



IBM CAPI SNAP framework **Version 1.0**

How is data managed in the SNAP environment?

This Guide describes how to use, declare and optimize data variables in the SNAP environment.

To understand without any ambiguity all the explanation provided in this document, let's recall the basics:

- an application is the program running on the host server and executed on a CPU
- an action is the program running on the FPGA

Overview

Even if you use a standard language like C or C++, porting a program on to an FPGA requires understanding some important concepts about how data flows or is accessed. The FPGA has no Operating System which means that standard use of data in code executed on a CPU will require, on an FPGA, a formal read or write access to the physical resource. Understanding this key point will help the user to define not only which **path** to use to exchange data between the application and the action, but also the **way** to access so that you can easily and efficiently access the data.

This document has been built using snap git release tag v1.3.2. It can certainly work also for new releases.

This document will successfully go through the following items:

- understand the ways for an action to exchange data with an application
- understand the ways for an action to exchange data with all different resources
- understand the control logic used by SNAP to manage different actions





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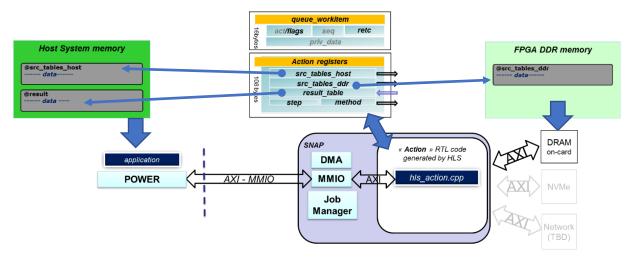
1. Actions interfaces

1.1 Configuration of an action by an application

The first key point to remember is that with CAPI, an action executed on the FPGA doesn't need a driver, which means that the application will NOT send the data to the action, since **the action will access data on its own**. The action will get information from the application as to where the required data is located.

A second key point to remember is that, thanks to CAPI, when an address is defined and used in the application, the same address will be usable without any modification in the action.

In other words, let's take the example where a user wants to process 2 tables of data in the action, one located in the host server memory and one in the FPGA DDR memory. The application will fill a structure (Action registers) to tell the action where the required data is located, and where the action should write the result. Any other parameters can be added to control the action if needed.



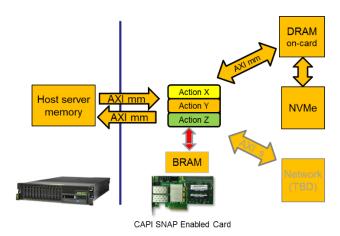
This structure exchanged between the application and the action is prepared by the application and sent before starting the processing. We will use the "MMIO" interface to read or write these settings. They can be dynamically changed during the processing since this interface is asynchronous to all other interfaces.





1.2 Transferring data from/to an action

The primary goal of the SNAP framework is to simplify the coding of the action which is implemented on the FPGA. Therefore, accessing several types of resources such as Host server memory, FPGA board DDR3 or DDR4, FPGA board Flash/NVMe memory needs to be standardized for the programmer by abstracting out the lower level details of each resource type. The other reason for this layer of abstraction is to prevent the user having to modify high level code when a new version of a resource type or a new resource type is added to the system.



A common standard and simple interface, named as AXI4 (Advanced eXtensible Interface release 4.0 developed by ARM), has been chosen to interface all the resources that an action needs to work with. This means that SNAP provides in its internal logic the driver to interface the different resources. If one (or more) of these resources is not used or not present on the board, the associated AXI driver to this resource can be disabled so that only used logic is implemented. Also, if a new resource is added to the system, the associated AXI driver can be added to the FPGA while the high-level user code stays as is, since it deals with the AXI abstraction layer.

As described in Xilinx UG761 document, this AXI4 protocol has 3 types:

- AXI (or AXI_mm) for high performance memory mapped requirements
- AXI-Lite for simple, low-throughput memory mapped communication (typically used for control and status registers)
- AXI-Stream for high-speed streaming data

In SNAP, **AXI_mm** is used for all data that is exchanged with external resources. **AXI-Lite** is used for status and control registers (MMIO). **AXI-stream** is not used yet but is meant for future usage such as connecting with the network driver attachment as soon as it is available.





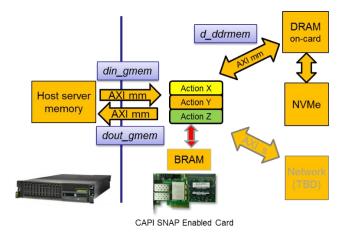
2. AXI data: the data "highway"

Let's start with the AXI data interface which allows data to be exchanged on a high-speed interface.

