

Hierarchy Measurement Revisited

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Abstract

Hierarchy permeates all aspects of human behavior and societal outcomes, making precise measurements fundamental to understanding its ramifications. However, existing hierarchical measures fail to capture key structural dimensions of hierarchy. This work addresses these limitations by significantly expanding the axiomatic foundation for hierarchical measurement. We develop precise measures of hierarchical depth that enhance comparability across different organizational structures. Our expanded axiom set captures dimensions that traditional measures overlook, revealing systematic violations in existing approaches. As a robust alternative, we propose the *average depth index*—a hierarchical measure that satisfies all axioms introduced in this work. This approach provides a more reliable and theoretically grounded tool for comparing organizational hierarchies.

Keywords: hierarchy measurement, hierarchical pre-order, hierarchical index, directed acyclic graph.

JEL classifications: D23, L22.

1. Introduction

Hierarchy is fundamental in shaping behavior across animal and human societies, serving as a powerful determinant of social outcomes. In animal groups, hierarchical structures regulate resource access, mating opportunities, and social alliances. Similarly, hierarchy influences behavior in human societies through power dynamics, social status, and cultural norms. Whether expressed through formal organizational charts, informal networks, or socially constructed status, hierarchical structures permeate nearly every aspect of human interaction.

High-status individuals often display greater confidence and assertiveness, while those lower in the hierarchy may conform to norms or strive for upward mobility—frequently driven by social comparison and competition, such as the phenomenon of “keeping up with the Joneses.”

Ultimately, hierarchy-induced behavior shapes outcomes in complex ways. Hierarchies can promote stability, adaptability, coordination, and efficiency in a complex and uncertain world (Simon, 1962, 1973; Holland, 1998; Lawrence and Lorsch, 1967; Thompson, 1967;

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Lincoln and Kalleberg, 1985). Informal hierarchical networks further contribute to organizational resilience by facilitating crisis management (Krackhardt and Stern, 1988), improving profitability beyond traditional financial metrics (Sarkar et al., 2010), and compensating for dysfunctional formal systems (Culen, 2017).

However, these same structures can generate unintended consequences that undermine their benefits. Hierarchies produce emergent collective behaviors and group dynamics that influence outcomes in unpredictable ways (Anderson, 1972). They can enable power abuse, entrench social and economic inequalities (Fix, 2017, 2018, 2019, 2021), and stifle innovation (Wright, 2024).

Hierarchy, therefore, is not merely a structural feature—it is an active force that shapes behavior and determines the consequences of social interaction.

To fully understand the impact of hierarchy on behavior and outcomes, precise and consistent measurement is essential. Without rigorous methods to quantify hierarchy, comparing structures across institutions or analyzing their dynamic evolution becomes difficult.

To address this need for precise measurement, this paper extends the axiomatic theory for hierarchy measurement developed in Carbonell-Nicolau (2025a,b) through a significant expansion of its axiom set, thereby enhancing our ability to compare hierarchical structures both across organizations and over time.

We model hierarchies as directed acyclic graphs in which nodes represent individuals (employees, managers, executives) and edges capture subordination relationships that flow from supervisor to subordinate.

Our theoretical framework rests on hierarchical pre-orders—reflexive and transitive binary relations defined across the universe of possible hierarchical structures. This axiomatic approach offers a rigorous foundation for measuring and comparing firm hierarchies, introducing mathematical precision and analytical flexibility to Organizational Economics. By establishing consistency through formal axioms, our framework addresses the fundamental limitations inherent in ad-hoc measurement approaches that have dominated the field.

The expanded axiom set we develop focuses on two fundamental dimensions of organizational structure: hierarchical depth and structural transformation dynamics. We formalize the relationship between structural modifications and overall hierarchical measurement, developing precise analytical tools for understanding organizational changes, including management layer additions, various forms of organigram integration, and real-world scenarios such as partial subunit mergers and acquisitions.

Our framework builds on a core hierarchical pre-order established in Carbonell-Nicolau (2025a,b). This pre-order is the only ordering relationship fully characterized by three fundamental axioms: Anonymity, the Replication Principle, and Subordination Removal.

Anonymity ensures that the essential hierarchical structure remains unchanged when we relabel the nodes within a hierarchy. In other words, the specific names or identifiers assigned to positions do not affect the underlying organizational relationships—what matters is the pattern of connections, not the labels themselves.

The Replication Principle addresses how we evaluate multiple copies of hierarchical structures. When we have several identical copies of a given hierarchy—essentially creating a structure with multiple disconnected components—this principle treats these replicated hierarchies as having the same degree of hierarchical organization as the original single hierarchy.

Subordination Removal captures a fundamental principle about how hierarchical relationships affect organizational structure. This axiom states that when we eliminate a direct subordination relationship—specifically by removing the link between a subordinate i and their immediate supervisor j , and instead connecting i directly to j 's supervisor—we necessarily reduce the overall “hierarchical degree” of the system. By shortening the chain of command and removing j 's direct control over i , we create a flatter, less hierarchical structure.

While the present work retains these fundamental principles of hierarchy measurement, we recognize their insufficiency when studying organizational transformations that affect structural depth, such as the addition of management layers and the integration of separate hierarchies.

Consider a simple example to illustrate this limitation. Hierarchy h consists of a two-node chain: a supervisor and a subordinate. Hierarchy h' introduces a “boss” node that directly supervises h 's supervisor from above. Crucially, hierarchy h cannot be derived from h' by removing subordination links alone, even when both hierarchies are made comparable in node count through replication. This means one cannot transition between these hierarchies purely through subordination removals and replication operations. Yet intuitively, h' incorporates a new top-management node that deepens the hierarchical structure. Despite this apparent increase in hierarchy depth, the core measure from [Carbonell-Nicolau \(2025a,b\)](#) renders h and h' incomparable.

To address this limitation, we propose several new axioms that formalize how hierarchical structures can be extended while increasing their hierarchical depth.

The Upward Extension axiom formalizes the intuitive notion that adding top managers to an existing hierarchy increases its hierarchical depth. When a hierarchy h is appended below another hierarchy h_o , the resulting compound hierarchy exhibits a greater hierarchical structure than h alone. For upward extension to be valid, each leaf node in the upper hierarchy h_o must connect to at least one root node in the lower hierarchy h . This connection requirement ensures a consistent increase in structural depth. It prevents scenarios where a weakly connected upper hierarchy introduces excessive “flatness” that could overshadow the lower hierarchy and reduce their connection to a negligible detail.

The Downward Extension axiom operates in reverse, stating that appending a hierarchy below another hierarchy results in a greater hierarchical structure, relative to the upper hierarchy, when every root node in the lower hierarchy connects to one or more leaf nodes at maximum depth in the upper hierarchy. Here, the “depth” of a node refers to the number of its direct or indirect supervisors.

The notion of “appending” differs between these two axioms: upward extension requires every leaf node in the upper hierarchy to link to at least one root node in the lower hierarchy, creating mandatory connections from top to bottom. In contrast, downward extension allows more flexibility by requiring that every root node in the lower hierarchy become a direct subordinate to a maximum-depth leaf node in the upper hierarchy while not requiring all upper leaf nodes to connect downward.

Beyond these axioms, additional variants capture different aspects of hierarchical extension under weaker connection requirements. The Upward Extension* and Downward Extension* axioms are formulated using a more relaxed notion of hierarchy appending, where at least one leaf node in the upper hierarchy connects to at least one root node in the lower hierarchy rather than requiring comprehensive connections.

Upward Extension* shares the conceptual foundation of standard Upward Extension but operates under different conditions. It asserts that adding a hierarchy h_o on top of h increases the hierarchical degree when two conditions are met: first, h_o must be weakly more hierarchical than h , and second, only partial linkage is required between the hierarchies. This weaker connection requirement is compensated by the stronger hierarchical structure of the appended hierarchy, ensuring that the compound structure achieves greater depth.

Similarly, Downward Extension* stipulates that adding h_o below h increases hierarchical degree when h_o is weakly more hierarchical than h . As with Upward Extension*, the connection between hierarchies is partial, requiring only that at least one leaf node in the upper hierarchy connects to at least one root node in the lower hierarchy.

These axioms with partial connection requirements, working in conjunction with those requiring strict connections, allow for more comprehensive comparisons between hierarchical structures while capturing different aspects of hierarchical composition.

Each extension principle gives rise to a canonical hierarchical pre-order—a reflexive and transitive binary relation on the set of all hierarchies—that fully characterizes the underlying extension axiom when combined with the Anonymity axiom and the Replication Principle. These individual hierarchical pre-orders can be combined into a composite pre-order by considering the transitive closure of their union. This approach creates a more complete measure by combining all extension pre-orders, though the composite framework can consider any subset of measures.

Every composite hierarchical pre-order defined in this manner represents the coarsest pre-order that aligns with the underlying axioms, providing a comprehensive framework for comparing hierarchical structures across multiple extension and connection requirements.

With our axiomatic groundwork in place, we shift our focus to the study of hierarchical indices, which, unlike our pre-order measures, give rise to complete rankings of hierarchies. Hierarchical indices present clear computational advantages due to their simplicity, contrasting with the complexity of composite hierarchical pre-orders defined as the transitive closure of a union of canonical measures. However, this computational convenience comes at a cost: every comparison between hierarchies receives a definitive ranking, even in ambiguous cases where none of the established axioms applies or provides clear guidance.

Our analysis of hierarchical indices begins with the *average depth index*, introduced in prior work (Carbonell-Nicolau, 2025a), which calculates the mean depth—the number of direct or indirect supervisors—across all nodes in a given hierarchy. We demonstrate that this index fulfills all the axioms considered in this paper, ensuring consistency with our hierarchical pre-orders whenever they produce a definitive ranking between two hierarchies. This alignment makes the average depth index a reliable computational tool that preserves the theoretical foundations established by our axiomatic framework.

We also examine other hierarchical indices proposed in the literature, contextualizing them within our axiomatic framework. Our analysis reveals that, unlike the average depth index, these alternative indices systematically violate our axioms. This finding highlights the importance of axiomatic validation in developing hierarchical measures and demonstrates why the average depth index stands out as both computationally efficient and theoretically sound.

The *global reaching centrality (grc) index*, introduced by Mones et al. (2012), is based on the average concentration of “reaching centrality,” which measures a node’s ability to reach

other nodes in the hierarchy. While this index satisfies the Anonymity axiom, it violates all other axioms in our framework.

The failure of the Subordination Removal axiom has been previously illustrated in [Carbonell-Nicolau \(2025a\)](#). Additionally, the replication operation does not leave the grc index unchanged, thus violating the Replication Principle. This violation occurs because replication proportionally increases the number of nodes without proportionally increasing the concentration of reaching centrality across those nodes, resulting in a lower average reaching centrality for the replicated structure.

We further demonstrate that global reaching centrality fails to capture the notion of vertical hierarchical growth underlying the various extension principles considered in this paper. The fundamental issue with the grc index's violations of these principles stems from the fact that deeper hierarchies may be associated with a less concentrated distribution of reaching centrality. This relationship between depth and centrality concentration illuminates the grc index's systematic failure to align with our axiomatic framework.

[Krackhardt \(1994\)](#) studied informal organizations—networks of relationships between individuals that are not necessarily captured by official organizational charts but are formed to accomplish tasks quickly and effectively. He identified four dimensions along which analysts can measure the similarities between an informal organization and what he termed formal organizational structures. Krackhardt assumed these formal structures take the form of traditional hierarchies, which he modeled as out-trees or arborescences, particular instances of directed acyclic graphs representing well-delineated hierarchical relationships.

For each dimension, Krackhardt defined an index quantifying a network's "hierarchical degree". We evaluate these measures, defined on the domain of directed acyclic graphs, within our axiomatic framework.

Krackhardt's *connectedness index* measures the extent of information flows between nodes in a hierarchy without regard for directionality, effectively quantifying linkages in the underlying undirected graph. Our analysis shows that the connectedness index violates all of our axioms except Anonymity.

The violation of the Replication Principle occurs because replication reduces connectedness by creating one or more copies of the original hierarchy, with each copy disconnected from its counterparts. The Subordination Removal axiom is violated because connectedness does not depend on the specific structure of subordination relations—these relations can be eliminated without breaking the linkages between nodes measured by the index.

Similarly, the connectedness index violates the various extension axioms because appending operations between hierarchies can be conducted without altering existing linkages between nodes within each hierarchy.

Krackhardt's *graph hierarchy index* proves trivial within the domain of directed acyclic graphs, as it assumes a constant value across all acyclic hierarchical structures. Consequently, this measure lacks discriminatory power when comparing different hierarchical configurations, making it unsuitable for our analytical purposes.

Krackhardt's *graph efficiency index* quantifies the gap between a hierarchy and its minimally connected counterpart, effectively measuring linkage "redundancy." Since Krackhardt evaluates efficiency within weakly connected components rather than across disconnected components, the graph efficiency index remains consistent with the Replication Principle—replication creates separate components without affecting internal efficiency measures.

However, this measure violates both the Subordination Removal axiom and the extension axioms. The Subordination Removal violation occurs because subordination operations can be conducted while maintaining minimal connections across nodes, leaving the efficiency measure unaltered despite fundamental changes to the hierarchical structure. Similarly, the extension axioms are violated because hierarchy appending can be achieved without adding redundant links, keeping the graph efficiency index constant even when the overall hierarchical structure has demonstrably increased in complexity and depth.

Krackhardt's final measure, the *least upper boundedness index*, assesses how closely a hierarchy approximates a pure arborescence. Like the efficiency index, it is calculated within weakly connected components rather than across disconnected ones, ensuring consistency with the Replication Principle.

However, the measure exhibits a fundamental limitation: it assigns the same value to all pure trees, regardless of the complexity or depth of their hierarchical structure. In other words, it is insensitive to differences in hierarchical depth or the specific patterns of subordination. As a result, the index cannot distinguish between shallow and deep hierarchies, thereby violating the Subordination Removal and extension axioms. While the index effectively identifies tree-like structures, it fails to capture the nuanced variations in hierarchical organization that our axiomatic framework is designed to measure.

Other hierarchical indices have been introduced in the literature, such as those proposed by Trusina et al. (2004), Luo and Magee (2011), Corominas-Murtra et al. (2013), and Czégel and Palla (2015).

The indices developed by Corominas-Murtra et al. (2013) and Trusina et al. (2004) are specifically designed to detect deviations from hierarchical structure that do not occur in directed acyclic graphs. As a result, these indices reduce to constant values when applied to directed acyclic graphs, making them ineffective for distinguishing between different topologies within this class of graphs.

A similar limitation applies to the measure introduced by Luo and Magee (2011). The flow hierarchy metric discussed in that work counts the number of links in a network that are not part of loops. Since directed acyclic graphs, by definition, contain no loops, every link is counted, and the metric always yields its maximum value. Consequently, this measure cannot differentiate between distinct directed acyclic graphs or provide meaningful comparisons of their hierarchical structures.

Czégel and Palla (2015) propose an algorithmic approach to measuring hierarchy based on random walks within a network. In their method, the stationary distribution of random walkers reflects each node's position in the hierarchy: a greater concentration of walkers at higher levels indicates a more pronounced hierarchical structure. The measure itself quantifies the inequality in this distribution.

An important feature of this approach is its dependence on a decay parameter, which the analyst must choose. The form of the stationary distribution—and thus the value of the hierarchy index—varies according to this parameter. While the measure intuitively captures hierarchy in multi-level structures, its reliance on the decay parameter can violate the extension axioms if it is not optimally set. This dependence limits the measure's ability to consistently reflect hierarchical growth through upward or downward extensions.

The paper is organized as follows. Section 2 formally defines hierarchies as directed acyclic graphs and introduces key concepts such as root nodes, leaf nodes, and node depth. Section 3 presents the axiomatic framework for hierarchy measurement, detailing the core axioms: Anonymity, Replication Principle, Subordination Removal, Upward Extension,

Downward Extension, and their relaxed variants. [Section 4](#) examines hierarchical indices, focusing on the average depth index and its compatibility with our axioms while discussing alternative measures from the literature. [Section 5](#) concludes with a summary of our contributions and suggestions for future research. Technical results are proven in the appendices.

2. Hierarchies

A *hierarchy* is formally defined as a *directed acyclic graph*, consisting of a set of nodes connected by directed edges. In this context, nodes represent individuals, while edges represent subordination relationships, with each edge originating from a supervisor and terminating at the supervisor's direct subordinate. The acyclicity of the graph ensures that traversing any path along edge directions will never return to the starting node, meaning there are no directed cycles.

This structure enforces a unidirectional flow of authority. Individuals may have both direct supervisors (immediate predecessors) and indirect supervisors (supervisors of their supervisors), creating a transitive chain of command. Authority propagates downward from higher-level supervisors through intermediate managers to subordinates, ensuring no ambiguity in reporting lines.

The acyclicity property imposes a strict hierarchical constraint: For any individual with k total supervisors (direct and indirect), all their supervisors must necessarily have fewer than k supervisors themselves. This guarantees a monotonically decreasing authority structure, where supervisory roles are strictly ordered.

A composite hierarchy, denoted as (h_1, \dots, h_k) where $k \geq 1$, consists of k independent hierarchies. In this context, independence means each h_i is a directed acyclic graph, and no directed edges exist between nodes of distinct hierarchies h_i and h_j for $i \neq j$.

We define the *size* of a hierarchy as the cardinality of its node set, which corresponds to the total number of individuals within the hierarchical structure.

A *root node* is a node with no supervisors, representing an apex of authority in the hierarchy. A node with no subordinates is called a *leaf node*.

The *depth* of a node is the total count of its supervisors (direct or indirect), measuring its distance from the hierarchy's topmost level. Nodes closer to the root have lower depth values, while those with longer supervisory chains have greater depth.

3. Axioms for hierarchy measurement

Let \mathcal{H}_n denote the set of all possible hierarchies comprising exactly n individuals, representing all valid organizational configurations adhering to our directed acyclic graph structure.

The universal set of all possible hierarchies is then defined as:

$$\mathcal{H} = \bigcup_{n=1}^{\infty} \mathcal{H}_n.$$

Definition 1. A *hierarchical pre-order* on \mathcal{H} is a reflexive and transitive binary relation on \mathcal{H} .

A hierarchical pre-order \succsim is a binary relation that systematically ranks hierarchies by their “degree of hierarchy,” where $h' \succsim h$ signifies “ h' is no less hierarchical than h .”

This relation splits into two components: its symmetric part (\sim), where $h' \sim h$ denotes mutual comparability (equivalent hierarchies in authority structure), and its asymmetric part ($>$), where $h' > h$ indicates strict dominance (one hierarchy is unambiguously superior to another).

A hierarchy h' is called a *relabeling* of h if it is obtained by reassigning labels (identifiers) to individuals in h while preserving the structure of authority relationships.

This can be formally stated as follows. There exists a bijection ϕ from the nodes in h to those in h' such that for all nodes i, j in h ,

$$j \text{ is an immediate supervisor of } i \text{ in } h \Leftrightarrow \phi(j) \text{ is an immediate supervisor of } \phi(i) \text{ in } h'.$$

This ensures h and h' are structurally identical, differing only in individual labels.

Anonymity (A). A hierarchical pre-order \succsim on \mathcal{H} satisfies **A** if for any two hierarchies h and h' in \mathcal{H} , $h \sim h'$ whenever h' is a relabeling of h .

Thus, h and h' are indistinguishable in their hierarchical degree when labels are irrelevant.

A *replication* of a hierarchy $h \in \mathcal{H}$ is a hierarchy in \mathcal{H} of the form (h, \dots, h) . By convention, h is a replication of itself.

The Replication Principle asserts that replicating a hierarchy does not change its hierarchical degree.

Replication Principle (RP). A hierarchical pre-order \succsim on \mathcal{H} satisfies **RP** if for any two hierarchies h and h' in \mathcal{H} , $h' \sim h$ whenever h' is a replication of h .

We now introduce an axiom concerning hierarchical extensions that occur when one hierarchy is positioned above another by connecting the leaf nodes of the upper hierarchy to the root nodes of the lower hierarchy.

