
MODULE *Config*

INSTANCE *Naturals*

INSTANCE *FiniteSets*

INSTANCE *Sequences*

INSTANCE *TLC*

An empty constant

CONSTANT *Nil*

Transaction type constants

CONSTANTS

TransactionChange,
TransactionRollback

Transaction status constants

CONSTANTS

TransactionPending,
TransactionValidating,
TransactionApplying,
TransactionComplete,
TransactionFailed

Configuration status constants

CONSTANTS

ConfigurationPending,
ConfigurationInitializing,
ConfigurationUpdating,
ConfigurationComplete,
ConfigurationFailed

The set of all nodes

CONSTANT *Node*

Target is the possible targets, paths, and values

Example: $Target \triangleq$ [
 $target1 \mapsto$ [
 $path1 \mapsto \{ "value1", "value2" \}$,
 $path2 \mapsto \{ "value2", "value3" \}$],
 $target2 \mapsto$ [
 $path2 \mapsto \{ "value3", "value4" \}$,
 $path3 \mapsto \{ "value4", "value5" \}$]]

CONSTANT *Target*

```

ASSUME Nil ∈ STRING

ASSUME TransactionPending ∈ STRING
ASSUME TransactionValidating ∈ STRING
ASSUME TransactionApplying ∈ STRING
ASSUME TransactionComplete ∈ STRING
ASSUME TransactionFailed ∈ STRING

ASSUME ConfigurationPending ∈ STRING
ASSUME ConfigurationInitializing ∈ STRING
ASSUME ConfigurationUpdating ∈ STRING
ASSUME ConfigurationComplete ∈ STRING
ASSUME ConfigurationFailed ∈ STRING

ASSUME ∧ IsFiniteSet(Node)
      ∧ ∀ n ∈ Node :
        ∧ n ∉ DOMAIN Target
        ∧ n ∈ STRING

ASSUME ∧ ∀ t ∈ DOMAIN Target :
      ∧ t ∉ Node
      ∧ t ∈ STRING
      ∧ ∀ p ∈ DOMAIN Target[t] :
        IsFiniteSet(Target[t][p])

```

```

TYPE TransactionType ::= type ∈
  { TransactionChange,
    TransactionRollback }

TYPE TransactionStatus ::= status ∈
  { TransactionPending,
    TransactionValidating,
    TransactionApplying,
    TransactionComplete,
    TransactionFailed }

TYPE Transaction  $\triangleq$  [
  type ::= type ∈ TransactionType,
  index ::= index ∈ Nat,
  revision ::= revision ∈ Nat,
  atomic ::= atomic ∈ BOOLEAN ,
  sync ::= sync ∈ BOOLEAN ,
  changes ::= [ target ∈ SUBSET (DOMAIN Target) ↦ [
    path ∈ SUBSET (DOMAIN Target[target]) ↦ [
      value ::= value ∈ STRING,
      delete ::= delete ∈ BOOLEAN ]],
  rollback ::= index ∈ Nat,
  status ::= status ∈ TransactionStatus]

```

TYPE *ConfigurationStatus* ::= *status* ∈

{ *ConfigurationPending*,
ConfigurationInitializing,
ConfigurationUpdating,
ConfigurationComplete,
ConfigurationFailed }

TYPE *Configuration* \triangleq [

id ::= *id* ∈ STRING,
revision ::= *revision* ∈ Nat,
target ::= *target* ∈ STRING,
paths ::= [*path* ∈ SUBSET (DOMAIN *Target*[*target*]) \mapsto [
value ::= *value* ∈ STRING,
index ::= *index* ∈ Nat,
deleted ::= *delete* ∈ BOOLEAN]],
txIndex ::= *txIndex* ∈ Nat,
syncIndex ::= *syncIndex* ∈ Nat,
term ::= *term* ∈ Nat,
status ::= *status* ∈ *ConfigurationStatus*]

A sequence of transactions

Each transactions contains a record of 'changes' for a set of targets

VARIABLE *transaction*

A record of target configurations

Each configuration represents the desired state of the target

VARIABLE *configuration*

A record of target states

VARIABLE *target*

A record of target masters

VARIABLE *master*

VARIABLE *history*

vars \triangleq $\langle transaction, configuration, master, target, history \rangle$

SetMaster(*n*, *t*) \triangleq

$\wedge master[t].master \neq n$

$\wedge master' = [master \text{ EXCEPT } ![t].term = master[t].term + 1,$
 $![t].master = n]$

