

## MATHEMATICAL TOOLS - PROBLEM SET 3

Due Sunday, December 4th, 23:59, either in the course mailbox or through the Moodle. You may submit scanned files through the moodle, but please make sure they are crystal clear!

**Problem 1.** In this problem you'll show that the threshold for  $G(n, n, p)$  containing a perfect matching is  $\frac{\ln n}{n}$ .

For a bipartite graph  $H = (U, V, E)$  with  $|U| = |V|$ , a *perfect matching* is a vertex-disjoint set  $M \subseteq E$  s.t.  $|M| = |U|$ . We say that  $H$  satisfies the *marriage condition* if for every  $W \subseteq U$ ,  $|N_H(W)| \geq |W|$  ( $N_H(W) = \{v \in V : \exists w \in W, \{w, v\} \in E\}$  is the set of  $W$ 's neighbors). Recall Hall's marriage theorem:  $H$  contains a perfect matching iff it satisfies the marriage condition.

- (1) Show that if  $c < 1$  and  $p \leq c \frac{\ln n}{n}$ , a.a.s.  $G(n, n, p)$  doesn't contain a perfect matching. You may rely on results from recitation.
- (2) Show that if  $H = (U, V, E)$  doesn't satisfy the marriage condition then there exists some  $W \subseteq U$  s.t.  $|N_H(W)| = |W| - 1$ .
- (3) Henceforth, let  $c > 1$  and assume  $p \geq c \frac{\ln n}{n}$ . Let  $(X, Y, E) = G \sim G(n, n, p)$ , and for every  $W \subseteq X$  let  $A_W$  be the event that  $|N_G(W)| = |W| - 1$ . What is the distribution of  $|N_G(W)|$ ? Use this knowledge to bound  $\mathbb{P}[A_W]$ .
- (4) **Note (1/12): No need to hand in this clause. The result is correct, but doesn't follow straightforwardly from the preceeding clauses.** Conclude that a.a.s.,  $G$  satisfies the marriage condition and therefore contains a perfect matching.

**Problem 2.** A fixed point of a permutation  $\pi : [n] \rightarrow [n]$  is a value for which  $\pi(x) = x$ . Find the expectation and variance of the number of fixed points of a permutation chosen uniformly at random from all permutations.

### Measure Concentration.

**Problem 3.** Consider the following setting:  $n$  balls are distributed into  $m$  bins, where each ball "chooses" a uniformly random bin independent of all other choices.

- (1) Define an appropriate probability space that captures the situation above. For each  $1 \leq i \leq m$ , let  $X_i$  be the number of balls in the  $i$ th bin.
- (2) How is  $X_i$  distributed?
- (3) Are  $X_1, X_2, \dots, X_m$  independent?
- (4) Assume  $3\sqrt{\frac{m \ln m}{n}} \leq 1$ . Show that  $\mathbb{P}\left[\forall i \in [m], \left|X_i - \frac{n}{m}\right| < 3\sqrt{\ln m} \sqrt{\frac{n}{m}}\right] \geq 1 - \frac{2}{m^2}$ .
- (5) No need to hand anything in: Consider what happens when  $n \rightarrow \infty$ . Part 4 implies that if  $m = \omega(1)$ , then a.a.s. the number of balls in every bin is tightly concentrated around the mean - the probability of even a single bin deviating by more than  $3\sqrt{\ln m} \sqrt{\frac{n}{m}}$  tends to 0.

### The Binary Entropy Function.

**Problem 4.** Define  $f : [0, 1] \rightarrow \mathbb{R}$  by:

$$f(x) = \begin{cases} -x \log_2 x - (1-x) \log_2 (1-x) & x \in (0, 1) \\ 0 & x = 0, 1 \end{cases}$$

- (1) Prove that  $f$  is continuous on  $[0, 1]$ .
- (2) Find  $\max_{x \in [0, 1]} f(x)$ . Where does  $f$  attain its maximum?

**An Exponential Bound on the Tail of Poisson Random Variables.**

**Problem 5.** Let  $\lambda > 0$  and  $X \sim Poi(\lambda)$ .

- (1) For  $t > 0$ , calculate  $\mathbb{E}[e^{tX}]$ .
- (2) Use Markov's inequality w.r.t. the random variable  $e^{tX}$ , to conclude that for any  $t > 0$  and any  $k \geq 0$ :

$$\mathbb{P}[e^{tX} \geq e^{tk}] \leq \exp(\lambda e^t - tk - \lambda)$$

- (3) Conclude that for any  $k \geq \lambda$ :

$$\mathbb{P}[X \geq k] \leq \left(\frac{\lambda e}{k}\right)^k e^{-\lambda}$$