The operation of *appending a hierarchy h' to another hierarchy h* , denoted $h \oplus h'$, involves positioning h' beneath h such that each leaf node in h connects to one or more root nodes in h' . This connection establishes a supervisory relationship in the resulting compound hierarchy: the leaf nodes in h become direct supervisors of their corresponding root nodes in h' .

Two important properties of this “addition” operator should be noted.

First, the operation $h \oplus h'$ is order-dependent and generally non-commutative—the sequence of hierarchies determines the outcome.

Second, since there isn’t a uniquely defined way to append one hierarchy to another, $h \oplus h'$ actually defines a *set* of hierarchies that can be obtained through the appending procedure.

Upward Extension (UE). A hierarchical pre-order \succsim on \mathcal{H} satisfies **UE** if for any $h, h' \in \mathcal{H}$, we have $h' > h$ whenever $h' \in h_0 \oplus h$ for some $h_0 \in \mathcal{H}$.

This axiom states that when a hierarchy h is appended to another hierarchy h_0 , the resulting compound hierarchy—a superstructure where h_0 is positioned “on top” of h —is more hierarchical than h alone. The terminology “Upward Extension” emphasizes the upward expansion of the authority structure through the addition of higher-level nodes,

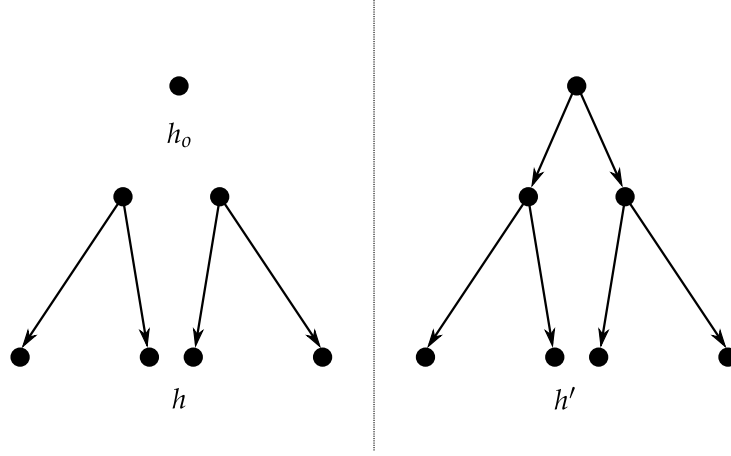


Figure 1: Illustrating UE.

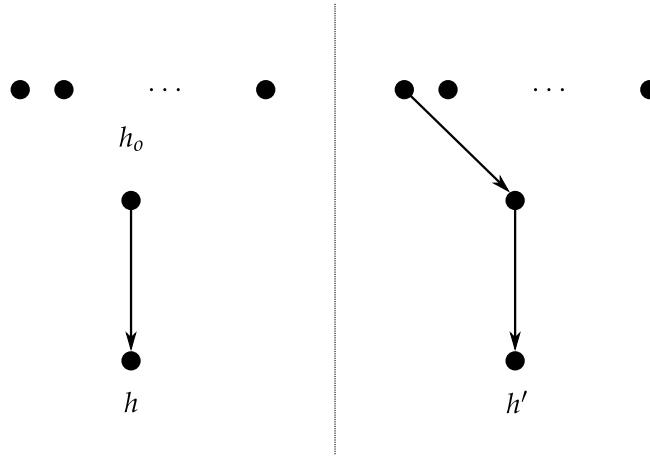


Figure 2: Illustrating UE.

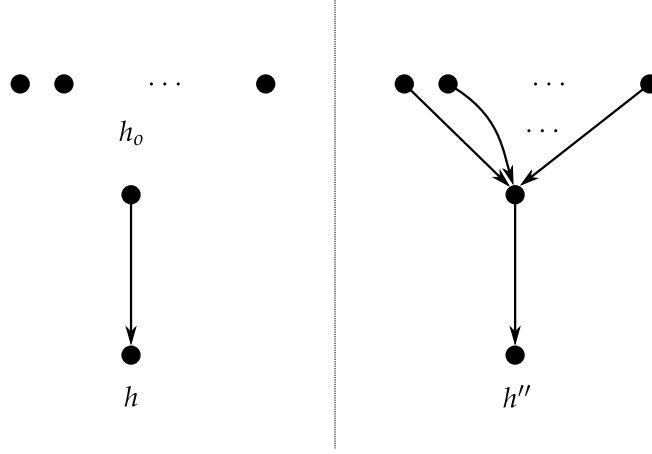
consistent with the axiom's focus on increasing hierarchical depth through vertical, upward extension.

A simple illustration of UE is provided in Figure 1, where the right panel demonstrates the result of appending hierarchy h to hierarchy h_o , which consists of a single isolated node. Note that h' belongs to the compound set $h_o \oplus h$. According to UE, h' is more hierarchical than h .

A critical requirement of this hierarchical structure is that each leaf node in h_o must connect to at least one root node in h . This constraint prevents configurations like the one shown in Figure 2, where hierarchy h is positioned below a predominantly flat hierarchy h_o , but connections between leaf nodes in h_o and root nodes in h are incomplete. Such incomplete connections create structural ambiguity, especially when h_o is extensive, as the resulting compound hierarchy h' maintains an essentially flat structure with only two isolated “anecdotal” supervisory levels at one boundary.

The proper method of appending hierarchy h to h_o is demonstrated in Figure 3, where h'' represents the resulting compound hierarchy in $h_o \oplus h$, with every leaf node in h_o connected to the (unique) root node in h . According to UE, h'' possesses a greater hierarchical degree than h alone—an intuitive conclusion given the increased depth of h'' .

Next, we establish a “more hierarchical than” pre-order based on the UE axiom. As a first step, we introduce a symmetric binary relation \sim^* on the set of n -person hierarchies,

Figure 3: Illustrating **UE**.

\mathcal{H}_n , defined as follows:

$$h \sim^* h' \Leftrightarrow h \text{ is a relabeling of } h'.$$

This relation can be extended to the domain \mathcal{H} of all hierarchies as follows: for $h, h' \in \mathcal{H}$,

$$h \sim^* h' \Leftrightarrow h_r \sim^* h'_r$$

for some equally-sized replicas h_r and h'_r of h and h' , respectively.

The proof of the following lemma is relegated to **Appendix A.1**.

Lemma 1. *The binary relation \sim^* defined on \mathcal{H} is symmetric, reflexive, and transitive.*

Next, define the following binary relation \succ_{UE} on \mathcal{H} :

$$h' \succ_{UE} h \Leftrightarrow \begin{cases} \exists h_o, H' \sim^* h', H \sim^* h : H' \in h_o \oplus H \\ \text{or} \\ h \sim^* h'. \end{cases}$$

The symmetric and asymmetric parts of \succ_{UE} are denoted by \sim_{UE} and $>_{UE}$, respectively.

We first verify that \succ_{UE} is a hierarchical pre-order satisfying **A**, **RP**, and **UE**.

Proposition 1. *The binary relation \succ_{UE} is a hierarchical pre-order on \mathcal{H} satisfying **A**, **RP**, and **UE**.*

The proof of **Proposition 1** is presented in **Appendix A.2**.

Given two hierarchical pre-orders \succ and \succ' on \mathcal{H} , \succ is \succ' -consistent if it satisfies the following conditions for all hierarchies $h, h' \in \mathcal{H}$:

- $h >' h' \Rightarrow h > h'$.
- $h \sim' h' \Rightarrow h \sim h'$.

The hierarchical pre-order \succ_{UE} stands as the unique coarsest hierarchical pre-order that satisfies the **A**, **RP**, and **UE** axioms.

Theorem 1. *A hierarchical pre-order on \mathcal{H} satisfies **A**, **RP**, and **UE** if and only if it is \succ_{UE} -consistent.*

Proof. (Sufficiency.) Assume that \succsim is a \succsim_{UE} -consistent hierarchical pre-order on \mathcal{H} . According to **Proposition 1**, \succsim_{UE} satisfies the axioms **A**, **RP**, and **UE**. Since \succsim is \succsim_{UE} -consistent, it follows that \succsim must also satisfy these three axioms.

(Necessity.) Suppose that \succsim is a hierarchical pre-order on \mathcal{H} satisfying axioms **A**, **RP**, and **UE**. We need to show that \succsim is \succsim_{UE} -consistent, which requires proving that for every pair of hierarchies $h, h' \in \mathcal{H}$:

- (a) $h \succ_{UE} h' \Rightarrow h \succ h'$.
- (b) $h \sim_{UE} h' \Rightarrow h \sim h'$.

Suppose first that $h \sim_{UE} h'$. Then there are two possibilities: either $h \sim^* h'$ or

$$\exists h_0, H' \sim^* h', H \sim^* h : H' \in h_0 \oplus H \quad \text{and} \quad \exists h'_0, H''' \sim^* h', H'' \sim^* h : H'' \in h'_0 \oplus H'''.$$

However, the second case is impossible. Indeed, $H' \in h_0 \oplus H$ implies that the depth of H' exceeds that of H , while $H'' \in h'_0 \oplus H'''$ implies that the depth of H'' exceeds that of H''' . Since

$$H' \sim^* h' \sim^* H''' \quad \text{and} \quad H \sim^* h \sim^* H'',$$

we have

$$H' \sim^* H''' \quad \text{and} \quad H \sim^* H''$$

(**Lemma 1**). Consequently, because replication and relabeling preserves depth, it follows that the depth of H' and H''' exceeds the depth of H and H'' , a contradiction.

Thus, $h \sim_{UE} h'$ implies $h \sim^* h'$, so there exist replicas h_r and h'_r of h and h' , respectively, such that h_r is a relabeling of h'_r .

Given that \succsim satisfies axioms **A** and **RP**, we have:

$$h \sim h_r \sim h'_r \sim h'. \tag{1}$$

Since \succsim is reflexive and transitive, \sim is transitive (**Sen, 2017**, Lemma 1*a, p. 56). Therefore, (1) implies $h \sim h'$, which establishes (b).

Now suppose that $h \succ_{UE} h'$. Then

$$\exists h_0, H' \sim^* h', H \sim^* h : H \in h_0 \oplus H'.$$

Consequently, because \succsim satisfies the **UE** axiom, we see that

$$H \succ H'. \tag{2}$$

Now, since $h \sim^* H$, there are replicas h_r and H_r of h and H , respectively, such that h_r is a relabeling of H_r . Similarly, there are replicas H'_r and h'_r of H' and h' , respectively, such that H'_r is a relabeling of h'_r .

Hence, since \succsim satisfies the **A** and **RP** axioms, it follows that

$$h \sim h_r \sim H_r \sim H \quad \text{and} \quad h' \sim h'_r \sim H'_r \sim H'.$$

Combining these relations with (2) yields

$$h \sim h_r \sim H_r \sim H \succ H' \sim H'_r \sim h'_r \sim h'.$$

Since \succsim is reflexive and transitive, this chain of relations implies $h \succ h'$ (Sen, 2017, Lemma 1*a, p. 56). This establishes (a) and completes the proof. ■

The hierarchical pre-order \succsim_{UE} can be compared to the core pre-order introduced in Carbonell-Nicolau (2025b), defined as follows.

For any two hierarchies h and h' in \mathcal{H}_n , $h \succsim_H h'$ if and only if there exists a bijection ϕ from the set of individuals in h to the set of individuals in h' satisfying the following condition:

- For every individual i in h such that $\phi(i)$ has at least one supervisor in h' , any immediate supervisor j' of $\phi(i)$ in h' must satisfy one of the following:
 - (a) If there exists an alternative indirect path from j' to $\phi(i)$ in h' besides the direct edge from j' to $\phi(i)$, then $\phi^{-1}(j')$ must be a direct supervisor of i in h .
 - (b) Otherwise, $\phi^{-1}(j')$ must be either a direct or indirect supervisor of i in h .

This relation can be extended to the domain \mathcal{H} as follows: for $h, h' \in \mathcal{H}$, $h' \succsim_H h$ if and only if there exist equally-sized replications h_r and h'_r of h and h' , respectively, such that $h'_r \succsim_H h_r$. This extension constitutes a properly defined hierarchical pre-order on \mathcal{H} (Carbonell-Nicolau, 2025b).

As shown in Carbonell-Nicolau (2025b), \succsim_H is intimately related to the notion of subordination removal, a concept first formalized in the context of hierarchical trees by Carbonell-Nicolau (2025a) and subsequently generalized by Carbonell-Nicolau (2025b), who recasts the mechanism within the broader mathematical structure of directed acyclic graphs.

Definition 2. We say that hierarchy h' is obtained from hierarchy h by *removing a subordination relation* if there exist a subordinate i in h and an immediate supervisor j of i satisfying one of the following conditions:

- If j has no supervisors in h , then h' is the hierarchy obtained by removing the directed edge from j to i while preserving all other relationships. Specifically:
 - The sub-hierarchy rooted at node i , $h(i)$, is no longer under j 's direct supervision.
 - Individual i loses exactly one direct supervisor (namely, j).
 - The structure of sub-hierarchy $h(i)$ remains unchanged.
 - All other supervisory relationships in h are preserved in h' .
- If j has at least one supervisor in h , then h' is the hierarchy obtained by:
 - Removing the directed edge from j to i .
 - For each direct supervisor j' of j in h , either
 - * adding a new directed edge from j' to i if no (direct or indirect) path from j' to i exists in h besides the path through j ; or
 - * making no change if there already exists an alternative path (besides the path through j) from j' to i in h .
 - Preserving the structure of sub-hierarchy $h(i)$.

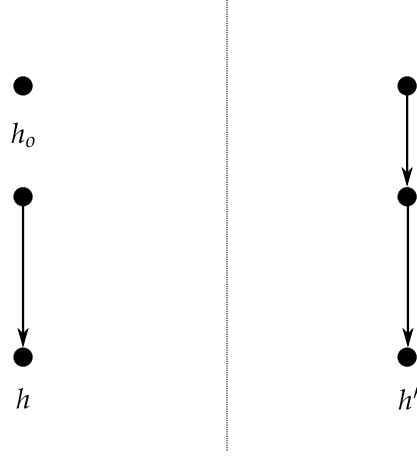


Figure 4: \succ_{UE} versus \succ_H .

- Preserving all other supervisory relationships in h .

This definition formalizes the process of subordination relation removal in two scenarios:

- If supervisor j has no superiors in the hierarchy, the subordination link $j \rightarrow i$ is completely eliminated.
- When j has superiors, then, for each immediate superior j' of i , a direct link $j' \rightarrow i$ is established (bypassing j entirely), only if there is no alternative path (besides the path through j) from j' to i in h .

Subordination Removal (SR). A hierarchical pre-order \succ on \mathcal{H} satisfies **SR** if for any two hierarchies h and h' in \mathcal{H} , $h > h'$ whenever h' is obtained from h by removing a subordination relation.

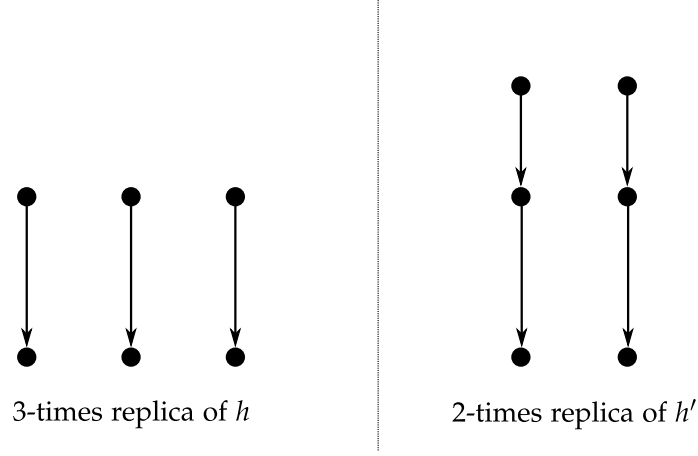
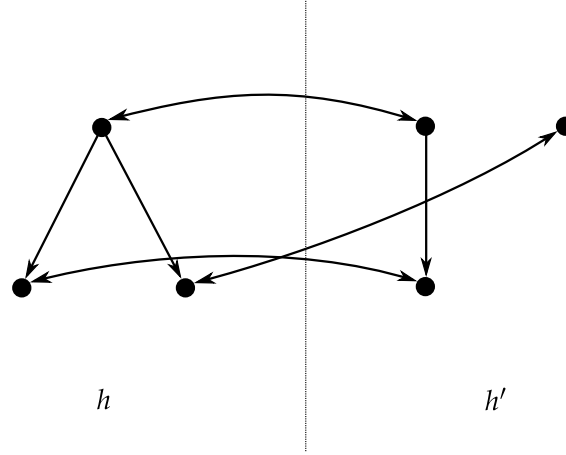
Carbonell-Nicolau (2025b) demonstrates that the strict dominance relation $h >_H h'$ precisely characterizes hierarchies obtainable through successive removals of subordination relationships. Additionally, \succ_H -consistent hierarchical pre-orders can be fully characterized by the **A**, **RP**, and **SR** axioms.

The core measures \succ_{UE} and \succ_H capture different aspects of hierarchical structure. While \succ_H focuses on the fundamental concept of subordination removal, \succ_{UE} specifically responds to the superposition of hierarchies through the appending operation, which enhances a structure's vertical dimension.

This complementarity between the two measures is illustrated in Figure 4, where h' results from appending h to h_o , establishing that $h' >_{UE} h$. However, h' and h are not comparable under \succ_H . To understand why, consider Figure 5, which shows a 3-times replica of h in the left panel and a 2-times replica of h' in the right panel.

Observe that it is impossible to construct a bijection between the nodes of these two replicas that would make the 2-times replica more hierarchical than the 3-times replica under \succ_H . This is because at least one root node in the latter hierarchy must connect to a subordinate in the former hierarchy whose immediate supervisor corresponds to a non-supervisor of that root node.

Conversely, situations exist where $h >_H h'$ while h and h' remain non-comparable under \succ_{UE} . For a simple illustration, consider the hierarchies depicted in Figure 6. Using

Figure 5: \succsim_{UE} versus \succsim_H .Figure 6: \succsim_{UE} versus \succsim_H .

the bijection indicated by the double-headed correspondence between nodes across hierarchies shown in the figure, we can easily verify that $h \succ_H h'$. However, h cannot be derived from h' through composition of measures in $h_0 \oplus h'$ for any h_0 , since such a compound measure would necessarily contain at least four nodes.

Since these two measures of hierarchy capture distinct aspects of hierarchical structure, combining them—by considering the transitive closure of their union—yields a more comprehensive metric.

Let $R = \succsim_H \cup \succsim_{UE}$ denote the union of the two relations, defined as:

$$hRh' \Leftrightarrow h \succsim_H h' \text{ or } h \succsim_{UE} h'.$$

Let I and P represent the symmetric and asymmetric parts of R , respectively.

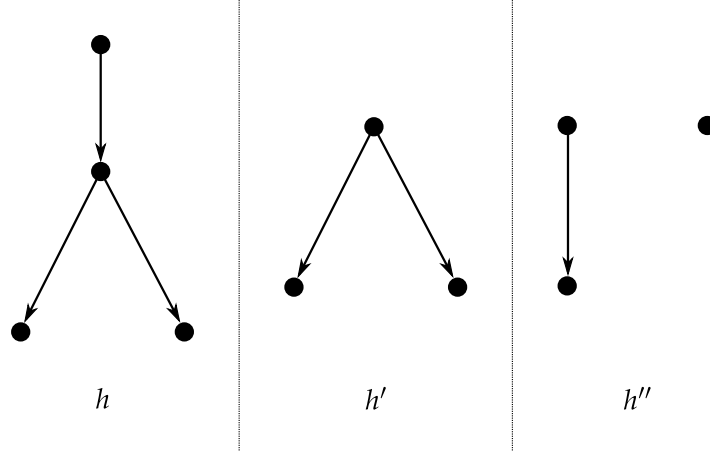
Lemma 2. *R has no cycles except those formed entirely by indifference (I).*

Proof. Suppose that

$$h_1 R \cdots R h_m R h_1.$$

Observe that

$$h_1 R \cdots R h_m$$

Figure 7: Non-transitivity of R .

implies $I_s(h_1) \geq I_s(h_m)$, where $I_s(h'')$ represents the average number of (direct or indirect) supervisors in h'' . This is clearly the case if $R = \succ_{UE}$. If $R = \succ_H$, the inequality follows from the fact (proven in Carbonell-Nicolau (2025b)) that the hierarchical order \succ_{I_s} on \mathcal{H} defined by

$$h \succ_{I_s} h' \Leftrightarrow I_s(h) \geq I_s(h'),$$

is \succ_H -consistent, so that $h \succ_H h'$ implies $I_s(h) \geq I_s(h')$.