$\wedge \text{UNCHANGED } \langle transaction, configuration, target, history \rangle$

UnsetMaster(*t*) \triangleq

$\wedge master[t].master \neq Nil$

$\wedge master' = [master \text{ EXCEPT } ![t].master = Nil]$

$\wedge \text{UNCHANGED } \langle transaction, configuration, target, history \rangle$

This section models the northbound *API* for the configuration service.

$$\text{ValidValues}(t, p) \triangleq \text{UNION } \{ \{ [value \mapsto v, delete \mapsto \text{FALSE}] : v \in \text{Target}[t][p] \}, \{ [value \mapsto \text{Nil}, delete \mapsto \text{TRUE}] \} \}$$

$$\text{ValidPaths}(t) \triangleq \text{UNION } \{ \{ v @@@ [path \mapsto p] : v \in \text{ValidValues}(t, p) \} : p \in \text{DOMAIN } \text{Target}[t] \}$$

$$\text{ValidTargets} \triangleq \text{UNION } \{ \{ p @@@ [target \mapsto t] : p \in \text{ValidPaths}(t) \} : t \in \text{DOMAIN } \text{Target} \}$$

$$\begin{aligned} \text{ValidPath}(s, t, p) &\triangleq \\ \text{LET } value &\triangleq \text{CHOOSE } v \in s : v.target = t \wedge v.path = p \\ \text{IN} & \\ &[value \mapsto value.value, \\ &delete \mapsto value.delete] \end{aligned}$$

$$\text{ValidTarget}(s, t) \triangleq [p \in \{ v.path : v \in \{ v \in s : v.target = t \} \} \mapsto \text{ValidPath}(s, t, p)]$$

$$\text{ValidChange}(s) \triangleq [t \in \{ v.target : v \in s \} \mapsto \text{ValidTarget}(s, t)]$$

$$\begin{aligned} \text{ValidChanges} &\triangleq \\ \text{LET } changeSets &\triangleq \{ s \in \text{SUBSET } \text{ValidTargets} : \\ &\quad \forall t \in \text{DOMAIN } \text{Target} : \\ &\quad \quad \forall p \in \text{DOMAIN } \text{Target}[t] : \\ &\quad \quad \quad \text{Cardinality}(\{ v \in s : v.target = t \wedge v.path = p \}) \leq 1 \} \\ \text{IN} & \\ &\{ \text{ValidChange}(s) : s \in changeSets \} \end{aligned}$$

$$\begin{aligned} \text{NextIndex} &\triangleq \\ \text{IF } \text{DOMAIN } transaction &= \{ \} \text{ THEN} \\ &1 \\ \text{ELSE} & \\ \text{LET } i &\triangleq \text{CHOOSE } i \in \text{DOMAIN } transaction : \\ &\quad \forall j \in \text{DOMAIN } transaction : \\ &\quad \quad i \geq j \\ \text{IN} & \quad i + 1 \end{aligned}$$

Add a set of changes to the transaction log

$$\begin{aligned} \text{Change}(c) &\triangleq \\ \wedge \quad transaction' &= transaction @@@ (\text{NextIndex} > [type \mapsto \text{TransactionChange}, \\ &\quad index \mapsto \text{NextIndex}, \\ &\quad atomic \mapsto \text{FALSE}, \\ &\quad sync \mapsto \text{FALSE}, \\ &\quad changes \mapsto c, \end{aligned}$$

$$\begin{aligned} & \text{sources} \mapsto \langle \rangle, \\ & \text{status} \mapsto \text{TransactionPending}) \\ \wedge \text{ UNCHANGED } \langle \text{configuration}, \text{master}, \text{target}, \text{history} \rangle \end{aligned}$$

Add a rollback to the transaction log

$$\begin{aligned} \text{Rollback}(t) &\triangleq \\ \wedge \text{ transaction}[t].\text{type} &= \text{TransactionChange} \\ \wedge \text{ transaction}' = \text{transaction} @ @ (\text{NextIndex} : > [& \text{type} \mapsto \text{TransactionRollback}, \\ & \text{index} \mapsto \text{NextIndex}, \\ & \text{atomic} \mapsto \text{FALSE}, \\ & \text{sync} \mapsto \text{FALSE}, \\ & \text{rollback} \mapsto t, \\ & \text{status} \mapsto \text{TransactionPending}]) \\ \wedge \text{ UNCHANGED } \langle \text{configuration}, \text{master}, \text{target}, \text{history} \rangle \end{aligned}$$

