
MODULE *Config*

INSTANCE *Naturals*

INSTANCE *FiniteSets*

INSTANCE *Sequences*

INSTANCE *TLC*

An empty constant

CONSTANT *Nil*

Transaction status constants

CONSTANTS

TransactionPending,
TransactionValidating,
TransactionApplying,
TransactionComplete,
TransactionFailed

Configuration status constants

CONSTANTS

ConfigurationPending,
ConfigurationInitializing,
ConfigurationUpdating,
ConfigurationComplete,
ConfigurationFailed

The set of all nodes

CONSTANT *Node*

Target is the possible targets, paths, and values

Example: $Target \triangleq$ [
 $target1 \mapsto$ [
 $path1 \mapsto \{“value1”, “value2”\}$,
 $path2 \mapsto \{“value2”, “value3”\}$],
 $target2 \mapsto$ [
 $path2 \mapsto \{“value3”, “value4”\}$,
 $path3 \mapsto \{“value4”, “value5”\}$]]

CONSTANT *Target*

ASSUME *Nil* ∈ STRING

ASSUME *TransactionPending* ∈ STRING

ASSUME *TransactionValidating* ∈ STRING

ASSUME *TransactionApplying* ∈ STRING

ASSUME $TransactionComplete \in \text{STRING}$
 ASSUME $TransactionFailed \in \text{STRING}$

 ASSUME $ConfigurationPending \in \text{STRING}$
 ASSUME $ConfigurationInitializing \in \text{STRING}$
 ASSUME $ConfigurationUpdating \in \text{STRING}$
 ASSUME $ConfigurationComplete \in \text{STRING}$
 ASSUME $ConfigurationFailed \in \text{STRING}$

 ASSUME $\wedge IsFiniteSet(Node)$
 $\wedge \forall n \in Node :$
 $\wedge n \notin \text{DOMAIN } Target$
 $\wedge n \in \text{STRING}$

 ASSUME $\wedge \forall t \in \text{DOMAIN } Target :$
 $\wedge t \notin Node$
 $\wedge t \in \text{STRING}$
 $\wedge \forall p \in \text{DOMAIN } Target[t] :$
 $IsFiniteSet(Target[t][p])$

TYPE $TransactionStatus ::= status \in$
 $\{TransactionPending,$
 $TransactionValidating,$
 $TransactionApplying,$
 $TransactionComplete,$
 $TransactionFailed\}$

 TYPE $Transaction \triangleq [$
 $id \quad ::= id \in \text{STRING},$
 $index ::= index \in Nat,$
 $revision ::= revision \in Nat,$
 $atomic ::= atomic \in \text{BOOLEAN} ,$
 $sync \quad ::= sync \in \text{BOOLEAN} ,$
 $changes ::= [target \in \text{SUBSET } (\text{DOMAIN } Target) \mapsto [$
 $path \in \text{SUBSET } (\text{DOMAIN } Target[target]) \mapsto [$
 $value ::= value \in \text{STRING},$
 $delete ::= delete \in \text{BOOLEAN }]],$
 $status ::= status \in TransactionStatus]$

 TYPE $ConfigurationStatus ::= status \in$
 $\{ConfigurationPending,$
 $ConfigurationInitializing,$
 $ConfigurationUpdating,$
 $ConfigurationComplete,$
 $ConfigurationFailed\}$

 TYPE $Configuration \triangleq [$
 $id \quad ::= id \in \text{STRING},$
 $revision ::= revision \in Nat,$

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target ::= target ∈ STRING,
paths ::= [ path ∈ SUBSET (DOMAIN Target[target]) ↦ [
    value ::= value ∈ STRING,
    index ::= index ∈ Nat,
    deleted ::= delete ∈ BOOLEAN ]],
txIndex ::= txIndex ∈ Nat,
syncIndex ::= syncIndex ∈ Nat,
term ::= term ∈ Nat,
status ::= status ∈ ConfigurationStatus]

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A sequence of transactions

Each transactions contains a record of 'changes' for a set of targets

VARIABLE *transaction*

A record of target configurations

Each configuration represents the desired state of the target

VARIABLE *configuration*

A record of target states

VARIABLE *targets*

A record of target masters

VARIABLE *masters*

VARIABLE *history*

$\text{vars} \triangleq \langle \text{transaction}, \text{configuration}, \text{masters}, \text{targets}, \text{history} \rangle$