Similarly,

$$h_m R h_1$$

implies $I_s(h_m) \geq I_s(h_1)$.

Consequently, h_1 and h_m have the same average depth, $I_s(h_1) = I_s(h_m)$, implying that

$$h_1 I \dots I h_m I h_1.$$

Indeed, if at least one of the relations were strict, one of the following would hold, since strict dominance in either \succ_H or \succ_{UE} forces a strict increase in I_s :¹

- $I_s(h_m) < I_s(h_1)$ and $I_s(h_1) \leq I_s(h_m)$. This is a contradiction.
- $I_s(h_m) \leq I_s(h_1)$ and $I_s(h_1) < I_s(h_m)$. This is a also contradiction. ■

Although both \succ_H and \succ_{UE} are individually transitive, their union R need not be. To restore transitivity, we take the transitive closure R^* , which is the smallest (coarsest) transitive relation containing R .

To see that the union relation R does not generally inherit transitivity from its transitive components, consider the three hierarchies illustrated in Figure 7. We can observe that:

$$h \succ_{UE} h' \succ_H h''.$$

However, it can be verified that neither $h \succ_H h''$ nor $h \succ_{UE} h''$ holds.

The transitive closure R^* can be defined as follows: $h R^* h'$ if and only if there exists a sequence of hierarchies

$$h = h_0, h_1, \dots, h_m = h'$$

¹In the case of \succ_H , this follows from \succ_H -consistency of \succ_{I_s} .

such that $h_k R h_{k+1}$ for each $k \in \{0, \dots, m-1\}$.

Since both \succ_H and \succ_{UE} are transitive relations, R^* can be equivalently defined as: $h R^* h'$ if and only if there exists a sequence of hierarchies

$$h = h_0, h_1, \dots, h_m = h'$$

where consecutive pairs in the sequence alternate between the relations \succ_H and \succ_{UE} , following either pattern:

$$h_0 \succ_H h_1 \succ_{UE} h_2 \succ_H \dots$$

or

$$h_0 \succ_{UE} h_1 \succ_H h_2 \succ_{UE} \dots$$

The symmetric and asymmetric parts of R^* are denoted by I^* and P^* respectively.

The transitive closure R^* is reflexive and transitive and satisfies the four axioms **A**, **RP**, **UE**, and **SR**.

Proposition 2. *The binary relation R^* is a hierarchical pre-order on \mathcal{H} satisfying **A**, **RP**, **UE**, and **SR**.*

Proof. R^* is transitive by definition of transitive closure and it is reflexive by reflexivity of \succ_H and \succ_{UE} . It satisfies **A**, **RP**, **UE**, and **SR** because it contains the union $R = \succ_H \cup \succ_{UE}$ and \succ_H satisfies **A**, **RP**, and **SR** (Carbonell-Nicolau, 2025b) and \succ_{UE} satisfies **A**, **RP**, and **UE** (Proposition 1). ■

R^* -consistency is necessary and sufficient for a hierarchical pre-order on \mathcal{H} to satisfy **A**, **RP**, **UE**, and **SR**.

Theorem 2. *A hierarchical pre-order on \mathcal{H} satisfies **A**, **RP**, **UE**, and **SR** if and only if it is R^* -consistent.*

The proof of this result is relegated to **Appendix A.3**

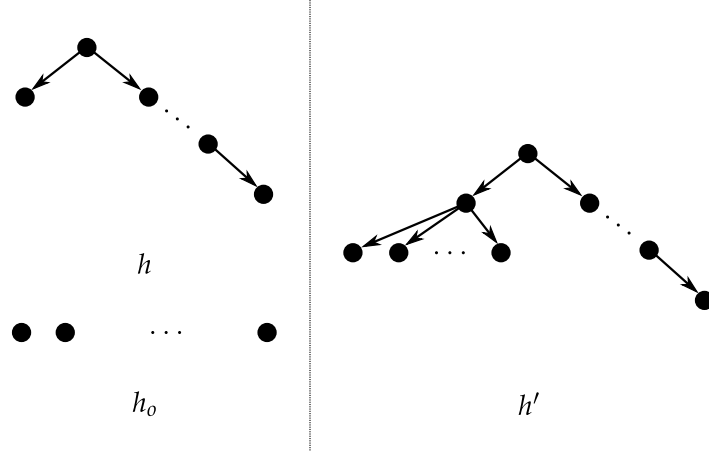
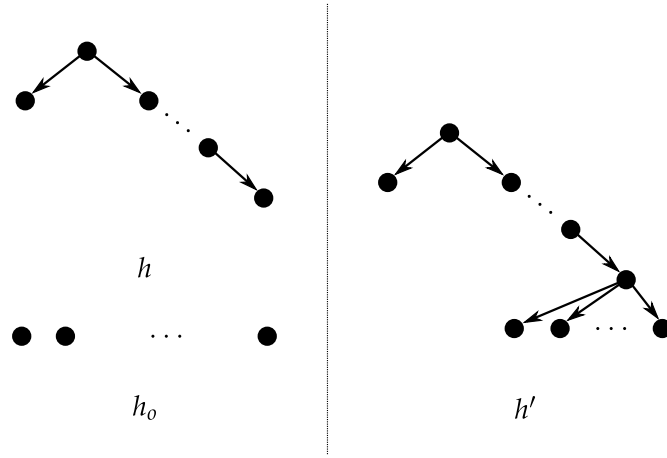
We now introduce a new axiom concerning hierarchical extensions that occur when one hierarchy is positioned below another by connecting certain leaf nodes of the upper hierarchy to the root nodes of the lower hierarchy.

This axiom requires a concept of “appending” that differs from the “ \oplus ” operator introduced earlier. A hierarchy obtained by positioning h' beneath h such that every root node in h' is connected to one or more leaf nodes in h with maximum depth within h , in such a way that the resulting compound structure is a directed acyclic graph, is said to result from the “addition” operation “ $h \boxplus h'$ ”.

Like the “ \oplus ” operator, this appending operation is order-dependent and generally non-commutative. Additionally, $h \boxplus h'$ defines a set of possible hierarchies that can result from this appending procedure, as there are multiple ways to connect the hierarchies.

Downward Extension (DE). A hierarchical pre-order \succ on \mathcal{H} satisfies **DE** if for any $h, h' \in \mathcal{H}$, we have $h' \succ h$ whenever $h' \in h \boxplus h_o$ for some $h_o \in \mathcal{H}$.

This axiom states that when a hierarchy h_o is appended to another hierarchy h via the operator \boxplus , the resulting compound hierarchy—a structure where h_o is positioned below h —is more hierarchical than h alone. The term “Downward Extension” reflects how the authority structure expands downward through the addition of lower-level nodes.

Figure 8: $h' \notin h \boxplus h_0$.Figure 9: $h' \in h \boxplus h_0$.

We illustrate the operation \boxplus in [Figure 8](#) and [Figure 9](#).

[Figure 8](#) demonstrates a case where a purely horizontal hierarchy, h_0 , is appended to hierarchy h by connecting to a leaf node in h that does not possess maximum depth within h . The resulting compound hierarchy, h' , is not a member of $h \boxplus h_0$, as it violates the maximum depth requirement.

In contrast, [Figure 9](#) shows the proper application of \boxplus , where hierarchy h_0 is appended specifically to a leaf node in h that has maximum depth within h . This creates a compound structure that is a valid member of $h \boxplus h_0$, satisfying the condition required by our axiom.

According to [DE](#), the hierarchy h' from [Figure 9](#) is definitively more hierarchical than h (i.e., $h' > h$), while the axiom makes no determination about the relative hierarchical nature of the structures in [Figure 8](#). This distinction highlights the specific scope of the Downward Extension axiom, which only applies when appending occurs at maximum-depth leaf nodes.

The binary relation associated with the [DE](#) axiom, denoted \succ_{DE} , can be formally defined on \mathcal{H} in a manner analogous to \succ_{UE} :

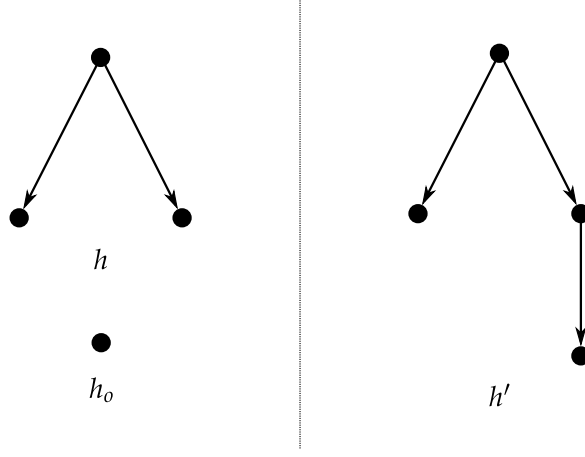


Figure 10: $h' >_{DE} h$ but $h' \not\approx_{UE} h$ and $h' \not\approx_H h$.

$$h' \succ_{DE} h \Leftrightarrow \begin{cases} \exists h_0, H' \sim^* h', H \sim^* h \text{ such that } H' \in H \boxplus h_0 \\ \text{or} \\ h \sim^* h'. \end{cases}$$

We denote the symmetric part of \succ_{DE} by \sim_{DE} and its asymmetric part by $>_{DE}$.

We note that there exist pairs of hierarchies in \mathcal{H} that are comparable under \succ_{DE} but not under either \approx_{UE} or \approx_H .

To demonstrate this, consider the hierarchies depicted in [Figure 10](#). We can observe that $h' \in h \boxplus h_0$, which by definition implies $h' >_{DE} h$. However, there exists no hierarchy h^* such that h' belongs to $h^* \oplus h$, meaning $h' \not\approx_{UE} h$.

Furthermore, examining all possible replicas of h and h' (denoted H and H' respectively), we can verify that H' cannot be obtained from H through any sequence of subordination removals. This implies $h' \not\approx_H h$.

Similarly, both hierarchies are uncomparable under R^* .

This example illustrates how the Downward Extension relation captures hierarchical comparisons that remain undetected by both \approx_{UE} and \approx_H .

The analysis of the properties of \succ_{DE} follows a parallel structure to our earlier examination of \approx_{UE} . We establish that \succ_{DE} is a hierarchical pre-order satisfying [A](#), [RP](#), and [DE](#).

Proposition 3. *The binary relation \succ_{DE} is a hierarchical pre-order on \mathcal{H} satisfying [A](#), [RP](#), and [DE](#).*

The complete proof of this proposition can be found in [Appendix A.4](#).

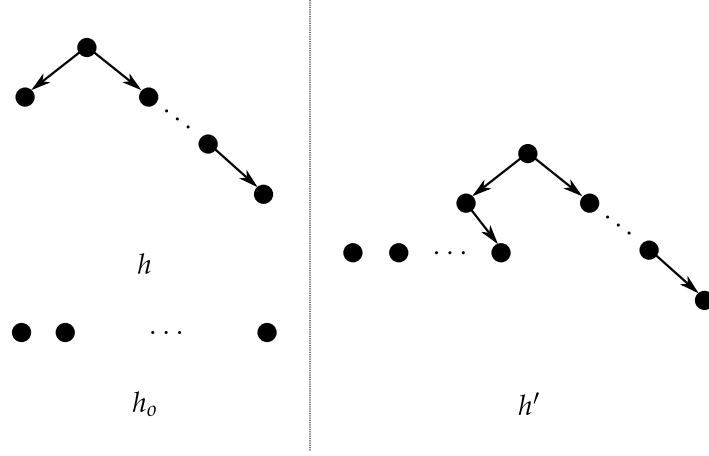
Furthermore, we can characterize \succ_{DE} as the unique coarsest hierarchical pre-order that simultaneously satisfies the [A](#), [RP](#), and [DE](#) axioms.

Theorem 3. *A hierarchical pre-order on \mathcal{H} satisfies [A](#), [RP](#), and [DE](#) if and only if it is \succ_{DE} -consistent.*

The detailed proof of [Theorem 3](#) is provided in [Appendix A.5](#).

Next, we introduce two additional axioms that serve as variants of [UE](#) and [DE](#). These new formulations capture distinct aspects of hierarchical structure that were not addressed by the original axioms.

Before presenting these axioms, we must define a more generalized notion of “appending.” For any two hierarchies $h, h' \in \mathcal{H}$, we denote by $h \uplus h'$ the set of all possible

Figure 11: $h' \in h \uplus h_0$.

hierarchies in \mathcal{H} that can be formed by positioning h' beneath h such that at least one leaf node in h connects to at least one root node in h' .

This operation can be illustrated by referring to [Figure 8](#), which demonstrates an instance where hierarchy h_0 is appended to hierarchy h via the \uplus operation, yielding $h' \in h \uplus h_0$. Similarly, [Figure 11](#) provides another example where h_0 is positioned below h , resulting in a compound hierarchy that belongs to the set $h \uplus h_0$.

It is important to observe that in neither of the examples mentioned does the resulting hierarchy h' belong to either $h \oplus h_0$ or $h \boxplus h_0$.

To clarify this distinction, consider [Figure 8](#). Here, $h' \notin h \oplus h_0$ because the \oplus operator requires *every* leaf node in h to connect to a root node in h_0 —a condition clearly violated in this figure, where only a subset of leaf nodes from h establish connections with h_0 .

Similarly, $h' \notin h \boxplus h_0$ since the connecting leaf node in h does not possess maximum depth within h , which is a defining requirement of the \boxplus operation.

We are now ready to formulate the last two axioms.

Upward Extension* (UE*). A hierarchical pre-order \succsim on \mathcal{H} satisfies **UE*** if for any $h, h' \in \mathcal{H}$, we have $h' \succ h$ whenever $h' \in h_0 \uplus h$ and $h_0 \succsim h$ for some $h_0 \in \mathcal{H}$.

Downward Extension* (DE*). A hierarchical pre-order \succsim on \mathcal{H} satisfies **DE*** if for any $h, h' \in \mathcal{H}$, we have $h' \succ h$ whenever $h' \in h \uplus h_0$ and $h_0 \succsim h$ for some $h_0 \in \mathcal{H}$.

It is instructive to compare **UE*** with the earlier **UE** axiom, as they represent different perspectives on hierarchical growth.

UE establishes that adding any hierarchy h_0 on top of h increases the hierarchical degree, provided that all leaf nodes of h_0 are connected to the roots of h . This can be conceptualized as stacking a new layer of authority (e.g., a new management tier) on top of an existing organization. Because every leaf of h_0 is linked to h , the resulting structure h' has greater depth—more levels of authority—making it more hierarchical. Importantly, no condition is placed on the hierarchical nature of h_0 itself; it can be flat, deep, or anything in between. The key requirement is the complete connection through all leaves of h_0 .

In contrast, **UE*** asserts that adding a hierarchy h_0 on top of h increases the hierarchical degree under the following two conditions:

- h_0 must itself be (weakly) more hierarchical than h .

- The connection between h_o and h requires only partial linkage (not all leaves of h_o need to connect to roots of h).

This can be visualized as attaching a highly hierarchical structure (e.g., a tall, well-organized management team) to an existing organization, even if the attachment is partial (e.g., only one manager oversees the original group). Because h_o is already more hierarchical than h , even a loose connection elevates the overall hierarchical nature of the resulting structure h' .

Thus, while **UE** focuses on vertical growth through complete attachment—where adding any new layer on top, fully connected, deepens the hierarchy—**UE*** emphasizes the quality of the added structure, allowing a sufficiently hierarchical h_o to enhance h even with partial linkage.

Similarly, we can compare **DE** with **DE***, which mirror the relationship between **UE** and **UE*** but in the downward direction.

DE states that extending a hierarchy h downward by adding any hierarchy h_o increases the hierarchical degree, provided that every root node of h_o connects to leaf nodes of h having maximum depth. This represents a downward extension of authority by attaching subordinates specifically to the lowest-level positions in the existing structure. The resulting hierarchy h' gains additional levels of subordination, thus becoming more hierarchical. As with **UE**, no restrictions are placed on the hierarchical nature of h_o itself—the critical factor is the complete connection to the deepest leaves of h .

Conversely, **DE*** takes a different approach, stipulating that adding h_o below h increases hierarchical degree when:

- h_o is (weakly) more hierarchical than h .
- Only partial connection is required between h and h_o (neither all roots of h_o need to connect, nor must connections be to maximum-depth leaves).

This corresponds to integrating a highly structured group of subordinates into an organization, even through limited points of attachment. The superior hierarchical quality of h_o ensures that even with minimal connection, the resulting structure h' becomes more hierarchical than h alone.

In essence, while **DE** prioritizes complete integration at the deepest level to achieve vertical growth, **DE*** focuses on the hierarchical quality of what is being appended, requiring only minimal attachment to enhance the overall structure.

Our next task is to develop essential hierarchical pre-orders that align with the **A**, **RP**, and **UE*** (respectively, **DE***) axioms.

Consider the binary relations on \mathcal{H} defined recursively as follows:

- $R_0 = \sim^*$.
- For each $k \in \{1, 2, \dots\}$,

$$R_k = R_{k-1} \cup \{(h', h) : \exists h_0, \dots, h_m = h : \forall l \in \{0, \dots, m-1\}, \\ h_l R_{k-1} h_{l+1} \ \& \ \exists H' \sim^* h', H \sim^* h : H' \in h_0 \uplus H\}.$$

We now define \succsim_{UE^*} as the transitive closure of the union

$$\bigcup_{k=0}^{\infty} R_k.$$

Equivalently, $h \succsim_{UE^*} h'$ if and only if there exists a finite sequence

$$h = h_0, h_1, h_2, \dots, h_m = h',$$

where each consecutive pair (h_l, h_{l+1}) satisfies $h_l R_k h_{l+1}$ for some $k \geq 0$.

As usual, the symmetric part of \succsim_{UE^*} is denoted by \sim_{UE^*} and its asymmetric part by $>_{UE^*}$.

The relation \succsim_{UE^*} is a proper hierarchical pre-order satisfying **A**, **RP**, and **UE***.

Proposition 4. *The binary relation \succsim_{UE^*} is a hierarchical pre-order on \mathcal{H} satisfying **A**, **RP**, and **UE***.*

The reader is referred to **Appendix A.6** for the proof of **Proposition 4**.

The relation \succsim_{UE^*} is the unique coarsest hierarchical pre-order that satisfies **A**, **RP**, and **UE***.

Theorem 4. *A hierarchical pre-order on \mathcal{H} satisfies **A**, **RP**, and **UE*** if and only if it is \succsim_{UE^*} -consistent.*

The proof of this result is relegated to **Appendix A.7**.

A similar analysis can be extended to the **DE*** axiom: we can construct the unique coarsest hierarchical pre-order that simultaneously satisfies the **A**, **RP**, and **DE*** axioms.

Define a sequence of binary relations on \mathcal{H} as follows:

- $S_0 = \sim^*$,
- For each $k \in \{1, 2, \dots\}$,

$$S_k = S_{k-1} \cup \{(h', h) : \exists h_0, \dots, h_m = h : \forall l \in \{0, \dots, m-1\}, \\ h_l S_{k-1} h_{l+1} \ \& \ \exists H' \sim^* h', H \sim^* h : H' \in H \uplus h_0\}.$$

Now define \succsim_{DE^*} as the transitive closure of the union

$$\bigcup_{k=0}^{\infty} S_k.$$

The characterizations of \succsim_{DE^*} presented below are analogous to those of \succsim_{UE^*} established in **Proposition 4** and **Theorem 4**. The proof of **Proposition 5** appears in **Appendix A.8**, while the proof of **Theorem 5** is provided in **Appendix A.9**.

