
MODULE *Config*

INSTANCE *Naturals*

INSTANCE *FiniteSets*

INSTANCE *Sequences*

INSTANCE *TLC*

GenerateTestCases \triangleq FALSE

Nil \triangleq "<nil>"

Change \triangleq "Change"

Rollback \triangleq "Rollback"

Commit \triangleq "Commit"

Apply \triangleq "Apply"

Pending \triangleq "Pending"

InProgress \triangleq "InProgress"

Complete \triangleq "Complete"

Aborted \triangleq "Aborted"

Failed \triangleq "Failed"

Done \triangleq { *Complete*, *Aborted*, *Failed* }

Node \triangleq { "node1" }

NumTransactions \triangleq 4

NumTerms \triangleq 2

NumConns \triangleq 2

Path \triangleq { "path1" }

Value \triangleq { "value1", "value2" }

A transaction log. Transactions may either request a set of changes to a set of targets or rollback a prior change.

VARIABLE *transaction*

A record of per-target proposals

VARIABLE *proposal*

A record of per-target configurations

VARIABLE *configuration*

A record of target masterhips

VARIABLE *mastership*

A record of node connections to the target

VARIABLE *conn*

The target state

VARIABLE *target*

A sequence of state changes used for model checking.

VARIABLE *history*

$vars \triangleq \langle transaction, proposal, configuration, mastership, conn, target, history \rangle$

LOCAL *Transaction* \triangleq INSTANCE *Transaction*

LOCAL *Proposal* \triangleq INSTANCE *Proposal*

LOCAL *Configuration* \triangleq INSTANCE *Configuration*

LOCAL *Mastership* \triangleq INSTANCE *Mastership*

LOCAL *Target* \triangleq INSTANCE *Target*

$RequestChange(p, v) \triangleq$
 $\wedge Transaction!RequestChange(p, v)$
 $\wedge UNCHANGED \langle mastership, conn, target, history \rangle$

$RequestRollback(i) \triangleq$
 $\wedge Transaction!RequestRollback(i)$
 $\wedge UNCHANGED \langle mastership, conn, target, history \rangle$

$ReconcileTransaction(n, i) \triangleq$
 $\wedge i \in \text{DOMAIN } transaction$
 $\wedge Transaction!ReconcileTransaction(n, i)$
 $\wedge UNCHANGED \langle mastership, conn, target, history \rangle$
 $\wedge GenerateTestCases \Rightarrow Transaction!Test!Log([node \mapsto n, index \mapsto i])$

$ReconcileProposal(n, i) \triangleq$
 $\wedge i \in \text{DOMAIN } proposal$
 $\wedge Proposal!ReconcileProposal(n, i)$
 $\wedge UNCHANGED \langle transaction \rangle$
 $\wedge GenerateTestCases \Rightarrow Proposal!Test!Log([node \mapsto n, index \mapsto i])$

$ReconcileConfiguration(n) \triangleq$
 $\wedge Configuration!ReconcileConfiguration(n)$
 $\wedge UNCHANGED \langle transaction, proposal, history \rangle$

$$\begin{aligned}
& \wedge \text{GenerateTestCases} \Rightarrow \text{Configuration!Test!Log}([node \mapsto n]) \\
\text{ReconcileMastership}(n) & \triangleq \\
& \wedge \text{Mastership!ReconcileMastership}(n) \\
& \wedge \text{UNCHANGED} \langle \text{transaction}, \text{proposal}, \text{configuration}, \text{target}, \text{history} \rangle \\
& \wedge \text{GenerateTestCases} \Rightarrow \text{Mastership!Test!Log}([node \mapsto n]) \\
\text{ConnectNode}(n) & \triangleq \\
& \wedge \text{Target!Connect}(n) \\
& \wedge \text{UNCHANGED} \langle \text{transaction}, \text{proposal}, \text{configuration}, \text{mastership}, \text{history} \rangle \\
\text{DisconnectNode}(n) & \triangleq \\
& \wedge \text{Target!Disconnect}(n) \\
& \wedge \text{UNCHANGED} \langle \text{transaction}, \text{proposal}, \text{configuration}, \text{mastership}, \text{history} \rangle \\
\text{StartTarget} & \triangleq \\
& \wedge \text{Target!Start} \\
& \wedge \text{UNCHANGED} \langle \text{transaction}, \text{proposal}, \text{configuration}, \text{mastership}, \text{history} \rangle \\
\text{StopTarget} & \triangleq \\
& \wedge \text{Target!Stop} \\
& \wedge \text{UNCHANGED} \langle \text{transaction}, \text{proposal}, \text{configuration}, \text{mastership}, \text{history} \rangle
\end{aligned}$$

Formal specification, constraints, and theorems.

