

Type Inference Rules For Container Types in CCL

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Abstract

We present the type inference rules introducing the notion of container types in the CCL programming language. This redesign required a number of substantial changes to the major aspects of the language, including the type system, the syntax, and the definitional semantics.

The $<::$ Relation

$$\frac{}{\mathcal{Box}(\mathit{Triv}) <:: \mathcal{Box}(\mathit{Triv})} \quad (1)$$

$$\frac{}{\mathcal{Box}(\mathit{Int}) <:: \mathcal{Box}(\mathit{Int})} \quad (2)$$

$$\frac{\mathcal{T} <:: \mathcal{Box}(\mathcal{U})}{\mathcal{T} <:: \mathcal{IBox}(\mathcal{U})} \quad (3)$$

$$\frac{\mathcal{Box}(\mathcal{T}) <:: \mathcal{IBox}(\mathcal{U})}{\mathcal{IBox}(\mathcal{T}) <:: \mathcal{IBox}(\mathcal{U})} \quad (4)$$

$$\frac{\mathcal{T} <:: \mathcal{U}}{\mathcal{Ref} \ \mathcal{T} <: \mathcal{Ref} \ \mathcal{U}} \quad (5)$$

$<::$ operator specifies the subtype relationship between two container types.

$\mathcal{T} <:: \mathcal{U}$ is read as “container type \mathcal{T} is a subtype of the container type \mathcal{U} .”

The $:$ and $::$ Relations

$$\frac{\text{triv } x}{x : \text{Triv}} \quad (6)$$

$$\frac{\text{triv } x}{x :: \text{Box}(\text{Triv})} \quad (7)$$

$$\frac{\text{int } x}{x : \text{Int}} \quad (8)$$

$$\frac{\text{int } x}{x :: \text{Box}(\text{Int})} \quad (9)$$

$$\frac{x :: \text{Box}(\mathcal{T})}{\text{immut } x :: \text{IBox}(\mathcal{T})} \quad (10)$$

$$\frac{x :: \mathcal{T}}{\text{ref } x :: \text{Box}(\text{Ref}(\mathcal{T}))} \quad (11)$$

$$\frac{x :: \mathcal{T}}{\& x : \text{Ref}(\mathcal{T})} \quad (12)$$

$$\frac{x : \text{Ref}(\text{Box}(\mathcal{T}))}{x @ : \mathcal{T}} \quad (13)$$

$$\frac{x : \text{Ref}(\text{IBox}(\mathcal{T}))}{x @ : \mathcal{T}} \quad (14)$$

$$\frac{x : \text{Ref}(\mathcal{T})}{x @ : \mathcal{T}} \quad (15)$$

$$\frac{x :: \mathcal{T} \quad y :: \mathcal{U} \quad \mathcal{T} <:: \mathcal{U} \quad x : \mathcal{V}}{x := y : \mathcal{V}} \quad (16)$$

$x : \mathcal{T}$ is read as “expression x is of type \mathcal{T} and is in an *r-context*.”

$x :: \mathcal{T}$ is read as “variable x is of type \mathcal{T} and is in an *l-context*.”

The operators could also be referred to as the “r-type of” and “l-type of” operators.

l-context denotes everything that is *assignable* (indicated as a storable memory). r-context, on the other hand, denotes everything that is *expressible* (can be produced by an expression).

There is no r-value (e.g. expression) of the type $\text{Box}(\mathcal{T})$.

For now, we omit rules for *Con* types as they only operate on r-values.

For now, we omit rules for *Fun* types as they only accept r-values. Any variable and/or primitive type has both r-value and l-value (when it comes to primitive types, only r-value). In all cases, the r-value part of the actual parameter is passed when the function is being called.

Resulting Relationships (A Short List)

<code>int i</code>	$\rightarrow i$	$\rightarrow \mathcal{B}\chi(Int)$
<code>immut int ii</code>	$\rightarrow ii$	$\rightarrow I\mathcal{B}\chi(Int)$
<code>ref int ri</code>	$\rightarrow ri$	$\rightarrow \mathcal{B}\chi(\mathcal{R}ef(\mathcal{B}\chi(Int)))$
<code>immut ref int iri</code>	$\rightarrow iri$	$\rightarrow I\mathcal{B}\chi(\mathcal{R}ef(\mathcal{B}\chi(Int)))$
<code>ref immut int rii</code>	$\rightarrow rii$	$\rightarrow \mathcal{B}\chi(\mathcal{R}ef((Immut(\mathcal{B}\chi(Int))))))$
<code>immut ref immut int irii</code>	$\rightarrow irii$	$\rightarrow I\mathcal{B}\chi(\mathcal{R}ef((Immut(\mathcal{B}\chi(Int))))))$

Type $Immut \mathcal{B}\chi(Immut \mathcal{B}\chi(Int))$ cannot exist. Nested $\mathcal{B}\chi$ types are only possible when there is at least one $\mathcal{R}ef$ type.

$$\mathcal{B}\chi(Triv) <: \mathcal{B}\chi(Triv)$$

$$\mathcal{B}\chi(Int) <: \mathcal{B}\chi(Int)$$

$$\mathcal{B}\chi(Triv) <: I\mathcal{B}\chi(Triv)$$

$$\mathcal{B}\chi(Int) <: I\mathcal{B}\chi(Int)$$

$$I\mathcal{B}\chi(Triv) <: I\mathcal{B}\chi(Triv)$$

$$I\mathcal{B}\chi(Int) <: I\mathcal{B}\chi(Int)$$

All the rules above should work with $\mathcal{R}ef$ types in the similar manner:

$$\mathcal{R}ef(\mathcal{B}\chi(Triv)) <: \mathcal{R}ef(\mathcal{B}\chi(Triv))$$

$$\mathcal{R}ef(\mathcal{B}\chi(Int)) <: \mathcal{R}ef(\mathcal{B}\chi(Int))$$

$$\mathcal{R}ef(\mathcal{B}\chi(Triv)) <: \mathcal{R}ef(I\mathcal{B}\chi(Triv))$$

$$\mathcal{R}ef(\mathcal{B}\chi(Int)) <: \mathcal{R}ef(I\mathcal{B}\chi(Int))$$

$$\mathcal{R}ef(I\mathcal{B}\chi(Triv)) <: \mathcal{R}ef(I\mathcal{B}\chi(Triv))$$

$$\mathcal{R}ef(I\mathcal{B}\chi(Int)) <: \mathcal{R}ef(I\mathcal{B}\chi(Int))$$