

Q.1.

insertion sort

Let

$$A = [10, 7, 3, 8, 1, 9, 0]$$

Step 1 : $A [\underbrace{10, 7}, \underbrace{3, 8, 1, 9, 0}]$

Step 2 : $[\underbrace{7, 10}, \underbrace{3, 8, 1, 9, 0}]$

Step 3 : $[3, \underbrace{7, 10}, \underbrace{8, 1, 9, 0}]$

Step 4 : $[3, 7, \underbrace{8, 10, 1}, 9, 0]$

Step 5 : $[1, 3, 7, 8, \underbrace{10, 9}, 0]$

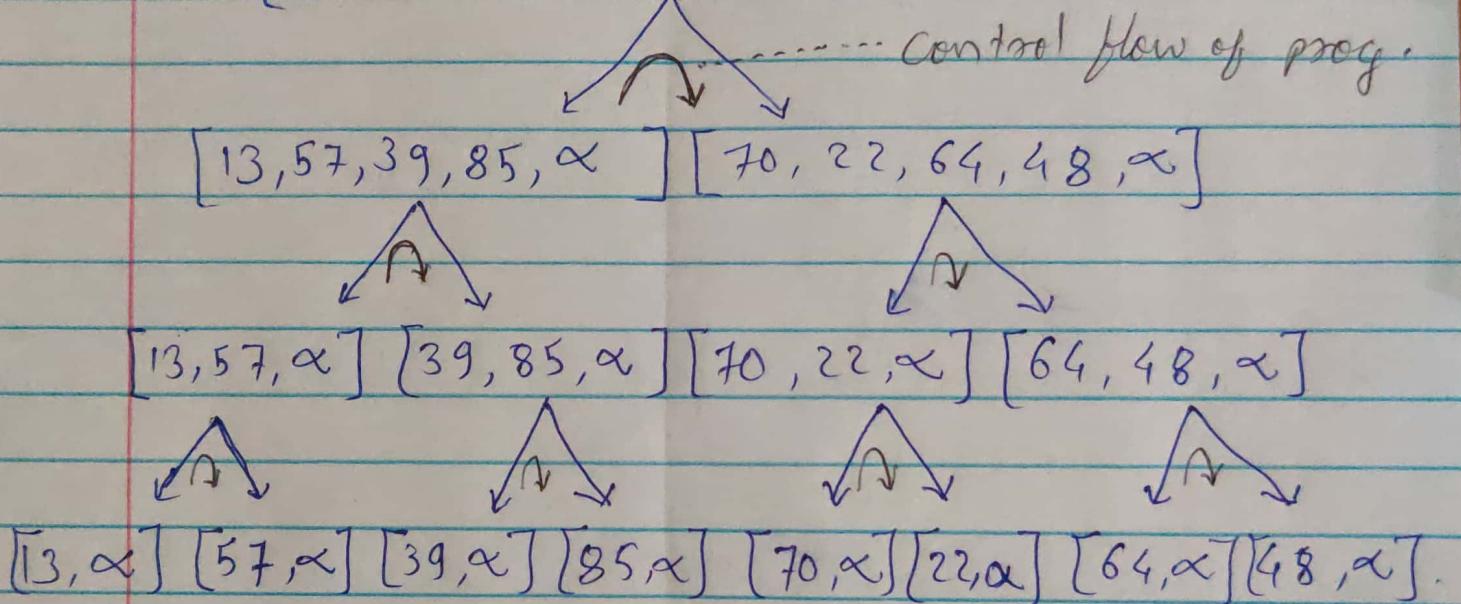
Step 6 : $[1, 3, 7, 8, 9, \underbrace{10, 0}]$

Final Array : $A [0, 1, 3, 7, 8, 9, 10]$

Q. 2. Merge sort algorithm.

1. Divide & conquer stage:

$$A = [13, 57, 39, 85, 70, 22, 64, 48]$$



2. Combine stage:

We will consider only the significant portion of array.
Step 1 : at each step.