$\text{ChangeMaster}(n, t) \triangleq$
 $\wedge \text{masters}[t].\text{master} \neq n$
 $\wedge \text{masters}' = [\text{masters} \text{ EXCEPT } ![t].\text{term} = \text{masters}[t].\text{term} + 1,$
 $\phantom{\wedge \text{masters}' = } ![t].\text{master} = n]$
 $\wedge \text{UNCHANGED } \langle \text{transaction}, \text{configuration}, \text{targets}, \text{history} \rangle$

This section models the northbound *API* for the configuration service.

$\text{ValidValues}(t, p) \triangleq$
 $\text{UNION } \{ \{ [value \mapsto v, delete \mapsto \text{FALSE}] : v \in \text{Target}[t][p] \}, \{ [value \mapsto \text{Nil}, delete \mapsto \text{TRUE}] \} \}$
 $\text{ValidPaths}(t) \triangleq$
 $\text{UNION } \{ \{ v @@@ [path \mapsto p] : v \in \text{ValidValues}(t, p) \} : p \in \text{DOMAIN Target}[t] \}$
 $\text{ValidTargets} \triangleq$
 $\text{UNION } \{ \{ p @@@ [target \mapsto t] : p \in \text{ValidPaths}(t) \} : t \in \text{DOMAIN Target} \}$
 $\text{ValidPath}(s, t, p) \triangleq$
 $\text{LET } value \triangleq \text{CHOOSE } v \in s : v.\text{target} = t \wedge v.\text{path} = p$

IN
 $[value \mapsto value.value,$
 $delete \mapsto value.delete]$

$ValidTarget(s, t) \triangleq$
 $[p \in \{v.path : v \in \{v \in s : v.target = t \wedge PrintT(t)\}\} \mapsto ValidPath(s, t, p)]$

$ValidChange(s) \triangleq$
 $[t \in \{v.target : v \in s\} \mapsto ValidTarget(s, t)]$

$ValidChanges \triangleq$
 LET $changeSets \triangleq \{s \in SUBSET ValidTargets :$
 $\quad \forall t \in DOMAIN Target :$
 $\quad \forall p \in DOMAIN Target[t] :$
 $\quad Cardinality(\{v \in s : v.target = t \wedge PrintT(t) \wedge v.path = p\}) \leq 1\}$

IN
 $\{ValidChange(s) : s \in changeSets\}$

$NextIndex \triangleq$
 IF $Len(transaction) = 0$ THEN
 1
 ELSE
 LET $i \triangleq CHOOSE i \in DOMAIN transaction :$
 $\quad \forall j \in DOMAIN transaction :$
 $\quad \quad transaction[j].index \leq transaction[i].index$
 IN $transaction[i].index + 1$

Add a set of changes to the transaction log
 $Change \triangleq$
 $\wedge \exists changes \in ValidChanges :$
 $\quad \wedge transaction' = Append(transaction, [index \mapsto NextIndex,$
 $\quad \quad \quad atomic \mapsto FALSE,$
 $\quad \quad \quad sync \mapsto FALSE,$
 $\quad \quad \quad changes \mapsto changes,$
 $\quad \quad \quad status \mapsto TransactionPending])$
 $\wedge UNCHANGED \langle configuration, masters, targets, history \rangle$

This section models the Transaction log reconciler.

Reconcile the transaction log
 $ReconcileTransaction(n, t) \triangleq$
 If the transaction is *Pending*, begin validation if the prior transaction
 has already been applied. This simplifies concurrency control in the controller
 and guarantees transactions are applied to the configurations in sequential order.
 $\wedge \forall \wedge transaction[t].status = TransactionPending$

