

Quantitative Types

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Introduction

Dependent types can refer to **values**

⇒ Enable advanced verification

e.g. List with length

```
data List : Type → Nat → Type where
```

```
  . . .
```

```
[1, 2] : List Nat 2
```

```
head []
```

Type error!

Introduction

Dependently-typed languages easily lead to inefficient programs

```
length : {n : Nat} -> List a n -> Nat
```

Is **n** is used only for the type definition?
Or used for the calculation?

Introduction : Quantitative Types

All variables have **multiplicity** $(0, 1, \omega)$

Multiplicity indicates how often the variable is used

- 0 : Unused at Runtime \Rightarrow Can be Erased
 - 1 : Used exactly once
 - ω : Unrestricted
- \Rightarrow Needed at runtime

Outline

- Background : A Role of Types
- Background : Dependent Types
- Quantitative Types
- Application Example of Quantitative Types

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Background : Type Systems

Incorrect programs cause bugs

```
print(1 + true)
```



Bug

Type systems can detect bugs before execution

```
print(1 + true)
```



type checker

`1 + true` is wrong!

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- Background : A Role of Types
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Dependent Types

Feature :

Dependent types can **refer to values** as well as types

⇒ Allows **advanced verification** and **automatic proof**

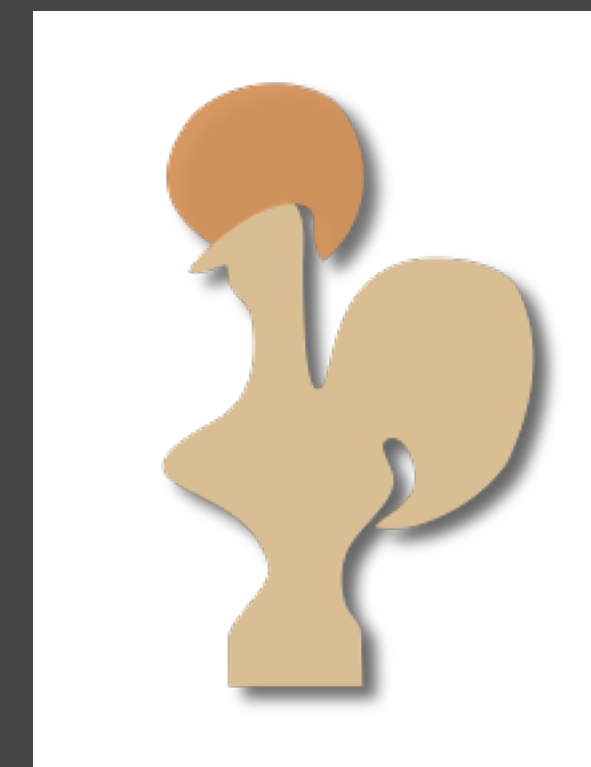
Dependently-typed languages :



Agda



Idris



Coq

Referring to Types (generics, type parameter)

In plain language, types can refer to types

e.g. Type of List in Scala

Type List receives a type

```
trait List[T]  
case class Nil[T] extends List[T]  
case class Cons[T](head : T, tail : List[T]) extends List[T]  
  
[true] : List[Bool]  
[1, 2] : List[Nat]
```

Limit of Verification in Plain Type Systems

The function **head** returns the top element of given list

```
head [1, 2, 3] = 1  
head [] = ?
```

head [] type checks but causes runtime error

head []  type checker   Bug

Hmm, OK!

Dependent Types

Dependent types can also refer to **values**

e.g. Type of List with length by Idris

```
data List : Type -> Nat -> Type where
  . . .

[]      : List a 0
[true]  : List Bool 1
[1, 2]  : List Nat 2
```

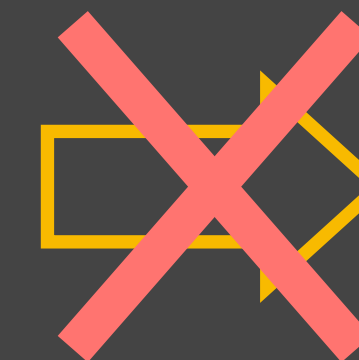
Advanced Verification

Recall : In plain language(like scala),

`head []` causes runtime error

```
safe_head : List a (1 + n) -> a  
[] : List a 0
```

`safe-head []`  type checker



It's wrong!

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Dependent Types lead to Inefficient code

Full Explicit Definition of `safe-head`

```
safe-head : {n : Nat} -> {a : Type} ->  
           List a (1 + n) -> a  
safe-head {n} {a} (x :: xs) = x
```

Too many arguments

`n` and `a` are needed only for type definition and
not relevant to the head calculation

Erasure of type information

```
safe-head : {n : Nat} -> {a : Type} ->  
           List a (1 + n) -> a  
safe-head {n} {a} (x :: xs) = x
```

⇒ Type Checker OK!

