
MODULE *Config*

INSTANCE *Naturals*

INSTANCE *FiniteSets*

INSTANCE *Sequences*

INSTANCE *TLC*

GenerateTestCases \triangleq FALSE

Nil \triangleq "<nil>"

Change \triangleq "Change"

Rollback \triangleq "Rollback"

Commit \triangleq "Commit"

Apply \triangleq "Apply"

Pending \triangleq "Pending"

Complete \triangleq "Complete"

Aborted \triangleq "Aborted"

Failed \triangleq "Failed"

Node \triangleq {"node1"}

NumTransactions \triangleq 3

NumTerms \triangleq 2

NumConns \triangleq 2

NumStarts \triangleq 2

Path \triangleq {"path1"}

Value \triangleq {"value1", "value2"}

A transaction log of changes and rollbacks.

VARIABLE *transaction*

A record of per-target configurations

VARIABLE *configuration*

A record of target masterhips

VARIABLE *mastership*

A record of node connections to the target

VARIABLE *conn*

The target state
 VARIABLE *target*

A sequence of state changes used for model checking.
 VARIABLE *history*

$\text{vars} \triangleq \langle \text{transaction}, \text{configuration}, \text{mastership}, \text{conn}, \text{target}, \text{history} \rangle$

LOCAL *Transaction* \triangleq INSTANCE *Transaction*

LOCAL *Configuration* \triangleq INSTANCE *Configuration*

LOCAL *Mastership* \triangleq INSTANCE *Mastership*

LOCAL *Target* \triangleq INSTANCE *Target*

$\text{AppendChange}(i) \triangleq$
 $\wedge \text{Transaction!AppendChange}(i)$

$\text{RollbackChange}(i) \triangleq$
 $\wedge \text{Transaction!RollbackChange}(i)$

$\text{ReconcileTransaction}(n, i) \triangleq$
 $\wedge \text{Transaction!ReconcileTransaction}(n, i)$
 $\wedge \text{GenerateTestCases} \Rightarrow \text{Transaction!Test!Log}([node \mapsto n, index \mapsto i])$

$\text{ReconcileConfiguration}(n) \triangleq$
 $\wedge \text{Configuration!ReconcileConfiguration}(n)$
 $\wedge \text{UNCHANGED } \langle \text{transaction}, \text{history} \rangle$
 $\wedge \text{GenerateTestCases} \Rightarrow \text{Configuration!Test!Log}([node \mapsto n])$

$\text{ReconcileMastership}(n) \triangleq$
 $\wedge \text{Mastership!ReconcileMastership}(n)$
 $\wedge \text{UNCHANGED } \langle \text{transaction}, \text{configuration}, \text{target}, \text{history} \rangle$
 $\wedge \text{GenerateTestCases} \Rightarrow \text{Mastership!Test!Log}([node \mapsto n])$

$\text{ConnectNode}(n) \triangleq$
 $\wedge \text{Target!Connect}(n)$
 $\wedge \text{UNCHANGED } \langle \text{transaction}, \text{configuration}, \text{mastership}, \text{history} \rangle$

$\text{DisconnectNode}(n) \triangleq$
 $\wedge \text{Target!Disconnect}(n)$
 $\wedge \text{UNCHANGED } \langle \text{transaction}, \text{configuration}, \text{mastership}, \text{history} \rangle$

$\text{StartTarget} \triangleq$
 $\wedge \text{Target!Start}$

$$\begin{aligned}
& \wedge \text{UNCHANGED } \langle \text{transaction}, \text{configuration}, \text{mastership}, \text{history} \rangle \\
\text{StopTarget} & \triangleq \\
& \wedge \text{Target!Stop} \\
& \wedge \text{UNCHANGED } \langle \text{transaction}, \text{configuration}, \text{mastership}, \text{history} \rangle
\end{aligned}$$

Formal specification, constraints, and theorems.

