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MODULE *Config*

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INSTANCE *Naturals*

INSTANCE *FiniteSets*

INSTANCE *Sequences*

INSTANCE *TLC*

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An empty constant

CONSTANT *Nil*

Transaction constants

CONSTANTS

*Pending*,  
*Validating*,  
*Applying*,  
*Complete*,  
*Failed*

The set of all nodes

CONSTANT *Node*

Target is the possible targets, paths, and values

Example:  $Target \triangleq$  [  
      $target1 \mapsto$  [  
          $path1 \mapsto \{ "value1", "value2" \}$ ,  
          $path2 \mapsto \{ "value2", "value3" \}$  ],  
      $target2 \mapsto$  [  
          $path2 \mapsto \{ "value3", "value4" \}$ ,  
          $path3 \mapsto \{ "value4", "value5" \}$  ] ]

CONSTANT *Target*

ASSUME *Nil* ∈ STRING

ASSUME *Pending* ∈ STRING

ASSUME *Validating* ∈ STRING

ASSUME *Applying* ∈ STRING

ASSUME *Complete* ∈ STRING

ASSUME *Failed* ∈ STRING

ASSUME  $\wedge IsFiniteSet(Node)$   
 $\wedge \forall n \in Node :$   
      $\wedge n \notin \text{DOMAIN } Target$   
      $\wedge n \in \text{STRING}$

ASSUME  $\wedge \forall t \in \text{DOMAIN } Target :$   
 $\wedge IsFiniteSet(Target[t])$   
 $\wedge t \notin Node$   
 $\wedge t \in \text{STRING}$

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TYPE Status ::= status  $\in$  {Pending, Validating, Applying, Complete, Failed}
TYPE Transaction  $\triangleq$  [
  id      ::= id  $\in$  STRING,
  index   ::= index  $\in$  Nat,
  revision ::= revision  $\in$  Nat,
  atomic  ::= atomic  $\in$  BOOLEAN ,
  sync    ::= sync  $\in$  BOOLEAN ,
  changes ::= [ target  $\in$  SUBSET (DOMAIN Target)  $\mapsto$  [
    path  $\in$  SUBSET (DOMAIN Target[target])  $\mapsto$  [
      value ::= value  $\in$  STRING,
      delete ::= delete  $\in$  BOOLEAN ]],
  status ::= status  $\in$  Status]
TYPE Configuration  $\triangleq$  [
  id      ::= id  $\in$  STRING,
  revision ::= revision  $\in$  Nat,
  target  ::= target  $\in$  STRING, paths ::= [
    path  $\in$  SUBSET (DOMAIN Target[target])  $\mapsto$  [
      value ::= value  $\in$  STRING,
      index ::= index  $\in$  Nat,
      deleted ::= delete  $\in$  BOOLEAN ]],
  transactionIndex ::= transactionIndex  $\in$  Nat,
  syncIndex        ::= syncIndex  $\in$  Nat,
  mastershipTerm   ::= mastershipTerm  $\in$  Nat]

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A sequence of transactions

Each transactions contains a record of 'changes' for a set of targets

VARIABLE *transactions*

A record of target configurations

Each configuration represents the desired state of the target

VARIABLE *configurations*

A record of target states

VARIABLE *targets*

A record of target masters

VARIABLE *masters*

$vars \triangleq \langle transactions, configurations, targets \rangle$

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This section models the northbound *API* for the configuration service.

This crazy thing returns the set of all possible sets of valid changes

$$\begin{aligned}
\text{ValidChanges} &\triangleq \\
&\text{LET } \text{allPaths} \triangleq \text{UNION } \{(\text{DOMAIN } \text{Target}[t]) : t \in \text{DOMAIN } \text{Target}\} \\
&\quad \text{allValues} \triangleq \text{UNION } \{\text{UNION } \{\text{Target}[t][p] : p \in \text{DOMAIN } \text{Target}[t]\} : t \in \text{DOMAIN } \text{Target}\} \\
&\text{IN} \\
&\quad \{\text{targetPathValues} \in \text{SUBSET } (\text{Target} \times \text{allPaths} \times \text{allValues} \times \text{BOOLEAN}) : \\
&\quad \quad \wedge \forall \text{target} \in \text{DOMAIN } \text{Target} : \\
&\quad \quad \quad \text{LET } \text{targetIndexes} \triangleq \{i \in 1 \dots \text{Len}(\text{targetPathValues}) : \wedge \text{targetPathValues}[i][1] = \text{target}\} \\
&\quad \quad \quad \text{IN } \quad \vee \wedge \text{Cardinality}(\text{targetIndexes}) = 0 \\
&\quad \quad \quad \quad \vee \wedge \text{Cardinality}(\text{targetIndexes}) = 1 \\
&\quad \quad \quad \quad \quad \wedge \text{LET } \text{targetPathValue} \triangleq \text{targetPathValues}[\text{CHOOSE } \text{index} \in \text{targetIndexes} : \text{TRUE}] \\
&\quad \quad \quad \quad \quad \text{IN} \\
&\quad \quad \quad \quad \quad \quad \wedge \text{targetPathValue}[2] \setminus (\text{DOMAIN } \text{Target}[\text{target}]) = \{\} \\
&\quad \quad \quad \quad \quad \quad \wedge \text{targetPathValue}[3] \in \text{Target}[\text{target}][\text{targetPathValue}[2]]\}
\end{aligned}$$