This section models the Transaction log reconciler.

$$\begin{aligned} \text{ReconcileChange}(n, i) &\triangleq \\ \text{If the transaction is Pending, begin validation if the prior transaction} & \\ \text{has already been applied. This simplifies concurrency control in the controller} & \\ \text{and guarantees transactions are applied to the configurations in sequential order.} & \\ \vee \wedge \text{ transaction}[i].\text{status} = \text{TransactionPending} & \\ \wedge \vee \wedge i - 1 \in \text{DOMAIN transaction} & \\ \wedge \text{ transaction}[i - 1].\text{status} \in \{ \text{TransactionComplete}, \text{TransactionFailed} \} & \\ \vee i - 1 \notin \text{DOMAIN transaction} & \\ \wedge \text{ transaction}' = [\text{transaction EXCEPT } ![i].\text{status} = \text{TransactionValidating}] & \\ \wedge \text{ UNCHANGED } \langle \text{configuration}, \text{history} \rangle & \\ \text{If the transaction is in the Validating state, compute and validate the} & \\ \text{Configuration for each target.} & \\ \vee \wedge \text{ transaction}[i].\text{status} = \text{TransactionValidating} & \\ \text{If validation fails any target, mark the transaction Failed.} & \\ \text{If validation is successful, proceed to Applying.} & \\ \wedge \exists \text{ valid} \in \text{BOOLEAN} : & \\ \text{IF valid THEN} & \\ \wedge \text{ transaction}' = [\text{transaction EXCEPT } ![i].\text{status} = \text{TransactionApplying}] & \\ \text{ELSE} & \\ \wedge \text{ transaction}' = [\text{transaction EXCEPT } ![i].\text{status} = \text{TransactionFailed}] & \\ \wedge \text{ UNCHANGED } \langle \text{configuration}, \text{history} \rangle & \\ \text{If the transaction is in the Applying state, update the Configuration for each} & \\ \text{target and Complete the transaction.} & \\ \vee \wedge \text{ transaction}[i].\text{status} = \text{TransactionApplying} & \\ \text{Update the target configurations, adding the transaction index to each updated path} & \\ \wedge \text{ configuration}' = [& \\ t \in \text{DOMAIN Target} \mapsto & \\ \text{IF } t \in \text{DOMAIN transaction}[i].\text{changes THEN} & \end{aligned}$$

```

[configuration[t] EXCEPT
  !.paths = [p ∈ DOMAIN transaction[i].changes[t] ↦
    [index ↦ transaction[i].index,
     value ↦ transaction[i].changes[t][p].value,
     deleted ↦ transaction[i].changes[t][p].delete]]
  @@ configuration[t].paths,
  !.txIndex = transaction[i].index,
  !.status = ConfigurationPending]
ELSE
  configuration[t]]
∧ history' = [r ∈ DOMAIN Target ↦ Append(history[r], configuration'[r])]
∧ transaction' = [transaction EXCEPT
  ![i].status = TransactionComplete,
  ![i].sources = [t ∈ DOMAIN transaction[i].changes ↦
    LET updatePaths ≜ {p ∈ DOMAIN transaction[i].changes[t] :
      ¬transaction[i].changes[t][p].delete}
  IN [p ∈ updatePaths ∩ DOMAIN configuration[t].paths ↦ configuration[t].paths[p]]]]]

```

ReconcileRollback(n, i) \triangleq

If the transaction is Pending, begin validation if the prior transaction has already been applied. This simplifies concurrency control in the controller and guarantees transactions are applied to the configurations in sequential order.

```

∨ ∧ transaction[i].status = TransactionPending
  ∧ ∨ ∧ i - 1 ∈ DOMAIN transaction
    ∧ transaction[i - 1].status ∈ {TransactionComplete, TransactionFailed}
    ∨ i - 1 ∉ DOMAIN transaction
  ∧ transaction' = [transaction EXCEPT ![i].status = TransactionValidating]
  ∧ UNCHANGED ⟨configuration, history⟩

```

If the transaction is in the Validating state, compute and validate the Configuration for each target.

```

∨ ∧ transaction[i].status = TransactionValidating
  ∧ ∨ ∧ transaction[transaction[i].rollback].status = TransactionComplete

```

A transaction can be rolled back if:

1. The source transaction is in the log
2. The source transaction is complete
3. For all changes in the transaction, no more recent change has been made

```

∧ ∨ ∧ transaction[i].rollback ∈ DOMAIN transaction
  ∧ LET canRollback ≜ ∀ t ∈ DOMAIN transaction[transaction[i].rollback].changes :
    ∀ p ∈ DOMAIN transaction[transaction[i].rollback].changes[t] :
      configuration[t].paths[p].index = transaction[i].rollback

```

IN

IF *canRollback* THEN

```

  ∧ transaction' = [transaction EXCEPT ![i].status = TransactionApplying]

```

ELSE

```

  ∧ transaction' = [transaction EXCEPT ![i].status = TransactionFailed]

```

$$\begin{aligned}
& \vee \wedge \text{transaction}[i].\text{rollback} \notin \text{DOMAIN } \text{transaction} \\
& \quad \wedge \text{transaction}' = [\text{transaction EXCEPT } ![i].\text{status} = \text{TransactionFailed}] \\
& \vee \wedge \text{transaction}[\text{transaction}[i].\text{rollback}].\text{status} = \text{TransactionFailed} \\
& \quad \wedge \text{transaction}' = [\text{transaction EXCEPT } ![i].\text{status} = \text{TransactionFailed}] \\
& \wedge \text{UNCHANGED } \langle \text{configuration}, \text{history} \rangle \\
& \text{If the transaction is in the Applying state, update the Configuration for each} \\
& \text{target and Complete the transaction.} \\
& \vee \wedge \text{transaction}[i].\text{status} = \text{TransactionApplying} \\
& \quad \text{Revert the target configurations} \\
& \quad \wedge \text{configuration}' = [\\
& \quad \quad t \in \text{DOMAIN } \text{Target} \mapsto \\
& \quad \quad \text{IF } t \in \text{DOMAIN } \text{transaction}[\text{transaction}[i].\text{rollback}].\text{changes} \text{ THEN} \\
& \quad \quad \quad \text{LET } \text{adds} \triangleq \{p \in \text{DOMAIN } \text{transaction}[\text{transaction}[i].\text{rollback}].\text{changes}[t] : \\
& \quad \quad \quad \quad \wedge p \notin \text{DOMAIN } \text{transaction}[\text{transaction}[i].\text{rollback}].\text{sources}[t] : \\
& \quad \quad \quad \quad \wedge \neg \text{transaction}[\text{transaction}[i].\text{rollback}].\text{changes}[t][p].\text{delete}\} \\
& \quad \quad \quad \text{updates} \triangleq \{p \in \text{DOMAIN } \text{transaction}[\text{transaction}[i].\text{rollback}].\text{changes}[t] : \\
& \quad \quad \quad \quad \wedge p \in \text{DOMAIN } \text{transaction}[\text{transaction}[i].\text{rollback}].\text{sources}[t] : \\
& \quad \quad \quad \quad \wedge \neg \text{transaction}[\text{transaction}[i].\text{rollback}].\text{changes}[t][p].\text{delete}\} \\
& \quad \quad \quad \text{removes} \triangleq \{p \in \text{DOMAIN } \text{transaction}[\text{transaction}[i].\text{rollback}].\text{changes}[t] : \\
& \quad \quad \quad \quad \wedge p \in \text{DOMAIN } \text{transaction}[\text{transaction}[i].\text{rollback}].\text{sources}[t] : \\
& \quad \quad \quad \quad \wedge \text{transaction}[\text{transaction}[i].\text{rollback}].\text{changes}[t][p].\text{delete}\} \\
& \quad \quad \quad \text{changes} \triangleq \text{adds} \cup \text{updates} \cup \text{removes} \\
& \quad \quad \quad \text{unchanges} \triangleq \text{DOMAIN } \text{configuration}[t].\text{paths} \setminus \text{changes} \\
& \quad \quad \text{IN} \\
& \quad \quad \quad [\text{configuration}[t] \text{ EXCEPT} \\
& \quad \quad \quad \quad !.\text{paths} = [p \in \text{unchanges} \mapsto \text{configuration}[t].\text{paths}[p]] \\
& \quad \quad \quad \quad @@ [p \in \text{updates} \cup \text{removes} \mapsto \\
& \quad \quad \quad \quad \quad \text{transaction}[\text{transaction}[i].\text{rollback}].\text{sources}[t][p]], \\
& \quad \quad \quad \quad !.\text{txIndex} = \text{transaction}[i].\text{index}, \\
& \quad \quad \quad \quad !.\text{status} = \text{ConfigurationPending}] \\
& \quad \quad \text{ELSE} \\
& \quad \quad \quad \text{configuration}[t] \\
& \quad \wedge \text{history}' = [r \in \text{DOMAIN } \text{Target} \mapsto \text{Append}(\text{history}[r], \text{configuration}'[r])] \\
& \quad \wedge \text{transaction}' = [\text{transaction EXCEPT } ![i].\text{status} = \text{TransactionComplete}] \\
& \quad \text{Reconcile the transaction log} \\
& \text{ReconcileTransaction}(n, i) \triangleq \\
& \quad \wedge \vee \wedge \text{transaction}[i].\text{type} = \text{TransactionChange} \\
& \quad \quad \wedge \text{ReconcileChange}(n, i) \\
& \quad \vee \wedge \text{transaction}[i].\text{type} = \text{TransactionRollback} \\
& \quad \quad \wedge \text{ReconcileRollback}(n, i) \\
& \quad \wedge \text{UNCHANGED } \langle \text{master}, \text{target} \rangle
\end{aligned}$$