2.1 Type of resources

Three types of resources can be listed as:

- Common resources:
 - Host server memory: 1 unidirectional port IN + 1 unidirectional port OUT
- FPGA specific resources:
 - FPGA memory, named as **BRAM**: for example, a Xilinx KU060 has 3.5MB of memory used for user code variables.
- Board specific resources (for example):
 - AlphaData ADM-PCIE-KU3:
 - 8GB DDR3 bidirectional port
 - 40Gb SQFP+ ports (future use)
 - Nallatech 250S-2T:
 - 4GB DDR4 bidirectional port
 - 1.92TB NVMe accessible via the DDR4



2 important points to highlight:

- Action code accesses all resources using a unique AXI protocol which abstracts out the specific protocols of each of the resources.
- **"Board resources" are NOT known or seen by the application** running on the server. These resources can be managed by the Host application or by the action but need to be managed by one or the other.





2.2 "Size" of the interfaces

To ease the use of these different resources by the action code, and reduce the latency to access them, the choice has been made in SNAP that all resources are accessed with a **data width set to a fixed value defined by "snap_membus_t"**. This data type is defined as follow:

~snap/actions/include/hls snap.H:

This means that a full word read from or written to the host server memory or the FPGA DDR will always be 512 bits / 64 Bytes data in physical size. We will see soon how to deal with this "constraint".

2.3 Ports declaration

All actions provided as examples in SNAP have the "**ports**" defined the same way. Remember that **hls action** is a reserved name that is used by SNAP logic to find the top file of the code.

~snap/actions/hls_*/hw/action_*.cpp:

```
//--- TOP LEVEL MODULE ----
void hls_action(
    snap_membus_t *din_gmem,
    snap_membus_t *dout_gmem,
    snap_membus_t *d_ddrmem, // CAN BE COMMENTED IF UNUSED
    snapu32_t *d_nvme, // CAN BE COMMENTED IF UNUSED
.../...
```

All unused ports can be commented except the ports of the host memory declared as "din_gmem" and "dout_gmem" which needs to remain active. "din_gmem" is the unidirectional port through which all data read from the Host server memory will enter the action. "dout_gmem" is the unidirectional port through which all data written to the Host server memory will exit the action.

The port "**d_ddrmem**" is bidirectional which means that data read from or written to the FPGA board DDR will flow through this port. The port can address DDR3 or DDR4 types of memory since the SNAP logic will adapt automatically the right driver to the right type of DDR (card dependent).

The port "**d_nvme**" is a control port used to define the access to the NVME. All data are read and written through the DDR memory, used as a cache. Therefore, if the action needs data from the NVMe, it sends a read request of X bytes at a NVMe address, and the DDR will be filled with a block containing the data requested. The reason of this architecture is that the NVMe has a large latency and accessing it directly would slow the action program execution.





More explanation can be found in the document provided in snap NVME example: https://github.com/open-

power/snap/blob/master/actions/hls_nvme_memcopy/doc/UG_SNAP_hls_nvme_memcopy_v1.0.pdf

2.4 Ports definition

In the top-level module *hls_action* of each hardware action, after their declaration, all the ports are defined as "#pragma HLS INTERFACE m_axi" ports. This type of definition is used for data path using the AXI4 protocol:

```
//--- TOP LEVEL MODULE -----
void hls action(
    // Host Memory AXI Interface
#pragma HLS INTERFACE m axi port=din gmem bundle=host mem offset=slave depth=512 \
 max_read_burst_length=64 max_write_burst_length=64
#pragma HLS INTERFACE s axilite port=din gmem bundle=ctrl reg offset=0x030
#pragma HLS INTERFACE m axi port=dout gmem bundle=host mem offset=slave depth=512 \
 max read burst length=64 max write burst length=64
#pragma HLS INTERFACE s axilite port=dout gmem bundle=ctrl reg offset=0x040
// DDR memory Interface
                            //CAN BE COMMENTED IF UNUSED
#pragma HLS INTERFACE m axi port=d ddrmem bundle=card mem0 offset=slave depth=512 \
 max read burst length=64 max write burst length=64
#pragma HLS INTERFACE s_axilite port=d_ddrmem bundle=ctrl_reg offset=0x050
    //NVME Config Interface //CAN BE COMMENTED IF UNUSED
#pragma HLS INTERFACE m axi port=d nvme bundle=nvme
```

Every port is defined with the name that will be used in the code to access it (din_gmem, dout_gmem, d_ddrmem, d_nvme). The declaration of the bundle is the prefix of the name of the signal which will be used by the logic to connect to the different physical interfaces through a wrapper. Modifying this bundle name will prevent your logic from being connected to the SNAP wrapper.

