

# Teoria dos Grafos

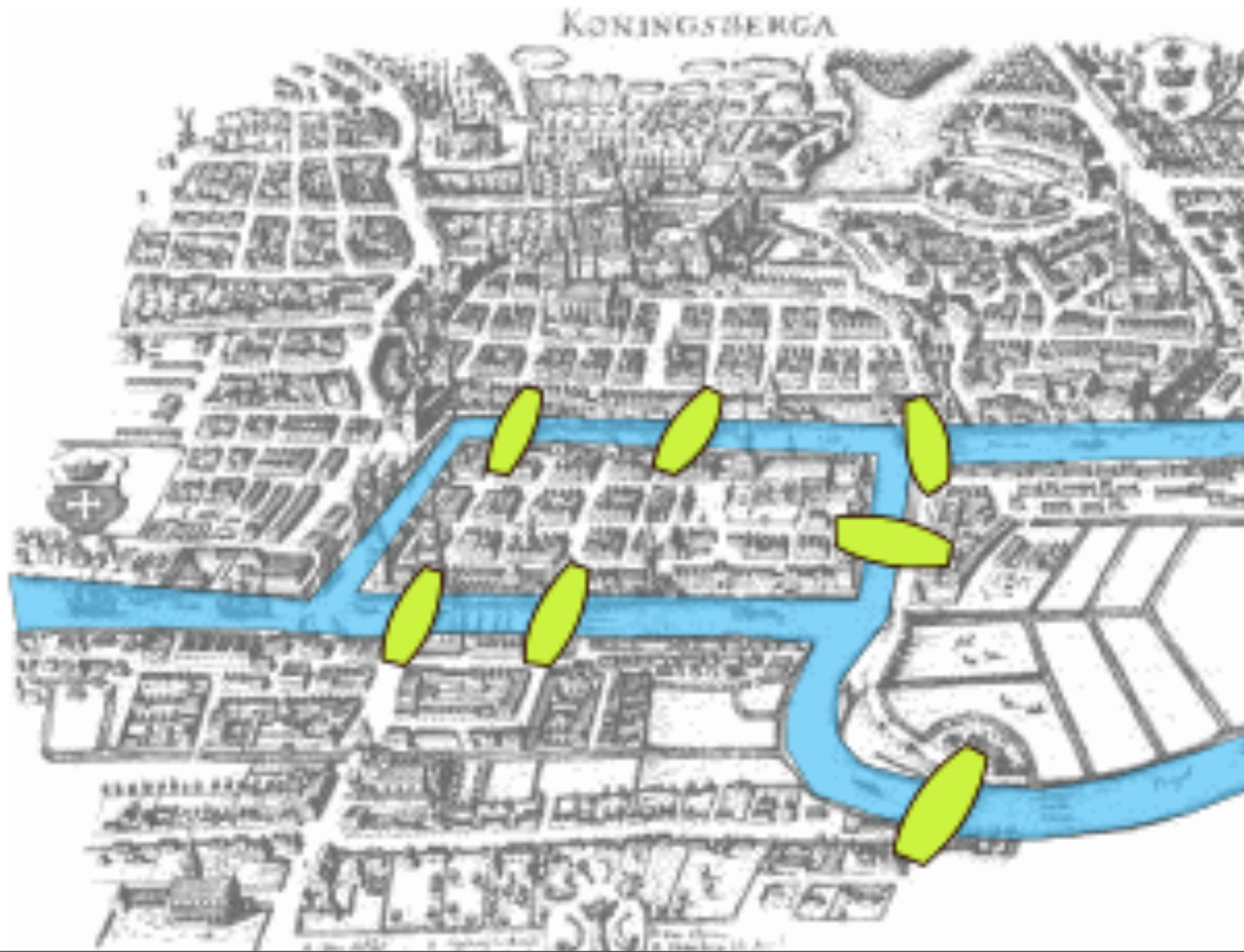
Prof. Leandro G. M. Alvim

# Agenda

- História
- Motivação
- Definições

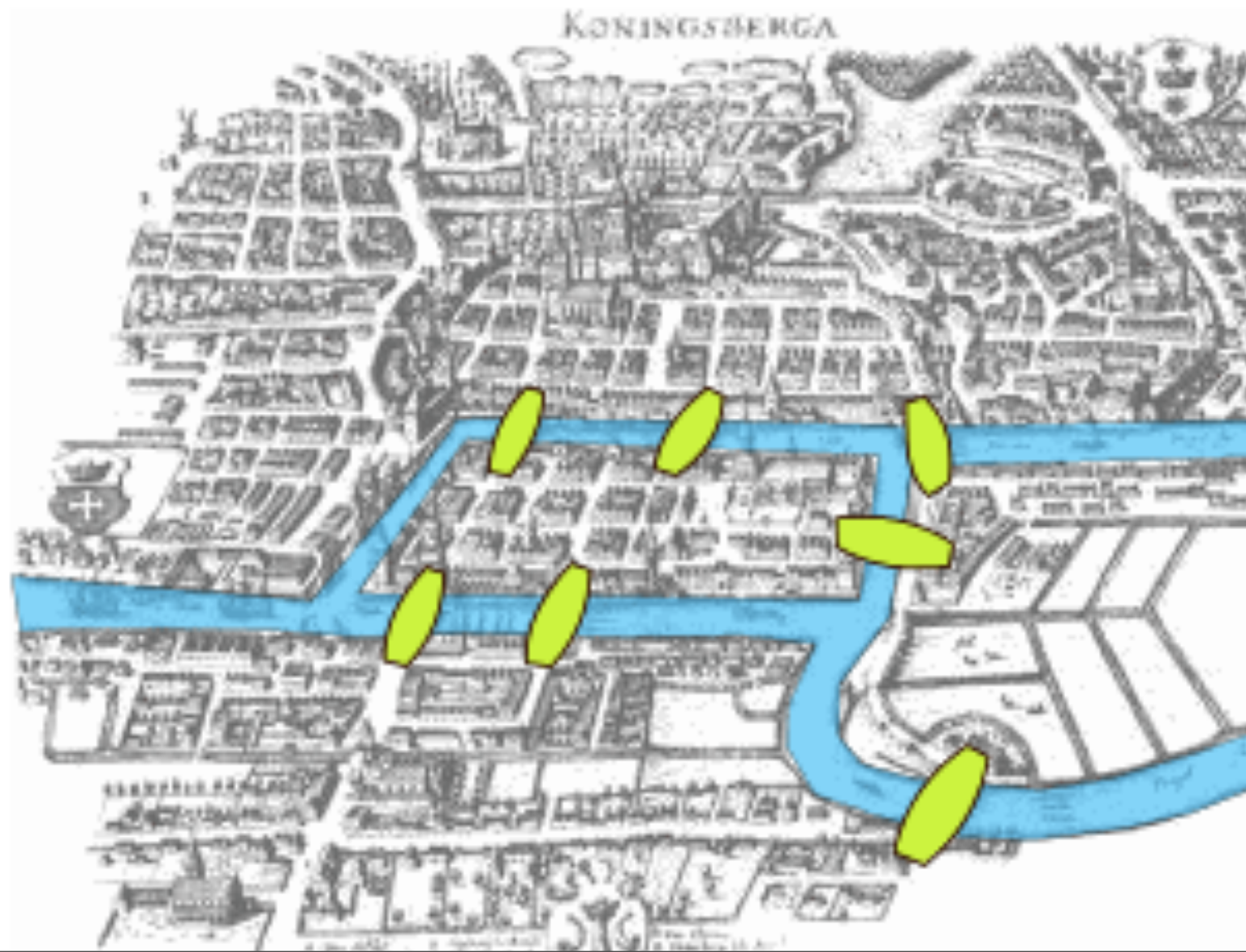
# As pontes de Königsberg

- Königsberg, Prússia (até 1945) - Kaliningrado, Rússia (Rússia)



# As pontes de Königsberg

- É possível caminhar pela cidade passando por todas as pontes uma única vez?





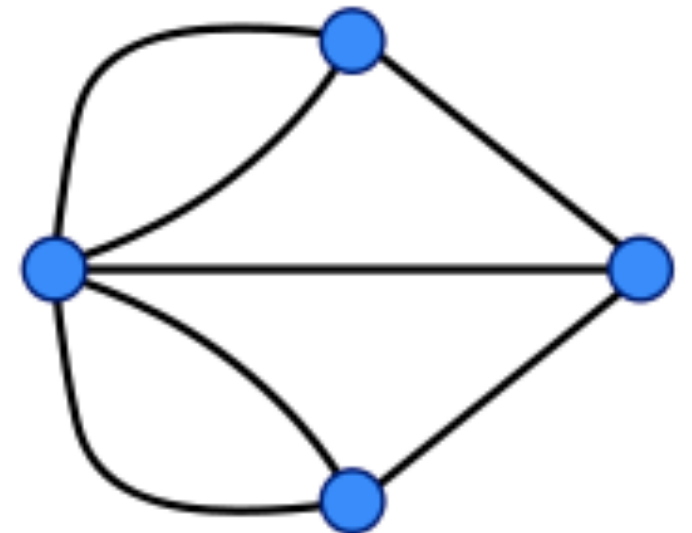
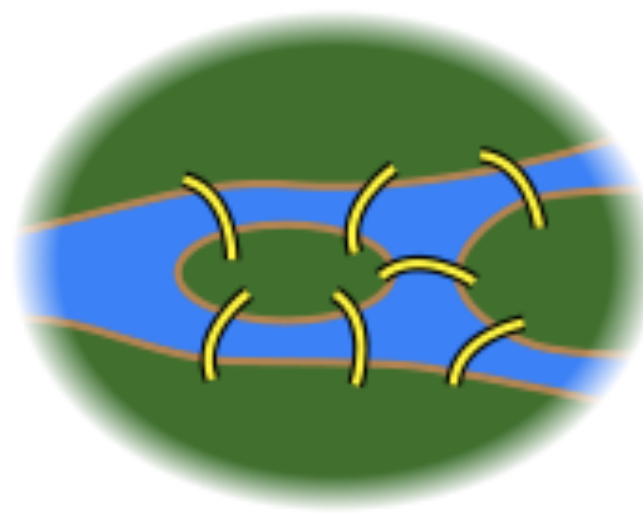
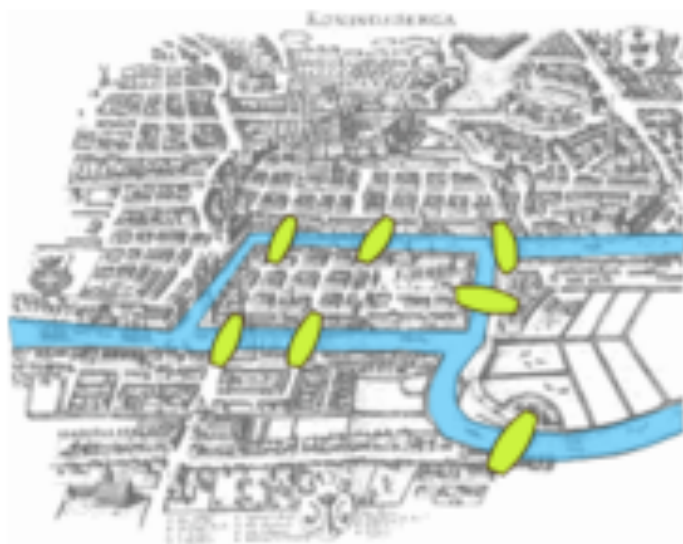
# As pontes de Königsberg

- Leonard Euler
  - Primeiro trabalho da Teoria dos Grafos (1735)
  - Curiosidades
    - 13 filhos
    - Cegueira



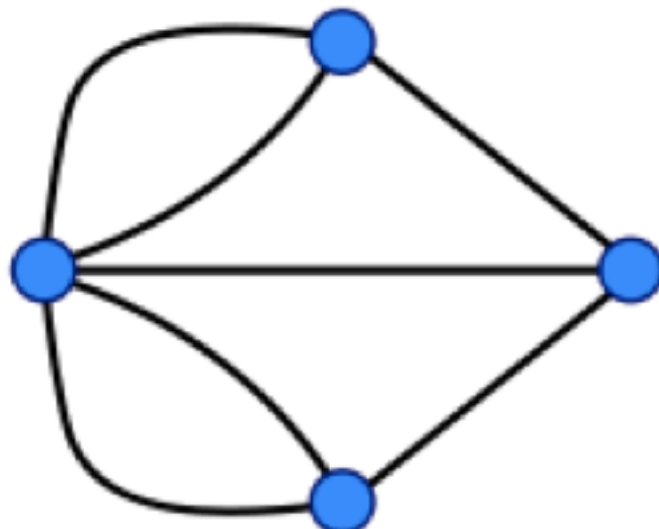
# As pontes de Königsberg

- Análise de Euler
- Primeiro passo



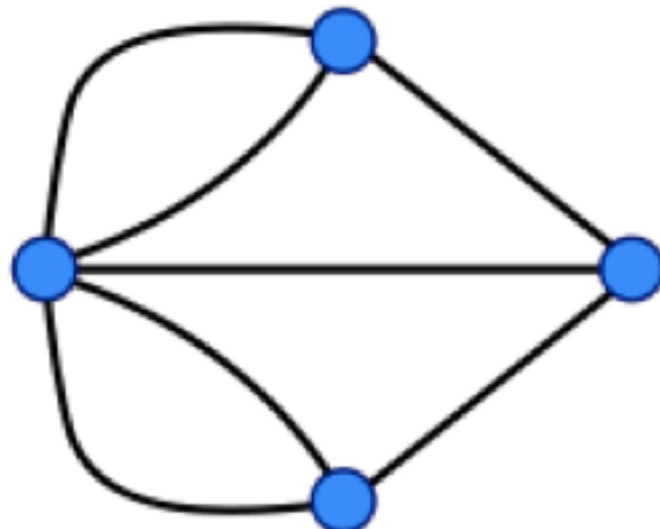
# As pontes de Königsberg

- Análise de Euler
- Segundo passo
- Para que tenha solução, devemos ter um número par de pontes em cada terreno



# As pontes de Königsberg

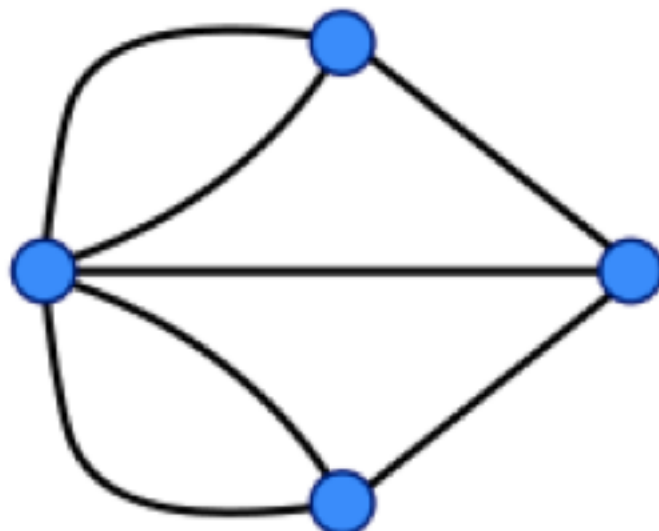
- Análise de Euler
  - Não há solução para este exemplo!





# As pontes de Königsberg

- Análise de Euler
  - Legado
    - Vértices, arestas, grau, Caminho Euleriano, Circuito Euleriano
  - Primeiro Teorema



# História

- Biggs, N.; Lloyd, E. and Wilson, R. (1986), *Graph Theory, 1736-1936*, Oxford University Press

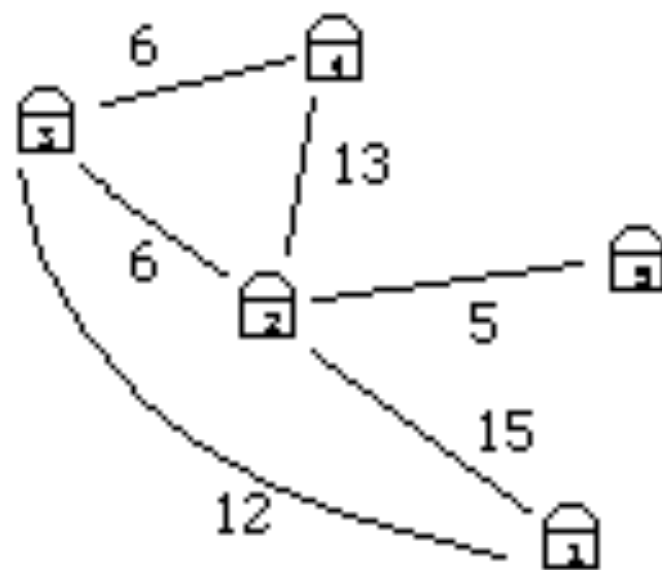
# Problema das cores



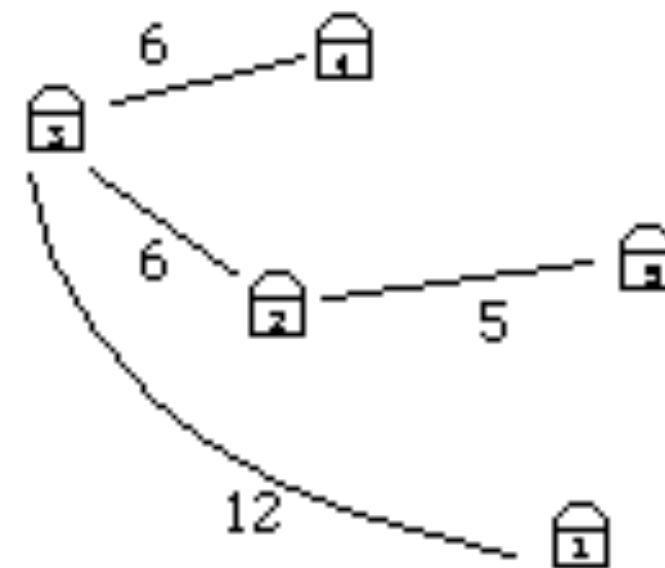
# Caminho Mínimo



# Cabeamento



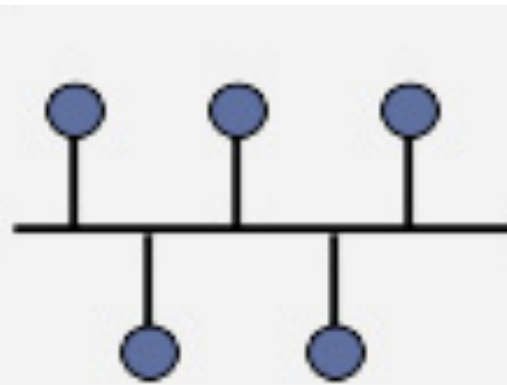
Ramos de rede possíveis com  
custo ambiental associado



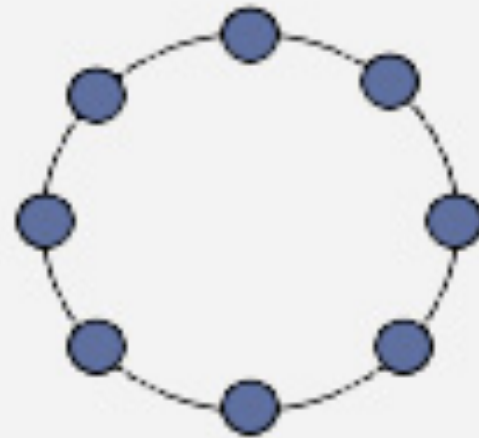
Interligação das tabas com  
menor custo ambiental