Add a set of changes to the transaction log

$$\begin{aligned}
\text{Change} &\triangleq \\
&\wedge \exists \text{changes} \in \text{ValidChanges} : \\
&\quad \wedge \text{transactions}' = \text{Append}(\text{transactions}, [\text{index} \mapsto \text{Len}(\text{transactions}) + 1, \\
&\quad \quad \quad \text{atomic} \mapsto \text{FALSE}, \\
&\quad \quad \quad \text{sync} \mapsto \text{FALSE}, \\
&\quad \quad \quad \text{changes} \mapsto \text{changes}, \\
&\quad \quad \quad \text{status} \mapsto \text{Pending}]) \\
&\wedge \text{UNCHANGED } \langle \text{configurations}, \text{targets} \rangle
\end{aligned}$$


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This section models the Transaction log reconciler.

Reconcile the transaction log

$$\begin{aligned}
\text{ReconcileTransaction}(n, tx) &\triangleq \\
&\quad \text{If the transaction is } \text{Pending}, \text{ begin validation if the prior transaction} \\
&\quad \text{has already been applied. This simplifies concurrency control in the controller} \\
&\quad \text{and guarantees transactions are applied to the configurations in sequential order.} \\
&\wedge \vee \wedge tx.\text{status} = \text{Pending} \\
&\quad \wedge \vee \wedge tx.\text{index} > 1 \\
&\quad \quad \wedge \text{transactions}[tx.\text{index} - 1].\text{status} \in \{\text{Complete}, \text{Failed}\} \\
&\quad \quad \vee tx.\text{index} = 1 \\
&\quad \wedge \text{transactions}' = [\text{transactions} \text{ EXCEPT } ![tx.\text{index}].\text{status} = \text{Validating}] \\
&\quad \wedge \text{UNCHANGED } \langle \text{configurations} \rangle \\
&\quad \text{If the transaction is in the } \text{Validating} \text{ state, compute and validate the} \\
&\quad \text{Configuration for each target.} \\
&\vee \wedge tx.\text{status} = \text{Validating} \\
&\quad \text{If validation fails any target, mark the transaction } \text{Failed}. \\
&\quad \text{If validation is successful, proceed to } \text{Applying}. \\
&\wedge \exists \text{valid} \in \text{BOOLEAN} : \\
&\quad \vee \wedge \text{valid}
\end{aligned}$$

$$\begin{aligned}
& \wedge \text{transactions}' = [\text{transactions} \text{ EXCEPT } ![tx.index].status = \text{Applying}] \\
& \vee \wedge \neg \text{valid} \\
& \wedge \text{transactions}' = [\text{transactions} \text{ EXCEPT } ![tx.index].status = \text{Failed}] \\
& \wedge \text{UNCHANGED } \langle \text{configurations} \rangle \\
& \text{If the transaction is in the } \text{Applying} \text{ state, update the Configuration for each} \\
& \text{target and } \text{Complete} \text{ the transaction.} \\
& \vee \wedge tx.status = \text{Applying} \\
& \wedge \vee \wedge tx.atomic \\
& \quad \text{TODO: Apply atomic transactions here} \\
& \wedge \text{transactions}' = [\text{transactions} \text{ EXCEPT } ![tx.index].status = \text{Complete}] \\
& \wedge \text{UNCHANGED } \langle \text{configurations} \rangle \\
& \wedge \vee \wedge \neg tx.atomic \\
& \quad \text{Add the transaction index to each updated path} \\
& \wedge \text{configurations}' = [ \\
& \quad t \in \text{DOMAIN } \text{Target} \mapsto [ \\
& \quad \quad \text{configurations}[t] \text{ EXCEPT} \\
& \quad \quad \quad !.paths = [path \in \text{DOMAIN } tx.changes \mapsto \\
& \quad \quad \quad \quad tx.changes[path] @@ [index \mapsto tx.index]] @@ \text{configurations}[t].paths, \\
& \quad \quad \quad !.transactionIndex = tx.index]] \\
& \wedge \text{transactions}' = [\text{transactions} \text{ EXCEPT } ![tx.index].status = \text{Complete}] \\
& \wedge \text{UNCHANGED } \langle \text{targets} \rangle
\end{aligned}$$


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This section models the Configuration reconciler.

