Final Year B.Tech. (CSE) - VII [2024-25]

6CS451: Cryptography and Network Security Lab (C&NS Lab)

Date: 26/08/2024

Assignment 7

PRN: 21510017 Name: Onkar Anand Yemul

1. Implementation of RSA Algorithm

Ans:

The RSA algorithm is one of the first public-key cryptosystems and is widely used for secure data transmission. It is an asymmetric cryptographic algorithm, meaning it uses a pair of keys: a public key for encryption and a private key for decryption. It relies on the mathematical properties of prime numbers.

How RSA Works:

1. Key Generation:

- Choose two large prime numbers p and q.
- \circ Compute n = p * q.
- Compute the totient $\phi(n) = (p-1) * (q-1)$.
- Choose an encryption key e such that $1 < e < \phi(n)$ and gcd(e, $\phi(n)$) = 1. The integer e is the public key exponent.
- o Calculate the decryption key d such that d * e \equiv 1 (mod φ (n)). The integer d is the private key exponent.

2. Encryption:

- The public key is (n, e).
- o Given a plaintext message M, the ciphertext C is computed as: $C = M \land e \mod n$.

3. **Decryption**:

- o The private key is (n, d).
- Given a ciphertext C, the plaintext M is recovered as:

```
M = C \wedge d \mod n
```

To implement the RSA algorithm using large prime numbers with 2048 bits and converting plaintext into numbers, we'll use the Crypto library in Python, which provides the necessary tools to handle such large prime numbers and perform RSA encryption and decryption.

The large primes and the strong key sizes make RSA secure against most attacks when implemented correctly.

Python Code:

```
import random
from sympy import isprime, mod_inverse
def generate prime candidate(length):
    """Generate an odd integer randomly."""
    p = random.getrandbits(length)
    # Ensure p is odd
    p = (1 \ll length - 1) \mid 1
    return p
def generate_prime_number(length):
    """Generate a prime number."""
    p = 4
    while not isprime(p):
        p = generate_prime_candidate(length)
    return p
def generate_keypair(keysize):
    """Generate RSA public and private keys."""
    # Generate two large primes p and q
    p = generate prime number(keysize)
```

```
q = generate prime number(keysize)
    print("\np: ", p)
    print("\nq: ", q)
    \# Compute n = p * q
    n = p * q
    # Compute Euler's Totient \phi(n) = (p-1)*(q-1)
    phi = (p - 1) * (q - 1)
    # Choose an integer e such that 1 < e < phi(n) and
gcd(e, phi(n)) = 1
    e = random.randrange(2, phi)
    g = gcd(e, phi)
    while g != 1:
        e = random.randrange(2, phi)
        g = gcd(e, phi)
    # Compute d, the modular inverse of e
    d = mod_inverse(e, phi)
    # Public key (e, n) and Private key (d, n)
    return ((e, n), (d, n))
def gcd(a, b):
    """Compute the greatest common divisor using Euclid's
algorithm."""
    while b != 0:
        a, b = b, a \% b
    return a
def encrypt(public key, plaintext):
    """Encrypt plaintext using the public key."""
    e, n = public key
    cipher = [pow(ord(char), e, n) for char in plaintext]
    return cipher
```

```
def decrypt(private key, ciphertext):
    """Decrypt ciphertext using the private key."""
    d, n = private key
    plain = [chr(pow(char, d, n)) for char in ciphertext]
    return ''.join(plain)
def main():
    """Run RSA algorithm with a menu-driven interface."""
    print("RSA Encryption/Decryption")
    # Take keysize as input from the user
    keysize = int(input("Enter key size (e.g., 1024, 2048):
"))
    # Generate public and private keys
    public key, private key = generate keypair(keysize)
    # Menu-driven system
    while True:
        print("\nMenu:")
        print("1. Show Public Key")
        print("2. Show Private Key")
        print("3. Enter a message to encrypt")
        print("4. Enter the encrypted text to decrypt")
        print("5. Exit")
        choice = input("Choose an option: ")
        if choice == '1':
            print(f"\nPublic key: {public_key}")
        elif choice == '2':
            print(f"\nPrivate key: {private key}")
        elif choice == '3':
            plaintext = input("\nEnter a message to encrypt:
            encrypted msg = encrypt(public key, plaintext)
```

```
print(f"\nEncrypted message: {encrypted msg}")
        elif choice == '4':
            encrypted_text = input("\nEnter the encrypted
text as a list of integers (e.g., [123, 456, ...]): ")
            try:
                # Convert the input string to a list of
integers
                encrypted_msg = [int(x) for x in]
encrypted_text.strip('[]').split(',')]
                decrypted_msg = decrypt(private_key,
encrypted_msg)
                print(f"\nDecrypted message:
{decrypted_msg}")
            except ValueError:
                print("Invalid input. Please provide the
encrypted text in the correct format.")
        elif choice == '5':
            print("Exiting...")
            break
        else:
            print("Invalid choice. Please try again.")
if name == " main ":
    main()
```

Output:

For Keysize = 2048:

problems output debug console $\underline{\mathsf{TERMINAL}}$ ports sql console \cdots $\underline{\triangleright}$ bash + \vee $\underline{\square}$ $\hat{\underline{\square}}$ \cdots \wedge \times

Onkar@LAPTOP-QRVSMK3C MINGW64 ~/Downloads/CNS Lab/Assignment 7

TERMINAL

TERMINAL

- \$ python rsa.py RSA Encryption/Decryption
- $p:\ \ 281906190582007192029765079276311565367397508427667048562319679283092016517352633083446203186342143549855783\\6080468807013767346958530944392339226154072595103524381415434843986849703825169611900936023052777716617462837577\\6613235243099543650163771789372583397314034518021824740578399791764078827040395781677457846437275487190995130551\\7273507056462482268184741268981740346737078725006858209667556546182212392899035976637602890029237565395119353084\\6661720343145114968222199870669210473147756731955607055564002218374315237586856959392765508140117407962631015943\\6509527259115164671203698923251367166220212601082275322534179$
- $q\colon\ 202303892491396233948343193134824573153477378015812034171853366824244635030890363647337454208999198024219131\\9187786350052463996457037038432181975080551510651109811651737681459820917046972666272645951853787907452060836183\\9319149065598563063887260011359218867961100215523348452156143427046450175965127862864117586143138389454522146327\\0451672757060090936268029233308810442059094171581281570316232833954809470064440540981318587331939466757257447164\\9772511363393956453895144446284731789412582580140137396525404401410158653077830632985641075204813132249622821346\\8618623795770314120612409499041935323061889732084110295536711$

∑ bash + ∨ □ · · · · · ×

∑ bash + ∨ □ ଢ ··· ^

Public key: (401255109074108980465921566640954685015789317040706324808294074793760149164137220756805659517308539 87962702766029, 570307196721614404983350840666356675627763256679820284549479911790970876115008475746256250223776

67013905446745269) Private key: (12568402183863569951896883230387076856416531860297403273406475503641345091092906621850466896701745 946515915266009, 57030719672161440498335084066635667562776325667982028454947991179097087611500847574625625022377



PROBLEMS OUTPUT DEBUG CONSOLE **TERMINAL** PORTS SQL CONSOLE ··· \(\sum_{\text{bash}} + \sim_{\text{li}} \frac{\text{ii}}{\text{li}} \cdots \sim_{\text{v}} \times \text{X}

 $6794302441260677255233484221102154386552818117109352807086485155278826834068716081348587537227994497193744231426\\ 7165084088703923535851860895753533434065080002383420478061520457019534053543276053565687962153328150120407485575\\ 7600898124413861455813151442882504905994635606789863454802496838436789210452784261428607479700845876196245157417\\ 0396991457178543029427276674076576720085511148431334640133149188918497523521468733306345873312114175427462328232\\ 0434883632608822683462013847253727501656879678455563175767620886746672568654197654288155113684530445972212518564\\ 2482505377068958275979061572601311089600391888309139060916160989670240852260548553559930114541332561344760501121\\ 7206116813148894527923859325824761143887574989242047517204922525955278176869094151670039164495073739396167640634\\ 3314763239473090019865705761528216139300297304281620687461341919910435075878455608838987052156986544555634921139\\ 067013905446745269)$

Enter a message to encrypt: We will meet tomorrow at 5pm at canteen.

TERMINAL

Encrypted message: [51240574717994847921488533151575126509457154796299845527107583540480756402690276029528277377 872883489696078139940, 27635090608112552992037187576122193821105195369672590996782823541108734144696845990626012

4509256603909973859995279717155382536306925704599351417810455855828678716644556816244268985677950687921216605337
4299242496072611624947361894130827882528057067123152785234097698096379546037989986332912975770903717632823142331
3560765816500757688740921001040736471281625346617973616385433879610866517280031283010125131455390660162721227309
2334394713328617606302804005819596050895569983314342312737970325239196546613923334214509257510758609134717750158
651703719871852585339278783565700442101287188676101942482842821068056455519685099170703324692845011594398743062

∑ bash + ∨ □ · · · · · ×

TERMINAL ∑ bash + ∨ □ • • · · · × 033719251332295433791504, 23705036455489225849849620179448513572350923525671099703765122632303400341950610664385 390970794536873163763570784, 55225466416240428439264103353995996065152439023630232771594383579607956153852159964

∑ bash + ∨ □ 🛍 ··· ∧ **TERMINAL** 012538141835731664941168430932 . 10563710283181769993643630292710553153977787763781062362663724928805445142603780 610704794668565231284120004666390, 13640175622460002922004732570640946224061761655339613211815867225859848188713

TERMINAL $433737567357161788429753521027903721,\ 13640175622460002922004732570640946224061761655339613211815867225859848188$ 077433737567357161788429753521027903721, 23705036455489225849849620179448513572350923525671099703765122632303400

∑ bash + ∨ □ 🛍 ··· ^ × TERMINAL 82932475958713645, 421368742534090852407339110525209976020477230572793808561639218983973557774442664299570749841 91085756018289911292] Decrypted message: We will meet tomorrow at 5pm at canteen. Onkar@LAPTOP-QRVSMK3C MINGW64 ~/Downloads/CNS Lab/Assignment 7

For Keysize = 16:

Onkar@LAPTOP-QRVSMK3C MINGW64 ~/Downloads/CNS Lab/Assignment 7 (main)

\$ python rsa.py
RSA Encryption/Decryption
Enter key size (e.g., 1024, 2048): 16

p: 42899

q: 63367

Menu:
1. Show Public Key
2. Show Private Key
3. Enter a message to encrypt
4. Enter the encrypted text to decrypt
5. Exit
Choose an option: 1

Public key: (899822465, 2718380933)

```
Menu:
1. Show Public Key
2. Show Private Key
3. Enter a message to encrypt
4. Enter the encrypted text to decrypt
5. Exit
Choose an option: 2
Private key: (2580611021, 2718380933)
1. Show Public Key
2. Show Private Key
3. Enter a message to encrypt
4. Enter the encrypted text to decrypt
5. Exit
Choose an option: 3
Enter a message to encrypt: We will cancel their order.
Encrypted message: [580383014, 1674094199, 1608554978, 1548141833, 2147988317, 2113262506, 21132625
06, 1608554978, 539500294, 324068690, 1171192377, 539500294, 1674094199, 2113262506, 1608554978, 12
31643643, 606790713, 1674094199, 2147988317, 1675740237, 1608554978, 2009642701, 1675740237, 191341
3955, 1674094199, 1675740237, 2518779932]
```

```
Menu:
1. Show Public Key
2. Show Private Key
3. Enter a message to encrypt
4. Enter the encrypted text to decrypt
5. Exit
Choose an option: 4
Enter the encrypted text as a list of integers (e.g., [123, 456, ...]): [580383014, 1674094199, 160
8554978, 1548141833, 2147988317, 2113262506, 2113262506, 1608554978, 539500294, 324068690, 11711923
77, 539500294, 1674094199, 2113262506, 1608554978, 1231643643, 606790713, 1674094199, 2147988317, 1
675740237, 1608554978, 2009642701, 1675740237, 1913413955, 1674094199, 1675740237, 2518779932]
Decrypted message: We will cancel their order.
Menu:
1. Show Public Key
2. Show Private Key
3. Enter a message to encrypt
4. Enter the encrypted text to decrypt
5. Exit
Choose an option: 5
Exiting...
Onkar@LAPTOP-QRVSMK3C MINGW64 ~/Downloads/CNS Lab/Assignment 7 (main)
```

Practical Applications of RSA

Secure Communication: Encrypting emails and messages.