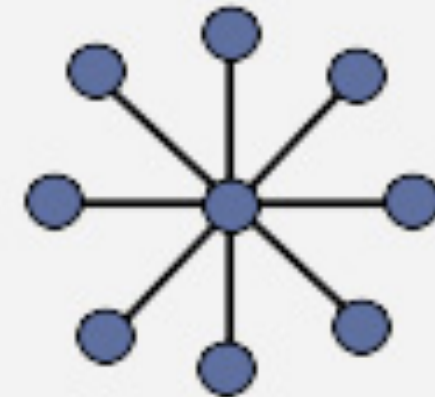
# Topologia de redes



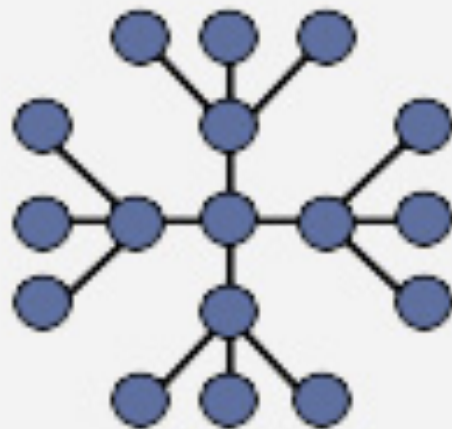
**Bus**



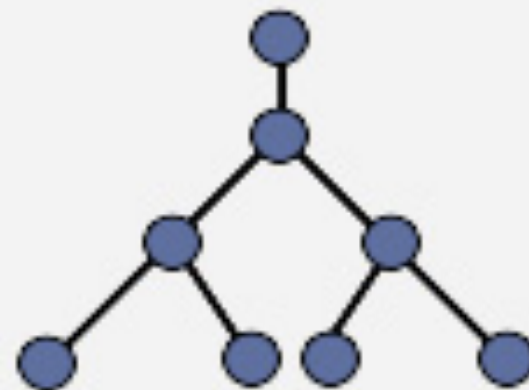
**Ring**



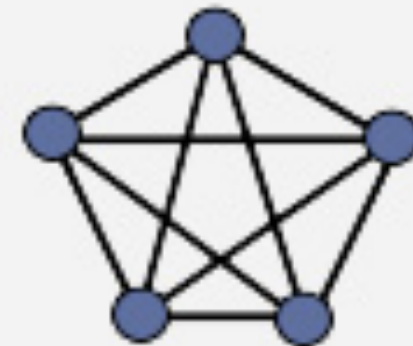
**Star**



**Extended Star**



**Hierarchical**



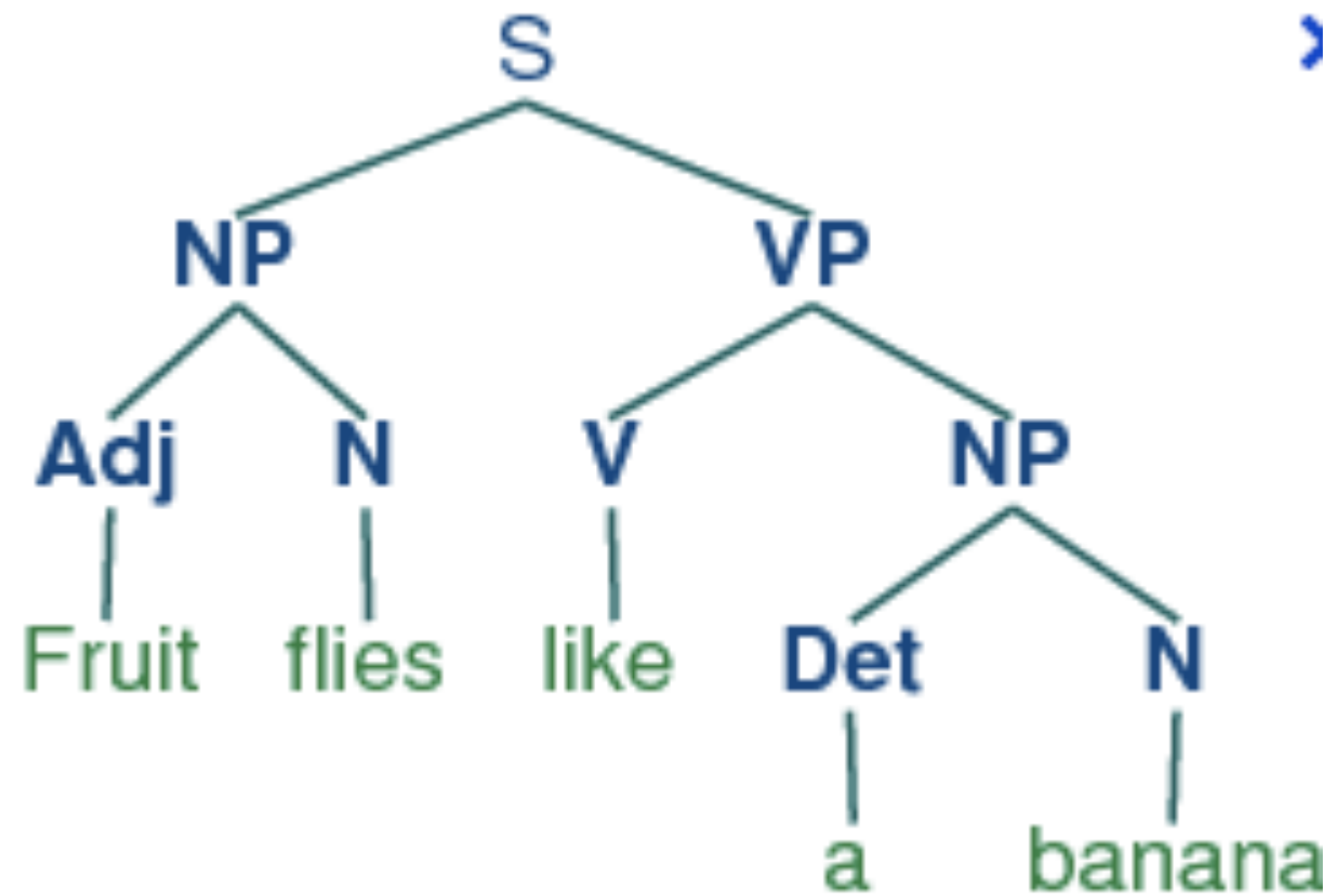
**Mesh**



# Rede Social

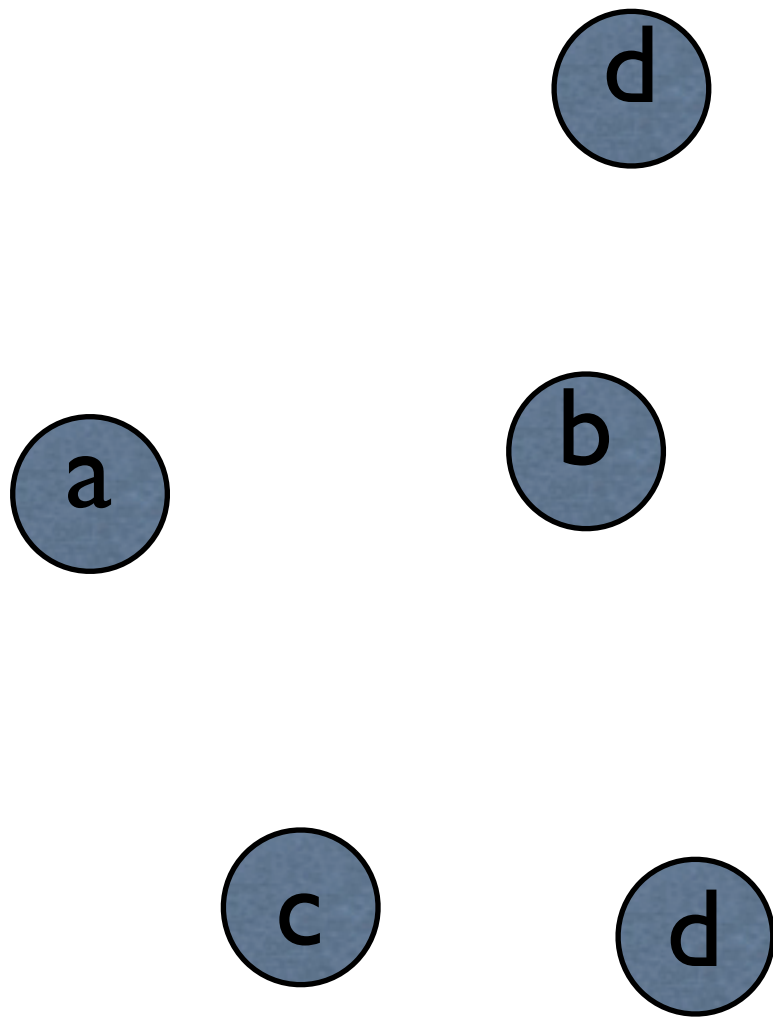


# Texto



# Definições

- Grafo Vazio

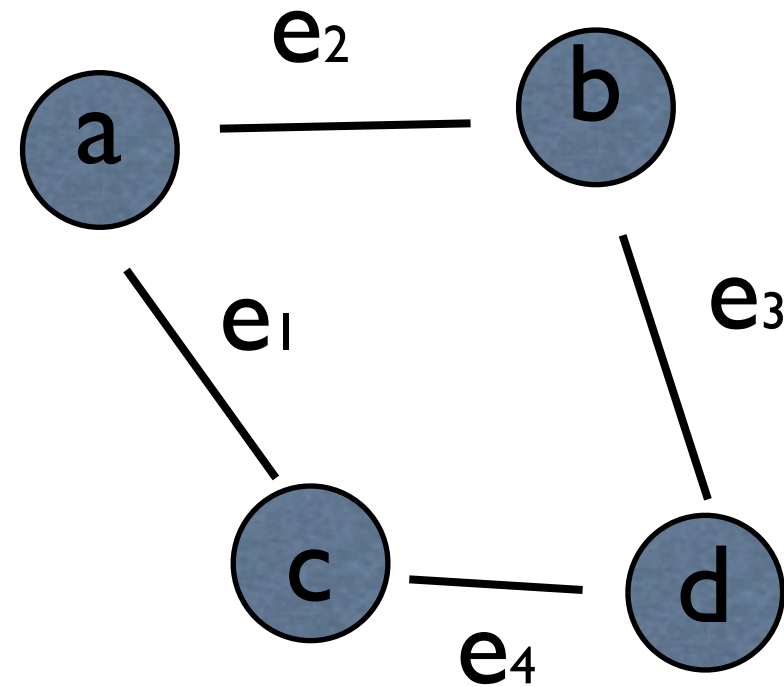


# Definições

- Grafo Nulo

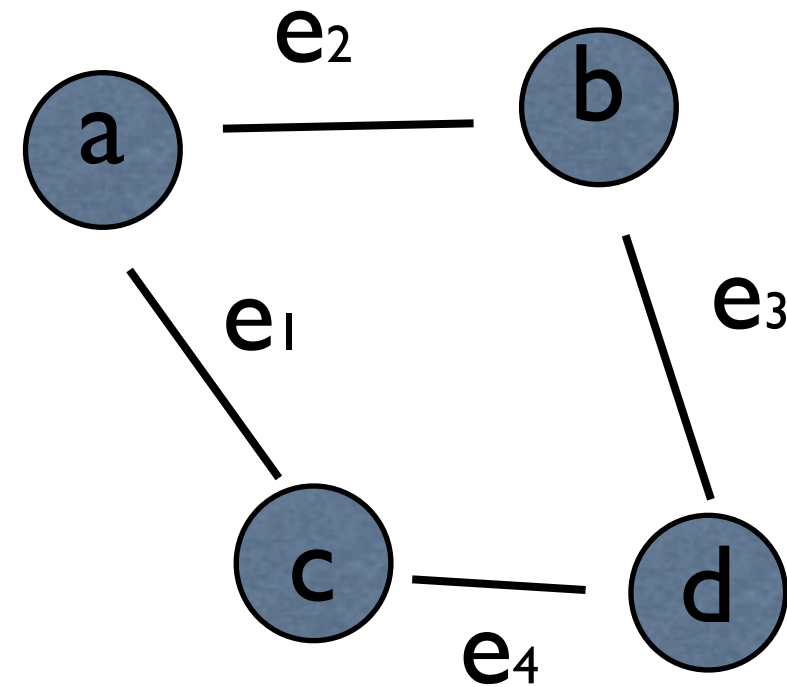
# Definições

- Grafo
  - $G = (V, E)$
  - $V = \{a, b, c, d\}$ ,  $|V| = 4$
  - $E = \{e_1, e_2, e_3, e_4\}$ ,  $|E| = 4$



# Definições

- Arestas
  - $e_1 = (a, c)$ 
    - não ordenado
    - $e_1$  incide em a e c





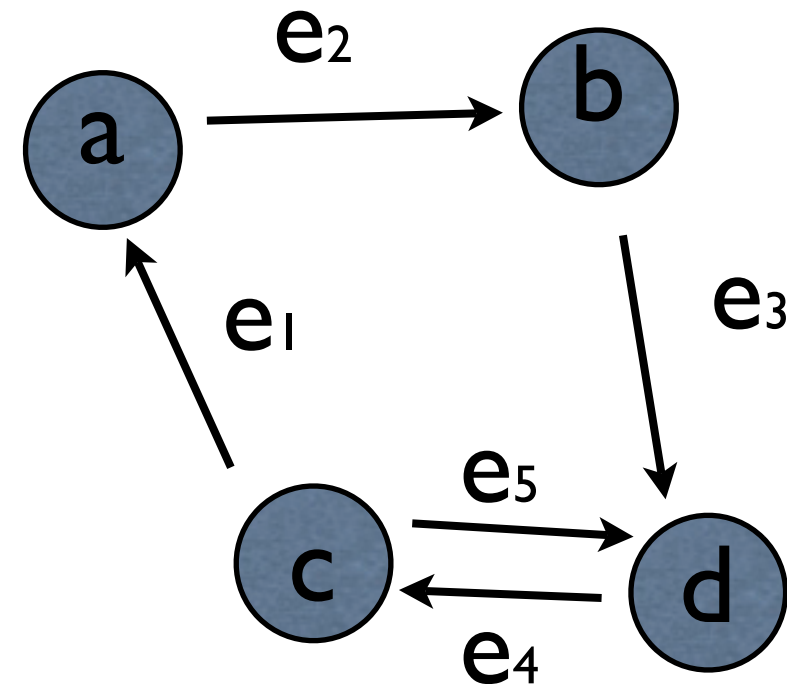
# Definições

- Grafo orientado (Digrafo)

- $G = (V, E)$

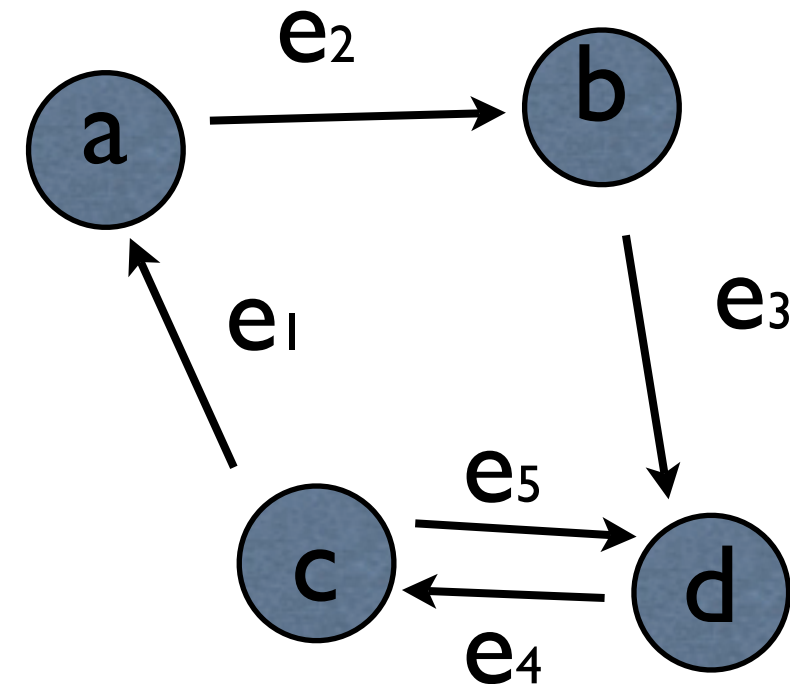
- $V = \{a, b, c, d\}, |V| = 4$

- $E = \{e_1, e_2, e_3, e_4, e_5\}, |E| = 5$



# Definições

- Arestas
  - $e_4 = (d,c) \Leftrightarrow e_5 = (c,d)$
  - $e_4$  incide em  $c$
  - $e_5$  incide em  $d$