⇒ Generate more efficient code

```
safe-head (x :: xs) = x
```

No longer need **n** and **a** ⇒ Erased

Problem

Difficulty : Distinguish between
necessary and unnecessary variables.

```
length : {n : Nat} -> {a : Type} ->  
        List a n -> a  
length {n} {a} xs = n
```

n is used in the calculation of length

as well as type definition  Cannot erase **{n}**

Quantitative Type System

Feature : Trace variable usage

All variables have **multiplicity** $(0, 1, \omega)$

Multiplicity indicates how often the variable is used

- 0 : Unused at Runtime \Rightarrow Can be Erased
 - 1 : Used exactly once
 - ω : Unrestricted
- \Rightarrow Needed at runtime

Idris2 : Language with Quantative Types

Multiplicity of variables can be specified
at function definitions

a is not used

x is used exactly once

```
id : {0 a : Type} -> (1 x : a) -> a
id {a} x = x
```

Multiplicity 0 : Unused at Runtime

Variable with multiplicity 0
can only be used freely in type definition

```
safe-head : {0 n : Nat} -> {0 a : Type} ->  
           List a (1 + n) -> a
```

```
safe-head {n} {a} (x :: xs) = x
```

Multiplicity 1 : Linear Usage

Variable with multiplicity 1
must be used exactly once in function body.

```
idNat : (1 x : Nat) -> Nat  
idNat x = x
```

Multiplicity ω : Unrestricted Usage

Variable with multiplicity ω
can be used freely anywhere.

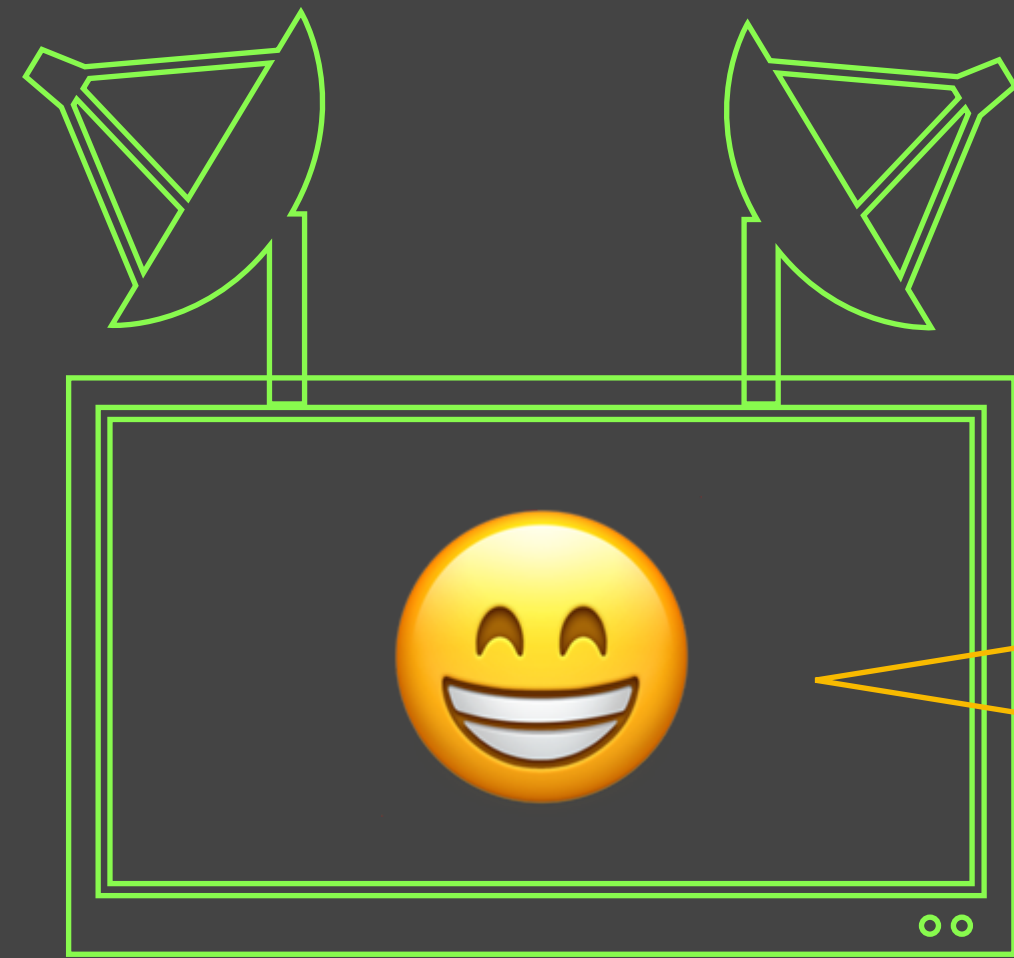
No annotation \Rightarrow Multiplicity ω

```
length : {n : Nat} -> List a n -> Nat
length {n} _ = n
```

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Application Example : Communication