Proposition 5. *The binary relation \succsim_{DE^*} is a hierarchical pre-order on \mathcal{H} satisfying **A**, **RP**, and **DE***.*

Theorem 5. *A hierarchical pre-order on \mathcal{H} satisfies **A**, **RP**, and **DE*** if and only if it is \succsim_{DE^*} -consistent.*

In the context of hierarchical structures, various axioms can be employed to capture different facets of hierarchy. We now develop a theorem that characterizes hierarchical pre-orders satisfying a collection of such axioms by relating them to a fundamental pre-order constructed as the transitive closure of the union of individual pre-orders, each corresponding to a specific axiom set. This result generalizes [Theorem 2](#).

A collection $\{\succsim_i\}_{i \in I}$ of hierarchical pre-orders in \mathcal{H} is called a *composite axiomatic collection of hierarchical pre-orders* if it satisfies the following conditions:

1. For each index i in some index set I , the pre-order \succsim_i satisfies a corresponding set of axioms A_i . Each axiom is formally defined as a conditional statement of the following form:

$$\forall h, h' \in \mathcal{H}, P(h, h', \mathcal{H}', \succsim_i) \Rightarrow Q(h, h', \succsim_i)$$

where:

- $P(h, h', \mathcal{H}', \succsim_i)$ is a predicate describing a relationship or transformation between h and h' . This predicate may depend not only on the pair (h, h') , but also on some subset \mathcal{H}' of \mathcal{H} and on the pre-order \succsim_i itself.
 - $Q(h, h', \succsim_i)$ represents one of the following relations: $h \succsim_i h'$, $h \sim_i h'$, $h' \succsim_i h$, $h >_i h'$, or $h' >_i h$.
2. Any cycle in the union $\bigcup_{i \in I} \succsim_i$ must be formed entirely by indifference in each individual \succsim_i .
 3. For some $i \in I$, the set A_i contains the [A](#) and [RP](#) axioms.
 4. For each $i \in I$:
 - If $h \sim_i h'$, then $h \sim^* h'$.
 - If $h >_i h'$, then there exists a finite sequence of hierarchies in \mathcal{H} ,

$$h = h_0, h_1, \dots, h_n = h',$$

such that for each $l \in \{0, \dots, n-1\}$, there exists some axiom $a \in A_i$ whose conditional statement for the tuple $(h_A, h'_A, \mathcal{H}', \succsim_i)$ is

$$P(h_A, h'_A, \mathcal{H}', \succsim_i) \Rightarrow Q(h_A, h'_A, \succsim_i),$$

the predicate $P(H_l, H_{l+1}, \mathcal{H}', \succsim_i)$ holds for some $H_l \sim_i h_l$ and $H_{l+1} \sim_i h_{l+1}$, and

$$Q(H_l, H_{l+1}, \succsim_i) = Q(H_l, H_{l+1}, \succsim)$$

for every hierarchical pre-order \succsim on \mathcal{H} satisfying axiom a , [A](#), and [RP](#).

Now define \tilde{R} as the transitive closure of the union of all \succsim_i , $\bigcup_{i \in I} \succsim_i$.

The following two results are proven in [Appendix A.10](#) and [Appendix A.11](#), respectively.

Proposition 6. *The binary relation \tilde{R} is a hierarchical pre-order on \mathcal{H} satisfying all the axioms in $\bigcup_{i \in I} A_i$.*

Theorem 6. A hierarchical pre-order \succsim on \mathcal{H} satisfies all A_i for $i \in I$ if and only if it is \tilde{R} -consistent.

As an example of a composite axiomatic collection of hierarchical pre-orders, consider

$$\{\succsim_H, \succsim_{UE}, \succsim_{DE}, \succsim_{UE^*}, \succsim_{DE^*}\}.$$

Proposition 7. The collection

$$\{\succsim_H, \succsim_{UE}, \succsim_{DE}, \succsim_{UE^*}, \succsim_{DE^*}\}$$

is a composite axiomatic collection of hierarchical pre-orders.

We prove this proposition in [Appendix A.12](#).

As a consequence of [Proposition 7](#) and [Theorem 6](#), we can identify the coarsest hierarchical pre-order on \mathcal{H} that satisfies the seven axioms [A](#), [RP](#), [SR](#), [UE](#), [DE](#), [UE*](#), and [DE*](#). This pre-order, denoted by \hat{R} , is defined as the transitive closure of the union

$$\succsim_H \cup \succsim_{UE} \cup \succsim_{DE} \cup \succsim_{UE^*} \cup \succsim_{DE^*}.$$

The pre-order \hat{R} is thus the minimal hierarchical relation that encompasses all ordering principles represented by these axioms.

Corollary 1. A hierarchical pre-order \succsim on \mathcal{H} satisfies [A](#), [RP](#), [SR](#), [UE](#), [DE](#), [UE*](#), and [DE*](#) if and only if it is \hat{R} -consistent.

Proof. The statement follows from [Theorem 6](#) and [Proposition 7](#). ■

4. Hierarchical indices

Having established our axiomatic theory of core pre-orders, we now extend our investigation to examine complete hierarchical orders, with a focus on one measure that quantifies average hierarchy depth. We demonstrate its compliance with all hierarchy measurement axioms introduced in this work, establishing its theoretical robustness.

We then analyze hierarchical pre-orders from prior research, situating these established measures within our theoretical framework.

Hierarchical orders offer computational advantages through simple formulae, avoiding the complexities of pre-orders that rely on transitive closures of binary relation unions. This makes them readily applicable to real-world datasets. However, by inducing a complete order on hierarchies, these measures necessarily make definitive comparisons between any two hierarchical structures—even in ambiguous cases where our axioms alone would not single out a more hierarchical structure.

A *hierarchical index* is a mapping $I : \mathcal{H} \rightarrow \mathbb{R}$ that assigns a “hierarchical degree” $I(h)$ to every hierarchy h in \mathcal{H} .

In the following subsections, we examine various hierarchical indices through the lens of our axiomatic theory of hierarchical measurement.

4.1. Average depth

A hierarchical index measuring the average depth of each hierarchy is formally defined as follows:

$$I_s(h) = \frac{1}{n} \sum_{i \in h} s_h(i),$$

where $s_h(i)$ denotes the number of (direct or indirect) supervisors of individual i in hierarchy h . This index represents the average number of supervisors per individual within the hierarchy h .

The average depth index induces a hierarchical order \succ_{I_s} defined by

$$h \succ_{I_s} h' \Leftrightarrow I_s(h) > I_s(h'),$$

which satisfies the **A**, **RP**, and **SR** axioms (Carbonell-Nicolau, 2025b).

The **UE** axiom applied to \succ_{I_s} states that for any $h, h' \in \mathcal{H}$, $h' \succ_{I_s} h$ whenever $h' \in h_o \oplus h$ for some $h_o \in \mathcal{H}$.

Observe that appending a hierarchy h to another hierarchy h_o (i.e., forming $h' \in h_o \oplus h$) strictly increases the average depth I_s . Indeed, when h is appended to h_o , every root node in h gains additional supervisors from the leaves of h_o . The depth of each node in h increases by the depth of its new supervisors in h_o plus one (for the direct edge). The total depth of h' equals the total depth of h_o plus the total depth of h adjusted upward due to the new supervisory relationships. Even if h_o is a single node, appending it increases the average depth. Hence, for any non-trivial h_o , the added supervisors increase the total depth of h' , ensuring $I_s(h') > I_s(h)$. Thus \succ_{I_s} satisfies **UE**.

The **DE** axioms asserts that if h' is obtained from a hierarchy h by attaching another hierarchy h_o below one of h 's deepest leaves, then h' is strictly more hierarchical than h (i.e. $h' \succ_{I_s} h$).

The index I_s measures the average number of supervisors per person, i.e., the average distance up the chart from each individual to the top.

Let D be the maximum depth in h . By definition, the average depth $I_s(h)$ cannot exceed D . When you attach a node of h_o under a leaf at depth D , each new node ends up at depth at least $D + 1$. Since $D + 1 > I_s(h)$, adding these nodes pulls the overall average up.

Thus any valid downward extension strictly increases the average-depth index, so \succ_{I_s} satisfies **DE**.

The **UE*** axiom requires $h' \succ_{I_s} h$ whenever $h, h' \in \mathcal{H}$ and $h' \in h_o \uplus h$ and $h_o \succ_{I_s} h$ for some $h_o \in \mathcal{H}$.

Let $h' \in h_o \uplus h$, where $h_o \succ_{I_s} h$. This means $I_s(h_o) \geq I_s(h)$, and h' is formed by appending h to h_o via \uplus (connecting at least one leaf of h_o to a root of h).

Let d be the depth of the attaching leaf node in h_o . Each node in h (originally with depth k in h) now has depth $d + 1 + k$ in h' .

Let n_o and n_h be the sizes of h_o and h . The total depth of h' is:

$$\sum_{i \in h'} s_{h'}(i) = \sum_{i \in h_o} s_{h_o}(i) + \sum_{i \in h} s_h(i) + (d + 1) \cdot n_h.$$

The average depth of h' is:

$$I_s(h') = \frac{1}{n_o + n_h} \sum_{i \in h'} s_{h'}(i) = \frac{1}{n_o + n_h} \left(\sum_{i \in h_o} s_{h_o}(i) + \sum_{i \in h} s_h(i) + (d+1) \cdot n_h \right).$$

Subtract $I_s(h)$ from $I_s(h')$:

$$I_s(h') - I_s(h) = \frac{I_s(h_o) \cdot n_o - I_s(h) \cdot n_o + (d+1) \cdot n_h}{n_o + n_h}.$$

Since $I_s(h_o) \geq I_s(h)$, the term $I_s(h_o) - I_s(h) \geq 0$. Additionally, $(d+1) \cdot n_h > 0$ (as $d \geq 0$ and $n_h \geq 1$). Thus, $I_s(h') > I_s(h)$, implying $h' \succ_{I_s} h$.

Thus, the average depth index \succ_{I_s} satisfies **UE*** because appending h to h_o (with $h_o \succ_{I_s} h$) strictly increases the average depth, ensuring $h' \succ_{I_s} h$.

One can similarly show that \succ_{I_s} satisfies **DE***.

Proposition 8. *The hierarchical order \succ_{I_s} satisfies **A**, **RP**, **SR**, **UE**, **DE**, **UE***, and **DE***.*

An immediate consequence of **Corollary 1** and **Proposition 8** is that \succ_{I_s} is \hat{R} -consistent, where \hat{R} is the transitive closure of the union

$$\succ_H \cup \succ_{UE} \cup \succ_{DE} \cup \succ_{UE*} \cup \succ_{DE*}.$$

4.2. Global reaching centrality

In this section, we situate the global reaching centrality index proposed by [Mones et al. \(2012\)](#) within our axiomatic framework, examining how this established measure aligns with our theoretical principles for hierarchy quantification.

For a hierarchy $h \in \mathcal{H}$ with a set of nodes $V(h)$ and total number of nodes $N = |V(h)|$, the *global reaching centrality* (grc) is defined as:

$$\text{grc}(h) = \frac{1}{N-1} \sum_{v \in V(h)} (C_R^{\max} - C_R(v)),$$

where:

- $C_R(v)$ represents the *reaching centrality* of node v , defined as follows:

$$C_R(v) = \frac{1}{N-1} |\{w \in V(h) \setminus \{v\} : \text{there exists a directed path from } v \text{ to } w \text{ in } h\}|.$$

It represents the fraction of other nodes that can be reached from v via directed paths.

- $C_R^{\max} = \max_{u \in V(h)} C_R(u)$ is the maximum reaching centrality among all nodes in h .

The grc index measures the degree of hierarchy within a directed acyclic graph by evaluating how concentrated the ability to reach others is among its nodes. In a highly hierarchical structure, such as an organization with a clear chain of command, a few nodes (e.g., those at the top) can reach many others through directed paths, while most nodes have limited or no reach. This results in a high grc value, as the differences $C_R^{\max} - C_R(v)$

are large for most nodes due to the significant disparity in reaching centralities. Conversely, in a flat structure where no node can reach others (e.g., a graph with no edges), all reaching centralities are zero, leading to a grc of zero, indicating the absence of hierarchy. Thus, the grc effectively quantifies the extent to which influence or control is centralized within the graph, providing a clear metric for the hierarchical nature of the system it represents.

As demonstrated in Carbonell-Nicolau (2025a), the hierarchical order \succ_{grc} induced by the grc index violates the **SR** axiom. While \succ_{grc} clearly satisfies the **A** axiom, the grc index, as defined, contradicts the Replication Principle, **RP**.

To illustrate this violation, consider hierarchy h with two nodes $v \rightarrow w$. We have:

$$\begin{aligned} N &= 2. \\ C_R(v) &= 1. \\ C_R(w) &= 0. \\ C_R^{\max} &= 1. \end{aligned}$$

Thus:

$$\text{grc}(h) = \frac{1}{1}(0 + 1) = 1$$

Now consider h' , a one-time replica of h with two copies: $v_1 \rightarrow w_1$ and $v_2 \rightarrow w_2$. For this hierarchy:

$$\begin{aligned} N' &= 4. \\ C_R(v_1) &= C_R(v_2) = \frac{1}{3}. \\ C_R(w_1) &= C_R(w_2) = 0. \\ C_R^{\max} &= \frac{1}{3}. \end{aligned}$$

Thus:

$$\text{grc}(h') = \frac{1}{3} \left(0 + \frac{1}{3} + 0 + \frac{1}{3} \right) = \frac{2}{9}.$$

Therefore, $\text{grc}(h') < \text{grc}(h)$. The **RP** axiom requires $\text{grc}(h') = \text{grc}(h)$, but the standard grc decreases with replication due to its normalization by the larger denominator.

Next, we show that \succ_{grc} also violates **UE**. Consider the hierarchy h with two nodes v and w and a single directed edge $v \rightarrow w$. Here $N = 2$. The reaching centralities are

$$C_R(v) = \frac{|\{w\}|}{N-1} = \frac{1}{1} = 1 \quad \text{and} \quad C_R(w) = 0,$$

so

$$C_R^{\max} = 1,$$

and hence

$$\text{grc}(h) = \frac{1}{N-1} \sum_{u \in h} (C_R^{\max} - C_R(u)) = \frac{1}{1} ((1-1) + (1-0)) = 1.$$

Form h'' by adding a new root x on top of h : $x \rightarrow v \rightarrow w$. Now $N' = 3$. Compute

$$C_R(x) = \frac{2}{2} = 1, \quad C_R(v) = \frac{1}{2} = 0.5, \quad C_R(w) = 0,$$

so again $C_R^{\max} = 1$, but

$$\text{grc}(h'') = \frac{1}{N' - 1} \sum_{u \in h''} (C_R^{\max} - C_R(u)) = \frac{1}{2} ((1 - 1) + (1 - 0.5) + (1 - 0)) = \frac{1.5}{2} = 0.75.$$

Since

$$\text{grc}(h'') = 0.75 < 1 = \text{grc}(h),$$

the grc-induced order \succ_{grc} violates the **UE** axiom.

This example reveals a fundamental limitation of the grc measure when assessing vertical hierarchical growth. While h'' contains a deeper vertical structure with three levels ($x \rightarrow v \rightarrow w$) compared to the simpler two-level hierarchy in h ($v \rightarrow w$), the grc measure counterintuitively decreases.

The issue stems from how grc quantifies hierarchical inequality through the distribution of reaching centrality. In hierarchy h , node v holds all possible reaching centrality (perfect concentration), reaching 100% of other nodes, while w reaches none. In contrast, in h'' , the reaching power becomes more distributed: x reaches all other nodes, but v still reaches 50% of the network. This distribution of reaching power is interpreted by grc as a decrease in hierarchical inequality, even though the hierarchy has actually become deeper and more stratified.

This counterintuitive behavior demonstrates why grc fails to consistently capture vertical hierarchical growth, prioritizing concentration of reaching power over the development of multi-level command chains.

A similar example can be constructed to demonstrate that \succ_{grc} also violates the **DE** axiom. When adding subordinates to lower-level nodes, grc can increase or decrease depending on how the addition affects the overall distribution of reaching centrality, rather than consistently capturing the deepening of hierarchical relationships.

Furthermore, slight modifications of these examples can be used to prove that \succ_{grc} fails to satisfy the **UE*** and **DE*** axioms.

Specifically for the case of **UE***, consider the following example.

For the hierarchy h , $v \rightarrow w$, we have $\text{grc}(h) = 1$.

Now form h''' by appending h to $h_0 = x \rightarrow y$, resulting in $x \rightarrow y \rightarrow v \rightarrow w$. Reaching centralities:

$$C_R(x) = 1 \text{ (reaches all nodes).}$$

$$C_R(y) = \frac{2}{3} \text{ (reaches } v, w \text{).}$$

$$C_R(v) = \frac{1}{3} \text{ (reaches } w \text{).}$$

$$C_R(w) = 0.$$

Hence,

$$\text{grc}(h''') = \frac{(1 - 1) + (1 - \frac{2}{3}) + (1 - \frac{1}{3}) + (1 - 0)}{4 - 1} = \frac{0 + \frac{1}{3} + \frac{2}{3} + 1}{3} = \frac{2}{3},$$

and so

$$\text{grc}(h''') = \frac{2}{3} < \text{grc}(h) = 1.$$

Since appending h_o (which is equally hierarchical to h) reduces grc , \succ_{grc} violates **UE***.

Proposition 9. *The hierarchical order \succ_{grc} satisfies **A** and violates **RP**, **SR**, **UE**, **DE**, **UE***, and **DE***.*

4.3. Krackhardt's graph theoretical dimensions

Krackhardt (1994) identified four distinct dimensions for measuring the “hierarchical degree” in networks of informal organizations. In the following sections, we examine each of these dimensions in detail.

4.3.1. Connectedness

The first dimension, connectedness, quantifies the extent to which the network is connected as a whole.

In Krackhardt's original definition, connectedness measures the degree to which all nodes in the network are connected through paths, regardless of direction. Specifically, it is the fraction of pairs of nodes that are mutually reachable when the directed graph is treated as undirected.²

Formally, within our framework, the *connectedness index* for a hierarchy $h \in \mathcal{H}$ can be formulated as follows:

$$C(h) = \frac{|\{(i, j) \in V \times V : i \neq j \text{ and there exists an undirected path between } i \text{ and } j\}|}{n(n-1)},$$

where:

- V is the set of all nodes in h .
- $n = |V|$ is the total number of nodes in h .
- The numerator counts all ordered pairs (i, j) where j is reachable from i via at least one path when direction is ignored.
- The denominator $n(n-1)$ represents the total number of possible ordered pairs (excluding self-loops).

When applying this definition to directed acyclic graphs (which represent hierarchies in our framework), we observe that connectedness measures the degree to which the hierarchy forms a connected structure rather than separate disconnected components. For

²Treating the graph as undirected for the connectedness index overlooks the directionality of information flows, leading to potential information leaks.

For example, consider a simple chain where $j \rightarrow i$. Information flows unidirectionally from j to i (e.g., from supervisor to subordinate), but not vice versa, reflecting the hierarchical nature of supervision. In directed acyclic graphs, information moves solely from supervisors to subordinates along defined paths. In contrast, networks with cycles allow bidirectional information flows. A refined connectedness measure should account for the directionality of these information pathways to better capture the network's structure.

non-trivial connected hierarchies (i.e., $n > 1$), $C(h) = 1$, while for disconnected hierarchies, $0 < C(h) < 1$.

The hierarchical order induced by the connectedness index is denoted by \succ_C .

The connectedness index $C(h)$ measures the proportion of node pairs (i, j) that are mutually reachable when treating the directed graph as undirected. Relabeling nodes does not affect the existence or count of such paths—it only changes node labels. Thus, relabeling preserves the connectedness index, so \succ_C treats structurally identical hierarchies as equivalent, fulfilling the **A** axiom.

To demonstrate that the connectedness-based hierarchical order \succ_C violates the **RP** axiom, we present the following counterexample.

