WHAT GRAPH NEURAL NETWORKS CANNOT LEARN: DEPTH VS WIDTH

Paper Author: Andreas Loukas

Presented by: Hauton J Tsang



Outline

- Introduction
- Background
- Importance of problem
- Limitations of previous methods
- Solutions of problem
- Interesting research questions



INTRODUCTION

Introduction

- Expressivity of a machine learning model is important to know
- Universal approximation for neural networks
 - A large enough neural network can solve any problem
 - How useful is this information?
- What a model cannot learn may be more useful
 - Establish lower bounds for hyperparameters
 - Example: proving a problem cannot be solved with fewer than f(n) layers for input size n



WHAT IS THE PROBLEM?

What is the problem?

- Authors want to analyze expressivity of message-passing GNNs (MPGNNs), particularly ones with node IDs
- Goals:
 - Formalize what problems MPGNNs can compute
 - Analyze what MPGNNs cannot compute under restrictions
 - Establish lower bounds for common problems



WHY IS IT IMPORTANT?

Why is it important?

- Formalizing what MPGNNs can compute:
 - Find blind spots of MPGNNs, if any
 - Provide a foundation to analyze limitations
- Formalizing what MPGNNs can't compute:
 - Theoretical lower bounds for hyperparameters for solving problems using GNNs
 - More informed design of MPGNN models



PRIOR WORK

Why don't previous methods address this problem?

- Small body of prior work on limitations in MPGNNs
 - Dehmamy et al. analyzed non-MPGNNs
 - Xu et al. and Morris et al. analyzed MPGNNs without node identification (ie. anonymous) using 1-WL
 - Sato et al. showed that partially-labelled MPGNNs are unable to approximate three NP-hard optimization problems well



Prior Work

- However, adding identifiers to nodes in a MPGNN significantly improves expressivity
- Identifiers can be added without violating permutation invariance/equivariance
- Authors analyze some problems that previous authors have not covered (decision, optimization, graph estimation)
- Depth and width of MPGNN directly connected with graph properties



WHAT IS THE SOLUTION?

- Turns out, MPGNNs are Turing-universal
- Proof sketch:
 - MPGNNs have many similarities to LOCAL, a distributed computing model
 - Differences between MPGNNs and LOCAL:
 - MPGNNs must sum received messages before computing
 - Arguments of messaging function are different
 - Information representation is different



- However, the authors have proven that despite these differences, each node's computation in a MPGNN and LOCAL have the same expressivity
 - Given messaging and update functions are Turing-complete
 - For GNNs, this means that functions in each layer should be sufficiently complex



- LOCAL is Turing-complete if number of rounds of the distributed algorithm is larger than the graph diameter if memory is not an issue
- Thus, each GNN node can compute any Turing computable function if:
 - Depth (number of layers) d must be at least as great as the graph diameter
 - Width (largest state of node across all layers) w must be unbounded
- Since this result considers computation per node, the node must be uniquely identifiable for this result to hold



- Sufficient conditions for universality:
 - Uniquely identifiable node
 - Messaging and update functions must be sufficiently complex
 - Depth must be at least as large as graph diameter
 - Width must be unbounded



- What happens if sufficient conditions for universality are relaxed?
- Analyze problems under constraints of depth and width of GNN
- It turns out restrictions on capacity significantly limit expressivity of MPGNNs
- Authors analyze relaxation of unique identification and depth/width conditions



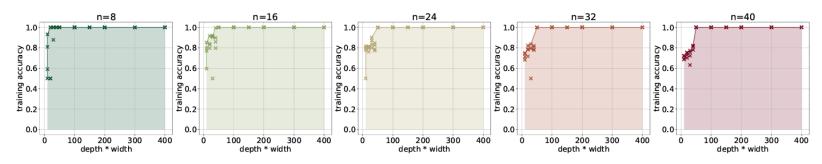
problem	bound	problem	bound
cycle detection (odd)	$dw = \Omega(n/\log n)$	shortest path	$d\sqrt{w} = \Omega(\sqrt{n}/\log n)$
cycle detection (even)	$dw = \Omega(\sqrt{n}/\log n)$	max. indep. set	$dw = \Omega(n^2/\log^2 n)$ for $w = O(1)$
subgraph verification*	$d\sqrt{w} = \Omega(\sqrt{n}/\log n)$	min. vertex cover	$dw = \Omega(n^2/\log^2 n)$ for $w = O(1)$
min. spanning tree	$d\sqrt{w} = \Omega(\sqrt{n}/\log n)$	perfect coloring	$dw = \Omega(n^2/\log^2 n)$ for $w = O(1)$
min. cut	$d\sqrt{w} = \Omega(\sqrt{n}/\log n)$	girth 2-approx.	$dw = \Omega(\sqrt{n}/\log n)$
diam. computation	$dw = \Omega(n/\log n)$	diam. ³ /2-approx.	$dw = \Omega(\sqrt{n}/\log n)$