- **Digital Signatures**: Verifying the authenticity of a message or document.
- **Key Exchange**: Securely exchanging keys for symmetric encryption algorithms.

RSA is widely used in various security protocols, including SSL/TLS for secure internet communications.

RSA ensures security through the difficulty of factoring large numbers. It is commonly used for securing sensitive data, digital signatures, and in SSL/TLS protocols.

Final Year B.Tech. (CSE) - VII [2024-25]

6CS451: Cryptography and Network Security Lab (C&NS Lab)

Date: 26/08/2024

Assignment 8

PRN: 21510017 Name: Onkar Anand Yemul

1. Implement the Diffie-Hellman Key Exchange algorithm for a given problem

Ans:

The Diffie-Hellman Key Exchange is a cryptographic algorithm that allows two parties to securely share a secret key over a public channel. This shared key can then be used for encrypted communication. The algorithm allows two parties to generate a shared secret key that can be used for subsequent encryption and decryption, even if the exchange itself is observed by an eavesdropper.

How Diffie-Hellman Works:

1. Public Parameters:

 Both parties agree on a large prime number p and a base g (a primitive root modulo p).

2. Key Exchange Process:

- Party A selects a private key 'a' and computes A = g^ a mod p, then sends A to Party B.
- o **Party B** selects a private key 'b' and computes $B = g \land b \mod p$, then sends B to Party A.

3. Shared Secret:

- o **Party A** computes the shared secret as $S = B^{\Lambda}$ a mod p.
- o **Party B** computes the shared secret as $S = A^h$ b mod p.

Since both calculations result in the same value, S becomes the shared secret key, even though an eavesdropper only knows p, q, A, and B.

The Diffie-Hellman algorithm securely establishes a shared secret key without transmitting it directly, making it fundamental for secure communications in protocols like SSL/TLS.

To implement the Diffie-Hellman Key Exchange algorithm for client-server communication across two different machines, we will create two Python programs: one for the client and one for the server. The server will generate its public key and share it with the client, and vice versa. Both will then calculate the shared secret key independently.

Python Code:

Client-side program:

```
import socket
import random

def generate_private_key(p):
    """Generate a private key."""
    private_key = random.randint(2, p - 2)
    print(f"\nGenerated Private Key: {private_key}")
    return private_key

def calculate_public_key(g, private_key, p):
    """Calculate the public key."""
    public_key = pow(g, private_key, p)
    print(f"\nCalculated Public Key: {public_key}")
    return public_key

def calculate_shared_secret(public_key, private_key, p):
    """Calculate the shared secret."""
    shared_secret = pow(public_key, private_key, p)
```

```
return shared secret
def start client(server host='localhost', server port=5000):
   # Take p and g as user input
    p = int(input("\nEnter a prime number (p): "))
    g = int(input("\nEnter a primitive root (g): "))
    # Generate client's private and public keys
    private key = generate private key(p)
    public_key = calculate_public_key(g, private_key, p)
    # Create client socket
    client_socket = socket.socket(socket.AF_INET,
socket.SOCK STREAM)
    client socket.connect((server host, server port))
    # Receive the server's public key
    server public key =
int(client socket.recv(1024).decode())
    print(f"\nReceived Server's Public Key:
{server public key}")
    # Send the client's public key to the server
    client_socket.sendall(str(public_key).encode())
    # Calculate the shared secret
    shared secret =
calculate shared secret(server public key, private key, p)
    print(f"\nShared Secret (Client): {shared secret}")
    client socket.close()
if name == " main ":
    start client()
```

<u>Server-side program:</u>

```
import socket
import random
def generate private key(p):
    """Generate a private key."""
    private key = random.randint(2, p - 2)
    print(f"\nGenerated Server's Private Key:
{private_key}")
    return private key
def calculate_public_key(g, private_key, p):
    """Calculate the public key."""
    public_key = pow(g, private_key, p)
    print(f"\nCalculated Server's Public Key: {public key}")
    return public key
def calculate_shared_secret(public_key, private_key, p):
    """Calculate the shared secret."""
    shared secret = pow(public key, private key, p)
    return shared secret
def start server(host='localhost', port=5000):
    # Take p and g as user input
    p = int(input("\nEnter a prime number (p): "))
    g = int(input("\nEnter a primitive root (g): "))
    # Generate server's private and public keys
    private key = generate private key(p)
    public key = calculate public key(g, private key, p)
    # Create server socket
    server socket = socket.socket(socket.AF INET,
socket.SOCK STREAM)
    server socket.bind((host, port))
    server socket.listen(1)
    print(f"\nServer started. Listening on {host}:{port}")
    conn, addr = server socket.accept()
```

```
print(f"\nConnected by {addr}")
   # Send the server's public key to the client
   conn.sendall(str(public key).encode())
    # Receive the client's public key
    client_public_key = int(conn.recv(1024).decode())
    print(f"\nReceived Client's Public Key:
{client_public_key}")
    # Calculate the shared secret
    shared secret =
calculate_shared_secret(client_public_key, private_key, p)
    print(f"\nShared Secret (Server): {shared secret}")
    conn.close()
    server socket.close()
if __name__ == "__main__":
    start server()
# For
# p: 1014273607262027361
# For
# g: 2
```

Output:

Server output-

```
Onkar@LAPTOP-QRVSMK3C MINGW64 ~/Downloads/CNS Lab/Assignment 8 (main)
$ python server_diffie_hellman.py

Enter a prime number (p): 1014273607262027361

Enter a primitive root (g): 2

Generated Server's Private Key: 710999113647365409

Calculated Server's Public Key: 261836127913730729

Server started. Listening on localhost:5000
```

Client output-

Practical Applications of Diffie-Hellman:

- **Secure Communication**: Establishing a shared secret for symmetric encryption over an insecure channel.
- VPNs: Secure key exchange for Virtual Private Networks.
- TLS/SSL: Part of the key exchange process in securing internet communications.

The Diffie-Hellman algorithm forms the basis of many modern cryptographic protocols and is crucial for secure communication in distributed systems.

Final Year B.Tech. (CSE) - VII [2024-25]

6CS451: Cryptography and Network Security Lab (C&NS Lab)

Date: 01/10/2024

Assignment 9

PRN: 21510017 Name: Onkar Anand Yemul

9. Calculate the message digest of a text using the SHA-1 algorithm Ans:

SHA-1 Algorithm:

SHA-1 (Secure Hash Algorithm 1) is a cryptographic hash function that produces a 160-bit hash value (20 bytes), often referred to as a **message digest**. It takes an input message of any size and outputs a fixed-size hash, which is commonly represented as a 40-character hexadecimal number. It was developed by the National Security Agency (NSA) and published by NIST in 1993.

Message Digest of a Text:

A message digest is a fixed-size numerical representation of the contents of a message. For SHA-1, this digest is 160 bits long, and any change in the input message, even a single bit, will result in a drastically different digest (this is known as the avalanche effect). The message digest ensures data integrity by allowing anyone to verify that the message hasn't been altered.

To implement the SHA-1 algorithm without using Python's *hashlib* library, we need to follow the algorithm's steps manually, which involves bitwise operations, padding the input message, and processing it in blocks.

Overview of SHA-1 Algorithm:

1. **Padding the message**: The message is padded so that its length becomes a multiple of 512 bits.

- 2. **Initialize hash values**: There are five constants (H0, H1, H2, H3, H4) initialized to specific values.
- 3. **Processing the message in blocks**: The message is processed in chunks of 512 bits, and the hash is updated after each chunk.
- 4. **Final output**: After all blocks are processed, the hash digest is formed by concatenating the values of H0, H1, H2, H3, and H4.

Python Code for SHA-1 Implementation without using Python's built-in hashlib library:

```
import struct
# Helper functions for bitwise operations
def left rotate(n, b):
    """Left rotate a 32-bit integer n by b bits."""
    return ((n << b) | (n >> (32 - b))) & 0xFFFFFFFF
def sha1_padding(message):
    """Pad the message to ensure the length is a multiple of
512 bits."""
    original byte len = len(message)
    original bit len = original byte len * 8
    # Padding with a '1' bit followed by '0's, and add the
original message length in bits at the end
    message += b' \times 80' # append the bit '1' (10000000 in
binary)
    # Pad with 0s so that the message length is 64 bits
short of a multiple of 512
    while (len(message) * 8) % 512 != 448:
        message += b' \times 00'
    # Append the length of the original message in bits (64-
bit big-endian integer)
```

```
message += struct.pack('>Q', original_bit_len)
    return message
def sha1(message):
    """Calculate the SHA-1 hash of a message."""
    # Initial hash values (first 32 bits of the fractional
parts of the square roots of the first 5 primes)
    h0 = 0x67452301
    h1 = 0 \times EFCDAB89
    h2 = 0 \times 98BADCFE
    h3 = 0 \times 10325476
    h4 = 0xC3D2E1F0
    # Preprocessing: padding the message
    message = sha1 padding(message)
    # Process the message in successive 512-bit chunks (64
bytes each)
    for i in range(0, len(message), 64):
        chunk = message[i:i + 64]
        # Break chunk into sixteen 32-bit big-endian words
        W = [0] * 80
        for j in range(16):
            w[j] = struct.unpack('>I',
chunk[j*4:(j*4)+4])[0]
        # Extend the sixteen 32-bit words into eighty 32-bit
words
        for j in range(16, 80):
            w[j] = left_rotate((w[j-3] ^ w[j-8] ^ w[j-14] ^
w[j-16]), 1)
        # Initialize hash value for this chunk
        a, b, c, d, e = h0, h1, h2, h3, h4
        # Main loop
```

```
for j in range(80):
            if 0 <= j <= 19:
                f = (b \& c) | ((\sim b) \& d)
                k = 0x5A827999
            elif 20 <= j <= 39:
                f = b ^ c ^ d
                k = 0x6ED9EBA1
            elif 40 <= j <= 59:
                f = (b \& c) | (b \& d) | (c \& d)
                k = 0x8F1BBCDC
            elif 60 <= j <= 79:
                f = b ^ c ^ d
                k = 0xCA62C1D6
            temp = (left_rotate(a, 5) + f + e + k + w[j]) &
0xFFFFFFF
            e = d
            d = c
            c = left_rotate(b, 30)
            b = a
            a = temp
        # Add this chunk's hash to the result so far
        h0 = (h0 + a) \& 0xFFFFFFFF
        h1 = (h1 + b) & 0xFFFFFFFF
        h2 = (h2 + c) \& 0 \times FFFFFFFF
        h3 = (h3 + d) \& 0xFFFFFFFF
        h4 = (h4 + e) \& 0 \times FFFFFFFF
    # Produce the final hash value (big-endian)
    return '{:08x}{:08x}{:08x}{:08x}'.format(h0, h1,
h2, h3, h4)
if __name__ == "__main__":
    # Input text
    text = input("Enter the text to calculate SHA-1 hash: ")
    # Convert the text to bytes and compute SHA-1 hash
```

```
sha1_digest = sha1(text.encode('utf-8'))

# Output the SHA-1 hash
print(f"SHA-1 Digest: {sha1_digest}")
```

Output:

```
PROBLEMS
            OUTPUT
                     DEBUG CONSOLE
                                     TERMINAL
                                               PORTS
                                                       SQL CONSOLE
                                                                    GITLENS
 Onkar@LAPTOP-QRVSMK3C MINGW64 ~/Downloads/CNS Lab/Assignment 9 (main)
$ python sha1 without using hashlib.py
 Enter the text to calculate SHA-1 hash: We will have meeting at 5pm tomorrow.
 SHA-1 Digest: 1d9744c575c698f088f2e0fa90c29311abaa1ffb
 Onkar@LAPTOP-QRVSMK3C MINGW64 ~/Downloads/CNS Lab/Assignment 9 (main)
$ python sha1 without using hashlib.py
 Enter the text to calculate SHA-1 hash: We will meet at CCF. *
 SHA-1 Digest: f36ccffcd2c4ef953b0ea93914c6e40f9c8b5afb
 Onkar@LAPTOP-QRVSMK3C MINGW64 ~/Downloads/CNS Lab/Assignment 9 (main)
```

To calculate the message digest of a text using the SHA-1 algorithm in Python, you can use the hashlib library, which provides easy access to various hash algorithms, including SHA-1.