$\text{ReconcileConfiguration}(n, c) \triangleq$

Only the master should reconcile the configuration

$\wedge \text{masters}[c.target].master = n$

If the configuration's mastership term is less than the current mastership term,  
assume the target may have restarted/reconnected and perform a full reconciliation  
of the target configuration from the root path.

$\wedge \vee \wedge \text{masters}[c.target].term > c.mastershipTerm$

Merge the configuration paths with the target paths, removing paths  
that have been marked deleted

$\wedge \text{targets}' = [\text{targets} \text{ EXCEPT } ![c.target] =$   
 $\quad [p \in \{p \in \text{DOMAIN } c.paths : \neg c.paths[p].deleted\} \mapsto [value \mapsto c.paths[p]]] @@$   
 $\quad [p \in \{p \in \text{DOMAIN } \text{targets}[c.target] : \neg c.paths[p].deleted\} \mapsto \text{targets}[c.target][p]]]$

Set the configuration's mastership term and *sync* index

$\wedge \text{configurations}' = [\text{configurations} \text{ EXCEPT } ![c.id].mastershipTerm = \text{masters}[c.target].term,$   
 $\quad \quad \quad ![c.id].syncIndex = c.transactionIndex]$

If the Configuration's transaction index is greater than the target index,  
reconcile the configuration with the target. Once the target has been updated,  
update the *sync* index to match the reconciled transaction index.

$\wedge \vee \wedge \text{masters}[c.target].term = c.mastershipTerm$

$\wedge c.transactionIndex > c.syncIndex$

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Compute the set of updated and deleted paths by comparing  
 their indexes to the target's last *sync* index.  
 $\wedge \text{LET } \text{updatedPaths} \triangleq \{p \in \text{DOMAIN } c.\text{paths} : c.\text{paths}[p].\text{index} > c.\text{syncIndex}\}$   
 $\text{deletedPaths} \triangleq \{p \in \text{updatedPaths} : c.\text{paths}[p].\text{deleted}\}$   
 IN  
 Update the target paths by adding/updating paths that have changed and  
 removing paths that have been deleted since the last *sync*.  
 $\wedge \text{targets}' = [\text{targets} \text{ EXCEPT } ![c.\text{target}] =$   
 $\quad [p \in \text{updatedPaths} \setminus \text{deletedPaths} \mapsto c.\text{paths}[p]] @@$   
 $\quad [p \in \text{DOMAIN } \text{targets}[c.\text{target}] \setminus \text{deletedPaths} \mapsto \text{targets}[c.\text{target}][p]]]$   
 $\wedge \text{configurations}' = [\text{configurations} \text{ EXCEPT } ![c.\text{id}].\text{syncIndex} = c.\text{transactionIndex}]$   
 $\wedge \text{UNCHANGED } \langle \text{transactions} \rangle$

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*Init* and next state predicates

$\text{Init} \triangleq$   
 $\wedge \text{transactions} = \langle \rangle$   
 $\wedge \text{configurations} = [t \in \text{Target} \mapsto [$   
 $\quad \text{id} \mapsto t,$   
 $\quad \text{config} \mapsto [\text{path} \in \{\} \mapsto [$   
 $\quad \quad \text{path} \mapsto \text{path},$   
 $\quad \quad \text{value} \mapsto \text{Nil},$   
 $\quad \quad \text{index} \mapsto 0,$   
 $\quad \quad \text{deleted} \mapsto \text{FALSE}]]]$   
 $\wedge \text{targets} = [t \in \text{Target} \mapsto$   
 $\quad [\text{path} \in \{\} \mapsto [$   
 $\quad \quad \text{value} \mapsto \text{Nil}]]]$   
 $\wedge \text{masters} = [t \in \text{Target} \mapsto [\text{master} \mapsto \text{Nil}, \text{term} \mapsto 0]]$   
 $\text{Next} \triangleq$   
 $\vee \text{Change}$   
 $\vee \exists n \in \text{Node} :$   
 $\quad \exists t \in \text{DOMAIN } \text{transactions} :$   
 $\quad \quad \text{ReconcileTransaction}(n, t)$   
 $\vee \exists n \in \text{Node} :$   
 $\quad \exists c \in \text{configurations} :$   
 $\quad \quad \text{ReconcileConfiguration}(n, c)$   
 $\text{Spec} \triangleq \text{Init} \wedge \square[\text{Next}]_{\text{vars}}$

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\ \* Modification History  
 \ \* Last modified *Thu Jan 13 15:45:19 PST 2022* by *jordanhalterman*  
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