$$A' = [13, 57, 39, 85, 70, 22, 64, 48]$$

$$L = [13, \alpha] \quad R = [57, \alpha]$$

Step 2: $A' = [13, 57]$

Step 3 : $A' = [13, 57]$

$$L = [13, \alpha], R = [57, \alpha]$$

$$\underline{\text{step 4}} : A' = [39, 85].$$

$$L = [39, \alpha] \quad R = [85, \alpha]$$

$$\underline{\text{Step 5}} : \quad A^1 = [39, 85]$$

$$L = [39, \alpha] \quad R = [85, \alpha]$$

$$\underline{\text{Step 6}} : A^{-1} = [39, 85]$$

$$L = [39, \infty) \quad R = [85, \infty)$$

Step 7 : $A' = [13, 57, 39, 85]$.

$$L = [13, 57, \alpha] \quad R = [39, 85, \alpha].$$

i *j*

Step 8 :

$$A' = [13, 57, 39, 85]$$

$$L = [13, 57, \alpha] \quad R = [39, 85, \alpha]$$

i *j*

Step 9 : $A' = [13, 39, 39, 85]$

$$L = [13, 57, \alpha] \quad R = [39, 85, \alpha]$$

i *j*

Step 10 : $A' = [13, 39, 57, 85]$

$$L = [13, 57, \alpha] \quad R = [39, 85, \alpha]$$

i *j* ***

Step 11 : $A' = [13, 39, 57, 85]$

$$L = [13, 57, \alpha] \quad R = [39, 85, \alpha]$$

i *j* ***

Step 12 : $A' = [70, 22]$.

$$L = [70, \alpha] \quad R = [22, \alpha]$$

i *j*

Step 13 : $A' = [22, 22]$

$$L = [70, \alpha] \quad R = [22, \alpha]$$

i *j*

Step 14 : $A' = [22, 70]$.

$$L = [70, \alpha] \quad R = [22, \alpha]$$

i *j*

Step 15 : $A' = [64, 48]$

$$L = [64, \alpha] \quad R = [48, \alpha]$$

i *j*

Step 16 $A' = [48, 48]$

$$L = [64, \alpha] \quad R = [48, \alpha]$$

i *j*

Step 17 : $A' = [48, 64]$

$$L = [64, \alpha] \quad R = [48, \alpha]$$

i *j*

Step 18 : $A' = [22, 70, 48, 64]$

$$L = [22, 70, \alpha] \quad R = [48, 64, \alpha]$$

i *j*

Step 19 $A' = [22, 70, 48, 64]$

$$L = [22, 70, \alpha] \quad R = [48, 64, \alpha]$$

i *j*

Step 20 $A' = [22, 48, 48, 64]$

$$L = [22, 70, \alpha] \quad R = [48, 64, \alpha]$$

i *j*

Step 21 : $A' = [22, 48, 64, 64]$

$$L = [22, 70, \alpha] \quad R = [48, 64, \alpha]$$

i *j*

Step 22 $A' = [22, 48, 64, 70]$

$$L = [22, 70] \quad R = [48, 64]$$

i *j*

Step 23 : $A' = A = [13, 39, 57, 85, 22, 48, 64, 70]$

$$L = [13, 39, 57, 85] \quad R = [22, 48, 64, 70]$$

i *j*

Step 24: $A' = [13, 39, 57, 85, 22, 48, 64, 70]$

$$L = [13, 39, 57, 85] \quad R = [22, 48, 64, 70]$$

i *j*

Step 25 $A' = [13, 22, 57, 85, 22, 48, 64, 70]$

$$L = [13, 39, 57, 85] \quad R = [22, 48, 64, 70]$$

i *j*

Step 26 $A = [13, 22, 39, 85, 22, 48, 64, 70]$

$$L = [13, 39, 57, 85] \quad R = [22, 48, 64, 70]$$

i *j*

step 27 $A = [13, 22, 39, 48, 22, 48, 64, 70]$

$L = [13, 39, 57, 85, \alpha]$ $R = [22, 48, 64, 70, \gamma]$

step 28 $A = [13, 22, 39, 48, 57, 48, 64, 70]$

$L = [13, 39, 57, 85, \alpha]$ $R = [22, 48, 64, 70, \gamma]$

step 29 $A = [13, 22, 39, 48, 57, 64, 64, 70]$

$L = [13, 39, 57, 85, \alpha]$ $R = [22, 48, 64, 70, \gamma]$

step 30 $A = [13, 22, 39, 48, 57, 64, 70, 70]$

$L = [13, 39, 57, 85, \alpha]$ $R = [22, 48, 64, 70, \gamma]$

\Rightarrow step 31

$A = [13, 22, 39, 48, 57, 64, 70, 85]$

$L = [13, 39, 57, 85, \alpha]$ $R = [22, 48, 64, 70, \gamma]$

Q.3. Let

$$A = \begin{bmatrix} 2 & 3 \\ 4 & 5 \end{bmatrix}; B = \begin{bmatrix} 5 & 8 \\ 1 & 3 \end{bmatrix}$$

