

Sorting Algorithms - II

Quick Sort

Merge Sort

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- Divide and Conquer algorithm.
- Picks an element as pivot and partitions the given array around the picked pivot, such that
 - The pivot is placed at its correct position
 - All elements smaller than the pivot are placed before the pivot.
 - All elements greater than the pivot are placed after the pivot.
- Several ways to pick a pivot.
 - The first element.
 - The last element.
 - Any random element.
 - The median.

Quick Sort

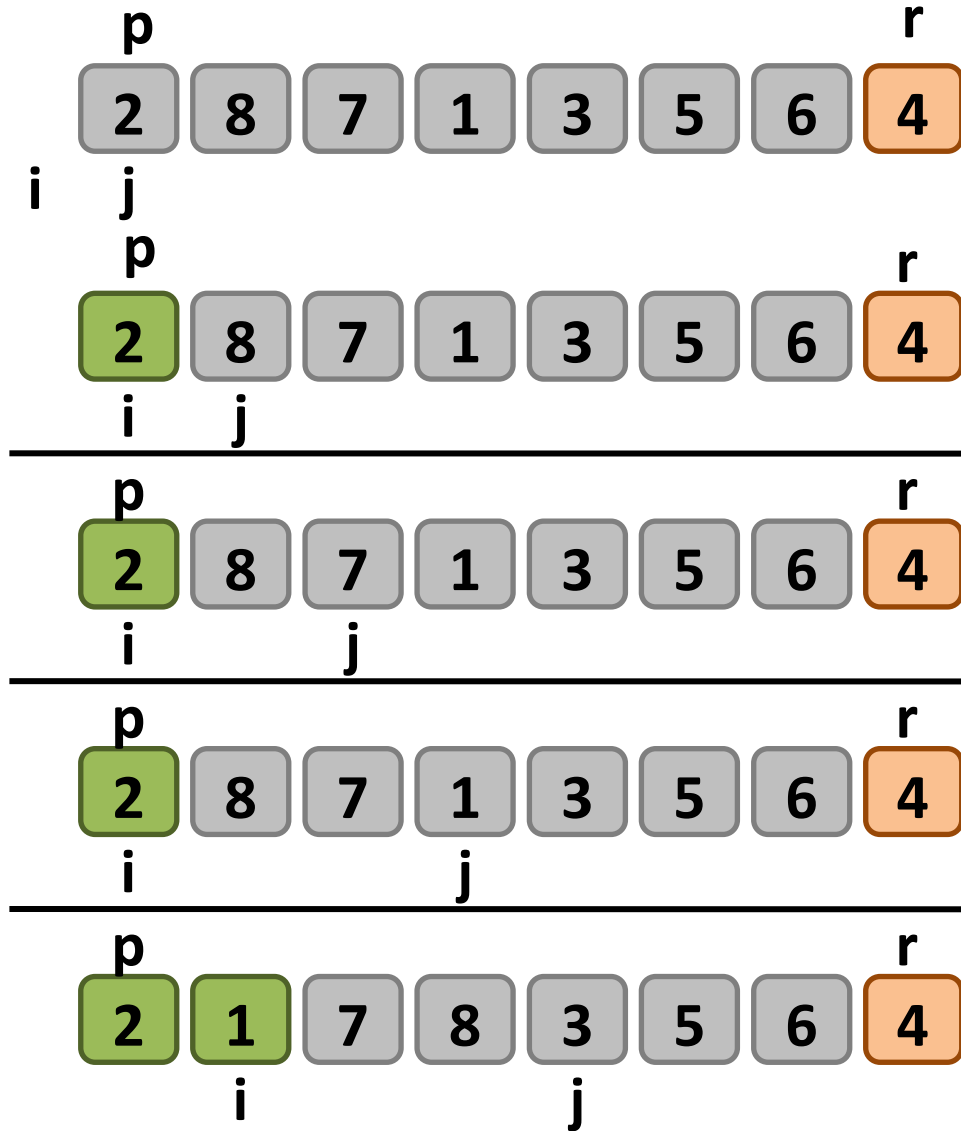


- QUICKSORT(A, p, r)
 1. if $p < r$
 2. $q = \text{PARTITION}(A, p, r)$
 3. $\text{QUICKSORT}(A, p, q - 1)$
 4. $\text{QUICKSORT}(A, q + 1, r)$