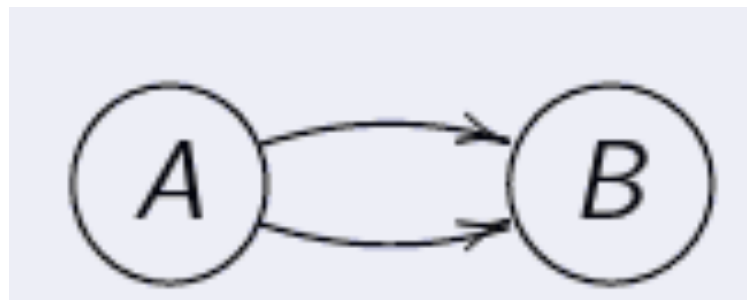


# Definições

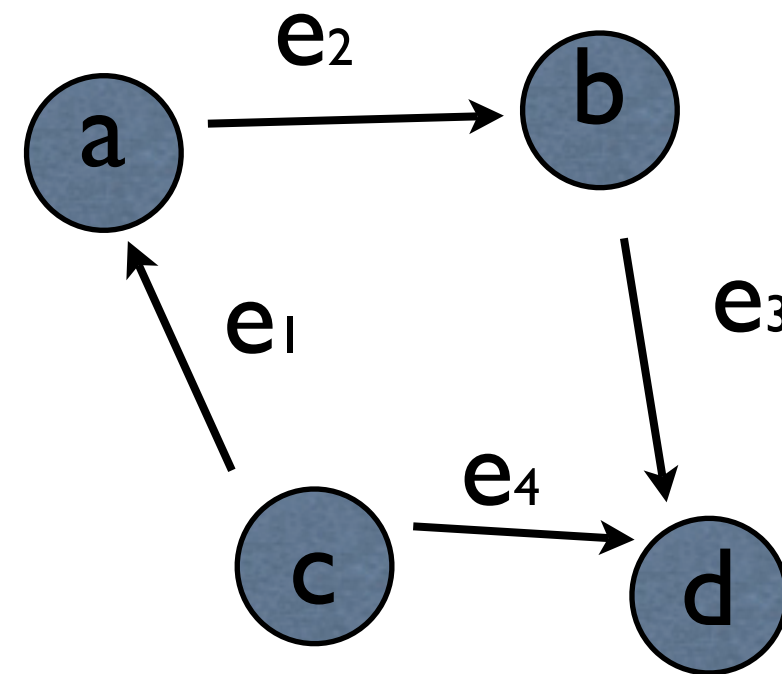
- Laço



- Aresta paralela

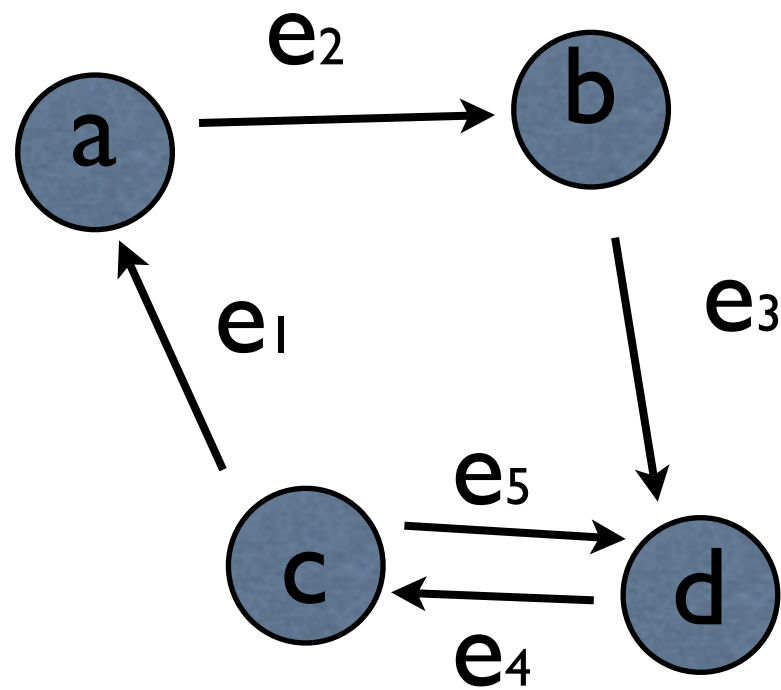


- Grafo Simples



# Definições

- Adjacência

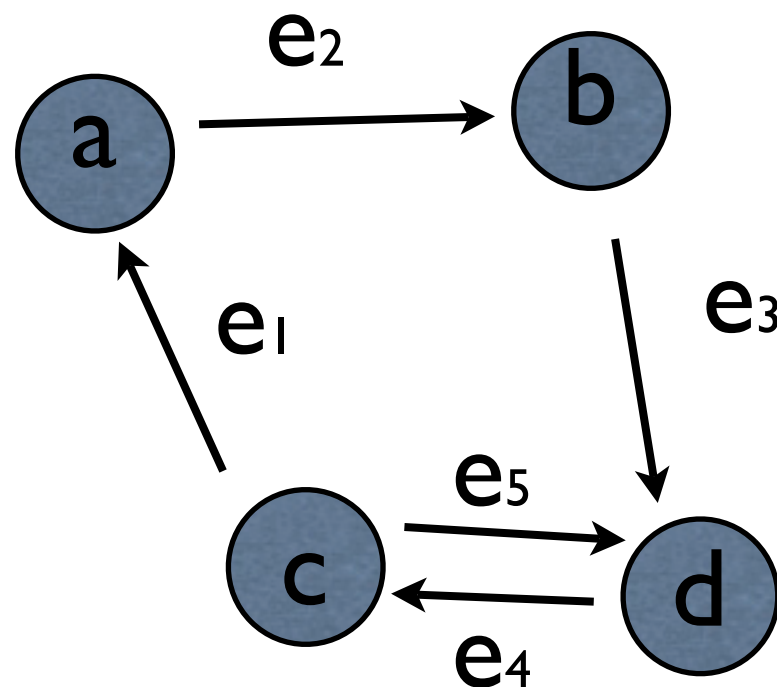


# Definições

- Grau

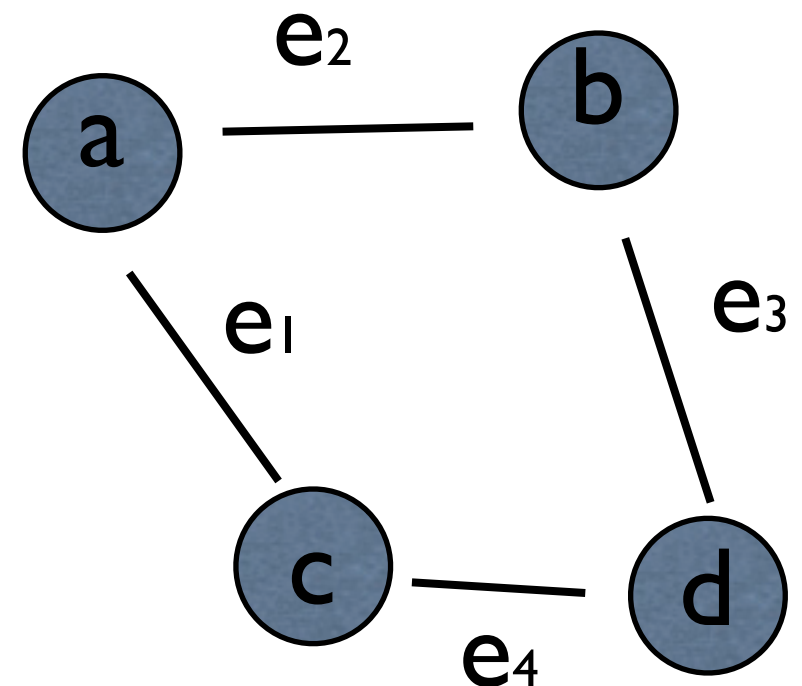
- $d^+(d) = 1$

- $d^-(d) = 2$



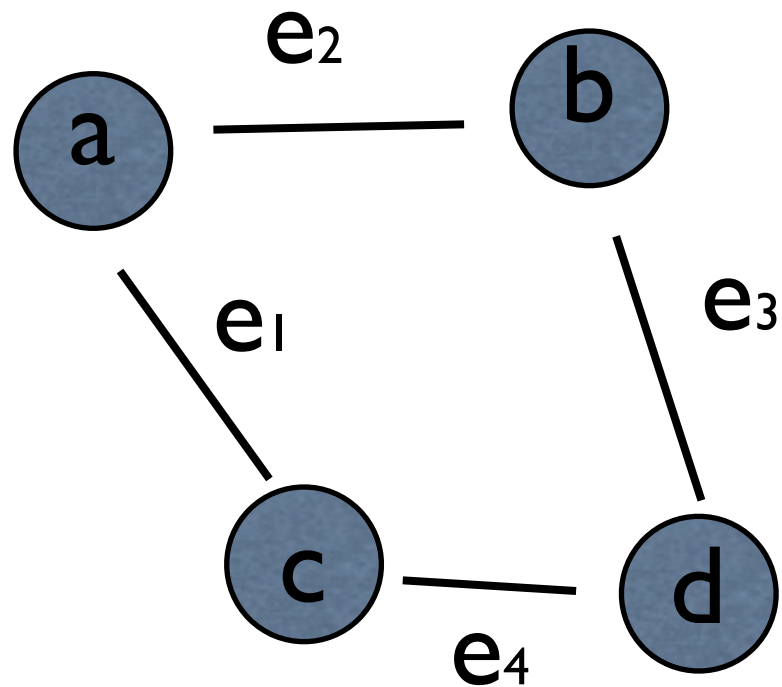
- Grau

- $d(d) = 2$



# Definições

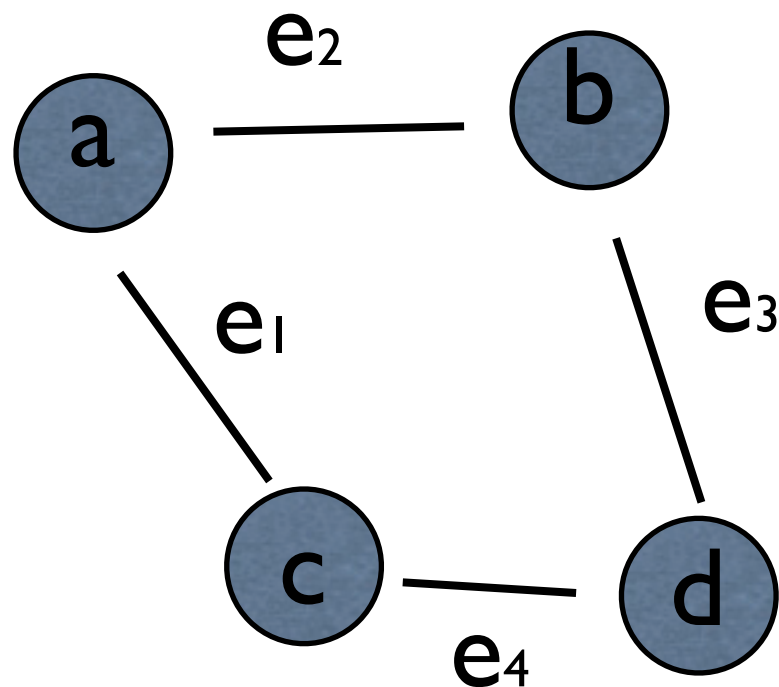
- Grau
- $\text{Sum}(d(v)) = 2m$   
(demostre)





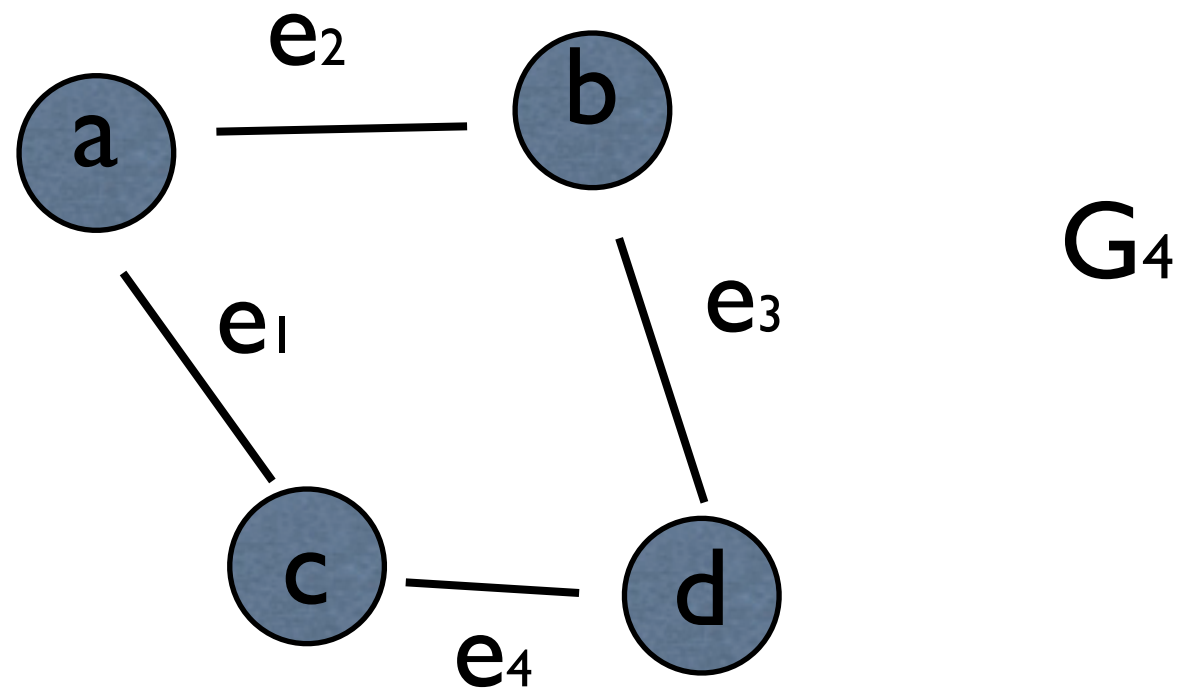
# Definições

- Grau
- O número de vértices de grau ímpar de um grafo sempre é par (demonstre)