Let the 10 strassen's matrices be:

$$\begin{aligned} S_1 &= B_{12} - B_{22} = 8 - 3 = 5 & S_6 &= B_{11} + B_{22} = 5 + 3 = 8 \\ S_2 &= A_{11} + A_{12} = 2 + 3 = 5 & S_7 &= A_{12} - A_{22} = 3 - 5 = -2 \\ S_3 &= A_{21} + A_{22} = 4 + 5 = 9 & S_8 &= B_{21} + B_{22} = 1 + 3 = 4 \\ S_4 &= B_{21} - B_{11} = 1 - 5 = -4 & S_9 &= A_{11} - A_{21} = 2 - 4 = -2 \\ S_5 &= A_{11} + A_{22} = 2 + 5 = 7 & S_{10} &= B_{11} + B_{12} = 5 + 8 = 13 \end{aligned}$$

Let

$$P_1 = A_{11} \cdot S_1 = 2 \times 5 = 10 \quad P_5 = S_5 \cdot S_6 = 7 \times 8 = 56$$

$$P_2 = S_2 \cdot B_{22} = 5 \times 3 = 15 \quad P_6 = S_7 \cdot S_8 = (-2) \times 4 = -8$$

$$P_3 = S_3 \cdot B_{11} = 9 \times 5 = 45 \quad P_7 = S_9 \cdot S_{10} = (-2) \times 13 = -26$$