The parameters (slave_depth, max_read_burst_length, max_write_burst_length) are set for the Xilinx HLS tool to overwrite default settings and get the maximum AXI burst (4KB).

Refer to Xilinx UG902 Vivado HLS user guide for more details on these options.





2.5 Unused ports

If the action code doesn't need any access to the DDR or the NVMe, then commenting these lines will remove the DDR driver and save place removing unnecessary logic.

The comment must also be done in the port definition:

Removing this port must **always** been done by also disabling the resource used in the **make snap_config** menu.

You will notice that all examples provided in SNAP (hls_helloworld, hls_memcopy, ...) have non-selectable resources in the **make snap_config** menu so that we cannot select or unselect a resource by mistake. These preselected choices are described in ~snap/scripts/Kconfig and also per card and per action in ~snap/defconfig/ADKU3.hls_helloworld.defconfig or ~snap/defconfig/ADKU3.hls memcopy.defconfig,...

When building your own action, the easiest way may be to pick as a template, an action which has the same interfaces that you need so that all the preselection job is already done.





2.6 Address of port access

You may have noticed in a previous paragraph that each port definition is followed by a #pragma HLS INTERFACE s_axilite line defining a register address for each port.

```
//--- TOP LEVEL MODULE ------
void hls_action(
.../...

// Host Memory AXI Interface

#pragma HLS INTERFACE m_axi port=din_gmem bundle=host_mem offset=slave depth=512 \
max_read_burst_length=64 max_write_burst_length=64

#pragma HLS INTERFACE s_axilite port=din_gmem bundle=ctrl_reg offset=0x030
.../...
```

This address offset should not be used by the user code since it is forced by SNAP internal logic to 0 for all ports.

Read and Write are considered from the application / software side								
Card init	Done on	ce at card initializati						
Write@	Read@	3	Typical Write value					
0x030			host_mem port din initialization (LSB)					
0x034			host_mem port	din initialization (MS	B)	0		
0x038			HLS reserved					
0x040			host_mem port dout initialization (LSB)					
0x044			host mem port dout initialization (MSB)					
0x048			HLS	reserved				
0x050			card_mem0 port ddr initialization (LSB)					
0x054			card_mem0 port	ddr initialization (M	SB)	0		
0x058			HLS	reserved				

The reason is that *Xilinx HLS uses only aligned address*. This means that with a data bus width set to 64 bytes, all the address passed through this register will be automatically converted to a 64 bytes address. To simplify things for the coder and reduce the latency taken by width conversions, all interfaces have been set to the same data bus width as explained in "2.2 "Size" of the interfaces". This means that the user will use same alignment for all address.

To prevent from losing this byte address information, it has been decided to use a specific register that allows the action code to know which byte to address when reading a 64 bytes word. The caveat is that the conversion of this Byte address need to be done by the user before giving it to the function accessing memories.

Let's explain deeper what is an aligned address.

Imagine that an application uses an address 0xABCD (in hexadecimal). This is a byte address. Every byte is associated to a specific address.

Reading a word at address **0xABCD** will give you a 1 Byte word **Byte X**





Byte	1-Byte
Address	Data
0xABD1	
0xABD0	
0xABCF	Byte_Z
0xABCE	Byte Y
0xABCD	Byte_X
0xABCD 0xABCC	Byte_X Byte_W
0xABCC	Byte_W
0xABCC 0xABCB	Byte_W Byte_V
0xABCC 0xABCB 0xABCA	Byte_W Byte_V

0xABD1	
0xABD0	
0xABCF	Byte_Z
0xABCE	Byte_Y
0xABCD	Byte_X
	_,
0xABCC	Byte_W
0xABCC 0xABCB	
	Byte_W
0xABCB	Byte_W Byte_V
0xABCB 0xABCA	Byte_W Byte_V