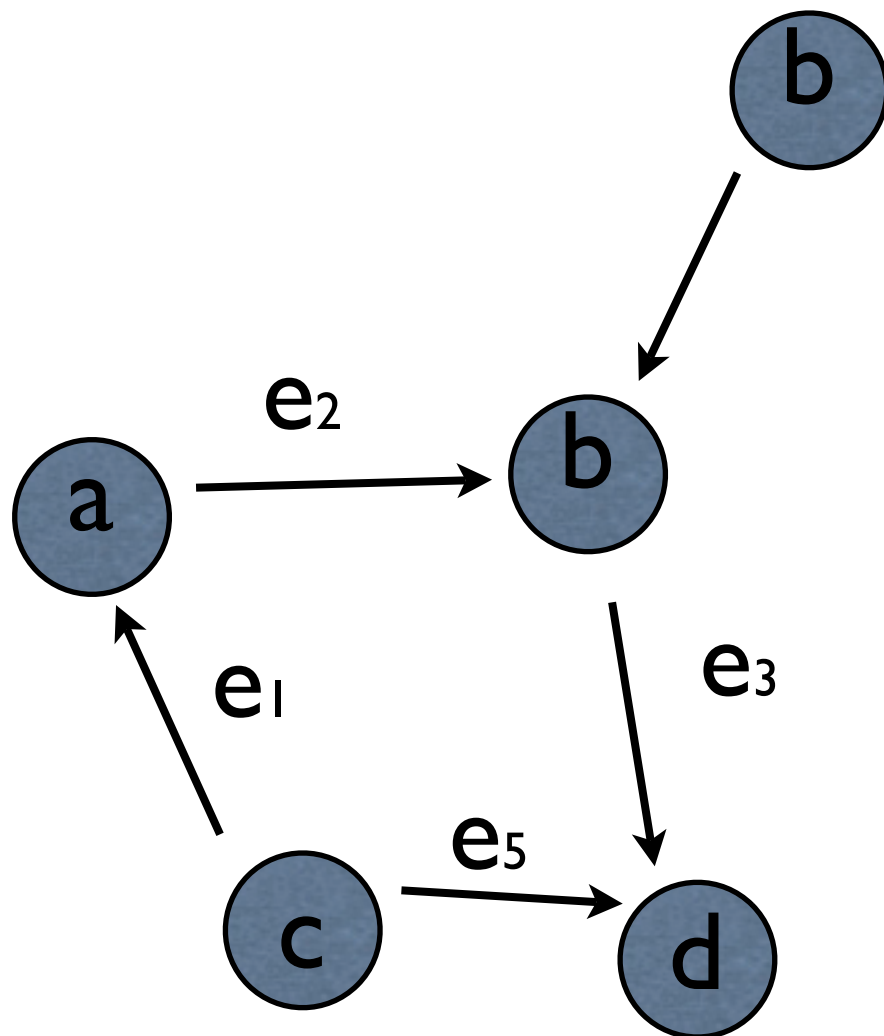
# Definições

- Ordem
  - O número de vértices de  $G$



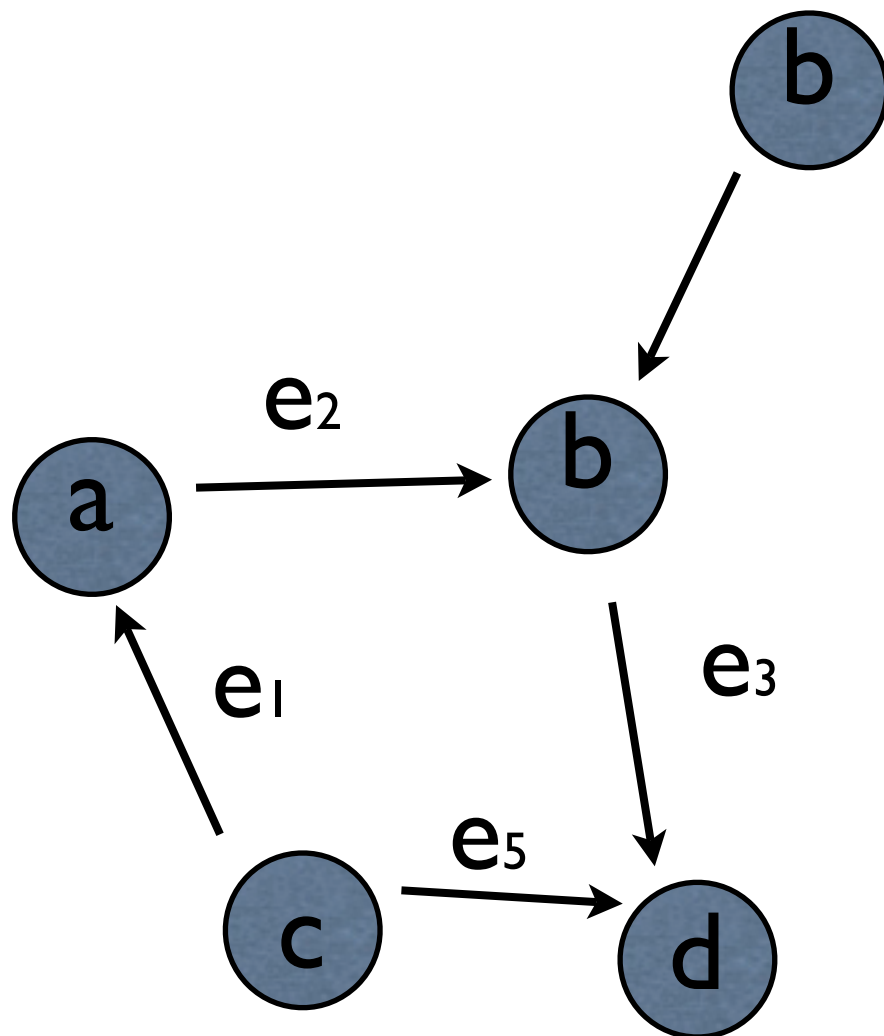
# Definições

- Fonte
  - vértice:  $d(v) = 0$



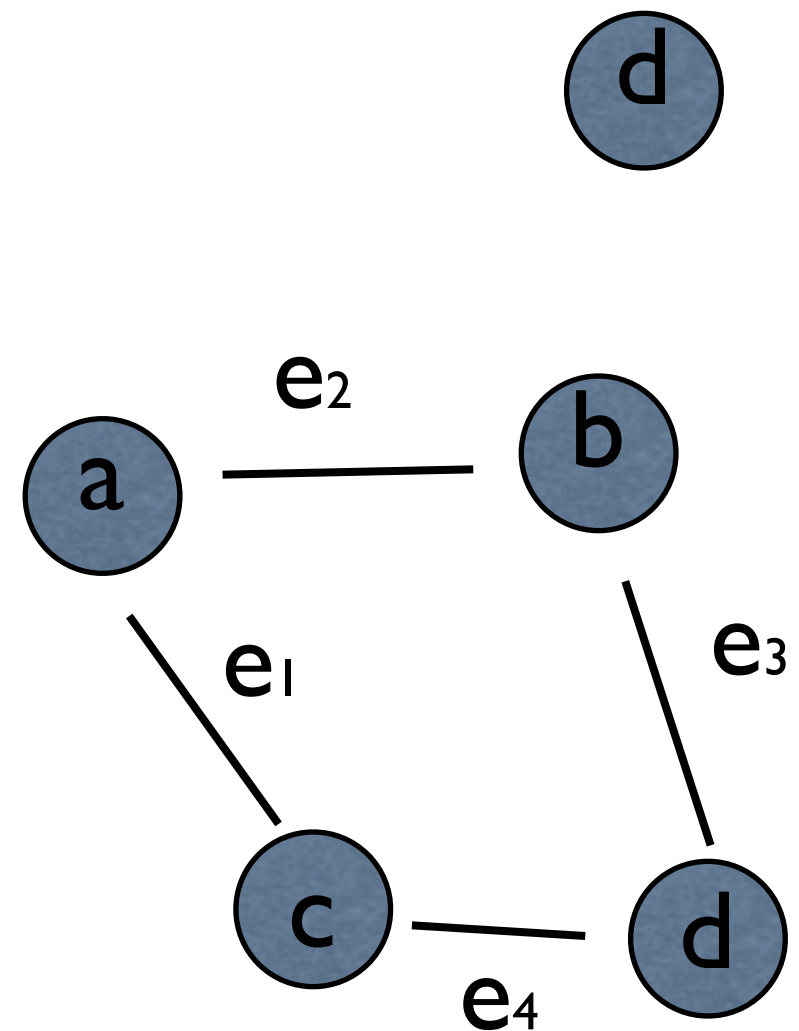
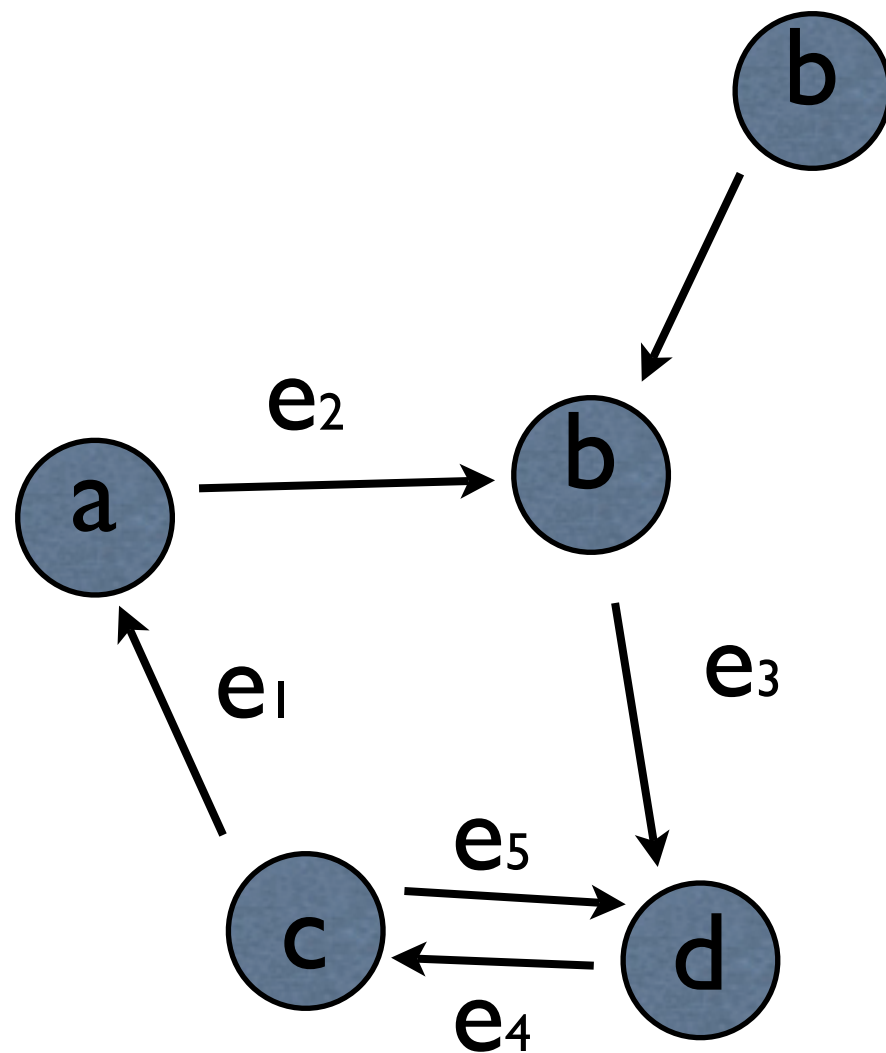
# Definições

- Sumidouro
  - vértice:  $d^+(v) = 0$



# Definições

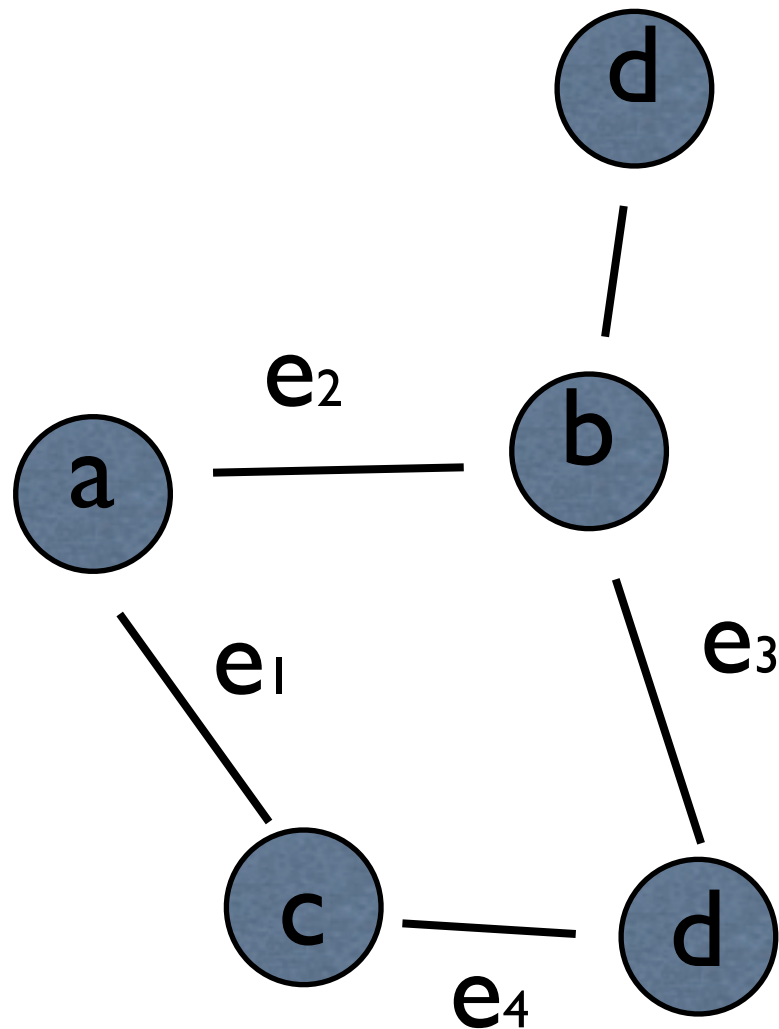
- Vértice Isolado
- nenhuma aresta incidente



# Definições

- Vértice Pendente

- $d(v) = 1$

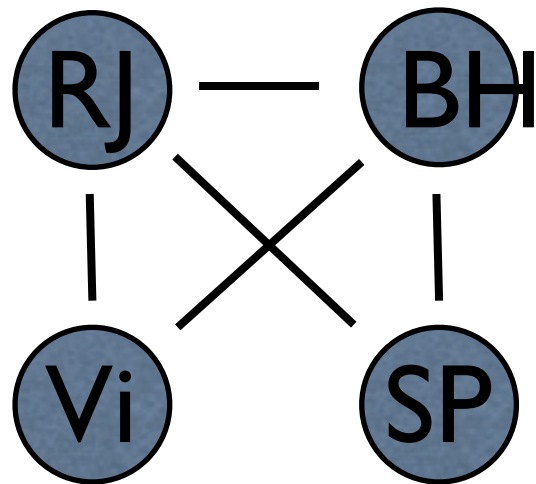




# Definições

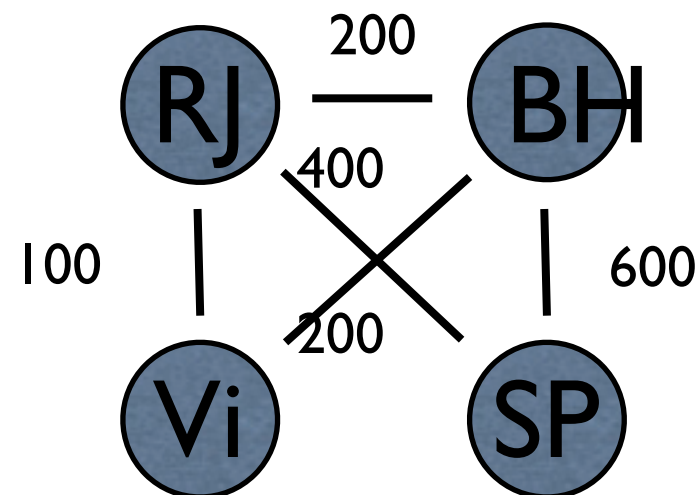
- Grafo Rotulado

- $V = \{v | v \text{ é uma capital}\}$



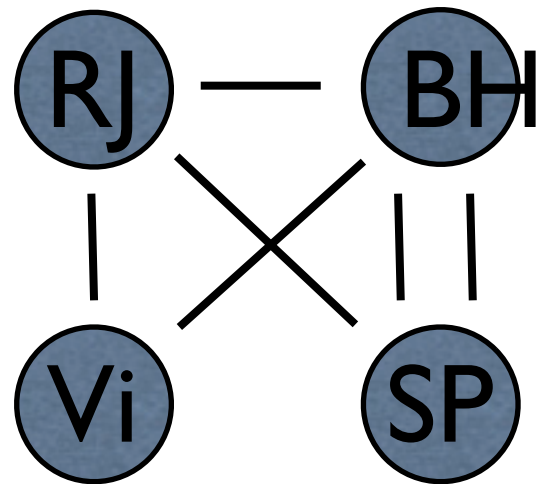
- Grafo Valorado

- $V = \{v | v \text{ é uma capital}\}$
- $E = \{(v1, v2, d) | \text{há uma vizinhança tal que a distância é } d \text{ km}\}$

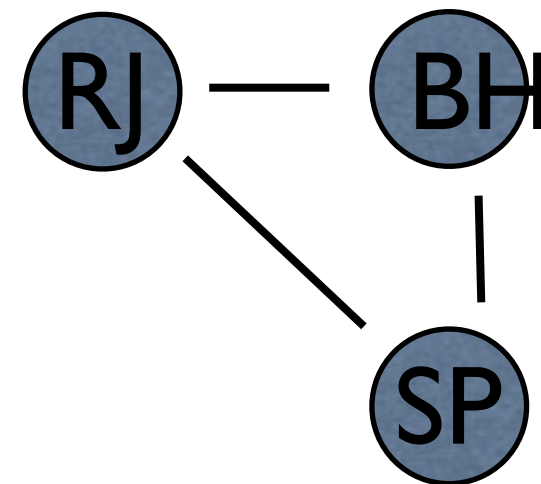


# Definições

- Multigrafo



- Subgrafo



# Definições

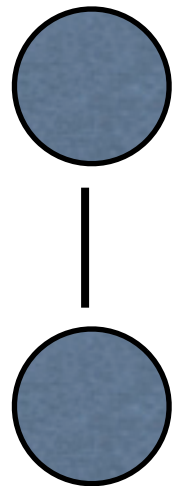
- Grafo Completo ( $k$  - regular)

G1



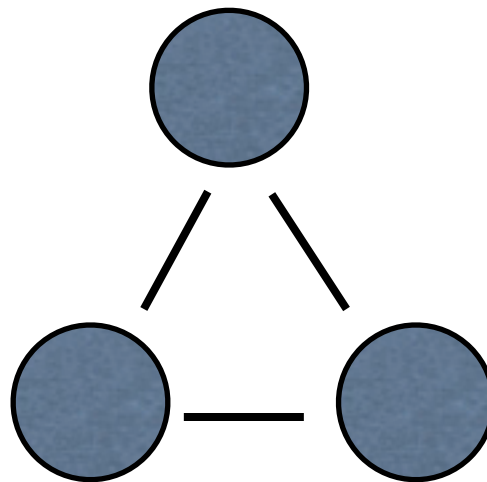
$k=0$

G2



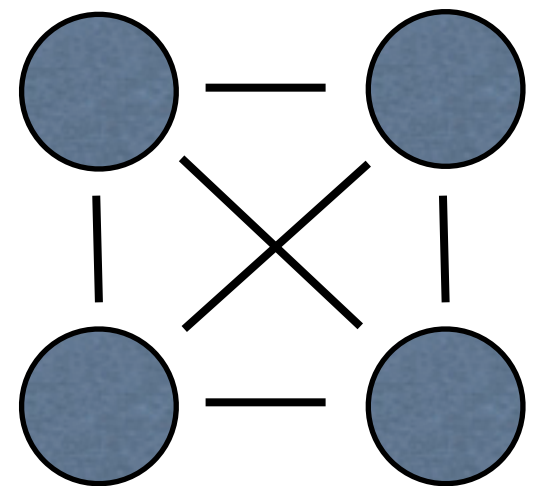
$k=1$

G3



$k=2$

G4



$k=3$

# Definições

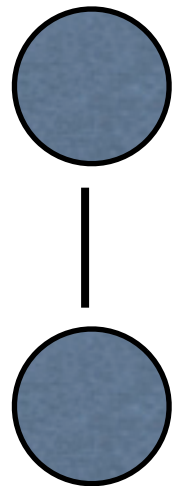
- $|E| = n(n-1)/2$  (demonstre)

G1



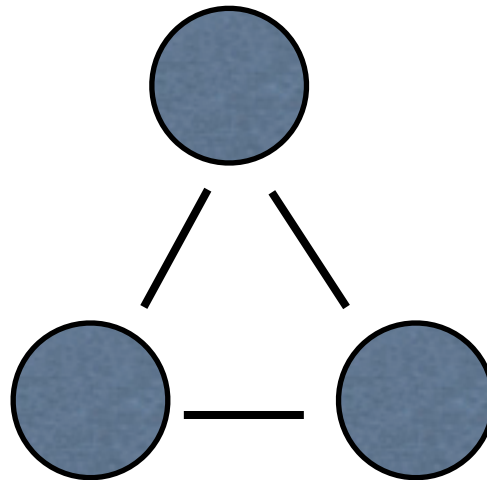
$k=0$

G2



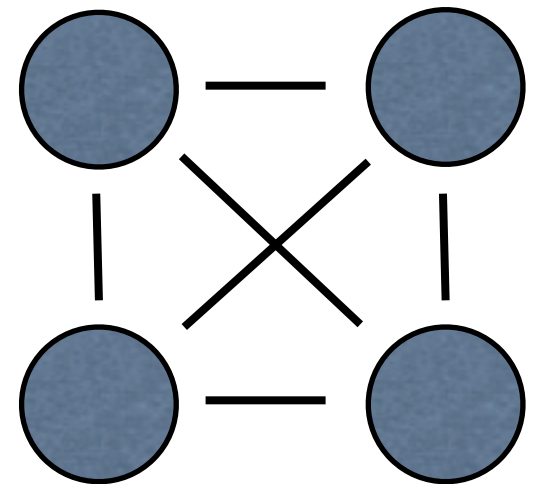
$k=1$

G3



$k=2$

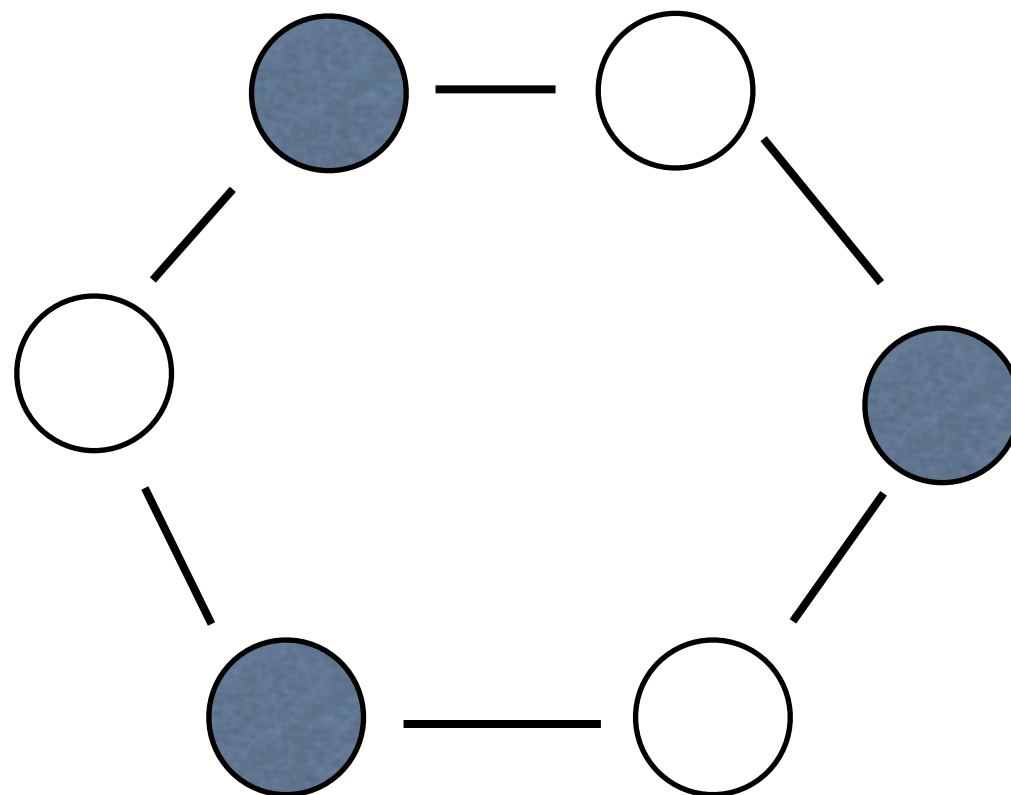
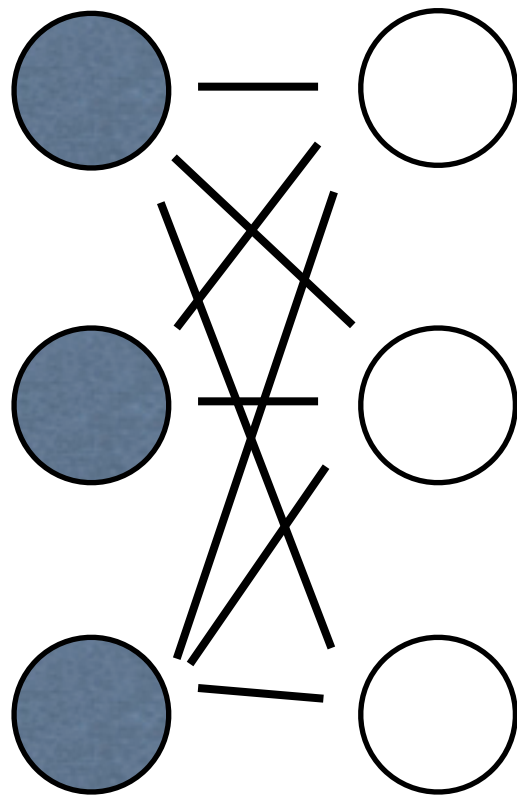
G4



