3 - digit Aleithmetic 元= 53+(5+57)=1.73要+(2.24+2.65)=6.62 x= 53+(55+57)= 6.61387009 (Exact Value) A.E= |x-x* |= 000612991 R.E= 121-21 = 0.0009 2682648987 (b) x=(121-0.327)-119=1.673 x = (121-0.327)-119=121-119=1-673 2 -- A-E= | 21-21 = 0.327 R.E = 121-21 = 0-19545726 (c) $x = -10n + 6e - \frac{3}{62} = -18 - 18462266$ nt = (-10n) + (6e) - (3/62) = (-31.4) + (16.3) - (0.05) = -15.2 A.E = |x-xx | = 0.0454 R.E = 12-2 /11 = 0.002995 $\mathcal{K} = \frac{\Lambda - 22/7}{1/12} = -0.021496318$ 20.0588 = 3.14 - 3.14 = 0.00 1x-xx = 0.021496318 = 0.0218 [n] = 0.021496318 = 1.0002 (4) Percentage event in A = Absolute event in A x 100 = 84 ×100 =0.2 $\Rightarrow SA = \frac{1}{240}A = \frac{1}{240}(nx^2)$ Now bercentage event in a = Sh x100 SA= JA x 82 Zn SA = SA ZA SA = loo x dh = loo x SA DA/Dh

 $\frac{21}{2} + \frac{123}{4} \times -\frac{1}{4} = 0$ => Que N=1, b=123, C=-1 Using 4-digit Abuithmetic, a=0.3333, b=30.75, c=-0.1667 71 = -b+ 562-4ac Jow J 62-4ac = J (30.75)(30.75) - (4.000)(0.3333) (-0.1667) $=\sqrt{945.6+0.2222}$ = $\sqrt{945.8}$ = 30.75 $\chi_1^* = -30.75 + 30.75 = 0.000$ (2.000)(0.33333) $2a = -b - \sqrt{b^2 - 4ac} = -30.75 - 30.75 = -61.50 = -92.26$ 2a = (2.000)(0.3333) = 0.6666Exact Solutions $\frac{\pi^2 + 123}{3} \pi - \frac{1}{6} = 0 = 7 + 369\pi - 2 = 0$ $\pi = -369 \pm \sqrt{(369)^2 - 4(4)(-2)} = -369 \pm \sqrt{136193}$ = -369 ± 369.0433579 = 0.005419735788, -92.25541974 ·· 71 = 0:00 5419 735788 25 = -92.25541974 A.E In M1 | m-M1 = 0.000001264212 $\frac{|x-x_1^*|}{|x_1|} = \frac{0.00000|264212}{0.005419735788} = 0.000233260817$ R.E du 21, A:E in 1/2 1/2- 1/2 | = 0:00458026 REinx 12-72*1 = 0.00458026 = 0.0000496475980 92-25541974 Rewrite M_1 , $M_1 = -b + \sqrt{b^2 - 4ac} = -b + \sqrt{b^2 - 4ac} \times \frac{-b - \sqrt{b^2 - 4ac}}{2a}$ $= \frac{-2C}{6+ (b^2+4a)} = \frac{(-2.000)(+0.1667)}{30.75 + 30.75} = \frac{0.3334}{61.50} = 0.00542438$

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8) $f(x) = \frac{1}{1-3 \times 2}$, $f'(x) = \frac{6 \times 1}{(1-3 \times 2)^2}$ Now At x = 0.577 and 3-digit chopping abithmetic, $f'(x) = \frac{6}{5} \frac{(6.00)(0.577)}{(0.00)(0.577)^2} = \frac{3.46}{(1.00-(3.00)(0.332))^2}$ $= \frac{3.46}{(1.00-0.996)^2} = \frac{3.46}{(0.004)^2} = \frac{3.46}{0.00} \rightarrow \infty$

```
At x =0.577 and 4-digit chopping abouthmetic
  f'(n) = 62 = (6.00)(0.5770)^{2}
(1-32)^{2} = (6.00)(0.5770)^{2}
     = 3.462
[1.000 - (3.000)(0.3329)] = [1.000 - 0.9987]<sup>2</sup>
        = 3.462 = 3.462 -> 0

\[ \text{0.0013} \]^2 = \[ \text{0.000} \]
  Yes, we have difficulties in evaluating the of (x) at x =0.577 using 3 and 4- digit anithmetic with chopping,
9) l=3am, w=4cm, h=5am
                                    { l > 2.5 => l=2.51 => l=3
   2.5 < 1 < 3.5
                                     1 < 3.5 => 1 = 3-4999 => 1=3
                                      10 73.5 => w=3.51.. > w=4
    3.5 < w < 4.5
                                      W<4.5 -> W= 4.4999 => W=4
     4.5 < h < 5.5
                                       h74.5 => h=4.51 => h=5
  Mohnue = luch
                                      -hcs.5=> h=5.4999 => h=5
· (2.5)(3.5)(4.5) < lwh < (3.5)(4.5)(5.5)
      39.375 < Volume < 86.625
      Subjuce Asea = 2[lw+lh+wh]
Now (2.5) (3.5) < lw < (3.5) (4.5) => 8.75 < lw < 15.75 =
         (2.5) (4.5) < lh < (3.5) (5.5) => 11.25 < lh < 19.25
          B.5) (4.5) < wh< (4.4) (5.5) => 15.75 < wh< 24.75
  Now 28 (8.75+11.25+15.75) < lw+lh+wh< 15.75+19.25+24.75
                      35.75 < lw+lh+wh< 59.75
         => 2(35.75) < 2(lw+lh+wh) < 2 (59.75)
                   171.5 < Surface ALEA < 119.5
: Best Lower bound for Volume = 39.371/
11 Upper 11 11 = 86.625
    " Lower " Surface Area = 71.5
    " Upper " " = 119.5
```

(a) Now
$$x = 26 - \frac{(x_1 - 26)y_0}{y_1 - y_0}$$
 $\frac{1}{y_1 - y_0}$ $\frac{1}{y_1 - y_0}$

(i) (a) Free & (2) (2) (2) (2) For each i, there are i multiplications in 1 . E have n iterations => Total Multiplication = \(\frac{2}{12} \) =1+2+3+...+n=n(m+1)There are i-1 additions for each i in 1) -. \(\left(\reft(\left(\left(\left(\left(\left(\left(\left(\left(\reft(\left(\left(\left(\left(\left(\left(\reft(\left(\left(\left(\left(\reft(\re Also (n-1) more additions (: of) - Total additions = M(m-1) + m-1 = (m+2)(m-1) (b) Equivalent form S (ai & bi) for each i, i-1 addition i. Total additions = $\sum_{i=1}^{m} (i-1) + (m-1) = (m+2) (m-1)$ none additions Total Multiplications = Z(1) = n .. New form $\sum_{i=1}^{N} (a_i \leq b_j)$ heduces no. of multiplication. P(n) = ann + an - 2 + ... + an x + ao (2) Algolithus 1(no) = au + no (a, + no (a2 + ··· + no (an, + no an))) Input: n, ao, a, ..., an, xo output y= P(no) Step1: Set y = an Step 2: for i= n-1, n-2, --, 0 Set y = >6 y + ai Step3: Output (y) STOP.

Suput: M, No, N1, ..., Nm, X

output: P

Step 1: Set P= x-xo; i=1

Step 2: while P to and U ≤ m,

Set P=P. (x-xi)

i=i+1

Step 3: Output (P)

STOP