$$P_4 = A_{22} \cdot S_4 = 5 \times (-4) = -20$$

Let $C_{11} = P_5 + P_4 + P_6 - P_2 = 56 - 20 - 8 - 15 =$

~~$C_{12} = P_1 + P_2 = 10 + 15 =$~~ ; $C_{21} = P_3 + P_4 = 45 - 20 =$

~~$C_{22} = P_5 + P_1 - P_3 - P_7 = 56 + 10 - 45 + 2 =$~~

let

$$C_{11} = P_5 + P_4 + P_6 - P_2$$

$$= 56 + (-20) + (-8) - 15 = \boxed{13 = C_{11}}$$

$$\begin{aligned} C_{12} &= P_1 + P_2 \\ &= 10 + 15 \end{aligned}$$

$$= \boxed{25 = C_{12}}.$$

$$\begin{aligned} C_{21} &= P_3 + P_4 \\ &= 45 - 20 \end{aligned}$$

$$= \boxed{25 = C_{21}}.$$

$$C_{22} = P_5 + P_1 - P_3 - P_7$$

$$= 56 + 10 - 45 - (-26) = \boxed{47 = C_{22}}$$

$$\therefore C = A \cdot B$$

$$\therefore \boxed{C = \begin{bmatrix} 13 & 25 \\ 25 & 47 \end{bmatrix}}$$

Q. 4.

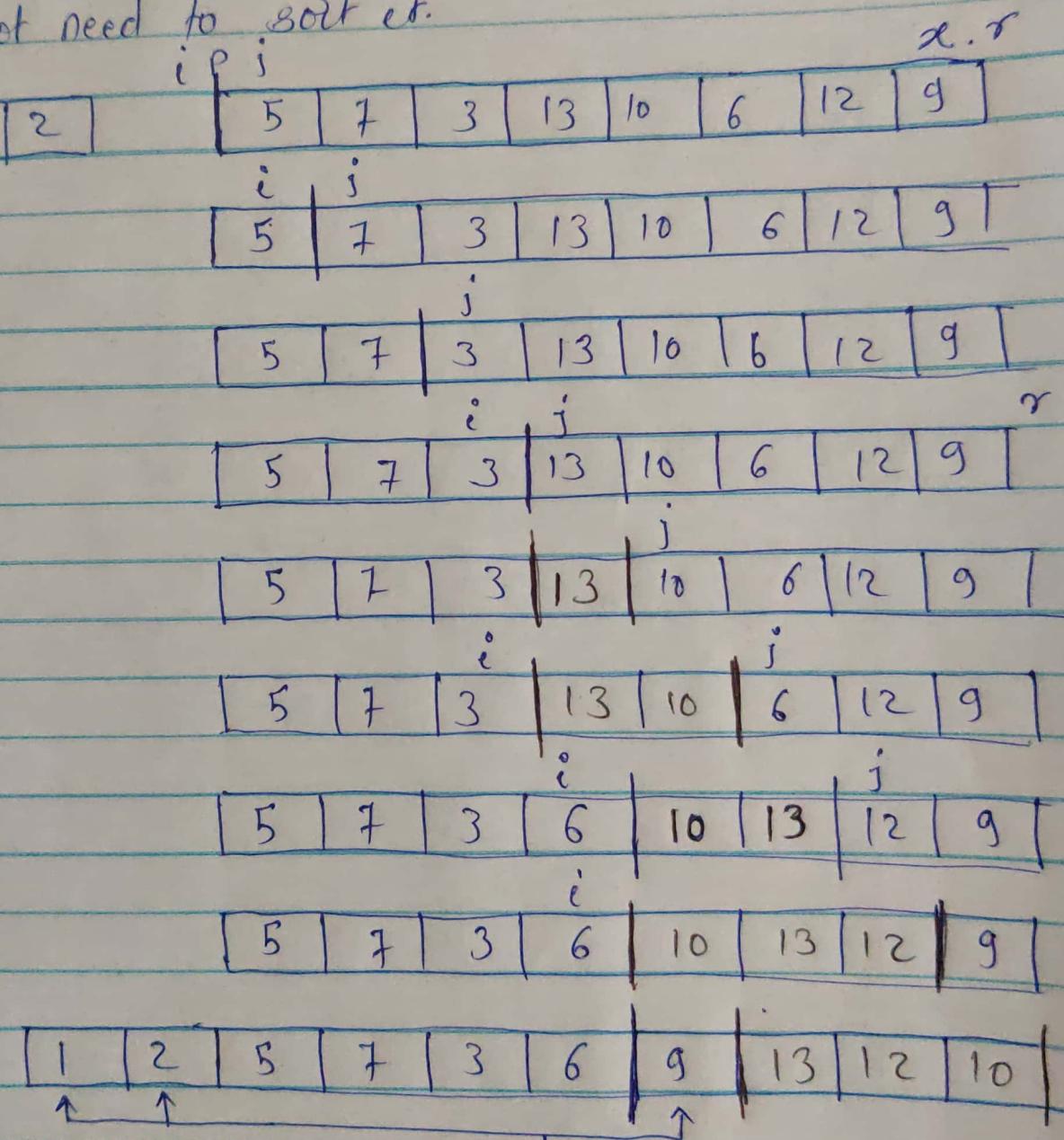
13	9	5	7	3	1	10	6	12	2
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Let

$$i = -1, p = 0, j = 0, r = 9, x = A[9] = 2$$

i	p	j	r						
13	9	5	7	3	1	10	6	12	2
i	j								
13	9	5	7	3	1	10	6	12	2
i		i							
13	9	5	7	3	1	10	6	12	2
i		j							
13	9	5	7	3	1	10	6	12	2
i			j						
13	9	5	7	3	1	10	6	12	2
i	p	-	j						
13	9	5	7	3	1	10	6	12	2
p	i								
1	9	5	7	3	13	10	6	12	2
i					j				r
1	9	5	7	3	13	10	6	12	2
i						j			
1	9	5	7	3	13	10	6	12	2
	A[i] < A[r] ← A[j] > A[r]. → unsorted A[r].								
1	9	5	7	3	13	10	6	12	2
1	2	5	7	3	13	10	6	12	9

