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MODULE *Config*

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INSTANCE *Naturals*

INSTANCE *FiniteSets*

INSTANCE *Sequences*

INSTANCE *TLC*

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An empty constant

CONSTANT *Nil*

Transaction type constants

CONSTANTS

*TransactionChange*,  
*TransactionRollback*

Transaction isolation constants

CONSTANTS

*IsolationDefault*,  
*IsolationSerializable*

Transaction status constants

CONSTANTS

*TransactionInitializing*,  
*TransactionInitialized*,  
*TransactionValidating*,  
*TransactionValidated*,  
*TransactionCommitting*,  
*TransactionCommitted*,  
*TransactionApplying*,  
*TransactionApplied*,  
*TransactionFailed*

*TransactionStatus*  $\triangleq$

$\langle$ *TransactionInitializing*,  
*TransactionInitialized*,  
*TransactionValidating*,  
*TransactionValidated*,  
*TransactionCommitting*,  
*TransactionCommitted*,  
*TransactionApplying*,  
*TransactionApplied*,  
*TransactionFailed* $\rangle$

Proposal type constants

CONSTANTS

*ProposalChange*,  
*ProposalRollback*

Proposal status constants

CONSTANTS

*ProposalInitializing*,  
*ProposalInitialized*,  
*ProposalValidating*,  
*ProposalValidated*,  
*ProposalCommitting*,  
*ProposalCommitted*,  
*ProposalApplying*,  
*ProposalApplied*,  
*ProposalFailed*

$ProposalStatus \triangleq$   
 $\langle ProposalInitializing,$   
 $ProposalInitialized,$   
 $ProposalValidating,$   
 $ProposalValidated,$   
 $ProposalCommitting,$   
 $ProposalCommitted,$   
 $ProposalApplying,$   
 $ProposalApplied,$   
 $ProposalFailed \rangle$

Configuration status constants

CONSTANTS

*ConfigurationUnknown*,  
*ConfigurationSynchronizing*,  
*ConfigurationSynchronized*,  
*ConfigurationPersisted*,  
*ConfigurationFailed*

The set of all nodes

CONSTANT *Node*

Target is the set of all targets and their possible paths and values.

Example:  $Target \triangleq$  [  
   $target1 \mapsto [persistent \mapsto FALSE, paths \mapsto [$   
     $path1 \mapsto \{“value1”, “value2”\},$   
     $path2 \mapsto \{“value2”, “value3”\}]],$   
   $target2 \mapsto [persistent \mapsto TRUE, paths \mapsto [$   
     $path2 \mapsto \{“value3”, “value4”\},$   
     $path3 \mapsto \{“value4”, “value5”\}]]]$

CONSTANT *Target*  
 $Phase(S, s) \triangleq \text{CHOOSE } i \in \text{DOMAIN } S : S[i] = s$   
 $TransactionPhase(s) \triangleq Phase(TransactionStatus, s)$   
 $ProposalPhase(s) \triangleq Phase(ProposalStatus, s)$   
 ASSUME *Nil* ∈ STRING  
  
 ASSUME *TransactionInitializing* ∈ STRING  
 ASSUME *TransactionInitialized* ∈ STRING  
 ASSUME *TransactionValidating* ∈ STRING  
 ASSUME *TransactionValidated* ∈ STRING  
 ASSUME *TransactionCommitting* ∈ STRING  
 ASSUME *TransactionCommitted* ∈ STRING  
 ASSUME *TransactionApplying* ∈ STRING  
 ASSUME *TransactionApplied* ∈ STRING  
 ASSUME *TransactionFailed* ∈ STRING  
  
 ASSUME *ProposalInitializing* ∈ STRING  
 ASSUME *ProposalInitialized* ∈ STRING  
 ASSUME *ProposalValidating* ∈ STRING  
 ASSUME *ProposalValidated* ∈ STRING  
 ASSUME *ProposalCommitting* ∈ STRING  
 ASSUME *ProposalCommitted* ∈ STRING  
 ASSUME *ProposalApplying* ∈ STRING  
 ASSUME *ProposalApplied* ∈ STRING  
 ASSUME *ProposalFailed* ∈ STRING  
  