Channel-kun

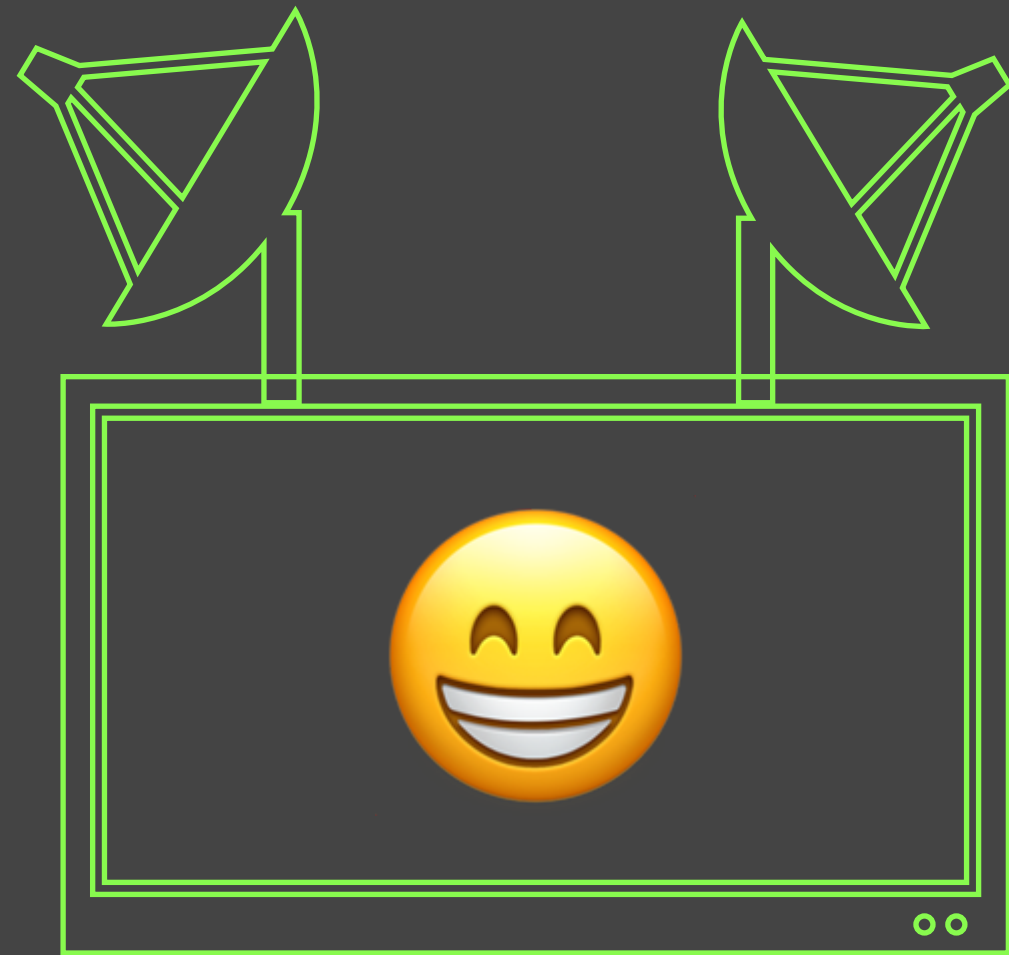
Hello client! I'm a channel.
You can **send one number** to the server
through me,
and **then receive the string** returned by
the server. OK?



OK. Write his action as

{Send Int ⇒ Recv Str ⇒ Close}

Dependently-Typed Channel



```
data Action : Type where
  Send  : Type -> (next : Action) -> Action
  Recv  : Type -> (next : Action) -> Action
  Close : Action
```

