CSE 4800: Thesis

# Binary Function Clone Detection Using MBA Simplification



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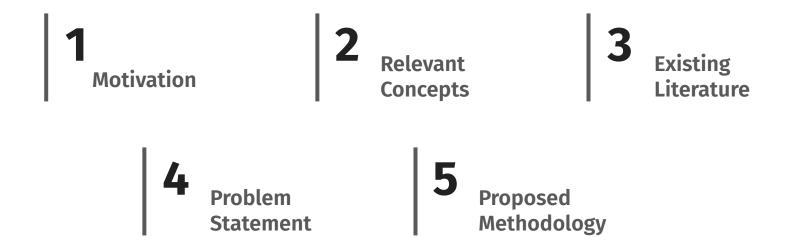
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### **Table of Contents**



#### **Motivation**

- **Copyright infringement** is a serious issue in the software development industry.
- Verification and code reuse identification in closed-source applications is significantly difficult.
- Cybersecurity challenges increase as attackers modify the existing malwares to evade antivirus detection.
- Existing methodologies for finding function clone is mostly Machine learning based which is difficult to scale.
- A systematic method to find function clones from binary files could be scalable and easy to use.

# **Function Cloning**

Function cloning is a method of **modifying** a function in such a way that the changes are limited to the internal code structure, **without altering** the function's overall **behavior**.

- Function cloning can be broadly categorized into two types:
  - → Optimization
  - → Obfuscation

# **Function Cloning (Optimization)**

```
func1:
     mov -20[rbp], edi
     mov -24[rbp], esi
                                                                       func1:
     mov -8[rbp], 10
                                  int func1(int a, int b){
                                                                                  eax, edi
     mov eax, -8[rbp]
                                                                             mov
                                        int x;
                                                                                  eax, 10
                                                                             xor
     xor eax, -20[rbp]
                                        int y;
                                                                                  eax, esi
                                                                             sub
     sub eax, -24[rbp]
                                       x = 10;
                                                                             imul esi, eax
     mov -4[rbp], eax
                                       y = (x ^a ) - b;
                                                                             lea eax, [rsi+rdi]
     mov eax, -24[rbp]
                                        return b * y + a;
                                                                             ret
     imul eax, -4[rbp]
     mov edx, eax
     mov eax, -20[rbp]
         eax, edx
     add
     ret
                                                                                    01
                                             Source
      No Optimization
                                            Function
                                                                               Optimization
```

# **Function Cloning (Obfuscation)**

**Equivalent Obfuscated** 

**Function** 

```
func1:
                                              edx, edi
                                         mov
                                              edx, 10
                                         xor
                                                                        func1:
int func1(int a, int b) {
                                              eax, esi
                                         mov
                                                                                   eax, edi
    int x;
                                                                              mov
                                         not
                                              eax
    int y;
                                         and eax, edx
                                                                              xor
                                                                                   eax, 10
   x = 10;
                                                                                   eax, esi
                                         add eax, eax
   y = (((x ^a a) & \sim b) << 1)
                                                                              imul esi, eax
                                         xor edx, esi
- ((x ^ a) ^ b);
                                                                              lea eax, [rsi+rdi]
                                         sub eax, edx
    return ((b * y - ~a) -1);
                                                                              ret
                                         imul esi, eax
                                              edi
                                         not
                                         sub esi, edi
                                         lea eax, -1[rsi]
                                         ret
```

Optimized and

obfuscated

01 Optimized

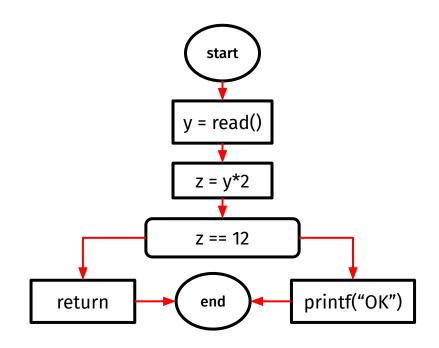
**Function** 

# Relevant Concepts

# **Control Flow Graph (CFG)**

```
int f() {
    ...
    y = read();
    z = y * 2;
    if (z == 12) {
        return;
    } else {
        printf("OK");
    }
}
```

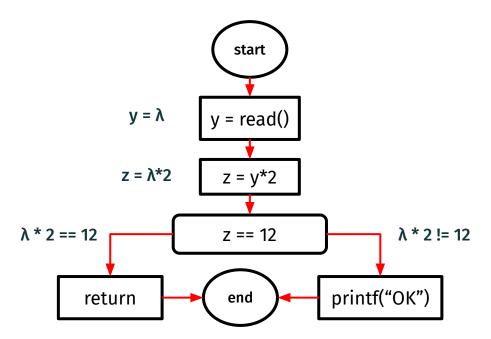
- Represents flow of execution
- **Nodes** → Basic block of codes
- **Edges** → Flow of control