1. hashlib library:

• The hashlib library provides various cryptographic hashing algorithms including SHA-1, SHA-256, MD5, etc.

2. SHA-1 Hash Object:

hashlib.sha1() creates a new SHA-1 hash object.

3. Updating the Hash:

• The update() method takes the input text (which is first encoded into bytes) and updates the hash object with that data.

4. Getting the Digest:

 The hexdigest() method returns the hash value as a hexadecimal string.

Python Code for SHA-1 Message Digest Calculation using hashlib library:

```
import hashlib
def calculate sha1(text):
    # Create a new SHA-1 hash object
    sha1 hash = hashlib.sha1()
    # Encode the input text to bytes and update the hash
object
    sha1 hash.update(text.encode('utf-8'))
    # Get the hexadecimal representation of the digest
    digest = sha1 hash.hexdigest()
    return digest
if __name__ == "__main__":
    # Input text
    text = input("Enter the text to calculate SHA-1 hash: ")
    # Calculate SHA-1 message digest
    sha1 digest = calculate sha1(text)
    # Output the result
    print(f"SHA-1 Digest: {sha1 digest}")
```

Output:

Advantages of SHA-1:

- **Speed and Efficiency**: SHA-1 was designed to be computationally efficient and can process large amounts of data quickly.
- **Widespread Use**: It has been widely adopted and used for various cryptographic applications, including digital signatures, certificates, and integrity checks.
- **Fixed-Length Output**: Regardless of the input size, the output is always 160 bits, making it convenient to use in various security protocols.

Disadvantages of SHA-1:

- Weakness to Collisions: SHA-1 is vulnerable to collision attacks, where two different inputs produce the same hash output. This reduces its effectiveness in ensuring data integrity and security.
- **Security Deprecation**: Due to these vulnerabilities, SHA-1 is no longer considered secure for cryptographic purposes. Most modern systems and protocols, including browsers and SSL certificates, have moved to stronger hash functions like SHA-256 or SHA-3.

Importance of SHA-1:

- Legacy Systems: Despite its vulnerabilities, SHA-1 was used for many years in security applications such as digital signatures and certificates.
- Data Integrity: SHA-1 can still be used to check the integrity of data, ensuring that files have not been altered during transmission.

Security Risks and Vulnerabilities of SHA-1:

- Collision Attacks: The primary vulnerability is the possibility of collision attacks. This means that an attacker could potentially create two different messages with the same hash, compromising the authenticity of the data.
- Birthday Attack: A specific type of attack known as a birthday attack
 makes it easier to find collisions in SHA-1 due to its 160-bit length,
 reducing the security level.
- Deprecation in Modern Systems: Due to these weaknesses, SHA-1 has been deprecated in most cryptographic protocols like TLS (Transport Layer Security) and digital certificates, where stronger algorithms like SHA-256 are preferred.

While SHA-1 played a significant role in the development of cryptographic standards, its vulnerabilities, especially to collision attacks, have made it unsuitable for modern security applications. Understanding SHA-1's purpose and limitations is important, especially when dealing with legacy systems or understanding the evolution of cryptographic hash functions.

<u>Practical Applications of SHA-1:</u>

1. **Digital Signatures**: SHA-1 was commonly used in creating digital signatures to ensure the authenticity and integrity of documents. It

- would generate a hash of the message, which is then signed by a private key.
- 2. **File Integrity Verification**: SHA-1 was used to generate checksums for files to verify that files were not altered during transfer or storage. The recipient could compare the hash of the received file with the original hash to ensure integrity.
- 3. **Version Control Systems**: In systems like Git, SHA-1 hashes were used to identify commits, ensuring the integrity and tracking of changes in code repositories.
- 4. **SSL Certificates**: Until 2017, SHA-1 was used in SSL/TLS certificates for secure web communications. The hash was part of the process to ensure a website's identity and secure data transmission.
- 5. **Password Hashing**: SHA-1 was once used for hashing passwords in databases, providing a layer of security by storing a hashed version of the password instead of the plaintext.

Despite these applications, most systems have transitioned to more secure alternatives due to SHA-1's vulnerabilities.

SHA-512 Algorithm:

SHA-512 (Secure Hash Algorithm 512-bit) is part of the SHA-2 family and produces a 512-bit message digest. It's a cryptographic hash function designed to provide higher security by generating a unique, fixed-length hash value from input data. SHA-512 is widely used for its robust security features.

Overview of SHA-512 Algorithm

SHA-512 is a cryptographic hash function that converts any input (such as a message or file) into a fixed 512-bit hash. The algorithm processes the input data in blocks and applies multiple rounds of complex operations to produce a unique hash value.

Key Steps in the SHA-512 Algorithm:

1. Padding the Message:

The input message is padded to ensure its length is congruent to 896 bits modulo 1024. Padding involves adding a single '1' bit followed by enough '0' bits, so that the message length becomes 1024 bits less than a multiple of 1024. The last 128 bits are used to store the original length of the message.

2. Breaking the Message into Blocks:

The padded message is divided into 1024-bit blocks for processing. Each block will undergo the hashing process individually.

3. Initialize Hash Values:

The algorithm uses eight 64-bit initial hash values, which are constant and specified by the SHA-512 standard. These values form the basis of the hash computation.

4. Processing Each Block:

For each 1024-bit block:

Message Expansion: The 1024-bit block is expanded into 80
 64-bit words. These words are used in the main hashing loop.

- Compression Function: The main loop processes the block with bitwise operations, modular additions, and logical functions (such as AND, XOR, OR). A set of 80 constant values is used along with the expanded message words.
- The hash values are updated after each round using these operations.

5. Update Hash Values:

After processing each block, the intermediate hash values are updated. These values accumulate the results of each block's computations.

6. Concatenate Final Hash:

After all the blocks have been processed, the final 512-bit hash is produced by concatenating the updated hash values from all rounds.

Python Code for SHA-512 Implementation using Python's built-in hashlib library:

```
import hashlib

# Function to hash a message using SHA-512

def sha512_encrypt(message):
    sha512_hash = hashlib.sha512()
    sha512_hash.update(message.encode('utf-8')) # Convert

the message to bytes
    return sha512_hash.hexdigest()

# Function to verify the hash (like a decryption process)

def verify_hash(original_message, provided_hash):
    original_hash = sha512_encrypt(original_message)
    return original_hash == provided_hash

# Menu-driven system

def menu():
    while True:
        print("\n===== SHA-512 Hashing System =====")
```

```
print("1. Encrypt a message using SHA-512")
        print("2. Verify a message against a given hash")
        print("3. Exit")
        choice = input("Enter your choice (1/2/3): ")
        if choice == '1':
            message = input("Enter the message to hash: ")
            hashed message = sha512 encrypt(message)
            print(f"\nSHA-512 Hash: {hashed_message}")
        elif choice == '2':
            original_message = input("Enter the original
message: ")
            provided hash = input("Enter the hash to verify
against: ")
            if verify hash(original message, provided hash):
                print("\nVerification successful! The
message matches the provided hash.")
            else:
                print("\nVerification failed! The message
does not match the provided hash.")
        elif choice == '3':
            print("Exiting the program...")
            break
        else:
            print("Invalid choice. Please choose a valid
option.")
if name == " main ":
    menu()
```

Output:

```
Onkar@LAPTOP-QRVSMK3C MINGW64 ~/Downloads/CNS Lab/Assignment 9 (main)
$ python sha512.py
 ==== SHA-512 Hashing System =====
1. Encrypt a message using SHA-512
 2. Verify a message against a given hash
 3. Exit
 Enter your choice (1/2/3): 1
 Enter the message to hash: We will attack on Pakistan on 11 November, 2019.
 SHA-512 Hash: 6d078e5048722b22c86e92e7390130d3a3b26ed6f36fee69e61eef3e6bccf4b59e42cd4645e142648396a83bea77b55317
 e50607742febfaba116c59eeee88be
 ==== SHA-512 Hashing System =====
 1. Encrypt a message using SHA-512
 2. Verify a message against a given hash
 3. Exit
 Enter your choice (1/2/3): 2
 Enter the original message: We will attack on Pakistan on 11 November, 2019.
 Enter the hash to verify against: 6d078e5048722b22c86e92e7390130d3a3b26ed6f36fee69e61eef3e6bccf4b59e42cd4645e142
 648396a83bea77b55317e50607742febfaba116c59eeee88be
 Verification successful! The message matches the provided hash.
```

```
∑ bash + ∨ □ ଢ ··· ∨ ×
Enter your choice (1/2/3): 2
Enter the original message: We will attack on Pakistan on 11 November, 2019.
Enter the hash to verify against: 6d078e5048722b22c86e92e7390130d3a3b26ed6f36fee69e61eef3e6bccf4b59e42cd4645e142
648396a83bea77b55317e50607742febfaba116c59eeee88be
Verification successful! The message matches the provided hash.
==== SHA-512 Hashing System ==
1. Encrypt a message using SHA-512
2. Verify a message against a given hash
3. Exit
Enter your choice (1/2/3): 2
Enter the original message: We will attack on Pakistan on 11th November, 2019.
Enter the hash to verify against: 6d078e5048722b22c86e92e7390130d3a3b26ed6f36fee69e61eef3e6bccf4b59e42cd4645e142
648396a83bea77b55317e50607742febfaba116c59eeee88be\\
Verification failed! The message does not match the provided hash.
==== SHA-512 Hashing System =====
1. Encrypt a message using SHA-512
2. Verify a message against a given hash
3. Exit
Enter your choice (1/2/3): 3
Exiting the program...
{\tt Onkar@LAPTOP-QRVSMK3C~MINGW64~~/Downloads/CNS~Lab/Assignment~9~(main)}
```

Advantages of SHA-512:

- 1. **High Security**: SHA-512 produces a 512-bit hash, making it much harder to break with brute force attacks compared to smaller hash sizes like SHA-1.
- 2. **Collision Resistance**: It offers strong resistance to collision attacks, meaning it's very unlikely two different inputs will produce the same hash.

- 3. **Efficiency**: Despite its large output size, SHA-512 is designed to be computationally efficient on modern hardware.
- 4. **Compatibility with SHA-2 Family**: SHA-512 shares its core design with other SHA-2 algorithms, making it easier to switch between different levels of security.

Disadvantages of SHA-512:

- 1. **Higher Computational Cost**: Because of its larger size, SHA-512 may require more processing power and memory, making it slower on less powerful devices.
- 2. **Large Hash Size**: The 512-bit hash is larger, which may not be necessary for all applications, particularly when storage or bandwidth is a concern.
- 3. **Overkill for Small Applications**: In some use cases, the high security provided by SHA-512 might be unnecessary, and a smaller hash size like SHA-256 might suffice.