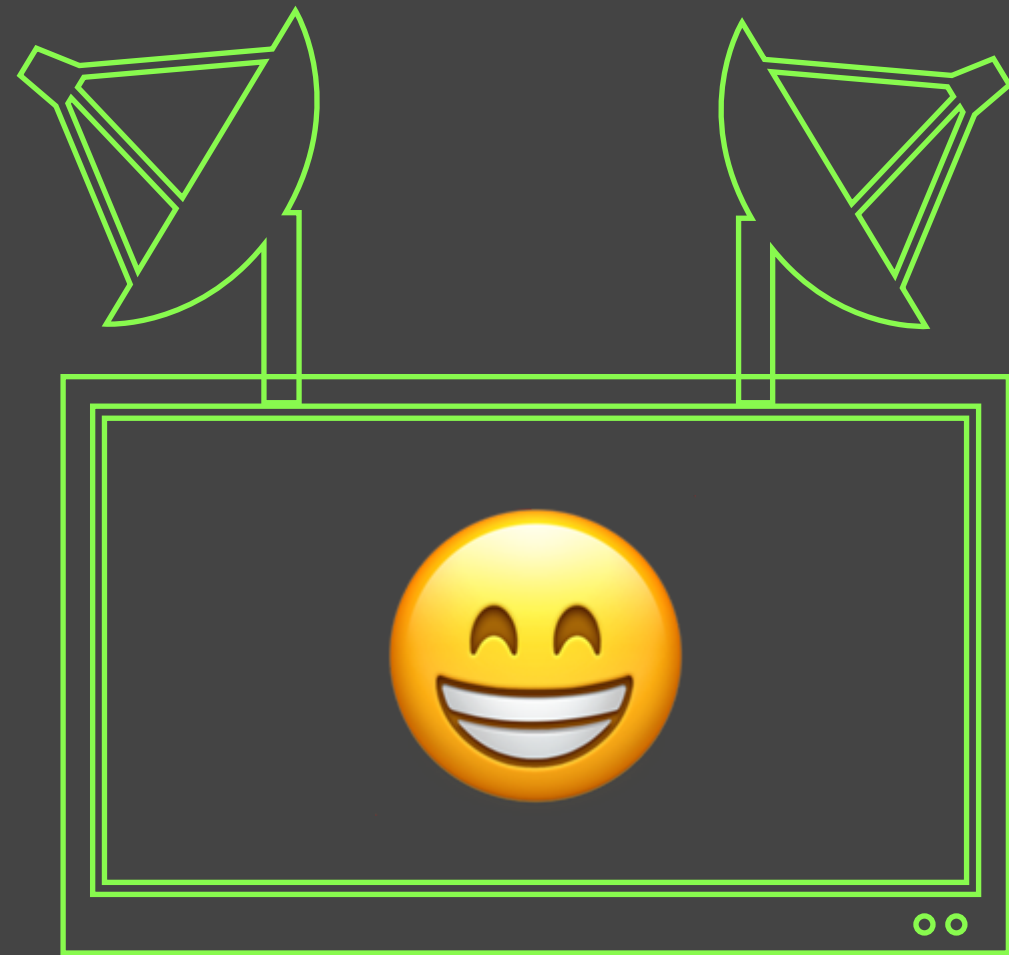
```
data Channel : Actions -> Type
```



His type is

```
Channel (Send Int (Recv Str (Close)))
```

Type-safe Channel Operations



```
send  : (1 chan : Channel (Send ty next))  
      -> (val : ty) -> Channel next
```

```
recv  : (1 chan : Channel (Recv ty next))  
      -> (ty , Channel next)
```

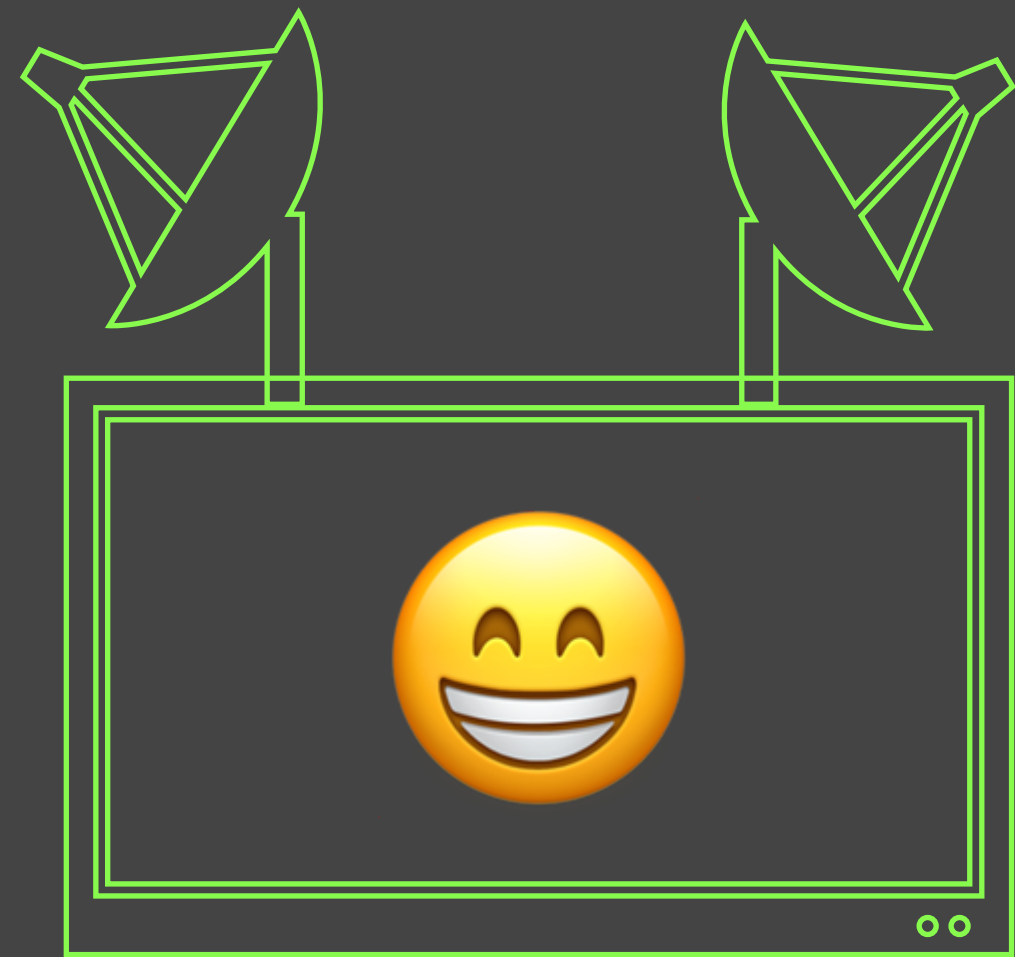
```
close : (1 chan : Channel Close) -> ()
```



