

# Privacy Proofs for OpenDP: Partition Map (Measurement)

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## 1 Algorithm Implementation

### 1.1 Code in Rust

The current OpenDP library contains the `make_map_partition_meas` function constructing a `Measurement` for partitioned data based on a list of `Measurements` (parallel composition). This is defined in lines 76-136 of the file `mod.rs` in the Git repository [https://github.com/opensdp/opensdp/blob/449222066a006a19d15bb68e62a87105ce49eb15/rust/src/comb/partition\\_map/mod.rs#L76-L136](https://github.com/opensdp/opensdp/blob/449222066a006a19d15bb68e62a87105ce49eb15/rust/src/comb/partition_map/mod.rs#L76-L136).

### 1.2 Pseudo Code in Python

#### Preconditions

To ensure the correctness of the output, we require the following preconditions:

- User-specified types:
  - Variable `measurements` must be a vector of elements of class `Measurement`.<sup>1</sup>
  - `output_measure::Distance` must have trait `TotalOrd`.
  - `input_domain`: must have trait `Domain`.
  - `output_domain`: must have trait `Domain`.
  - `input_metric`: must have trait `Metric`.
  - `output_measure`: must have trait `Measure`.

#### Postconditions

- A `Measurement` is returned (i.e., if a `Measurement` cannot be returned successfully, then an error should be returned).

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<sup>1</sup>To be defined in the pseudocode defs doc.

## Pseudo Code

```

1 def make_map_partition_meas(measurements: List[Measurement]) -> Measurement
2 :
3   input_domain = ProductDomain(m.input_domain for m in measurements)
4   output_domain = ProductDomain(m.output_domain for m in measurements)
5   def function(data: Vec<DI::Carrier>) -> Vec<DO::Carrier>:
6     output = []
7     for part, m in zip(data, measurements):
8       output.append(m.function(part))
9   input_metric = ProductMetric(measurements[0].input_metric)
10  output_measure = measurements[0].output_measure
11  def privacy_map(d_in: input_metric::Distance) -> output_measure::
12    Distance:
13    return max(m.map(d_in) for m in measurements)
14  return Measurement(input_domain, output_domain, function, input_metric,
15    output_measure, privacy_map)

```

## 2 Proof

The necessary definitions for the proof can be found at [“List of definitions used in the proofs”](#).

**Theorem 2.1.** *For every setting of the input parameter `measurements` to `make_map_partition_meas` such that the given preconditions hold, `make_map_partition_meas` raises an exception (at compile time or runtime) or returns a valid `Measurement` with the following privacy guarantee:*

1. **(Domain-metric compatibility.)** *The domain `input_domain` matches one of the possible domains listed in the definition of `input_metric`.*
2. **(Privacy guarantee.)** *For every pair of elements  $v, w$  in `input_domain` and for any `d_in`, where `d_in` has the associated type for `input_metric`, if  $v, w$  are `d_in`-close under `input_metric`, then `function(v)`, `function(w)` are `privacy_map(d_in)`-close under `output_measure`.*

*Proof.*

1. **(Domain-metric compatibility.)** The `input_domain` of `make_map_partition_meas` is `ProductDomain` and the `input_metric` is `ProductMetric`. Since each component `m` of the variable `measurements` is a `Measurement`, `m.input_domain` matches one of the possible domains listed in the definition of `m.input_metric`. Therefore, the product of `m.input_domain` is compatible with the product of `m.input_metric`.
2. **(Privacy guarantee.)** Let  $m_i$  be the  $i$ th element of the variable `measurements`, thus  $m_i$  is a `Measurement`. Let `MI` and `MO` denote the `input_metric` and the `output_measure` of  $m_i$ , respectively. For any  $i$  and  $v_i, w_i$  in the input domain of  $m_i$ ,

$$d_{MI}(v_i, w_i) \leq d\_in$$

implies

$$d_{MO}(f_i(v_i), f_i(w_i)) \leq m_i.privacy\_map(d\_in),$$

where  $f_i$  denotes the function in  $\mathbf{m}_i$ , and  $v_i$  and  $w_i$  are the  $i$ th component of  $v$  and  $w$ , respectively. Let  $d_{\mathbf{PM}, \mathbf{MI}}$  denote the distance under `ProductMetric(MI)`. Then by Definition 1(i),

$$d_{\mathbf{PM}, \mathbf{MI}}(v, w) = \sum_i d_{\mathbf{MI}}(v_i, w_i).$$

Note that, given  $d_{\mathbf{PM}, \mathbf{MI}}(v, w) \leq \mathbf{d\_in}$ , by the nature of partitioned data, all  $d_{\mathbf{MI}}(v_i, w_i)$  are zeros except for at most one. We then have

$$\begin{aligned} d_{\mathbf{MO}}(g(v), g(w)) &\leq \max_i(d_{\mathbf{MO}}(f_i(v_i), f_i(w_i))) \\ &\leq \max_i(\mathbf{m}_i.\text{privacy\_map}(\mathbf{d\_in})), \end{aligned}$$

where  $g$  is the function in `make_map_partition_meas`.

Therefore, we have shown that  $g(v)$  and  $g(w)$  are `privacy_map(d_in)`-close if  $v$  and  $w$  are `d_in`-close.

□

**Definition 1** (Distance under `ProductMetric`). Let  $d_{\mathbf{PM}, M}$  denote the distance under `ProductMetric(M)` where  $M$  is a valid metric. Then  $d_{\mathbf{PM}, M}$  is defined as the sum of distance under each  $M$ . Specifically, for any  $v, w$  in the input domain and  $v_i, w_i$  denote their  $i$ th entry, respectively,

(i) for input metric  $\mathbf{MI}$ ,

$$d_{\mathbf{PM}, \mathbf{MI}}(v, w) = \sum_i d_{\mathbf{MI}}(v_i, w_i).$$

(ii) for output metric  $\mathbf{MO}$ ,

$$d_{\mathbf{PM}, \mathbf{MO}}(g(v), g(w)) = \sum_i d_{\mathbf{MO}}(f_i(v_i), f_i(w_i)),$$

where  $g$  and  $f_i$  denote the function in their corresponding *Transformation*.