Importance of SHA-512:

SHA-512 is critical in contexts where strong security is essential, particularly in environments that demand protection against sophisticated attacks. Its high bit-length and resistance to common cryptographic attacks make it crucial for protecting sensitive data.

Practical Applications of SHA-512:

- 1. **Digital Signatures and Certificates**: SHA-512 is often used in creating digital signatures and securing SSL/TLS certificates to verify the authenticity and integrity of data.
- 2. **File Integrity**: It's used in verifying file integrity by generating checksums to ensure that files have not been tampered with during transfer or storage.

- 3. **Cryptocurrency**: SHA-512 is used in blockchain technology to secure transactions and validate blocks.
- 4. **Password Hashing**: It's commonly used in securely hashing passwords in databases, making stored passwords difficult to reverse-engineer.
- 5. **Secure Communication**: SHA-512 plays a role in securing communications over networks by being part of cryptographic protocols such as TLS.

Security Risks and Vulnerabilities:

- Larger Hash Size Overhead: While SHA-512 is more secure than smaller hashes, the added size may introduce performance issues for systems that don't need this level of security.
- Quantum Computing Threat: In the future, quantum computing might pose a threat to even robust algorithms like SHA-512, necessitating the development of quantum-resistant algorithms.

Despite its high security, SHA-512 is still vulnerable to advances in technology, but for now, it remains one of the strongest cryptographic hash functions available.

Final Year B.Tech. (CSE) - VII [2024-25]

6CS451: Cryptography and Network Security Lab (C&NS Lab)

Date: 08/10/2024

Assignment 10

PRN: 21510017 Name: Onkar Anand Yemul

10. Implement the SIGNATURE SCHEME - Digital Signature Standard

Ans:

To implement the **Digital Signature Standard (DSS)**, we need to understand its process, which involves the **Digital Signature Algorithm (DSA)**. The DSS is a Federal Information Processing Standard (FIPS) for digital signatures, and it involves three main stages:

- 1. **Key Generation**: Generate a public and private key pair.
- 2. **Signature Generation**: Use the private key to sign a message.
- 3. **Signature Verification**: Use the public key to verify the authenticity of the message.

Overview of the DSA Algorithm:

- DSA involves a pair of keys: private key (used for signing) and public key (used for verification).
- The signature is generated by applying a **hashing algorithm** (such as SHA-1 or SHA-256) to the message, which is then signed using the private key.
- The signature is verified using the public key.

Python Implementation of the DSA (Digital Signature Algorithm) using Python's *cryptography* library for cryptographic operations:

- Using a library here simplifies the task

Steps:

- 1. **Key Generation**: Generate the DSA keys.
 - A DSA private key is generated using dsa.generate_private_key(key_size=2048). The key size can be adjusted (1024, 2048, or 3072 bits), but 2048 is commonly used.
 - The corresponding public key is derived from the private key and returned.
- 2. **Signing**: Use the private key to sign a message.
 - The private_key.sign function is used to sign the message using a specified hash function (SHA-256 here).
 - This produces a signature that can later be verified.
- 3. **Verification**: Use the public key to verify the signature.
 - The public_key.verify function checks whether the signature matches the original message using the public key.
 - If the verification fails, an InvalidSignature exception is raised, indicating that the signature is not valid.

Python Code: DSA-based Digital Signature Implementation:

```
from cryptography.hazmat.backends import default_backend
from cryptography.hazmat.primitives.asymmetric import rsa, padding
from cryptography.hazmat.primitives import hashes, serialization
from cryptography.exceptions import InvalidSignature
import hashlib

# Function to generate RSA private and public keys
def generate_keys():
    private_key = rsa.generate_private_key(
        public_exponent=65537,
        key_size=2048,
        backend=default_backend()
    )
    public_key = private_key.public_key()
```

```
return private key, public key
# Function to sign a message and print its hash
def sign message(private key, message):
    # Hash the message using SHA-256
    message_hash = hashlib.sha256(message.encode()).hexdigest()
    # Sign the hashed message using RSA private key
    signature = private key.sign(
        bytes.fromhex(message hash), # Convert the hex hash to
bytes
        padding.PSS(
            mgf=padding.MGF1(hashes.SHA256()),
            salt length=padding.PSS.MAX LENGTH
        ),
        hashes.SHA256()
    print(f"Hash of the message: {message_hash}") # Print the
message hash
    return signature, message hash
# Function to verify the signature using the provided hash
def verify signature(public key, message hash, signature):
    try:
        # Verify the signature using RSA public key
        public_key.verify(
            signature,
            bytes.fromhex(message hash), # Convert the hex hash
back to bytes
            padding.PSS(
                mgf=padding.MGF1(hashes.SHA256()),
                salt length=padding.PSS.MAX LENGTH
            hashes.SHA256()
        return True
    except InvalidSignature:
        return False
# Function to save keys to files
def save keys to file(private key, public key):
```

```
# Save private key
    with open("private key.pem", "wb") as private file:
        private file.write(
            private key.private bytes(
                encoding=serialization.Encoding.PEM,
                format=serialization.PrivateFormat.PKCS8,
                encryption_algorithm=serialization.NoEncryption()
    # Save public key
    with open("public_key.pem", "wb") as public_file:
        public file.write(
            public key.public bytes(
                encoding=serialization.Encoding.PEM,
                format=serialization.PublicFormat.SubjectPublicKeyIn
fo
    print("Keys saved to files: private key.pem and public key.pem")
# Menu for the digital signature system
def menu():
    private_key, public_key = None, None
    signature = None
    message_hash = None
    while True:
        print("\n===== Digital Signature System =====")
        print("1. Generate RSA Keys")
        print("2. Sign a Message")
        print("3. Verify Signature")
        print("4. Save Keys to Files")
        print("5. Exit")
        choice = input("Enter your choice (1/2/3/4/5): ")
        if choice == '1':
            # Generate RSA private and public keys
            private_key, public_key = generate_keys()
            print("\nRSA Keys Generated!")
        elif choice == '2':
            # Sign a message
```

```
if private key is None:
                print("You need to generate RSA keys first.")
            else:
                message = input("Enter the message to sign: ")
                signature, message hash = sign message(private key,
message)
                print("\nMessage signed successfully!")
        elif choice == '3':
            # Verify the signature
            if public key is None or message hash is None or
signature is None:
                print("You need to sign a message first.")
            else:
                input hash = input("Enter the hash of the message to
verify: ")
                verification_result = verify_signature(public_key,
input hash, signature)
                if verification result:
                    print("\nSignature verified successfully! The
message is authentic.")
                else:
                    print("\nSignature verification failed! The
message is not authentic.")
        elif choice == '4':
            # Save RSA keys to files
            if private_key is None or public_key is None:
                print("You need to generate RSA keys first.")
            else:
                save_keys_to_file(private_key, public_key)
        elif choice == '5':
            print("Exiting the program...")
            break
        else:
            print("Invalid choice. Please try again.")
if name == " main ":
    menu()
```

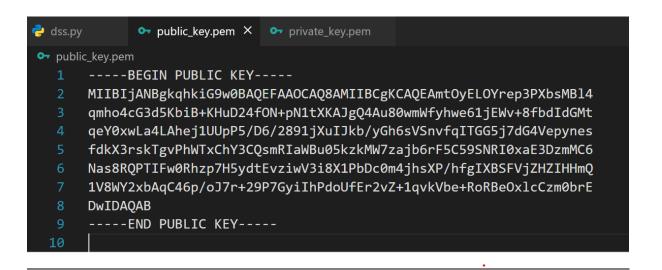
Output:

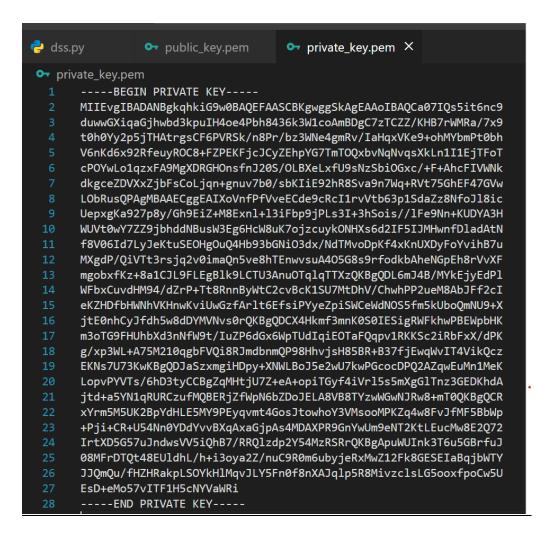
```
MINGW64:/c/Users/Onkar/Downloads
Onkar@LAPTOP-QRVSMK3C MINGW64 ~/Downloads/CNS Lab/Assignment 10 (main)
$ python dss.py
==== Digital Signature System =====
1. Generate RSA Keys
2. Sign a Message
3. Verify Signature
4. Save Keys to Files
Exit
Enter your choice (1/2/3/4/5): 1
RSA Keys Generated!
==== Digital Signature System =====
1. Generate RSA Keys
2. Sign a Message
3. Verify Signature
4. Save Keys to Files
5. Exit
Enter your choice (1/2/3/4/5): 2
Enter the message to sign: Tomorrow's test is resheduled on 25th october, 2024.
Hash of the message: a17d3a1a76be18c39063139bc4d9df98aa9d707d4a1eae5c4cc2a512f53bc371
Message signed successfully!
```

```
==== Digital Signature System =====
1. Generate RSA Keys
2. Sign a Message
Verify Signature
4. Save Keys to Files
5. Exit
Enter your choice (1/2/3/4/5): 3
Enter the hash of the message to verify: a17d3a1a76be18c39063139bc4d9df98aa9d707d4a1eae5c4cc2a512f53bc371
Signature verified successfully! The message is authentic.
==== Digital Signature System =====
1. Generate RSA Keys
2. Sign a Message
3. Verify Signature
4. Save Keys to Files
5. Exit
Enter your choice (1/2/3/4/5): 3
Enter the hash of the message to verify: a17d3a1a76be18c39063139bc4d9df98aa9d707d4a1eae5c4cc2a512f53bc361
Signature verification failed! The message is not authentic.
==== Digital Signature System =====
1. Generate RSA Keys
2. Sign a Message
3. Verify Signature
4. Save Keys to Files
5. Exit
Enter your choice (1/2/3/4/5): 4
Keys saved to files: private_key.pem and public_key.pem
==== Digital Signature System =====
1. Generate RSA Keys

    Sign a Message
    Verify Signature

4. Save Keys to Files
5. Exit
Enter your choice (1/2/3/4/5): 5
Exiting the program...
```





Virtual labs:



Digitally sign the plaintext with Hashed RSA.

