B + C \cdot D' + B \cdot C' \\ \text{or } A' \cdot B + C \cdot D' + A \cdot C'$$

TC



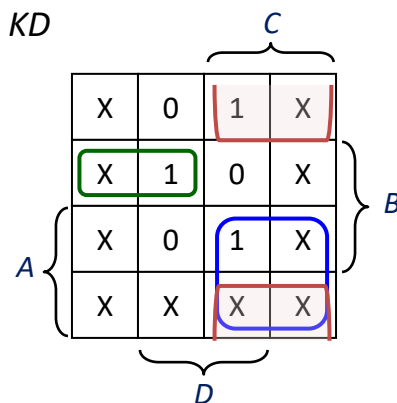
$$TC = A \cdot D + B \cdot C \cdot D + B' \cdot C' \cdot D$$

JD



$$JD = A + B \cdot C + B' \cdot C'$$

KD



$$KD = A \cdot C + B' \cdot C + A' \cdot B \cdot C'$$

Q5. Tested on QTSpm

Data	Text
Text	
User Text Segment [00400000]..[00440000]	
[00400000] 8fa40000	lw \$4, 0(\$29) ; 183: lw \$a0 0(\$sp) # argc
[00400004] 27a50004	addiu \$5, \$29, 4 ; 184: addiu \$a1 \$sp 4 # argv
[00400008] 24a60004	addiu \$6, \$5, 4 ; 185: addiu \$a2 \$a1 4 # envp
[0040000c] 00041080	sll \$2, \$4, 2 ; 186: sll \$v0 \$a0 2
[00400010] 00c23021	addu \$6, \$6, \$2 ; 187: addu \$a2 \$a2 \$v0
[00400014] 0c100009	jal 0x00400024 [main] ; 188: jal main
[00400018] 00000000	nop ; 189: nop
[0040001c] 3402000a	ori \$2, \$0, 10 ; 191: li \$v0 10
[00400020] 0000000c	syscall ; 192: syscall # syscall 10 (exit)
[00400024] 3c101001	lui \$16, 4097 [A] ; 8: la \$s0, A # \$s0 is the base address of array A
[00400028] 3c011001	lui \$1, 4097 [B] ; 9: la \$s1, B # \$s1 is the base address of array B
[0040002c] 34310020	ori \$17, \$1, 32 [B]
[00400030] 3c011001	lui \$1, 4097 [n] ; 10: la \$t0, n # \$t0 is the address of n (size of array)
[00400034] 34280040	ori \$8, \$1, 64 [n]
[00400038] 8d120000	lw \$18, 0(\$8) ; 11: lw \$s2, 0(\$t0) # \$s2 is the content of n
[0040003c] 12400011	beq \$18, \$0, 68 [End-0x0040003c]
[00400040] 2258ffff	addi \$24, \$18, -1 ; 16: addi \$t8, \$s2, -1 # \$t8 = n-1
[00400044] 0018c080	sll \$24, \$24, 2 ; 17: sll \$t8, \$t8, 2 # \$t8 = 4*(n-1)
[00400048] 02184020	add \$8, \$16, \$24 ; 18: add \$t0, \$s0, \$t8
[0040004c] 02384820	add \$9, \$17, \$24 ; 19: add \$t1, \$s1, \$t8
[00400050] 8d0a0000	lw \$10, 0(\$8) ; 20: lw \$t2, 0(\$t0) # \$t2 = A[i]
[00400054] 8d2b0000	lw \$11, 0(\$9) ; 21: lw \$t3, 0(\$t1) # \$t3 = B[i]
[00400058] 316c0003	andi \$12, \$11, 3 ; 22: andi \$t4, \$t3, 3 # \$t4 = B[i]%4
[0040005c] 218cffff	addi \$12, \$12, -3 ; 23: addi \$t4, \$t4, -3 # \$t4 = (B[i]%4)-3
[00400060] 11800003	beq \$12, \$0, 12 [A1-0x00400060] ; 24: beq \$t4, \$zero, A1 # if (B[i]%4 == 3) goto A1
[00400064] 014b5020	add \$10, \$10, \$11 ; 25: add \$t2, \$t2, \$t3 # else A[i] += B[i]
[00400068] 0810001c	j 0x00400070 [A2] ; 26: j A2
[0040006c] 214a0001	addi \$10, \$10, 1 ; 27: addi \$t2, \$t2, 1 # A[i]++
[00400070] ad0a0000	sw \$10, 0(\$8) ; 28: sw \$t2, 0(\$t0)
[00400074] 2318ffff	addi \$24, \$24, -8 ; 29: addi \$t8, \$t8, -8 # i = i - 2
[00400078] 0300782a	slt \$15, \$24, \$0 ; 30: slt \$t7, \$t8, \$zero # if (\$t8
[0040007c] 11e0fff3	beq \$15, \$0, -52 [Loop-0x0040007c]
[00400080] 3402000a	ori \$2, \$0, 10 ; 32: li \$v0, 10 # system call code for exit
[00400084] 0000000c	syscall ; 33: syscall