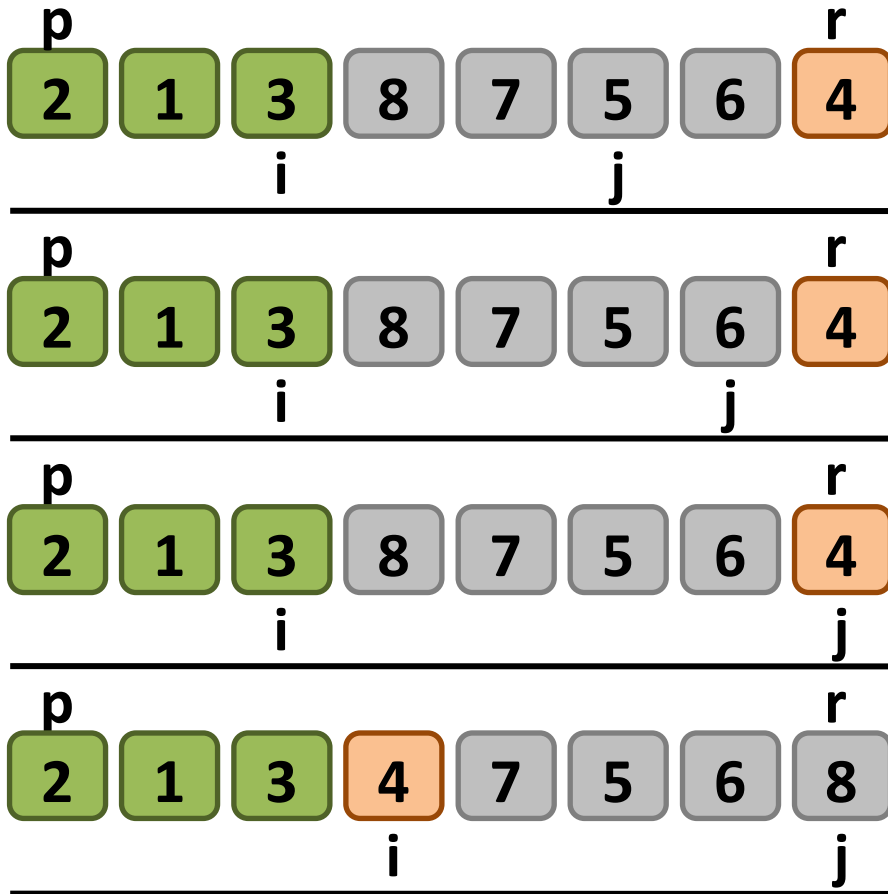
To sort an array A with n elements, the first call to QUICKSORT is made with $p = 0$ and $r = n - 1$.

Algorithm

1. PARTITION(A, p, r)
2. $x = A[r]$
3. $i = p - 1$
4. for $j = p$ to $r - 1$
5. if $A[j] \leq x$
6. $i = i + 1$
7. Exchange $A[i]$ with $A[j]$
8. Exchange $A[i + 1]$ with $A[r]$
9. return $i + 1$