			•
1-R\	/te	inte	rface

	Byte			2-Bytes Address
	Address	Data_Byte0	Data_Byte1	(Byte address >> 1)
	0xABD0	Byte_Z		0x55E8
	0xABCE	Byte_Y	Byte_Z	0x55E7
	0xABCC	Byte_W	Byte_X	0x55E6
	0xABCA	Byte_U	Byte_V	0x55E5
	0xABC8			0x55E4
_				

2-Bytes interface

Now if you use a 2-Bytes interface, meaning that you can read 2 Bytes with a single address, then you have 2 choices:

- Read 2 Bytes at **0xABCC** will give you a 2 Bytes word **Byte W Byte X**
- Read 2 Bytes at OxABCD will give you a 2 Bytes word <u>Byte_X</u> Byte_Y

To save place and if we use only 2 Bytes access then removing all **odd** addresses will still let you have access to all Bytes. This is what we call a 2-Bytes aligned address.

Thus, the new 2-Bytes address would be 0xABCC or 0xABCD shifted 1 bit to the right becomes 0x55E6. Having the Byte address information can let you work with the byte you expect in the 2-Bytes word read.

Extending that to our case, a 64-Bytes address needs the Byte address shifted 6 to the right (address 00h to 63h = address 00_0000b to 11_1111b)

Thus, the new 64-Bytes address would be 0xABCC or 0xABCD shifted 6 bits to the right becomes 0x2AF

In the action code, you will so understand why an address sent by the application as a Byte address, will need to be aligned with the port width (ADDR RIGHT SHIFT is set to 6 and defined in file ~snap/actions/include/hls snap.H)

```
/* byte address received need to be aligned with port width */
input address = act reg->Data.in.addr >> ADDR RIGHT SHIFT;
output_address = act_reg->Data.out.addr >> ADDR_RIGHT_SHIFT;
```

The key point is that the coder can still work in the action and in the application with a Byte address. If he wants to access a word at a specific Byte address, then he has the information and can access the Xth Byte of the 64 words.

2.7 How to use ports

Now that the ports are declared with a name, data that flows through them need to be read or written at an address associated with the access request.

To read a data from host server memory, then the address used will be din_gmem + input_address. To write a data to host server memory, then the address used will be **dout gmem + output address**. To read or write a data from/to FPGA DDR memory, then the address used will be **d_ddrmem + address**.





Several ways can be used to read a data. The easiest way for a read:

Same for a write:

All data are stored in a buffer which can be used as a standard array of data.

You can also use type casting as for any C code:

```
char text[BPERDW]; /*BPERDW (Byte per Data Word is defined as 64*/

/* Read in one 64B word */
memcpy((char*) text, din_gmem + input_address, BPERDW);
```

The function **memcpy** called is a HLS function which guarantee the most efficient access to the data in the way that:

- It allows to access any number of bytes you need not depending on the size of the word defined by snap_membus_t type. Requesting 1 byte will give you 1 byte!
- It will initiate a **burst** so that accessing a big amount of data stored at contiguous address will be done in a minimum amount of time. Accessing 4K bytes of data (maximum burst of AXI) in one shot will be much quicker than accessing 64 times 64 byte words.

There are other ways to access data which can be as efficient than the HLS **memcpy** but may lose one of these 2 above properties.

The first option is to use the closest translation of the **memcpy** function which keeps the **burst** property if no other processing is done within the loop, but force to read full **words of 64 bytes**:





```
for (int k=0; k<size_in_words; k++)
#pragma HLS PIPELINE
    buffer[k] = (din_gmem + input_address)[k];</pre>
```

By extension, if you want to read or write a single word of 64 bytes, you can just use

```
buffer[0] = (din_gmem + input_address)[0];
or
  (dout_gmem + input_address)[0] = buffer[0];
```

2.8 Optimizing data access

It is good to understand that Host memory access is done through a 128Bytes cache line. This means that:

- Read 2 words from Host memory will give user the maximum bandwidth
- Write less than 2 words to host memory will use a "read-modify-write" access that may be not as efficient as writing 2 words which do not need a read access and the partial writes.

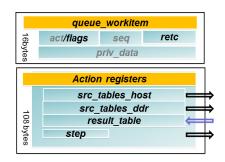
It may be more efficient to buffer data you need to use in the code prior to use them, rather than doing an access when you need the data. This would act as a cache and optimize the access.





3. AXI Lite data: the registers interface

Let's understand how to specify and use the MMIO interface used to exchange registers. As you may remember, AXI Lite resources are used for registers status and control using the MMIO interface.