Plaintext (string):

Ex. Minister of Maharashtra got assassinated yest SHA-1

Hash output(hex):

041d610065f6a89ecd2059032221ec4e6e2e42f0

Input to RSA(hex):
041d610065f6a89ecd2059032221ec4e6e2e42f0
Digital Signature(hex):
/2
Digital Signature(base64):
Status:
RSA public key
Public exponent (hex, F4=0x10001):
Modulus (hex):
ABC30681295774F7CECA691EC17F4E762DA6DE70F198EAEE3CCE3A435FC006B9 71DC24E55904F1D2705758C041C2B0B18E8BFAE2C9CD96B50082D7D8C7342CBA B7F6E0622DA53B8B56DBDB24174F00173263CFECAE604795CDA2A037BC3A69B7
C0090AA2DE1568998BCD6D70CC2E0574755B9F7986AE01CE8714A26144279CDB 1024 bit 1024 bit (e=3) 512 bit 512 bit (e=3)





Input to RSA(hex):	
041d610065f6a89ecd2059032221ec4e6e2e42f0 Apply RSA	
Digital Signature(hex): 3a6b5f4c871da1863c4aff0ffed6a61bfb12a602b7379f509a98ba611ae80193 dc4a608f5061a861fdf10c1e7c13064e3370fa1d84f3ca198d15034dd872ee45 86d9555b8b3118f17e1dc20c2cd34aa4cc73ba77d324ccdb5f7b60044f1c2d42	
9d175f379b16263c14c93e3d4858b1e06ad9365287fb4ef7017f0f5cd41ecf2b Digital Signature(base64):	<u></u>
OmtfTIcdoYY8Sv8P/tamG/sSpgK3N59Qmpi6YRroAZPcSmCPUGGoYf3xDB58EwZO M3D6HYTzyhmNFQNN2HLuRYbZVVuLMRjxfh3CDCzTSqTMc7p30yTM2197YARPHC1C nRdfN5sWJjwUyT49SFix4GrZN1KH+073AX8PXNQezys=	6
Status: Time: 3ms	
RSA public key	
Public exponent (hex, F4=0x10001):	
Modulus (hex): ABG30681295774F7CECA691EC17F4E762DA6DE70F198EAEE3CCE3A435FC006B9 71DC24E55904F1D2705758C041C2B0B18E8BFAE2C9CD96B50082D7D8C7342CBA B7F6E0622DA53B8B56DBDB24174F00173263CFECAE604795CDA2A037BC3A69B7 C0090AA2DE1568998BCD6D70CC2E0574755B9F7986AE01CE8714A26144279CDB	de la companya de la
1024 bit 1024 bit (e=3) 512 bit 512 bit (e=3)	

How It Works:

- The **private key** signs the message, creating a unique signature.
- The **public key** is used by anyone to verify that the message has not been tampered with and that it was signed by the owner of the private key.

Practical Applications:

- **Digital Signatures** ensure the authenticity and integrity of a message or document.
- **Message Authentication**: In secure communication, the sender can sign the message, and the receiver can verify the signature to ensure the message is authentic.

Advantages of DSS (Digital Signature Standard, based on DSA):

- **Authentication**: Ensures the identity of the sender, as only the sender's private key can create a valid signature.
- Integrity: Guarantees that the message hasn't been altered, as any change would invalidate the signature.
- Non-repudiation: The sender cannot deny having signed the message, as the signature is unique to the private key.
- **Efficiency**: DSA is optimized for creating signatures, making it relatively fast for signing compared to some other algorithms.
- **Security**: Provides a high level of security, especially with modern key sizes (2048 bits or more).

Disadvantages of DSS:

- **Slower Verification**: DSA is slower at verifying signatures compared to some alternatives like RSA, making it less suitable for environments that require frequent verification.
- **Key Management**: Requires secure management of private keys; if the private key is compromised, the entire security system is at risk.
- Message Size Limitations: DSA only signs the hash of a message, so very large messages require hashing before signature generation.
- **Limited Use Cases**: DSA is primarily designed for digital signatures and not for encryption, unlike algorithms like RSA.

Importance of DSS:

- Government Standard: DSS (and DSA) is an official standard used by governments and organizations worldwide for secure digital signatures.
- Legal Recognition: Digital signatures created using DSS are often legally recognized, making them suitable for contracts and other legal documents.

•	Widely Used in Secure Communications : DSS is crucial in protocols like SSL/TLS, ensuring secure web communications, digital certificates, an more.
nc	summary, DSS is important for ensuring the authenticity, integrity, and on-repudiation of digital communications, despite some performance-lated limitations.

Final Year B.Tech. (CSE) - VII [2024-25]

6CS451: Cryptography and Network Security Lab (C&NS Lab)

Date: 15/10/2024

Assignment 11

PRN: 21510017 Name: Onkar Anand Yemul

11. Demonstration of SSL using Wireshark.

Ans:

SSL (Secure Sockets Layer) and TLS (Transport Layer Security) are cryptographic protocols designed to secure communication over a network, especially the internet. They ensure that data transmitted between a client (e.g., web browser) and a server (e.g., website) is encrypted, authenticated, and tamper-proof.

TLS is the more recent and secure version of SSL. SSL is now considered obsolete, and TLS is widely used today (referred to as "SSL/TLS" in many contexts).

SSL vs TLS:

- SSL: Older protocol, now considered insecure due to vulnerabilities like POODLE and BEAST.
- TLS: Successor to SSL, currently in use with versions like TLS 1.2 and TLS 1.3. TLS 1.3 is the most secure and efficient version available today.

Goals of SSL/TLS:

- 1. Confidentiality: Encrypts data to prevent unauthorized access.
- 2. Integrity: Ensures data has not been altered during transmission.

3. **Authentication**: Verifies the identity of the server, and optionally the client.

How SSL/TLS Works:

SSL/TLS secures communication through a series of steps that establish a **secure connection** before data exchange. These steps are known as the **SSL/TLS handshake**. The handshake is a process where the client and server agree on encryption methods and share keys for secure communication.

Steps Involved in SSL/TLS Handshake:

1. Client Hello:

- The client (e.g., a web browser) sends a Client Hello message to the server.
- The message includes:
 - SSL/TLS version the client supports.
 - A list of supported **cipher suites** (encryption algorithms).
 - A randomly generated value called the Client Random.
 - Optional information like session resumption data.

2. Server Hello:

- The server responds with a **Server Hello** message.
- It selects the SSL/TLS version and cipher suite from the list provided by the client.
- o It also generates and sends a Server Random value.

3. Server Certificate:

- o The server sends its SSL/TLS certificate to the client.
- This certificate contains the server's public key and is signed by a trusted Certificate Authority (CA).
- The client uses this certificate to verify the server's identity.

4. Key Exchange:

- The client and server exchange information that allows both parties to generate a shared session key.
- The session key is used to encrypt data exchanged during the session.
- The key exchange can use algorithms like RSA (Rivest-Shamir-Adleman) or Diffie-Hellman.

5. Client Key Exchange:

- The client generates a Pre-master Secret (a random number) and encrypts it using the server's public key from the certificate.
- This pre-master secret is sent to the server.

6. Session Key Generation:

- Both the client and the server use the pre-master secret, along with the Client Random and Server Random values, to independently generate the same session key.
- The session key is a symmetric key used for encrypting and decrypting the communication during the session.

7. Change Cipher Spec:

 The client and server send a Change Cipher Spec message to inform each other that future messages will be encrypted using the session key.

8. Finished

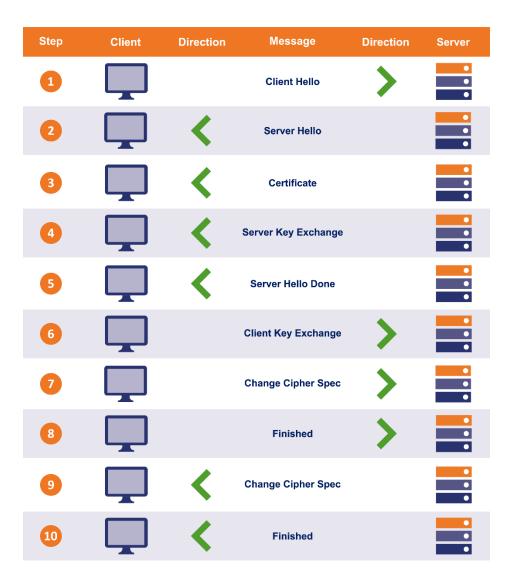
- Both the client and the server send a Finished message, encrypted with the session key, to indicate that the handshake is complete.
- If both messages are successfully decrypted and verified, the SSL/TLS connection is established.

9. Encrypted Communication:

- After the handshake, the client and server use the session key to encrypt and decrypt all subsequent data exchanged during the session.
- This ensures that sensitive data, such as login credentials, is transmitted securely.

Summary of the SSL/TLS Handshake:

- 1. **Client Hello** (Client proposes SSL/TLS version, cipher suite, and sends a random number)
- 2. **Server Hello** (Server selects version, cipher suite, and sends a random number)
- 3. **Server Certificate** (Server sends its certificate for client to verify)
- 4. **Key Exchange** (Client and server exchange information to generate a shared session key)
- 5. **Change Cipher Spec** (Both agree to switch to encrypted communication)
 - 1. **Finished** (Handshake is complete, and encrypted communication begins)



Importance of SSL/TLS:

- Web Security: SSL/TLS ensures that sensitive information (such as passwords, credit card details) is transmitted securely over the internet.
- **Authentication**: Verifies the identity of the website to prevent man-in-the-middle (MITM) attacks.
- **Data Privacy**: Prevents eavesdropping and tampering during communication.

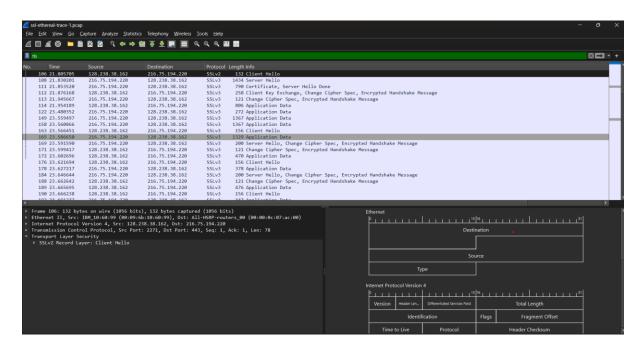
In conclusion, SSL/TLS is crucial for ensuring secure and trusted communication over networks, particularly for websites and online services.

Wireshark Lab: SSL v8.0

Download the zip file http://gaia.cs.umass.edu/wireshark-labs/wireshark- traces.zip and extract the ssl-ethereal trace-1 packet trace.

1. Capturing packets in an SSL session

It displays only the Ethernet frames that contain SSL records sent from and received by your host



An SSL record is the same thing as an SSL message.

2. A look at the captured trace

An Ethernet frame may contain one or more SSL records. (This is very different from HTTP, for which each frame contains either one complete HTTP message or a portion of a HTTP message.) Also, an SSL record may not

completely fit into an Ethernet frame, in which case multiple frames will be needed to carry the record.

1. For each of the first 8 Ethernet frames, specify the source of the frame (client or server), determine the number of SSL records that are included in the frame, and list the SSL record types that are included in the frame. Draw a timing diagram between client and server, with one arrow for each SSL record.

Ans:

Ē	ile	l-ethereal-trace	<u>G</u> o <u>C</u> apture <u>A</u> nalyze <u>S</u> tati	stics Telephony Wireless		_
	L tis		<u>— M ⊠ Ø</u> ९⇔=		લ્લ્લ	· ==
L	tis	S				
No		Time	Source	Destination	Protocol	Length Info
		106 21.805	705 128.238.38.162	216.75.194.220	SSLv2	132 Client Hello
		108 21.830	216.75.194.220	128.238.38.162	SSLv3 •	1434 Server Hello
		111 21.853	216.75.194.220	128.238.38.162	SSLv3	790 Certificate, Server Hello Done
		112 21.876	.68 128.238.38.162	216.75.194.220	SSLv3	258 Client Key Exchange, Change Cipher Spec, Encrypted Handshake Message
		113 21.945	667 216.75.194.220	128.238.38.162	SSLv3	121 Change Cipher Spec, Encrypted Handshake Message
		114 21.954	89 128.238.38.162	216.75.194.220	SSLv3	806 Application Data
		122 23.480	352 216.75.194.220	128.238.38.162	SSLv3	272 Application Data
		149 23.559	197 216.75.194.220	128.238.38.162	SSLv3	1367 Application Data
		158 23.560	366 216.75.194.220	128.238.38.162	SSLv3	1367 Application Data

Frame 1: (Frame 106)

• Source: Client (128.238.38.162)

Number of SSL Records: 1

• SSL Record Type: Client Hello (SSLv2)

Frame 2: (Frame 108)

• Source: Server (216.75.194.220)

Number of SSL Records: 1

• SSL Record Type: Server Hello (SSLv3)

Frame 3: (Frame 111)

• Source: Server (216.75.194.220)

• Number of SSL Records: 2

SSL Record Types:

1. Certificate (SSLv3)

2. Server Hello Done (SSLv3)

Frame 4: (Frame 112)

- Source: Client (128.238.38.162)
- Number of SSL Records: 3
- SSL Record Types:
 - 1. Client Key Exchange (SSLv3)
 - 2. Change Cipher Spec (SSLv3)
 - 3. Encrypted Handshake Message (SSLv3)

Frame 5: (Frame 113)

- Source: Server (216.75.194.220)
- Number of SSL Records: 2
- SSL Record Types:
 - 1. Change Cipher Spec (SSLv3)
 - 2. Encrypted Handshake Message (SSLv3)

Frame 6: (Frame 114)

- Source: Client (128.238.38.162)
- Number of SSL Records: 1
- SSL Record Type: Application Data (SSLv3)

Frame 7: (Frame 122)

- Source: Server (216.75.194.220)
- Number of SSL Records: 1
- SSL Record Type: Application Data (SSLv3)

Frame 8: (Frame 149)

- Source: Server (216.75.194.220)
- Number of SSL Records: 1
- SSL Record Type: Application Data (SSLv3)

2. Each of the SSL records begins with the same three fields (with possibly different values). One of these fields is "content type" and has length of one byte. List all three fields and their lengths.