 ASSUME *ConfigurationUnknown* ∈ STRING  
 ASSUME *ConfigurationSynchronizing* ∈ STRING  
 ASSUME *ConfigurationSynchronized* ∈ STRING  
 ASSUME *ConfigurationPersisted* ∈ STRING  
 ASSUME *ConfigurationFailed* ∈ STRING  
  
 ASSUME  $\wedge IsFiniteSet(Node)$   
      $\wedge \forall n \in Node :$   
          $\wedge n \notin \text{DOMAIN } Target$   
          $\wedge n \in \text{STRING}$   
  
 ASSUME  $\wedge \forall t \in \text{DOMAIN } Target :$   
      $\wedge t \notin Node$   
      $\wedge t \in \text{STRING}$   
      $\wedge Target[t].persistent \in \text{BOOLEAN}$   
      $\wedge \forall p \in \text{DOMAIN } Target[t].paths :$   
          $IsFiniteSet(Target[t].paths[p])$

Configuration update/rollback requests are tracked and processed through two data types. Transactions represent the lifecycle of a single configuration change request and are stored in an append-only log. Configurations represent the desired configuration of a *gNMI* target based on the aggregate of relevant changes in the Transaction log.

```

TYPE TransactionType ::= type ∈
  { TransactionChange,
    TransactionRollback }

TYPE TransactionStatus ::= status ∈
  { TransactionInitializing,
    TransactionInitialized,
    TransactionValidating,
    TransactionValidated,
    TransactionCommitting,
    TransactionCommitted,
    TransactionApplying,
    TransactionApplied,
    TransactionFailed }

TYPE Transaction  $\triangleq$  [
  type      ::= type ∈ TransactionType,
  index     ::= index ∈ Nat,
  isolation ::= isolation ∈ { IsolationDefault, IsolationSerializable }
  values ::= [
    target ∈ SUBSET (DOMAIN Target)  $\mapsto$  [ path ∈ SUBSET (DOMAIN Target[target].paths)  $\mapsto$ 
    [
      value ::= value ∈ STRING,
      delete ::= delete ∈ BOOLEAN ]],
    rollback ::= index ∈ Nat,
    targets ::= targets ∈ SUBSET (DOMAIN Target)
    status ::= status ∈ TransactionStatus]

TYPE ProposalStatus ::= status ∈
  { ProposalInitializing,
    ProposalInitialized,
    ProposalValidating,
    ProposalValidated,
    ProposalCommitting,
    ProposalCommitted,
    ProposalApplying,
    ProposalApplied,
    ProposalFailed }

TYPE Proposal  $\triangleq$  [
  index      ::= index ∈ Nat,
  values     ::= [ path ∈ SUBSET (DOMAIN Target[target].paths)  $\mapsto$  [
    value ::= value ∈ STRING,
    delete ::= delete ∈ BOOLEAN ]],
  rollback   ::= index ∈ Nat,
  prevIndex  ::= prevIndex ∈ Nat,

```

```

nextIndex      ::= nextIndex ∈ Nat,
rollbackIndex ::= rollbackIndex ∈ Nat,
rollbackValues ::= [ path ∈ SUBSET (DOMAIN Target[target].paths) ↦ [
    value ::= value ∈ STRING,
    delete ::= delete ∈ BOOLEAN ]],
status         ::= status ∈ ProposalStatus]

TYPE ConfigurationStatus ::= status ∈
{ ConfigurationUnknown,
  ConfigurationSynchronizing,
  ConfigurationSynchronized,
  ConfigurationPersisted,
  ConfigurationFailed}

TYPE Configuration  $\triangleq$  [
  id          ::= id ∈ STRING,
  target      ::= target ∈ STRING,
  values      ::= [ path ∈ SUBSET (DOMAIN Target[target]) ↦ [
    value ::= value ∈ STRING,
    index ::= index ∈ Nat,
    deleted ::= delete ∈ BOOLEAN ]],
  configIndex ::= index ∈ Nat,
  proposedIndex ::= proposedIndex ∈ Nat,
  committedIndex ::= committedIndex ∈ Nat,
  appliedIndex ::= appliedIndex ∈ Nat,
  appliedTerm ::= appliedTerm ∈ Nat,
  appliedValues ::= [ path ∈ SUBSET (DOMAIN Target[target]) ↦ [
    value ::= value ∈ STRING,
    index ::= index ∈ Nat,
    deleted ::= delete ∈ BOOLEAN ]],
  status ::= status ∈ ConfigurationStatus]

```

A transaction log. Transactions may either request a set of changes to a set of targets or rollback a prior change.

