Rebooting Supercompilation for Haskell

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■ An overview of supercompilation.

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- What's interesting about it in the context of Haskell? Current state-of-the-art.
- Overview of how it works.
- "But where's my supercompiler for Haskell?" My preliminary work and research goals.

Supercompilation: An overview

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- Evaluate programs in compile time.
- Make the most out of known inputs and definitions.
- Evaluate open terms.

Supercompilation in the context of Haskell

- Why is it interesting?
- In a sense, it's the "ultimate" optimization. ("-O99")
- This optimizes in the sense that: If we have a programs \mathcal{P}_1 and \mathcal{P}_2 , and $\mathcal{P}_1 \Downarrow v$ in N steps and $\mathcal{P}_2 \Downarrow v$ in M steps, we consider \mathcal{P}_2 optimized if M < N.
- An approximation, but works well in practice.
 (i.e. if M<N then usually M is a faster program)

Supercompilation in the context of Haskell

It generalizes:

- Deforestation(Wadler [1988])
- Partial evaluation
- Call-pattern specialization(Peyton Jones [2007])
- Ad-hoc optimizations via rewrite rules, e.g. shortcut fusion (Gill et al. [1993]) or library-specific rewrite rules
- "Optimizing SYB is Easy!" (Adams et al. [2014]) and
 "Optimizing Generics is Easy!" (Magalhães et al. [2010]) style
 "domain-specific" partial evaluators
- Function specialization(SPECIALIZE pragmas)
- ... and many more

Current state-of-the-art

- Bolingbroke [2013] shows some great potential:
 - Up to 20x faster runtime.
 - Up to 100% reduction in allocation.
- But it also suffers from problems that are inherent to supercompilation:
 - "We do not attempt to supercompile the full Nofib suite because the other Nofib benchmarks are considerably more complicated and generally suffer from extremely long supercompilation times."
 - (Jonsson [201?] focuses on compilation performance, and reports < 3 seconds for all the small programs from Nofib)
 - Up to 132x compile time.
 - Up to 2.8x generated code size.

Bolingbroke [2013] laid out a great framework for supercompiling Haskell:

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- Matching: Evaluating open terms lead to loops. Matcher tries to detect loops, returns information about how to refer to this new loop.
- Termination checking: Because perfect matcher is not possible, and some programs just loop.

```
mapOfMap f g = map f . map g
h1 fga = map f (map g a)
h4 fga =
  case (case a of
         [] -> []
        h1 : t1 -> g h1 : map g t1) of
    [] -> []
    h0 : t0 -> f h0 : map f t0
h6 fga =
 case a of
   [] -> []
   h: t \rightarrow f(gh): map f(map g t)
h7 f g a =
  case a of
   [] -> []
   h : t -> f(gh) : h7 fgt
```

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Propagating too much information may lead to work duplication.

```
let n = fib 100
    b = n + 1
    c = n + 2
in (b, c)
```

```
let b =
    let f = <fib, unrolled a few times>
    in f + 1
    c =
    let f = <fib, unrolled a few times>
    in f + 2
in (b, c)
```

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```
Splitter: (from Bolingbroke [2013])
```

Propagating too little information may lead to missing optimization opportunities.

```
let map = ...
    ys = map f zs
    xs = map g ys
in Just xs
```

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Matcher: Injectivity of substitutions effect optimizations.

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```
(from Bolingbroke [2013])
```

```
xor x y = case x of True -> not y; False -> y
goal = (xor a b, xor c c)
```

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Matcher: Injectivity of substitutions effect optimizations. (from Bolingbroke [2013])

```
xor x y = case x of True -> not y; False -> y
goal = (xor a b, xor c c)
```

```
h0 \times y = case \times of True \rightarrow not y; False \rightarrow y

h1 = (h0 a b, h0 c c)
```

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Matcher: Injectivity of substitutions effect optimizations. (from Bolingbroke [2013])

```
h0 x y = case x of True -> not y; False -> y
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```

```
h0 x y = case x of True -> not y; False -> y
h1 x = case x of True -> False; False -> False
h2 = (h0 a b, h1 c)
```

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```
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loop n = loop (n + 1)
countFrom n = n : countFrom (n + 1)
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```
Some programs just loop.
loop n = loop (n + 1)
countFrom n = n : countFrom (n + 1)
Sometimes detecting loops is not so easy: (growing arguments)
reverse_acc [] acc = acc
reverse_acc (h : t) acc = reverse_acc t (h : acc)
goal lst = reverse_acc (reverse_acc lst []) []
. . .
h_ lst = ... reverse_acc t1 (h1 : []) ...
. . .
h_ lst = ... reverse_acc t2 (h2 : h1 : []) ...
. . .
```

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- But we still don't have something that we can use today.
- I'm rebooting the supercompiler!
- The goal here is to distribute it as a package, downloadable from Hackage.
- Then the research will follow.

Conclusions

Have a working implementation of supercompiler described in Bolingbroke [2013].

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Collecting benchmark programs - send yours! (with expected optimizations)

We can create a benchmark suite like Nofib, but for supercompilation-specific problems. (pathological cases, programs with lots of intermediate data structures)

Once we have a working implementation:

- Focus on specific parts(matcher, splitter etc.). Try other ideas from the literature(e.g. homeomorphic embedding for matching)
- Work on some of the obvious improvements, like parallelizing the matcher.
- More experimental ideas:
 - Can we formulate it as a search problem and apply ideas from the literature?
 - Is profile-driven decision making possible?
 - Can we make use of existing rewrite rules mechanism?
 - Can we make use of free theorems?

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Thanks!

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