$k=3$

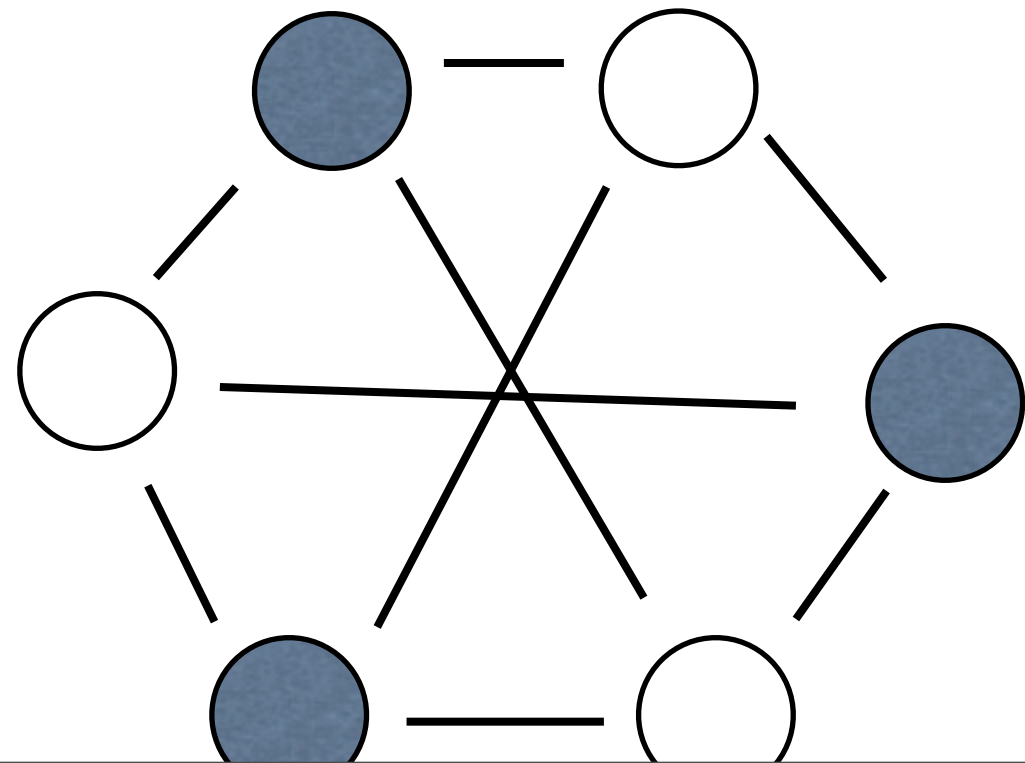
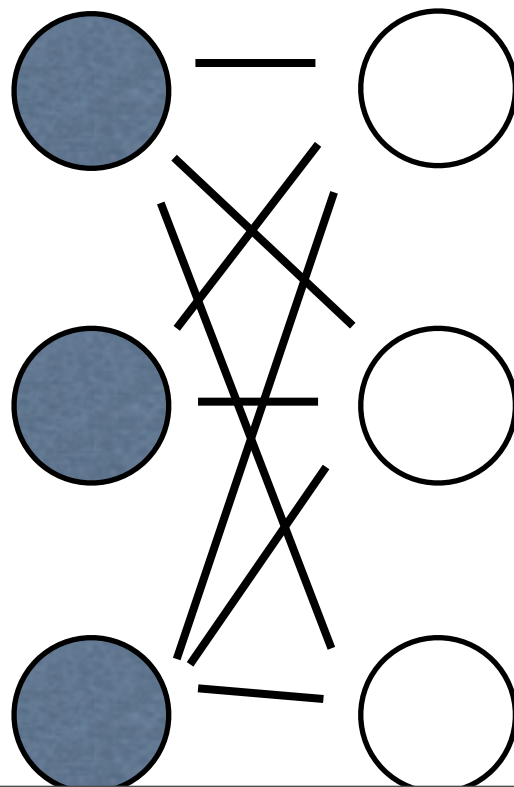
# Definições

- Grafo Bipartido
  - Se existem duas partições  $V_1$  e  $V_2$  de forma que qualquer aresta pertença a  $V_1$  e  $V_2$



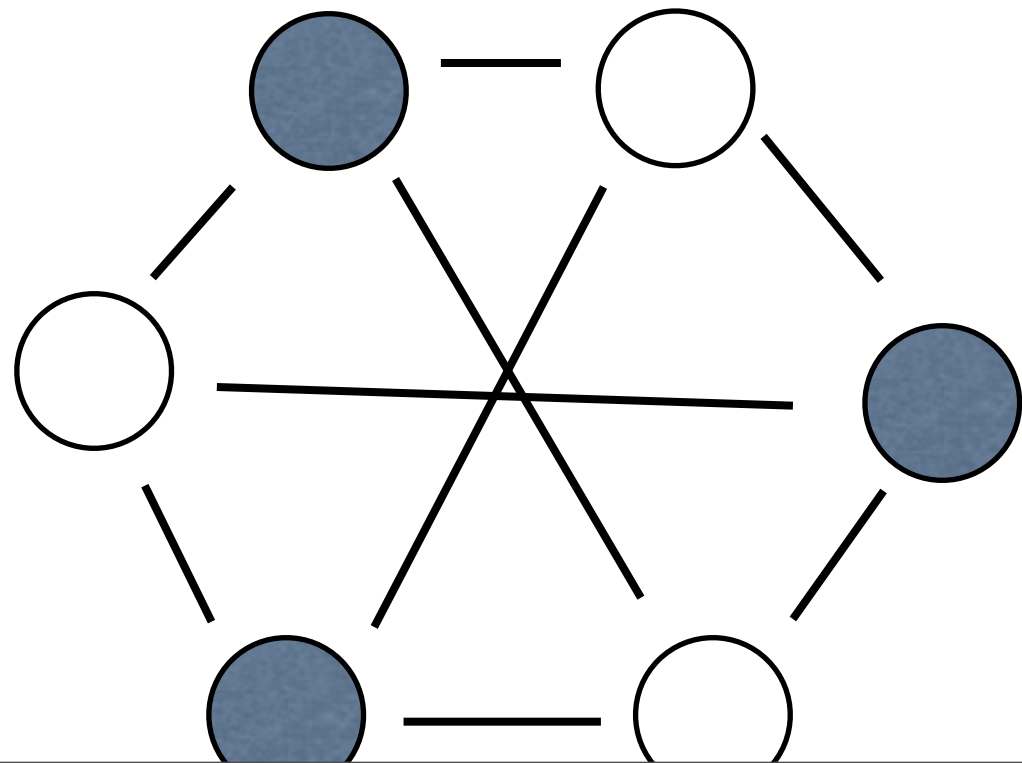
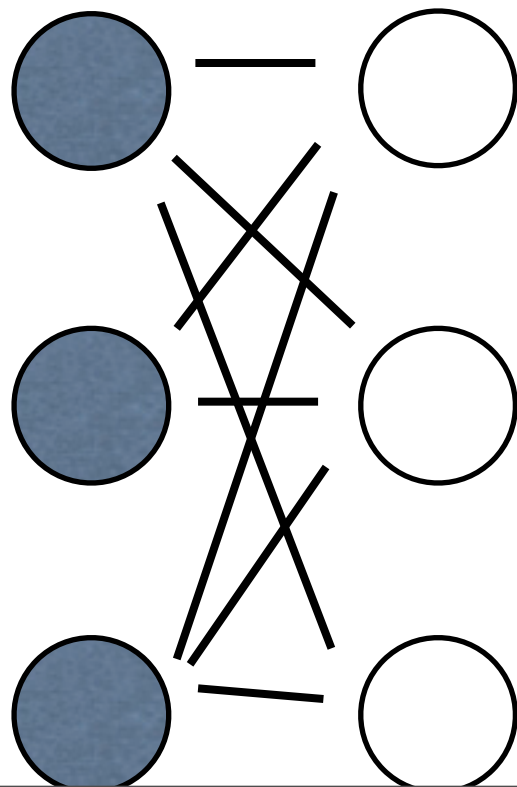
# Definições

- Grafo Bipartido Completo
- Se existem duas partições  $V_1$  e  $V_2$  de forma que qualquer aresta pertença a  $V_1$  e  $V_2$
- Todo vértice de  $V_1$  é adjacente a todo vértice de  $V_2$



# Definições

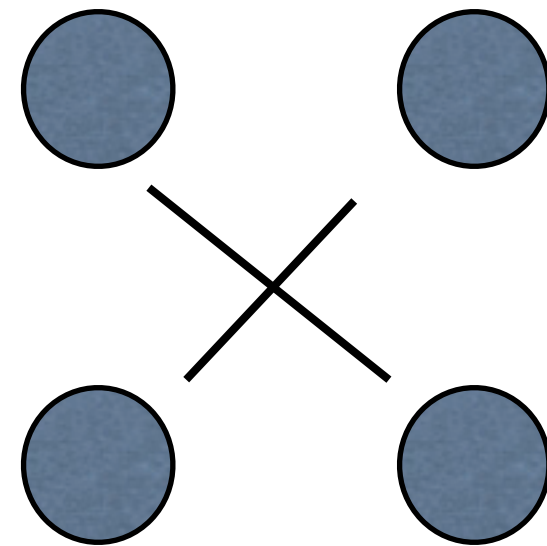
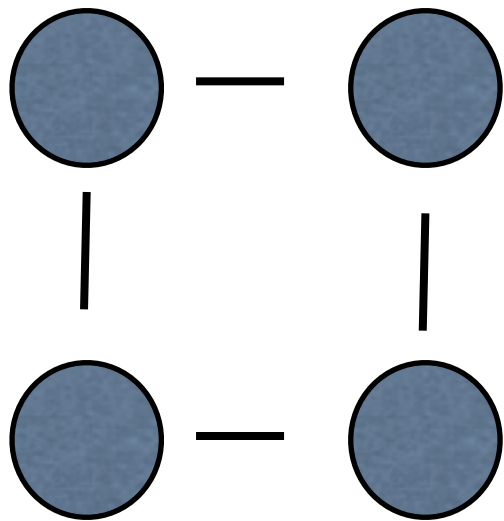
- Grafo Bipartido Completo
- $|E| = n \times m$  (demonstre)





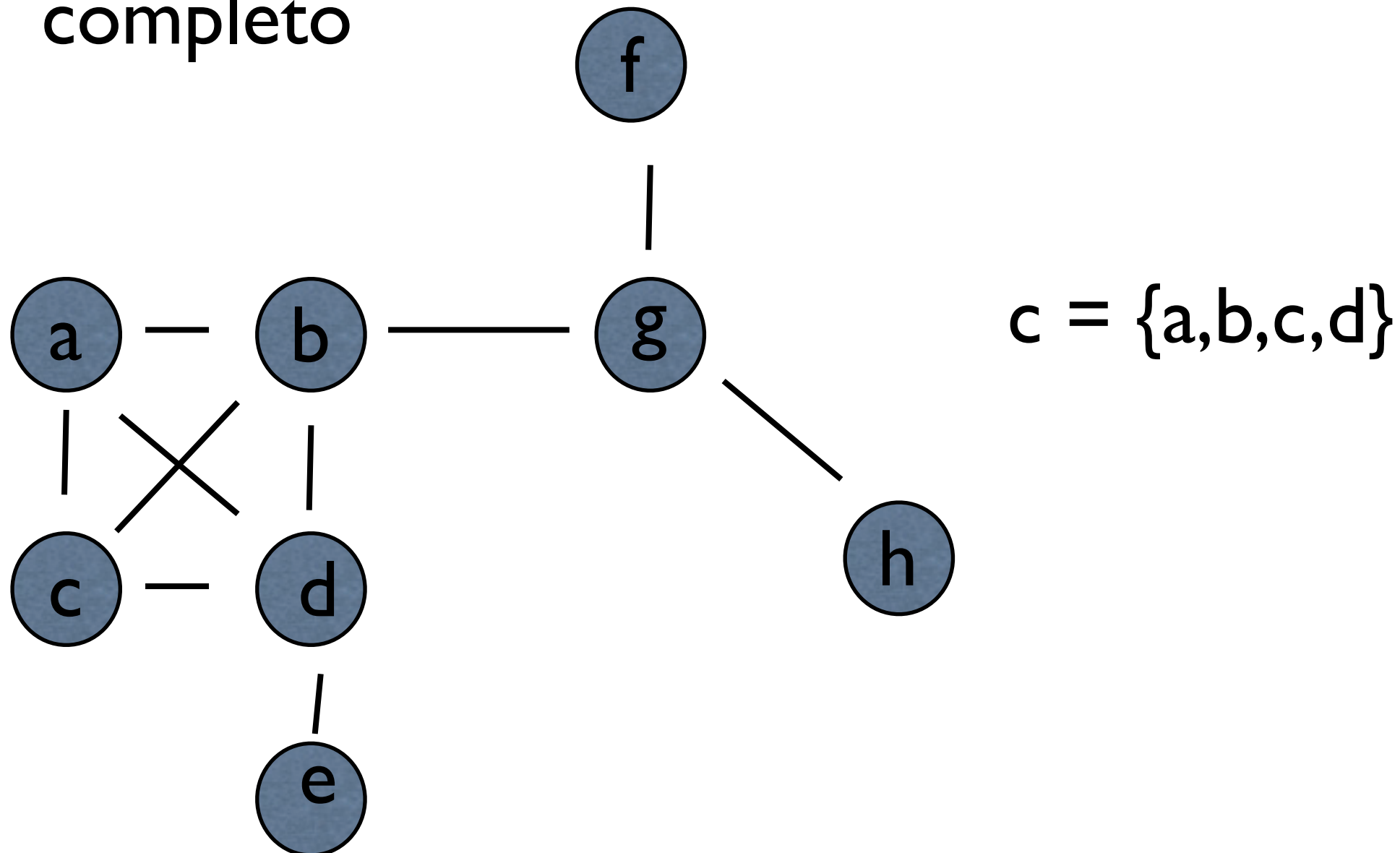
# Definições

- Grafos Complementares
  - Arestas que faltam para um grafo completo



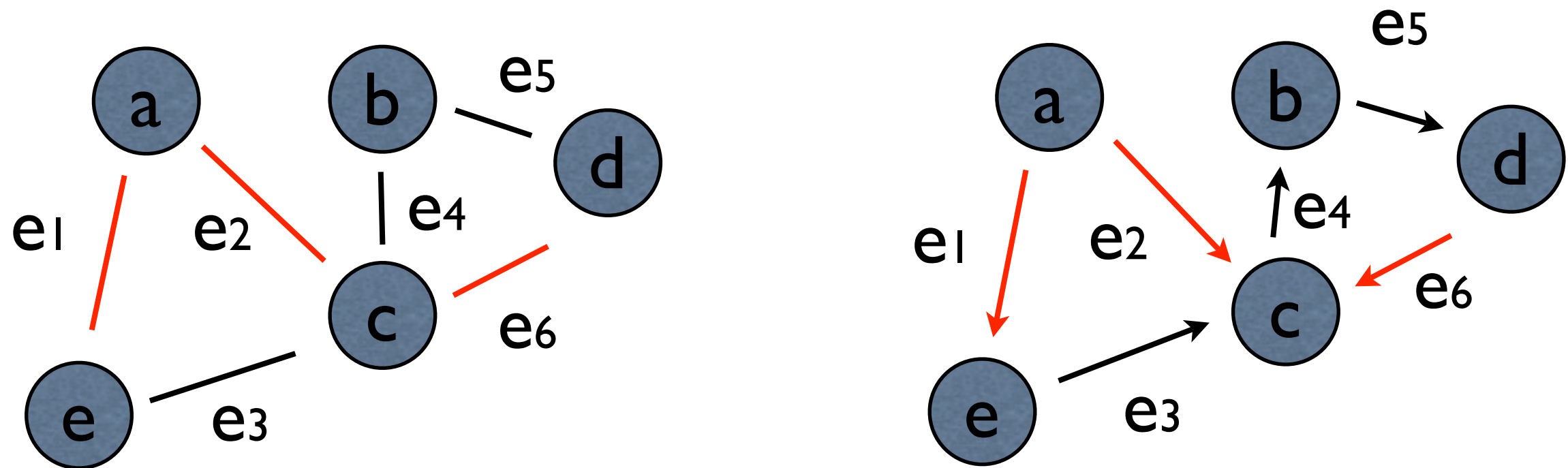
# Definições

- Clique
- Subconjunto de vértices que induz grafo completo



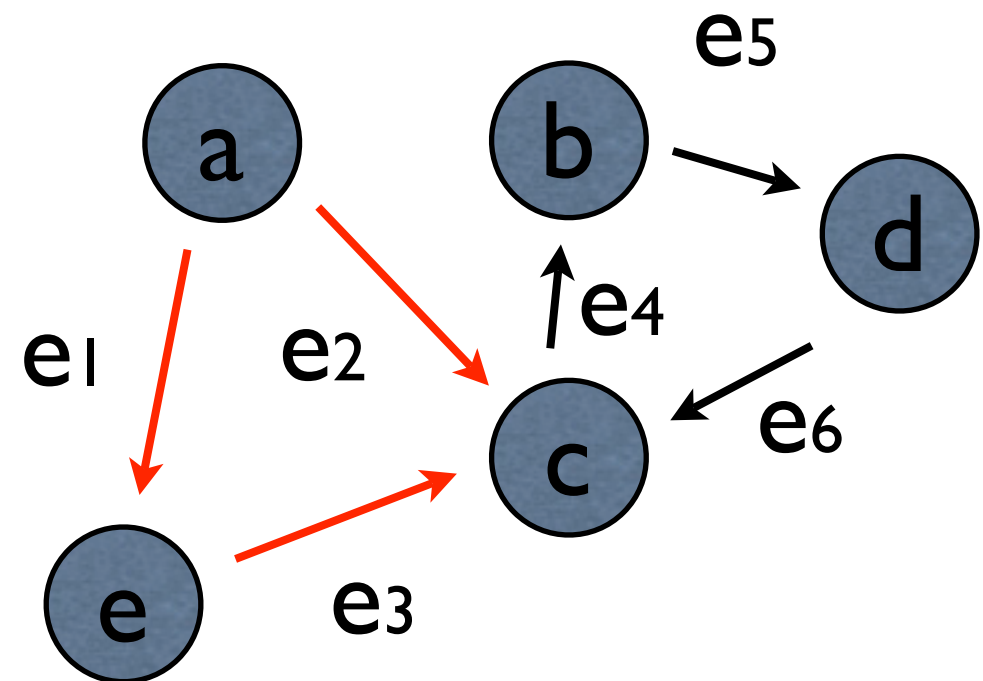
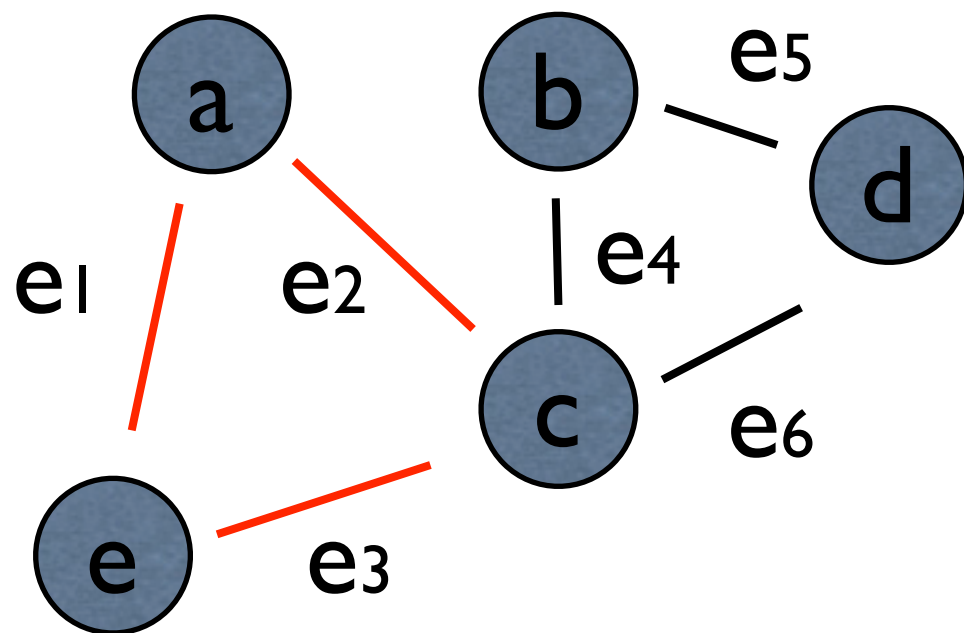
# Definições