Ans:

Each SSL record starts with the following three fields:

• Content Type: 1 byte

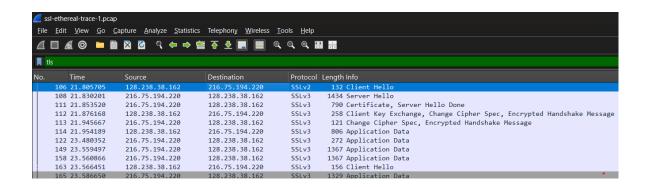
Version: 2 bytesLength: 2 bytes

How to Find These Fields:

If you are using packet capture software like **Wireshark**, you can find these fields in the packet capture by:

- 1. **Open Wireshark** and load the captured SSL/TLS packet data (the one you listed).
- 2. **Select an SSL/TLS packet** from the list and expand the **"Secure Sockets Layer"** or **"Transport Layer Security"** section in the detailed packet view.
- 3. You will see the **Record Layer** header information, where these fields will be listed:
 - Content Type: Displays the type of SSL/TLS record (Handshake, Application Data, etc.)
 - Version: The protocol version (e.g., TLS 1.2)
 - Length: The size of the encrypted data.

Selecting Client Hello packet:



```
Frame 106: 132 bytes on wire (1056 bits), 132 bytes captured (1056 bits)

Ethernet II, Src: IBM_10:60:99 (00:09:6b:10:60:99), Dst: All-HSRP-routers_00 (00:00:0c:07:ac:00)

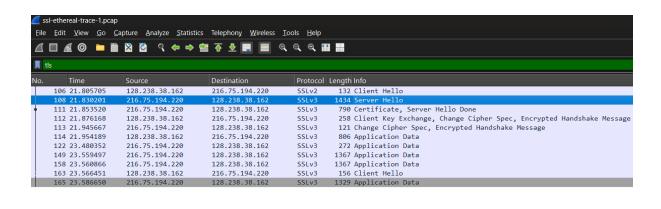
Internet Protocol Version 4, Src: 128.238.38.162, Dst: 216.75.194.220

Transmission Control Protocol, Src Port: 2271, Dst Port: 443, Seq: 1, Ack: 1, Len: 78

Transport Layer Security

▼ SSLv2 Record Layer: Client Hello
    [Version: SSL 2.0 (0x0002)]
    Length: 76
    Handshake Message Type: Client Hello (1)
    Version: SSL 3.0 (0x0300)
    Cipher Spec Length: 51
    Session ID Length: 0
    Challenge Length: 16
    Cipher Specs (17 specs)
    Challenge
```

Selecting Server Hello packet:



```
Frame 108: 1434 bytes on wire (11472 bits), 1434 bytes captured (11472 bits)

Ethernet II, Src: Cisco_83:e4:54 (00:b0:8e:83:e4:54), Dst: IBM_10:60:99 (00:09:6b:10:60:99)

Internet Protocol Version 4, Src: 216.75.194.220, Dst: 128.238.38.162

Transmission Control Protocol, Src Port: 443, Dst Port: 2271, Seq: 1, Ack: 79, Len: 1380

Transport Layer Security

SSLv3 Record Layer: Handshake Protocol: Server Hello
Content Type: Handshake (22)
Version: SSL 3.0 (0x0300)
Length: 74

Handshake Protocol: Server Hello
TLS segment data (1301 bytes)

Frame 108: 1434 bytes captured (11472 bits)

Frame 108: 1434 bytes on wire (11472 bits), 1434 bytes captured (11472 bits)

Ethernet II, Src: Cisco_83:e4:54 (00:b0:8e:83:e4:54), Dst: IBM_10:60:99 (00:09:6b:10:60:99)

Frame 108: 1434 bytes on wire (11472 bits), 1434 bytes captured (11472 bits)

Frame 108: 1434 bytes on wire (11472 bits), 1434 bytes captured (11472 bits)

Frame 108: 1434 bytes on wire (11472 bits), 1434 bytes captured (11472 bits)

Frame 108: 1434 bytes on wire (11472 bits), 1434 bytes captured (11472 bits)

Frame 108: 1434 bytes on wire (11472 bits), 1434 bytes captured (11472 bits)

Frame 108: 1434 bytes on wire (11472 bits)

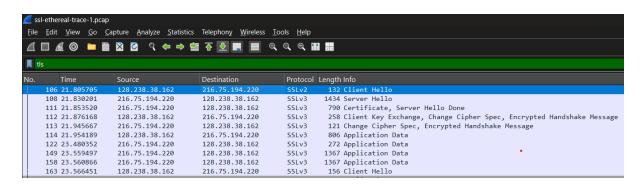
Frame 108:
```

ClientHello Record:

3. Expand the ClientHello record. (If your trace contains multiple ClientHello records, expand the frame that contains the first one.) What is the value of the content type?

Ans:

The **ClientHello** record in **Frame 106** is an SSLv2 message with a handshake message type of **Client Hello (1)**.



```
Frame 106: 132 bytes on wire (1056 bits), 132 bytes captured (1056 bits)

Ethernet II, Src: IBM_10:60:99 (00:09:6b:10:60:99), Dst: All-HSRP-routers_00 (00:00:0c:07:ac:00)

Internet Protocol Version 4, Src: 128.238.38.162, Dst: 216.75.194.220

Transmission Control Protocol, Src Port: 2271, Dst Port: 443, Seq: 1, Ack: 1, Len: 78

Transport Layer Security

* SSLv2 Record Layer: Client Hello
    [Version: SSL 2.0 (0x0002)]
    Length: 76

Handshake Message Type: Client Hello (1)

Version: SSL 3.0 (0x0300)
    Cipher Spec Length: 51
    Session ID Length: 0
    Challenge Length: 16

* Cipher Specs (17 specs)
    Challenge
```

4. Does the ClientHello record contain a nonce (also known as a "challenge")? If so, what is the value of the challenge in hexadecimal notation?

Ans:

Yes, the ClientHello record contains a nonce (also known as a "challenge").

```
Frame 106: 132 bytes on wire (1056 bits), 132 bytes captured (1056 bits)

Ethernet II, Src: IBM_10:60:99 (00:09:6b:10:60:99), Dst: All-HSRP-routers_00 (00:00:0c:07:ac:00)

Internet Protocol Version 4, Src: 128.238.38.162, Dst: 216.75.194.220

Transmission Control Protocol, Src Port: 2271, Dst Port: 443, Seq: 1, Ack: 1, Len: 78

Transport Layer Security

* SSLv2 Record Layer: Client Hello
    [Version: SSL 2.0 (0x0002)]
    Length: 76
    Handshake Message Type: Client Hello (1)
    Version: SSL 3.0 (0x0300)
    Cipher Spec Length: 51
    Session ID Length: 0
    Challenge Length: 16
    Cipher Specs (17 specs)

Challenge
```

```
| Frame 106: 132 bytes on wire (1800 6 Mix), 312 bytes captured (1806 bits) | 120 bytes captured (1
```

5. Does the ClientHello record advertise the cyber suites it supports? If so, in the first listed suite, what are the public-key algorithm, the symmetric-key algorithm, and the hash algorithm?

Ans:

Yes, the ClientHello record does advertise the cipher suites it supports.

```
▼ Cipher Specs (17 specs)

Cipher Spec: TLS_RSA_WITH_RC4_128_MD5 (0x000004)

Cipher Spec: TLS RSA WITH RC4 128 SHA (0x000005)
```

In the first listed cipher suite, which is TLS_RSA_WITH_RC4_128_MD5 (0x000004), the following algorithms are used:

- Public-key algorithm: RSA
- Symmetric-key algorithm: RC4 (with a key length of 128 bits)
- Hash algorithm: MD5

This combination indicates that the client supports this suite for secure communication.

```
Frame 106: 132 bytes on wire (1056 bits), 132 bytes captured (1056 bits)
 Ethernet II, Src: IBM_10:60:99 (00:09:6b:10:60:99), Dst: All-HSRP-routers_00 (00:00:0c:07:ac:00)
▶ Internet Protocol Version 4, Src: 128.238.38.162, Dst: 216.75.194.220
> Transmission Control Protocol, Src Port: 2271, Dst Port: 443, Seq: 1, Ack: 1, Len: 78
 Transport Layer Security
    SSLv2 Record Layer: Client Hello
        [Version: SSL 2.0 (0x0002)]
        Length: 76
        Handshake Message Type: Client Hello (1)
        Version: SSL 3.0 (0x0300)
        Cipher Spec Length: 51
        Session ID Length: 0
        Challenge Length: 16
       Cipher Specs (17 specs)
           Cipher Spec: TLS RSA WITH RC4 128 MD5 (0x000004)
           Cipher Spec: TLS_RSA_WITH_RC4_128_SHA (0x0000005)
           Cipher Spec: TLS_RSA_WITH_3DES_EDE_CBC_SHA (0x00000a)
           Cipher Spec: SSL2_RC4_128_WITH_MD5 (0x010080)
           Cipher Spec: SSL2_DES_192_EDE3_CBC_WITH_MD5 (0x0700c0)
           Cipher Spec: SSL2_RC2_128_CBC_WITH_MD5 (0x030080)
           Cipher Spec: TLS_RSA_WITH_DES_CBC_SHA (0x000009)
           Cipher Spec: SSL2 DES 64 CBC WITH MD5 (0x060040)
           Cipher Spec: TLS_RSA_EXPORT1024_WITH_RC4_56_SHA (0x000064)
           Cipher Spec: TLS_RSA_EXPORT1024_WITH_DES_CBC_SHA (0x000062)
           Cipher Spec: TLS_RSA_EXPORT_WITH_RC4_40_MD5 (0x0000003)
           Cipher Spec: TLS_RSA_EXPORT_WITH_RC2_CBC_40_MD5 (0x0000006)
           Cipher Spec: SSL2 RC4 128 EXPORT40 WITH MD5 (0x020080)
           Cipher Spec: SSL2_RC2_128_CBC_EXPORT40_WITH_MD5 (0x040080)
           Cipher Spec: TLS_DHE_DSS_WITH_3DES_EDE_CBC_SHA (0x000013)
Cipher Spec: TLS_DHE_DSS_WITH_DES_CBC_SHA (0x000012)
           Cipher Spec: TLS_DHE_DSS_EXPORT1024_WITH_DES_CBC_SHA (0x0000063)
        Challenge
```

ServerHello Record:

6. Locate the ServerHello SSL record. Does this record specify a chosen cipher suite? What are the algorithms in the chosen cipher suite?