Consider hierarchy h , a simple two-person structure with a single directed edge $j \rightarrow i$. With $n = 2$ nodes, these two nodes are connected when direction is ignored, giving us one unordered pair with an undirected path. The connectedness index of h is therefore:

$$C(h) = \frac{|\{(i, j) \in V \times V : i \neq j \text{ and there exists an undirected path between } i \text{ and } j\}|}{n(n-1)} = \frac{2}{2(2-1)} = \frac{2}{2} = 1.$$

Now consider hierarchy h' , formed by creating two independent copies of h : $j_1 \rightarrow i_1$ and $j_2 \rightarrow i_2$. This replicated hierarchy contains $n' = 4$ nodes. When treated as an undirected graph, h' has two connected components. The node pairs with undirected paths are (i_1, j_1) , (j_1, i_1) , (i_2, j_2) and (j_2, i_2) , giving a connectedness index of:

$$C(h') = \frac{4}{n'(n'-1)} = \frac{4}{4(4-1)} = \frac{4}{12} = \frac{1}{3}.$$

Since $C(h) = 1 > \frac{1}{3} = C(h')$, we must conclude that $h \succ_C h'$ according to the connectedness-based ordering. However, the **RP** axiom explicitly requires that $h \sim_C h'$ (the original hierarchy and its replication should be equivalent). This contradiction demonstrates that \succ_C violates the **RP** axiom.

The **SR** axiom stipulates that whenever a hierarchy h' is derived from h by removing a subordination relation, then h must be considered more hierarchical than h' .

To demonstrate that the hierarchical order \succ_C based on Krackhardt's connectedness index violates the **SR** axiom, consider the following example:

- Hierarchy \hat{h} : A chain structure with nodes $j \rightarrow i \rightarrow k$.
- Hierarchy \tilde{h} : A star structure with edges $j \rightarrow i$ and $j \rightarrow k$.

Hierarchy \tilde{h} is obtained from \hat{h} by removing the subordination relation $i \rightarrow k$ and adding the relation $j \rightarrow k$. Since this transformation involves removing a subordination relation, the **SR** axiom requires that $\hat{h} \succ_C \tilde{h}$.

Let us calculate the connectedness index for both hierarchies:

For hierarchy \hat{h} (chain structure):

- When treated as an undirected graph, all three nodes form a single connected component.

- All 6 possible ordered pairs have undirected paths between them: (j, i) , (i, j) , (i, k) , (k, i) , (j, k) , and (k, j) .
- Therefore, $C(\hat{h}) = \frac{6}{3(3-1)} = \frac{6}{6} = 1$.

For hierarchy \tilde{h} (star structure):

- When treated as an undirected graph, all three nodes also form a single connected component.
- All 6 possible ordered pairs have undirected paths between them: (j, i) , (i, j) , (j, k) , (k, j) , (i, k) , and (k, i) .
- Therefore, $C(\tilde{h}) = \frac{6}{3(3-1)} = \frac{6}{6} = 1$.

Since $C(\hat{h}) = C(\tilde{h}) = 1$, we have $\hat{h} \sim_C \tilde{h}$ according to the connectedness-based ordering. However, the **SR** axiom requires $\hat{h} \succ_C \tilde{h}$. This contradiction demonstrates that \succ_C violates the **SR** axiom.

The fundamental issue is that Krackhardt's connectedness index measures only whether nodes can reach each other when direction is ignored, not the specific structure of subordination relationships. In both hierarchies, all nodes can reach all other nodes when direction is ignored, resulting in identical connectedness values despite their different hierarchical structures.

This violation makes intuitive sense: connectedness, as defined by Krackhardt, primarily captures whether a network forms a cohesive whole rather than the degree of hierarchical subordination within that network. The example clearly shows that significant changes to subordination relations can occur without affecting the overall connectedness of the structure.

Let us now examine why the connectedness-based hierarchical order \succ_C may violate the **UE** axiom.

Consider hierarchy h , a simple two-person structure with one directed edge $j \rightarrow i$. When treated as an undirected graph, the connectedness index of h is:

$$C(h) = \frac{2}{2(2-1)} = \frac{2}{2} = 1.$$

Now consider hierarchy h' formed by appending h to a single-node hierarchy k , resulting in a three-node chain $k \rightarrow j \rightarrow i$. In h' , when treated as an undirected graph, all three nodes are connected. The total number of ordered pairs with undirected paths is 6: (i, j) , (j, i) , (i, k) , (k, i) , (j, k) and (k, j) . This gives us:

$$C(h') = \frac{6}{3(3-1)} = \frac{6}{6} = 1.$$

According to the **UE** axiom, appending h to k should make h' *strictly* more hierarchical than h . However, we observe that $C(h') = C(h) = 1$, which implies $h' \sim_C h$ (the hierarchies are equivalent under the connectedness-based ordering).

This equality contradicts the requirement of the **UE** axiom, which demands a strict increase in hierarchical measure after upward extension. Krackhardt's original connectedness

index fails to capture the increased depth of hierarchy in h' compared to h , demonstrating that \succ_C violates the **UE** axiom.

Similar examples can be used to demonstrate that \succ_C fails to satisfy **DE**, **UE***, and **DE***.

Proposition 10. *The hierarchical order \succ_C satisfies **A** and violates **RP**, **SR**, **UE**, **DE**, **UE***, and **DE***.*

4.3.2. Graph hierarchy

We now briefly discuss Krackhardt's graph hierarchy measure, noting that this measure proves uninformative in our framework of directed acyclic graphs because it remains constant across this structural domain.

Krackhardt's measure quantifies how closely a structure adheres to a pure hierarchical pattern with clear top-down authority. A high hierarchy score indicates an unambiguous chain of command where influence flows from upper nodes to subordinates without cycles or lateral connections.

Formally, Krackhardt's *graph hierarchy index* $H(h)$ is defined as:

$$H(h) = 1 - \frac{P}{\max P},$$

where

- P denotes number of *unordered pairs* $\{i, j\}$ with mutual reachability (i.e., paths exist both $i \rightarrow j$ and $j \rightarrow i$),
- $\max P$ represents the total *unordered pairs* $\{i, j\}$ with at least one directed path (either $i \rightarrow j$ or $j \rightarrow i$).

The interpretation of this index is straightforward:

- $H(h) = 1$ indicates “perfect hierarchy” with no mutual reachability between any node pairs.
- $H(h) = 0$ indicates “maximal cyclicity” where every reachable pair is mutually reachable.

For directed acyclic graphs—which by definition contain no cycles—mutual reachability is impossible, meaning $P = 0$. Consequently, in any directed acyclic graph with at least one directed path ($\max P > 0$), we always have $H(h) = 1$.

This invariance renders Krackhardt's graph hierarchy index constant across our domain of hierarchical structures, making it unsuitable for distinguishing between different acyclic hierarchical configurations.

4.3.3. Graph efficiency

Krackhardt's *graph efficiency index* $E(h)$ measures the structural efficiency of a hierarchy by quantifying the redundancy of supervisory links relative to the theoretical maximum possible redundancy. This index evaluates how densely the hierarchy deviates from a minimally connected structure.

For a hierarchy h with k weakly connected components (where the i -th component has n_i nodes) and a total of $m = \sum_{i=1}^k m_i$ edges, the graph efficiency index is defined as:³

$$E(h) = 1 - \frac{\text{actual redundancy}}{\text{maximum possible redundancy}} = 1 - \frac{\sum_{i=1}^k (m_i - (n_i - 1))}{\sum_{i=1}^k \frac{(n_i - 1)(n_i - 2)}{2}},$$

where:

- actual redundancy: $\sum_{i=1}^k (m_i - (n_i - 1))$ represents the total number of supervisory links exceeding the minimum required to keep all components weakly connected ($n_i - 1$ links per component);
- maximum possible redundancy: $\sum_{i=1}^k \frac{(n_i - 1)(n_i - 2)}{2}$ represents the theoretical maximum number of redundant links allowed in each component before it becomes a complete graph.

This measure can be given the following interpretation:

- $E(h) = 1$: A perfectly efficient hierarchy with no redundant links (i.e., an outtree).
- $E(h) < 1$: Increasing redundancy in supervisory relationships, deviating from the minimally connected structure.
- $E(h) = 0$: A maximally redundant hierarchy where every possible supervisory link exists.

Let \succ_E represent the hierarchical order induced by the graph efficiency index E .

The graph efficiency index $E(h)$ is defined entirely by structural properties of the hierarchy—specifically, the number of nodes and edges within each weakly connected component. These parameters, n_i (the size of the i -th component) and m_i (its edge count), remain invariant under any relabeling of nodes.

When renaming nodes to create hierarchy h' from h , the component sizes n_i and edge counts m_i are preserved. This ensures that structurally identical hierarchies receive identical efficiency scores, regardless of node labels.

Consequently, the efficiency-based hierarchical ordering \succ_E satisfies the **A** axiom.

We now show that \succ_E satisfies the **RP** axiom.

Let h' be a k -times replication of h . This means each component in h is replicated into k identical copies in h' . If h has c components, then h' will have $k \cdot c$ components. For each original component i in h with n_i nodes and m_i edges, we now have k identical copies in h' .

Therefore, for h' :

- The total number of edges becomes $\sum_{i=1}^{k \cdot c} m'_i = k \cdot \sum_{i=1}^c m_i$.
- The total number of nodes becomes $\sum_{i=1}^{k \cdot c} n'_i = k \cdot \sum_{i=1}^c n_i$.

³A directed acyclic graph is weakly connected if, when all directed edges are converted to undirected edges, the resulting undirected graph is connected. Equivalently, for every pair of vertices i and j , there exists an undirected path connecting them.

The actual redundancy in h' is:

$$\text{actual redundancy}(h') = \sum_{i=1}^{k \cdot c} (m'_i - (n'_i - 1)) = k \cdot \sum_{i=1}^c (m_i - (n_i - 1)).$$

And the maximum possible redundancy becomes:

$$\text{maximum possible redundancy}(h') = \sum_{i=1}^{k \cdot c} \frac{(n'_i - 1)(n'_i - 2)}{2} = k \cdot \sum_{i=1}^c \frac{(n_i - 1)(n_i - 2)}{2}.$$

Thus, the graph efficiency measure E for h' is:

$$\begin{aligned} E(h') &= 1 - \frac{\text{actual redundancy}(h')}{\text{maximum possible redundancy}(h')} \\ &= 1 - \frac{k \cdot \sum_{i=1}^c (m_i - (n_i - 1))}{k \cdot \sum_{i=1}^c \frac{(n_i - 1)(n_i - 2)}{2}} \\ &= 1 - \frac{\sum_{i=1}^c (m_i - (n_i - 1))}{\sum_{i=1}^c \frac{(n_i - 1)(n_i - 2)}{2}} \\ &= E(h). \end{aligned}$$

The scaling factor k cancels out in the calculation, demonstrating that $E(h') = E(h)$. Since $E(h') = E(h)$, it follows that $h' \sim_E h$. Therefore, \succsim_E satisfies **RP**.

Next, we show that \succsim_E violates the **SR** axiom.

Consider hierarchy h with a chain structure $j \rightarrow i \rightarrow k$. By eliminating the subordination relation between i and k , we obtain hierarchy h' with a star structure where $j \rightarrow i$ and $j \rightarrow k$.

For hierarchy h :

- Number of nodes: $n = 3$.
- Number of edges: $m = 2$.
- Minimum edges required for weak connectivity: $n - 1 = 2$.
- Actual redundancy: $m - (n - 1) = 2 - 2 = 0$.
- Maximum possible redundancy: $\frac{(n-1)(n-2)}{2} = \frac{2 \cdot 1}{2} = 1$.

Therefore:

$$E(h) = 1 - \frac{\text{actual redundancy}}{\text{maximum possible redundancy}} = 1 - \frac{0}{1} = 1.$$

For hierarchy h' :

- Number of nodes: $n = 3$.
- Number of edges: $m = 2$.
- Minimum edges required for weak connectivity: $n - 1 = 2$.

- Actual redundancy: $m - (n - 1) = 2 - 2 = 0$.
- Maximum possible redundancy: $\frac{(n-1)(n-2)}{2} = \frac{2 \cdot 1}{2} = 1$.

Therefore:

$$E(h') = 1 - \frac{\text{actual redundancy}}{\text{maximum possible redundancy}} = 1 - \frac{0}{1} = 1.$$

Consequently:

$$h \sim_E h'.$$

This demonstrates that the graph efficiency index E fails to distinguish between hierarchies that differ solely by subordination removal. Although h' is obtained from h by removing a subordination relation, we have $E(h) = E(h')$, which means $h \sim_E h'$. This contradicts the requirement of the **SR** axiom, proving that \succsim_E violates **SR**.

Next, we show that \succsim_E fails to satisfy the **UE** axiom.

Consider the following hierarchies:

- h : a chain structure with $j \rightarrow i$.
- h_o : an isolated node k .
- h' : a chain structure with $k \rightarrow j \rightarrow i$.

Note that h' can be obtained by extending hierarchy h with hierarchy h_o , formally expressed as:

$$h' \in h_o \oplus h.$$

Because neither h nor h' have any redundant links,

$$E(h) = 1 \quad \text{and} \quad E(h') = 1.$$

Hence,

$$h \sim_E h'.$$

However, according to the **UE** axiom, we should have $h' \succ_E h$ whenever $h' \in h_o \oplus h$. This contradiction demonstrates that \succsim_E violates the **UE** axiom.

One can similarly show that \succsim_E also violates **DE**, **UE***, and **DE***.

Proposition 11. *The hierarchical order \succsim_E satisfies **A** and **RP** and violates **SR**, **UE**, **DE**, **UE***, and **DE***.*

4.3.4. Least upper boundedness

The least upper boundedness index measures the extent to which a hierarchy approximates a pure tree structure by quantifying the presence of a unique common boss for each pair of nodes within the same component.

For a hierarchy $h \in \mathcal{H}$ with node set V and k weakly connected components, where each component i contains n_i nodes, the *least upper boundedness index* $\text{LUB}(h)$ is defined as:

$$\text{LUB}(h) = \frac{|\{(i, j) \in V \times V : i \neq j, i \text{ and } j \text{ are in the same component, and } \exists! \text{ LUB}(i, j)\}|}{\sum_{i=1}^k n_i(n_i - 1)},$$

where:

- The numerator counts ordered pairs of distinct nodes within the same component that have a unique least upper bound.
- The denominator represents the total number of ordered pairs of distinct nodes within all components.

To establish whether a unique least upper bound exists for a pair of nodes (i, j) (where i and j belong to the same component):

1. Identify the set of common bosses:

$$B(i, j) = \{i' \in V : \text{both } i \text{ and } j \text{ have directed paths from } i'\}.$$

By convention, each i' has a directed path to itself.

2. Determine the subset of minimal elements $M(i, j) \subseteq B(i, j)$: A node $i' \in B(i, j)$ is minimal if there is no other node $j' \in B(i, j)$ such that there exists a directed path from i' to j' .
3. A unique least upper bound exists if and only if $M(i, j)$ is a singleton.

The LUB index represents the proportion of node pairs within the same component that have a unique least upper bound in the hierarchy. This measures how closely each component resembles a pure tree structure:

- $\text{LUB}(h) = 1$: Each component has a perfect tree-like structure where every pair of nodes has exactly one least common boss.
- $\text{LUB}(h) = 0$: No pair of nodes in any component has a unique least common boss.
- $0 < \text{LUB}(h) < 1$: Components exhibit varying degrees of tree-like characteristics.

The LUB index thus captures the unity of command principle in organizational theory for each component of the hierarchy.

Let \succ_L represent the hierarchical order induced by the LUB index.

Observe that $\text{LUB}(h)$ depends solely on the structure of h , not on the specific labels of the nodes. If h' is a relabeling of h , then:

- For every pair of nodes (i, j) in h , all structural relationships (paths, reachability, boss relationships) are preserved in the corresponding pair in h' .
- The set of pairs with unique least upper bounds remains unchanged under relabeling.

- The total number of such pairs is preserved.
- The denominator (total number of pairs within components) remains constant.

Therefore, $\text{LUB}(h) = \text{LUB}(h')$, which implies $h \sim_L h'$. This demonstrates that \succsim_L satisfies the **A** axiom.

Next, we argue that the order \succsim_L satisfies **RP**.

Let h' be a k -times replication of h . By definition, h' consists of k disjoint copies of h . For the LUB index:

- Numerator: Each copy of h in h' has $x(h)$ pairs with a unique least upper bound. The total is $k \cdot x(h)$.
- Denominator: Each copy contributes $y(h)$ pairs. The total is $k \cdot y(h)$.

Thus,

$$\text{LUB}(h') = \frac{k \cdot x(h)}{k \cdot y(h)} = \text{LUB}(h).$$

Since $\text{LUB}(h') = \text{LUB}(h)$, we have $h' \sim_L h$. Hence, \succsim_L satisfies **RP**.

Next, we demonstrate that \succsim_L violates the **SR** axiom.

Consider hierarchy h , a chain structure $j \rightarrow i \rightarrow k$. By removing the subordination relation between i and k , we obtain hierarchy h' , a star structure with relations $j \rightarrow i$ and $j \rightarrow k$.

For hierarchy h (the chain structure), with $n = 3$ nodes, there are $n(n - 1) = 3 \cdot 2 = 6$ ordered pairs of distinct nodes. For each pair, a unique least upper bound (LUB) exists:

Node Pair	Unique LUB
(j, i)	$\text{LUB} = j$
(j, k)	$\text{LUB} = j$
(i, j)	$\text{LUB} = j$
(i, k)	$\text{LUB} = i$
(k, j)	$\text{LUB} = j$
(k, i)	$\text{LUB} = i$

For hierarchy h' (the star structure), there are also 6 ordered pairs of distinct nodes. For each pair, a unique LUB exists, but in this case, the unique LUB for all pairs is node j :

Node Pair	Unique LUB
(j, i)	$\text{LUB} = j$
(j, k)	$\text{LUB} = j$
(i, j)	$\text{LUB} = j$
(i, k)	$\text{LUB} = j$
(k, j)	$\text{LUB} = j$
(k, i)	$\text{LUB} = j$

Therefore:

$$\text{LUB}(h) = \frac{6}{6} = 1 = \frac{6}{6} = \text{LUB}(h').$$

This implies $h \sim_L h'$, even though h' is obtained from h by removing the subordination relation $i \rightarrow k$. According to the **SR** axiom, we should have $h \succ_L h'$, which contradicts our finding. Hence, \succsim_L violates the **SR** axiom.

Similar examples can be used to illustrate that \succsim_L violates **UE**, **DE**, **UE***, and **DE***.

Proposition 12. *The hierarchical order \succ_L satisfies **A** and **RP** and violates **SR**, **UE**, **DE**, **UE***, and **DE***.*

5. Concluding remarks

This paper establishes a robust axiomatic foundation for hierarchy measurement by introducing novel axioms centered on structural depth. These new axioms complement existing principles—the Anonymity axiom, Replication Principle, and Subordination Removal axiom from Carbonell-Nicolau (2025a,b)—to create a more comprehensive theoretical framework.

Our expanded axiom set enables the development of hierarchical measures with significantly enhanced discriminatory power for ranking pairs of hierarchies. We formulate canonical hierarchical pre-orders uniquely characterized by these underlying principles and identify the average depth hierarchical index as a computationally efficient, full completion of these pre-orders. This index measures hierarchy by averaging each node’s number of (direct or indirect) supervisors, providing both axiom consistency and practical applicability.

The theoretical framework developed here strengthens the conceptual basis for measuring hierarchies and identifies several avenues for future research.

A natural avenue for investigation involves the empirical study of hierarchical transformations during organizational processes such as mergers and acquisitions. The extension axioms developed here, which address upward and downward structural integration, capture stylized outcomes commonly observed in such contexts, including the addition of management layers and integration of distinct organizational charts. This connection suggests a testable hypothesis: real-world mergers and acquisitions tend to produce deeper, more hierarchical structures. This claim can be assessed using the measures introduced in this study, thereby linking axiomatic theory with organizational data. The average depth index provides a particularly promising tool for such inquiry due to its computational simplicity and complete alignment with our expanded axiom set.

Two important caveats warrant attention in pursuing this empirical agenda. First, the scarcity of publicly available data on internal organizational hierarchies may pose a significant barrier to large-scale validation. Second, while the average depth index provides a complete ordering of hierarchies, it may deliver definitive comparisons even in cases where our axioms yield no such conclusion. This mismatch can lead to potentially overstated or non-compelling rankings in edge cases not covered by the axiomatic structure.

To address these limitations, one may employ the composite axiomatic collections constructed in this paper—specifically, the transitive closures of the union of canonical hierarchical pre-orders. These offer a more principled approach by ensuring that every comparison is grounded in specific axioms. However, such methods are computationally intensive, particularly for large organizational networks.