# **Symbolic Execution**

- Explores a single control flow path within a CFG.
- Execution is carried out using **symbolic inputs** rather than specific, concrete values.

```
int f() {
    ...
    y = read();
    z = y * 2;
    if (z == 12) {
        return;
    } else {
        printf("OK");
    }
}
```



# **Mixed Boolean Arithmetic (MBA)**

#### **MBA Expression:**

Expression containing both traditional arithmetic operations (+, -, ×, . . .) and bitwise boolean operators (∧, ∨, ¬, ⊕, . . .)

$$\sum_{i \in I} a_i e_i(x_1, \dots, x_t)$$
 where,  
 
$$x_t \to \text{variables}$$
 
$$a_i \to \text{constant coefficient}$$
 
$$e_i(x_1, \dots, x_t) \to \text{bitwise expressions}$$

Figure: Formal Definition of MBA Expression<sup>1</sup>

# Mixed Boolean Arithmetic (MBA)

#### **MBA Expression:**

◆ A sequence of assembly instructions can be represented as MBA expressions.

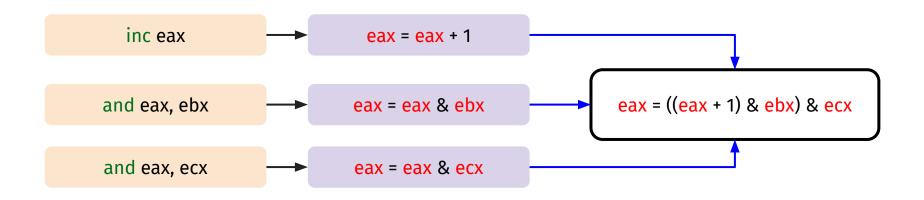


Figure: Assembly instructions to MBA expressions.

**Existing Literature** 

# Binary Function Clone Search in the Presence of Code Obfuscation and Optimization over Multi-CPU Architectures

ASIA CCS '23: Proceedings of the 2023 ACM Asia Conference on Computer and Communications Security, July 2023

Abdullah Qasem, Concordia University

Mourad Debbabi, Concordia University

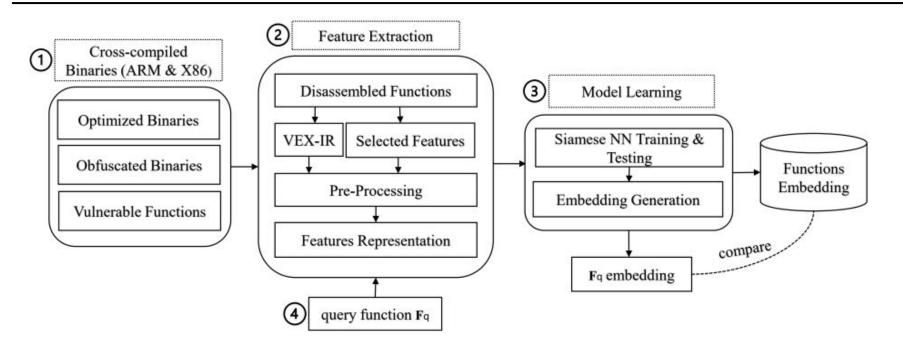
Bernard Lebel, Thales Research and Technologies

Marthe Kassouf, Hydro-Québec Research Institute

#### **Contribution**

- Developed and implemented BinFinder.
- Identified resilient, interpretable binary function features against code optimizations and obfuscations on multi-CPU architectures.
- Designed Siamese neural network architecture for training a model using proposed features to generate binary function embeddings.
- Evaluated BinFinder in 3 scenarios- x86 with different optimization and obfuscation, multi-cpu architecture with and without optimization and obfuscation.

# **BinFinder - Methodology**



[2] Binary Function Clone Search in the Presence of Code Obfuscation and Optimization over Multi-CPU Architectures ACM Asia Conference on Computer and Communications Security, July 2023

#### **BinFinder - Evaluation**

A		AUG	C		XM					
Approach	XC	XC+XB	XA	XM	small	medium	large	MRR10	Recall@1	
BinFinder	0.98	0.97	0.98	0.98	0.98	0.98	0.93	0.8	0.73	
GMN_OPC-200_e16	0.86	0.85	0.86	0.86	0.89	0.82	0.79	0.53	0.45	
GNN-s2v_GeminiNN_OPC-200_e5	0.78	0.81	0.82	0.81	0.84	0.77	0.79	0.36	0.28	
SAFE_ASM-list_e5	0.8	0.8	0.81	0.81	0.83	0.77	0.77	0.17	0.27	
Zeek	0.84	0.84	0.85	0.84	0.85	0.83	0.87	0.28	0.13	
asm2vec	0.62	0.81	0.74	0.69	0.63	0.7	0.78	0.12	0.07	