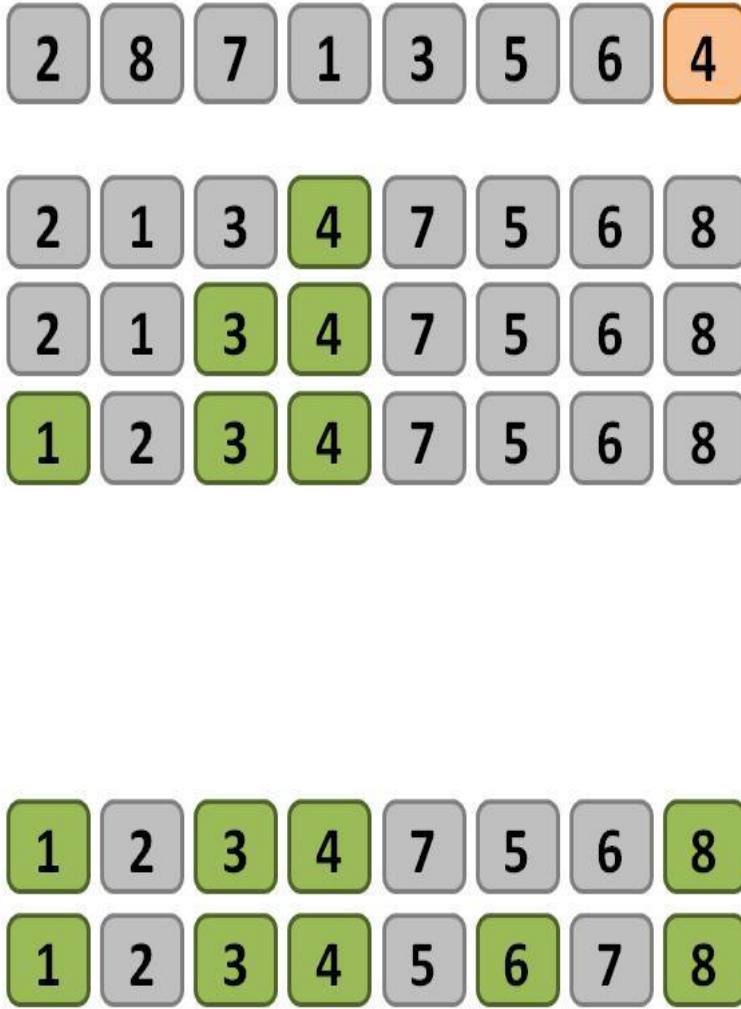
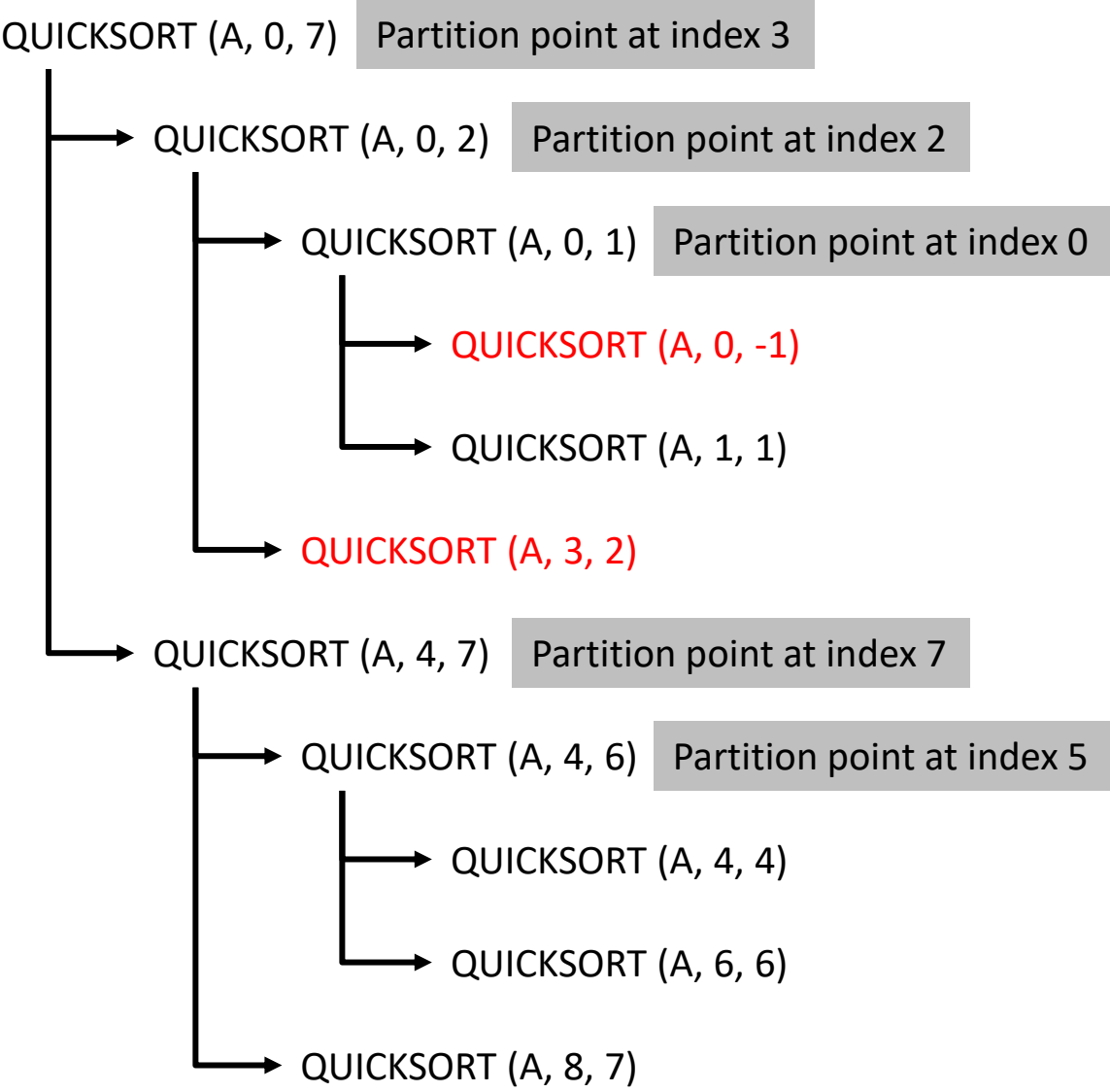