\$ACTION_ROOT/hw/hw_action.H

3.1 Action registers: Control structure

The *queue_workitem* defined as the "*Control*" structure is the header of the action registers. It is managed by SNAP job manager and defined as follow:

\$SNAP_ROOT/actions/include/hls_snap.H

```
typedef struct {
    snapu8_t sat; // short action type
    snapu8_t flags;
    snapu16_t seq;
    snapu32_t Retc;
    snapu64_t Reserved; // Priv_data
} CONTROL;
```

The action uses the *flags* field as a read only register to know if the action has been set already once (this is what we call the discovery mode).

The *retc* field is used as a write only register by the action to report the return code of the action.

\$ACTION_ROOT/hw/hw_action.cpp

```
void hls_action(snap_membus_t *din_gmem,....
.../...
switch (act_reg->Control.flags) {
  case 0:
    Action_Config->action_type = HELLOWORLD_ACTION_TYPE; //TO BE ADAPTED
    Action_Config->release_level = RELEASE_LEVEL;
    act_reg->Control.Retc = 0xe00f;
    return;
    break;
  default:
        process_action(din_gmem, dout_gmem, act_reg);
    break;
}
```





On the application side, there is no access to the Control flag, but the **Return code** can be tested at any time:

\$ACTION ROOT/sw/snap action.cpp

3.2 Action registers: Data structure

The "Actions registers" **Data** structure is where the user defines all parameters that will be used for a specific action. This structure is filled by the application for example as follows:

\$ACTION ROOT/include/action.H

```
typedef struct action_job {
    struct snap_addr src_tables_host;
    struct snap_addr src_tables_ddr;
    struct snap_addr result_table;
    uint16_t step;
} action_job_t;

/* input text in HOST: 128 bits*/
    /* output result in HOST: 128 bits*/
```

The smartest and easiest way is to use **snap_addr**, a SNAP predefined 16 Bytes (= 128bits) type, defined to specify in a minimum of bytes the exact location of the data to use. This type is defined as follow:

\$SNAP_ROOT/software/include/snap_types.h

```
typedef struct snap_addr {
    uint64_t addr;
    uint32_t size;
    snap_addrtype_t type;
    snap_addrflag_t flags;
} /* DRAM, NVME, ... */

Snap_addr_t;

** SRC, DST, EXT, ... */
```

Rules to follow when defining the 108 Bytes of registers structure:

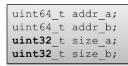
1. It has been experienced that the C compiler used by HLS prefer some simple definitions:

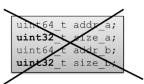
```
snap_addr table_src0;
snap_addr table_src1;
is preferable to
snap_addr table_src[2];
...
```

2. As the registers are defined as 64 bit registers, prefer 64 bit alignments rather than mixing sizes:









' is preferred at

3. C compiler forces 64 bits alignment which generates SNAP errors

```
typedef struct helloworld_job {
  struct snap_addr in; // input data => 16 Bytes
  struct snap_addr out; // offset table => 16 Bytes
  char is_the_problem; // => 1 Byte
  } helloworld_job_t;
  j = sizeof(helloworld_job_t); => reports 40 Bytes while it should report 33 Bytes!
```

Have a look at the issue described in https://github.com/open-power/snap/issues/488 which provides a way to circumvent this

```
typedef struct helloworld_job {
  struct snap_addr in; // input data => 16 Bytes
  struct snap_addr out; // offset table => 16 Bytes
  char is_the_problem; // => 1 Byte
  char padding[7]; // => 7 Bytes
} helloworld_job_t;
j = sizeof(helloworld_job_t); => reports 40 Bytes
```

Maximum Size:

The maximum size of the "Action registers" **Data** structure defined here as "action_job" has been set to 108 bytes.

\$ACTION ROOT/hw/hw action.H

Checking is done when compiling the (hardware) action: attempting to go beyond this limit will generate a compilation error.

```
Checking for reserved MMIO area during HLS synthesis ... /home/.../snap//actions/hls.mk:69: recipe for target 'check' failed
```

If the definition of your structure is less than these 108 bytes a **padding** will be done so that this structure is always 108 Bytes. This constraint is due to the way Xilinx HLS tool manages registers, and the way we were constrained to use it so that it could friendly enough to use it.

In case you need more than 108 Bytes, it's very easy to define an area in host memory that will contain all the information you need to exchange between the application and the action.