- Cadeia
  - Sequência de arestas  $(e_1, e_2, \dots, e_n)$
  - $e_i$  vértice comum a  $e_{i-1}$  e  $e_{i+1}$



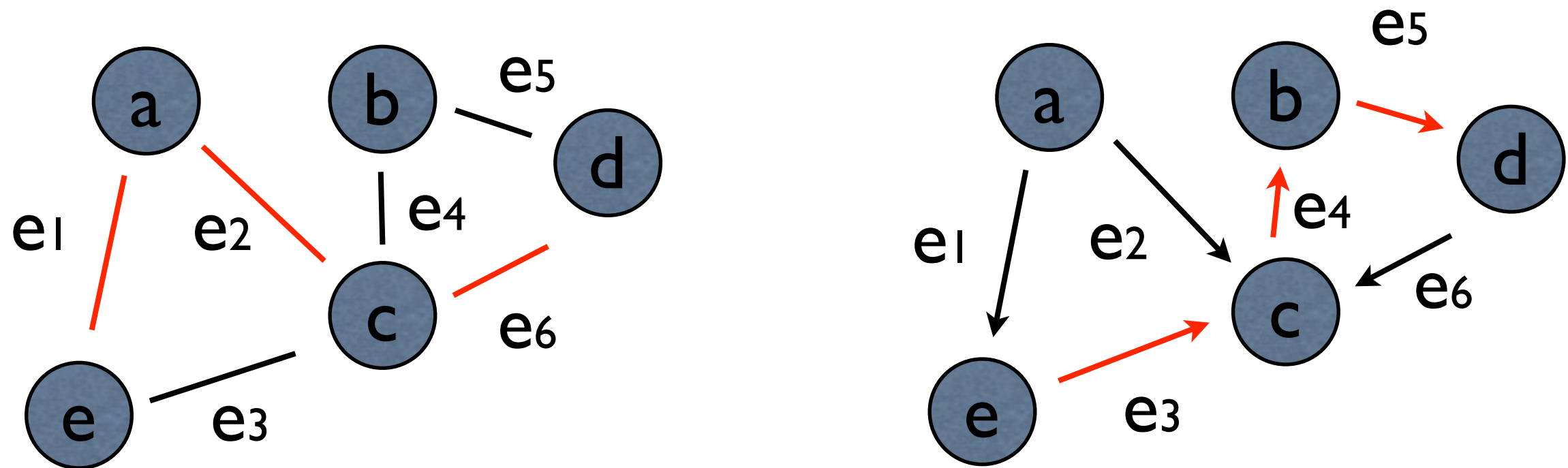
# Definições

- Ciclo
- Cadeia  $(e_1, e_2, \dots, e_n)$
- Vértice comum  $e_1$  e  $e_n$



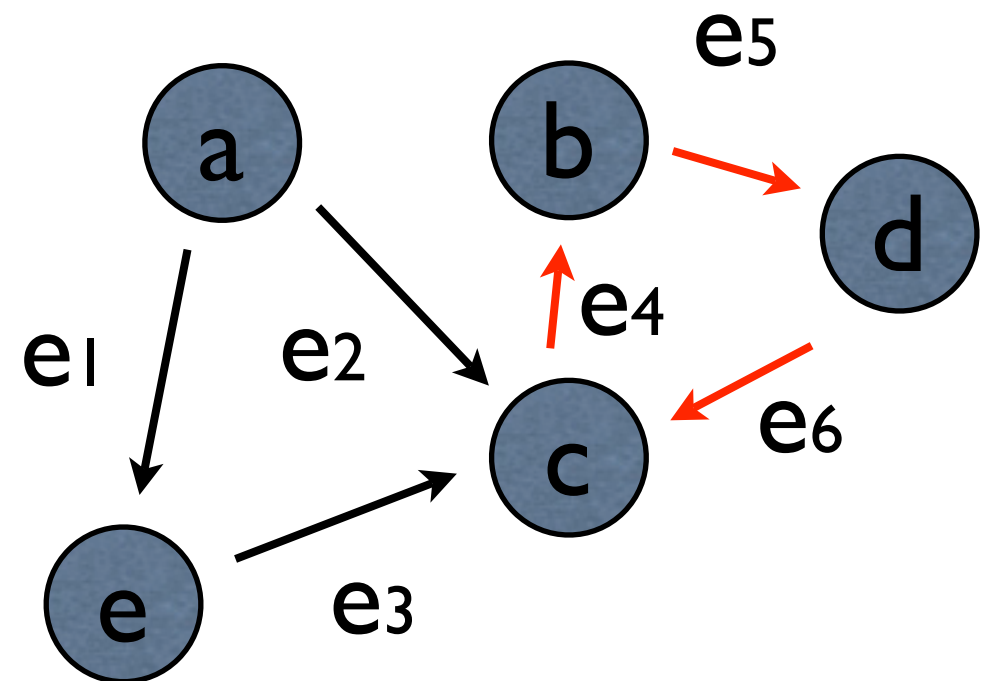
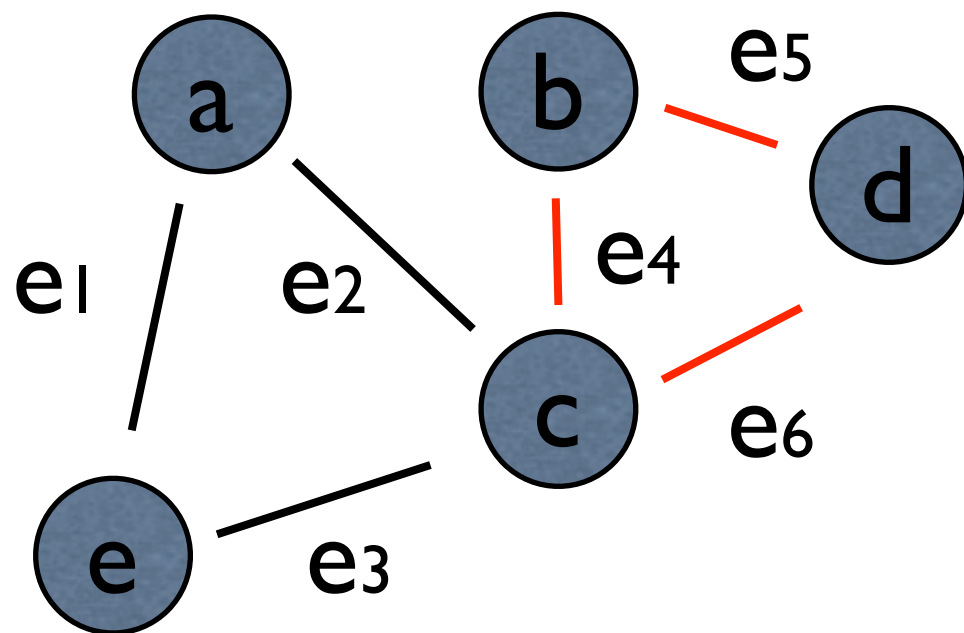
# Definições

- Caminho
- Cadeia ( $e_1, e_2, \dots, e_n$ )
- Obedece sentido das arestas (Digrafo)



# Definições

- Circuito
- Caminho  $(e_1, e_2, \dots, e_n)$
- Vértice comum  $e_1$  e  $e_n$



# Definições

- Caminho Euleriano
  - Caminho sem repetições de arestas
  - Deve passar por todos os vértices
- Circuito Euleriano
  - Caminho Euleriano Fechado



# Definições

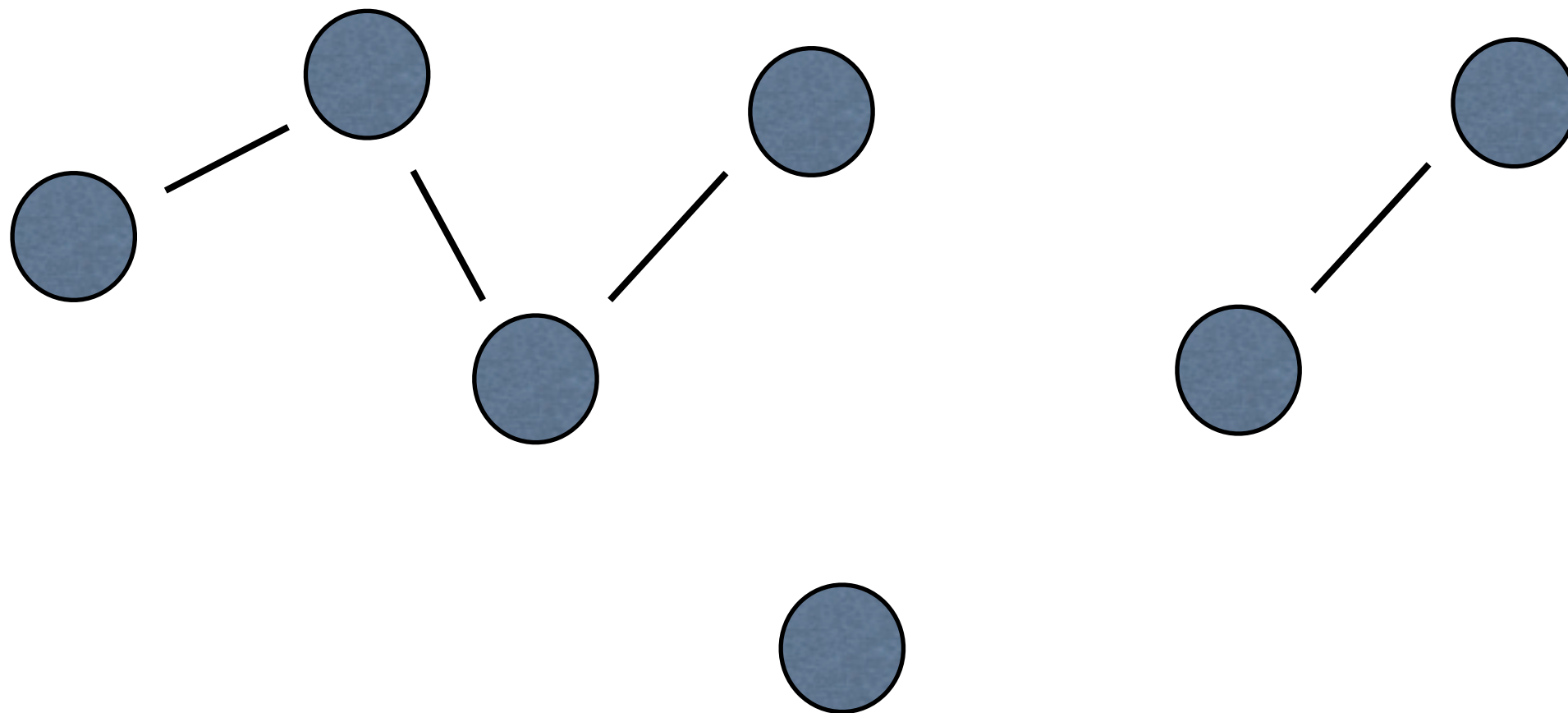
- Caminho Hamiltoniano
  - Caminho sem repetições de vértices
  - Passa por todos os vértices
- Circuito Hamiltoniano
  - Caminho Hamiltoniano Fechado

# Definições

- Fecho Transitivo Direto
  - Conjunto de todos os vértices atingidos por um caminho a partir de  $v$
- Fecho Transitivo Inverso
  - Conjunto de todos os vértices que podem atingir  $v$  por algum caminho

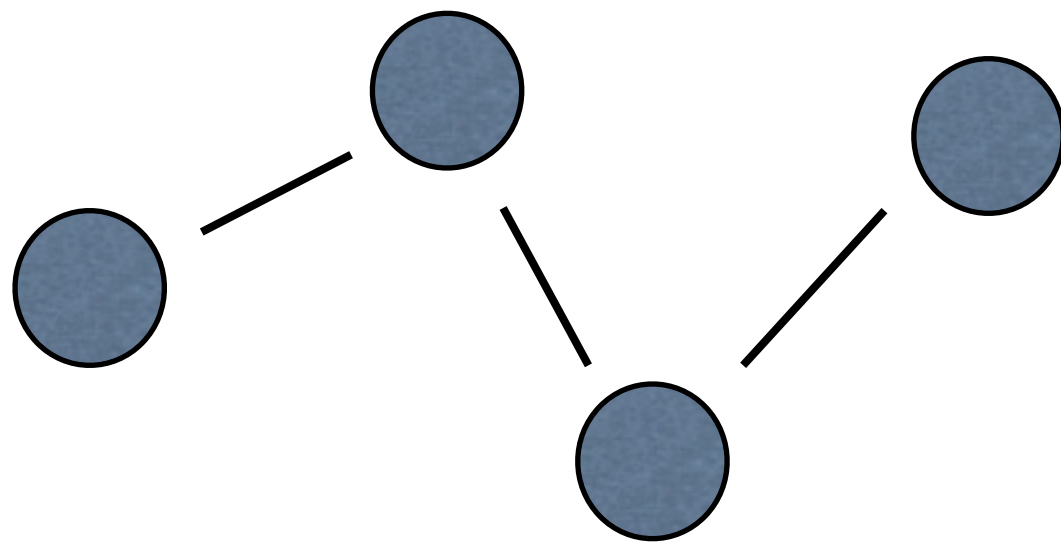
# Definições

- Componente Conexa
  - Existe cadeia entre qualquer par de vértices



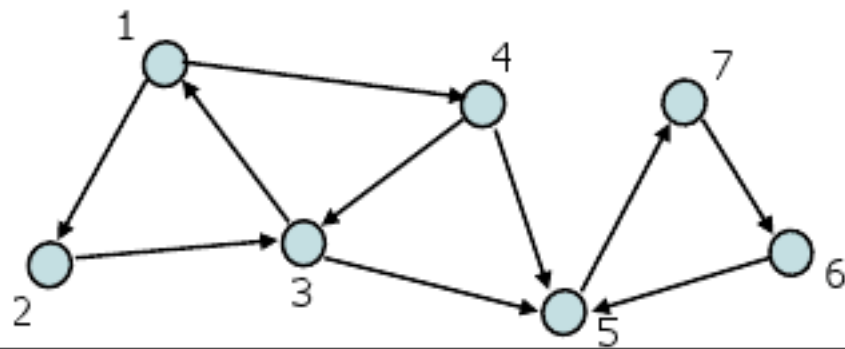
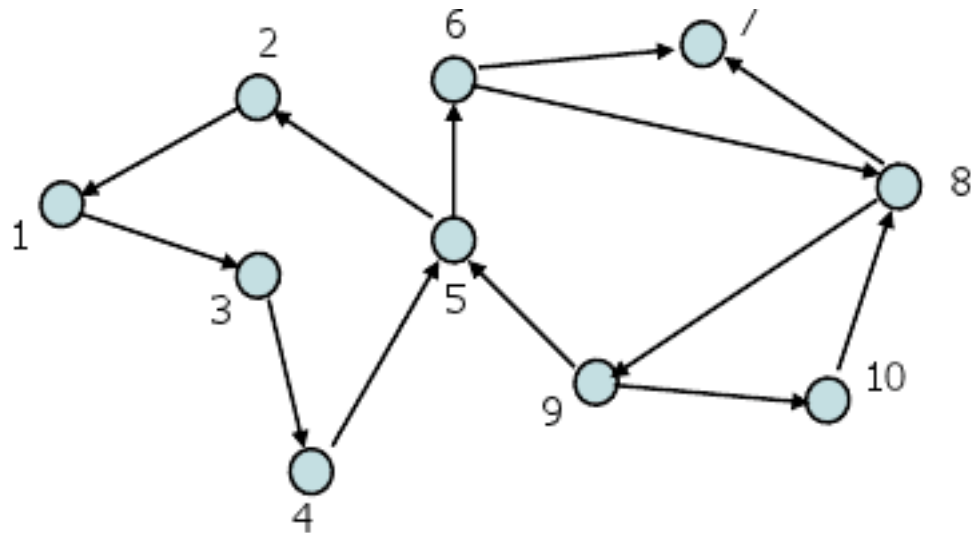
# Definições

- Grafo Conexo
- Pelo menos uma cadeia que liga cada par de vértices



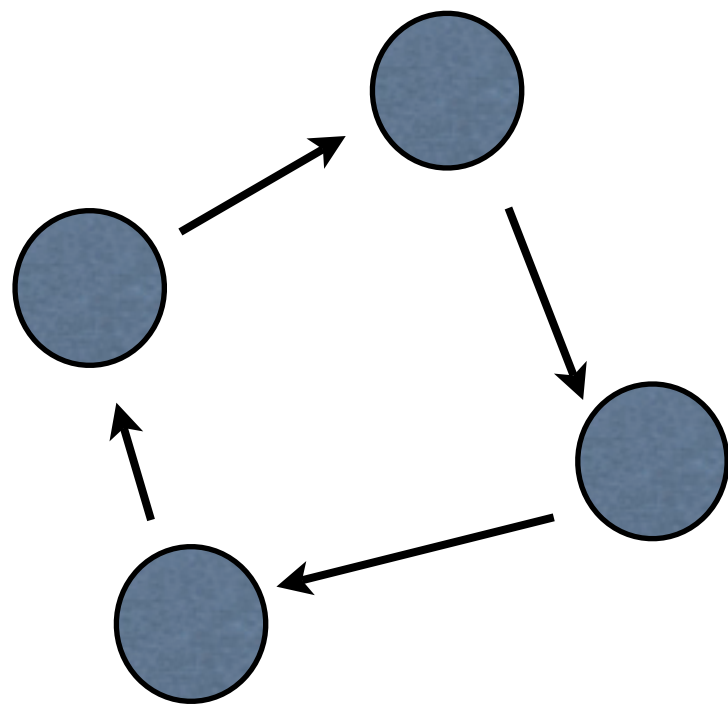
# Definições

- Componente Fortemente Conexa
  - Caminho  $v_i$  a  $v_j$
  - Caminho  $v_j$  a  $v_i$