$\wedge \vee \wedge \text{transaction}[t].\text{index} - 1 \in \text{DOMAIN } \text{transaction}$
 $\wedge \text{transaction}[\text{transaction}[t].\text{index} - 1].\text{status} \in \{\text{TransactionComplete}, \text{TransactionFailed}\}$
 $\vee \text{transaction}[t].\text{index} - 1 \notin \text{DOMAIN } \text{transaction}$
 $\wedge \text{transaction}' = [\text{transaction} \text{ EXCEPT } ![t].\text{status} = \text{TransactionValidating}]$
 $\wedge \text{UNCHANGED } \langle \text{configuration}, \text{history} \rangle$
 If the transaction is in the *Validating* state, compute and validate the
 Configuration for each target.
 $\vee \wedge \text{transaction}[t].\text{status} = \text{TransactionValidating}$
 If validation fails any target, mark the transaction *Failed*.
 If validation is successful, proceed to *Applying*.
 $\wedge \exists \text{valid} \in \text{BOOLEAN} :$
 $\vee \wedge \text{valid}$
 $\wedge \text{transaction}' = [\text{transaction} \text{ EXCEPT } ![t].\text{status} = \text{TransactionApplying}]$
 $\vee \wedge \neg \text{valid}$
 $\wedge \text{transaction}' = [\text{transaction} \text{ EXCEPT } ![t].\text{status} = \text{TransactionFailed}]$
 $\wedge \text{UNCHANGED } \langle \text{configuration}, \text{history} \rangle$
 If the transaction is in the *Applying* state, update the Configuration for each
 target and *Complete* the transaction.
 $\vee \wedge \text{transaction}[t].\text{status} = \text{TransactionApplying}$
 $\wedge \vee \wedge \text{transaction}[t].\text{atomic}$
TODO: Apply atomic transactions here
 $\wedge \text{transaction}' = [\text{transaction} \text{ EXCEPT } ![t].\text{status} = \text{TransactionComplete}]$
 $\wedge \text{UNCHANGED } \langle \text{configuration}, \text{history} \rangle$
 $\vee \wedge \neg \text{transaction}[t].\text{atomic}$
 Add the transaction index to each updated path
 $\wedge \text{configuration}' = [$
 $\quad r \in \text{DOMAIN } \text{Target} \mapsto [$
 $\quad \quad \text{configuration}[r] \text{ EXCEPT}$
 $\quad \quad \quad !.\text{paths} = [\text{path} \in \text{DOMAIN } \text{transaction}[t].\text{changes} \mapsto$
 $\quad \quad \quad \quad \text{transaction}[t].\text{changes}[\text{path}] @@ [\text{index} \mapsto \text{transaction}[t].\text{index}]$
 $\quad \quad \quad @@ \text{configuration}[t].\text{paths},$
 $\quad \quad \quad !.\text{txIndex} = \text{transaction}[t].\text{index},$
 $\quad \quad \quad !.\text{status} = \text{ConfigurationPending}]$
 $\wedge \text{history}' = [r \in \text{DOMAIN } \text{Target} \mapsto \text{Append}(\text{history}[r], \text{configuration}'[r])]$
 $\wedge \text{transaction}' = [\text{transaction} \text{ EXCEPT } ![t].\text{status} = \text{TransactionComplete}]$
 $\wedge \text{UNCHANGED } \langle \text{masters}, \text{targets} \rangle$

This section models the Configuration reconciler.

$\text{ReconcileConfiguration}(n, c) \triangleq$

$\wedge \vee \wedge \text{configuration}[c].\text{status} = \text{ConfigurationPending}$

If the configuration is marked *ConfigurationPending* and mastership
 has changed (indicated by an increased mastership term), mark the
 configuration *ConfigurationInitializing* to force full re-synchronization.

