

Student Information

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Answer 1

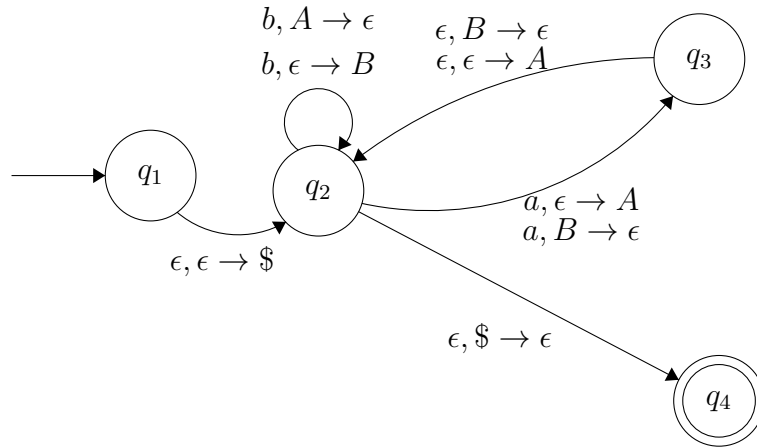
a) $G=(V, \Sigma, R, S)$ with set of variables $V=S$, where S is start variable; set of terminals $\Sigma = \{a, b\}$; and rules:

$$S \rightarrow bSbSa \mid aSbSb \mid bSaSb \mid SS \mid \epsilon$$

b) $G=(V, \Sigma, R, S)$ with set of variables $V=S$, where S is start variable; set of terminals $\Sigma = \{a, b\}$; and rules:

$$S \rightarrow aaSb \mid aSb \mid \epsilon$$

c) The corresponding Pushdown automata $M = (\{q_1, q_2, q_3, q_4\}, \{a, b\}, \{A, B, \$\}, \Delta, q_1, Z, \{q_4\})$ where:



d) Let context-free grammar for L_1 is G_1 , likewise let G_2 to be a context-free grammar for L_2 Then:

$$G_1 = (V_1, \Sigma_1, R_1, S_1)$$

$$G_2 = (V_2, \Sigma_2, R_2, S_2)$$

Let S be a new symbol and let $G = (V_1 \cup V_2 \cup \{S\}, \Sigma_1 \cup \Sigma_2, R, S)$ where:

$R = (R_1 \cup R_2 \cup \{S \rightarrow S_1, S \rightarrow S_2\})$ Then;

we can claim that $L(G) = L(G_1) \cup L(G_2)$

$G = (\{S_1, S_2, S\}, \{a, b\}, R, S)$ where;
 $R = \{S \rightarrow S_1 | S_2, S_1 \rightarrow bSbSa | aSbSb | bSaSb | SS | \epsilon, S_2 \rightarrow aaSb | aSb | \epsilon\}$

Answer 2

a) A context-free grammar G is ambiguous if \exists a string $w \in L(G)$ to have at least two different left parse trees.

Let try string $w = "0001111"$ to reach two different parse trees:

The First Parse tree;

$S \rightarrow AS$
 $\rightarrow A1S$
 $\rightarrow 0A11S$
 $\rightarrow 00A111S$
 $\rightarrow 0001111S$
 $\rightarrow 0001111$ and

The Second Parse Tree;

$S \rightarrow AS$
 $\rightarrow 0A1S$
 $\rightarrow 00A11S$
 $\rightarrow 00A111S$
 $\rightarrow 0001111S$
 $\rightarrow 0001111$

Hence;

We find a two different parse trees for a same string "0001111" we can conclude that this grammar is ambiguous.

b) $G_1 = (V, \Sigma, R, S)$, where;

$V = \{0, 1, S, A, X\}$, $\Sigma = \{0, 1\}$, and the rules are:

$R = \{S \rightarrow AS | \epsilon, A \rightarrow 01 | X, X \rightarrow 01 | 0A1\}$

disambiguate the G_1

c) Leftmost derivation of string $w = "00111"$ is:

$S \rightarrow AS$
 $\rightarrow A1S$
 $\rightarrow X1S$
 $\rightarrow 0A11S$
 $\rightarrow 00111S$
 $\rightarrow 00111$ and the corresponding parse tree;

