Lecture 9-10 Stacks

Introduction

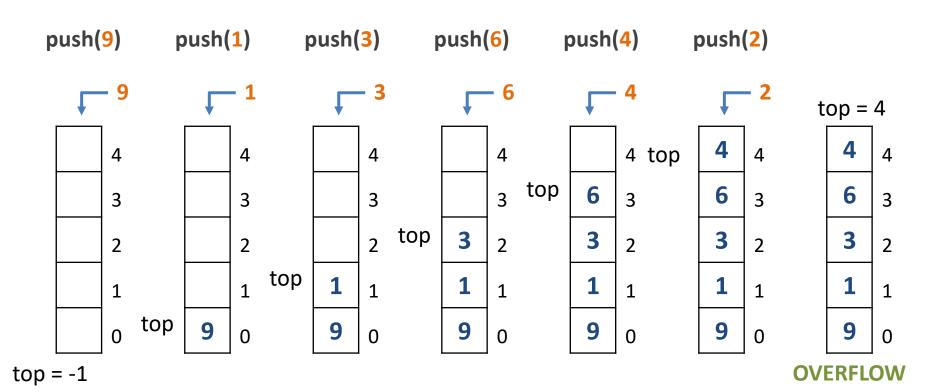
- Non primitive linear data structure.
- Allows operations at one end only.
- The top element can only be accessed at any time.
 - LIFO (Last-in-firstout) data structure.

Data Element n **Data Element 3** Data Element 2 Data Element 1

Operations

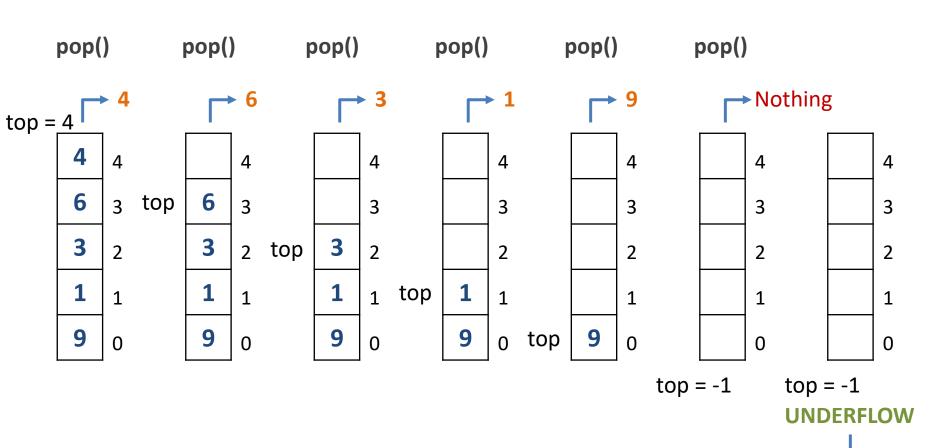
- Two primary operations:
 - push() Pushing (storing) an element on the stack.
 - pop() Removing (accessing) an element from the stack.
- Other operations for effective functionality:
 - peek() Get the top data element of the stack,top() without removing it.
 - isFull() Check if stack is full.OVERFLOW
 - isEmpty() Check if stack is empty. UNDERFLOW

Stack - Push



The stack is full, no more elements can be added. **OVERFLOW**

Stack - Pop



The stack is empty, no element can be removed. **UNDERFLOW**

Stack as an ADT

- A stack is an ordered list of elements of same data type.
- Elements are always inserted and deleted at one end.
- Following are its basic operations:
 - -S = init() Initialize an empty stack.
 - -isEmpty(S) Returns "true" if and only if the stack S is empty, i.e., contains no elements.

Stack as an ADT

- -isFull(S) Returns "true" if and only if the stack S has a bounded size and holds the maximum number of elements it can.
- -top(S) Returns the element at the top of the stack S, or error if the stack is empty.
- -S = push(S,x) Push an element x at the top of the stack S.
- -S = pop(S) Pop an element from the top of the stack S.
- -print(S) Prints the elements of the stack S from top to bottom.

Implementation

- Using static arrays
 - Realizes stacks of a maximum possible size.
 - Top is taken as the maximum index of an element in the array.
- Using dynamic linked lists
 - Choose beginning of the list as the top of the stack.