VARIABLE *transaction*

A record of per-target proposals

VARIABLE *proposal*

A record of per-target configurations

VARIABLE *configuration*

A record of target states

VARIABLE *target*

A record of target masterhips

VARIABLE *mastership*

$\text{vars} \triangleq \langle \text{transaction}, \text{proposal}, \text{configuration}, \text{mastership}, \text{target} \rangle$

This section models *mastership* for the configuration service.

Mastership is used primarily to track the lifecycle of individual configuration targets and react to state changes on the southbound. Each target is assigned a master from the *Node* set, and masters can be unset when the target disconnects.

Set node  $n$  as the master for target  $t$   
 $SetMaster(n, t) \triangleq$   
 $\wedge mastership[t].master \neq n$   
 $\wedge mastership' = [mastership \text{ EXCEPT } ![t].term = mastership[t].term + 1,$   
 $![t].master = n]$   
 $\wedge \text{UNCHANGED } \langle transaction, proposal, configuration, target \rangle$

UnsetMaster( $t$ )  $\triangleq$   
 $\wedge mastership[t].master \neq Nil$   
 $\wedge mastership' = [mastership \text{ EXCEPT } ![t].master = Nil]$   
 $\wedge \text{UNCHANGED } \langle transaction, proposal, configuration, target \rangle$

This section models configuration changes and rollbacks. Changes are appended to the transaction log and processed asynchronously.

$Value(s, t, p) \triangleq$   
 $\text{LET } value \triangleq \text{CHOOSE } v \in s : v.target = t \wedge v.path = p$   
 $\text{IN}$   
 $[value \mapsto value.value,$   
 $delete \mapsto value.delete]$

$Paths(s, t) \triangleq$   
 $[p \in \{v.path : v \in \{v \in s : v.target = t\}\} \mapsto Value(s, t, p)]$

$Changes(s) \triangleq$   
 $[t \in \{v.target : v \in s\} \mapsto Paths(s, t)]$

$ValidValues(t, p) \triangleq$   
 $\text{UNION } \{[value \mapsto v, delete \mapsto \text{FALSE}] : v \in Target[t][p]\}, \{[value \mapsto Nil, delete \mapsto \text{TRUE}]\}$

$ValidPaths(t) \triangleq$   
 $\text{UNION } \{[v @@@ [path \mapsto p] : v \in ValidValues(t, p)] : p \in \text{DOMAIN } Target[t]\}$

$ValidTargets \triangleq$   
 $\text{UNION } \{[p @@@ [target \mapsto t] : p \in ValidPaths(t)] : t \in \text{DOMAIN } Target\}$

The set of all valid sets of changes to all targets and their paths.

The set of possible changes is computed from the *Target* model value.

$ValidChanges \triangleq$   
 $\text{LET } changeSets \triangleq \{s \in \text{SUBSET } ValidTargets :$   
 $\forall t \in \text{DOMAIN } Target :$   
 $\forall p \in \text{DOMAIN } Target[t] :$   
 $Cardinality(\{v \in s : v.target = t \wedge v.path = p\}) \leq 1\}$

IN  
 $\{ \text{Changes}(s) : s \in \text{changeSets} \}$

The next available index in the transaction log.

This is computed as the max of the existing indexes in the log to allow for changes to the log (*e.g.* log compaction) to be modeled.