$$\begin{aligned}
\text{Init} & \triangleq \\
& \wedge \text{transaction} = [\\
& \quad i \in \{\} \mapsto [\\
& \quad \quad \text{type} \mapsto \text{Change}, \\
& \quad \quad \text{index} \mapsto 0, \\
& \quad \quad \text{values} \mapsto [p \in \{\} \mapsto \text{Nil}], \\
& \quad \quad \text{commit} \mapsto \text{Pending}, \\
& \quad \quad \text{apply} \mapsto \text{Pending}] \\
& \wedge \text{proposal} = [\\
& \quad i \in \{\} \mapsto [\\
& \quad \quad \text{change} \mapsto [\\
& \quad \quad \quad \text{phase} \mapsto \text{Nil}, \\
& \quad \quad \quad \text{state} \mapsto \text{Nil}, \\
& \quad \quad \quad \text{values} \mapsto [\\
& \quad \quad \quad \quad p \in \{\} \mapsto [\\
& \quad \quad \quad \quad \quad \text{index} \mapsto 0, \\
& \quad \quad \quad \quad \quad \text{value} \mapsto \text{Nil}]]], \\
& \quad \text{rollback} \mapsto [\\
& \quad \quad \text{phase} \mapsto \text{Nil}, \\
& \quad \quad \text{state} \mapsto \text{Nil}, \\
& \quad \quad \text{values} \mapsto [
\end{aligned}$$

$$\begin{aligned}
& p \in \{\} \mapsto [\\
& \quad index \mapsto 0, \\
& \quad value \mapsto Nil]]]]] \\
\wedge configuration = [\\
& \quad state \mapsto InProgress, \\
& \quad term \mapsto 0, \\
& \quad committed \mapsto [\\
& \quad \quad index \mapsto 0, \\
& \quad \quad revision \mapsto 0, \\
& \quad \quad values \mapsto [\\
& \quad \quad \quad p \in \{\} \mapsto [\\
& \quad \quad \quad \quad index \mapsto 0, \\
& \quad \quad \quad \quad value \mapsto Nil]]], \\
& \quad applied \mapsto [\\
& \quad \quad index \mapsto 0, \\
& \quad \quad revision \mapsto 0, \\
& \quad \quad target \mapsto 0, \\
& \quad \quad values \mapsto [\\
& \quad \quad \quad p \in \{\} \mapsto [\\
& \quad \quad \quad \quad index \mapsto 0, \\
& \quad \quad \quad \quad value \mapsto Nil]]]]] \\
\wedge target = [\\
& \quad id \mapsto 0, \\
& \quad running \mapsto FALSE, \\
& \quad values \mapsto [\\
& \quad \quad p \in \{\} \mapsto [\\
& \quad \quad \quad index \mapsto 0, \\
& \quad \quad \quad value \mapsto Nil]]] \\
\wedge mastership = [\\
& \quad master \mapsto Nil, \\
& \quad term \mapsto 0, \\
& \quad conn \mapsto 0] \\
\wedge conn = [\\
& \quad n \in Node \mapsto [\\
& \quad \quad id \mapsto 0, \\
& \quad \quad connected \mapsto FALSE]] \\
\wedge history = \langle \rangle \\
Next \triangleq \\
& \vee \exists p \in Path, v \in Value : \\
& \quad RequestChange(p, v) \\
& \vee \exists i \in DOMAIN transaction : \\
& \quad RequestRollback(i) \\
& \vee \exists n \in Node : \\
& \quad \exists i \in DOMAIN transaction :
\end{aligned}$$