Algorithm



1. PARTITION(A, p, r)
2. $x = A[r]$
3. $i = p - 1$
4. for $j = p$ to $r - 1$
5. if $A[j] \leq x$
6. $i = i + 1$
7. Exchange $A[i]$ with $A[j]$
8. Exchange $A[i + 1]$ with $A[r]$
9. return $i + 1$

Call sequence for QUICKSORT



Merge Sort

Merge Sort

- Based on the divide-and-conquer paradigm.
- To sort an array $A[p \dots r]$, (initially $p = 0$ and $r = n-1$)

1. Divide Step

- If a given array A has zero or one element, then return as it is already sorted.
- Otherwise, split $A[p \dots r]$ into two subarrays $A[p \dots q]$ and $A[q + 1 \dots r]$, each containing about half of the elements of $A[p \dots r]$. That is, q is the halfway point of $A[p \dots r]$.

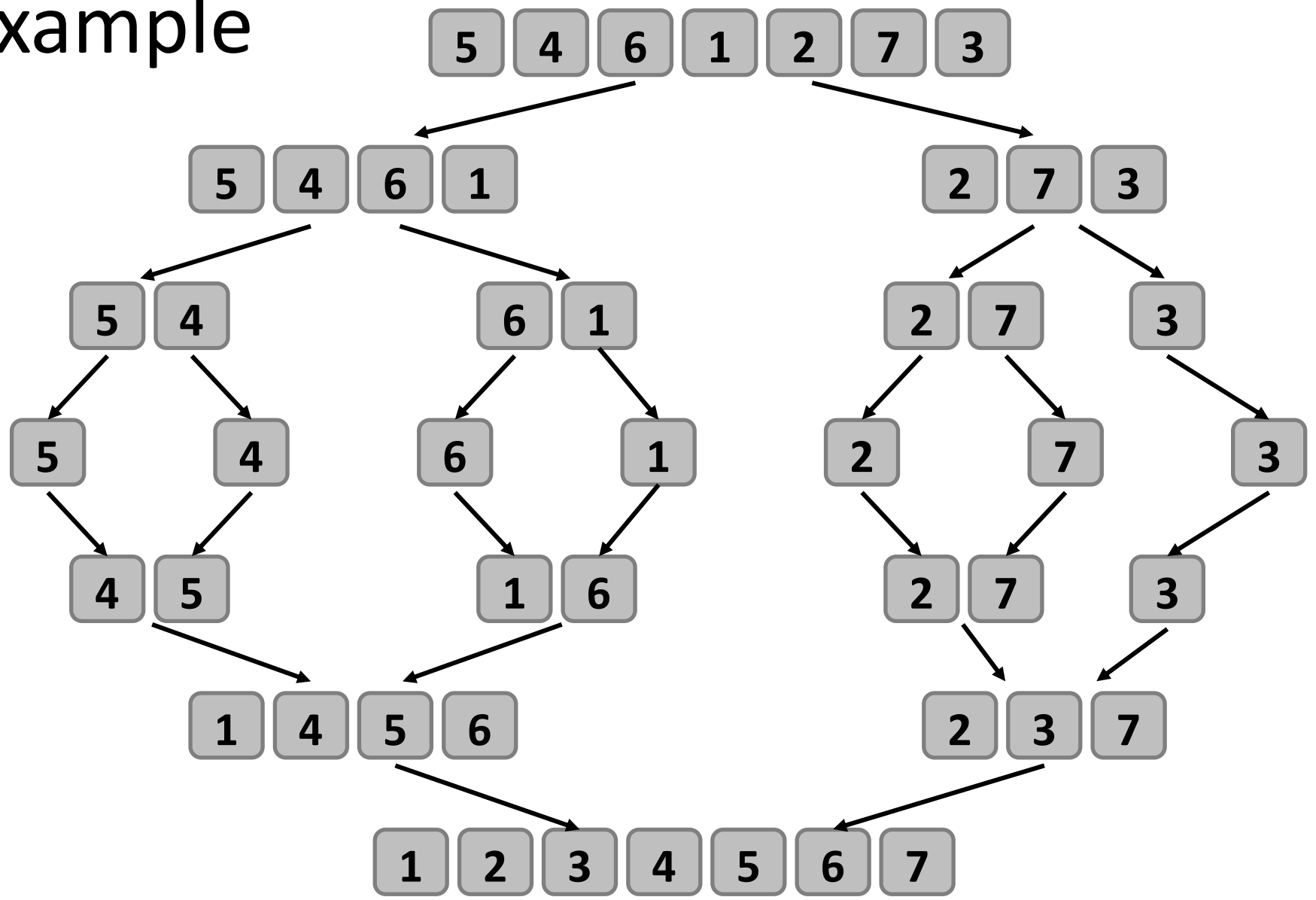
2. Conquer Step

- Recursively sort the two subarrays $A[p \dots q]$ and $A[q + 1 \dots r]$.

3. Combine Step

- Combine the elements back in $A[p \dots r]$ by merging the two sorted subarrays $A[p \dots q]$ and $A[q + 1 \dots r]$ into a sorted sequence.

Example



Merge Two Sorted Arrays

n1 - #Elements in L
n2 - #Elements in R

L:

1	4	5	6
---	---	---	---

i i i i i

A:

1	2	3	4	5	6	7
---	---	---	---	---	---	---

k k k k k k k k

R:

2	3	7
---	---	---

j j j j

8. $i = 0, j = 0$, and $k = p$.

9. while $i < n1$ and $j < n2$

10. if $L[i] \leq R[j]$

11. $A[k] = L[i]$

12. $i = i + 1$

13. else

14. $A[k] = R[j]$

15. $j = j + 1$

16. $k++$

17. while $i < n1$

18. $A[k] = L[i]$

19. $i++$

20. $k++$

21. while $j < n2$

22. $A[k] = R[j]$

23. $j++$

24. $k++$

Algorithm

- MERGE-SORT (A, p, r)
 1. if $p < r$
 2. $q = \text{FLOOR}[(p + r)/2]$
 3. MERGE-SORT(A, p, q)
 4. MERGE-SORT($A, q + 1, r$)
 5. MERGE (A, p, q, r)
- To sort an array A with n elements, the first call to MERGE-SORT is made with $p = 0$ and $r = n - 1$.

Contd...

- Algorithm MERGE (A, p, q, r)
- Input: Array A and indices p, q, r such that $p \leq q \leq r$.
Subarrays $A[p \dots q]$ and $A[q + 1 \dots r]$ are sorted.
- Output: The two subarrays are merged into a single sorted subarray in $A[p \dots r]$.
 1. $n1 = q - p + 1$
 2. $n2 = r - q$
 3. Create arrays $L[n1]$ and $R[n2]$
 4. for $i = 0$ to $n1 - 1$
 5. $L[i] = A[p + i]$
 6. for $j = 0$ to $n2 - 1$
 7. $R[j] = A[q + 1 + j]$

Contd...