$\text{NextIndex} \triangleq$

IF DOMAIN  $\text{transaction} = \{\}$  THEN  
 1  
 ELSE  
 LET  $i \triangleq$  CHOOSE  $i \in \text{DOMAIN } \text{transaction} :$   
 $\forall j \in \text{DOMAIN } \text{transaction} : i \geq j$   
 IN  $i + 1$

Add a set of changes 'c' to the transaction log

$\text{Change}(c) \triangleq$

$\wedge \exists \text{isolation} \in \{ \text{IsolationDefault}, \text{IsolationSerializable} \} :$   
 $\wedge \text{transaction}' = \text{transaction} @ @ (\text{NextIndex} :> [ \text{type} \mapsto \text{TransactionChange},$   
 $\text{index} \mapsto \text{NextIndex},$   
 $\text{isolation} \mapsto \text{isolation},$   
 $\text{values} \mapsto c,$   
 $\text{targets} \mapsto \{\},$   
 $\text{status} \mapsto \text{TransactionInitializing} ] )$   
 $\wedge \text{UNCHANGED } \langle \text{proposal}, \text{configuration}, \text{mastership}, \text{target} \rangle$

Add a rollback of transaction 't' to the transaction log

$\text{Rollback}(t) \triangleq$

$\wedge \exists \text{isolation} \in \{ \text{IsolationDefault}, \text{IsolationSerializable} \} :$   
 $\wedge \text{transaction}' = \text{transaction} @ @ (\text{NextIndex} :> [ \text{type} \mapsto \text{TransactionRollback},$   
 $\text{index} \mapsto \text{NextIndex},$   
 $\text{isolation} \mapsto \text{isolation},$   
 $\text{rollback} \mapsto t,$   
 $\text{targets} \mapsto \{\},$   
 $\text{status} \mapsto \text{TransactionInitializing} ] )$   
 $\wedge \text{UNCHANGED } \langle \text{proposal}, \text{configuration}, \text{mastership}, \text{target} \rangle$

This section models the Transaction log reconciler.

Transactions come in two flavors : – *Change* transactions contain a set of changes to be applied to a set of *targets* – *Rollback* transactions reference a prior change transaction to be reverted to the previous state

Both types of transaction are reconciled in stages:

- \* Pending - waiting for prior transactions to complete
- \* Validating - validating the requested changes
- \* Applying - applying the changes to target configurations
- \* Complete - completed applying changes successfully

