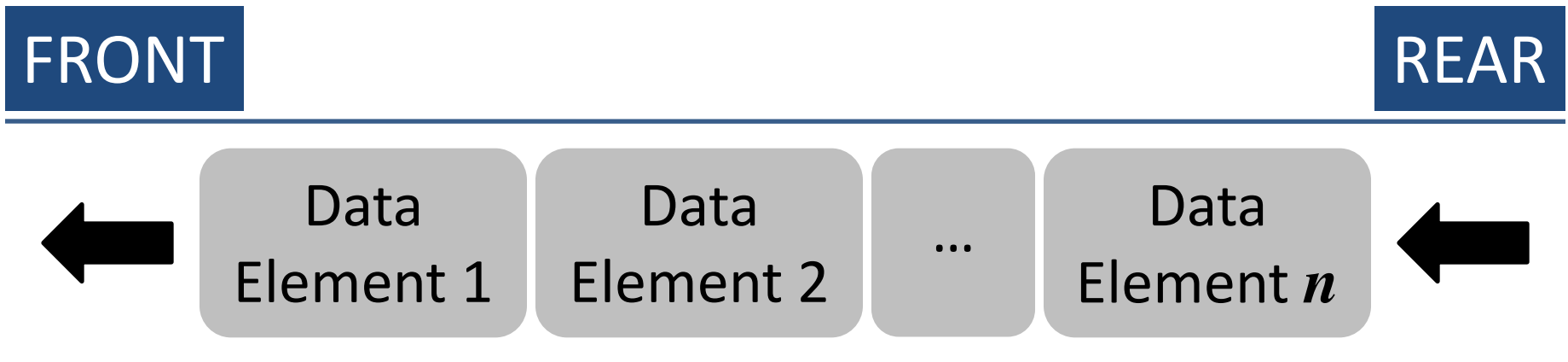


Queues

Introduction

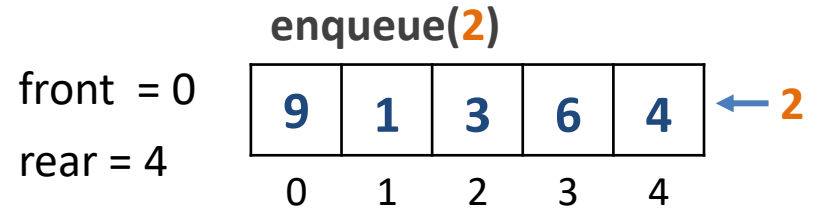
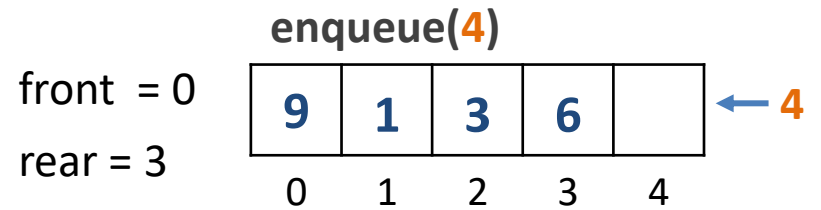
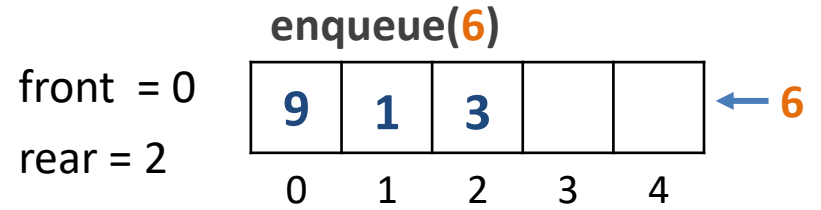
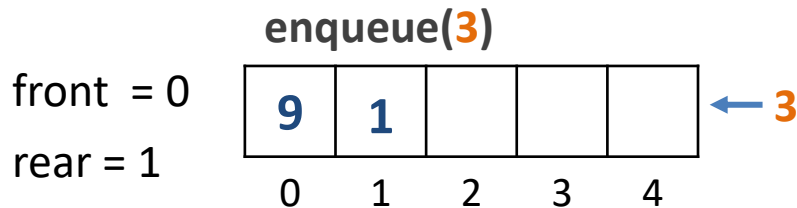
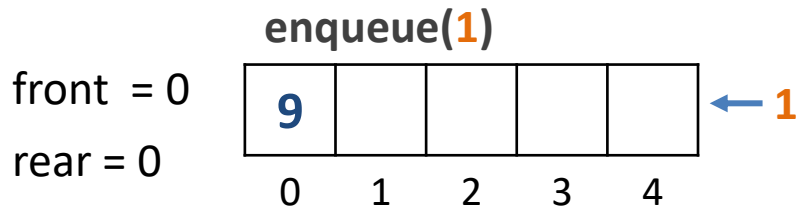
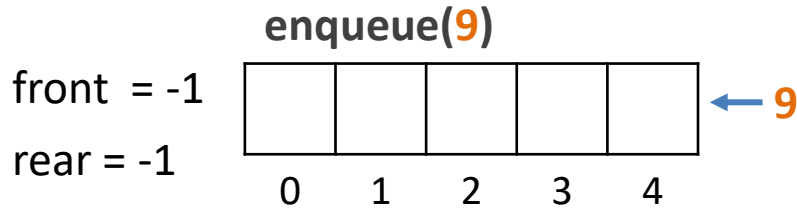
- First in, first out (FIFO) structure (equivalent to Last in, last out (LIFO) structure).
- An ordered list of homogeneous elements in which
 - Insertions take place at one end (REAR).
 - Deletions take place at the other end (FRONT).