Apply send , recv , close
to him in that order

Type safe operations

Dependent types can detect wrong manipulation.



```
send 10155 channel_kun
```



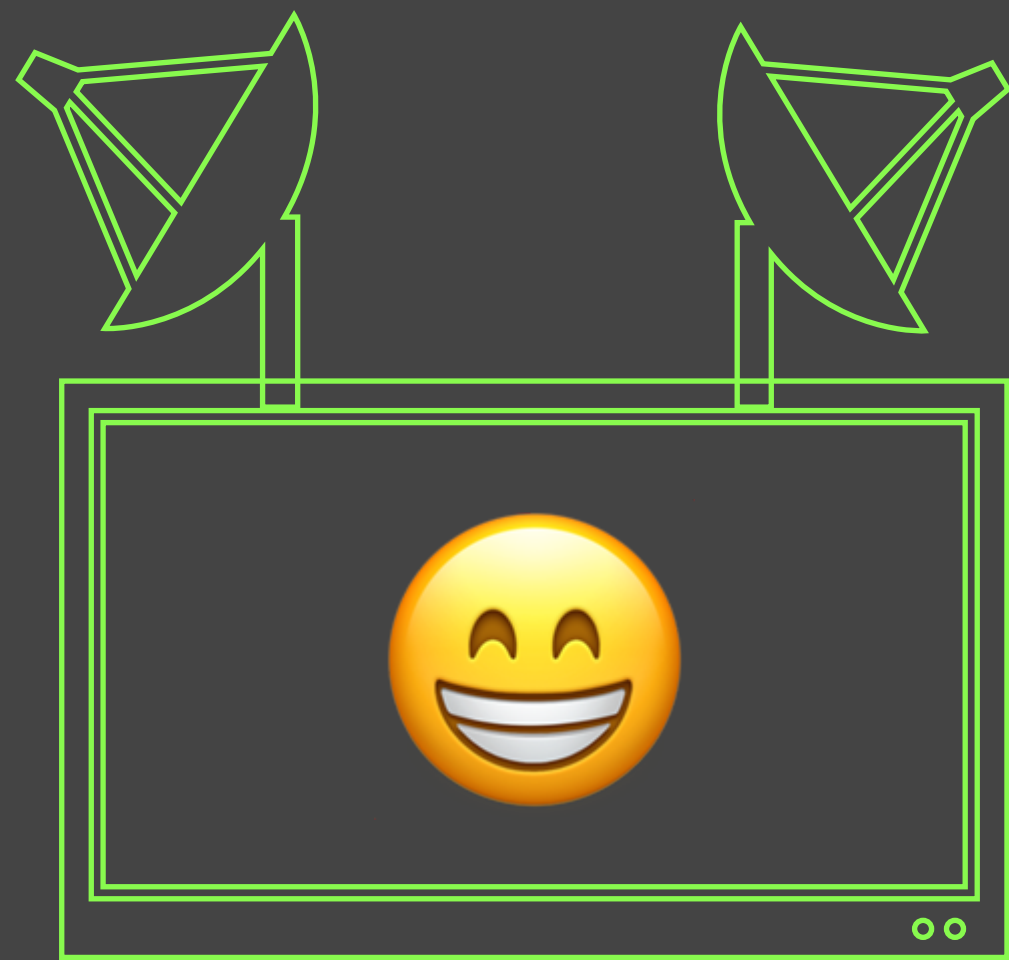
```
recv channel_kun
```



```
close channel_kun
```

{Send Int ⇒
Recv Str ⇒ Close}

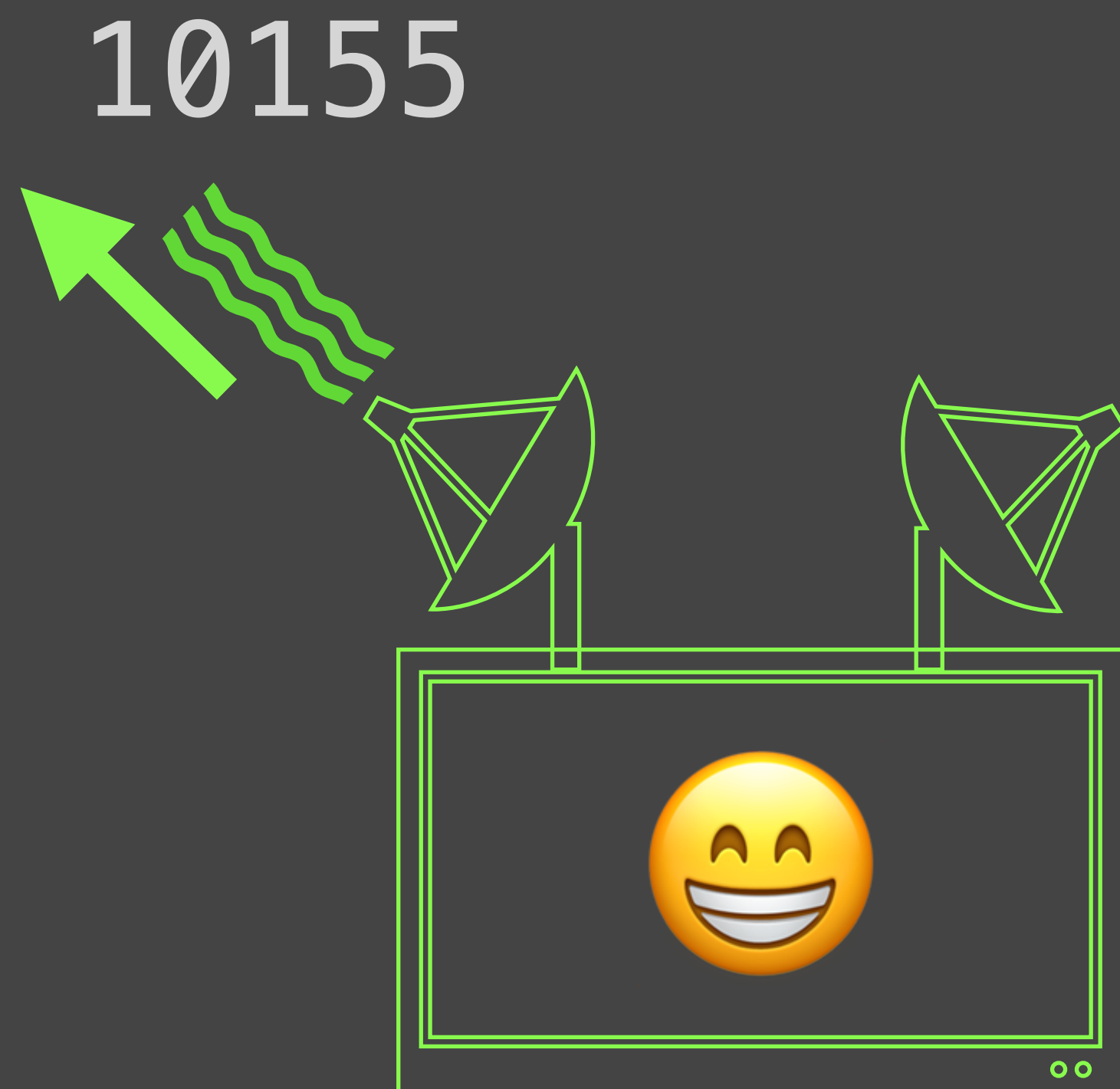
Correct Program



```
let ch1 = send 10155 channel_kun in  
let v , ch2 = recv ch1 in  
close ch2
```