\* Failed - failed applying changes

Reconcile a transaction

$ReconcileTransaction(n, i) \triangleq$

$$\begin{aligned}
& \wedge \vee \wedge transaction[i].status = TransactionInitializing \\
& \wedge i - 1 \in transaction \Rightarrow \\
& \quad TransactionPhase(transaction[i - 1].status) > TransactionPhase(TransactionInitializing) \\
& \wedge \vee \wedge transaction[i].targets = \{\} \\
& \quad \wedge \vee \wedge transaction[i].type = TransactionChange \\
& \quad \quad \wedge transaction' = [transaction \text{ EXCEPT } ![i].targets = \text{DOMAIN } transaction[i].values] \\
& \quad \quad \wedge proposal' = [t \in \text{DOMAIN } proposal \mapsto proposal[t] @@ \\
& \quad \quad \quad (i :> [type \mapsto ProposalChange, \\
& \quad \quad \quad \quad index \mapsto i, \\
& \quad \quad \quad \quad values \mapsto transaction[i].changes[t], \\
& \quad \quad \quad \quad status \mapsto ProposalInitializing])]] \\
& \quad \vee \wedge transaction[i].type = TransactionRollback \\
& \quad \quad \wedge \vee \wedge transaction[i].rollback \in \text{DOMAIN } transaction \\
& \quad \quad \quad \wedge transaction[transaction[i].rollback].type = TransactionChange \\
& \quad \quad \quad \wedge transaction' = [transaction \text{ EXCEPT } ![i].targets = \\
& \quad \quad \quad \quad \text{DOMAIN } transaction[transaction[i].rollback].values] \\
& \quad \quad \quad \wedge proposal' = [t \in \text{DOMAIN } proposal \mapsto proposal[t] @@ \\
& \quad \quad \quad \quad (i :> [type \mapsto ProposalRollback, \\
& \quad \quad \quad \quad \quad index \mapsto i, \\
& \quad \quad \quad \quad \quad rollback \mapsto transaction[i].rollback, \\
& \quad \quad \quad \quad \quad status \mapsto ProposalInitializing])]] \\
& \quad \vee \wedge \vee \wedge transaction[i].rollback \in \text{DOMAIN } transaction \\
& \quad \quad \quad \wedge transaction[transaction[i].rollback].type = TransactionRollback \\
& \quad \quad \quad \vee transaction[i].rollback \notin \text{DOMAIN } transaction \\
& \quad \quad \quad \wedge transaction' = [transaction \text{ EXCEPT } ![i].status = TransactionFailed] \\
& \quad \quad \quad \wedge \text{UNCHANGED } \langle proposal \rangle \\
& \quad \vee \wedge transaction[i].targets \neq \{\} \\
& \quad \quad \wedge \vee \wedge \exists t \in transaction[i].targets : \\
& \quad \quad \quad \wedge proposal[t][i].status = ProposalFailed \\
& \quad \quad \quad \wedge transaction' = [transaction \text{ EXCEPT } ![i].status = TransactionFailed] \\
& \quad \quad \vee \wedge \forall t \in transaction[i].targets : \\
& \quad \quad \quad \wedge proposal[t][i].status = ProposalInitialized \\
& \quad \quad \quad \wedge transaction' = [transaction \text{ EXCEPT } ![i].status = TransactionInitialized] \\
& \vee \wedge transaction[i].status = TransactionInitialized \\
& \quad \wedge \forall t \in transaction[i].targets : \\
& \quad \quad proposal[t][i].prevIndex \neq 0 \Rightarrow \\
& \quad \quad \quad (transaction[proposal[t][i].prevIndex].isolation = IsolationSerializable \Rightarrow \\
& \quad \quad \quad \quad TransactionPhase(transaction[proposal[t][i].prevIndex].status) \geq \\
& \quad \quad \quad \quad \quad TransactionPhase(TransactionValidated)) \\
& \quad \quad \wedge transaction' = [transaction \text{ EXCEPT } ![i].status = TransactionValidating] \\
& \quad \quad \wedge \text{UNCHANGED } \langle proposal \rangle \\
& \vee \wedge transaction[i].status = TransactionValidating
\end{aligned}$$



