



# Aaron Wilson

## Teaching experience

- Mentoring and development colleagues
- Problem classes and tutorials during my Ph.D

## Topics / Modules

- Information / Computer Security
- Cryptography
- Programing and software development

# A distraction

You are in Wales.

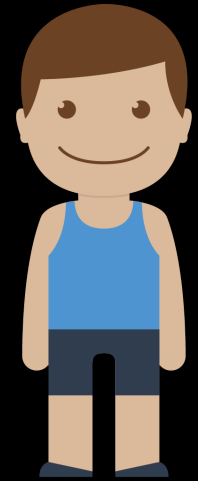
A bookworm eats from the first page of volume 1 to the last page of volume 3 in a 3-volume encyclopedia, which are stacked on the library shelves. Each volume is 3 centimeters wide.

How many centimeters has the bookworm eaten through ? More importantly, *why* is this the answer ?

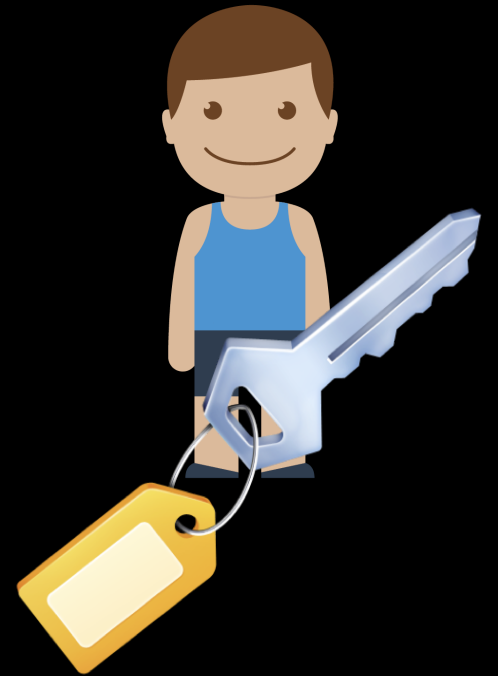
# Cryptographic Key Exchange

- Understand why we would want a key exchange mechanism
- Understand how the Diffie Hellman mechanism works
- Understand why the Diffie Hellman mechanism is secure