Operations

- Two primary operations:
 - **enqueue()** – Adds an element to the *rear* of a queue.
 - **dequeue()** – Removes the *front* element of the queue.
- Other operations for effective functionality:
 - **isFull()** – Check if queue is full. *OVERFLOW*
 - **isEmpty()** – Check if queue is empty. *UNDERFLOW*
 - **size()** – Returns the number of elements in the queue.
 - **peek()** – Returns the element at the front of the queue.

Queue – Enqueue

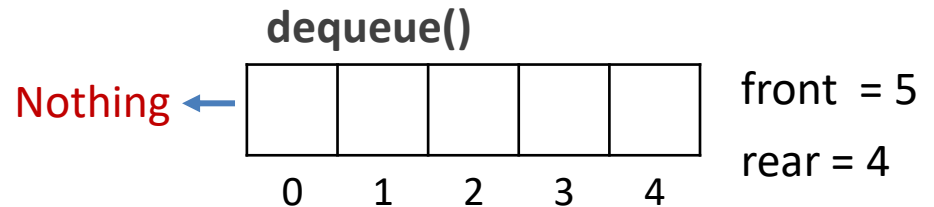
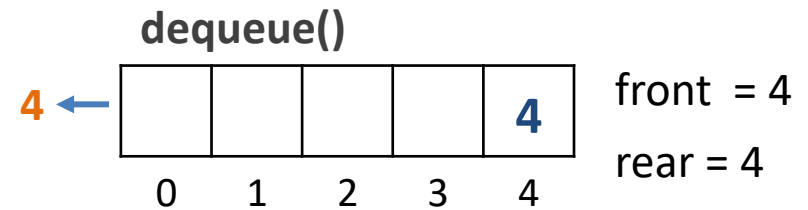
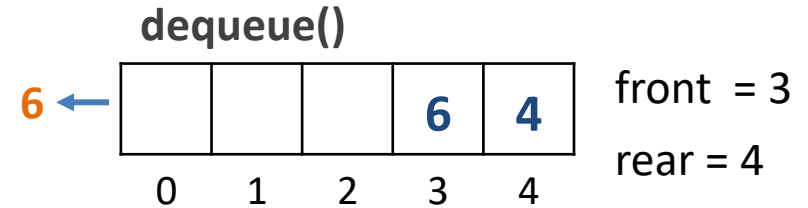
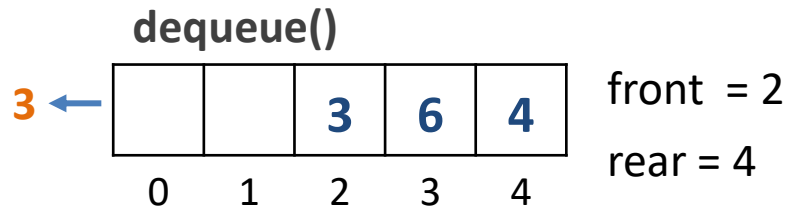
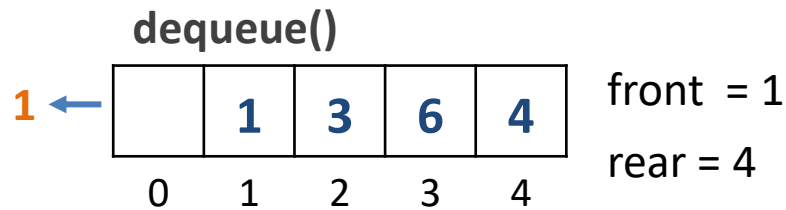
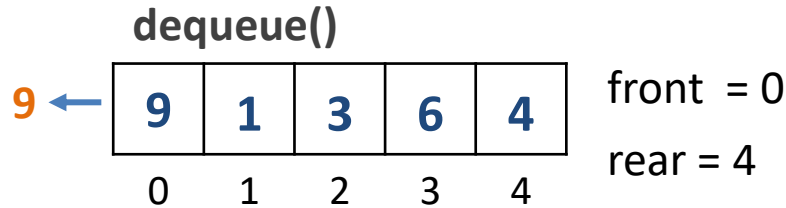


The queue is full, no more elements
can be added. **OVERFLOW**

OVERFLOW

←

Queue – Dequeue



The queue is empty, no element
can be removed. **UNDERFLOW**

UNDERFLOW

Queue as an ADT

- A queue is an ordered list of elements of same data type.
- Elements are always inserted at one end (rear) and deleted from another end (front).
- Following are its basic operations:
 - $Q = \textit{init}()$ – Initialize an empty queue.
 - $\textit{size}()$ – Returns the number of elements in the queue.
 - $\textit{isEmpty}(Q)$ – Returns "true" if and only if the queue Q is empty, i.e., contains no elements.

Contd...

- *isFull(Q)* – Returns "true" if and only if the queue Q has a bounded size and holds the maximum number of elements it can.
- *front(Q)* – Returns the element at the front of the queue Q.
- *Q = enqueue(Q,x)* – Inserts an element x at the rear of the queue Q.
- *Q = dequeue(Q)* – Removes an element from the front of the queue Q.
- *print(Q)* – Prints the elements of the queue Q from front to rear.

Implementation

- Using static arrays
 - Realizes queues of a maximum possible size.
 - Front is maintained at the smallest index and rear at the maximum index values in the array.
- Using dynamic linked lists
 - Choose beginning of the list as the front and tail as rear of the queue.

Static Array Implementation

Enqueue Operation

- Let,
 - QUEUE be an array with N locations.
 - FRONT and REAR points to the front and rear of the QUEUE.
 - ITEM is the value to be inserted.
- 1. If ($\text{REAR} == N - 1$)
- 2. Print[Overflow]
- 3. Else
- 4. If ($\text{FRONT} == -1 \ \&\& \ \text{REAR} == -1$)
- 5. Set $\text{FRONT} = 0$ and $\text{REAR} = 0$.
- 6. Else
- 7. Set $\text{REAR} = \text{REAR} + 1$.
- 8. $\text{QUEUE}[\text{REAR}] = \text{ITEM}$.