$$\begin{aligned}
& \wedge \vee \wedge \exists t \in \text{transaction}[i].\text{targets} : \text{proposal}[t][i].\text{status} \neq \text{ProposalValidating} \\
& \wedge \text{proposal}' = [t \in \text{DOMAIN } \text{proposal} \mapsto \\
& \quad \text{IF } t \in \text{transaction}[i].\text{targets} \text{ THEN} \\
& \quad \quad [\text{proposal}[t] \text{ EXCEPT } ![i].\text{status} = \text{ProposalValidating}] \\
& \quad \text{ELSE} \\
& \quad \quad \text{proposal}[t]] \\
& \wedge \text{UNCHANGED } \langle \text{transaction} \rangle \\
\vee & \wedge \text{transaction}[i].\text{status} = \text{TransactionValidated} \\
& \wedge \forall t \in \text{transaction}[i].\text{targets} : \\
& \quad \text{proposal}[t][i].\text{prevIndex} \neq 0 \Rightarrow \\
& \quad \quad (\text{transaction}[\text{proposal}[t][i].\text{prevIndex}].\text{isolation} = \text{IsolationSerializable} \Rightarrow \\
& \quad \quad \quad \text{TransactionPhase}(\text{transaction}[\text{proposal}[t][i].\text{prevIndex}].\text{status}) \geq \\
& \quad \quad \quad \text{TransactionPhase}(\text{TransactionCommitted})) \\
& \wedge \text{transaction}' = [\text{transaction} \text{ EXCEPT } ![i].\text{status} = \text{TransactionCommitting}] \\
& \wedge \text{UNCHANGED } \langle \text{proposal} \rangle \\
\vee & \wedge \text{transaction}[i].\text{status} = \text{TransactionCommitting} \\
& \wedge \vee \wedge \exists t \in \text{transaction}[i].\text{targets} : \text{proposal}[t][i].\text{status} \neq \text{ProposalCommitting} \\
& \wedge \text{proposal}' = [t \in \text{DOMAIN } \text{proposal} \mapsto \\
& \quad \text{IF } t \in \text{transaction}[i].\text{targets} \text{ THEN} \\
& \quad \quad [\text{proposal}[t] \text{ EXCEPT } ![i].\text{status} = \text{ProposalCommitting}] \\
& \quad \text{ELSE} \\
& \quad \quad \text{proposal}[t]] \\
& \wedge \text{UNCHANGED } \langle \text{transaction} \rangle \\
\vee & \wedge \text{transaction}[i].\text{status} = \text{TransactionCommitted} \\
& \wedge \forall t \in \text{transaction}[i].\text{targets} : \\
& \quad \text{proposal}[t][i].\text{prevIndex} \neq 0 \Rightarrow \\
& \quad \quad (\text{transaction}[\text{proposal}[t][i].\text{prevIndex}].\text{isolation} = \text{IsolationSerializable} \Rightarrow \\
& \quad \quad \quad \text{TransactionPhase}(\text{transaction}[\text{proposal}[t][i].\text{prevIndex}].\text{status}) \geq \\
& \quad \quad \quad \text{TransactionPhase}(\text{TransactionApplied})) \\
& \wedge \text{transaction}' = [\text{transaction} \text{ EXCEPT } ![i].\text{status} = \text{TransactionApplying}] \\
& \wedge \vee \wedge \exists t \in \text{transaction}[i].\text{targets} : \text{proposal}[t][i].\text{status} \neq \text{ProposalApplying} \\
& \wedge \text{proposal}' = [t \in \text{DOMAIN } \text{proposal} \mapsto \\
& \quad \text{IF } t \in \text{transaction}[i].\text{targets} \text{ THEN} \\
& \quad \quad [\text{proposal}[t] \text{ EXCEPT } ![i].\text{status} = \text{ProposalApplying}] \\
& \quad \text{ELSE} \\
& \quad \quad \text{proposal}[t]] \\
& \wedge \text{UNCHANGED } \langle \text{transaction} \rangle \\
& \wedge \text{UNCHANGED } \langle \text{proposal} \rangle \\
\vee & \wedge \text{transaction}[i].\text{status} = \text{TransactionApplying} \\
\vee & \wedge \text{transaction}[i].\text{status} = \text{TransactionApplied} \\
& \wedge \text{UNCHANGED } \langle \text{configuration}, \text{mastership}, \text{target} \rangle
\end{aligned}$$

Reconcile a proposal

$\text{ReconcileProposal}(n, t, i) \triangleq$

$$\wedge \vee \wedge \text{proposal}[t][i].\text{status} = \text{ProposalInitializing}$$