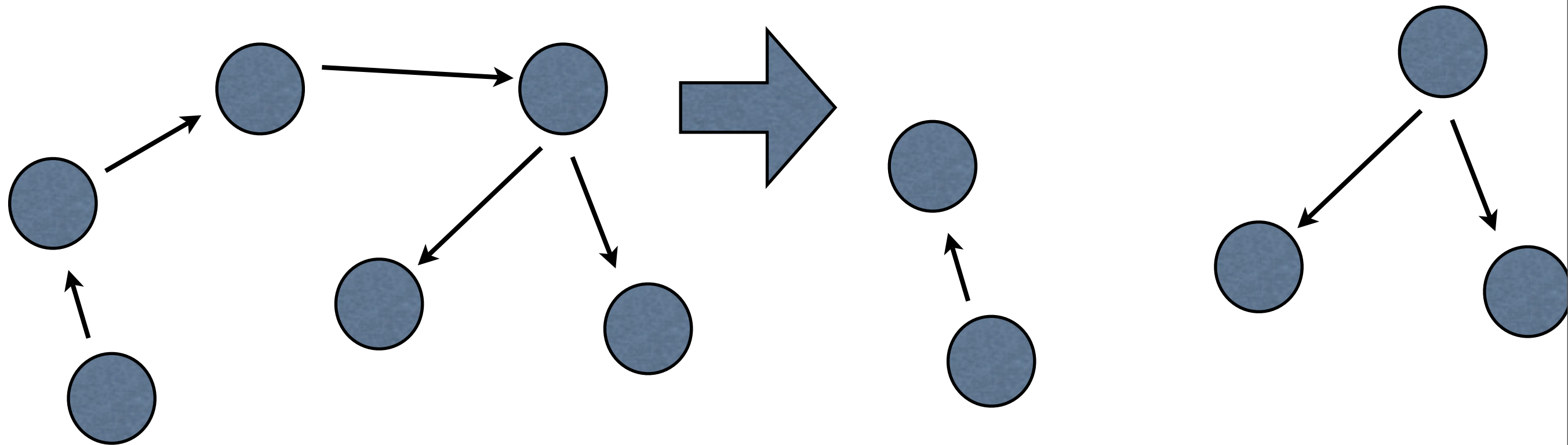
# Definições

- Grafo Fortemente Conexo
  - Caminho de  $v_i$  a  $v_j$
  - Caminho de  $v_j$  a  $v_i$



# Definições

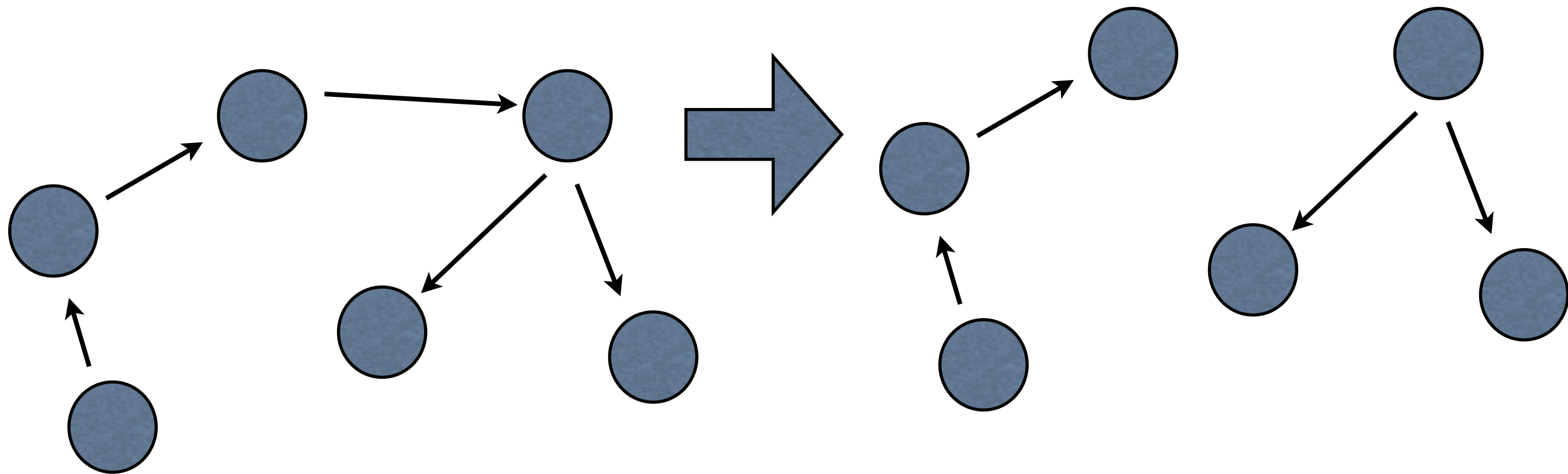
- Ponto de Articulação





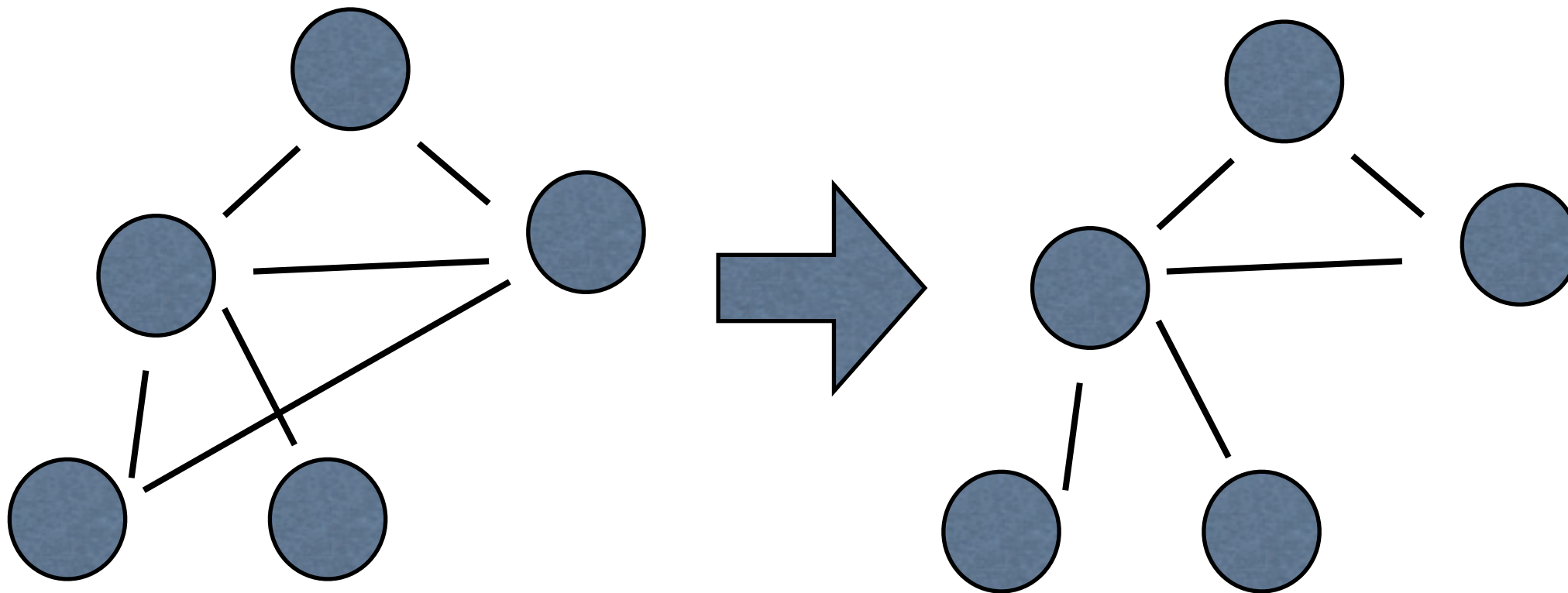
# Definições

- Ponte



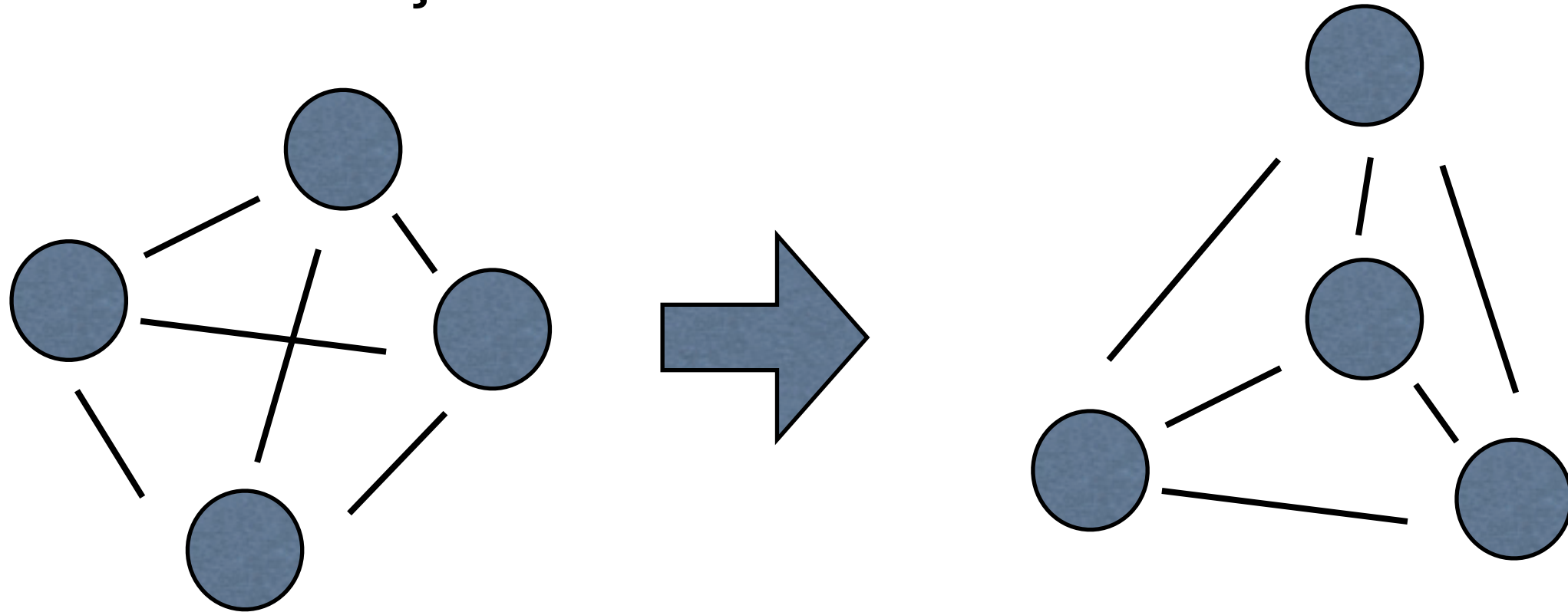
# Definições

- Grafo Gerador
  - $G=(V,E)$ ,  $G'=(V',E')$  e  $V=V'$



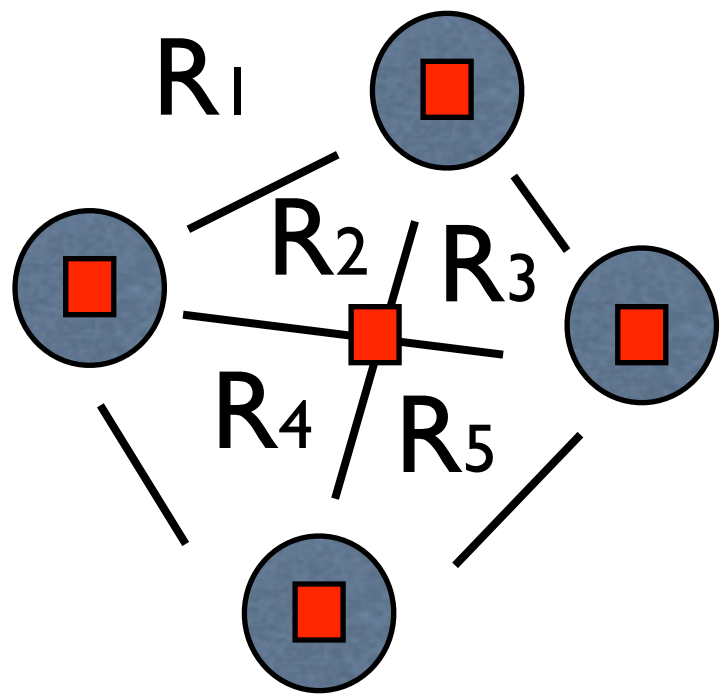
# Definições

- Planar
- **Pode ser representado** no plano sem interseção de arestas



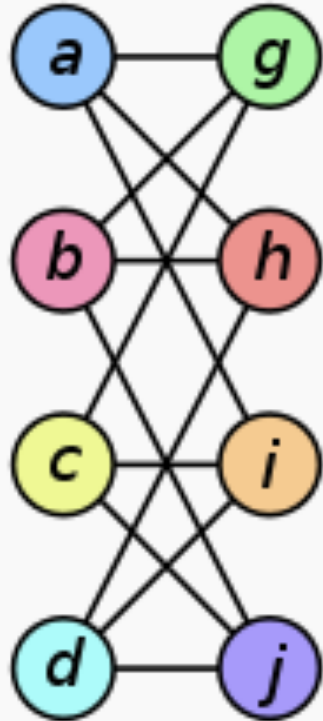
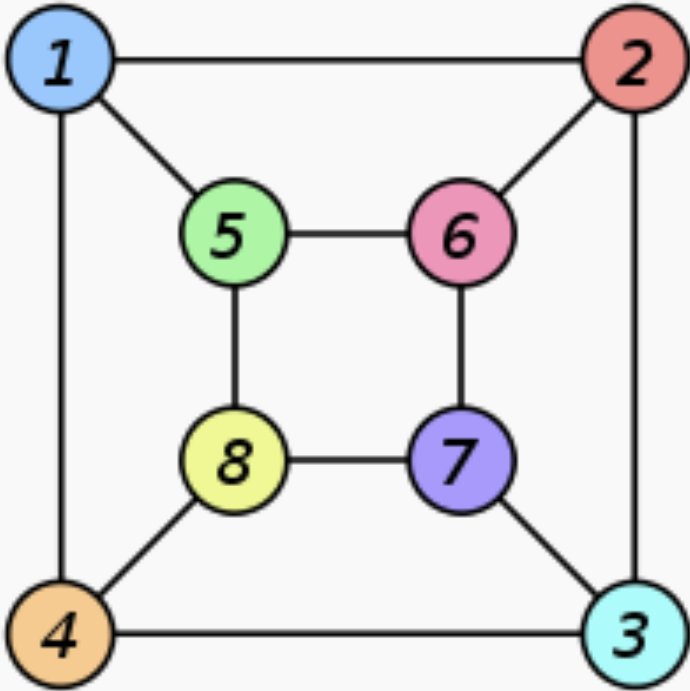
# Definições

- Planar
- $|R| = |E| - |V| + 2$  (demonstre)



# Definições

- Grafos Isomorfos
  - $G = (V, E)$  e  $G' = (V', E')$
  - $f: V \rightarrow V' \mid (f(u'), f(v')) \in E' \Leftrightarrow (u, v) \in E$

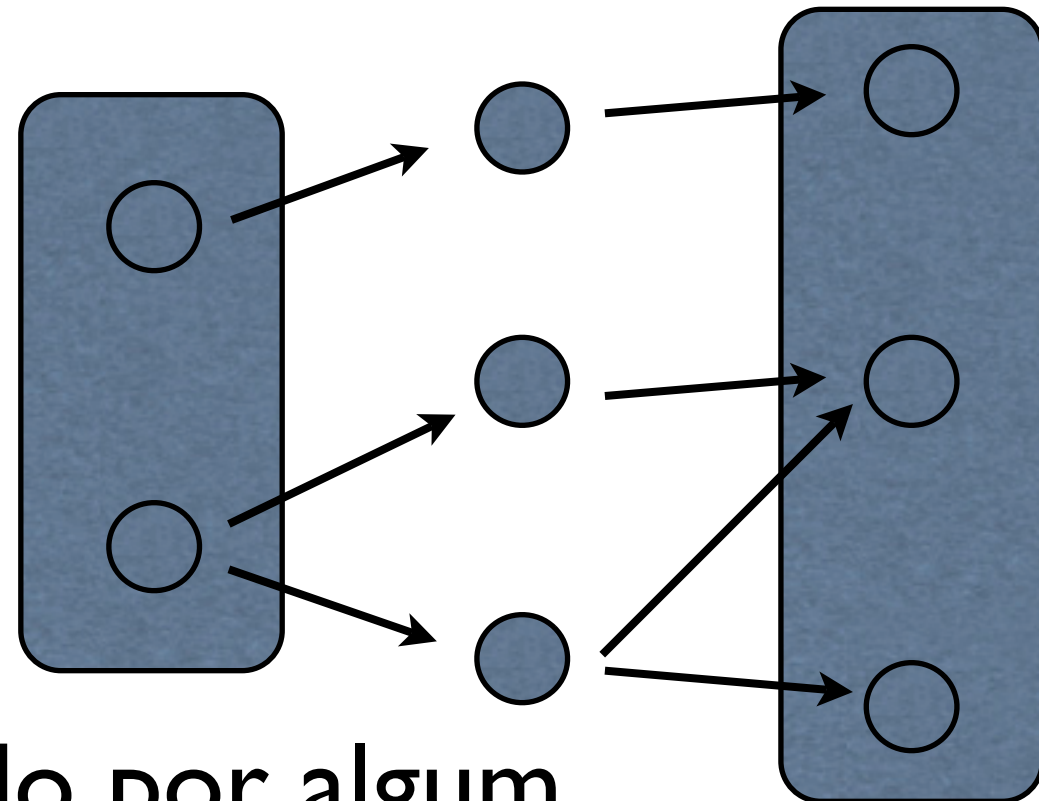
Graph G	Graph H	An isomorphism between G and H
		$\begin{aligned} f(a) &= 1 \\ f(b) &= 6 \\ f(c) &= 8 \\ f(d) &= 3 \\ f(g) &= 5 \\ f(h) &= 2 \\ f(i) &= 4 \\ f(j) &= 7 \end{aligned}$

# Definições

- Grafos Isomorfos
  - Encontre um grafo  $G_5$  isomorfo a seu complementar
  - Não há algoritmo eficiente

# Definições

- Base
  - Vértices em  $A$
  - Nenhuma aresta
- Vértices fora de  $A$ 
  - Todo vértice é atingido por algum caminho que parte de  $A$