- **XA** → Different Architecture, Same Obfuscation, Different Optimization.
- **XB** → Same Architecture, Different Obfuscation.
- **XC** → Different Architecture, Same Obfuscation, Same Optimization.
- **XM** → Different Architecture, Different Obfuscation, Different Optimization.
- [2] Binary Function Clone Search in the Presence of Code Obfuscation and Optimization over Multi-CPU Architectures
  ACM Asia Conference on Computer and Communications Security, July 2023

# QsynthA Program Synthesis based Approach for Binary Code Deobfuscation

NDSS Symposium Binary Analysis Research Workshop, 2020

Robin David, Quarkslab

Luigi Coniglio, University of Trento

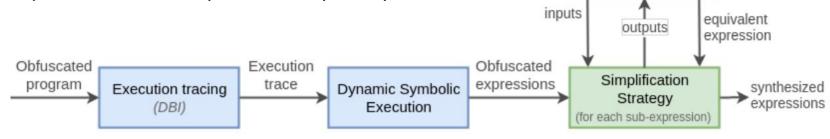
Mariano Ceccato, University of Verona

#### **Contribution**

- Introduced a combination of dynamic symbolic execution, dataflow graph extraction, and program synthesis against obfuscation.
- Proposed a black-box synthesis method based on offline enumerative search.
- Utilized precomputed lookup tables for near constant-time expression synthesis.
- Iteratively simplified complex obfuscated expressions.
- Empirical validation through a comparison with the deobfuscation tool Syntia, demonstrating superior accuracy and speed (20x faster).

# **Qsynth - Methodology**

- 1 Uses dynamic binary instrumentation (DBI) to trace paths in the binary.
- 2 From the **traces**, compute expressions at particular offsets using dynamic symbolic execution **(DSE)**.
- **3** A synthesis **oracle** simplifies a complex equation.



Enumerative Synthesis Oracle (generated once for all)

[3] Qsynth- A Program Synthesis based Approach for Binary Code Deobfuscation NDSS Symposium Binary Analysis Research Workshop, 2020

# **Qsynth - Evaluation**

	Mean expr. size		Simplification		M	Sem.	Time							
	Orig	Obf <sub>B</sub>	Synt	Ø	Partial	Full	Obf <sub>S</sub> /Orig	Synt/Obf <sub>B</sub>	Synt/Orig		Sym.Ex	Synthesis	Total	per fun.
Syntia	/	/	1	52	0	448	/	/	/	/	/	/	34 min	4.08s
<b>QSynth</b>	3.97	203.19	3.71	0	500	500	x35.03	x0.02	x0.94	500	1m20s	15s	1m35s	0.19s

Orig, Obf<sub>S</sub>, Obf<sub>B</sub>, Synt are respectively original, obfuscated (source, binary level) and synthesized expressions

#### [3] Qsynth- A Program Synthesis based Approach for Binary Code Deobfuscation

NDSS Symposium Binary Analysis Research Workshop, 2020

# Simplifying Mixed Boolean-Arithmetic Obfuscation by Program Synthesis and Term Rewriting

CCS '23: Proceedings of the 2023 ACM SIGSAC Conference on Computer and Communications Security, November 2023. Pages 2351–2365

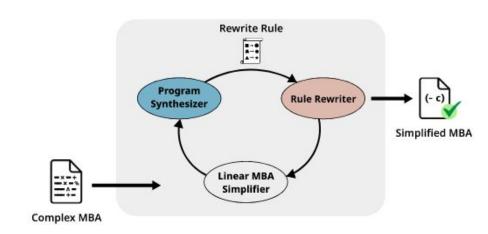
Jaehyung Lee, Hanyang University, Ansan, Korea Woosuk Lee, Hanyang University, Ansan, Korea

#### **Contributions**

- Proposes a novel and versatile method called **ProMBA** for deobfuscating MBA expressions.
- Combines *program synthesis*, *term rewriting*, and *algebraic simplification* methods.
- Unlike existing techniques, ProMBA can deobfuscate a much broader class of MBA expressions.
- Implemented the proposed method in an open source tool.
- The experimental results show that ProMBA outperforms the state-of-the-art deobfuscation method, **MBASolver** by a large margin.

## **ProMBA - Methodology**

- 1 Proposed method first simplifies linear MBA sub-expressions using an off-the-shelf deobfuscator.
- 2 Recursively simplifies non-linear sub-expressions by synthesizing simpler sub-expressions.
- 3 Applies the resulting rewrite rules to other sub-expressions until no further simplification is possible.