$$\begin{aligned}
& \wedge \vee \wedge \text{configuration}[t].\text{proposedIndex} > 0 \\
& \wedge \text{proposal}' = [\text{proposal} \text{ EXCEPT } ![t] = [\text{proposal}[t] \text{ EXCEPT} \\
& \quad ![i] = [\text{status} \mapsto \text{ProposalInitialized}, \\
& \quad \text{prevIndex} \mapsto \text{configuration}[t].\text{proposedIndex}] @@ \text{proposal}[t][i], \\
& \quad ![\text{configuration}[t].\text{proposedIndex}] = [\text{nextIndex} \mapsto i] @@ \\
& \quad \text{proposal}[t][\text{configuration}[t].\text{proposedIndex}]] \\
& \vee \wedge \text{configuration}[t].\text{proposedIndex} = 0 \\
& \wedge \text{proposal}' = [\text{proposal} \text{ EXCEPT } ![t] = [\text{proposal}[t] \text{ EXCEPT } ![i].\text{status} = \text{ProposalInitialized}]] \\
& \wedge \text{configuration}' = [\text{configuration} \text{ EXCEPT } ![t].\text{proposedIndex} = i] \\
& \wedge \text{UNCHANGED } \langle \text{target} \rangle \\
\vee & \wedge \text{proposal}[t][i].\text{status} = \text{ProposalValidating} \\
& \wedge \text{configuration}[t].\text{committedIndex} = \text{proposal}[t][i].\text{prevIndex} \\
& \wedge \vee \wedge \text{proposal}[t][i].\text{type} = \text{ProposalChange} \\
& \quad \wedge \text{LET } \text{rollbackIndex} \triangleq \text{configuration}[t].\text{configIndex} \\
& \quad \text{rollbackValues} \triangleq [p \in \text{DOMAIN } \text{proposal}[t][i].\text{values} \mapsto [ \\
& \quad \quad p \mapsto \text{IF } p \in \text{DOMAIN } \text{configuration}[t].\text{config} \text{ THEN} \\
& \quad \quad \text{configuration}[t].\text{values}[p] \\
& \quad \quad \text{ELSE} \\
& \quad \quad [\text{delete} \mapsto \text{TRUE}]]] \\
& \text{IN} \\
& \wedge \text{proposal}' = [\text{proposal} \text{ EXCEPT } ![t] = [ \\
& \quad \text{proposal}[t] \text{ EXCEPT } ![i].\text{rollbackIndex} = \text{rollbackIndex}, \\
& \quad \quad ![i].\text{rollbackValues} = \text{rollbackValues}, \\
& \quad \quad ![i].\text{status} = \text{ProposalValidated}]] \\
& \vee \wedge \text{proposal}[t][i].\text{type} = \text{ProposalRollback} \\
& \wedge \vee \wedge \text{configuration}[t].\text{index} = \text{proposal}[t][i].\text{rollback} \\
& \wedge \vee \wedge \text{proposal}[t][i].\text{rollback} \in \text{DOMAIN } \text{proposal}[t] \\
& \quad \wedge \vee \wedge \text{proposal}[t][\text{proposal}[t][i].\text{rollback}].\text{type} = \text{ProposalChange} \\
& \quad \wedge \text{LET } \text{rollbackIndex} \triangleq \text{proposal}[t][\text{proposal}[t][i].\text{rollback}].\text{rollbackIndex} \\
& \quad \text{rollbackValues} \triangleq \text{proposal}[t][\text{proposal}[t][i].\text{rollback}].\text{rollbackValues} \\
& \quad \text{IN } \text{proposal}' = [\text{proposal} \text{ EXCEPT } ![t] = [ \\
& \quad \text{proposal}[t] \text{ EXCEPT } ![i].\text{rollbackIndex} = \text{rollbackIndex}, \\
& \quad \quad ![i].\text{rollbackValues} = \text{rollbackValues}, \\
& \quad \quad ![i].\text{status} = \text{ProposalValidated}]] \\
& \vee \wedge \text{proposal}[t][\text{proposal}[t][i].\text{rollback}].\text{type} = \text{ProposalRollback} \\
& \wedge \text{configuration}' = [\text{configuration} \text{ EXCEPT } ![t].\text{committedIndex} = i] \\
& \wedge \text{proposal}' = [\text{proposal} \text{ EXCEPT } ![t] = [ \\
& \quad \text{proposal}[t] \text{ EXCEPT } ![i].\text{status} = \text{ProposalFailed}]] \\
& \vee \wedge \text{proposal}[t][i].\text{rollback} \notin \text{DOMAIN } \text{proposal}[t] \\
& \wedge \text{configuration}' = [\text{configuration} \text{ EXCEPT } ![t].\text{committedIndex} = i] \\
& \wedge \text{proposal}' = [\text{proposal} \text{ EXCEPT } ![t] = [ \\
& \quad \text{proposal}[t] \text{ EXCEPT } ![i].\text{status} = \text{ProposalFailed}]] \\
& \vee \wedge \text{configuration}[t].\text{index} \neq \text{proposal}[t][i].\text{rollback} \\
& \wedge \text{configuration}' = [\text{configuration} \text{ EXCEPT } ![t].\text{committedIndex} = i] \\
& \wedge \text{proposal}' = [\text{proposal} \text{ EXCEPT } ![t] = [\text{proposal}[t] \text{ EXCEPT } ![i].\text{status} = \text{ProposalFailed}]]
\end{aligned}$$