$$\begin{aligned}
\text{Init} & \triangleq \\
& \wedge \text{transaction} = [\\
& \quad i \in \{\} \mapsto [\\
& \quad \quad \text{type} \mapsto \text{Nil}, \\
& \quad \quad \text{index} \mapsto 0, \\
& \quad \quad \text{revision} \mapsto 0, \\
& \quad \quad \text{commit} \mapsto \text{Nil}, \\
& \quad \quad \text{apply} \mapsto \text{Nil}, \\
& \quad \quad \text{change} \mapsto [\\
& \quad \quad \quad \text{index} \mapsto 0, \\
& \quad \quad \quad \text{revision} \mapsto 0, \\
& \quad \quad \quad \text{values} \mapsto [\\
& \quad \quad \quad \quad p \in \{\} \mapsto [\\
& \quad \quad \quad \quad \quad \text{index} \mapsto 0, \\
& \quad \quad \quad \quad \quad \text{value} \mapsto \text{Nil}]]], \\
& \quad \quad \text{rollback} \mapsto [\\
& \quad \quad \quad \text{index} \mapsto 0, \\
& \quad \quad \quad \text{revision} \mapsto 0, \\
& \quad \quad \quad \text{values} \mapsto [\\
& \quad \quad \quad \quad p \in \{\} \mapsto [\\
& \quad \quad \quad \quad \quad \text{index} \mapsto 0, \\
& \quad \quad \quad \quad \quad \text{value} \mapsto \text{Nil}]]]]], \\
& \wedge \text{configuration} = [\\
& \quad \text{state} \mapsto \text{Pending}, \\
& \quad \text{term} \mapsto 0, \\
& \quad \text{committed} \mapsto [\\
& \quad \quad \text{index} \mapsto 0, \\
& \quad \quad \text{revision} \mapsto 0, \\
& \quad \quad \text{values} \mapsto [\\
& \quad \quad \quad p \in \{\} \mapsto [\\
& \quad \quad \quad \quad \text{index} \mapsto 0, \\
& \quad \quad \quad \quad \text{value} \mapsto \text{Nil}]]], \\
& \quad \text{applied} \mapsto [\\
& \quad \quad \text{target} \mapsto 0, \\
& \quad \quad \text{index} \mapsto 0, \\
& \quad \quad \text{revision} \mapsto 0, \\
& \quad \quad \text{values} \mapsto [
\end{aligned}$$

$$\begin{aligned}
& p \in \{\} \mapsto [\\
& \quad index \mapsto 0, \\
& \quad value \mapsto Nil]]]] \\
\wedge target = [\\
& \quad id \mapsto 1, \\
& \quad running \mapsto \text{TRUE}, \\
& \quad values \mapsto [\\
& \quad \quad p \in \{\} \mapsto [\\
& \quad \quad \quad index \mapsto 0, \\
& \quad \quad \quad value \mapsto Nil]]] \\
\wedge mastership = [\\
& \quad master \mapsto \text{CHOOSE } n \in Node : \text{TRUE}, \\
& \quad term \mapsto 1, \\
& \quad conn \mapsto 1] \\
\wedge conn = [\\
& \quad n \in Node \mapsto [\\
& \quad \quad id \mapsto 1, \\
& \quad \quad connected \mapsto \text{TRUE}] \\
\wedge history = \langle \rangle \\
Next \triangleq \\
& \vee \exists i \in 1 \dots NumTransactions : \\
& \quad \vee AppendChange(i) \\
& \quad \vee RollbackChange(i) \\
& \vee \exists n \in Node : \\
& \quad \exists i \in \text{DOMAIN } transaction : \\
& \quad \quad ReconcileTransaction(n, i) \\
& \vee \exists n \in Node : \\
& \quad ReconcileConfiguration(n) \\
& \vee \exists n \in Node : \\
& \quad ReconcileMastership(n) \\
& \vee \exists n \in Node : \\
& \quad \vee ConnectNode(n) \\
& \quad \vee DisconnectNode(n) \\
& \vee StartTarget \\
& \vee StopTarget \\
Spec \triangleq \\
& \wedge Init \\
& \wedge \Box [Next]_{vars} \\
& \wedge \forall i \in 1 \dots NumTransactions : \\
& \quad \text{WF}_{\langle transaction \rangle} (Transaction! RollbackChange(i)) \\
& \wedge \forall n \in Node : \\
& \quad \text{WF}_{vars} (\exists i \in \text{DOMAIN } transaction : Transaction! ReconcileTransaction(n, i)) \\
& \wedge \forall n \in Node :
\end{aligned}$$

$$\begin{aligned}
& \text{WF}_{\langle \text{configuration}, \text{mastership}, \text{conn}, \text{target} \rangle} (\text{Configuration! ReconcileConfiguration}(n)) \\
& \wedge \forall n \in \text{Node} : \\
& \quad \text{WF}_{\langle \text{mastership}, \text{conn} \rangle} (\text{Mastership! ReconcileMastership}(n)) \\
& \wedge \forall n \in \text{Node} : \\
& \quad \text{WF}_{\langle \text{conn}, \text{target} \rangle} (\text{Target! Connect}(n) \vee \text{Target! Disconnect}(n)) \\
& \wedge \text{WF}_{\langle \text{conn}, \text{target} \rangle} (\text{Target! Start} \vee \text{Target! Stop})
\end{aligned}$$