This computational challenge suggests another direction for future research: developing scalable algorithms that can efficiently compute composite pre-orders. Such advances would enable the application of theoretically sound methods to real-world data without compromising either computational feasibility or axiomatic consistency.

A. Proofs

A.1. Proof of Lemma 1

Lemma 1. *The binary relation \sim^* defined on \mathcal{H} is symmetric, reflexive, and transitive.*

Proof. Symmetry is implied by the symmetry of the relabeling relation itself. If $h' \sim^* h$, then there exist replicas of h' and h that are relabelings of each other, which implies $h \sim^* h'$.

Reflexivity follows from the fact that each hierarchy is trivially a relabeling of itself. Therefore, $h \sim^* h$ for any hierarchy h .

To prove transitivity, suppose that

$$h'' \sim^* h' \sim^* h.$$

By definition of \sim^* , there exist replicas h''_r and h'_r of h'' and h' , respectively, such that h''_r is a relabeling of h'_r . Similarly, there exist replicas h'_{rr} and h_{rr} of h' and h , respectively, such that h'_{rr} is a relabeling of h_{rr} .

Suppose that h'_r and h'_{rr} are l and m -dimensional replicas of h' , respectively. Then the m -replica of h'_r , h'_{rrr} , equals the l -replica of h'_{rr} . Therefore, we have:

- The m -replica of h''_r is a relabeling of h'_{rrr} .
- h'_{rrr} is a relabeling of the l -replica of h_{rr} .

Consequently, by transitivity of the relabeling relation, the m -replica of h''_r is a relabeling of the l -replica of h_{rr} . Since the m -replica of h''_r is a replica of h'' and the l -replica of h_{rr} is a replica of h , it follows that $h'' \sim^* h$. ■

A.2. Proof of Proposition 1

Proposition 1. *The binary relation \succ_{UE} is a hierarchical pre-order on \mathcal{H} satisfying **A**, **RP**, and **UE**.*

Proof. Reflexivity follows from the fact that each hierarchy is a relabeling of itself.

To prove transitivity, suppose that

$$h'' \succ_{UE} h' \succ_{UE} h.$$

We consider four cases.

Case 1. We have

$$H'' \in h'_o \oplus H' \quad \text{and} \quad \tilde{H} \in h_o \oplus H, \tag{3}$$

for $H'' \sim^* h''$, $H' \sim^* h'$, and $H \sim^* h$. By Lemma 1, $H' \sim^* \tilde{H}$, so there exist replicas H'_r and \tilde{H}_r of H' and \tilde{H} , respectively, such that H'_r is a relabeling of \tilde{H}_r . Therefore, there exist h_a and h_b such that h_a is a relabeling of h_b and

$$H'_r = (h_a, \dots, h_a) \quad \text{and} \quad \tilde{H}_r = (h_b, \dots, h_b).$$

By (3), we have

$$\begin{aligned} (H'', \dots, H'') &\in (h'_o \oplus H') \times \dots \times (h'_o \oplus H') \\ &\subseteq (h'_o, \dots, h'_o) \oplus (H', \dots, H') = (h'_o, \dots, h'_o) \oplus H'_r = (h'_o, \dots, h'_o) \oplus (h_a, \dots, h_a) \end{aligned}$$

and

$$(h_b, \dots, h_b) = \tilde{H}_r = (\tilde{H}, \dots, \tilde{H}) \in (h_o \oplus H) \times \dots \times (h_o \oplus H) \subseteq (h_o, \dots, h_o) \oplus (H, \dots, H).$$

Therefore, since (h_a, \dots, h_a) is a relabeling of (h_b, \dots, h_b) and h_a is a relabeling of h_b , there exists a relabeling (H''', \dots, H''') of (H'', \dots, H'') such that

$$(H''', \dots, H''') \in (h'_o, \dots, h'_o) \oplus (h_o, \dots, h_o) \oplus (H, \dots, H).$$

Since

$$(H''', \dots, H''') \sim^* (H'', \dots, H'') \sim^* h'' \quad \text{and} \quad (H, \dots, H) \sim^* h,$$

it follows that $h'' \succ_{UE} h$.

Case 2. $h' \sim^* h$ and

$$H'' \in h'_o \oplus H' \tag{4}$$

for $H'' \sim^* h''$ and $H' \sim^* h'$. By **Lemma 1**, $H' \sim^* h$, so there exist replicas H'_r and h_r of H' and h , respectively, such that H'_r is a relabeling of h_r . Hence,

$$H'_r = (H', \dots, H') = (h_a, \dots, h_a) = (h_b, \dots, h_b) = (h, \dots, h) = h_r$$

for some h_a and h_b such that h_a is a relabeling of h_b , and (4) implies

$$\begin{aligned} (H'', \dots, H'') &\in (h'_o \oplus H') \times \dots \times (h'_o \oplus H') \\ &\subseteq (h'_o, \dots, h'_o) \oplus (H', \dots, H') = (h'_o, \dots, h'_o) \oplus (h_a, \dots, h_a). \end{aligned}$$

Since (h_a, \dots, h_a) is a relabeling of (h_b, \dots, h_b) and h_a is a relabeling of h_b , it follows that there exists a relabeling (H''', \dots, H''') of (H'', \dots, H'') such that

$$(H''', \dots, H''') \in (h'_o, \dots, h'_o) \oplus (h_b, \dots, h_b) = (h'_o, \dots, h'_o) \oplus (h, \dots, h).$$

Consequently, since

$$(H''', \dots, H''') \sim^* (H'', \dots, H'') \sim^* h'' \quad \text{and} \quad (h, \dots, h) \sim^* h,$$

we have $h'' \succ_{UE} h$.

Case 3. $h'' \sim^* h'$ and

$$\tilde{H} \in h_o \oplus H$$

for $\tilde{H} \sim^* h'$, and $H \sim^* h$. This case can be handled as Case 2.

Case 4. $h'' \sim^* h' \sim^* h$. By **Lemma 1**, $h'' \sim^* h$, implying $h'' \succ_{UE} h$.

To see that \succ_{UE} satisfies **A**, pick any two hierarchies h and h' . Suppose that h is a relabeling of h' . Then $h \sim^* h'$ and $h' \sim^* h$, implying $h \sim_{UE} h'$.

The relation \succ_{UE} also satisfies the **RP** axiom, since, $(h, \dots, h) \sim^* h$ and $h \sim^* (h, \dots, h)$ for any hierarchy h , implying $h' \sim_{UE} h$ whenever h' is a replication of h .

The relation \succ_{UE} also adheres to the **UE** axiom. Indeed, suppose that $h' \in h_o \oplus h$ for some $h_o \in \mathcal{H}$. Then, the definition of \succ_{UE} implies $h' \succ_{UE} h$.

To ensure that $h' \succ_{UE} h$, we must confirm that $h \succ_{UE} h'$ does not hold. Suppose, for contradiction, that $h \succ_{UE} h'$. This implies either $h \sim^* h'$ or

$$\exists h'_0, H \sim^* h, H' \sim^* h' : H \in h'_0 \oplus H'.$$

Case A. $h \sim^* h'$. If $h \sim^* h'$, there exist equally-sized replicas h_r of h and h'_r of h' such that h_r is a relabeling of h'_r . However, since $h' \in h_0 \oplus h$, the depth of h' exceeds that of h . Since replicas preserve depth, it follows that h_r and h'_r have different depths. Relabeling preserves structure, including depth, so no relabeling can make h_r and h'_r equivalent. Thus, $h \not\sim^* h'$.

Case B. $H \in h'_0 \oplus H'$ with $H \sim^* h$, $H' \sim^* h'$. This case is impossible due to depth differences. Since $h' \in h_0 \oplus h$, h' 's depth exceeds h 's depth. Replicas and relabelings preserve depth so $H \sim^* h$ and $H' \sim^* h'$ imply H' 's depth exceeds H 's depth. But if $H \in h'_0 \oplus H'$, H 's depth would need to exceed that of H' , a contradiction.

Therefore, neither case holds, confirming $h \not\succ_{UE} h'$. ■

A.3. Proof of Theorem 2

Theorem 2. A hierarchical pre-order on \mathcal{H} satisfies **A**, **RP**, **UE**, and **SR** if and only if it is R^* -consistent.

Proof. (Sufficiency.) Assume that \succ is R^* -consistent. By Proposition 2, R^* satisfies the axioms **A**, **RP**, **UE**, and **SR**. Since \succ is R^* -consistent, it follows that \succ must also satisfy these axioms.

(Necessity.) Suppose that \succ is a hierarchical pre-order on \mathcal{H} satisfying axioms **A**, **RP**, **UE**, and **SR**. We must show that \succ is R^* -consistent, i.e., that for every pair of hierarchies $h, h' \in \mathcal{H}$:

$$(a) \quad hP^*h' \Rightarrow h > h'.$$

$$(b) \quad hI^*h' \Rightarrow h \sim h'.$$

Suppose first that hI^*h' . Then there are finitely many hierarchies in \mathcal{H} ,

$$h_0, \dots, h_k, h'_0, \dots, h'_l,$$

such that

$$(h = h_0)Rh_1R \cdots R(h_k = h' = h'_0)Rh'_1R \cdots R(h'_l = h).$$

Since this represents a cycle, Lemma 2 implies

$$(h = h_0)Ih_1I \cdots I(h_k = h').$$

We claim that for each $\kappa \in \{0, \dots, k-1\}$, there exist replicas H_κ and $H_{\kappa+1}$ of h_κ and $h_{\kappa+1}$ respectively such that H_κ is a relabeling of $H_{\kappa+1}$. To see this, fix $\kappa \in \{0, \dots, k-1\}$ and note that $h_\kappa Ih_{\kappa+1}$ implies

$$[h_\kappa \succ_H h_{\kappa+1} \text{ or } h_\kappa \succ_{UE} h_{\kappa+1}] \quad \text{and} \quad [h_{\kappa+1} \succ_H h_\kappa \text{ or } h_{\kappa+1} \succ_{UE} h_\kappa]. \quad (5)$$

There are four exhaustive cases to consider.

Case 1. $h_\kappa \sim_H h_{\kappa+1}$. Then there exist equally-sized replicas H_κ and $H_{\kappa+1}$ of h_κ and $h_{\kappa+1}$, respectively, such that $H_\kappa \sim_H H_{\kappa+1}$. By Lemma 4 in Carbonell-Nicolau (2025b), it follows that H_κ is a relabeling of $H_{\kappa+1}$.

Case 2. $h_\kappa \sim_{UE} h_{\kappa+1}$. Using the argument from the proof of Theorem 1, we can conclude that $h_\kappa \sim^* h_{\kappa+1}$, which gives the desired conclusion.

Case 3. $h_\kappa >_H h_{\kappa+1}$. In this case, (5) implies $h_{\kappa+1} \succ_{UE} h_\kappa$. The case where $h_{\kappa+1} >_{UE} h_\kappa$ is impossible, since $h_\kappa >_H h_{\kappa+1}$ implies that the average depth in h_κ exceeds that of $h_{\kappa+1}$, contradicting $h_{\kappa+1} >_{UE} h_\kappa$. Indeed, the order \succ_{I_s} on \mathcal{H} defined by

$$h \succ_{I_s} h' \Leftrightarrow I_s(h) \geq I_s(h'),$$

where $I_s(h'')$ represents the average number of (direct or indirect) supervisors in h'' , is \succ_H -consistent (Carbonell-Nicolau, 2025b). Consequently, $h_\kappa >_H h_{\kappa+1}$ implies $I_s(h_\kappa) > I_s(h_{\kappa+1})$.

Thus, we have $h_\kappa >_H h_{\kappa+1}$ and $h_{\kappa+1} \sim_{UE} h_\kappa$, so we can proceed as in Case 2.

Case 4. $h_\kappa >_{UE} h_{\kappa+1}$. In this case, (5) implies $h_{\kappa+1} \succ_H h_\kappa$. The case where $h_{\kappa+1} >_H h_\kappa$ is impossible, since $h_\kappa >_{UE} h_{\kappa+1}$ implies that the average depth in h_κ exceeds that of $h_{\kappa+1}$, contradicting $h_{\kappa+1} >_{UE} h_\kappa$.

Thus, we have $h_\kappa >_{UE} h_{\kappa+1}$ and $h_{\kappa+1} \sim_H h_\kappa$, so we can proceed as in Case 1.

We have seen that for each $\kappa \in \{0, \dots, k-1\}$, there exist replicas H_κ and $H_{\kappa+1}$ of h_κ and $h_{\kappa+1}$ respectively such that H_κ is a relabeling of $H_{\kappa+1}$. Consequently, because \succ satisfies A and RP, we have

$$h \sim H_0 \sim \dots \sim H_k \sim h'. \quad (6)$$

By transitivity of \succ , \sim is transitive, so (6) implies $h \sim h'$. This establishes (b).

It remains to prove (a). Suppose that hP^*h' . Then there are finitely many hierarchies in \mathcal{H} ,

$$h''_0, \dots, h''_l,$$

such that

$$(h = h''_0)Rh''_1R \dots R(h''_l = h'), \quad (7)$$

with at least one strict relation.

Suppose that $h''_\kappa Ph''_{\kappa+1}$, where $\kappa \in \{0, \dots, l-1\}$. Then

$$h''_\kappa >_H h''_{\kappa+1} \quad \text{or} \quad h''_\kappa >_{UE} h''_{\kappa+1}.$$

Suppose first that $h''_\kappa >_H h''_{\kappa+1}$. Then by Theorem 1 in Carbonell-Nicolau (2025b), $h''_{\kappa+1}$ can be obtained from some relabeling of h''_κ by successive removals of subordination relations. Consequently, since \succ satisfies SR, there are finitely many hierarchies in \mathcal{H} ,

$$\tilde{h}_0, \dots, \tilde{h}_\ell,$$

such that

$$\tilde{h}_0 > \tilde{h}_1 > \dots > (\tilde{h}_\ell = h''_{\kappa+1}),$$

where \tilde{h}_0 is a relabeling of h''_κ . Hence, because \succ satisfies A,

$$h''_\kappa \sim \tilde{h}_0 > \tilde{h}_1 > \dots > (\tilde{h}_\ell = h''_{\kappa+1}).$$

Since \succ is reflexive and transitive, it follows that $h''_\kappa > h''_{\kappa+1}$ (Sen, 2017, Lemma 1*a, p. 56).

Next, suppose that $h''_{\kappa} \succ_{UE} h''_{\kappa+1}$. Then

$$\exists h_o, H''_{\kappa} \sim^* h''_{\kappa}, H''_{\kappa+1} \sim^* h''_{\kappa+1} : H''_{\kappa} \in h_o \oplus H''_{\kappa+1}.$$

Consequently, because \succ satisfies the **UE** axiom, we see that

$$H''_{\kappa} \succ H''_{\kappa+1},$$

and, since \succ satisfies **A** and **RP**, it follows that $h''_{\kappa} \succ h''_{\kappa+1}$.

We conclude that $h''_{\kappa} P h''_{\kappa+1}$, where $\kappa \in \{0, \dots, l-1\}$, implies $h''_{\kappa} \succ h''_{\kappa+1}$.

Furthermore, if $h''_{\kappa} I h''_{\kappa+1}$, where $\kappa \in \{0, \dots, l-1\}$, then, applying the argument used in the proof of **(b)**, we see that there exist replicas \tilde{H}_{κ} and $\tilde{H}_{\kappa+1}$ of h''_{κ} and $h''_{\kappa+1}$, respectively, such that \tilde{H}_{κ} is a relabeling of $\tilde{H}_{\kappa+1}$. Therefore, since \succ satisfies **A** and **RP**, we have

$$h''_{\kappa} \sim \tilde{H}_{\kappa} \sim \tilde{H}_{\kappa+1} \sim h''_{\kappa+1},$$

implying $h''_{\kappa} \sim h''_{\kappa+1}$ by transitivity of \sim .

Thus, $h''_{\kappa} I h''_{\kappa+1}$, where $\kappa \in \{0, \dots, l-1\}$, implies $h''_{\kappa} \sim h''_{\kappa+1}$. Moreover, we have also established that $h''_{\kappa} P h''_{\kappa+1}$, where $\kappa \in \{0, \dots, l-1\}$, implies $h''_{\kappa} \succ h''_{\kappa+1}$. Consequently, (7), which holds with at least one strict dominance, implies

$$(h = h''_0) \succ h''_1 \succ \dots \succ (h'_l = h'),$$

with at least one strict dominance. Since \succ is reflexive and transitive, Lemma 1*a in Sen (2017, p. 56) implies $h \succ h'$. This establishes **(a)**. ■

A.4. Proof of Proposition 3

Proposition 3. The binary relation \succ_{DE} is a hierarchical pre-order on \mathcal{H} satisfying **A**, **RP**, and **DE**.

Proof. Reflexivity follows from the fact that each hierarchy is a relabeling of itself.

To prove transitivity, suppose that

$$h'' \succ_{DE} h' \succ_{DE} h.$$

We consider four cases.

Case 1. We have

$$H'' \in H' \boxplus h'_o \quad \text{and} \quad \tilde{H} \in H \boxplus h_o, \quad (8)$$

for $H'' \sim^* h''$, $H' \sim^* h'$, and $\tilde{H} \sim^* h$. By Lemma 1, $H' \sim^* \tilde{H}$, so there exist replicas H'_r and \tilde{H}_r of H' and \tilde{H} , respectively, such that H'_r is a relabeling of \tilde{H}_r . Therefore, there exist h_a and h_b such that h_a is a relabeling of h_b and

$$H'_r = (h_a, \dots, h_a) \quad \text{and} \quad \tilde{H}_r = (h_b, \dots, h_b).$$

By (8), we have

$$\begin{aligned} (H'', \dots, H'') &\in (H' \boxplus h'_o) \times \dots \times (H' \boxplus h'_o) \\ &\subseteq (H', \dots, H') \boxplus (h'_o, \dots, h'_o) = H'_r \boxplus (h'_o, \dots, h'_o) = (h_a, \dots, h_a) \boxplus (h'_o, \dots, h'_o) \end{aligned}$$

and

$$(h_b, \dots, h_b) = \tilde{H}_r = (\tilde{H}, \dots, \tilde{H}) \in (H \boxplus h_o) \times \dots \times (H \boxplus h_o) \subseteq (H, \dots, H) \boxplus (h_o, \dots, h_o).$$

Therefore, since (h_a, \dots, h_a) is a relabeling of (h_b, \dots, h_b) and h_a is a relabeling of h_b , there exists a relabeling (H''', \dots, H''') of (H'', \dots, H'') such that

$$(H''', \dots, H''') \in (H, \dots, H) \boxplus (h_o, \dots, h_o) \boxplus (h'_o, \dots, h'_o).$$

Since

$$(H''', \dots, H''') \sim^* (H'', \dots, H'') \sim^* h'' \quad \text{and} \quad (H, \dots, H) \sim^* h,$$

it follows that $h'' \succ_{DE} h$.

Case 2. $h' \sim^* h$ and

$$H'' \in H' \boxplus h'_o \tag{9}$$

for $H'' \sim^* h''$ and $H' \sim^* h'$. By **Lemma 1**, $H' \sim^* h$, so there exist replicas H'_r and h_r of H' and h , respectively, such that H'_r is a relabeling of h_r . Hence,

$$H'_r = (H', \dots, H') = (h_a, \dots, h_a) = (h_b, \dots, h_b) = (h, \dots, h) = h_r$$

for some h_a and h_b such that h_a is a relabeling of h_b , and (9) implies

$$\begin{aligned} (H'', \dots, H'') &\in (H' \boxplus h'_o) \times \dots \times (H' \boxplus h'_o) \\ &\subseteq (H', \dots, H') \boxplus (h'_o, \dots, h'_o) = (h_a, \dots, h_a) \boxplus (h'_o, \dots, h'_o). \end{aligned}$$

Since (h_a, \dots, h_a) is a relabeling of (h_b, \dots, h_b) and h_a is a relabeling of h_b , it follows that there exists a relabeling (H''', \dots, H''') of (H'', \dots, H'') such that

$$(H''', \dots, H''') \in (h_b, \dots, h_b) \boxplus (h'_o, \dots, h'_o) = (h, \dots, h) \boxplus (h'_o, \dots, h'_o).$$

Consequently, since

$$(H''', \dots, H''') \sim^* (H'', \dots, H'') \sim^* h'' \quad \text{and} \quad (h, \dots, h) \sim^* h,$$

we have $h'' \succ_{DE} h$.