Ans:

<u></u>	ssl-ethereal-trace-1.pcap e <u>E</u> dit <u>V</u> iew <u>G</u> o <u>C</u>		cs Telephony <u>W</u> ireless	<u>T</u> ools <u>H</u> elp	
		🛚 🖺 🚨 🥄 ← ⇒	🤷 春 🕹 🔲 🔳 🧐		
	tls				
No.	Time	Source	Destination	Protocol	Length Info
	106 21.805705	128.238.38.162	216.75.194.220	SSLv2	132 Client Hello
	108 21.830201	216.75.194.220	128.238.38.162	SSLv3	1434 Server Hello
+	111 21.853520	216.75.194.220	128.238.38.162	SSLv3	790 Certificate, Server Hello Done
	112 21.876168	128.238.38.162	216.75.194.220	SSLv3	258 Client Key Exchange, Change Cipher Spec, Encrypted Handshake Message
	113 21.945667	216.75.194.220	128.238.38.162	SSLv3	121 Change Cipher Spec, Encrypted Handshake Message
	114 21.954189	128.238.38.162	216.75.194.220	SSLv3	806 Application Data
	122 23.480352	216.75.194.220	128.238.38.162	SSLv3	272 Application Data
	149 23.559497	216.75.194.220	128.238.38.162	SSLv3	1367 Application Data
	158 23.560866	216.75.194.220	128.238.38.162	SSLv3	1367 Application Data
	163 23.566451	128.238.38.162	216.75.194.220	SSLv3	156 Client Hello
	105 22 500050	217 75 104 222	100 000 00 100	CCLUS	1300 Auditoritor Bata

```
Frame 108: 1434 bytes on wire (11472 bits), 1434 bytes captured (11472 bits)
  Ethernet II, Src: Cisco_83:e4:54 (00:b0:8e:83:e4:54), Dst: IBM_10:60:99 (00:09:6b:10:60:99)
▶ Internet Protocol Version 4, Src: 216.75.194.220, Dst: 128.238.38.162
Transmission Control Protocol, Src Port: 443, Dst Port: 2271, Seq: 1, Ack: 79, Len: 1380
 Transport Layer Security
    SSLv3 Record Layer: Handshake Protocol: Server Hello
       Content Type: Handshake (22)
        Version: SSL 3.0 (0x0300)
        Length: 74
       Handshake Protocol: Server Hello
          Handshake Type: Server Hello (2)
          Length: 70
           Version: SSL 3.0 (0x0300)
        ▶ Random: 0000000042dbed248b8831d04cc98c26e5badc4e267c391944f0f070ece57745
           Session ID Length: 32
           Session ID: 1bad05faba02ea92c64c54be4547c32f3e3ca63d3a0c86ddad694b45682da22f
           Cipher Suite: TLS RSA WITH RC4 128 MD5 (0x0004)
          Compression Method: null (0)
          [JA3S Fullstring: 768,4,]
[JA3S: 1f8f5a3d2fd435e36084db890693eafd]
     TLS segment data (1301 bytes)
```

Yes, the ServerHello SSL record specifies a chosen cipher suite. The chosen cipher suite is TLS_RSA_WITH_RC4_128_MD5 (0x0004).

The algorithms in this chosen cipher suite are:

- Public-key algorithm: RSA
- **Symmetric-key algorithm**: RC4 (with a key length of 128 bits)
- Hash algorithm: MD5

This indicates the server's selected encryption method for the session.

7. Does this record include a nonce? If so, how long is it? What is the purpose of the client and server nonces in SSL?

Ans:

Locate the Nonce:

- The ServerHello response may not explicitly list a nonce like the ClientHello does, but it usually includes a Session ID and potentially a Server Random value (which acts similarly to a nonce).
- Look for fields labeled Session ID Length, Session ID, and Random.

```
Frame 108: 1434 bytes on wire (11472 bits), 1434 bytes captured (11472 bits)
▶ Ethernet II, Src: Cisco_83:e4:54 (00:b0:8e:83:e4:54), Dst: IBM_10:60:99 (00:09:6b:10:60:99)
▶ Internet Protocol Version 4, Src: 216.75.194.220, Dst: 128.238.38.162
 Transmission Control Protocol, Src Port: 443, Dst Port: 2271, Seq: 1, Ack: 79, Len: 1380
 Transport Layer Security
    SSLv3 Record Layer: Handshake Protocol: Server Hello
       Content Type: Handshake (22)
       Version: SSL 3.0 (0x0300)
       Length: 74
     ▼ Handshake Protocol: Server Hello
          Handshake Type: Server Hello (2)
          Length: 70
          Version: SSL 3.0 (0x0300)
                  0000000042dbed248b8831d04cc98c26e5badc4e267c391944f0f070ece57745
             GMT Unix Time: Jan 1, 1970 05:30:00.000000000 India Standard Time
             Random Bytes: 42dbed248b8831d04cc98c26e5badc4e267c391944f0f070ece57745
          Session ID Length: 32
          Session ID: 1bad05faba02ea92c64c54be4547c32f3e3ca63d3a0c86ddad694b45682da22f
          Cipher Suite: TLS_RSA_WITH_RC4_128_MD5 (0x0004)
          Compression Method: null (0)
          [JA3S Fullstring: 768,4,]
          [JA3S: 1f8f5a3d2fd435e36084db890693eafd]
     TLS segment data (1301 bytes)
```

Yes, the ServerHello record includes a nonce, which is referred to as the "Random" value. In this case, the nonce is 32 bytes long (256 bits).

Purpose of Nonces in SSL:

- 1. **Prevent Replay Attacks**: Nonces ensure that each session is unique, preventing attackers from reusing old messages to impersonate a user or a session.
- 2. **Key Generation**: Nonces are used in the key generation process during the SSL handshake. They contribute to creating session keys that are unique for each session, ensuring that even if the same keys were used in different sessions, the actual keys derived will be different due to the unique nonces.

By using nonces, SSL enhances the security and integrity of the communication between the client and server.

Purpose of Nonce in the ServerHello Record:

- 1. Session Uniqueness:
 - Similar to the ClientHello, the Server Random value helps ensure that the session is unique. It differentiates this session from previous ones.

2. Key Derivation:

The Server Random value is combined with the Client Random value (from the ClientHello) during the key derivation process to create session keys for encrypting the data exchanged in the session. This ensures that the keys are unique for each session.

3. Preventing Replay Attacks:

- Just as with the client, the server's nonce (or Server Random)
 helps protect against replay attacks, ensuring that each session
 is independent and cannot be reused maliciously.
- 8. Does this record include a session ID? What is the purpose of the session ID?

Ans:

```
Frame 108: 1434 bytes on wire (11472 bits), 1434 bytes captured (11472 bits)
Ethernet II, Src: Cisco_83:e4:54 (00:b0:8e:83:e4:54), Dst: IBM_10:60:99 (00:09:6b:10:60:99)
Internet Protocol Version 4, Src: 216.75.194.220, Dst: 128.238.38.162
Transmission Control Protocol, Src Port: 443, Dst Port: 2271, Seq: 1, Ack: 79, Len: 1380
Transport Layer Security
  SSLv3 Record Layer: Handshake Protocol: Server Hello
      Content Type: Handshake (22)
      Version: SSL 3.0 (0x0300)
      Length: 74
    ▼ Handshake Protocol: Server Hello
        Handshake Type: Server Hello (2)
         Length: 70
        Version: SSL 3.0 (0x0300)
        Random: 0000000042dbed248b8831d04cc98c26e5badc4e267c391944f0f070ece57745
           GMT Unix Time: Jan 1, 1970 05:30:00.000000000 India Standard Time
           Random Bytes: 42dbed248b8831d04cc98c26e5badc4e267c391944f0f070ece57745
        Session ID Length: 32
         Session ID: 1bad05faba02ea92c64c54be4547c32f3e3ca63d3a0c86ddad694b45682<u>da22f</u>
         Cipher Suite: TLS_RSA_WITH_RC4_128_MD5 (0x0004)
         Compression Method: null (0)
         [JA3S Fullstring: 768,4,]
         [JA3S: 1f8f5a3d2fd435e36084db890693eafd]
   TLS segment data (1301 bytes)
```

Yes, the ServerHello record includes a session ID, which has a length of 32 bytes in this case.

Purpose of the Session ID:

- 1. **Session Resumption**: The session ID allows clients and servers to resume a previous SSL/TLS session without needing to perform a full handshake again. This speeds up the connection process for subsequent sessions between the same client and server.
- 2. **State Management**: The session ID helps manage session states on the server. It allows the server to recognize and retrieve session parameters (like the cipher suite and keys) associated with that ID, facilitating quicker reconnections.
- 3. **Security**: By using a session ID, SSL/TLS can provide continuity and consistency for secure communications, ensuring that established session parameters are reused securely.

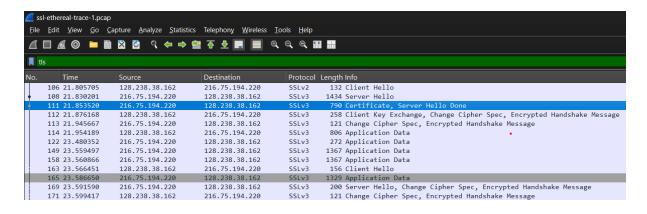
Overall, the session ID is crucial for improving performance and maintaining security during repeated connections.

9. Does this record contain a certificate, or is the certificate included in a separate record. Does the certificate fit into a single Ethernet frame?

Ans:

The ServerHello record itself does not contain a certificate; certificates are included in a separate record, specifically in the Certificate message that follows the ServerHello during the handshake process.

Regarding the size of the certificate, whether it fits into a single Ethernet frame depends on the actual size of the certificate being transmitted. Ethernet frames typically have a maximum payload size of around 1500 bytes. If the certificate's size is within this limit, it can fit into a single frame; otherwise, it will be fragmented across multiple frames.



```
▶ Frame 111: 790 bytes on wire (6320 bits), 790 bytes captured (6320 bits)
Ethernet II, Src: Cisco_83:e4:54 (00:b0:8e:83:e4:54), Dst: IBM_10:60:99 (00:09:6b:10:60:99)
▶ Internet Protocol Version 4, Src: 216.75.194.220, Dst: 128.238.38.162
> Transmission Control Protocol, Src Port: 443, Dst Port: 2271, Seq: 2049, Ack: 79, Len: 736
 [3 Reassembled TCP Segments (2696 bytes): #108(1301), #109(668), #111(727)]
 Transport Layer Security
   ▼ SSLv3 Record Layer: Handshake Protocol: Certificate
       Content Type: Handshake (22)
        Version: SSL 3.0 (0x0300)
        Length: 2691

    Handshake Protocol: Certificate

          Handshake Type: Certificate (11)
          Length: 2687
          Certificates Length: 2684
        ▼ Certificates (2684 bytes)
             Certificate Length: 1352
           Certificate [...]: 308205443082042ca003020102021066a50f1630ded7949e62be443164f4a1300d06092a..
              ▼ signedCertificate
                  version: v3 (2)
                   serialNumber: 0x66a50f1630ded7949e62be443164f4a1
                > signature (sha1WithRSAEncryption)
                ▶ issuer: rdnSequence (0)
                ▶ validity
                subject: rdnSequence (0)
                subjectPublicKeyInfo
                ▶ extensions: 8 items
              algorithmIdentifier (sha1WithRSAEncryption)
                   Algorithm Id: 1.2.840.113549.1.1.5 (sha1WithRSAEncryption)
                Padding: 0
                encrypted [...]: c5874d64289b79189349ab06412f4083c873da91831e7e535677f6d009cfba23bab89b3...
             Certificate Length: 1326
           Certificate [...]: 3082052a30820493a00302010202040200029a300d06092a864886f70d01010505003075...
             signedCertificate
              algorithmIdentifier (sha1WithRSAEncryption)
                Padding: 0
                encrypted [...]: 3a1e246fdadb366cfe3a7339ddcd6a1c69a1001f6fd4aed51fd48829a521c257f6557b3...
 Transport Layer Security

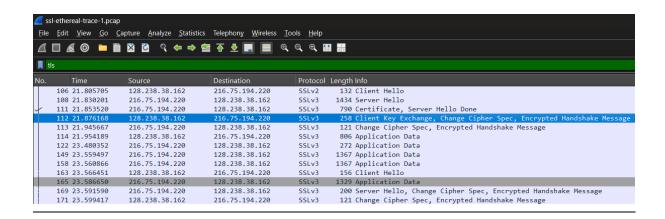
    SSLv3 Record Layer: Handshake Protocol: Server Hello Done

        Content Type: Handshake (22)
        Version: SSL 3.0 (0x0300)
        Length: 4
       Handshake Protocol: Server Hello Done
```

Client Key Exchange Record:

10. Locate the client key exchange record. Does this record contain a premaster secret? What is this secret used for? Is the secret encrypted? If so, how? How long is the encrypted secret?

