## **Programming Project 5 Report**

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Author: Oscar Lopez

Class: CS 433 Operating Systems | Section 1

Professor: Dr. Zhang

## **Summary**

This report documents the implementation of page replacement algorithms for Assignment 5. The project includes implementations of FIFO (First-In-First-Out), LIFO (Last-In-First-Out), and LRU (Least Recently Used). These algorithms are tested against common page reference traces and evaluated for runtime efficiency and correctness. All implementations conform to a unified interface and adhere to object-oriented design principles.

## **Implementation Details**

The simulation is written in C++ and follows an object-oriented structure. A base class defines the common interface.

while each page replacement strategy inherits and implements algorithm-specific logic.

#### **Architecture Overview**

All algorithms derive from an abstract class `Replacement`, which provides a consistent interface.

This design supports clean modularity and easy substitution of replacement strategies.

Class Hierarchy:

class Replacement (base class)

**FIFOReplacement** 

LIFOReplacement

LRUReplacement

## **Key Components**

- Page Frame Table: Stores active pages

- Access Tracker: Used by LRU to track usage

- Eviction Logic: Specific to each algorithm

- Statistics: Tracks page faults and replacements

# **Algorithm-Specific Implementations**

## FIFOReplacement:

- Evicts the oldest page in a queue
- Time Complexity: O(1)

#### LIFOReplacement:

- Evicts the most recently inserted page using a stack
- Time Complexity: O(1)

LRUReplacement (Clock-based):

- Approximates LRU using a reference-bit clock buffer
- Uses a circular array and map for efficient lookup
- Time Complexity: O(n) in worst-case, O(1) amortized

## Implementation-Specific Details

The following outlines the exact implementation techniques used in the C++ source files for each page replacement algorithm.

FIFOReplacement (fifo\_replacement.cpp):

- Uses `std::queue<int> fifo\_queue` to store the order of page arrivals.
- Pages are inserted using `fifo\_queue.push()` and evicted using `fifo\_queue.front()` then popped.
- A `std::unordered\_set<int> page\_table` is used for fast membership checking.
- Eviction occurs only when page table reaches max frame size.

LIFOReplacement (lifo replacement.cpp):

- Implements a stack structure using `std::vector<int> lifo\_stack`.
- New pages are pushed to the back of the vector; the last page is evicted when needed.
- A `std::unordered\_set<int> page\_table` ensures O(1) lookup and removal.
- This structure leads to very high turnover under certain workloads.

LRUReplacement (Iru\_replacement.h / Iru\_replacement.cpp):

- Combines a fixed-size `int\* clock\_buffer` with a clock hand pointer for LRU approximation.
- A parallel 'bool\* reference bits' tracks recent usage of pages.
- `std::unordered\_map<int, int> page\_to\_index` maps pages to indices in the buffer.
- On access: updates the page's reference bit to true.
- On eviction: rotates clock hand and clears bits until it finds one with a 0.
- This approximates true LRU behavior with lower computational overhead.

Statistics and Profiling:

- Each algorithm logs page faults and evictions using counters defined in the base class `Replacement`.
- Execution time is captured using `std::chrono::high\_resolution\_clock` to ensure compliance with < 1s runtime for large inputs.
- All results are printed in a standardized format via the main simulator.

Unified Interface (replacement.h):

- All strategies implement `insert(int page)`, `evict()`, `exists(int page)`, and `access(int page)`.
- This design ensures algorithms are interchangeable at runtime and simplifies testing.

These implementation choices optimize each algorithm for clarity, performance, and ease of debugging.

#### **Performance Metrics**

Each algorithm tracks total references, page faults, and evictions.

Timing is captured using <chrono> to verify compliance with runtime requirements (< 1s on large inputs).

Sample Output - Small Trace:

Logical address: 1048576, page number: 1024, frame = 0, fault? 1 Logical address: 1049600, page number: 1025, frame = 1, fault? 1

Sample Output - Large Trace (LRU):

Number of references: 20000

Page faults: 1500 Replacements: 1490

Elapsed time: 1.1342 seconds

#### **Features**

- Multiple page replacement algorithms
- Clean inheritance-based design
- Precise statistics for evaluation
- Modular architecture

## **Advanced Features**

- Clock-based approximation for LRU
- Runtime profiling using <chrono>
- Dynamic frame sizing via CLI
- Efficient memory and container management

#### **Limitations and Drawbacks**

- LRU implementation is not true LRU (uses Clock approximation)
- No modeling of I/O cost or dirty pages
- All pages assumed to be equal cost
- No demand-based prefetching

## **Compilation and Execution**

Compilation Command:

g++ -std=c++17 main.cpp fifo\_replacement.cpp lifo\_replacement.cpp lru\_replacement.cpp pagetable.cpp replacement.cpp -o pager\_sim

**Execution Command:** 

./pager sim 1024 16

Input files required:

- small\_refs.txt
- large\_refs.txt

#### Conclusion

This project highlights key trade-offs in page replacement strategies. While FIFO and LIFO offer low runtime overhead,

they tend to underperform in real-world scenarios. The Clock-based LRU implementation provides a good balance between

runtime efficiency and replacement accuracy. Overall, the simulation demonstrates foundational OS memory management concepts.