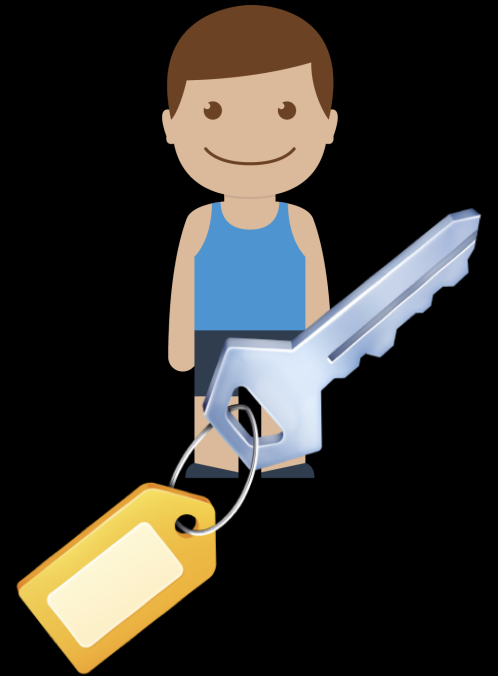
# Symmetric Cryptography



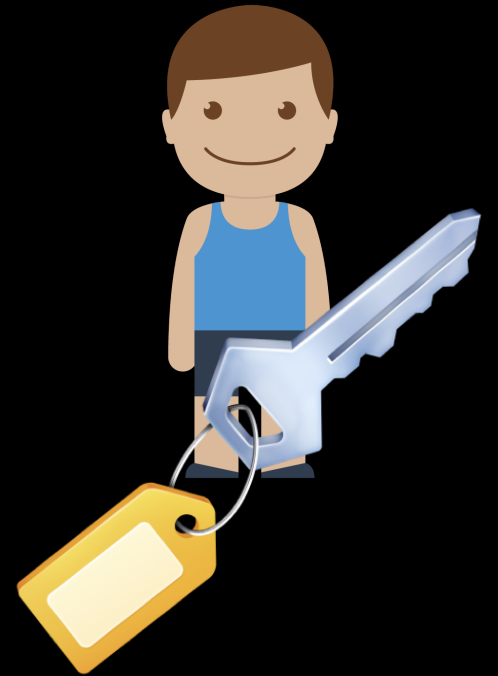
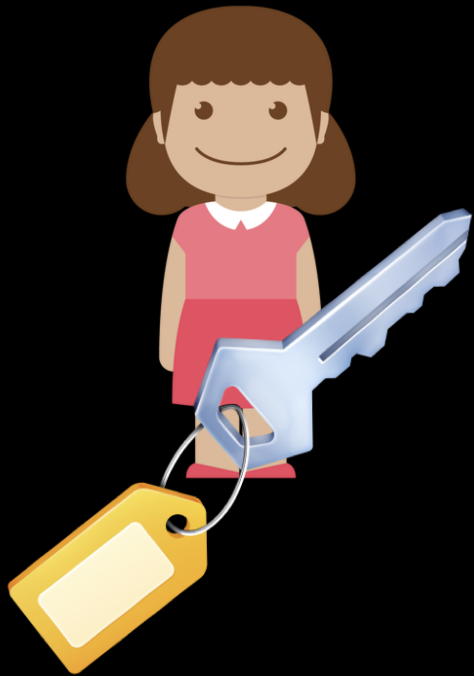
# Symmetric Cryptography



# Symmetric Cryptography

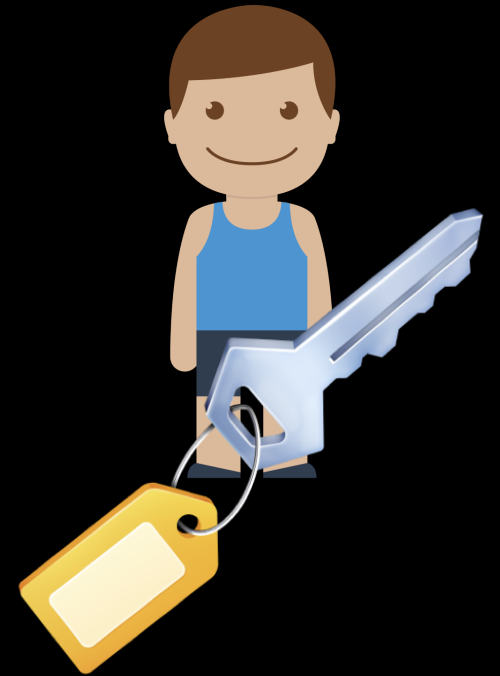
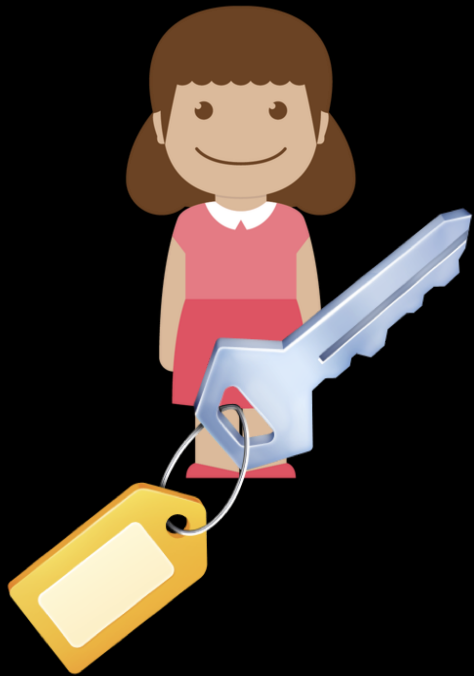


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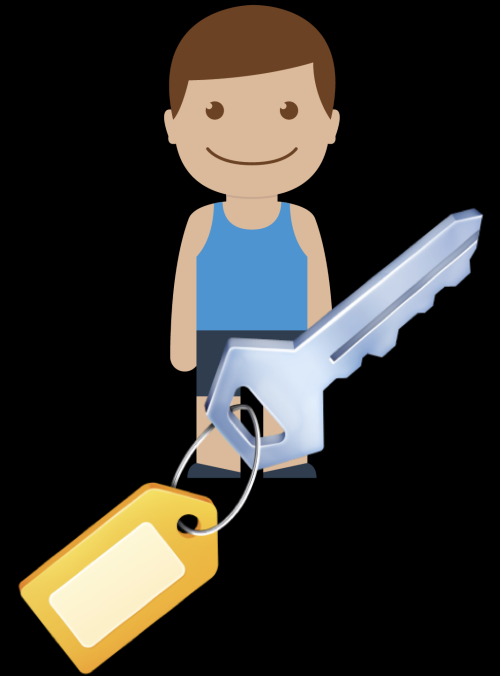