Case 3. $h'' \sim^* h'$ and

$$\tilde{H} \in H \boxplus h_o$$

for $\tilde{H} \sim^* h'$, and $H \sim^* h$. This case can be handled as Case 2.

Case 4. $h'' \sim^* h' \sim^* h$. By **Lemma 1**, $h'' \sim^* h$, implying $h'' \succ_{DE} h$.

To see that \succ_{DE} satisfies **A**, pick any two hierarchies h and h' . Suppose that h is a relabeling of h' . Then $h \sim^* h'$ and $h' \sim^* h$, implying $h \sim_{DE} h'$.

The relation \succ_{DE} also satisfies the **RP** axiom, since, $(h, \dots, h) \sim^* h$ and $h \sim^* (h, \dots, h)$ for any hierarchy h , implying $h' \sim_{DE} h$ whenever h' is a replication of h .

The relation \succ_{DE} also adheres to the **DE** axiom. Indeed, suppose that $h' \in h \boxplus h_o$ for some $h_o \in \mathcal{H}$. Then, the definition of \succ_{DE} implies $h' \succ_{DE} h$.

To ensure that $h' \succ_{DE} h$, we must confirm that $h \succ_{DE} h'$ does not hold. Suppose, for contradiction, that $h \succ_{DE} h'$. This implies either $h \sim^* h'$ or

$$\exists h'_0, H \sim^* h, H' \sim^* h' : H \in H' \boxplus h'_0.$$

Case A. $h \sim^* h'$. If $h \sim^* h'$, there exist equally-sized replicas h_r of h and h'_r of h' such that h_r is a relabeling of h'_r . However, since $h' \in h \boxplus h_0$, the depth of h' exceeds that of h . Since replicas preserve depth, it follows that h_r and h'_r have different depths. Relabeling preserves structure, including depth, so no relabeling can make h_r and h'_r equivalent. Thus, $h \not\sim^* h'$.

Case B. $H \in H' \boxplus h'_0$ with $H \sim^* h, H' \sim^* h'$. This case is impossible due to depth differences. Since $h' \in h \boxplus h_0$, h' 's depth exceeds h 's depth. Replicas and relabelings preserve depth so $H \sim^* h$ and $H' \sim^* h'$ imply H 's depth exceeds H' 's depth. But if $H \in h'_0 \boxplus H'$, H 's depth would need to exceed that of H' , a contradiction.

Therefore, neither case holds, confirming $h \not\succ_{DE} h'$. ■

A.5. Proof of Theorem 3

Theorem 3. A hierarchical pre-order on \mathcal{H} satisfies **A**, **RP**, and **DE** if and only if it is \succ_{DE} -consistent.

Proof. (Sufficiency.) Assume that \succ is a \succ_{DE} -consistent hierarchical pre-order on \mathcal{H} . According to **Proposition 3**, \succ_{DE} satisfies the axioms **A**, **RP**, and **DE**. Since \succ is \succ_{DE} -consistent, it follows that \succ must also satisfy these three axioms.

(Necessity.) Suppose that \succ is a hierarchical pre-order on \mathcal{H} satisfying axioms **A**, **RP**, and **DE**. We need to show that \succ is \succ_{DE} -consistent, which requires proving that for every pair of hierarchies $h, h' \in \mathcal{H}$:

$$(a) \ h \succ_{DE} h' \Rightarrow h \succ h'.$$

$$(b) \ h \sim_{DE} h' \Rightarrow h \sim h'.$$

Suppose first that $h \sim_{DE} h'$. Then there are two possibilities: either $h \sim^* h'$ or

$$\exists h_0, H' \sim^* h', H \sim^* h : H' \in H \boxplus h_0 \quad \text{and} \quad \exists h'_0, H''' \sim^* h', H'' \sim^* h : H'' \in H''' \boxplus h'_0.$$

However, the second case is impossible. Indeed, $H' \in H \boxplus h_0$ implies that the depth of H' exceeds that of H , while $H'' \in H''' \boxplus h'_0$ implies that the depth of H'' exceeds that of H''' . Since

$$H' \sim^* h' \sim^* H''' \quad \text{and} \quad H \sim^* h \sim^* H'',$$

we have

$$H' \sim^* H''' \quad \text{and} \quad H \sim^* H''$$

(**Lemma 1**). Consequently, because replication and relabeling preserves depth, it follows that the depth of H' and H''' exceeds the depth of H and H'' , a contradiction.

Thus, $h \sim_{DE} h'$ implies $h \sim^* h'$, so there exist replicas h_r and h'_r of h and h' , respectively, such that h_r is a relabeling of h'_r .

Given that \succ satisfies axioms **A** and **RP**, we have:

$$h \sim h_r \sim h'_r \sim h'. \tag{10}$$

Since \succsim is reflexive and transitive, \sim is transitive (Sen, 2017, Lemma 1*a, p. 56). Therefore, (10) implies $h \sim h'$, which establishes (b).

Now suppose that $h \succ_{DE} h'$. Then

$$\exists h_o, H' \sim^* h', H \sim^* h : H \in H' \boxplus h_o.$$

Consequently, because \succsim satisfies the DE axiom, we see that

$$H \succ H'. \quad (11)$$

Now, since $h \sim^* H$, there are replicas h_r and H_r of h and H , respectively, such that h_r is a relabeling of H_r . Similarly, there are replicas H'_r and h'_r of H' and h' , respectively, such that H'_r is a relabeling of h'_r .

Hence, since \succsim satisfies the A and RP axioms, it follows that

$$h \sim h_r \sim H_r \sim H \quad \text{and} \quad h' \sim h'_r \sim H'_r \sim H'.$$

Combining these relations with (11) yields

$$h \sim h_r \sim H_r \sim H \succ H' \sim H'_r \sim h'_r \sim h'.$$

Since \succsim is reflexive and transitive, this chain of relations implies $h \succ h'$ (Sen, 2017, Lemma 1*a, p. 56). This establishes (a) and completes the proof. ■

A.6. Proof of Proposition 4

Proposition 4. *The binary relation \succsim_{UE^*} is a hierarchical pre-order on \mathcal{H} satisfying A, RP, and UE^* .*

Proof. Reflexivity and transitivity are immediate.

If h' is a relabeling or a replica of h , then $h' \sim^* h$, implying $h' \sim_{UE^*} h$. Thus, \succsim_{UE^*} satisfies A and RP.

Suppose that $h' \in h_o \uplus h$ and $h_o \succ_{UE^*} h$ for some $h_o \in \mathcal{H}$. Then there exists a finite sequence

$$h_o = h_0, \dots, h_m = h$$

such that

$$h_0 R_{k_0} h_1 R_{k_1} \dots R_{k_{m-1}} h_m,$$

implying

$$(h_o = h_0) R_{k^*} h_1 R_{k^*} \dots R_{k^*} (h_m = h), \quad (12)$$

where $k^* = \max\{k_0, \dots, k_{m-1}\}$.

Since $h' \in h_o \uplus h$ and (12) holds, we have $h' R_{k^*+1} h$, implying $h' \succ_{UE^*} h$.

Now, since $h' \in h_o \uplus h$, h' has a higher maximum depth than h , so $h \not\succ_{UE^*} h'$.

Consequently, $h' \succ_{UE^*} h$. We conclude that \succsim_{UE^*} satisfies UE^* . ■

A.7. Proof of Theorem 4

The proof of Theorem 4 relies on the following intermediate result.

Let I_k and P_k denote the symmetric and asymmetric parts of R_k , respectively.

Lemma 3. Suppose that

$$h_0 R_{k_1} \cdots R_{k_{l+1}} h_l R_k h_0.$$

Then

$$h_0 I_{k_1} \cdots I_{k_{l+1}} h_l I_k h_0.$$

Proof. Suppose that

$$h_0 R_{k_1} \cdots R_{k_{l+1}} h_l R_k h_0.$$

Observe that

$$h_0 R_{k_1} \cdots R_{k_{l+1}} h_l$$

implies $I_s(h_0) \geq I_s(h_l)$, where $I_s(h'')$ represents the average number of (direct or indirect) supervisors in h'' .

Similarly, $h_l R_k h_0$ implies $I_s(h_l) \geq I_s(h_0)$.

Consequently, h_0 and h_l have the same average depth, $I_s(h_0) = I_s(h_l)$, implying that

$$h_0 I_{k_1} \cdots I_{k_{l+1}} h_l I_k h_0.$$

Indeed, if at least one of the relations were strict, one of the following would hold, since strict dominance in R_k forces a strict increase in I_s :

- $I_s(h_l) < I_s(h_0)$ and $I_s(h_0) \leq I_s(h_l)$. This is a contradiction.
- $I_s(h_l) \leq I_s(h_0)$ and $I_s(h_0) < I_s(h_l)$. This is a also contradiction. ■

We are now ready to prove **Theorem 4**

Theorem 4. A hierarchical pre-order on \mathcal{H} satisfies **A**, **RP**, and **UE*** if and only if it is \succsim_{UE^*} -consistent.

Proof. (Sufficiency.) Assume that \succsim is \succsim_{UE^*} -consistent. By **Proposition 4**, \succsim_{UE^*} satisfies the axioms **A**, **RP**, and **UE***. Since \succsim is \succsim_{UE^*} -consistent, it follows that \succsim must also satisfy these axioms.

(Necessity.) Suppose that \succsim is a hierarchical pre-order on \mathcal{H} satisfying axioms **A**, **RP**, and **UE***. We must show that \succsim is \succsim_{UE^*} -consistent, i.e., that for every pair of hierarchies $h, h' \in \mathcal{H}$:

- (a) $h \succ_{UE^*} h' \Rightarrow h \succ h'$.
- (b) $h \sim_{UE^*} h' \Rightarrow h \sim h'$.

Suppose first that $h \sim_{UE^*} h'$. Then there are finitely many hierarchies in \mathcal{H} ,

$$h_0, \dots, h_k, h'_0, \dots, h'_L,$$

such that

$$(h = h_0) R_{k_1} h_1 R_{k_2} \cdots R_{k_k} (h_k = h' = h'_0) R_{k'_1} h'_1 R_{k'_2} \cdots R_{k'_L} (h'_L = h).$$

Since this represents a cycle, **Lemma 3** implies

$$(h = h_0) I_{k_1} h_1 I_{k_2} \cdots I_{k_k} (h_k = h').$$

Hence,

$$(h = h_0)I_{k^*}h_1I_{k^*}\cdots I_{k^*}(h_k = h'),$$

where $k^* = \max\{k_1, \dots, k_k\}$. Note that this relational chain implies

$$(h = h_0) \sim^* h_1 \sim^* \cdots \sim^* (h_k = h').$$

Consequently, because \succsim satisfies **A** and **RP**, we have

$$h \sim h_0 \sim \cdots \sim h_k \sim h'. \quad (13)$$

By transitivity of \succsim , \sim is transitive, so (13) implies $h \sim h'$. This establishes (b).

It remains to prove (a). First, we show that if (h, h') belongs to the transitive closure of R_k then $h \succsim h'$. We proceed by induction on k .

If (h, h') belongs to the transitive closure of R_0 then hR_0h' (**Lemma 1**), implying $h \sim h'$ because \succsim satisfies **A** and **RP**.

Now suppose that, for each $k \in \{0, \dots, \kappa\}$, if (h, h') belongs to the transitive closure of R_k then $h \succsim h'$. We will show that if (h, h') belongs to the transitive closure of $R_{\kappa+1}$ then $h \succsim h'$.

If (h, h') belongs to the transitive closure of $R_{\kappa+1}$, then there exist finitely many hierarchies

$$\bar{h}_0, \dots, \bar{h}_\ell$$

in \mathcal{H} such that

$$\bar{h}_0 = hR_{\kappa+1} \cdots R_{\kappa+1}\bar{h}_\ell = h'. \quad (14)$$

Observe that $\bar{h}_l R_{\kappa+1} \bar{h}_{l+1}$ implies that either $\bar{h}_l \sim^* \bar{h}_{l+1}$ or

$$\exists \tilde{h}_0 R_\kappa \cdots R_\kappa \tilde{h}_m = \bar{h}_{l+1}, H_l \sim^* \bar{h}_l, H_{l+1} \sim^* \bar{h}_{l+1} : H_l \in \tilde{h}_0 \uplus H_{l+1}. \quad (15)$$

In the first case, we see that $\bar{h}_l \sim \bar{h}_{l+1}$ because \succsim satisfies **A** and **RP**. In the second case, (15) implies that $(\tilde{h}_0, \bar{h}_{l+1})$ belongs to the transitive closure of R_κ , so the induction hypothesis gives $\tilde{h}_0 \succsim \bar{h}_{l+1}$. Because

$$H_l \in \tilde{h}_0 \uplus H_{l+1} \quad \text{and} \quad \tilde{h}_0 \succsim \bar{h}_{l+1} \sim^* H_{l+1},$$

we have

$$H_l \in \tilde{h}_0 \uplus H_{l+1} \quad \text{and} \quad \tilde{h}_0 \succsim \bar{h}_{l+1} \sim H_{l+1}$$

(since \succsim satisfies **A** and **RP**). Therefore, by transitivity of \succsim , we have

$$H_l \in \tilde{h}_0 \uplus H_{l+1} \quad \text{and} \quad \tilde{h}_0 \succsim H_{l+1},$$

implying $H_l > H_{l+1}$, since \succsim satisfies the **UE*** axiom. Thus,

$$\bar{h}_l \sim^* H_l > H_{l+1} \sim^* \bar{h}_{l+1},$$

whence

$$\bar{h}_l \sim H_l > H_{l+1} \sim \bar{h}_{l+1} \quad (16)$$

(since \succsim satisfies **A** and **RP**). Because \succsim is reflexive and transitive, (16) yields $\bar{h}_l > \bar{h}_{l+1}$ (**Sen, 2017, Lemma 1*a**, p. 56).

We have seen that $\bar{h}_l R_{\kappa+1} \bar{h}_{l+1}$ implies

$$\bar{h}_l \sim \bar{h}_{l+1} \quad \text{or} \quad \bar{h}_l > \bar{h}_{l+1}.$$

Consequently, (14) implies

$$\bar{h}_0 = h \succ \cdots \succ \bar{h}_\ell = h'.$$

By transitivity of \succ , we obtain $h \succ h'$, as we sought.

We are now ready to prove (a). Suppose that $h \succ_{UE^*} h'$. Then there are finitely many hierarchies in \mathcal{H} ,

$$h''_0, \dots, h''_l,$$

such that

$$(h = h''_0) R_{k_1} h''_1 R_{k_2} \cdots R_{k_l} (h''_l = h'), \quad (17)$$

with at least one strict relation.

Suppose that $h''_\kappa P_{\kappa+1} h''_{\kappa+1}$, where $\kappa \in \{0, \dots, l-1\}$. Then

$$\exists \hat{h}_0 R_{k_{\kappa+1}-1} \cdots R_{k_{\kappa+1}-1} \hat{h}_M = h''_{\kappa+1}, H''_\kappa \sim^* h''_\kappa, H''_{\kappa+1} \sim^* h''_{\kappa+1} : H''_\kappa \in \hat{h}_0 \uplus H''_{\kappa+1}.$$

Since $(\hat{h}_0, h''_{\kappa+1})$ belongs to the transitive closure of $R_{k_{\kappa+1}-1}$, our previous result yields $\hat{h}_0 \succ h''_{\kappa+1}$. Consequently, because \succ satisfies the **UE*** axiom, we see that

$$H''_\kappa > H''_{\kappa+1},$$

and, since \succ satisfies **A** and **RP**, it follows that $h''_\kappa > h''_{\kappa+1}$.

We conclude that $h''_\kappa P_{\kappa+1} h''_{\kappa+1}$, where $\kappa \in \{0, \dots, l-1\}$, implies $h''_\kappa > h''_{\kappa+1}$.

Furthermore, if $h''_\kappa I_{\kappa+1} h''_{\kappa+1}$, where $\kappa \in \{0, \dots, l-1\}$, then, applying the argument used in the proof of (b), we see that $h''_\kappa \sim h''_{\kappa+1}$.

Consequently, (17), which holds with at least one relation being strict, implies

$$(h = h''_0) \succ h''_1 \succ \cdots \succ (h''_l = h'),$$

with at least one strict dominance. Since \succ is reflexive and transitive, Lemma 1*a in Sen (2017, p. 56) implies $h > h'$. This establishes (a). ■

A.8. Proof of Proposition 5

Proposition 5. *The binary relation \succ_{DE^*} is a hierarchical pre-order on \mathcal{H} satisfying **A**, **RP**, and **DE***.*

Proof. Reflexivity and transitivity are immediate by the definition of \succ_{DE^*} : \succ_{DE^*} is defined as the transitive closure of $\bigcup_{k=0}^{\infty} S_k$, where $S_0 = \sim^*$ ensures $h \succ_{DE^*} h$ for all $h \in \mathcal{H}$, and the transitive closure guarantees transitivity.

To verify **A** and **RP**, consider any two hierarchies $h, h' \in \mathcal{H}$. If h' is a relabeling or a replica of h , then $h' \sim^* h$. Since $S_0 = \sim^*$ is included in \succ_{DE^*} , it follows that $h'S_0h$ and hS_0h' , implying $h' \succ_{DE^*} h$ and $h \succ_{DE^*} h'$. Thus, $h' \sim_{DE^*} h$, so \succ_{DE^*} satisfies **A** and **RP**.

Suppose that $h' \in h \uplus h_0$ and $h_0 \succ_{DE^*} h$ for some $h_0 \in \mathcal{H}$. Since $h_0 \succ_{DE^*} h$, there exists a finite sequence of hierarchies

$$h_0 = h'_0, h'_1, \dots, h'_m = h$$

such that

$$h'_0 S_{k_0} h'_1 S_{k_1} \dots S_{k_{m-1}} h'_m,$$

where each S_{k_l} is from the sequence defining \succ_{DE^*} . Define

$$k^* = \max\{k_0, \dots, k_{m-1}\},$$

so the sequence can be rewritten as

$$(h_0 = h'_0) S_{k^*} h'_1 S_{k^*} \dots S_{k^*} (h'_m = h).$$

Given $h' \in h \uplus h_0$ and $h_0 S_{k^*} h$, the definition of S_{k^*+1} implies

$$(h', h) \in S_{k^*+1},$$

since there exist $H' \sim^* h'$, $H \sim^* h$ such that $H' \in H \uplus h_0$, and $h_0 S_{k^*} h$. Thus, $h' \succ_{DE^*} h$.

To confirm strict dominance, note that since $h' \in h \uplus h_0$, h' has a higher maximum depth than h . Suppose, for contradiction, that $h \succ_{DE^*} h'$. Then there exists a sequence

$$h = h''_0, h''_1, \dots, h''_n = h'$$

with $h''_l S_{k_l} h''_{l+1}$. However, each step in the sequence via S_{k_l} either preserves depth (via \sim^*) or increases maximum depth (via the \uplus operation), making it impossible for h (with lower depth) to be obtained from h' (with higher depth). Thus, $h \not\succ_{DE^*} h'$, and hence $h' \succ_{DE^*} h$.

Consequently, \succ_{DE^*} satisfies **DE***. Therefore, \succ_{DE^*} is a hierarchical pre-order on \mathcal{H} satisfying **A**, **RP**, and **DE***. ■

A.9. Proof of Theorem 5

Theorem 5. A hierarchical pre-order on \mathcal{H} satisfies **A**, **RP**, and **DE*** if and only if it is \succ_{DE^*} -consistent.

Proof. (Sufficiency.) Assume that \succ is \succ_{DE^*} -consistent. By **Proposition 5**, \succ_{DE^*} satisfies **A**, **RP**, and **DE***. Since \succ is \succ_{DE^*} -consistent, it follows that \succ also satisfies these axioms.