$$\begin{aligned}
& \text{ReconcileTransaction}(n, i) \\
\vee \exists n \in \text{Node} : \\
& \quad \exists i \in \text{DOMAIN } \text{proposal} : \\
& \quad \quad \text{ReconcileProposal}(n, i) \\
\vee \exists n \in \text{Node} : \\
& \quad \text{ReconcileConfiguration}(n) \\
\vee \exists n \in \text{Node} : \\
& \quad \text{ReconcileMastership}(n) \\
\vee \exists n \in \text{Node} : \\
& \quad \vee \text{ConnectNode}(n) \\
& \quad \vee \text{DisconnectNode}(n) \\
\vee \text{StartTarget} \\
\vee \text{StopTarget} \\
\text{Spec} \triangleq & \\
& \wedge \text{Init} \\
& \wedge \Box[\text{Next}]_{\text{vars}} \\
& \wedge \forall p \in \text{Path}, v \in \text{Value} : \\
& \quad \text{WF}_{\langle \text{transaction}, \text{proposal}, \text{configuration}, \text{mastership}, \text{target} \rangle}(\text{Transaction!RequestChange}(p, v)) \\
& \wedge \forall i \in 1 \dots \text{NumTransactions} : i \in \text{DOMAIN } \text{transaction} \Rightarrow \\
& \quad \text{WF}_{\langle \text{transaction}, \text{proposal}, \text{configuration}, \text{mastership}, \text{target} \rangle}(\text{Transaction!RequestRollback}(i)) \\
& \wedge \forall n \in \text{Node}, i \in 1 \dots \text{NumTransactions} : \\
& \quad \text{WF}_{\langle \text{transaction}, \text{proposal}, \text{configuration}, \text{mastership}, \text{target} \rangle}(\text{Transaction!ReconcileTransaction}(n, i)) \\
& \wedge \forall n \in \text{Node}, i \in 1 \dots \text{NumTransactions} : \\
& \quad \text{WF}_{\langle \text{proposal}, \text{configuration}, \text{mastership}, \text{conn}, \text{target}, \text{history} \rangle}(\text{Proposal!ReconcileProposal}(n, i)) \\
& \wedge \forall n \in \text{Node} : \\
& \quad \text{WF}_{\langle \text{configuration}, \text{mastership}, \text{conn}, \text{target} \rangle}(\text{Configuration!ReconcileConfiguration}(n)) \\
& \wedge \forall n \in \text{Node} : \\
& \quad \text{WF}_{\langle \text{mastership}, \text{conn} \rangle}(\text{Mastership!ReconcileMastership}(n)) \\
& \wedge \forall n \in \text{Node} : \\
& \quad \text{WF}_{\langle \text{conn}, \text{target} \rangle}(\text{Target!Connect}(n) \vee \text{Target!Disconnect}(n)) \\
& \wedge \text{WF}_{\langle \text{conn}, \text{target} \rangle}(\text{Target!Start} \vee \text{Target!Stop})
\end{aligned}$$

$$\text{LimitTransactions} \triangleq \text{Len}(\text{transaction}) \leq \text{NumTransactions}$$

$$\begin{aligned}
\text{LimitTerms} & \triangleq \\
& \vee \text{mastership.term} < \text{NumTerms} \\
& \vee \wedge \text{mastership.term} = \text{NumTerms} \\
& \wedge \text{mastership.master} \neq \text{Nil}
\end{aligned}$$

$$\begin{aligned}
\text{LimitConns} & \triangleq \\
& \forall n \in \text{DOMAIN } \text{conn} : \\
& \quad \vee \text{conn}[n].\text{id} < \text{NumConns} \\
& \quad \vee \wedge \text{conn}[n].\text{id} = \text{NumConns}
\end{aligned}$$

$\wedge \text{conn}[n].\text{connected}$

$\text{TypeOK} \triangleq$

$\wedge \text{Transaction! TypeOK}$
 $\wedge \text{Proposal! TypeOK}$
 $\wedge \text{Configuration! TypeOK}$
 $\wedge \text{Mastership! TypeOK}$

LOCAL $\text{IsOrderedChange}(p, i) \triangleq$

$\wedge \text{history}[i].\text{type} = \text{Change}$
 $\wedge \text{history}[i].\text{phase} = p$
 $\wedge \neg \exists j \in \text{DOMAIN } \text{history} :$
 $\quad \wedge j < i$
 $\quad \wedge \text{history}[j].\text{type} = \text{Change}$
 $\quad \wedge \text{history}[j].\text{phase} = p$
 $\quad \wedge \text{history}[j].\text{index} \geq \text{history}[i].\text{index}$

LOCAL $\text{IsOrderedRollback}(p, i) \triangleq$

$\wedge \text{history}[i].\text{type} = \text{Rollback}$
 $\wedge \text{history}[i].\text{phase} = p$
 $\wedge \neg \exists j \in \text{DOMAIN } \text{history} :$
 $\quad \wedge j < i$
 $\quad \wedge \text{history}[j].\text{type} = \text{Change}$
 $\quad \wedge \text{history}[j].\text{phase} = p$
 $\quad \wedge \text{history}[j].\text{index} > \text{history}[i].\text{index}$
 $\wedge \neg \exists k \in \text{DOMAIN } \text{history} :$
 $\quad \wedge k > j$
 $\quad \wedge k < i$
 $\quad \wedge \text{history}[k].\text{type} = \text{Rollback}$
 $\quad \wedge \text{history}[k].\text{phase} = p$
 $\quad \wedge \text{history}[k].\text{index} = \text{history}[j].\text{index}$