Ans:



Yes, the Client Key Exchange record contains a pre-master secret. This secret is used to derive the session keys for encryption and integrity in the SSL/TLS session.

Details:

- 1. **Use of the Pre-Master Secret**: The pre-master secret is combined with the client and server nonces to generate the session keys used for symmetric encryption and message authentication during the session.
- 2. **Encryption of the Pre-Master Secret**: The pre-master secret is encrypted using RSA encryption. This ensures that only the intended recipient (the server) can decrypt it using its private key.
- 3. **Length of the Encrypted Secret**: The length of the encrypted premaster secret is 128 bytes (1024 bits). This length corresponds to the size of the RSA-encrypted data, which typically is larger than the original pre-master secret due to the encryption overhead.

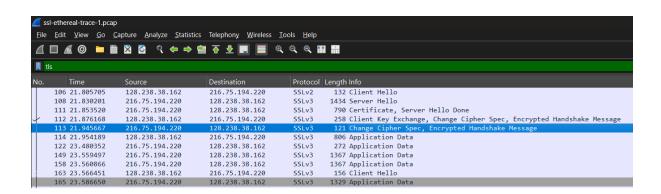
In summary, the Client Key Exchange record does contain an encrypted premaster secret, which is essential for establishing secure communication between the client and server.

```
▶ Frame 112: 258 bytes on wire (2064 bits), 258 bytes captured (2064 bits)
▶ Ethernet II, Src: IBM_10:60:99 (00:09:6b:10:60:99), Dst: All-HSRP-routers_00 (00:00:0c:07:ac:00)
▶ Internet Protocol Version 4, Src: 128.238.38.162, Dst: 216.75.194.220
> Transmission Control Protocol, Src Port: 2271, Dst Port: 443, Seq: 79, Ack: 2785, Len: 204
▼ Transport Layer Security
     SSLv3 Record Layer: Handshake Protocol: Client Key Exchange
Content Type: Handshake (22)
        Version: SSL 3.0 (0x0300)
        Length: 132
     ▼ Handshake Protocol: Client Key Exchange
          Handshake Type: Client Key Exchange (16)
          Length: 128
        ▼ RSA Encrypted PreMaster Secret
             Encrypted PreMaster [...]: bc49494729aa2590477fd059056ae78956c77b12af08b47c609e61f104b0fbf83e...
   ▼ SSLv3 Record Layer: Change Cipher Spec Protocol: Change Cipher Spec
        Content Type: Change Cipher Spec (20)
        Version: SSL 3.0 (0x0300)
        Length: 1
        Change Cipher Spec Message
   SSLv3 Record Layer: Handshake Protocol: Encrypted Handshake Message
        Content Type: Handshake (22)
        Version: SSL 3.0 (0x0300)
        Length: 56
        Handshake Protocol: Encrypted Handshake Message
```

Change Cipher Spec Record (sent by client) and Encrypted Handshake Record:

11. What is the purpose of the Change Cipher Spec record? How many bytes is the record in your trace?

Ans:



Purpose: The Change Cipher Spec record indicates that the sender is ready to switch to encrypted communication using the new cipher suite and keys.

The **Change Cipher Spec** record signals that the sender will start using the newly negotiated cipher suite and keys for subsequent messages. This record is part of the SSL/TLS handshake process and indicates that the sender is ready to switch from using the previous cipher settings to the new ones established during the handshake.

Size: The Change Cipher Spec record in above trace is 121 bytes in total. This includes the Change Cipher Spec message itself, which is 1 byte, plus additional bytes related to the transmission (such as the Ethernet and IP headers).

This transition ensures that the communication is now encrypted using the new parameters established during the handshake.

12. In the encrypted handshake record, what is being encrypted? How?

Ans:

In the Encrypted Handshake Message record, the data being encrypted is the handshake messages that were exchanged during the SSL/TLS handshake process after the Client Key Exchange. This typically includes messages such as the ServerHelloDone message and any other necessary messages required to complete the handshake.

How is it encrypted?

- 1. Encryption Method: The messages are encrypted using the symmetric encryption algorithm specified by the negotiated cipher suite (in this case, TLS_RSA_WITH_RC4_128_MD5).
- 2. Keying Material: The encryption uses the session keys derived from the pre-master secret, along with the client and server nonces. This ensures that the data can only be decrypted by the intended recipient (the server in this case) who possesses the corresponding session key.
- 3. Integrity Check: Additionally, a Message Authentication Code (MAC) is often applied to ensure the integrity and authenticity of the messages, helping to prevent tampering.

Using symmetric-key algorithms determined by the negotiated cipher suite, leveraging session keys derived from the pre-master secret. The messages are encrypted and often accompanied by a MAC for integrity and authenticity.

In summary, the Encrypted Handshake Message record contains encrypted handshake messages that are protected using the newly established symmetric keys and algorithms.

13. Does the server also send a change cipher record and an encrypted handshake record to the client? How are those records different from those sent by the client?

Ans:

Yes, the server also sends a Change Cipher Spec record and an Encrypted Handshake Message record to the client. Here's how these records differ from those sent by the client:

Change Cipher Spec Record:

• Client: The Change Cipher Spec record sent by the client indicates that it is ready to switch to the new cipher suite and keys established during the handshake.

 Server: The server sends a Change Cipher Spec record to notify the client that it is also switching to the new cipher settings. The purpose is the same: to confirm that the server will now use the new encryption parameters.

Encrypted Handshake Message:

- Client: The Encrypted Handshake Message sent by the client contains handshake messages that have been encrypted using the negotiated cipher suite and session keys.
- Server: The server sends its own Encrypted Handshake Message that contains its handshake messages (such as ServerHelloDone) encrypted in the same way.

Key Differences:

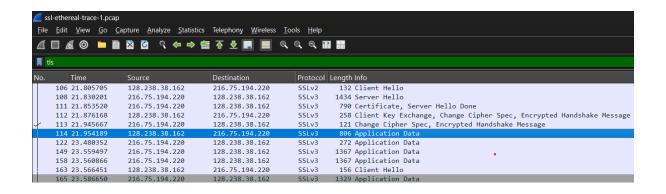
- 1. Content: The content of the records differs. The client's records will contain its specific handshake messages, while the server's records will contain its own messages.
- 2. Direction: The client's Change Cipher Spec and Encrypted Handshake records are sent from the client to the server, while the server's records are sent in the opposite direction.

Overall, while the structure and purpose of the Change Cipher Spec and Encrypted Handshake Message records are consistent between client and server, the actual content and direction of communication differ.

Application Data:

14. How is the application data being encrypted? Do the records containing application data include a MAC? Does Wireshark distinguish between the encrypted application data and the MAC?

Ans:



```
Frame 114: 806 bytes on wire (6448 bits), 806 bytes captured (6448 bits)

Ethernet II, Src: IBM_10:60:99 (00:09:6b:10:60:99), Dst: All-HSRP-routers_00 (00:00:0c:07:ac:00)

Internet Protocol Version 4, Src: 128.238.38.162, Dst: 216.75.194.220

Transmission Control Protocol, Src Port: 2271, Dst Port: 443, Seq: 283, Ack: 2852, Len: 752

Transport Layer Security

SSLv3 Record Layer: Application Data Protocol: Hypertext Transfer Protocol

Content Type: Application Data (23)

Version: SSL 3.0 (0x0300)

Length: 747

Encrypted Application Data [...]: 7e8cdc7fe71d6d59c45ecae7bad064ec705ea592d4b82b35cfc48675c16e461e22-

[Application Data Protocol: Hypertext Transfer Protocol]
```

How is the Application Data being Encrypted?

The application data is encrypted using the symmetric encryption algorithm specified by the negotiated cipher suite (in this case, likely TLS_RSA_WITH_RC4_128_MD5). The encryption is done using the session keys derived during the handshake process, which are based on the premaster secret and the nonces.

Does the Record Containing Application Data Include a MAC?

Yes, records containing application data include a Message Authentication Code (MAC). The MAC ensures the integrity and authenticity of the data, protecting it from tampering and ensuring that it comes from a legitimate source.

Does Wireshark Distinguish Between the Encrypted Application Data and the MAC?

Wireshark typically does not display the MAC as a separate field in the decrypted view of the application data. Instead, the MAC is included in the

encrypted payload. When the encrypted application data is analyzed, Wireshark shows it as a single encrypted block. The MAC is calculated over the plaintext data and is used during decryption to verify the integrity of the received data.

In summary, the application data is encrypted with session keys, includes a MAC for integrity, and Wireshark displays the data as a single encrypted entity without separating the MAC in its output.

15. Comment on and explain anything else that you found interesting in the trace.

Ans:

Use of Different SSL Versions:

The trace indicates a transition from SSLv2 to SSLv3. It's interesting to note the evolution of the SSL protocol versions, as SSLv2 is considered outdated and insecure. Modern applications primarily use TLS, which is the successor to SSL. The presence of SSLv2 could indicate compatibility settings or legacy systems.

Cipher Suite Negotiation:

The ClientHello message lists multiple cipher suites supported by the client. The server chooses one from this list for the session, which can reveal insights into the security posture and configurations of both the client and server. Observing this negotiation process can be critical for understanding potential vulnerabilities.

Challenge and Nonce Usage:

The ClientHello message includes a nonce (challenge), which is a random value used to prevent replay attacks. This is an interesting feature of SSL/TLS that enhances security by ensuring that each

session is unique. The presence of nonces shows the protocols' design to handle specific security threats effectively.

Certificate Exchange:

The certificate exchange step during the ServerHello message and subsequent records is crucial for establishing trust. This trace shows the server providing its certificate, which may be signed by a trusted Certificate Authority (CA). The ability to verify this certificate is essential for the client to ensure that it is communicating with the legitimate server.

Packet Sizes and Performance:

Analyzing the sizes of the packets in the trace could provide insights into network performance. Larger packets may indicate bulk data transfers, while smaller packets might signify many small transactions. Identifying patterns in packet sizes could help in optimizing application performance and network resource utilization.

Timing of Records:

Observing the timing between records can provide insights into latency and performance issues. For example, if there are significant delays between the ClientHello and ServerHello messages, it could indicate network congestion or processing delays.

Application Data Records:

The presence of application data records after the handshake signals that secure communication has commenced. Analyzing the types of application data exchanged can provide insights into the nature of the application traffic, whether it's HTTP requests, file transfers, etc.

Network Security Considerations:

The trace can help identify potential security concerns, such as unencrypted traffic, or weak cipher suites. It is important to ensure

	that strong encryption practices are followed, as vulnerabilities in these areas could lead to exposure of sensitive data.
and	se points provide a deeper understanding of the SSL handshake process the resulting secure communication, illustrating both the complexity and ortance of cryptographic protocols in modern network security.
•	,, 3