Deque Operation

- Let,
 - QUEUE be an array with N locations.
 - FRONT and REAR points to the front and rear of the QUEUE.
 - ITEM holds the value to be deleted.
 - 1. If (FRONT == -1 || FRONT > REAR)
 - 2. Print[Underflow]
 - 3. Else
 - 4. ITEM = QUEUE[FRONT]
 - 5. Set FRONT = FRONT + 1

Static Array Implementation

```
1. #define MAXLEN 100
2. typedef struct
3. { int element[MAXLEN];
4.   int front, rear; } queue;
5. queue init ()
6. { queue Q;
7.   Q.front = Q.rear = -1;
8.   return Q; }
9. int size( queue Q )
10. { return ( Q.rear - Q.front + 1 ); }
11. int isEmpty ( queue Q )
12. { return ((Q.front == -1) ||
           (Q.front > Q.rear)); }
13. int isFull ( queue Q )
14. { return (Q.rear == MAXLEN - 1); }
15. int front ( queue Q )
16. { if (isEmpty(Q))
17.   printf("Empty queue\n");
18.   else
19.   return Q.element[Q.front]; }
```

Contd...

```
20. queue enqueue ( queue Q , int x )
21. {   if (isFull(Q))
22.     printf("OVERFLOW\n");
23.     else if (isEmpty(Q))
24.     {   Q.front = Q.rear = 0;
25.         Q.element[Q.rear] = x;
26.     }
27.     else
28.     {   ++Q.rear;
29.         Q.element[Q.rear] = x;
30.     }
31.     return Q; }

32. queue dequeue ( queue Q )
33. {   if (isEmpty(Q))
34.     printf("UNDERFLOW\n");
35.     else
36.     Q.front++;
37.     return Q; }
```

Contd...

```
38. void print ( queue Q )
39. {   int i;
40.     for (i = Q.front; i <= Q.rear; i++)
41.         printf("%d ",Q.element[i]); }

42. int main ()
43. {   queue Q;
44.     Q = init();
45.     Q = enqueue(Q,5);
46.     Q = enqueue(Q,3);
47.     Q = dequeue(Q);
48.     Q = enqueue(Q,7);
49.     Q = dequeue(Q);

50.     printf("Current queue : "); print(Q);
51.     printf(" with front = %d.\n", front(Q));
52.     Q = enqueue(Q,9);
53.     Q = enqueue(Q,3);
54.     Q = enqueue(Q,1);
55.     printf("Current queue : "); print(Q);
56.     printf(" with front = %d.\n", front(Q));
57.     printf("Size is %d.",size(Q));
58.     return 0; }
```

Dynamic Linked List Implementation

Enqueue Operation

- Let,
 - FRONT and REAR points to the front and rear of the QUEUE.
 - ITEM is the value to be inserted.
 - 1. Create a node pointer (temp).
 - 2. temp[data] = ITEM.
 - 3. temp[next] = NULL.
 - 4. If FRONT == NULL
 - 5. FRONT = REAR = temp.
 - 6. Else
 - 7. REAR[next] = temp.
 - 8. REAR = temp.

Dequeue Operation

- Let,
 - FRONT and REAR points to the front and rear of the QUEUE.
 - temp points to the element deleted from the front of the queue.
- 1. if (FRONT == NULL)
- 2. Print [Underflow]
- 3. else
- 4. Initialize a node pointer (temp) with FRONT.
- 5. if (FRONT == REAR)
- 6. FRONT = REAR = NULL
- 7. else
- 8. FRONT = FRONT[next]
- 9. Release the memory location pointed by temp.

Dynamic Linked List Implementation

```
1. struct node
2. { int data;
3.   struct node *next, *prev;
4. } *front, *rear;

5. void init()
6. { front = rear = NULL; }
```

Contd...

```
7. void enqueue(int num)
8. { struct node *temp = (struct node *) malloc (sizeof(struct node));
9.   temp -> data = num;
10.  if(front == NULL)
11.  {   temp -> prev = temp;
12.      temp -> next = temp;
13.      front = rear = temp; }
14.  else
15.  {   temp -> prev = rear;
16.      rear -> next = temp;
17.      temp -> next = front;
18.      front -> prev = temp;
19.      rear = temp; } }
```

Contd...

```
20. void dequeue()
21. {  if(front == NULL)
22.     printf("\nQueue is empty.\n");
23.     else
24.     {  struct node *temp = front;
25.         if (front == rear)
26.             front = rear = NULL;
27.         else
28.         {  rear -> next = front -> next;
29.             front = front -> next;
30.             front -> prev = rear;  }
31.     free(temp);    }    }
```

Contd...

```
33. void print()
34. { printf("\nfront --> ");
35.   if (front != NULL)
36.   { struct node *temp = front;
37.     while(temp != rear)
38.     { printf("%d --> ",temp->data);
39.       temp = temp->next; }
40.     printf("%d --> ",temp->data);
41.   } printf("rear\n");
42. }
```