• where d=depth, w=width, n=# of nodes in the graph

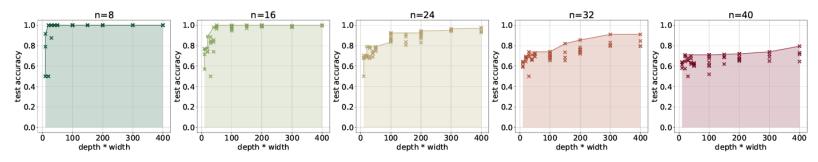
- Proof sketch:
 - Show equivalence of expressivity limits between MPGNN and CONGEST, a variant of LOCAL
 - For each problem, find equivalent in CONGEST
 - Established limits of expressivity of CONGEST can be translated into limits of MPGNN



- Problem: 4-cycle classification
- Empirical results with constrained capacity dw



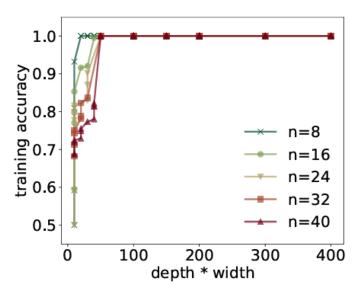
(a) training accuracy of all trained networks



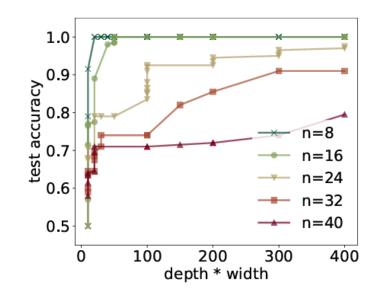
(b) test accuracy of all trained networks



Empirical results with constrained capacity dw



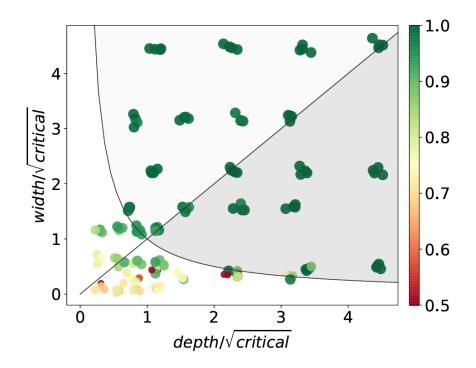
(c) best training accuracy



(d) best test accuracy

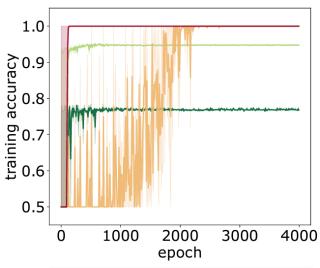


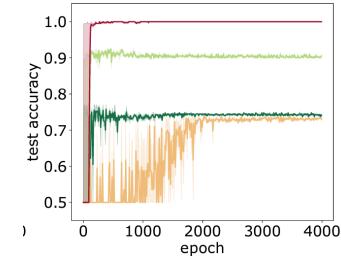
• Empirical results with independently constrained depth d and width w





- Empirical results with various levels of anonymity:
 - Anonymous: no graph IDs
 - Degree: ID is node degree
 - Random unique ID: Inconsistent node ID between graphs
 - Unique ID: consistent node ID between graphs
 - ----- anonymous ----- degree
 - random unique id unique id







Limitations of results

- Analyzed lower bounds are worst-case
 - One impossible graph is enough to prove lower bound
 - Does not mean they are all impossible
- Lower bounds were found with the assumption of universal layers and nodes with discriminative attributes
- However, lower bound results will still be applicable for graphs with less expressivity (e.g. computationally limited layers, anonymous graphs)



FUTURE DIRECTIONS

Future Directions

- Non-worst case bounds?
 - How tight are the derived bounds?
- Loosening assumptions on complexity of message and update functions