8. $i = 0, j = 0$, and $k = p$.

9. while $i < n1$ and $j < n2$

10. if $L[i] \leq R[j]$

11. $A[k] = L[i]$

12. $i = i + 1$

13. else

14. $A[k] = R[j]$

15. $j = j + 1$

16. $k++$

17. while $i < n1$

18. $A[k] = L[i]$

19. $i++$

20. $k++$

21. while $j < n2$

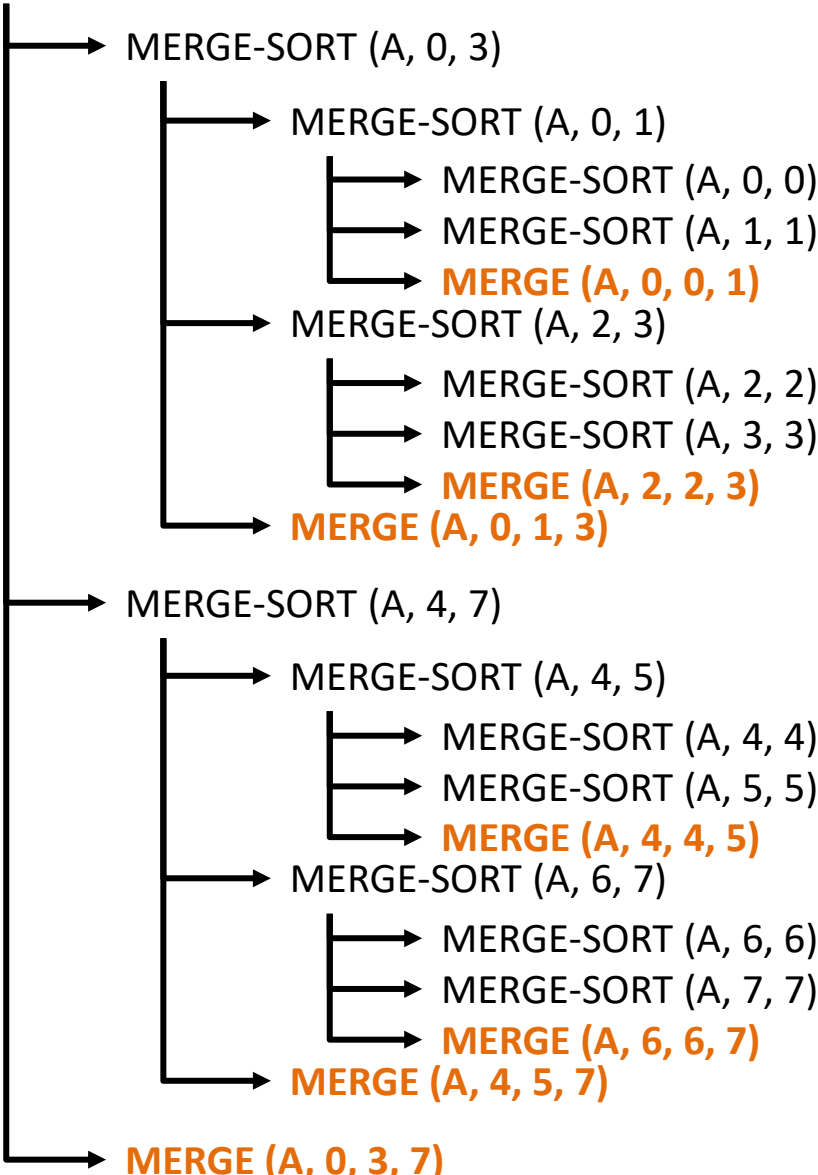
22. $A[k] = R[j];$

23. $j++;$

24. $k++;$

Call sequence for an array with size 8

MERGE-SORT (A, 0, 7)

[illegible]

Thank You