

Rebooting Supercompilation for Haskell

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Rebooting Supercompilation for Haskell - Talk outline

- An overview of supercompilation.

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- What's interesting about it in the context of Haskell? Current state-of-the-art.

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- What's interesting about it in the context of Haskell? Current state-of-the-art.
- Overview of how it works.
- "But where's my supercompiler for Haskell?" My preliminary work and research goals.

Supercompilation: An overview

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- Evaluate programs in compile time.
- Make the most out of known inputs and definitions.
- Evaluate open terms.

Supercompilation in the context of Haskell

- Why is it interesting?
- In a sense, it's the "ultimate" optimization. ("-O99")
- An evaluator-based supercompiler optimizes in the sense that:
If we have programs \mathcal{P}_1 and \mathcal{P}_2 , and
 $\mathcal{P}_1 \Downarrow v$ in N steps and
 $\mathcal{P}_2 \Downarrow v$ in M steps,
we consider \mathcal{P}_2 optimized if $M < N$.
- An approximation, but works well in practice.
(i.e. if $M < N$ then usually M is a faster program)

Supercompilation in the context of Haskell

It generalizes:

- Deforestation(Wadler [1988])
- Partial evaluation
- Call-pattern specialization(Peyton Jones [2007])
- Ad-hoc optimizations via rewrite rules, e.g. shortcut fusion (Gill et al. [1993]) or library-specific rewrite rules
- "Optimizing SYB is Easy!" (Adams et al. [2014]) and "Optimizing Generics is Easy!" (Magalhães et al. [2010]) style "domain-specific" partial evaluators
- Function specialization(SPECIALIZE pragmas)
- ... and many more

Current state-of-the-art for Haskell

- Bolingbroke [2013] shows some great potential:
 - Up to 20x faster runtime.
 - Up to 100% reduction in allocation.
- But it also suffers from problems that are inherent to supercompilation:
 - "We do not attempt to supercompile the full Nofib suite because the other Nofib benchmarks are considerably more complicated and generally suffer from extremely long supercompilation times."
(Jonsson [201?] focuses on compilation performance, and reports *<3 seconds* for all the small programs from Nofib)
 - Up to 132x compile time.
 - Up to 2.8x generated code size.

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- **Matching:** Evaluating open terms lead to loops. Matcher tries to detect loops, returns information about how to refer to this new loop.

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- **Matching:** Evaluating open terms lead to loops. Matcher tries to detect loops, returns information about how to refer to this new loop.
- **Termination enforcement:** Because perfect matcher is not possible, and some programs just loop.

mapOfMap f g = map f . map g

h1 f g a = map f (map g a)

...

h4 f g a =

case (case a of

 [] -> []

 h1 : t1 -> g h1 : map g t1) of

 [] -> []

 h0 : t0 -> f h0 : map f t0

...

h6 f g a =

case a of

 [] -> []

 h : t -> f (g h) : map f (map g t)

h7 f g a =

case a of

 [] -> []

 h : t -> f (g h) : h7 f g t

Problems with supercompilation operations

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Propagating too much information may lead to work duplication.

```
let n = fib 100
    b = n + 1
    c = n + 2
in (b, c)
```

```
let b =
    let f = <fib, unrolled a few times>
    in f + 1
    c =
    let f = <fib, unrolled a few times>
    in f + 2
in (b, c)
```

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Splitter: (from Bolingbroke [2013])

Propagating too little information may lead to missing optimization opportunities.

```
let map = ...  
    ys = map f zs  
    xs = map g ys  
in Just xs
```

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```
xor x y = case x of True -> not y; False -> y
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```
goal = (xor a b, xor c c)
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```
xor' x y = case x of True -> not y; False -> y
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```
goal' = (xor' a b, xor' c c)
```


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goal = (xor a b, xor c c)
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xor' x y = case x of True -> not y; False -> y
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goal' = (xor' a b, xor' c c)
```

```
xor' x y = case x of True -> not y; False -> y
```

```
xor'' x = case x of True -> False; False -> False
```

```
goal' = (xor' a b, xor'' c)
```

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Termination enforcement:

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Termination enforcement:

Some programs just loop.

```
loop n = loop (n + 1)
```

```
countFrom n = n : countFrom (n + 1)
```

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Termination enforcement:

Some programs just loop.

```
loop n = loop (n + 1)
countFrom n = n : countFrom (n + 1)
```

Sometimes detecting loops is not so easy: (growing arguments)

```
reverse_acc [] acc = acc
reverse_acc (h : t) acc = reverse_acc t (h : acc)
goal lst = reverse_acc (reverse_acc lst []) []
...
h_ lst = ... reverse_acc t1 (h1 : []) ...
...
h_ lst = ... reverse_acc t2 (h2 : h1 : []) ...
...
```

"Where's my supercompiler for Haskell?"

- Bolingbroke [2013] has some solutions, and it documents and implements it nicely.
- But we still don't have something that we can use *today*.
- I'm rebooting the supercompiler!
- The goal here is to distribute it as a package, downloadable from Hackage.
- Then the research will follow.

Conclusions

Have a working implementation of supercompiler described in Bolingbroke [2013].

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Collecting benchmark programs - send yours! (with expected optimizations)

Create a benchmark suite like Nofib, but for supercompilation-specific problems. (pathological cases, programs with lots of intermediate data structures)

Once we have a working implementation:

- Focus on specific parts(matcher, splitter etc.). Try other ideas from the literature(e.g. homeomorphic embedding for matching)
- Work on some of the obvious improvements, like parallelizing the matcher.
- More experimental ideas:
 - Can we formulate it as a search problem and apply ideas from the literature?
 - Is profile-driven decision making possible?
 - Can we make use of existing rewrite rules mechanism?
 - Can we make use of free theorems?

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Thanks!

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