Privacy Proofs for OpenDP: Bounded Sum with Unknown n

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Contents

| 1 | Algorithm Implementation | 1 |
|---|--------------------------|---|
| | 1.1 Code in Rust | 1 |
| | 1.2 Pseudocode in Python | 2 |
| 2 | Proof | 2 |
| | 2.1 Symmetric Distance | 2 |

1 Algorithm Implementation

1.1 Code in Rust

The current OpenDP library contains the transformation $make_bounded_sum_n$ implementing bounded sum unknown n. This is defined in lines 18-33 of the file sum_rs in the Git repository (https://github.com/opendp/opendp/blob/53e8d67b8dde4425930fb8bc $397c126cd4f18370/rust/opendp/src/trans/sum_rs\#L18-L33$).

¹As of July 20, 2021.

1.2 Pseudocode in Python

We present a simplified Python-like pseudocode of the Rust implementation below. The necessary definitions for the pseudocode can be found at "List of definitions used in the pseudocode".

Preconditions

To ensure the correctness of the output, we require the following preconditions:

• User-specified types:

- Type T must implement DistanceConstant(IntDistance), TotalOrd,² Abs,
 Sub(Output=T), SaturatingAdd, and Zero.
- IntDistance must have trait InfCast(T). Question: Same point that Connor raised in make_count this is not needed for the proof.

Postconditions

• Either a valid Transformation is returned or an error is returned.

```
def MakeBoundedSum(L: T, U: T):
      input_domain = VectorDomain(IntervalDomain(L, U))
      output_domain = AllDomain(T)
3
      input_metric = SymmetricDistance()
4
      output_metric = AbsoluteDistance(T)
      def relation(d_in: u32, d_out: T) -> bool:
          return d_out >= inf_cast(d_in, T) * max(abs(U), abs(L))
9
      def function(data: Vec[T]) -> T:
10
          result:T = 0
          for i in data:
12
              result = saturating_add(result, i)
13
          return result
14
      return Transformation(input_domain, output_domain, function,
      input_metric, output_metric, stability_relation = relation)
```

2 Proof

2.1 Symmetric Distance

Theorem 1. For every setting of the input parameters (L, U) to MakeBoundedSum, the transformation returned by MakeBoundedSum has the following properties:

1. (Appropriate output domain). For every vector v in the input domain, function(v) is in the output domain.

²For now, the OpenDP library only implements PartialOrd, but TotalOrd will soon be implemented. Then, TotalOrd will be redundant, since the trait TotalOrd is part of the trait DistanceConstant.

- 2. (Domain-metric compatibility). The domain input_domain matches one of the possible domains listed in the definition of input_metric, and likewise output_domain matches one of the possible domains listed in the definition of output_metric.
- 3. (Stability guarantee). For every pair of elements v, w in $input_domain$ and for every pair (d_in, d_out) , where d_in is of the associated type for $input_metric$ and d_out is the associated type for $output_metric$, if v, w are d_in -close under $input_metric$ and $relation(d_in, d_out) = True$, then function(v), function(v) are d_out -close under $output_metric$.

Proof. (Appropriate output domain). In the case of MakeBoundedSum, this corresponds to showing that for every vector v in VectorDomain(IntervalDomain(L, U)), where L and U have type T, the element function(v) belongs to AllDomain(T). The type signature of function as defined in line 10 automatically enforces that function(v) has type T. Since the Rust code successfully compiles, by the type signature the appropriate output domain property must hold. Otherwise, the code will raise an exception for incorrect input type. It is also necessary to check that function(v) is contained within the interval [get_min_value(T), get_max_value(T)]. This is enforced by the use of the function saturating_add in line 13, as described in "List of definitions used in the pseudocode".

If the sum of all the vector elements in data is greater than get_max_value(T), then result will be equal to get_max_value(T). If the sum of all the vector elements in data is less than get_min_value(T), then result will be equal to get_min_value(T). Otherwise, result will be equal to the sum of all the vector elements in data, and it will be contained within the interval [get_min_value(T), get_max_value(T)]. Therefore, function(v) is guaranteed to be in output_domain in all cases.

(Domain-metric compatibility). For MakeBoundedSum, this corresponds to showing that VectorDomain(IntervalDomain (L, U)) is compatible with symmetric distance, and that AllDomain(T) is compatible with absolute distance. Both follow directly from the definition of symmetric distance and absolute distance, as stated in "List of definitions used in the pseudocode", along with the appropriate output domain property shown above, which ensures that output_domain is indeed AllDomain(T).

(Stability guarantee). Throughout the stability guarantee proof, we can assume that function(v) and function(w) are in the correct output domain, by the appropriate output domain property shown above.

Since by assumption $\operatorname{relation}(\operatorname{d_in},\operatorname{d_out}) = \operatorname{True}$, by the MakeBoundedSum stability relation (as defined in line 7 in the pseudocode), we have that $\operatorname{d_out} \ge \operatorname{d_in} \cdot \max(|\operatorname{U}|,|\operatorname{L}|)$. Moreover, v,w are assumed to be $\operatorname{d_in_close}$. By the definition of the symmetric difference metric, this is equivalent to stating that $d_{Sym}(v,w) = |\operatorname{MultiSet}(v)\Delta\operatorname{MultiSet}(w)| \le \operatorname{d_in}$.