Data: before and after

User data segment [10000000]..[10040000]

```

[10000000]..[1000ffff] 00000000
[10010000] 0000000011 0000000009 0000000031 0000000002
[10010010] 0000000009 0000000001 0000000006 0000000010
[10010020] 0000000003 0000000007 0000000002 0000000012
[10010030] 0000000011 0000000041 0000000019 0000000035
[10010040] 0000000008 0000000000 0000000000 0000000000
[10010050]..[1003ffff] 00000000

```

User data segment [10000000]..[10040000]

```

[10000000]..[1000ffff] 00000000
[10010000] 0000000011 0000000010 0000000031 0000000014
[10010010] 0000000009 0000000042 0000000006 0000000011
[10010020] 0000000003 0000000007 0000000002 0000000012
[10010030] 0000000011 0000000041 0000000019 0000000035
[10010040] 0000000008 0000000000 0000000000 0000000000
[10010050]..[1003ffff] 00000000

```

Q6. (c)

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30
I1 beq	F	D	E	M	W																									
I2 addi			F	D	E	M	W																							
I3 sll				F	D	E	M	W																						
I4 add					F	D	E	M	W																					
I5 add						F	D	E	M	W																				
I6 lw							F	D	E	M	W																			
I7 lw								F	D	E	M	W																		
I8 andi									F	D		E	M	W																
I9 addi										F	D		E	M	W															
I10 beq A1											F			D	E	M	W													
I11 add																														
I12 J A2																														
I13 A1: addi															F	D	E	M	W											
I14 A2: sw																F	D	E	M	W										
I15 addi																	F	D	E	M	W									
I16 slt																		F	D	E	M	W								
I17 beq																			F		D	E	M	W						

Q6. (d)

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30
I1 beq	F	D	E	M	W																									
I2 addi			F	D	E	M	W																							
I3 sll				F	D	E	M	W																						
I4 add					F	D	E	M	W																					
I5 add						F	D	E	M	W																				
I6 lw							F	D	E	M	W																			
I7 lw								F	D	E	M	W																		
I8 andi									F	D		E	M	W																
I9 addi										F	D		E	M	W															
I10 beq A1											F			D	E	M	W													
I11 add															F	D	E	M	W											
I12 J A2																F	D													
I13 A1: addi																														
I14 A2: sw																	F	D	W	M	W									
I15 addi																		F	D	E	M	W								
I16 slt																			F	D	E	M	W							
I17 beq																				F		D	E	M	W					

Q7.(a) Number of blocks = $128/4 = 32$; Index: 5 bits; Offset: 4 bits

(b) $A[0]$ at $0x10001000 \rightarrow A[1023]$ at $0x10001FFC$ (because $A[1024]$ at $0x10002000$)
 $10001FFC \rightarrow \dots 111\underline{1\ 1111}\ 1100 \rightarrow \text{Index } 31$

$B[0]$ at $0x1003F100 \rightarrow B[1023]$ at $0x100400FC$ (because $B[1024]$ at $0x10040100$)
 $100400FC \rightarrow \dots 000\underline{0\ 1111}\ 1100 \rightarrow \text{Index } 15$

(c) 512 elements for each array are accessed: $[1023], [1021], \dots, [3], [1]$.

A : 1024 accesses (1 read and 1 write per element); B : 512 accesses

(d) Since $A[1023]$ and $B[1023]$ are mapped to different indices, there will not be racing.

A : 256 misses, 768 hits, hit rate = 75%

B : 256 misses, 256 hits, hit rate = 50%

(e) 6147 instructions ($3 + 512 \times 12$). 9 misses, 6138 hits.

Index field: 3 bits; Offset: 3 bits

The first **beq** instruction is at addr $0x0040003c \rightarrow \dots 00\underline{11\ 1}\ 100 \rightarrow \text{index } 7, 2^{\text{nd}}$ word

The following shows the cache content for the first iteration:

Index		
0	I2	I3
1	I4	I5
2	I6	I7
3	I8	I9
4	I10	
5		I13
6	I14	I15
7	I16	I17

The misses are I1, I2, I4, I6, I8, I10, I13, I14 and I16. After that, all instructions in the loop are present in the cache.