$\text{Order} \triangleq$

$\wedge \forall i \in \text{DOMAIN } \text{history} :$
 $\quad \vee \text{IsOrderedChange}(\text{Commit}, i)$
 $\quad \vee \text{IsOrderedChange}(\text{Apply}, i)$
 $\quad \vee \text{IsOrderedRollback}(\text{Commit}, i)$
 $\quad \vee \text{IsOrderedRollback}(\text{Apply}, i)$
 $\wedge \forall i \in \text{DOMAIN } \text{proposal} :$
 $\quad \wedge \text{proposal}[i].\text{change.phase} = \text{Apply}$
 $\quad \wedge \text{proposal}[i].\text{change.state} = \text{Failed}$
 $\quad \wedge \text{proposal}[i].\text{rollback.phase} = \text{Apply} \Rightarrow \text{proposal}[i].\text{rollback.state} \neq \text{Complete}$
 $\Rightarrow \forall j \in \text{DOMAIN } \text{proposal} : (j > i \Rightarrow$
 $\quad (\text{proposal}[j].\text{change.phase} = \text{Apply} \Rightarrow$

$$proposal[j].change.state \in \{Nil, Pending, Aborted\}))$$

$$Consistency \triangleq$$

$$\begin{aligned}
& \wedge \forall i \in \text{DOMAIN } proposal : \\
& \quad \vee configuration.committed.index < i \\
& \quad \vee configuration.committed.revision < i \\
& \quad \Rightarrow \neg \exists p \in \text{DOMAIN } configuration.committed.values : \\
& \quad \quad configuration.committed.values[p].index = i \\
& \wedge \forall i \in \text{DOMAIN } proposal : \\
& \quad \vee configuration.applied.index < i \\
& \quad \vee configuration.applied.revision < i \\
& \quad \Rightarrow \wedge \neg \exists p \in \text{DOMAIN } configuration.applied.values : \\
& \quad \quad configuration.applied.values[p].index = i \\
& \quad \wedge \neg \exists p \in \text{DOMAIN } target.values : \\
& \quad \quad target.values[p].index = i \\
& \wedge \wedge target.running \\
& \wedge configuration.applied.target = target.id \\
& \wedge configuration.state = Complete \\
& \Rightarrow \forall i \in \text{DOMAIN } proposal : \\
& \quad \wedge configuration.applied.index \geq i \\
& \quad \wedge configuration.applied.revision \geq i \\
& \quad \Rightarrow \forall p \in \text{DOMAIN } proposal[i].change.values : \\
& \quad \quad \wedge \neg \exists j \in \text{DOMAIN } proposal : \\
& \quad \quad \quad \wedge j > i \\
& \quad \quad \quad \wedge configuration.applied.index \geq j \\
& \quad \quad \quad \wedge configuration.applied.revision \geq j \\
& \quad \Rightarrow \wedge p \in \text{DOMAIN } target.values \\
& \quad \quad \wedge target.values[p].value = proposal[i].change.values[p].value \\
& \quad \quad \wedge target.values[p].index = proposal[i].change.values[p].index
\end{aligned}$$

$$Safety \triangleq \Box (Order \wedge Consistency)$$

$$\text{THEOREM } Spec \Rightarrow Safety$$

$$Terminates(i) \triangleq$$

$$\begin{aligned}
& \wedge i \in \text{DOMAIN } transaction \\
& \wedge transaction[i].commit \in Done \\
& \wedge transaction[i].apply \in Done \\
& \wedge transaction[i].index \in \text{DOMAIN } proposal \\
& \wedge \vee \wedge transaction[i].type = Change \\
& \quad \wedge \vee \wedge proposal[transaction[i].index].change.phase = Commit \\
& \quad \quad \wedge proposal[transaction[i].index].change.state \in \{Aborted, Failed\} \\
& \quad \vee \wedge proposal[transaction[i].index].change.phase = Apply \\
& \quad \quad \wedge proposal[transaction[i].index].change.state \in Done \\
& \vee \wedge transaction[i].type = Rollback \\
& \quad \wedge \vee \wedge proposal[transaction[i].index].rollback.phase = Commit
\end{aligned}$$

$$\begin{aligned}
& \wedge \text{proposal}[\text{transaction}[i].\text{index}].\text{rollback.state} \in \{\text{Aborted}, \text{Failed}\} \\
& \vee \wedge \text{proposal}[\text{transaction}[i].\text{index}].\text{rollback.phase} = \text{Apply} \\
& \wedge \text{proposal}[\text{transaction}[i].\text{index}].\text{rollback.state} \in \text{Done}
\end{aligned}$$

$$\begin{aligned}
\text{Termination} & \triangleq \\
& \forall i \in 1 \dots \text{NumTransactions} : \Diamond \text{Terminates}(i)
\end{aligned}$$

$$\text{Liveness} \triangleq \text{Termination}$$

THEOREM $\text{Spec} \Rightarrow \text{Liveness}$