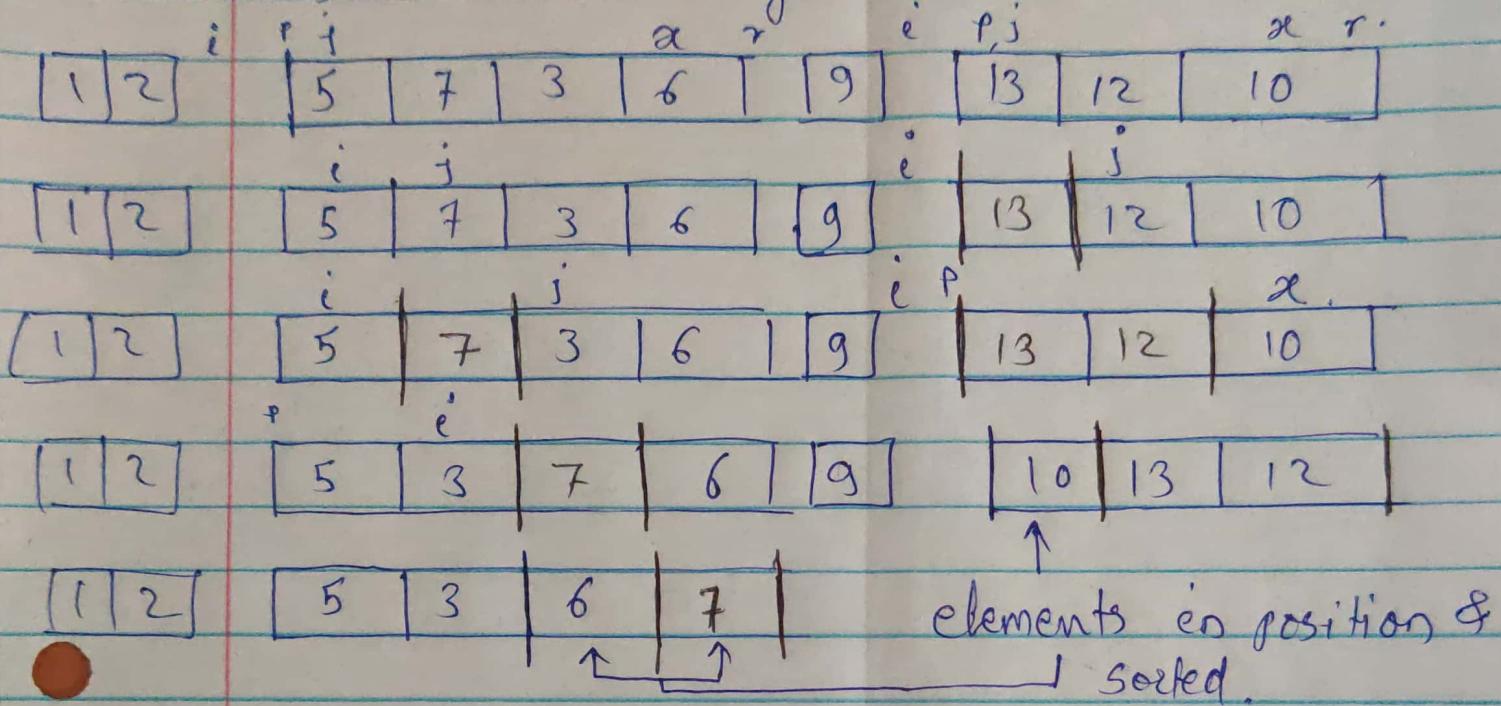
As there's only one item element on left side, we do not need to sort it.



