

HOMEWORK III

Due day: midnight December. 13, 2010

In this homework, you are to complete the design of a simplified pipeline CPU (at least 3 stages), referred to as mini-pipeline CPU (mPC). It could be a CPU derived from SISC taught in class or developed by you own. The most distinct feature of mPC is its basic hazard resolution capability. Your mPC must meet the specified requirements.

There will be bonus points for doing extra work for those who like challenges. Partial credits will be given based on level of completeness for bonus parts. Points for this homework are depicted as follows.

Baseline (300 points): Complete a working synthesized CPU

Bonus-I (+150 points): Complete a working synthesized CPU with a cache

Bonus-II (+250 points): Complete a working synthesized CPU with a cache and interrupt (or exception) mechanism

NOTE: Suppose you complete a working synthesized CPU with a cache, your grade will be the summation of baseline + Bonus-I, i.e., 450 points. You only need to deliver one set of code and verification. Therefore, you shall specify in your report what you have achieved.

General rules for deliverables

- This homework needs to be completed by INDIVIDUAL student.
- Compress all files described in the problem statements into one zip or rar.
- Submit the compress file to the course website before the due day. **Warning!**
AVOID submit in the last minute. Late submission is not accepted.

Grading Notes

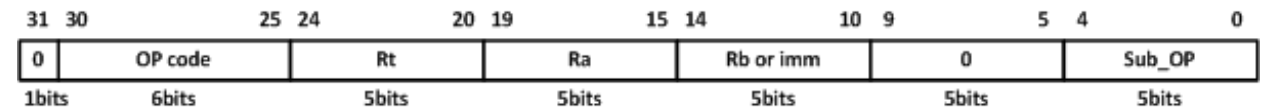
- **Important!** DO remember to include your Verilog code. NO code, NO grades. Also, if your code can not be recompiled by TA successfully using tools in SoC Lab, you receive NO credit.
- Write your report seriously and professionally. Incomplete description and information will reduce your chances to get more credits.
- If extra works (like synthesis, post-simulation or additional instructions) are done, please describe them in your final report clearly for bonus points.
- Please follow course policy.

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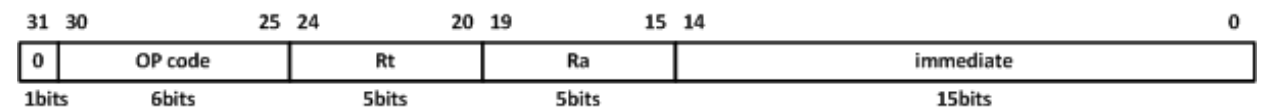
Exercise

mPC has the following instruction format basically

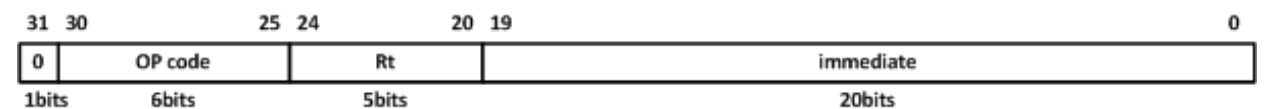
☞ Data-processing (SE:sign-extended, ZE:zero-extended)



OP code	Mnemonics	DST	SRC1	SRC2	Sub OP	Description
100000	NOP	00000	00000	00000	01001	No operation
100000	ADD	\$Rt	\$Ra	\$Rb	00000	$Rt = Ra + Rb$
100000	SUB	\$Rt	\$Ra	\$Rb	00001	$Rt = Ra - Rb$
100000	AND	\$Rt	\$Ra	\$Rb	00010	$Rt = Ra \& Rb$
100000	OR	\$Rt	\$Ra	\$Rb	00100	$Rt = Ra Rb$
100000	XOR	\$Rt	\$Ra	\$Rb	00011	$Rt = Ra \wedge Rb$
100000	SRLI	\$Rt	\$Ra	imm	01001	Shift right : $Rt = Ra \gg imm$
100000	SLLI	\$Rt	\$Ra	imm	01000	Shift left : $Rt = Ra \ll imm$
100000	ROTRI	\$Rt	\$Ra	imm	01011	Rotate right : $Rt = Ra \gg imm$



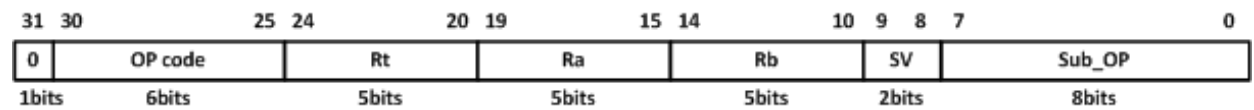
OP code	Mnemonics	DST	SRC1	SRC2	Description
101000	ADDI	\$Rt	\$Ra	imm	$Rt = Ra + SE(imm)$
101100	ORI	\$Rt	\$Ra	imm	$Rt = Ra ZE(imm)$
101011	XORI	\$Rt	\$Ra	imm	$Rt = Ra \wedge ZE(imm)$
000010	LWI	\$Rt	\$Ra	imm	$Rt = Mem[Ra + (imm \ll 2)]$
001010	SWI	\$Rt	\$Ra	imm	$Mem[Ra + (imm \ll 2)] = Rt$



OP code	Mnemonics	DST	SRC1	Description
100010	MOVI	\$Rt	imm	$Rt = SE(immediate)$

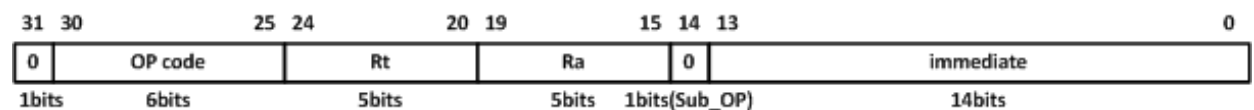
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☞ Load and store