Using Static Arrays

Algorithm for Push

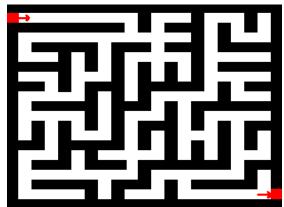
- Let,
 - STACK[SIZE] is a one dimensional array that will hold the stack elements.
 - TOP is the pointer that points to the top most element of the stack.
 - DATA is the data item to be pushed.
 - 1. If TOP == SIZE 1
 - 2. Display "Overflow condition"
 - 3. Else
 - 4. TOP = TOP + 1
 - 5. STACK [TOP] = DATA

Algorithm for Pop

- Let,
 - STACK[SIZE] is a one dimensional array that will hold the stack elements.
 - TOP is the pointer that points to the top most element of the stack.
 - DATA is the element popped from the top of the stack.
 - 1. If TOP < 0 (or TOP == -1)
 - 2. Display "Underflow condition."
 - 3. Else
 - 4. DATA = STACK[TOP]
 - 5. TOP = TOP 1
 - 6. Return DATA

Applications

- Reverse a word.
 - Push a word letter by letter and then pop letters from the stack.
- "UNDO" mechanism in text editors.
- Backtracking.
 - Game playing, finding paths, exhaustive searching.
- Parsing.
- Recursive function calls.
- Calling a function.
- Expression Evaluation.
- Expression Conversion.



Stack Applications

Expression Representation

- Infix Operator is in-between the operands.
- Prefix Operator is before the operands. Also known as polish notation.
- Postfix Operator is after the operands. Also known as suffix or reverse polish notation.

Infix	Prefix	Postfix
a + b	+ a b	a b +
a + b * c	+ a * b c	a b c * +
(a + b) * (c - d)	* + a b - c d	a b + c d - *

Example (Infix to Postfix)

```
• 5 + 3 * 2
         • 532*+
• 3 + 4 * 5 / 6
         • 345 * 6/+
• (300 + 23) * (43 - 21) / (84 + 7)

 300 23 + 43 21 - * 84 7 + /

• (4+8)*(6-5)/((3-2)*(2+2))
         • 48+65-*32-22+*/
```

Infix to Postfix Conversion Algorithm

Let Q be any infix expression to be converted in to a postfix expression P.

- 1. Push left parenthesis onto STACK and add right parenthesis at the end of Q.
- 2. Scan Q from left to right and repeat step 3 to 6 for each element of Q until the STACK is empty.
- 3. If an operand is encountered add it to P.
- 4. If a left parenthesis is encountered push it onto the STACK.
- 5. If an operator is encountered, then
 - i. Repeatedly pop from STACK and add to P each operator which has same precedence as or higher precedence than the operator encountered.
 - ii. Push the encountered operator onto the STACK.
- 6. If a right parenthesis is encountered, then
 - i. Repeatedly pop from the STACK and add to P each operator until a left parenthesis is encountered.
 - ii. Remove the left parenthesis; do not add it to P.
- 7. Exit.

Input	Stack	Output			
A + (B * (C – D) / E)					
A + (B * (C – D) / E))	(
+ (B * (C – D) / E))	(Α			
(B * (C – D) / E))	(+	Α	Input	Stack	Output
B * (C – D) / E))	(+(Α	– D) / E))	(+(*(АВС
* (C – D) / E))	(+(A B	D) / E))	(+(*(-	АВС
(C – D) / E))	(+(*	АВ) / E))	(+(*(-	ABCD
C – D) / E))	(+(*(АВ	/ E))	(+(*	A B C D –
			E))	(+(/	A B C D - *
))	(+(/	A B C D - * E
)	(+	A B C D - * E
					A B C D - * E

Postfix Evaluation Algorithm

- 1. Initialize empty stack
- 2. For every token in the postfix expression (scanned from left to right):
 - a. If the token is an operand (number), push it on the stack
 - b. Otherwise, if the token is an operator (or function):
 - i. Check if the stack contains the sufficient number of values (usually two) for given operator
 - ii. If there are not enough values, finish the algorithm with an error
 - iii. Pop the appropriate number of values from the stack
 - iv. Evaluate the operator using the popped values and push the single result on the stack
- 3. If the stack contains only one value, return it as a final result of the calculation
- 4. Otherwise, finish the algorithm with an error

Example

- Evaluate 2 3 4 + * 6 -
- 2*(3+4)-6=8.

	Input token	Operation	Stack contents (top on the right)	Details
234+*6-	2	Push on the stack	2	
34+*6-	3	Push on the stack	2, 3	
4 + * 6 -	4	Push on the stack	2, 3, 4	
+ * 6 -	+	Add	2, 7	Pop two values: 3 and 4 and push the result 7 on the stack
* 6 –	*	Multiply	14	Pop two values: 2 and 7 and push the result 14 on the stack
6 -	6	Push on the stack	14, 6	
-	-	Subtract	8	Pop two values: 14 and 6 and push the result 8 on the stack
	(End of tokens)	(Return the result)	8	Pop the only value 8 and return it

Thank You