## HW9 (CSCI-C241)

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- 1. Question One
  - (a)  $\forall x ((A(x) \lor B(x)) \to (C(x) \land \neg D(x)))$
  - (b)  $\neg \exists x (A(x) \land B(x))$
- 2. Question Two
  - (a)  $U_1, U_2, U_5, U_6, U_7$
  - (b)  $U_1, U_6, U_7, U_8$
  - (c)  $\mathcal{U}_7$
  - (d)  $\mathcal{U}_3, \mathcal{U}_4, \mathcal{U}_8$
  - (e)  $U_2, U_3, U_4, U_5$
  - (f)  $\mathcal{U}_2, \mathcal{U}_3, \mathcal{U}_4, \mathcal{U}_5$
  - (g)  $\mathcal{U}_3, \mathcal{U}_4, \mathcal{U}_8$
  - (h) Part (h)
    - i.  $\neg \forall x (L(x) \to V(x))$  and  $\exists x (L(x) \land \neg V(x))$
    - ii.  $\forall x(L(x) \to \neg V(x))$  and  $\neg \exists x(L(x) \land V(x))$
- 3. Question Three



- (b) This is not possible as,  $\exists y \forall x P(x,y)$  states that there exists a shape x such that every shape points to, but on the other hand  $\forall x \exists y P(x,y)$  which must be false, states that for every shape, there exists a shape y such that the first shape points to, meaning not every shape can point to another. With that being said, if not every shape can point to another, then a shape that has every shape pointing to itself can't exist.
- (c)  $\mathcal{M}_1, \mathcal{M}_2, \mathcal{M}_3$
- (d)  $\mathcal{M}_2$
- (e)  $M_{10}$
- 4. Question Four
  - (a)  $\forall x \forall y \forall z (\exists w R(x, w) \land ((R(x, y) \land R(x, z)) \rightarrow y = z))$
- 5. Question Five
  - (a)  $\forall x (A(x) \rightarrow ((\exists y (B(y) \land R(x,y))) \land (\forall y \forall z ((B(y) \land B(z)) \rightarrow ((R(x,y) \land R(x,z)) \rightarrow y = z)))))$
- 6. Question Six
  - (a)  $5! = 1 \cdot 2 \cdot 3 \cdot 4 \cdot 5 = 120$
  - (b) n-1
  - (c) n+1
  - (d) n

(e) 
$$(k+1) \cdot k!$$

## 7. Question Seven

(a) Claim:  $n \ge 2, 3^n > 2^{n+1}$ 

Proof. (induction on n) (Base Step, n = 2):

$$3^2 = 9 \tag{1}$$

$$> 8$$
 (2)

$$= 2^3 (3)$$

$$= 2^{2+1}$$
 (4)

(Inductive Step):

Assume  $3^k > 2^{k+1}$  for some  $k \ge 2$ 

$$3^{k+1} = 3^k \cdot 3^1 \tag{1}$$

$$= 3^k \cdot 3 \tag{2}$$

$$> 2^{k+1} \cdot 2$$
 (by the induction hypothesis, and  $3 > 2$ ) (3)

$$= 2^{k+1+1}$$
 (4)

(b) Claim:  $n \ge 9$ ,  $3^n < (n-1)!$ 

Proof. (induction on n) (Base Step, n = 9):

$$3^9 = 19683 \tag{1}$$

$$< 40320$$
 (2)

$$= 8! (3)$$

$$= (9-1)!$$
 (4)

(Inductive Step):

Assume  $3^k < (k-1)!$  for some  $k \ge 9$ 

$$3^{k+1} = 3^k \cdot 3 \tag{1}$$

$$(k-1)! \cdot k$$
 (by the induction hypothesis, and  $k > 3$ )

$$= k! \tag{3}$$

$$= (k-1+1)! (4)$$

(c) Claim:  $n \ge 2, 3^n > n^2$ 

Proof. (induction on n) (Base Step, n = 2):

$$3^2 = 9 \tag{1}$$

$$> 4$$
 (2)

$$= 2^2 \tag{3}$$

(Inductive Step):

Assume  $3^k > k^2$  for some  $k \ge 2$ 

$$3^{k+1} = 3^k \cdot 3 \tag{1}$$

$$> k^2 \cdot 3$$
 (by the induction hypothesis) (2)

$$= k^2 + k^2 + k^2 \tag{3}$$

$$\geq k^2 + 2k + 1$$
  $(k^2 \geq 2k > 1 \text{ for all } k \geq 2)$  (4)

$$= (k+1)(k+1) (5)$$

$$= (k+1)^2 \tag{6}$$

(d) Claim:  $n^2 - 3n$  is even for all  $n \in \mathbb{N}$ 

*Proof.* (induction on n) (Base Step, n = 0):

$$0^{2} - 3(0) = 0 - 0$$

$$= 0$$
(1)
$$= 0$$
(2)

$$= 0$$
 (2)

0 is even

(Inductive Step):

Assume  $k^2 - 3k$  is even for all  $k \in \mathbb{N}$ 

 $k^2 - 3k = 2c_1$ 

$$(k+1)^2 - 3(k+1) = (k+1)(k+1) - 3k - 3$$
(1)

$$= k^2 + 2k + 1 - 3k - 3 \tag{2}$$

$$= (k^2 - 3k) + (2k - 2) (3)$$

$$= (k^2 - 3k) + 2(k - 1) \tag{4}$$

= 
$$2c_1 + 2(k-1)$$
 (by the induction hypothesis) (5)

$$= 2c_1 + 2c_2 \qquad \text{(Let } c_2 = k - 1\text{)} \tag{6}$$

$$= 2(c_1 + c_2) (7)$$

$$= 2c_3 \qquad \text{(Let } c_3 = c_1 + c_2) \tag{8}$$

 $(2n \text{ is even, where } n \in \mathbb{N})$ (9) $2c_3$  is even

(10)

Since  $2c_3$  is even,  $(k+1)^2 - 3(k+1)$  is even

(e) Claim: There are  $2^n$  binary string of length n for all  $n \in \mathbb{N}$ 

*Proof.* (induction on n) (Base Step, n = 0):

$$2^0 = 1 \tag{1}$$

There is only one possible binary string, of length 0, the empty string. (Inductive Step):

Assume for length k, there are  $2^k$  binary strings

For every binary string of length k, there are  $2^k$  binary strings, (1)

for each binary string, there are two new possible binary strings of length k+1, (2)

So, by the induction hypothesis there are  $2^k \cdot 2$  binary strings of length k+1(3)

(f) Claim: For any set of characters, with the length of a and any  $n \in \mathbb{N}$ , there are  $a^n$  possible strings of length n that explicitly use the said alphabet

*Proof.* (induction on n) (Base Step, n=0):

$$a^0 = 1 (1)$$

Regardles of the possible characters, there is only one possible string, of length 0, the empty string.

(Inductive Step):

Assume for length k, there are  $a^k$  strings, where a is total amount of different characters

For every string of length k, there are a possible different characters, (1)

and 
$$a^k$$
 possible strings (2)

For every string of length k+1, with the same a possible characters, (3)

for each string, there would a possible new strings of length k+1(4)

So, by the induction hypothesis there are  $a^k \cdot a$  strings of length k+1(5)