$$\begin{aligned}
\text{LimitTerms} & \triangleq \\
& \vee \text{mastership.term} < \text{NumTerms} \\
& \vee \wedge \text{mastership.term} = \text{NumTerms} \\
& \quad \wedge \text{mastership.master} \neq \text{Nil}
\end{aligned}$$

$$\begin{aligned}
\text{LimitConns} & \triangleq \\
& \forall n \in \text{DOMAIN conn} : \\
& \quad \vee \text{conn}[n].\text{id} < \text{NumConns} \\
& \quad \vee \wedge \text{conn}[n].\text{id} = \text{NumConns} \\
& \quad \quad \wedge \text{conn}[n].\text{connected}
\end{aligned}$$

$$\begin{aligned}
\text{LimitStarts} & \triangleq \\
& \vee \text{target.id} < 2 \\
& \vee \wedge \text{target.id} = 2 \\
& \quad \wedge \text{target.running}
\end{aligned}$$

$$\begin{aligned}
\text{TypeOK} & \triangleq \\
& \wedge \text{Transaction! TypeOK} \\
& \wedge \text{Configuration! TypeOK} \\
& \wedge \text{Mastership! TypeOK}
\end{aligned}$$

$$\begin{aligned}
\text{LOCAL } \text{IsOrderedChange}(p, i) & \triangleq \\
& \wedge \text{history}[i].\text{type} = \text{Change} \\
& \wedge \text{history}[i].\text{phase} = p \\
& \wedge \neg \exists j \in \text{DOMAIN history} : \\
& \quad \wedge j < i \\
& \quad \wedge \text{history}[j].\text{type} = \text{Change} \\
& \quad \wedge \text{history}[j].\text{phase} = p \\
& \quad \wedge \text{history}[j].\text{revision} \geq \text{history}[i].\text{revision}
\end{aligned}$$

$$\begin{aligned}
\text{LOCAL } \text{IsOrderedRollback}(p, i) & \triangleq \\
& \wedge \text{history}[i].\text{type} = \text{Rollback} \\
& \wedge \text{history}[i].\text{phase} = p \\
& \wedge \exists j \in \text{DOMAIN history} : \\
& \quad \wedge j < i \\
& \quad \wedge \text{history}[j].\text{type} = \text{Change}
\end{aligned}$$

$$\begin{aligned}
& \wedge \text{history}[j].\text{revision} = \text{history}[i].\text{revision} \\
\wedge \quad & \neg \exists j \in \text{DOMAIN } \text{history} : \\
& \quad \wedge j < i \\
& \quad \wedge \text{history}[j].\text{type} = \text{Change} \\
& \quad \wedge \text{history}[j].\text{phase} = p \\
& \quad \wedge \text{history}[j].\text{revision} > \text{history}[i].\text{revision} \\
& \wedge \neg \exists k \in \text{DOMAIN } \text{history} : \\
& \quad \wedge k > j \\
& \quad \wedge k < i \\
& \quad \wedge \text{history}[k].\text{type} = \text{Rollback} \\
& \quad \wedge \text{history}[k].\text{phase} = p \\
& \quad \wedge \text{history}[k].\text{revision} = \text{history}[j].\text{revision}
\end{aligned}$$

Order \triangleq

$$\begin{aligned}
& \wedge \forall i \in \text{DOMAIN } \text{history} : \\
& \quad \vee \text{IsOrderedChange}(\text{Commit}, i) \\
& \quad \vee \text{IsOrderedChange}(\text{Apply}, i) \\
& \quad \vee \text{IsOrderedRollback}(\text{Commit}, i) \\
& \quad \vee \text{IsOrderedRollback}(\text{Apply}, i) \\
& \wedge \forall i \in \text{DOMAIN } \text{transaction} : \\
& \quad \wedge \text{transaction}[i].\text{type} = \text{Change} \\
& \quad \wedge \text{transaction}[i].\text{apply} = \text{Failed} \\
& \quad \wedge \neg \exists j \in \text{DOMAIN } \text{transaction} : \\
& \quad \quad \wedge \text{transaction}[j].\text{type} = \text{Rollback} \\
& \quad \quad \wedge \text{transaction}[j].\text{rollback}.\text{revision} = \text{transaction}[i].\text{change}.\text{revision} \\
& \quad \quad \wedge \text{transaction}[j].\text{apply} = \text{Complete} \\
& \Rightarrow \forall j \in \text{DOMAIN } \text{transaction} : (j > i \Rightarrow \\
& \quad (\text{transaction}[j].\text{type} = \text{Change} \Rightarrow \text{transaction}[j].\text{apply} \neq \text{Complete}))
\end{aligned}$$