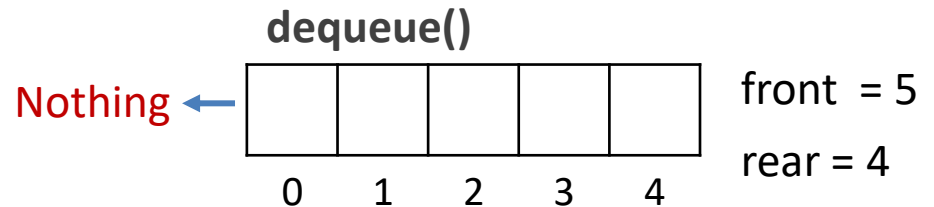
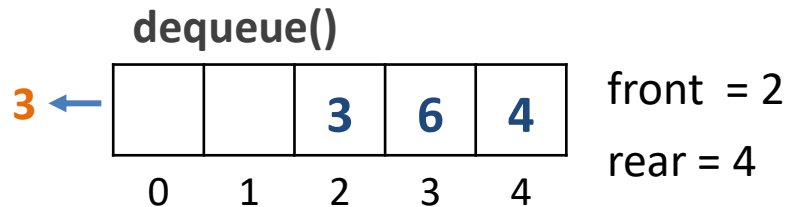
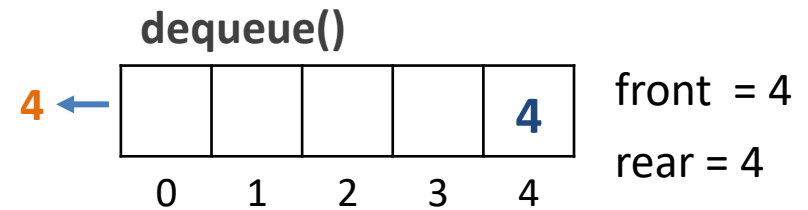
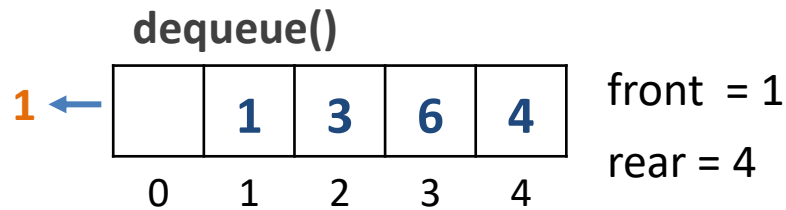
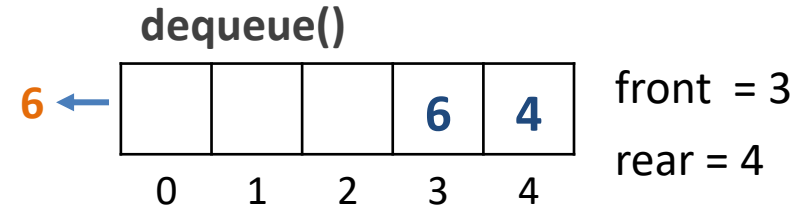
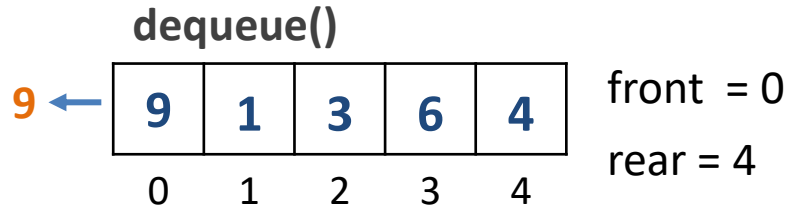
```
Current queue :
front --> 7 --> rear
Current queue :
front --> 7 --> 9 --> 3 --> 1 --> rear
```

```
43. int main ()
44. {   init();
45.     enqueue(5); enqueue(3);
46.     dequeue();
47.     enqueue(7);
48.     dequeue();
49.     printf("Current queue : ");
50.     print();
51.     enqueue(9); enqueue(3);
52.     enqueue(1);
53.     printf("Current queue : ");
54.     print();
55.     return 0;
56. }
```

Queues

Variants

Problem with Simple Queues



UNDERFLOW

The queue is empty, still no
element can be added as
REAR = N-1 (Queue Full)



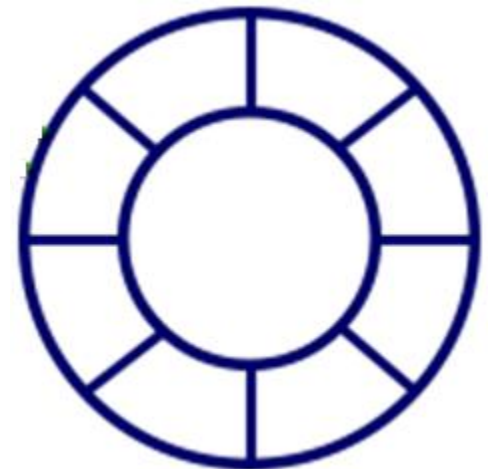
Circular Queues

- The front and rear ends of a queue are joined to make the queue circular.
- Also known as circular buffer, circular queue, cyclic buffer or ring buffer.
- Overflow

$\text{front} == (\text{rear} + 1) \% \text{MAXLEN}$

- Underflow

$(\text{front} == \text{rear}) \ \&\& \ (\text{rear} == -1)$



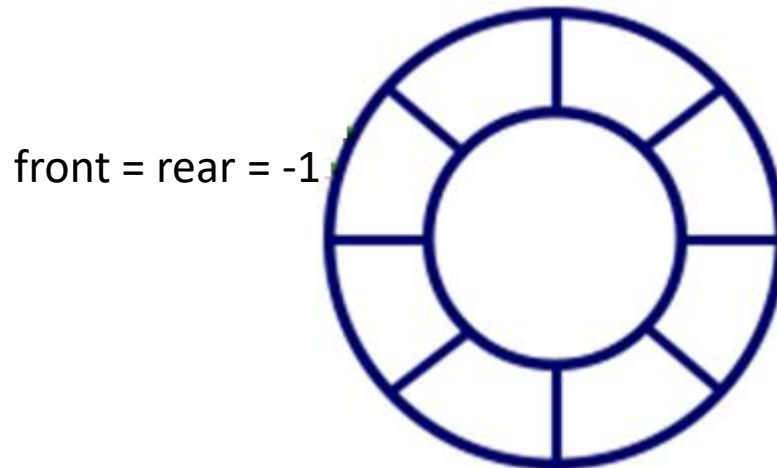
Contd...