Action's access to these registers:

The action will see these 108 Bytes of Data as a structure defined in \$ACTION_ROOT/include/action.H and accessed as for any C structure:

\$ACTION_ROOT/hw/hw_action.cpp

```
/* byte address received need to be aligned with port width */
i_idx = act_reg->Data.src_table_host.addr >> ADDR_RIGHT_SHIFT;
o_idx = act_reg->Data.result.addr >> ADDR_RIGHT_SHIFT;
size = act_reg->Data.src_table_host.size;
```

Application's access to these registers:

The application fills the Data structure using predefined SNAP APIs as follow:

\$ACTION_ROOT/sw/snap_action.cpp

```
static void snap prepare action(struct snap job *cjob...
    snap addr set (&mjob->in, table host addr, size table host, type table host,
              SNAP ADDRFLAG ADDR | SNAP ADDRFLAG SRC);
.../...
    snap job set(cjob, mjob, sizeof(*mjob), NULL, 0);
int main(int argc, char *argv[])
{
.../ ...
   // prepare params to be written in MMIO registers for action
   type table host = SNAP ADDRTYPE HOST DRAM;
    table host addr = (unsigned long)ibuff;
.../ ...
    snap prepare action(&cjob, &mjob,
         (void *) table host addr, size table host, type table host,
         (void *) table ddr addr, size table ddr, type table ddr,
         (void *) result addr, size result, type result);
.../...
}
```

It is important to notice that these **snap_addr_set** and **snap_job_set** APIs are not sending the data to the FPGA but just filling the data structure.

To read the result, just get it from the address you allocated and gave to the action:





\$ACTION_ROOT/sw/snap_action.cpp

SNAP management of Registers

SNAP is in-between the application and the action and its logic manages all the registers address so that it becomes transparent for the user.

For each action, SNAP gives a specific offset to add to the registers structure so that each action can have its set of registers independently.

Each register structure is then used as follows:

- The action accesses the registers as a single read/write structure.
- The application accesses the registers as 1 read only structure and 1 write only structure.

For example, when accessing the registers of an action, the SNAP application's logic writes to:

action_offset + 0x100 + register address

but will read the value at address

action_offset + 0x180 + register address.

Read and	Write are c	onsidered f	from the application /	oftware side						
act_re	g.Control	This head	This header is initialized by the SNAP job manager. The action will update the Return code and read the flags value.							
coi	VTROL	If the flags value is 0, then action sends only the action_RO_config_reg value and exit the action, otherwise it will process the action								
Simu - WR	Write@	Read@	Read@ 3 2 1			0	Typical Write value		Typical Read value	
0x3C40	0x100	0x180	sequ	sequence flags short action type						
0x3C41	0x104	0x184		Retc (return co	de 0x102/0x104)		0		0x102 - 0x104	SUCCESS/FAILURE
0x3C42	C42 0x108 0x188 Private Data						c0febabe			
0x3C43	0x10C	0x18C		Privat	te Data		deadbeef			
action	req.Data	Action specific - user defined - need to stay in 108 Bytes								
		This is the way for application and action to exchange information through this set of register								
memco	py_job_t					h this set of registers				
memco	write@					h this set of registers	Typical W	/rite value	Typica	l Read value
<i>memco</i> 0x3C44		This is the	e way for application a	nd action to exchang 2		1 1 1	Typical W	/rite value	Typica	l Read value
	Write@	This is the	e way for application a	nd action to exchang 2 snap_addr.	e information through	1 1 1	Typical W	/rite value	Туріса	l Read value
0x3C44	Write@ 0x110	This is the Read@ 0x190	e way for application a	nd action to exchang 2 snap_addr. snap_addr.a	e information through 1 addr_in (LSB)	1 1 1	Typical W	/rite value	Туріса	I Read value
0x3C44 0x3C45	0x110 0x114	This is the Read@ 0x190 0x194	e way for application a	nd action to exchang 2 snap_addr.asnap_addr.asnap_ac	e information through 1 addr_in (LSB) addr_in (MSB) ddr_in.size	1 1 1	Typical W	/rite value	Туріса	l Read value
0x3C44 0x3C45 0x3C46	Write@ 0x110 0x114 0x118	This is the Read@ 0x190 0x194 0x198	e way for application a	2 snap_addr. snap_addr. snap_addr. snap_adgr. snap_acgs (SRC, DST,)	e information through 1 addr_in (LSB) addr_in (MSB) ddr_in.size	0	Typical W	/rite value	Туріса	l Read value
0x3C44 0x3C45 0x3C46 0x3C47	Write@ 0x110 0x114 0x118 0x11C	This is the Read@ 0x190 0x194 0x198 0x19C	e way for application a	2 snap_addr snap_addr snap_addr snap_addr.a snap_ac	e information throug 1 addr_in (LSB) addr_in (MSB) ddr_in.size snap.addr_in.type (0	Typical W	/rite value	Туріса	l Read value
0x3C44 0x3C45 0x3C46 0x3C47 0x3C48	Write@ 0x110 0x114 0x118 0x11C 0x120	This is the Read@ 0x190 0x194 0x198 0x19C 0x1A0	e way for application a	and action to exchang 2 snap_addr. snap_addr.e snap_addr.e snap_addr.a gs (SRC, DST,) snap_addr.a	e information through 1 addr_in (LSB) addr_in (MSB) ddr_in.size snap.addr_in.type (iddr_out (LSB)	0	Typical W	/rite value	Туріса	l Read value