{Send Int \Rightarrow
Recv Str \Rightarrow Close}

Correct Program



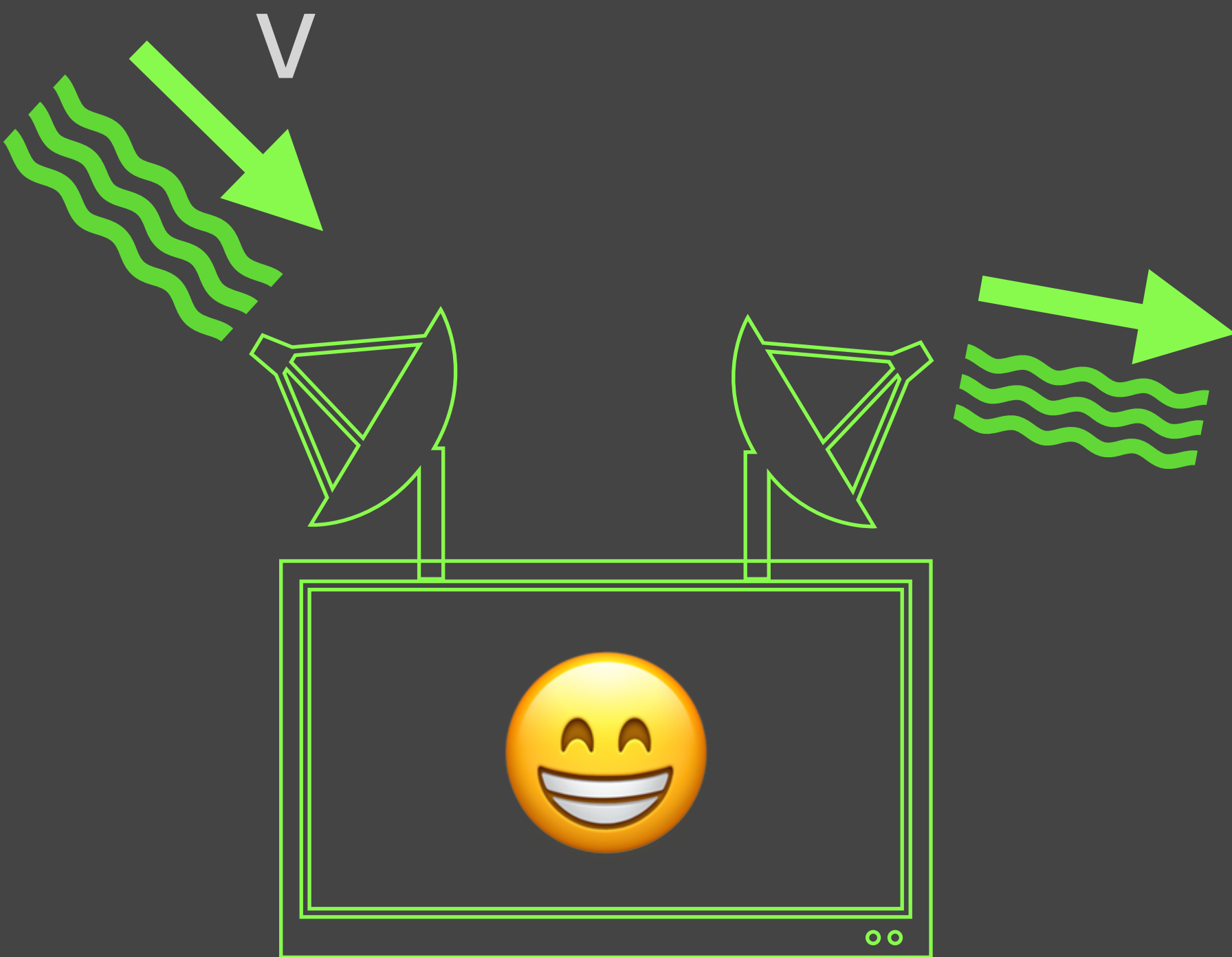
Send number 10155 to server

```
let ch1 = send 10155 channel_kun in
```