Further, applying the histogram notation,³ it follows that

$$d_{Sym}(v,w) = \|h_v - h_w\|_1 = \sum_z |h_v(z) - h_w(z)| \leq \texttt{d_in}.$$

We want to show that

$$d_{Abs}(\texttt{function}(v), \texttt{function}(w)) \leq d_{Sym}(v, w) \cdot \max(|\mathtt{U}|, |\mathtt{L}|).$$

³See A Programming Framework for OpenDP, footnote 1 in page 3. Note that there is a bijection between multisets and histograms, which is why the proof can be carried out with either notion. For further details, please consult https://www.overleaf.com/project/60d214e390b337703d200982.

This would imply that

 $d_{Abs}(\texttt{function}(v), \texttt{function}(w)) \leq d_{Sym}(v, w) \cdot \max(|\mathtt{U}|, |\mathtt{L}|) \leq \mathtt{d_{-in}} \cdot \max(|\mathtt{U}|, |\mathtt{L}|),$ and by the stability relation this will imply that

$$d_{Abs}(function(v), function(w)) \leq d_{out},$$

as we want to see.

Let u denote the vector formed by all the elements of v and w without multiplicities (i.e., u contains exactly once each of the elements in MultiSet $(v) \cup$ MultiSet(w), in any order). Let u_i denote the i-th element of u, and similarly for v and w, and let m denote the length of u. Then, by definition,

$$d_{Sym}(v, w) = \sum_{z} |h_v(z) - h_w(z)| = \sum_{i} |h_v(u_i) - h_w(u_i)|;$$

 $d_{Abs}(\texttt{function}(v), \texttt{function}(w)) = |\texttt{function}(v) - \texttt{function}(w)| \leq \Big|\sum_i v_i - \sum_i w_i\Big| = \sum_i v_i - \sum_$

$$= \Big| \sum_{i} u_i \cdot h_v(u_i) - \sum_{i} u_i \cdot h_w(u_i) \Big| = \Big| \sum_{i} u_i \cdot (h_v(u_i) - h_w(u_i)) \Big|.$$

We remark that we have the inequality $| \mathtt{function}(v) - \mathtt{function}(w) | \leq |\sum_i v_i - \sum_i w_i|$ above (instead of an equality) due to the definition of $\mathtt{saturating_add}$. The equality case holds whenever $\sum_i v_i \in [\mathtt{get_min_value}(\mathtt{T})$, $\mathtt{get_max_value}(\mathtt{T})]$ and $\sum_i w_i \in [\mathtt{get_min_value}(\mathtt{T})$, $\mathtt{get_max_value}(\mathtt{T})]$. In any of the possible cases where $\sum_i v_i > \mathtt{get_max_value}(\mathtt{T})$ or $\sum_i v_i < \mathtt{get_min_value}(\mathtt{T})$ and $\sum_i w_i > \mathtt{get_max_value}(\mathtt{T})$ or $\sum_i w_i < \mathtt{get_min_value}(\mathtt{T})$, the difference $|\sum_i v_i - \sum_i w_i|$ will always upper bound the value $|\mathtt{function}(v) - \mathtt{function}(w)|$, and hence it is sufficient to carry our proof by only considering the quantity $|\sum_i v_i - \sum_i w_i|$.

(silvia) Might be easier to just state this in terms of absolute clamp and refer to that transformation proof.

By the definition of absolute distance and symmetric distance, and by applying the triangle inequality, we obtain:

$$d_{Abs}(\texttt{function}(v),\texttt{function}(w)) \leq \Big|\sum_i u_i \cdot (h_v(u_i) - h_w(u_i))\Big| \leq |u_i| \cdot \sum_i |h_v(u_i) - h_w(u_i)|.$$

By the appropriate output domain property $u_i \in [L, U] \ \forall i \ \text{it follows that} \ |u_i| \leq \max(|U|, |L|)$ for all i. Hence,

$$d_{Abs}(\mathtt{function}(v),\mathtt{function}(w)) \leq |u_i| \cdot \sum_i |h_v(u_i) - h_w(u_i)| \leq$$

$$\leq \max\left(|\mathbf{U}|,|\mathbf{L}|\right) \cdot \sum_{i} |h_v(u_i) - h_w(u_i)| \leq \max\left(|\mathbf{U}|,|\mathbf{L}|\right) \cdot d_{Sym}(v,w).$$

Lastly, since by assumption v and w are d_in-close, by the defined relation(d_in, d_out) (line 10 in the pseudocode) it follows that

$$d_{Abs}(\texttt{function}(v),\texttt{function}(w)) \leq \max\left(|\mathtt{U}|,|\mathtt{L}|\right) \cdot d_{Sym}(v,w) \leq$$

$$\leq \max(|U|, |L|) \cdot d_{in} \leq d_{out},$$

as we wanted to show.

(silvia) Flag: need to account for rounding errors in the stability relation given the non-closure of float addition, as discussed on the week of August 2nd. We are figuring this out.