# Definições

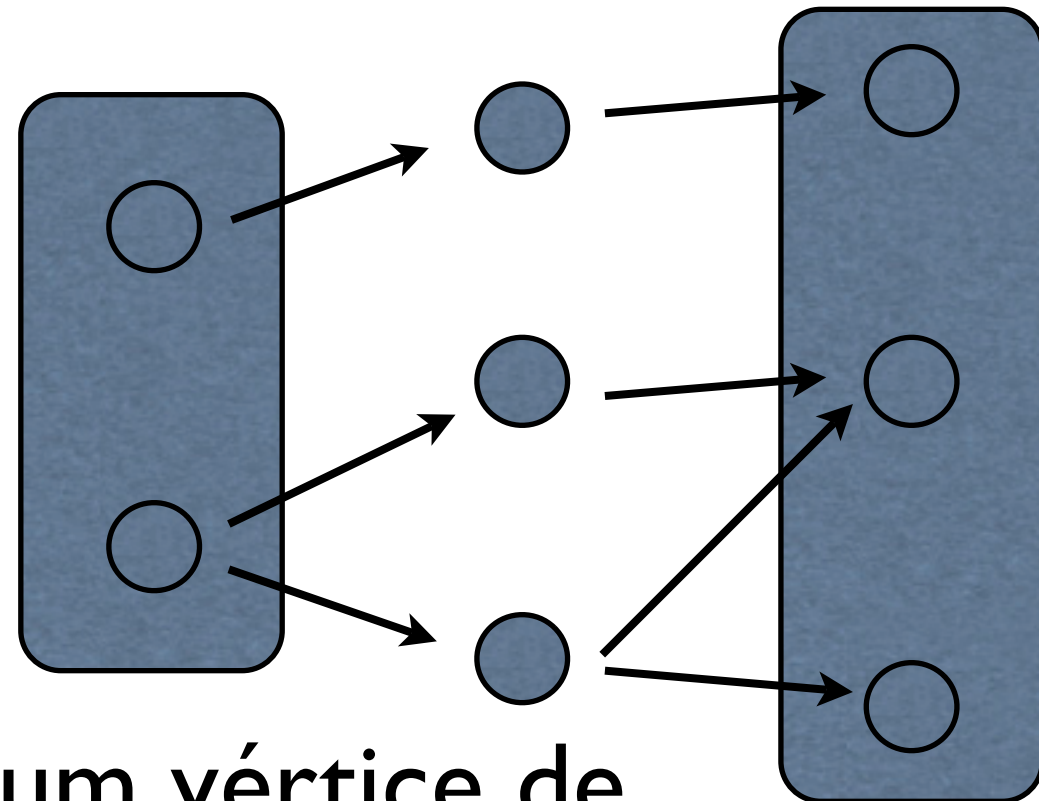
- Anti-Base

- Vértices em B

- Nenhuma aresta

- Vértices fora de B

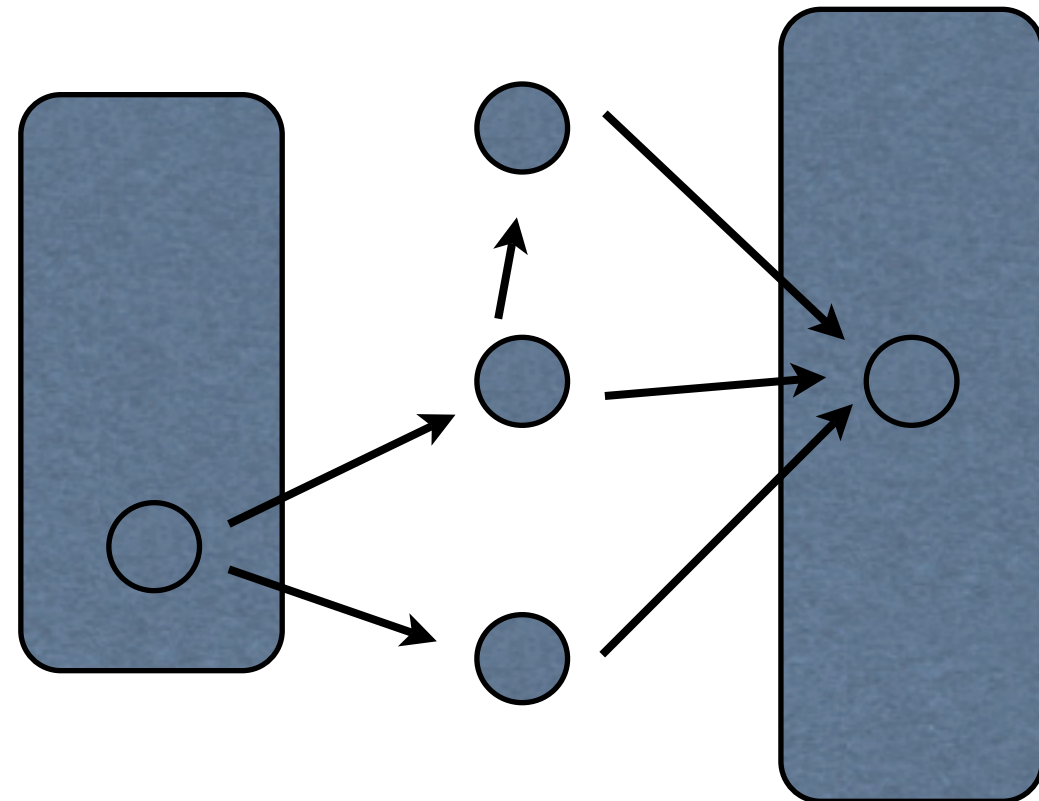
- Atingem pelo menos um vértice de B por algum caminho





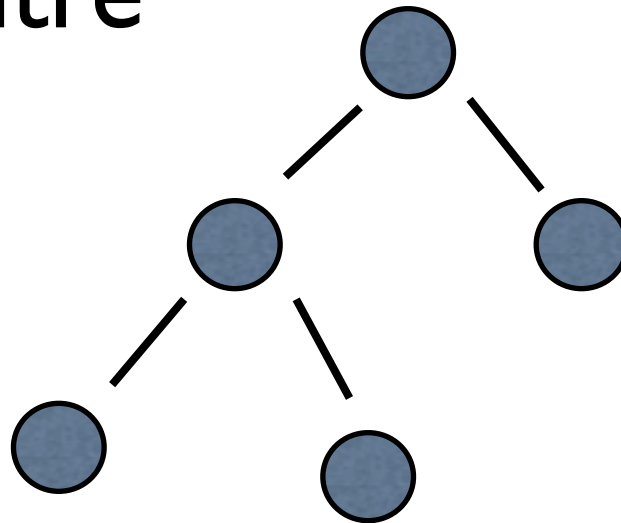
# Definições

- Raiz
  - Base unitária
- Anti-Raiz
  - Anti-Base unitária



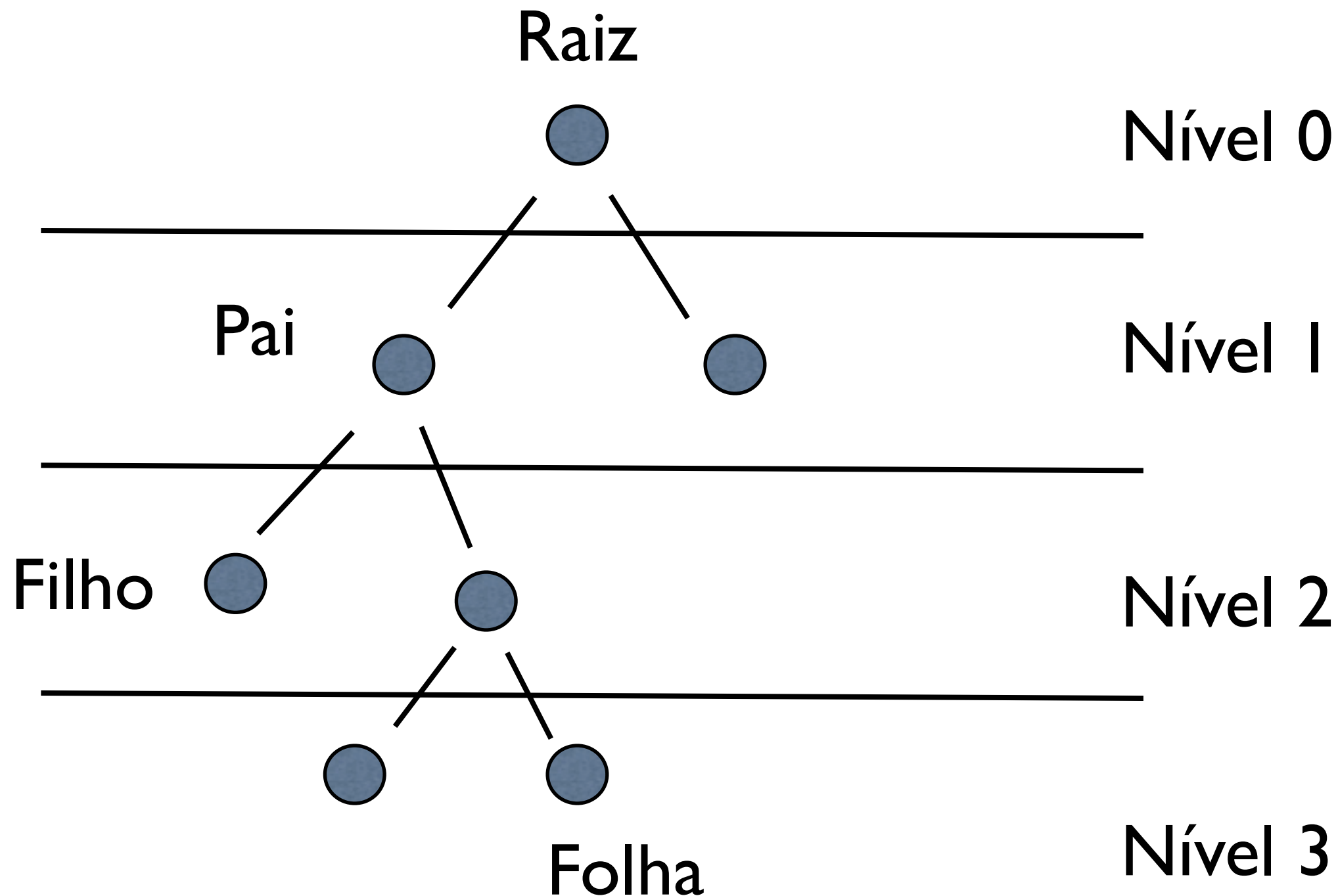
# Definições

- Árvore
  - Conexo
  - Sem ciclos
  - Único caminho entre dois vértices
- $|E| = |V| - 1$



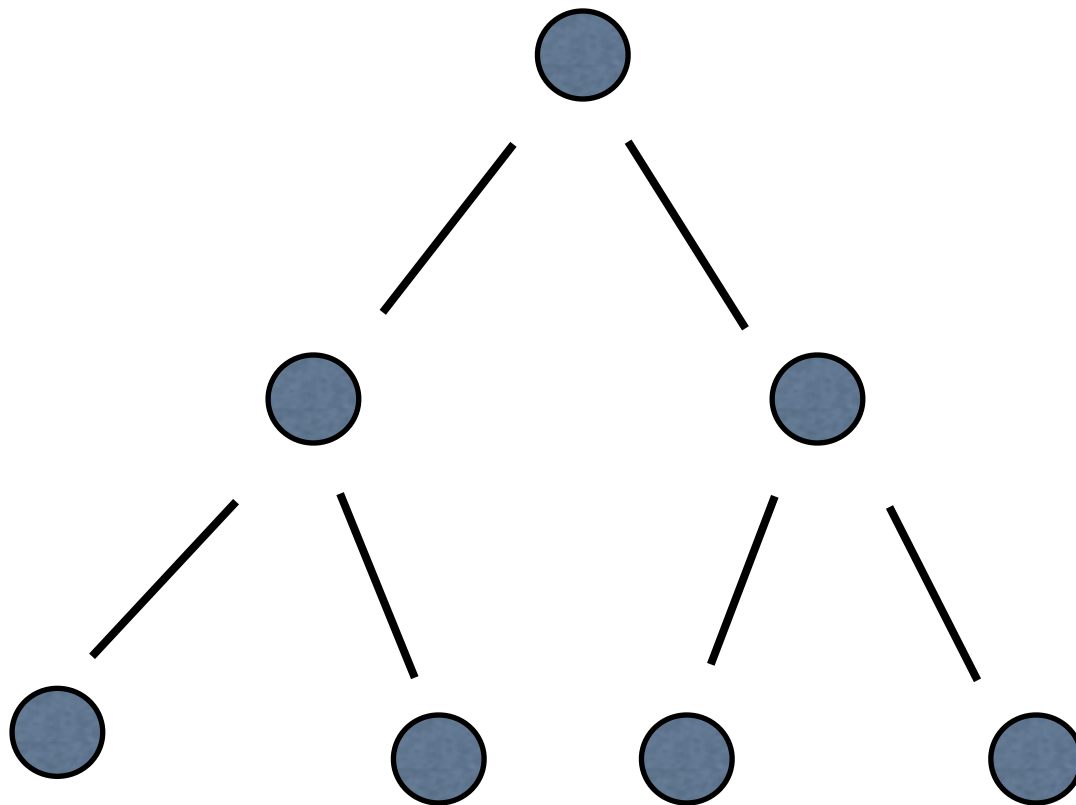
# Definições

- Árvore



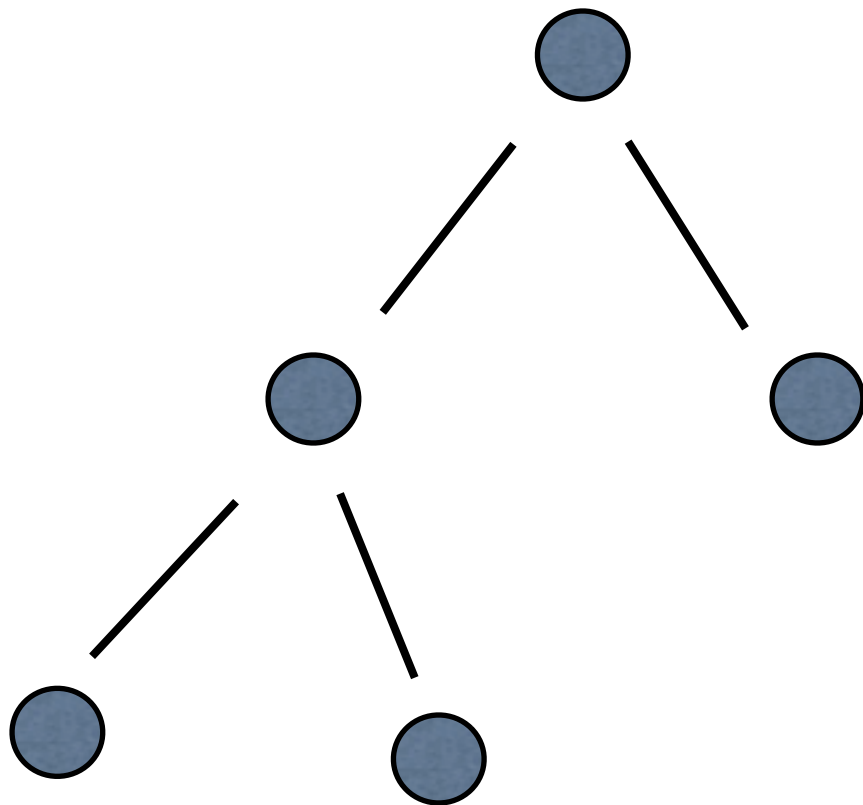
# Definições

- Árvore Cheia
- Subárvore vazia apenas no último nível



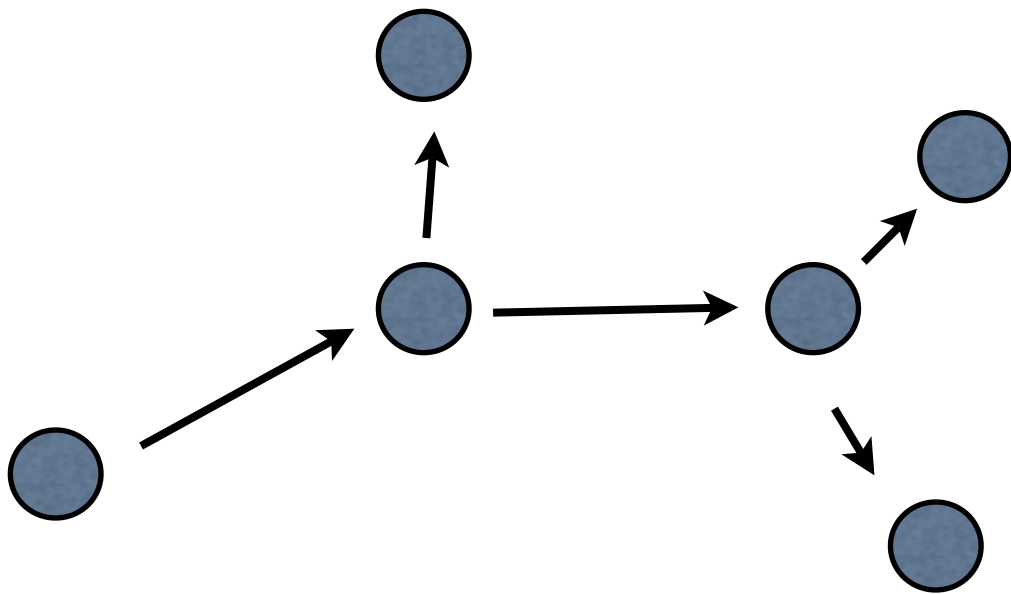
# Definições

- Árvore Completa
- Subárvore vazia apenas no último nível ou penúltimo nível



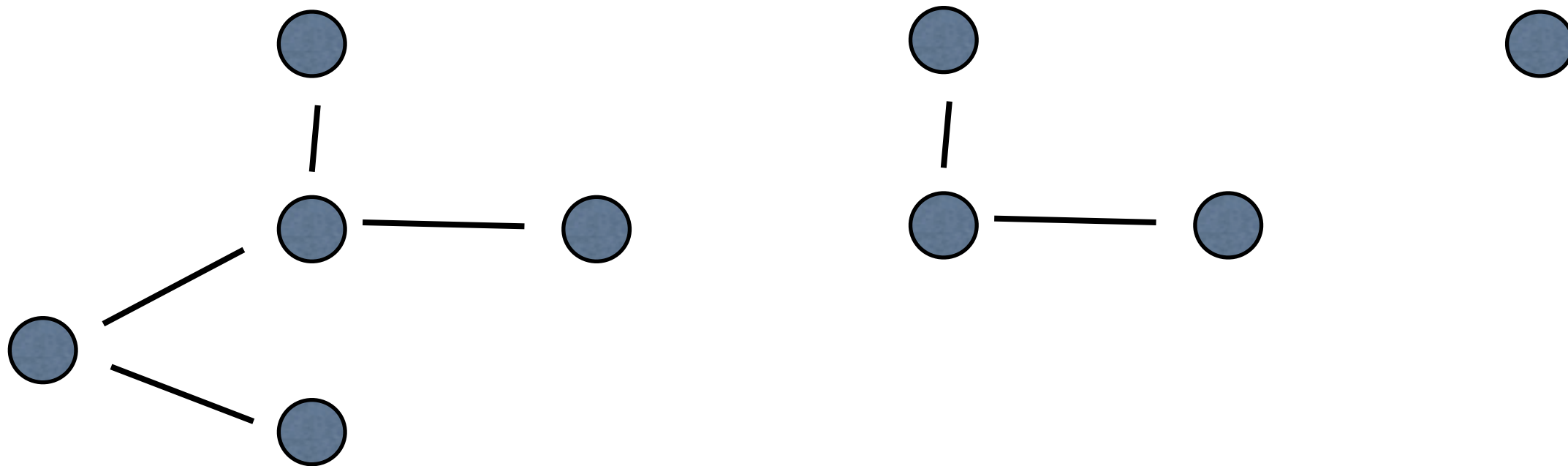
# Definições

- Arborescência
- Árvore Orientada



# Definições

- Floresta



# Definições

- Árvore Geradora
  - Árvore de  $G$
  - $G=(V,E)$ ,  $G'=(V',E')$  e  $V=V'$