Consistency \triangleq

$$\begin{aligned}
& \wedge \forall i \in \text{DOMAIN } \text{transaction} : \\
& \quad \wedge \text{transaction}[i].\text{commit} = \text{Complete} \\
& \quad \wedge \neg \exists j \in \text{DOMAIN } \text{transaction} : \\
& \quad \quad \wedge j > i \\
& \quad \quad \wedge \text{transaction}[j].\text{commit} = \text{Complete} \\
& \Rightarrow \forall p \in \text{DOMAIN } \text{transaction}[i].\text{change}.\text{values} : \\
& \quad \wedge \text{configuration}.\text{committed}.\text{values}[p] = \text{transaction}[i].\text{change}.\text{values}[p] \\
& \wedge \forall i \in \text{DOMAIN } \text{transaction} : \\
& \quad \wedge \text{transaction}[i].\text{apply} = \text{Complete} \\
& \quad \wedge \neg \exists j \in \text{DOMAIN } \text{transaction} : \\
& \quad \quad \wedge j > i \\
& \quad \quad \wedge \text{transaction}[j].\text{apply} = \text{Complete} \\
& \Rightarrow \forall p \in \text{DOMAIN } \text{transaction}[i].\text{change}.\text{values} : \\
& \quad \wedge \text{configuration}.\text{applied}.\text{values}[p] = \text{transaction}[i].\text{change}.\text{values}[p] \\
& \quad \wedge \text{target}.\text{running}
\end{aligned}$$

$$\begin{aligned}
& \wedge \text{configuration.applied.target} = \text{target.id} \\
& \wedge \text{configuration.state} = \text{Complete} \\
& \Rightarrow \text{target.values}[p] = \text{transaction}[i].\text{change.values}[p]
\end{aligned}$$

$$\text{Safety} \triangleq \Box(\text{Order} \wedge \text{Consistency})$$

THEOREM $\text{Spec} \Rightarrow \text{Safety}$

$$\begin{aligned}
\text{LOCAL } \text{IsChanging}(i) & \triangleq \\
& \exists j \in \text{DOMAIN } \text{transaction} : \\
& \quad \wedge \text{transaction}[j].\text{type} = \text{Change} \\
& \quad \wedge \text{transaction}[j].\text{change.revision} = i
\end{aligned}$$

$$\begin{aligned}
\text{LOCAL } \text{IsChanged}(i) & \triangleq \\
& \exists j \in \text{DOMAIN } \text{transaction} : \\
& \quad \wedge \text{transaction}[j].\text{type} = \text{Change} \\
& \quad \wedge \text{transaction}[j].\text{change.revision} = i \\
& \quad \wedge \text{transaction}[j].\text{commit} \neq \text{Pending} \\
& \quad \wedge \text{transaction}[j].\text{apply} \neq \text{Pending}
\end{aligned}$$

$$\begin{aligned}
\text{LOCAL } \text{IsRollingBack}(i) & \triangleq \\
& \exists j \in \text{DOMAIN } \text{transaction} : \\
& \quad \wedge \text{transaction}[j].\text{type} = \text{Rollback} \\
& \quad \wedge \text{transaction}[j].\text{rollback.revision} = i
\end{aligned}$$

$$\begin{aligned}
\text{LOCAL } \text{IsRolledBack}(i) & \triangleq \\
& \exists j \in \text{DOMAIN } \text{transaction} : \\
& \quad \wedge \text{transaction}[j].\text{type} = \text{Rollback} \\
& \quad \wedge \text{transaction}[j].\text{rollback.revision} = i \\
& \quad \wedge \text{transaction}[j].\text{commit} \neq \text{Pending} \\
& \quad \wedge \text{transaction}[j].\text{apply} \neq \text{Pending}
\end{aligned}$$

$$\begin{aligned}
\text{Terminates}(i) & \triangleq \\
& \wedge \text{IsChanging}(i) \leadsto \text{IsChanged}(i) \\
& \wedge \text{IsRollingBack}(i) \leadsto \text{IsRolledBack}(i)
\end{aligned}$$

$$\begin{aligned}
\text{Termination} & \triangleq \\
& \forall i \in 1 \dots \text{NumTransactions} : \text{Terminates}(i)
\end{aligned}$$

$$\text{Liveness} \triangleq \text{Termination}$$

THEOREM $\text{Spec} \Rightarrow \text{Liveness}$