The offset 0x100 is defined in:





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. The user doesn't need to understand the offset management to read or write a register. This offset is "built-in" in SNAP and shouldn't be modified, or this will prevent access of registers





4. Control Registers

Some control registers have also been set to manage the action. These explanations are provided to be exhaustive but are not needed to program an action.

4.1 Control registers used for discovery mode

The SNAP can handle several actions. To be able to know what to manage, before the first time SNAP can work with an action, it needs first to do the inventory of all the actions of a card.

On an application point of view, the user must run almost once the command **snap_maint -v -CO** to get a list of all the actions related to the card plugged in slot CO.

The processing done is a 3 steps process which will:

- Query available action types
- Query number of actions
- Enable actions

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The resulting table is kept in memory to connect an action to an application

Index			
0	MEMCOPY	0x21	0
1	MEMCOPY	0x21	0
2	CYPHER	0x2	1
3	SEARCH	0x4	2
15			

On the action side, the code is pre-written to return the Write-only **Action_Config** registers values if the action registers have never been set (Control.flags= 0).

These **Action_Config** registers must <u>never be read by the action</u> or this will generate a discrepancy in address definition.





\$ACTION_ROOT/hw/hw_action.cpp

4.2 Control registers used for action internal logic

Xilinx HLS generates some internal logic and registers to handle the action. This logic is not accessible within the action code but will be used by SNAP and the application code to handle the action.

As for previous access, SNAP has a read and a write address to access these registers. They are defined in the *hls_action_ctrl_reg_s_axi.vhd* file generated by HLS as

```
-- -----Address Info-----
-- 0x000 : Control signals
        bit 0 - ap start (Read/Write/COH)
         bit 1 - ap_done (Read/COR)
         bit 2 - ap_idle (Read)
        bit 3 - ap_ready (Read)
__
         bit 7 - auto restart (Read/Write)
         others - reserved
-- 0x004 : Global Interrupt Enable Register
         bit 0 - Global Interrupt Enable (Read/Write)
         others - reserved
-- 0x008 : IP Interrupt Enable Register (Read/Write)
         bit 0 - Channel 0 (ap done)
         bit 1 - Channel 1 (ap ready)
         others - reserved
-- 0x00c : IP Interrupt Status Register (Read/TOW)
        bit 0 - Channel 0 (ap done)
         bit 1 - Channel 1 (ap_ready)
         others - reserved
```

That can also be represented as follow:





HLS_action_control	Control signals to start an Action and/or set/clear an IRQ Generated by HLS in => /snap/actions/hls helloworld/hw/hlsUppercasexxx/helloworld/syn/vhdl/hls action ctrl reg s axi.vhd								
Write@							Write value	Тур	ical Read value
0x00	0x00	Control signals				xxxxxxx1	Action Start	xxxxxxx2	Action done
0x04	0x04		Global Interrupt Enable Register						
0x08	0x08	IP Interrupt Enable Register (Read/Write)			xxxxxxx1	"Done" IRQ			
0x0C	0x0C		IP Interrupt Status R	Register (Read/TOW)				xxxxxxx1	Clear "Done" IRQ

These control signals are generated by the port=return instruction that has been added to the act_reg register definition.

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The application will set and manage the start of the action, the set and clear of the interrupt with specific APIs defined in "snap/sotware/lib/snap.c such as snap_action_start" or snap_action_completed