- Overflow

$\text{front} == (\text{rear} + 1) \% 8$

- Underflow

$(\text{front} == \text{rear}) \ \&\& \ (\text{rear} == -1)$



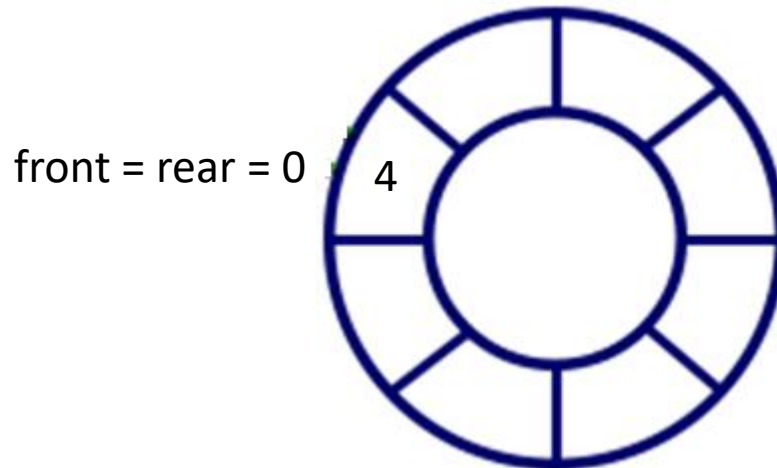
Contd...

- Overflow

$\text{front} == (\text{rear} + 1) \% 8$

- Underflow

$(\text{front} == \text{rear}) \ \&\& \ (\text{rear} == -1)$



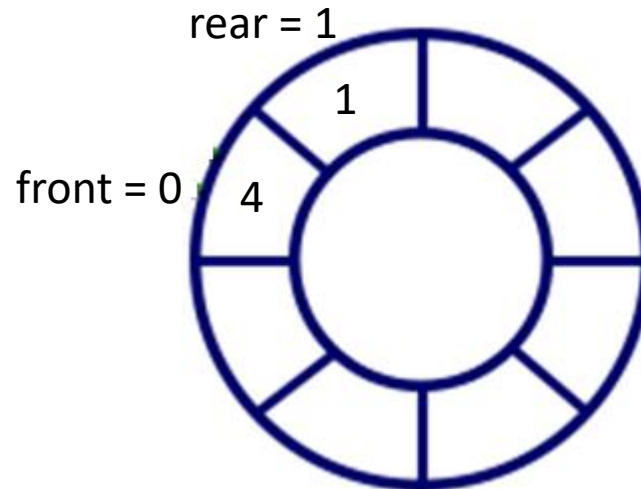
Contd...

- Overflow

$\text{front} == (\text{rear} + 1) \% 8$

- Underflow

$(\text{front} == \text{rear}) \ \&\& \ (\text{rear} == -1)$



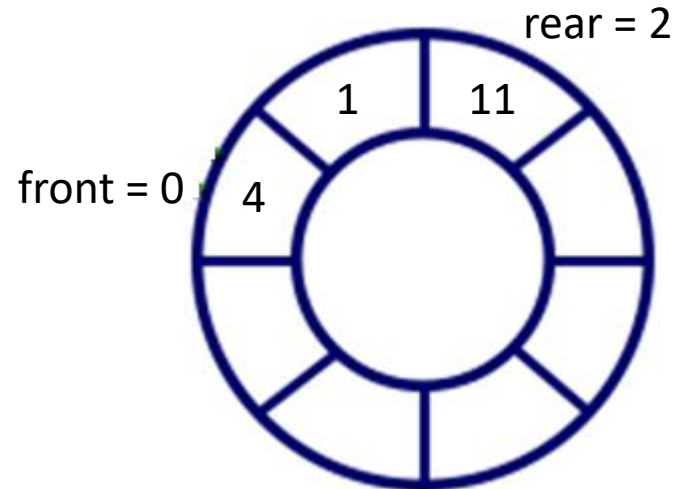
Contd...

- Overflow

$\text{front} == (\text{rear} + 1) \% 8$

- Underflow

$(\text{front} == \text{rear}) \ \&\& \ (\text{rear} == -1)$



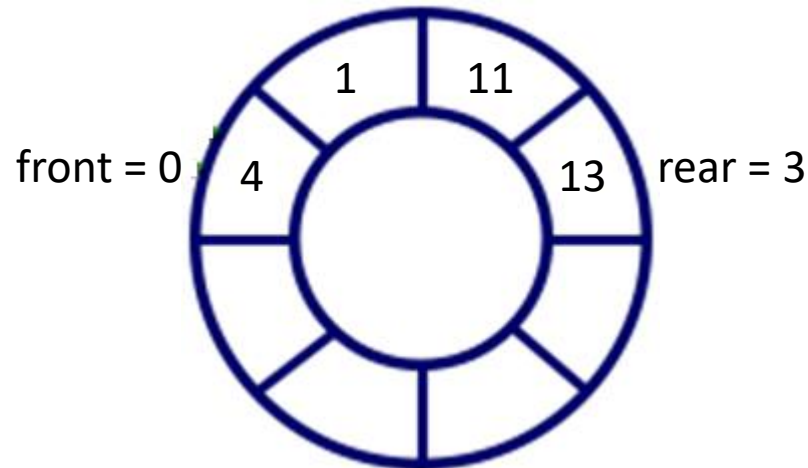
Contd...