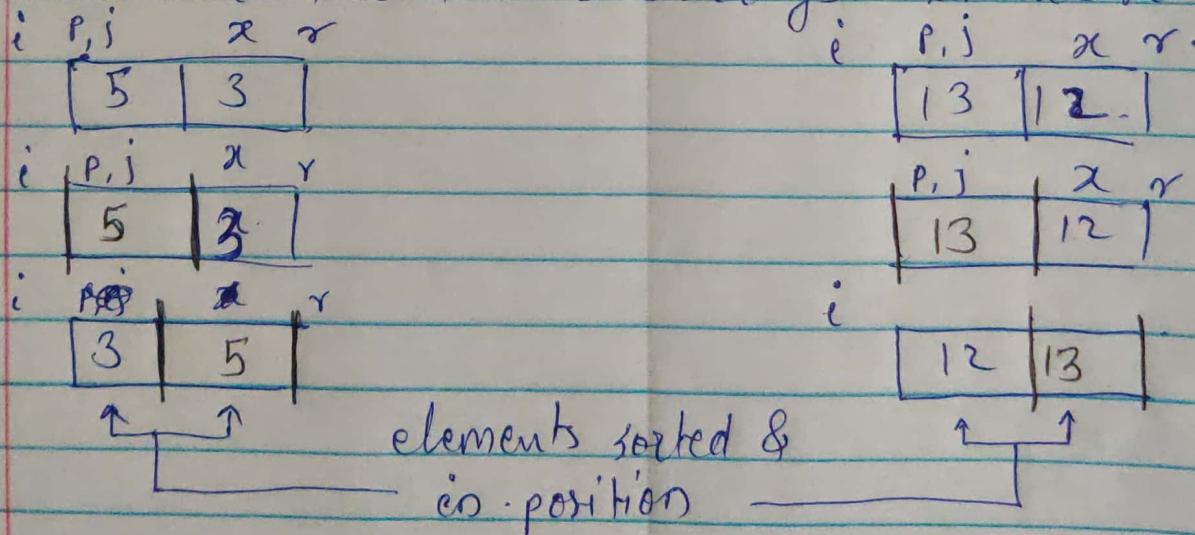
sorted.

elements in position

Consider two sub arrays that are not sorted



Consider the two sub arrays that are not sorted



\therefore Final Array:

1	2	3	5	6	7	9	10	12	13
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Q.6. Let x be the number on the face of the die.

It is given that

$$P(X=x) \propto \frac{1}{x}.$$

$$\therefore P(X=x) = K \cdot \frac{1}{x}.$$

Also.

$$P(X > 3) = P(X = \{4, 5, 6\})$$

$$\therefore P(X > 3) = P(4) + P(5) + P(6)$$

$$= K \cdot \frac{1}{4} + K \cdot \frac{1}{5} + K \cdot \frac{1}{6}$$

$$= K \cdot \left(\frac{1}{4} + \frac{1}{5} + \frac{1}{6} \right)$$

$$= K \left(\frac{15}{60} + \frac{12}{60} + \frac{10}{60} \right).$$

$$= K \cdot \left(\frac{37}{60} \right). \quad - \textcircled{1}.$$

we also know that $\sum P(X=x) = 1$

$$\therefore \sum_{x=0}^1 P(X=x) = 1$$

$$\therefore K \cdot \left(\frac{1}{1} + \frac{1}{2} + \frac{1}{3} + \frac{1}{4} \right) + \frac{1}{5} + \frac{1}{6} = 1$$

$$\therefore K \left(\frac{60}{60} + \frac{30}{60} + \frac{20}{60} + \frac{15}{60} + \frac{12}{60} + \frac{10}{60} \right) = 1$$

$$\therefore \boxed{K = \frac{60}{147}} \quad -\textcircled{2}$$

putting eqⁿ $\textcircled{2}$ in $\textcircled{1}$.

$$\therefore \frac{60}{147} \times \frac{37}{60}$$

$$\therefore \boxed{P(X \geq 3) = 0.25170068}$$

Q. 5. Let A be the sorted arry such that

$$A = [a_1, a_2, a_3, \dots, a_n].$$

$$\text{size}[A] = n \quad \& \quad a_1 \leq a_2 \leq a_3 \leq \dots \leq a_n.$$

In order to get the algorithmic complexity $O(\log n)$, i will use binary search.

Algorithm:

for every index i in array:

$$\text{middle} = \lfloor (\text{size}/2) \rfloor$$

if size of array = 1

check if $A[\text{middle}] = \text{i}$

return 1.

else, if size > 1

if $A[\text{middle}] > \text{i}$

search left subarray $[a_1 \text{ to } a_{\text{middle}}]$
middle

if $A[\text{middle}] < \text{i}$:

search subarray $[a_{\text{middle}+1} \text{ to } a_n]$

if $A[\text{middle}] = \text{i}$:

if middle = ~~key~~ index
return index.

* Pseudocode:

for i = 1 to n

left = 1

right = 1

mid = $\lfloor (left + right) / 2 \rfloor$

while (left <= right)

: if (A[mid] < key)

: left = mid + 1

: elseif (A[mid] = i AND mid = i)

: found = i

: else

: : right = mid - 1

: : mid = $\lfloor (left + right) / 2 \rfloor$

: return found.