$\wedge \vee \wedge \text{masters}[\text{configuration}[c].\text{target}].\text{term} > \text{configuration}[c].\text{term}$
 $\wedge \text{configuration}' = [\text{configuration} \text{ EXCEPT } ![c].\text{status} = \text{ConfigurationInitializing},$
 $\hspace{15em} ![c].\text{term} = \text{masters}[\text{configuration}[c].\text{target}].\text{term}]$
 $\wedge \text{history}' = [\text{history} \text{ EXCEPT } ![c] = \text{Append}(\text{history}[c], \text{configuration}'[c])]$
 If the configuration is marked *ConfigurationPending* and the values have
 changed (determined by comparing the transaction index to the last *sync*
 index), mark the configuration *ConfigurationUpdating* to push the changes
 to the target.
 $\vee \wedge \text{configuration}[c].\text{syncIndex} < \text{configuration}[c].\text{txIndex}$
 $\wedge \text{configuration}' = [\text{configuration} \text{ EXCEPT } ![c].\text{status} = \text{ConfigurationUpdating}]$
 $\wedge \text{history}' = [\text{history} \text{ EXCEPT } ![c] = \text{Append}(\text{history}[c], \text{configuration}'[c])]$
 $\wedge \text{UNCHANGED } \langle \text{targets} \rangle$
 $\vee \wedge \text{configuration}[c].\text{status} = \text{ConfigurationInitializing}$
 $\wedge \text{masters}[\text{configuration}[c].\text{target}].\text{master} = n$
 Merge the configuration paths with the target paths, removing paths
 that have been marked deleted
 $\wedge \text{LET } \text{deletePaths} \triangleq \{p \in \text{DOMAIN } \text{configuration}[c].\text{paths} : \text{configuration}[c].\text{paths}[p].\text{deleted}\}$
 $\hspace{2em} \text{configPaths} \triangleq \text{DOMAIN } \text{configuration}[c].\text{paths} \setminus \text{deletePaths}$
 $\hspace{2em} \text{targetPaths} \triangleq \text{DOMAIN } \text{targets}[\text{configuration}[c].\text{target}] \setminus \text{deletePaths}$
 IN
 $\wedge \text{targets}' = [\text{targets} \text{ EXCEPT } ![c].\text{target} =$
 $\hspace{2em} [p \in \text{configPaths} \mapsto [\text{value} \mapsto \text{configuration}[c].\text{paths}[p]]]$
 $\hspace{2em} @ @ [p \in \text{targetPaths} \mapsto \text{targets}[\text{configuration}[c].\text{target}][p]]]$
 Set the configuration's status to *Complete*
 $\wedge \text{configuration}' = [\text{configuration} \text{ EXCEPT } ![c].\text{status} = \text{ConfigurationComplete},$
 $\hspace{15em} ![c].\text{syncIndex} = \text{configuration}[c].\text{txIndex}]$
 $\wedge \text{history}' = [\text{history} \text{ EXCEPT } ![c] = \text{Append}(\text{history}[c], \text{configuration}'[c])]$
 If the configuration is marked *ConfigurationUpdating*, we only need to
 push paths that have changed since the target was initialized or last
 updated by the controller. The set of changes made since the last
 synchronization are identified by comparing the index of each path-value
 to the last synchronization index, *syncIndex*
 $\vee \wedge \text{configuration}[c].\text{status} = \text{ConfigurationUpdating}$
 $\wedge \text{masters}[\text{configuration}[c].\text{target}].\text{master} = n$
 Compute the set of updated and deleted paths by comparing
 their indexes to the target's last sync index.
 $\wedge \text{LET } \text{updatePaths} \triangleq \{p \in \text{DOMAIN } \text{configuration}[c].\text{paths} :$
 $\hspace{10em} \text{configuration}[c].\text{paths}[p].\text{index} > \text{configuration}[c].\text{syncIndex}\}$
 $\hspace{2em} \text{deletePaths} \triangleq \{p \in \text{updatePaths} : \text{configuration}[c].\text{paths}[p].\text{deleted}\}$
 $\hspace{2em} \text{configPaths} \triangleq \text{updatePaths} \setminus \text{deletePaths}$
 $\hspace{2em} \text{targetPaths} \triangleq \text{DOMAIN } \text{targets}[\text{configuration}[c].\text{target}] \setminus \text{deletePaths}$
 IN
 Update the target paths by adding/updated paths that have changed and
 removing paths that have been deleted since the last *sync*.
 $\wedge \text{targets}' = [\text{targets} \text{ EXCEPT } ![c].\text{target} =$