Updated channel_kun
with action {Recv Str ⇒ Close}

{Send Int ⇒
Recv Str ⇒ Close}

Correct Program



```
let ch1 = send 10155 channel_kun in
```

Recieve a value from ch1

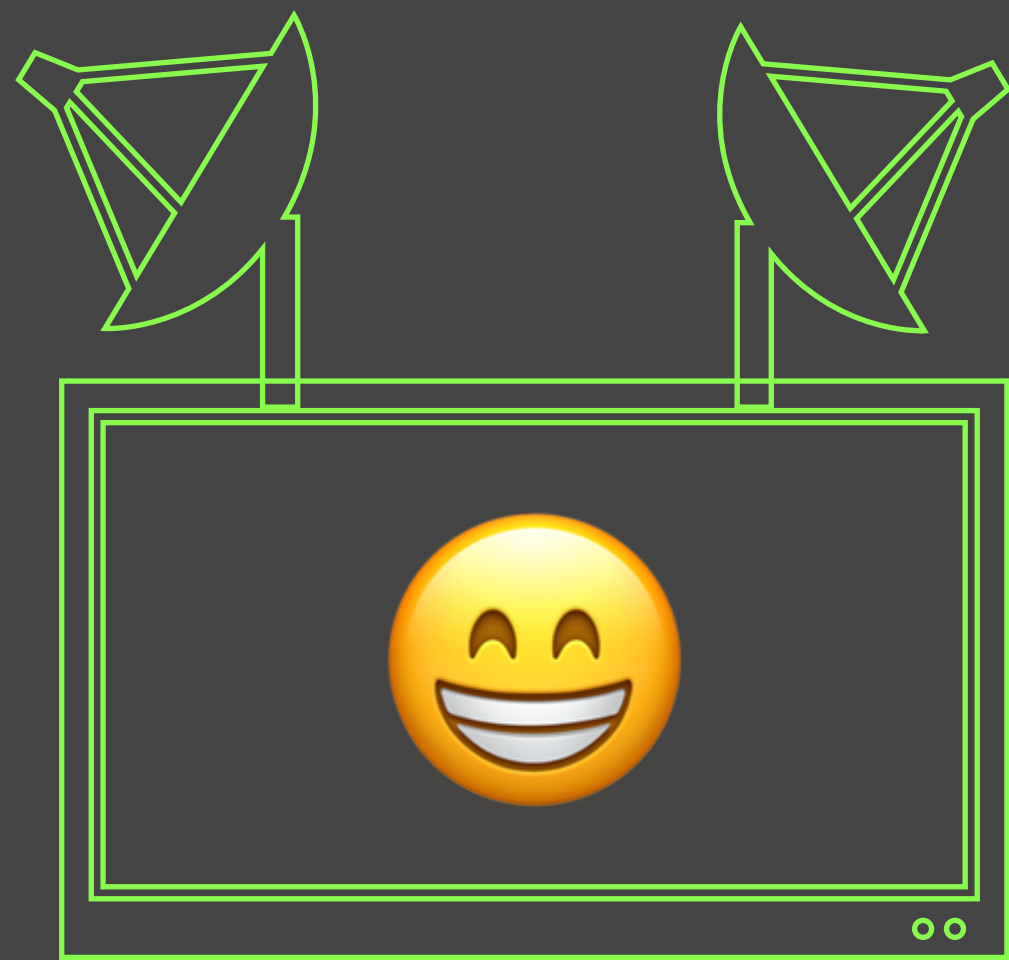
```
let v , ch2 = recv ch1 in
```

Updated channel_kun
with action {Close}

{Send Int ⇒
Recv Str ⇒ Close}

Correct Program

```
let ch1 = send 10155 channel_kun in  
let v , ch2 = recv ch' in
```