- Overflow

$\text{front} == (\text{rear} + 1) \% 8$

- Underflow

$(\text{front} == \text{rear}) \ \&\& \ (\text{rear} == -1)$



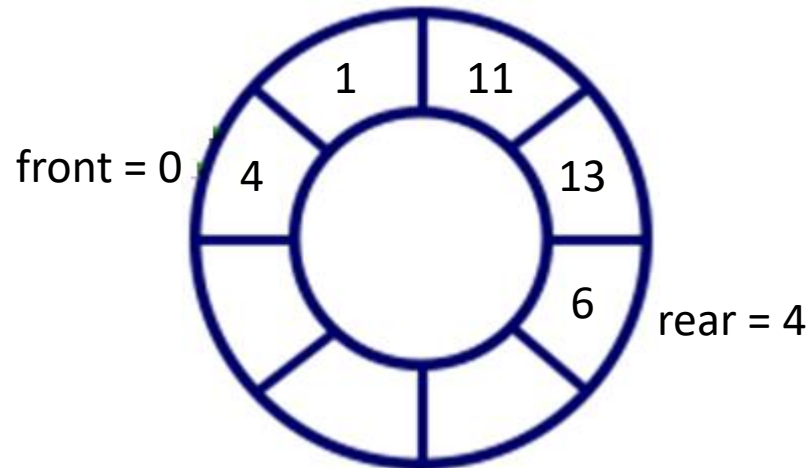
Contd...

- Overflow

$\text{front} == (\text{rear} + 1) \% 8$

- Underflow

$(\text{front} == \text{rear}) \ \&\& \ (\text{rear} == -1)$



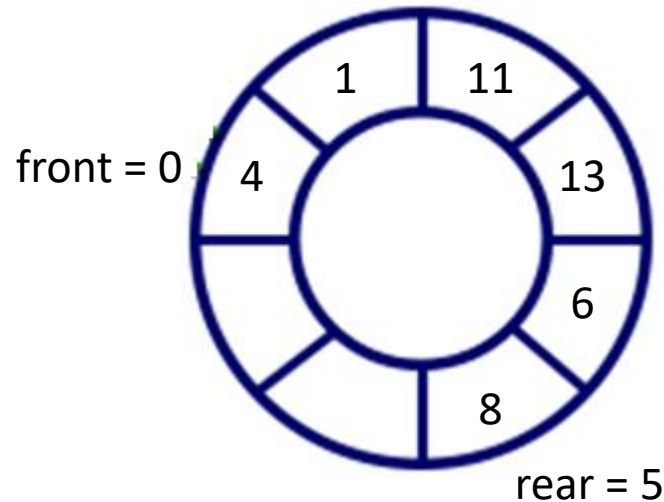
Contd...

- Overflow

$\text{front} == (\text{rear} + 1) \% 8$

- Underflow

$(\text{front} == \text{rear}) \ \&\& \ (\text{rear} == -1)$



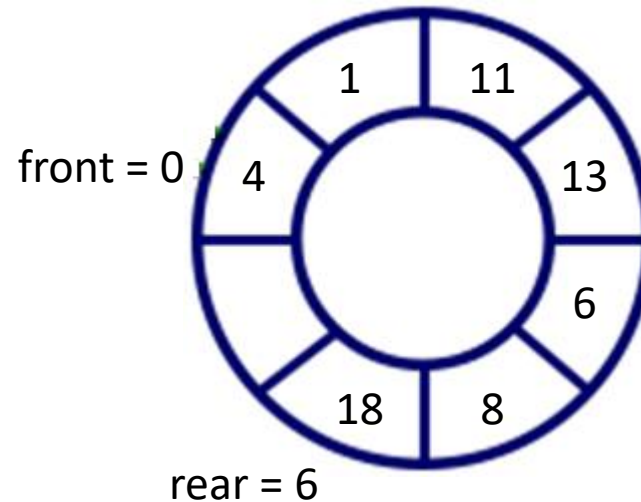
Contd...

- Overflow

$\text{front} == (\text{rear} + 1) \% 8$

- Underflow

$(\text{front} == \text{rear}) \ \&\& \ (\text{rear} == -1)$



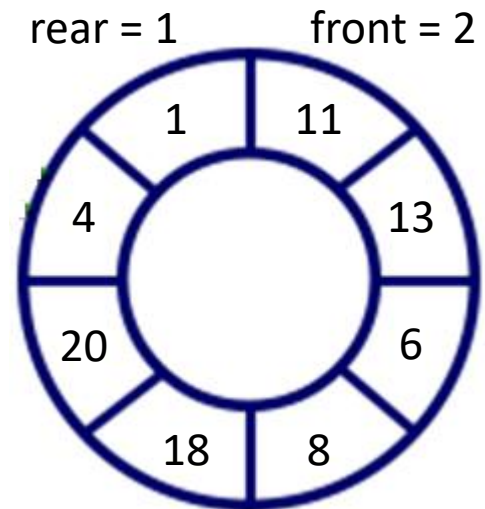
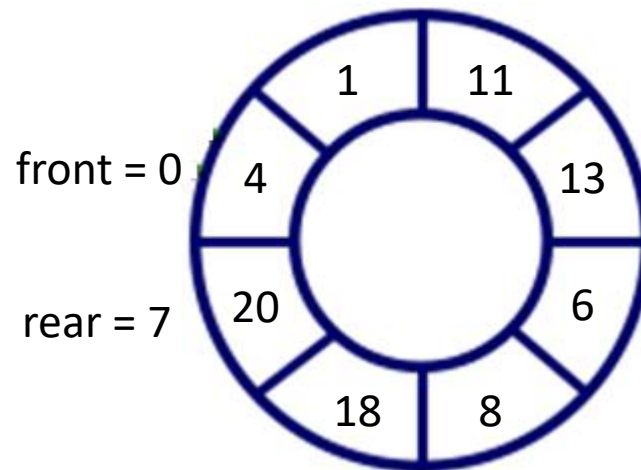
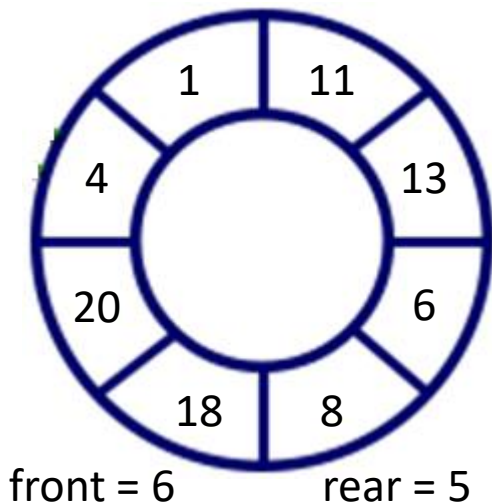
Contd...