(Necessity.) Suppose that \succ satisfies **A**, **RP**, and **DE***. We need to show that \succ is \succ_{DE^*} -consistent, i.e., for all $h, h' \in \mathcal{H}$:

- If $h \succ_{DE^*} h'$, then $h \succ h'$.
- If $h \sim_{DE^*} h'$, then $h \sim h'$.

First, suppose $h \sim_{DE^*} h'$. Then there exists a cycle in \succ_{DE^*} , which, by a lemma analogous to **Lemma 3**, implies that all relations in the cycle are indifferences. Specifically, there are hierarchies such that

$$(h = h_0) S_{k_1} h_1 S_{k_2} \dots S_{k_m} (h_m = h' = h'_0) S_{k'_1} h'_1 S_{k'_2} \dots S_{k'_n} (h'_n = h),$$

and thus

$$(h = h_0)I_{k_1}h_1I_{k_2}\cdots I_{k_m}(h_m = h'),$$

where I_k is the symmetric part of S_k . Since I_k implies \sim^* , we have

$$h \sim^* h_1 \sim^* \cdots \sim^* h',$$

and thus $h \sim h'$ because \succsim satisfies **A** and **RP**.

Next, suppose $h \succ_{DE^*} h'$. Then there exists a sequence

$$(h = h''_0)S_{k_1}h''_1S_{k_2}\cdots S_{k_l}(h''_l = h'), \quad (18)$$

with at least one strict relation. We need to show that this implies $h \succ h'$.

First, we show that if (h, h') is in the transitive closure of S_k , then $h \succsim h'$. We proceed by induction on k .

For $k = 0$, since $S_0 = \sim^*$, and \succsim satisfies **A** and **RP**, hS_0h' implies $h \sim h'$.

Assume that for all $k \leq \kappa$, if (h, h') is in the transitive closure of S_k , then $h \succsim h'$.

For $k = \kappa + 1$, if (h, h') is in the transitive closure of $S_{\kappa+1}$, there exists a sequence

$$(h = \bar{h}_0)S_{\kappa+1}\bar{h}_1S_{\kappa+1}\cdots S_{\kappa+1}(\bar{h}_\ell = h'). \quad (19)$$

For each $\bar{h}_{l'}S_{\kappa+1}\bar{h}_{l'+1}$, either $\bar{h}_{l'}S_{\kappa}\bar{h}_{l'+1}$, which by induction implies $\bar{h}_{l'} \succsim \bar{h}_{l'+1}$, or there exist a sequence

$$\tilde{h}_0S_{\kappa}\cdots S_{\kappa}(\tilde{h}_{m'} = \bar{h}_{l'+1})$$

and $H_{l'} \sim^* \bar{h}_{l'}$, $H_{l'+1} \sim^* \bar{h}_{l'+1}$ such that $H_{l'} \in H_{l'+1} \uplus \tilde{h}_0$.

Since $(\tilde{h}_0, \tilde{h}_{m'} = \bar{h}_{l'+1})$ is in the transitive closure of S_{κ} , by induction, $\tilde{h}_0 \succsim \bar{h}_{l'+1}$.

Because $\tilde{h}_0 \succsim \bar{h}_{l'+1} \sim^* H_{l'+1}$, we have $\tilde{h}_0 \succsim \bar{h}_{l'+1} \sim H_{l'+1}$, since \succsim satisfies **A** and **RP**.

The relations $\tilde{h}_0 \succsim \bar{h}_{l'+1} \sim H_{l'+1}$ imply $\tilde{h}_0 \succsim H_{l'+1}$. This follows from the reflexivity and transitivity of \succsim (Sen, 2017, Lemma 1*a, p. 56).

Then, since $H_{l'} \in H_{l'+1} \uplus \tilde{h}_0$ and $\tilde{h}_0 \succsim H_{l'+1}$, and \succsim satisfies **DE***, we have $H_{l'} \succ H_{l'+1}$.

Thus,

$$\bar{h}_{l'} \sim^* H_{l'} \succ H_{l'+1} \sim^* \bar{h}_{l'+1},$$

so

$$\bar{h}_{l'} \sim H_{l'} \succ H_{l'+1} \sim \bar{h}_{l'+1},$$

and this chain implies $\bar{h}_{l'} \succ \bar{h}_{l'+1}$.

We have shown that $\bar{h}_{l'}S_{\kappa+1}\bar{h}_{l'+1}$ implies $\bar{h}_{l'} \succ \bar{h}_{l'+1}$. Consequently, (19) implies

$$h = \bar{h}_0 \succ \bar{h}_1 \succ \cdots \succ \bar{h}_\ell = h',$$

whence $h \succ h'$ by transitivity of \succ .

We conclude that if (h, h') is in the transitive closure of S_k , then $h \succ h'$. Hence, using (18), which holds with at least one strict relation, we see that

$$(h = h''_0) \succ h''_1 \succ \cdots \succ (h''_l = h'),$$

with at least one \succ . Hence, $h \succ h'$, as required. ■

A.10. Proof of Proposition 6

Proposition 6. *The binary relation \tilde{R} is a hierarchical pre-order on \mathcal{H} satisfying all the axioms in $\bigcup_{i \in I} A_i$.*

Proof. For any hierarchy $h \in \mathcal{H}$, each pre-order \succsim_i is reflexive. Thus, $h \succsim_i h$ for all $i \in I$, implying $h\tilde{R}h$. Hence, \tilde{R} is reflexive.

By construction, \tilde{R} is the transitive closure of $\bigcup_{i \in I} \succsim_i$, so \tilde{R} is transitive.

Let $A \in \bigcup_{i \in I} A_i$ be an arbitrary axiom. By assumption, each \succsim_i satisfies its corresponding axioms in A_i . For any $h, h' \in \mathcal{H}$, A imposes a condition $P(h, h') \Rightarrow Q(h, h', \succsim_i)$ for some $i \in I$. Since \tilde{R} contains \succsim_i , the condition $Q(h, h', \tilde{R})$ holds. Thus, \tilde{R} inherits all axioms in $\bigcup_{i \in I} A_i$.

Therefore, \tilde{R} is a reflexive and transitive hierarchical pre-order satisfying all axioms in $\bigcup_{i \in I} A_i$. ■

A.11. Proof of Theorem 6

Theorem 6. *A hierarchical pre-order \succsim on \mathcal{H} satisfies all A_i for $i \in I$ if and only if it is \tilde{R} -consistent.*

Proof. (Sufficiency.) Assume that \succsim is \tilde{R} -consistent. By Proposition 6, \tilde{R} satisfies all axioms in $\bigcup_i A_i$. Pick any axiom $A \in \bigcup_i A_i$ and suppose that $P(h, h') \Rightarrow Q(h, h', \tilde{R})$ is its conditional statement for \tilde{R} . To see that \succsim satisfies A , assume $P(h, h')$. Then $Q(h, h', \tilde{R})$ holds and so $Q(h, h', \succsim)$ must also hold by \tilde{R} -consistency. For example, if $Q(h, h', \tilde{R})$ represents the statement “ $h\tilde{R}h'$,” then either $h\tilde{P}h'$ or $h\tilde{I}h'$, where \tilde{P} and \tilde{I} denote the asymmetric and symmetric parts of \tilde{R} , respectively. If $h\tilde{P}h'$, then $h \succ_i h'$ by \tilde{R} -consistency, and if $h\tilde{I}h'$, then $h \sim_i h'$ by \tilde{R} -consistency.

(Necessity.) Suppose that \succsim is a hierarchical pre-order on \mathcal{H} satisfying all the axioms in $\bigcup_i A_i$. We must show that \succsim is \tilde{R} -consistent, i.e., that for every pair of hierarchies $h, h' \in \mathcal{H}$:

$$(a) \quad h\tilde{P}h' \Rightarrow h \succ h'.$$

$$(b) \quad h\tilde{I}h' \Rightarrow h \sim h'.$$

Suppose first that $h\tilde{I}h'$. Then there are finitely many hierarchies in \mathcal{H} ,

$$h_0, \dots, h_k, h'_0, \dots, h'_L,$$

such that

$$(h = h_0)(\cup_i \succsim_i)h_1(\cup_i \succsim_i) \cdots (\cup_i \succsim_i)(h_k = h' = h'_0)(\cup_i \succsim_i)h'_1(\cup_i \succsim_i) \cdots (\cup_i \succsim_i)(h'_L = h).$$

Since this chain forms a cycle, by the second condition in the definition of a composite axiomatic collection of hierarchical pre-orders, the cycle must be formed entirely by indifference in each individual \succsim_i :

$$(h = h_0) \sim_{i_1} h_1 \sim_{i_2} \cdots \sim_{i_k} (h_k = h' = h'_0) \sim_{i'_1} h'_1 \sim_{i'_2} \cdots \sim_{i'_L} (h'_L = h).$$

By assumption, this implies

$$(h = h_0) \sim^* h_1 \sim^* \cdots \sim^* (h_k = h')$$

(see the first bullet point in the fourth condition of the definition of a composite axiomatic collection of hierarchical pre-orders). Since \succsim satisfies all the axioms in $\bigcup_i A_i$, and since A_i includes the **A** and **RP** axioms for some i (see the third condition in the definition of a composite axiomatic collection of hierarchical pre-orders), we see that \succsim satisfies the **A** and **RP** axioms. Hence,

$$(h = h_0) \sim h_1 \sim \dots \sim (h_k = h'),$$

which yields $h \sim h'$ by transitivity of \sim . This establishes (b).

It remains to prove (a). Suppose that $h \tilde{P} h'$. Then there are finitely many hierarchies in \mathcal{H} ,

$$h''_0, \dots, h''_l,$$

such that

$$(h = h''_0)(\cup_i \succsim_i)h''_1(\cup_i \succsim_i) \dots (\cup_i \succsim_i)(h''_l = h'), \quad (20)$$

with at least one strict relation.

Suppose that $h''_\kappa(\cup_i \succsim_i)h''_{\kappa+1}$ and not $h''_{\kappa+1}(\cup_i \succsim_i)h''_\kappa$, where $\kappa \in \{0, \dots, l-1\}$. Then

$$h''_\kappa \succ_j h''_{\kappa+1}, \quad \text{some } j. \quad (21)$$

Hence, there exist finitely many hierarchies in \mathcal{H} ,

$$h''_\kappa = \tilde{h}_0, \dots, \tilde{h}_m = h''_{\kappa+1}, \quad (22)$$

such that

$$P_\ell(\tilde{H}_\ell, \tilde{H}_{\ell+1}, \mathcal{H}_\ell, \succsim_j) \quad \text{holds for each } \ell \in \{0, \dots, m-1\} \text{ and some } \tilde{H}_\ell \sim_j \tilde{h}_\ell \text{ and } \tilde{H}_{\ell+1} \sim_j \tilde{h}_{\ell+1}, \quad (23)$$

where each $P_\ell(\tilde{H}_\ell, \tilde{H}_{\ell+1}, \mathcal{H}_\ell, \succsim_j)$ is a predicate corresponding to axiom $a_\ell \in A_j$.

Since \succsim_j satisfies the axiom a_ℓ , we have

$$P_\ell(\tilde{H}_\ell, \tilde{H}_{\ell+1}, \mathcal{H}_\ell, \succsim_j) \Rightarrow Q_\ell(\tilde{H}_\ell, \tilde{H}_{\ell+1}, \succsim_j). \quad (24)$$

By (23) and (24), we see that $Q_\ell(\tilde{H}_\ell, \tilde{H}_{\ell+1}, \succsim_j)$ holds for each $\ell \in \{0, \dots, m-1\}$. This, together with (21) and (22), implies

$$h''_\kappa \sim_j \tilde{H}_0 \succsim_j \dots \succsim_j \tilde{H}_m \sim_j h''_{\kappa+1},$$

with at least one strict dominance. Since \succsim satisfies axiom a_ℓ for each ℓ , we have

$$Q_\ell(\tilde{H}_\ell, \tilde{H}_{\ell+1}, \succsim_j) = Q_\ell(\tilde{H}_\ell, \tilde{H}_{\ell+1}, \succsim), \quad \text{for each } \ell.$$

Consequently,

$$h''_\kappa \sim_j \tilde{H}_0 \succsim \dots \succsim \tilde{H}_m \sim_j h''_{\kappa+1},$$

with at least one strict dominance. Hence, from the fourth condition in the definition of a composite axiomatic collection of hierarchical pre-orders, we see that

$$h''_\kappa \sim^* \tilde{H}_0 \succsim \dots \succsim \tilde{H}_m \sim^* h''_{\kappa+1},$$

with at least one strict dominance. Since \succsim satisfies **A** and **RP**, it follows that

$$h''_{\kappa} \sim \tilde{H}_0 \succsim \cdots \succsim \tilde{H}_m \sim h''_{\kappa+1},$$

with at least one strict dominance. Consequently, because \succsim is reflexive and transitive, we obtain $h''_{\kappa} > h''_{\kappa+1}$ (Sen, 2017, Lemma 1*a, p. 56).

We conclude that $h''_{\kappa}(\cup_i \succsim_i)h''_{\kappa+1}$ and not $h''_{\kappa+1}(\cup_i \succsim_i)h''_{\kappa}$, where $\kappa \in \{0, \dots, l-1\}$, implies $h''_{\kappa} > h''_{\kappa+1}$.

Furthermore, if $h''_{\kappa}(\cup_i \succsim_i)h''_{\kappa+1}$ and $h''_{\kappa+1}(\cup_i \succsim_i)h''_{\kappa}$, where $\kappa \in \{0, \dots, l-1\}$, then, applying the argument used in the proof of (b), we see that $h''_{\kappa} \sim h''_{\kappa+1}$.

Hence, (20), which holds with at least one strict dominance, implies

$$(h = h''_0) \succsim h''_1 \succsim \cdots \succsim (h'_l = h''),$$

with at least one strict dominance. Since \succsim is reflexive and transitive, Lemma 1*a in Sen (2017, p. 56) implies $h > h'$. This establishes (a). ■

A.12. Proof of Proposition 7

Proposition 7. *The collection*

$$\{\succsim_H, \succsim_{UE}, \succsim_{DE}, \succsim_{UE^*}, \succsim_{DE^*}\}$$

is a composite axiomatic collection of hierarchical pre-orders.

Proof. We verify that the given collection of hierarchical pre-orders satisfies the conditions required by the definition of a composite axiomatic collection of hierarchical pre-orders.

1. Each hierarchical pre-order in the collection satisfies the following axioms:
 - \succsim_H : **A**, **RP**, and **SR**.
 - \succsim_{UE} : **A**, **RP**, and **UE**.
 - \succsim_{DE} : **A**, **RP**, and **DE**.
 - \succsim_{UE^*} : **A**, **RP**, and **UE***.
 - \succsim_{DE^*} : **A**, **RP**, and **DE***.
2. Let us denote the collection of hierarchical pre-orders by $\{\succsim_i\}_{i \in I}$. Suppose there exists a cycle in the union of these pre-orders:

$$h_0(\cup_i \succsim_i)h_1(\cup_i \succsim_i) \cdots (\cup_i \succsim_i)h_n(\cup_i \succsim_i)h_0.$$

For each pre-order in the collection, the average depth weakly decreases along this chain, and strictly decreases for any strict relation. Therefore, this cycle can only be consistent with indifference in each individual pre-order:

$$h_0 \sim_{i_1} h_1 \sim_{i_2} \cdots \sim_{i_n} h_n \sim_{i_{n+1}} h_0.$$

3. Each member of the collection satisfies **A** and **RP**.

4. For each $i \in I$:

- If $h \sim_i h'$ then $h \sim^* h'$. This follows from:
 - Lemma 4 in [Carbonell-Nicolau \(2025b\)](#) if $\sim_i = \sim_H$.
 - The proof of [Theorem 1](#) if $\sim_i = \sim_{UE}$.
 - The proof of [Theorem 3](#) if $\sim_i = \sim_{DE}$.
 - The proof of [Theorem 4](#) if $\sim_i = \sim_{UE^*}$.
 - The proof of [Theorem 5](#) if $\sim_i = \sim_{DE^*}$.
- Suppose that $h >_i h'$. Then:
 - If $>_i = >_H$, then h' can be obtained from some relabeling H of h by successive removals of subordination relations ([Carbonell-Nicolau, 2025b](#), Theorem 1). Hence, there is a finite sequence

$$H = h_0, h_1, \dots, h_K = h'$$

such that $h_l > h_{l+1}$ for each $l \in \{0, \dots, K-1\}$ and every hierarchical pre-order \succsim satisfying [SR](#).

- If $>_i = >_{UE}$, then

$$\exists h_0, H' \sim^* h', H \sim^* h : H \in h_0 \oplus H'.$$

Therefore, $H > H'$ for every hierarchical pre-order \succsim satisfying [UE](#).

- If $>_i = >_{DE}$, then

$$\exists h_0, H' \sim^* h', H \sim^* h : H \in H' \boxplus h_0.$$

Therefore, $H > H'$ for every hierarchical pre-order \succsim satisfying [DE](#).

- If $>_i = >_{UE^*}$, then there are finitely many hierarchies in \mathcal{H} ,

$$h''_0, \dots, h''_l,$$

such that

$$(h = h''_0)R_{k_1}h''_1R_{k_2}\dots R_{k_l}(h''_l = h'),$$

with at least one strict relation.

From the proof of [Theorem 4](#), we know that the following holds:

- * If $h''_{\kappa}P_{k_{\kappa+1}}h''_{\kappa+1}$, where $\kappa \in \{0, \dots, l-1\}$, then

$$\begin{aligned} \exists \hat{h}_0 R_{k_{\kappa+1}-1} \dots R_{k_{\kappa+1}-1} \hat{h}_M = h''_{\kappa+1}, H''_{\kappa} \sim^* h''_{\kappa}, H''_{\kappa+1} \sim^* h''_{\kappa+1} \\ : H''_{\kappa} \in \hat{h}_0 \uplus H''_{\kappa+1}, \end{aligned}$$

so that $H''_{\kappa} >_{UE^*} H''_{\kappa+1}$, and $H''_{\kappa} > H''_{\kappa+1}$ for every hierarchical pre-order \succsim satisfying [A](#), [RP](#), and [UE*](#).

- * If $h''_{\kappa}I_{k_{\kappa+1}}h''_{\kappa+1}$, where $\kappa \in \{0, \dots, l-1\}$, then $h''_{\kappa} \sim^* h''_{\kappa+1}$, so that $h''_{\kappa} \sim_{UE^*} h''_{\kappa+1}$, and $h''_{\kappa} \sim h''_{\kappa+1}$ for every \succsim satisfying [A](#) and [RP](#).

– If $\succ_i = \succ_{DE^*}$, then there exists a sequence

$$(h = h''_0)S_{k_1}h''_1S_{k_2}\cdots S_{k_l}(h''_l = h'),$$

with at least one strict relation.

From the proof of **Theorem 5**, we know that the following holds:

* If $h''_\kappa S_{k_{\kappa+1}} h''_{\kappa+1}$ and not $h''_{\kappa+1} S_{k_{\kappa+1}} h''_\kappa$, where $\kappa \in \{0, \dots, l-1\}$, then

$$\begin{aligned} \exists \hat{h}_0 S_{k_{\kappa+1}-1} \cdots S_{k_{\kappa+1}-1} \hat{h}_M = h''_{\kappa+1}, H''_\kappa \sim^* h''_\kappa, H''_{\kappa+1} \sim^* h''_{\kappa+1} \\ : H''_\kappa \in H''_{\kappa+1} \uplus \hat{h}_0, \end{aligned}$$

so that $H''_\kappa \succ_{DE^*} H''_{\kappa+1}$, and $H''_\kappa \succ H''_{\kappa+1}$ for every hierarchical pre-order \succ satisfying **A**, **RP**, and **DE***.

* If $h''_\kappa S_{k_{\kappa+1}} h''_{\kappa+1}$ and $h''_{\kappa+1} S_{k_{\kappa+1}} h''_\kappa$, where $\kappa \in \{0, \dots, l-1\}$, then $h''_\kappa \sim^* h''_{\kappa+1}$, so that $h''_\kappa \sim_{DE^*} h''_{\kappa+1}$, and $h''_\kappa \sim h''_{\kappa+1}$ for every \succ satisfying **A** and **RP**. ■

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