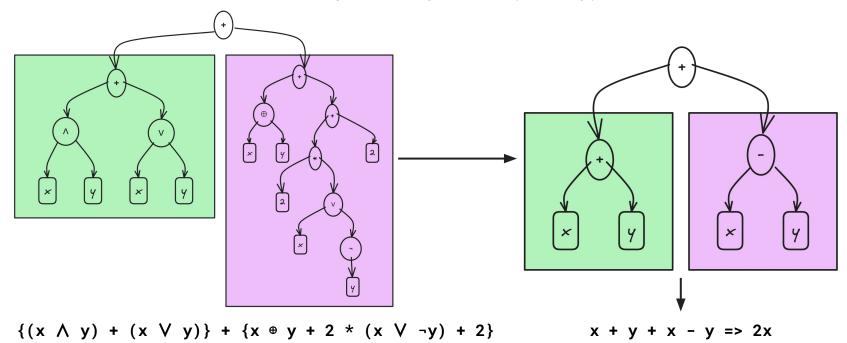


[4] Simplifying Mixed Boolean-Arithmetic Obfuscation by Program Synthesis and Term Rewriting ACM SIGSAC Conference on Computer and Communications Security, November 2023

# **MBA Simplification**

#### **Known Rules:**

- 1.  $x + y == (x \land y) + (x \lor y)$ 2.  $x y == x \oplus y + 2 * (x \lor \neg y) + 2$



# **ProMBA - Evaluation**

Dataset		Size (Avg.)			Success Rate		Time (Avg.)			
	ProMBA	MBASolver	SYNTIA	ProMBA	MBASolver	SYNTIA	ProMBA	MBASolver	SYNTIA	
MBA-Solver	11.76	21.67	4.61	80.31%	25.16%	17.45%	65.64s	6.3s	12.95s	
QSynth	17.48	77.71	4.72	62.8%	4.2%	22.8%	241s	64.83s	12.37s	
Loki	3.51	866.25	3.1	97.2%	0.13%	74.4%	100.03s	347.77s	2.07s	
Total	9.39	344.5	3.55	84.44%	13.19%	39.42%	100.36s	141.29s	8.81s	

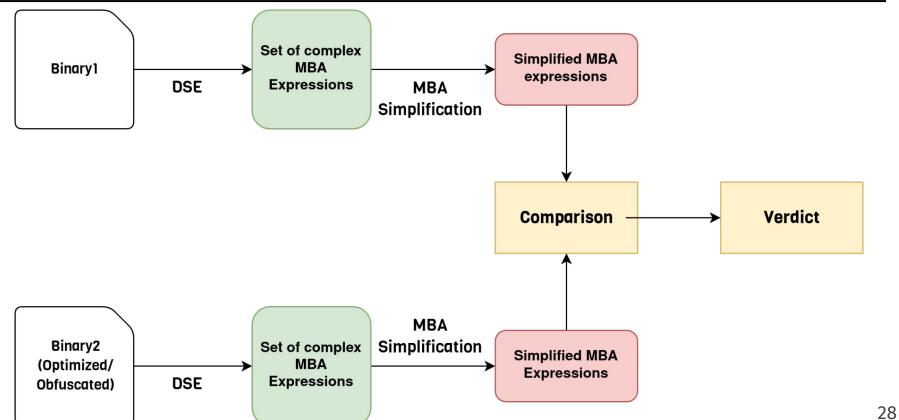
# **Limitations of existing literature**

- Machine Learning based solutions incur unnecessary overhead.
- Existing synthesis methods spend time **deriving** the same expressions **again and again**.
- Works properly for only small to medium sized binaries.

#### **Problem Statement**

Given two binary files identify any function equivalency through extraction and simplification of MBA expressions

# **Proposed Methodology**



# **Proposed Methodology**

- **1** Create obfuscated binaries from **TREX** binary database.
- **2** Extraction of complex MBA expressions from obfuscated binary through dynamic symbolic execution.
- **3** Simplify the set of MBA expressions.
- 4 Check for **equivalency** between set of expressions using possible means e.g. SMT, fuzzing, program synthesis, graph isomorphism etc.

# **Challenges**

- Extracting MBA expressions from branches and loops.
- Comparing equivalence of set of simplified expressions
- Choosing the appropriate form of verdict (binary vs similarity metrics).

#### Reference

- **1 MBA-Blast: Unveiling and Simplifying Mixed Boolean-Arithmetic Obfuscation** in the Proceedings of USENIX Security Symposium, 2021.
- 2 Binary Function Clone Search in the Presence of Code Obfuscation and Optimization over Multi-CPU Architectures
  - ACM Asia Conference on Computer and Communications Security, July 2023
- **Qsynth- A Program Synthesis based Approach for Binary Code Deobfuscation**NDSS Symposium Binary Analysis Research Workshop, 2020
- 4 Simplifying Mixed Boolean-Arithmetic Obfuscation by Program Synthesis and Term Rewriting ACM SIGSAC Conference on Computer and Communications Security, November 2023

# **Thank You**