Terminate Connection

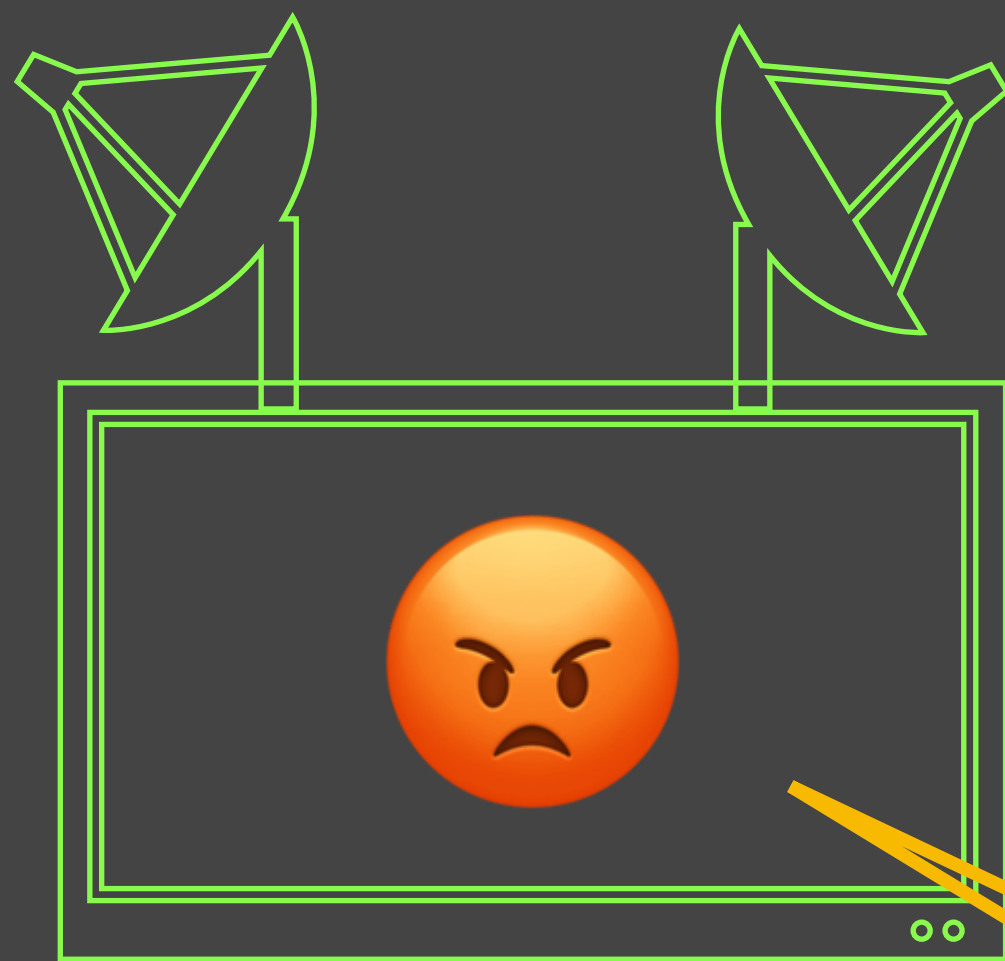
```
close ch2
```

{Send Int \Rightarrow
Recv Str \Rightarrow Close}

Communication : Incorrect Program

Use send twice to the same channel

```
let ch' = send 10155 channel_kun in  
let ch'' = send 30221 channel_kun in
```

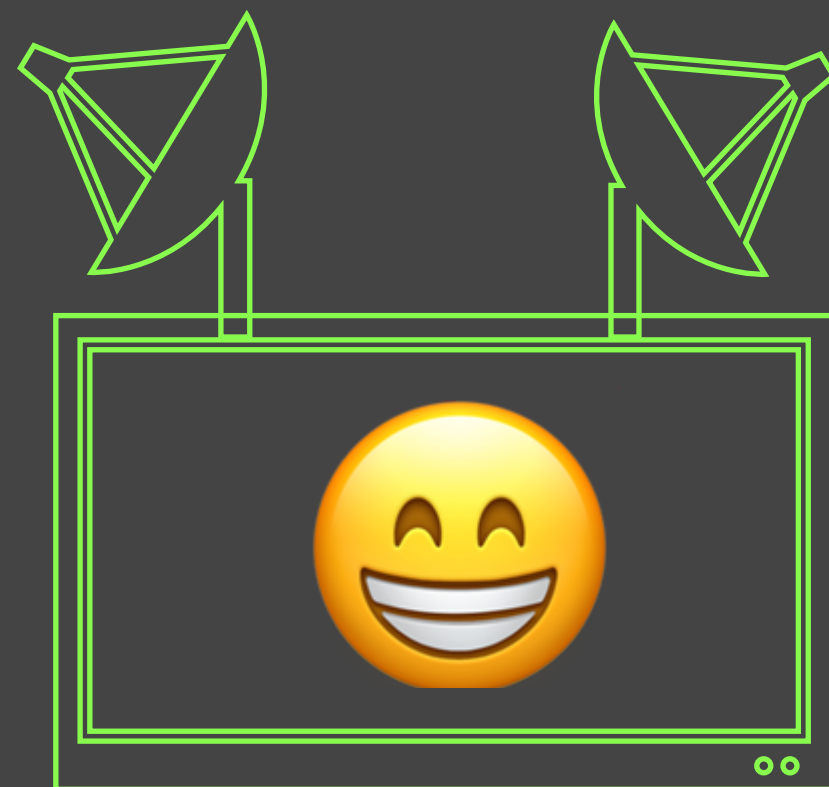


You can send **only one** number

Multiplicity of the Channel : 1

Quantitative types can detect **duplicate use** of channel.

```
main : (1 channel_kun : Channel (...)) -> ()  
✗ main channel_kun =  
    let ch'   = send 10155 channel_kun in  
    let ch''  = send 30221 channel_kun in
```



Summary

- Quantitative types make dependently-typed program more efficient
- Quantitative types allow to verify about how often the variable is used

Quantitative Type Theory (QTT)

Linear Types

Feature : Variables must be used **exactly once**

多くの言語にはない変数の使用状況の制約

Linear Haskell

Illegal Program in Linearly-Typed Language

Linear resource must be used

```
add1 : (1 x : Nat) -> Nat  
add1 x = 1
```

x must be used in right hand side!

Illegal Program in Linearly-Typed Language

Linear resource must not be used more than once

```
add1 : (1 x : Nat) -> Nat  
add1 x = x + x
```

x must not be used more than once
in right hand side!