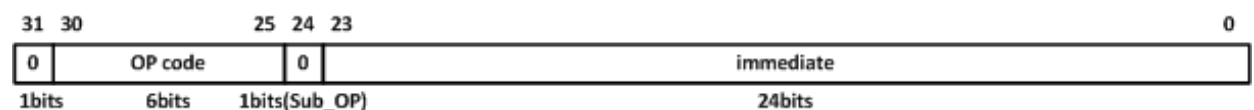
OP code	Mnemonics	DST	SRC1	SRC2	Sub OP	Description
011100	LW	\$Rt	\$Ra	\$Rb	00000010	$Rt = Mem[Ra + (Rb \ll sv)]$
011100	SW	\$Rt	\$Ra	\$Rb	00001010	$Mem[Ra + (Rb \ll sv)] = Rt$

☞ Branch (SE:sign-extended)



OP code	Mnemonics	DST	SRC1	Sub OP	Description
100110	BEQ	\$Rt	\$Ra	0	if($Rt == Ra$) $PC = PC + SE(imm \ll 1)$ else $PC = PC + 4$

☞ Jump (SE:sign-extended)



OP code	Mnemonics	Sub OP	Description
100100	J	0	$PC = PC + SE(imm \ll 1)$

Baseline (300 points) Design a pipeline CPU (either extending the SISC as the baseline design or developing by your own) with the following features:

- Nineteen instructions equivalent to above basically
- Bypassing mechanism to resolve the interlock problem
- Data and instruction memory has zero delay
- No need to synthesize the data and instruction memory
- Revise the example of the constraint file given in HWII(Part2) Problem1 to do synthesis (Note that the output load shall not be changed!!)
- Verification
 - Perform verification for individual instruction
 - Develop 5 distinct testbenches for hazard problems
 - Read a data array, whose size is 10, from data memory and restore them back with the sorted order in the same data array. Please give 3 examples to show that the machine code is working fine.

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- iv. Design a Fibonacci sequence generator that can give at least first 10, i.e., $F_0 \dots F_{10}$. In math term, $F_n = F_{n-1} + F_{n-2}$, where $F_0=0$ and $F_1=1$.

Bonus I (+150 points) Extend the synthesizable pipeline CPU of HW3 Problem1 to equip with a cache. Note that the same verification in HW3 Problem 1 shall be applied and make them work.

Bonus II (+250 points) Extend the pipeline CPU to equip with a cache and interrupt. Note that the same verification in HW3 Problem 1 shall be applied and make them work.

Deliverables

1. All CPU Verilog codes including components, testbenches and machine code for each lab exercise.
2. A homework report that includes
 - a. A summary in the beginning to state what has been done (such as baseline mPC, synthesis, post-synthesis simulation)
 - b. A block diagram for your completed mPC indicating all necessary components and I/O pins
 - c. Instruction set, instruction format, and addressing modes
 - d. Simulated waveforms for at least required test benches with proper explanation
 - e. Report your maximum simulation speed
 - f. Synthesis reports and log file
 - g. Verification of post-synthesis simulation waveforms with proper explanation
 - h. Conclusion

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//////////////////////////////////// Reference //////////////////////////////////////

☞ SMILE CPU has the following instruction format (Andes ISA)

The Andes 32bit instruction formats and the meaning of each filed are described below:

➤ Type-0 Instruction Format

0	Opc_6	{sub_1, imm_24}
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➤ Type-1 Instruction Format

0	Opc_6	rt_5	imm_20	
0	Opc_6	rt_5	sub_4	imm_16

➤ Type-2 Instruction Format

0	Opc_6	rt_5	ra_5	imm_15
0	Opc_6	rt_5	ra_5	{sub_1, imm_14}

➤ Type-3 Instruction Format

0	Opc_6	rt_5	ra_5	rb_5	sub_10
0	Opc_6	rt_5	ra_5	imm_5	sub_10

➤ Type-4 Instruction Format

0	Opc_6	rt_5	ra_5	rb_5	rd_5	sub_5
0	Opc_6	rt_5	ra_5	imm1_5	imm2_5	sub_5

- ◆ opc_6: 6-bit opcode
- ◆ rt_5: target register in 5-bit index register set
- ◆ ra_5: source register in 5-bit index register set
- ◆ rb_5: source register in 5-bit index register set
- ◆ rd_5: destination register in 5-bit index register set
- ◆ sub_10: 10-bit sub-opcode
- ◆ sub_5: 5-bit sub-opcode
- ◆ sub_4: 4-bit sub-opcode
- ◆ sub_1: 1-bit sub-opcode
- ◆ imm_24: 24-bit immediate value, for unconditional jump instructions (J, JAL). The immediate value is used as the lower 24-bit offset of same 32MB memory block (new
- ◆ PC[31:0] = {current PC[31:25], imm_24, 1'b0}
- ◆ imm_20: 20-bit immediate value. Sign-extended to 32-bit for MOVI operations.
- ◆ imm_16: signed PC relative address displacement for branch instructions.
- ◆ imm_15: 15-bit immediate value. Zero extended to 32-bit for unsigned operations, while sign extended to 32-bit for signed operations.
- ◆ imm_14: signed PC relative address displacement for branch instructions.
- ◆ imm_5, imm1_5, imm2_5: 5-bit unsigned count value or index value

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