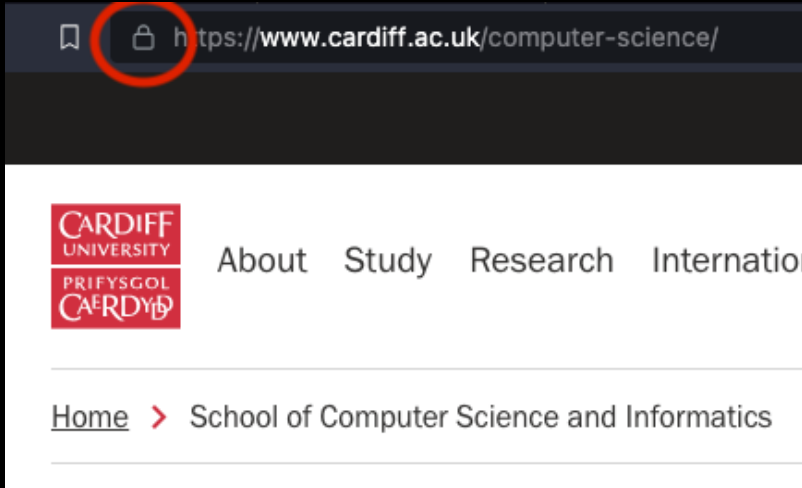
# Symmetric Cryptography



# Symmetric Cryptography



# Transport Layer Security



# What properties should a mechanism have ?

- Not rely on previously shared *secret* information
- Agreement
- Possibly also: authentication
- Computationally efficient

# Diffie Hellman Key Exchange

- Based on a concept of [Ralph Merkle](#)
- First published in 1976 by [Whitfield Diffie](#) and [Martin Hellman](#)
- One of the first public key protocols
- Non-authenticated key agreement protocol

# In action (simple)

▶ 0:00 / 0:34



# How does it work

- Factorisation is computationally hard
  - $23 * 97 = ?$
  - Factor 4321
  - Quantum computers
- Or equivalently the Discrete Logarithm problem is hard

# Algorithim (Simplest form, slide 1 of 2)

- Agree on a public prime  $p$  as a modulus, and a primitive root to use as a base  $g$ 
  - e.g let's pick  $p = 23$ , and  $g = 5$
- The first person picks a secret  $a$  and calculates  $g^a \mod p$ 
  - Let's have  $a = 4$  so that
$$g^a = 5^4 = 625 = 23 * 27 + 4 = 4 \mod 23$$
- The second person picks a secret  $b$  and calculates  $g^b \mod p$ 
  - Let's have  $b = 3$  so that
$$g^b = 5^3 = 125 = 23 * 5 + 10 = 10 \mod 23$$



# Algorithm (Simplest form, slide 2 of 2)

- The parties exchange their values
- The first person now calculates the key,  $g^{ab} = (g^b)^a \mod p$ 
  - That is the first party gets 10, so calculates
$$10^4 = 10000 = 23 * 434 + 11 = 18 \mod 23$$
- The second person now calculates the key,  $g^{ab} = (g^a)^b \mod p$ 
  - That is the second party receives 4, so calculates
$$4^3 = 64 = 2 + 18 \mod 23 = 18 \mod 24$$

# Public Information

- The modulus  $p$  and primitive root  $g$  (in our example 23 and 5)
- The transmitted values between Alice and Bob -  $g^a$  and  $g^b$  (in our example 4 and 10)
- Thus we can work out  $g^{a+b}$ ,  $g^{a-b}$
- **BUT NOT**  $g^{ab}$

# In action

▶ 0:00 / 0:48



# Security Properties

- Solve the discrete logarithm problem
- person-in-the-middle-attacks
- [Logjam](#)
  - first made public 2015-05-20

# Exercises / things to discuss

- Go through the example in the text above, and confirm it works
- Come up with other examples, and then try to break them
- person-in-the-middle attack: how does this work ?
- Could this work with more than two people involved ?

# Thank you !

Source code and slides are at [https://github.com/optimism-marjoram/diffie\\_hellman](https://github.com/optimism-marjoram/diffie_hellman)

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