$$\begin{aligned}
& [p \in \text{configPaths} \mapsto \text{configuration}[c].\text{paths}[p]] \\
& @@ [p \in \text{targetPaths} \mapsto \text{targets}[\text{configuration}[c].\text{target}][p]] \\
& \wedge \text{configuration}' = [\text{configuration} \text{ EXCEPT } ![c].\text{status} = \text{ConfigurationComplete}, \\
& \quad \quad \quad ![c].\text{syncIndex} = \text{configuration}[c].\text{txIndex}] \\
& \wedge \text{history}' = [\text{history} \text{ EXCEPT } ![c] = \text{Append}(\text{history}[c], \text{configuration}'[c])] \\
& \text{If the configuration is not already } \text{ConfigurationPending} \text{ and mastership} \\
& \text{has been lost revert it. This can occur when the connection to the} \\
& \text{target has been lost and the mastership is no longer valid.} \\
& \text{TODO: We still need to model mastership changes} \\
& \vee \wedge \text{configuration}[c].\text{status} \neq \text{ConfigurationPending} \\
& \quad \wedge \text{masters}[\text{configuration}[c].\text{target}].\text{master} = \text{Nil} \\
& \quad \wedge \text{configuration}' = [\text{configuration} \text{ EXCEPT } ![c].\text{status} = \text{ConfigurationPending}] \\
& \quad \wedge \text{history}' = [\text{history} \text{ EXCEPT } ![c] = \text{Append}(\text{history}[c], \text{configuration}'[c])] \\
& \quad \wedge \text{UNCHANGED } \langle \text{targets} \rangle \\
& \wedge \text{UNCHANGED } \langle \text{transaction}, \text{masters} \rangle
\end{aligned}$$

Init and next state predicates

$$\begin{aligned}
\text{Init} & \triangleq \\
& \wedge \text{transaction} = \langle \rangle \\
& \wedge \text{configuration} = [t \in \text{DOMAIN } \text{Target} \mapsto \\
& \quad [target \mapsto t, \\
& \quad \quad paths \mapsto \\
& \quad \quad [path \in \{\} \mapsto \\
& \quad \quad \quad [path \mapsto path, \\
& \quad \quad \quad \quad value \mapsto \text{Nil}, \\
& \quad \quad \quad \quad index \mapsto 0, \\
& \quad \quad \quad \quad deleted \mapsto \text{FALSE}], \\
& \quad \quad txIndex \mapsto 0, \\
& \quad \quad syncIndex \mapsto 0, \\
& \quad \quad term \mapsto 0, \\
& \quad \quad status \mapsto \text{ConfigurationPending}]] \\
& \wedge \text{targets} = [t \in \text{DOMAIN } \text{Target} \mapsto \\
& \quad [path \in \{\} \mapsto \\
& \quad \quad [value \mapsto \text{Nil}]]] \\
& \wedge \text{masters} = [t \in \text{DOMAIN } \text{Target} \mapsto [master \mapsto \text{Nil}, term \mapsto 0]] \\
& \wedge \text{history} = [t \in \text{DOMAIN } \text{Target} \mapsto \langle \rangle]
\end{aligned}$$

$$\begin{aligned}
\text{Next} & \triangleq \\
& \vee \text{Change} \\
& \vee \exists n \in \text{Node} : \\
& \quad \vee \exists t \in \text{DOMAIN } \text{Target} : \\
& \quad \quad \text{ChangeMaster}(n, t) \\
& \quad \vee \exists t \in \text{DOMAIN } \text{transaction} : \\
& \quad \quad \text{ReconcileTransaction}(n, t)
\end{aligned}$$

$$\forall \exists t \in \text{DOMAIN } \text{configuration} : \\ \text{ReconcileConfiguration}(n, t)$$

$$\text{Spec} \triangleq \text{Init} \wedge \Box[\text{Next}]_{\text{vars}}$$

$$\text{Inv} \triangleq \\ \wedge \forall a, b \in \text{DOMAIN } \text{transaction} : \\ \text{transaction}[a].\text{index} > \text{transaction}[b].\text{index} \Rightarrow \\ (\text{transaction}[a].\text{status} \in \{ \text{TransactionComplete}, \text{TransactionFailed} \} \Rightarrow \\ \text{transaction}[b].\text{status} \in \{ \text{TransactionComplete}, \text{TransactionFailed} \}) \\ \wedge \forall t \in \text{DOMAIN } \text{Target} : \\ \forall c \in \text{DOMAIN } \text{history}[t] : \\ \wedge \text{configuration}[t].\text{txIndex} \geq \text{history}[t][c].\text{txIndex} \\ \wedge \text{configuration}[t].\text{syncIndex} \geq \text{history}[t][c].\text{syncIndex}$$

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