- **Overflow**

$$\text{front} == (\text{rear} + 1) \% 8$$

- Underflow

$$(\text{front} == \text{rear}) \ \&\& \ (\text{rear} == -1)$$



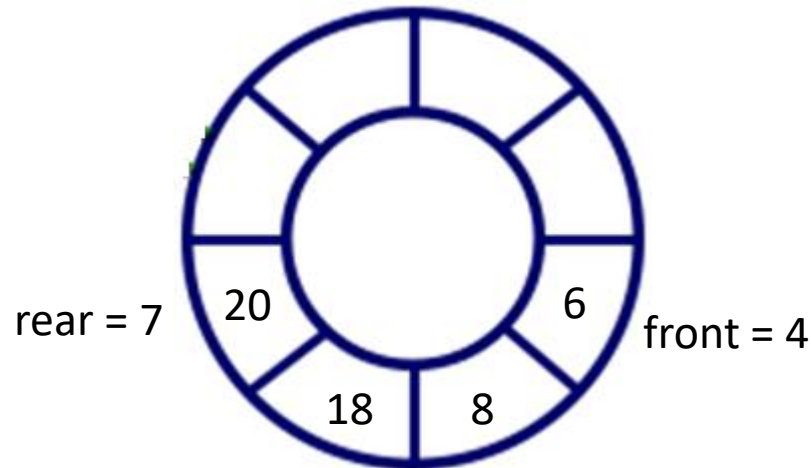
Contd...

- Overflow

$\text{front} == (\text{rear} + 1) \% 8$

- Underflow

$(\text{front} == \text{rear}) \ \&\& \ (\text{rear} == -1)$



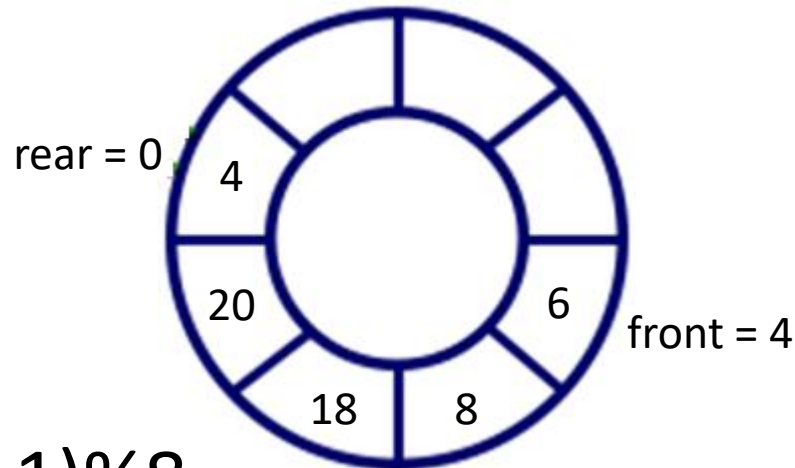
Contd...

- Overflow

$\text{front} == (\text{rear} + 1) \% 8$

- Underflow

$(\text{front} == \text{rear}) \ \&\& \ (\text{rear} == -1)$



$\text{REAR} = (\text{REAR} + 1) \% 8$

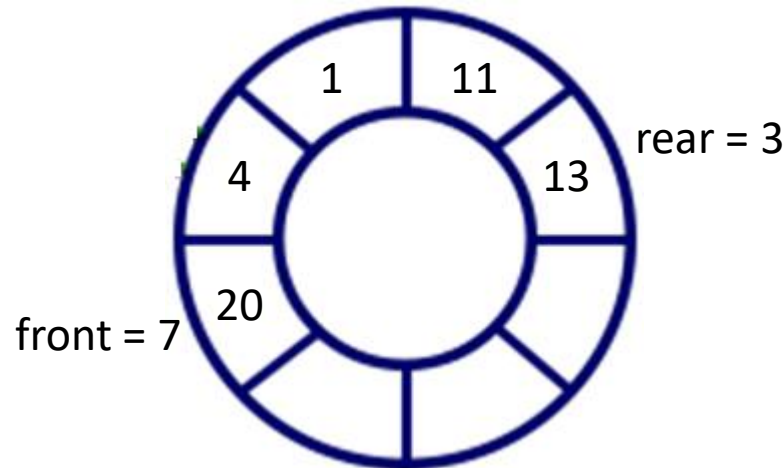
Contd...