$\wedge \text{UNCHANGED } \langle \text{proposal}, \text{transaction}, \text{mastership} \rangle$

---

*Init* and next state predicates

*Init*  $\triangleq$

$\wedge \text{transaction} = \langle \rangle$   
 $\wedge \text{proposal} = [t \in \text{DOMAIN } \text{Target} \mapsto$   
 $\quad [p \in \{\} \mapsto [\text{status} \mapsto \text{ProposalInitializing}]]]$   
 $\wedge \text{configuration} = [t \in \text{DOMAIN } \text{Target} \mapsto$   
 $\quad [target \mapsto t,$   
 $\quad \text{status} \mapsto \text{ConfigurationUnknown},$   
 $\quad \text{values} \mapsto$   
 $\quad \quad [path \in \{\} \mapsto$   
 $\quad \quad \quad [path \mapsto path,$   
 $\quad \quad \quad \text{value} \mapsto \text{Nil},$   
 $\quad \quad \quad \text{index} \mapsto 0,$   
 $\quad \quad \quad \text{deleted} \mapsto \text{FALSE}]],$   
 $\quad \text{configIndex} \mapsto 0,$   
 $\quad \text{proposedIndex} \mapsto 0,$   
 $\quad \text{committedIndex} \mapsto 0,$   
 $\quad \text{appliedIndex} \mapsto 0,$   
 $\quad \text{appliedTerm} \mapsto 0,$   
 $\quad \text{appliedValues} \mapsto$   
 $\quad \quad [path \in \{\} \mapsto$   
 $\quad \quad \quad [path \mapsto path,$   
 $\quad \quad \quad \text{value} \mapsto \text{Nil},$   
 $\quad \quad \quad \text{index} \mapsto 0,$   
 $\quad \quad \quad \text{deleted} \mapsto \text{FALSE}]]]]]$   
 $\wedge \text{target} = [t \in \text{DOMAIN } \text{Target} \mapsto$   
 $\quad [path \in \{\} \mapsto$   
 $\quad \quad [\text{value} \mapsto \text{Nil}]]]$   
 $\wedge \text{mastership} = [t \in \text{DOMAIN } \text{Target} \mapsto [\text{master} \mapsto \text{Nil}, \text{term} \mapsto 0]]$

*Next*  $\triangleq$

$\vee \exists c \in \text{ValidChanges} :$   
 $\quad \text{Change}(c)$   
 $\vee \exists t \in \text{DOMAIN } \text{transaction} :$   
 $\quad \text{Rollback}(t)$   
 $\vee \exists n \in \text{Node} :$   
 $\quad \exists t \in \text{DOMAIN } \text{Target} :$   
 $\quad \quad \text{SetMaster}(n, t)$   
 $\vee \exists t \in \text{DOMAIN } \text{Target} :$   
 $\quad \text{UnsetMaster}(t)$   
 $\vee \exists n \in \text{Node} :$   
 $\quad \exists t \in \text{DOMAIN } \text{transaction} :$

$$\begin{aligned}
& \text{ReconcileTransaction}(n, t) \\
\forall \exists n \in \text{Node} : \\
& \quad \exists c \in \text{DOMAIN configuration} : \\
& \quad \quad \text{ReconcileConfiguration}(n, c) \\
\text{Spec} & \triangleq \text{Init} \wedge \Box[\text{Next}]_{\text{vars}} \\
\text{Order} & \triangleq \text{TRUE} \quad \text{TODO redefine order spec} \\
\text{THEOREM Safety} & \triangleq \text{Spec} \Rightarrow \Box \text{Order} \\
\text{Completion} & \triangleq \forall i \in \text{DOMAIN transaction} : \\
& \quad \text{transaction}[i].\text{status} \in \{\text{TransactionApplied}, \text{TransactionFailed}\} \\
\text{THEOREM Liveness} & \triangleq \text{Spec} \Rightarrow \Diamond \text{Completion}
\end{aligned}$$


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