This section models the Configuration reconciler.

$ReconcileConfiguration(n, c) \triangleq$
 $\wedge \vee \wedge configuration[c].status = ConfigurationPending$
 $\wedge master[configuration[c].target].master \neq Nil$
 If the configuration is marked *ConfigurationPending* and mastership has changed (indicated by an increased mastership term), mark the configuration *ConfigurationInitializing* to force full re-synchronization.
 $\wedge \vee \wedge master[configuration[c].target].term > configuration[c].term$
 $\wedge configuration' = [configuration \text{ EXCEPT } ![c].status = ConfigurationInitializing,$
 $![c].term = master[configuration[c].target].term]$
 $\wedge history' = [history \text{ EXCEPT } ![c] = Append(history[c], configuration'[c])]$
 If the configuration is marked *ConfigurationPending* and the values have changed (determined by comparing the transaction index to the last *sync* index), mark the configuration *ConfigurationUpdating* to push the changes to the target.
 $\vee \wedge master[configuration[c].target].term = configuration[c].term$
 $\wedge configuration[c].syncIndex < configuration[c].txIndex$
 $\wedge configuration' = [configuration \text{ EXCEPT } ![c].status = ConfigurationUpdating]$
 $\wedge history' = [history \text{ EXCEPT } ![c] = Append(history[c], configuration'[c])]$
 $\wedge UNCHANGED \langle target \rangle$
 $\vee \wedge configuration[c].status = ConfigurationInitializing$
 $\wedge master[configuration[c].target].master = n$
 Merge the configuration paths with the target paths, removing paths that have been marked deleted
 $\wedge LET \ deletePaths \triangleq \{p \in DOMAIN \ configuration[c].paths : configuration[c].paths[p].deleted\}$
 $\quad \quad \quad configPaths \triangleq DOMAIN \ configuration[c].paths \setminus deletePaths$
 $\quad \quad \quad targetPaths \triangleq DOMAIN \ target[configuration[c].target] \setminus deletePaths$
 IN
 $\wedge target' = [target \text{ EXCEPT } ![configuration[c].target] =$
 $\quad [p \in configPaths \mapsto [value \mapsto configuration[c].paths[p]]]$
 $\quad @@ [p \in targetPaths \mapsto target[configuration[c].target][p]]]$
 Set the configuration's status to Complete
 $\wedge configuration' = [configuration \text{ EXCEPT } ![c].status = ConfigurationComplete,$
 $\quad \quad \quad ![c].syncIndex = configuration[c].txIndex]$
 $\wedge history' = [history \text{ EXCEPT } ![c] = Append(history[c], configuration'[c])]$
 If the configuration is marked *ConfigurationUpdating*, we only need to push paths that have changed since the target was initialized or last updated by the controller. The set of changes made since the last synchronization are identified by comparing the index of each path-value to the last synchronization index, *syncIndex*
 $\vee \wedge configuration[c].status = ConfigurationUpdating$
 $\wedge master[configuration[c].target].master = n$
 Compute the set of updated and deleted paths by comparing their indexes to the target's last sync index.
 $\wedge LET \ updatePaths \triangleq \{p \in DOMAIN \ configuration[c].paths :$
 $\quad \quad \quad configuration[c].paths[p].index > configuration[c].syncIndex\}$

$$\begin{aligned}
deletePaths &\triangleq \{p \in updatePaths : configuration[c].paths[p].deleted\} \\
configPaths &\triangleq updatePaths \setminus deletePaths \\
targetPaths &\triangleq \text{DOMAIN } target[configuration[c].target] \setminus deletePaths
\end{aligned}$$