- Overflow

$\text{front} == (\text{rear} + 1) \% 8$

- Underflow

$(\text{front} == \text{rear}) \ \&\& \ (\text{rear} == -1)$



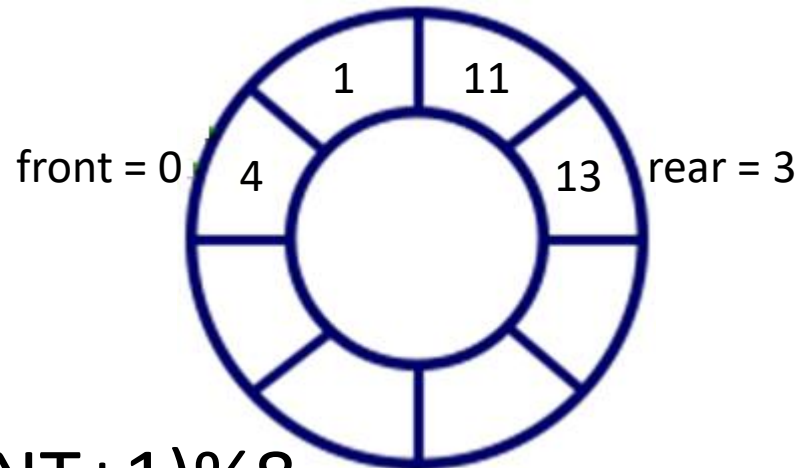
Contd...

- Overflow

$\text{front} == (\text{rear} + 1) \% 8$

- Underflow

$(\text{front} == \text{rear}) \ \&\& \ (\text{rear} == -1)$



$\text{FRONT} = (\text{FRONT} + 1) \% 8$

Enqueue Operation

- Let,
 - QUEUE be an array with MAX locations.
 - FRONT and REAR points to the front and rear of the QUEUE.
 - ITEM is the value to be inserted.
- 1. if $(\text{FRONT} == (\text{REAR} + 1) \% \text{MAX})$
- 2. Print [Overflow]
- 3. else
- 4. Set $\text{REAR} = (\text{REAR} + 1) \% \text{MAX}$
- 5. Set $\text{QUEUE}[\text{REAR}] = \text{element}$
- 6. If $(\text{FRONT} == -1)$
- 7. Set $\text{FRONT} = 0$

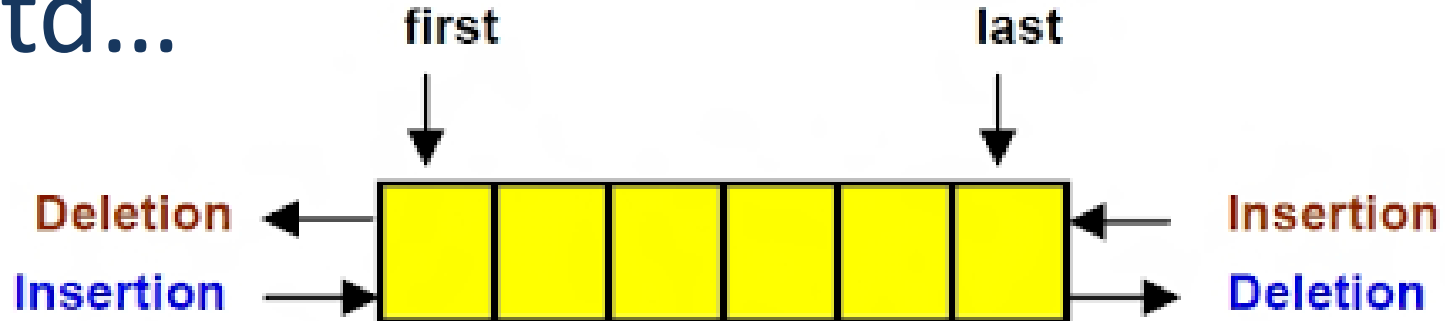
Dequeue Operation

- Let,
 - QUEUE be an array with MAX locations.
 - FRONT and REAR points to the front and rear of the QUEUE.
 - ITEM holds the value to be deleted.
- 1. if ((FRONT == REAR) && (REAR == -1))
- 2. Print [Underflow]
- 3. else
- 4. ITEM = Q[FRONT]
- 5. If (FRONT == REAR)
- 6. FRONT = REAR = -1
- 7. Else
- 8. FRONT = (FRONT + 1) % MAX

Deque

- **Double-ended queue.**
- Generalization of queue data structure.
- Elements can be added to or removed from either of the two ends.
- A hybrid linear structure that provides all the capabilities of stacks and queues in a single data structure.
- Does not require the LIFO and FIFO orderings.

Contd...



Types

- Input-restricted deque.
 - Deletion can be made from both ends, but insertion can be made at one end only.
- Output-restricted deque.
 - Insertion can be made at both ends, but deletion can be made from one end only.

Priority Queues

- Another variant of queue data structure.
- Each element has an associated priority.
- Insertion may be performed based on the priority.
- Deletion is performed based on the priority.
- Elements having the same priority are served or deleted according to first come first serve order.
- Two types:
 - Min-priority queues (Ascending priority queues)
 - Max-priority queues (Descending priority queues)

Implementation

- Array representation: Unordered and Ordered
- Linked-list representations: Unordered and Ordered
- Unordered does not consider priority during insertion, instead insertion takes place at the end.
- Ordered considers priority during insertion and inserts an element at correct place as per min or max priority.
- **Note**
 - Either insertion or deletion take linear time in the worst case.
 - Priority queues are often implemented with heaps.

Example

Element to be deleted is replaced with the last array element.

Operation	Argument	Return Value	Size	Contents	
				Unordered	Ordered (Ascending)
Insert	P		1	P	P
Insert	Q		2	P Q	P Q
Insert	E		3	P Q E	E P Q
Remove MAX		Q	2	P E	E P
Insert	X		3	P E X	E P X
Insert	A		4	P E X A	A E P X
Insert	M		5	P E X A M	A E M P X
Remove MIN		A	4	P E X M	E M P X