IN

Update the target paths by adding/updating paths that have changed and removing paths that have been deleted since the last *sync*.

$$\begin{aligned}
&\wedge target' = [target \text{ EXCEPT } ![configuration[c].target] = \\
&\quad [p \in configPaths \mapsto configuration[c].paths[p]] \\
&\quad @@ [p \in targetPaths \mapsto target[configuration[c].target][p]] \\
&\wedge configuration' = [configuration \text{ EXCEPT } ![c].status = ConfigurationComplete, \\
&\quad ![c].syncIndex = configuration[c].txIndex] \\
&\wedge history' = [history \text{ EXCEPT } ![c] = Append(history[c], configuration'[c])]
\end{aligned}$$

If the configuration is not already *ConfigurationPending* and mastership has been lost revert it. This can occur when the connection to the target has been lost and the mastership is no longer valid.

TODO: We still need to model mastership changes

$$\begin{aligned}
&\vee \wedge configuration[c].status \neq ConfigurationPending \\
&\quad \wedge master[configuration[c].target].master = Nil \\
&\quad \wedge configuration' = [configuration \text{ EXCEPT } ![c].status = ConfigurationPending] \\
&\quad \wedge history' = [history \text{ EXCEPT } ![c] = Append(history[c], configuration'[c])] \\
&\quad \wedge \text{UNCHANGED } \langle target \rangle \\
&\quad \wedge \text{UNCHANGED } \langle transaction, master \rangle
\end{aligned}$$

Init and next state predicates

$$\begin{aligned}
Init &\triangleq \\
&\wedge transaction = \langle \rangle \\
&\wedge configuration = [t \in \text{DOMAIN } Target \mapsto \\
&\quad [target \mapsto t, \\
&\quad paths \mapsto \\
&\quad \quad [path \in \{\} \mapsto \\
&\quad \quad \quad [path \mapsto path, \\
&\quad \quad \quad value \mapsto Nil, \\
&\quad \quad \quad index \mapsto 0, \\
&\quad \quad \quad deleted \mapsto FALSE]], \\
&\quad txIndex \mapsto 0, \\
&\quad syncIndex \mapsto 0, \\
&\quad term \mapsto 0, \\
&\quad status \mapsto ConfigurationPending]] \\
&\wedge target = [t \in \text{DOMAIN } Target \mapsto \\
&\quad [path \in \{\} \mapsto \\
&\quad \quad [value \mapsto Nil]]] \\
&\wedge master = [t \in \text{DOMAIN } Target \mapsto [master \mapsto Nil, term \mapsto 0]] \\
&\wedge history = [t \in \text{DOMAIN } Target \mapsto \langle \rangle]
\end{aligned}$$

$$\begin{aligned}
Next &\triangleq \\
&\vee \exists c \in ValidChanges : \\
&\quad Change(c) \\
&\vee \exists t \in DOMAIN \ transaction : \\
&\quad Rollback(t) \\
&\vee \exists n \in Node : \\
&\quad \exists t \in DOMAIN \ Target : \\
&\quad \quad SetMaster(n, t) \\
&\vee \exists t \in DOMAIN \ Target : \\
&\quad \quad UnsetMaster(t) \\
&\vee \exists n \in Node : \\
&\quad \exists t \in DOMAIN \ transaction : \\
&\quad \quad ReconcileTransaction(n, t) \\
&\vee \exists n \in Node : \\
&\quad \exists c \in DOMAIN \ configuration : \\
&\quad \quad ReconcileConfiguration(n, c) \\
Spec &\triangleq Init \wedge \Box [Next]_{vars} \\
Inv &\triangleq \\
&\wedge \forall a, b \in DOMAIN \ transaction : \\
&\quad transaction[a].index > transaction[b].index \Rightarrow \\
&\quad (transaction[a].status \in \{TransactionComplete, TransactionFailed\} \Rightarrow \\
&\quad \quad transaction[b].status \in \{TransactionComplete, TransactionFailed\}) \\
&\wedge \forall t \in DOMAIN \ Target : \\
&\quad \forall c \in DOMAIN \ history[t] : \\
&\quad \quad \wedge configuration[t].txIndex \geq history[t][c].txIndex \\
&\quad \quad \wedge configuration[t].syncIndex \geq history[t][c].syncIndex
\end{aligned}$$

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