

Oscar Lopez

CSUSM: CS 433 : Section 1

Dr. Xiaoyu Zhang

02-06-2025

## Overview

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This report analyzes the **Process Control Block (PCB) Table and Ready Queue System**, detailing its efficiency, optimizations, and computational complexity. The system utilizes a **priority queue (max heap) for process scheduling** and an **array-based PCB Table for process management**. The goal is to maintain an optimal balance between **speed, memory usage, and scalability**.

## Data Structures and Optimizations

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### PCB Table (`PCBTable`)

The `PCBTable` stores an array of pointers to PCB objects and provides constant time ( $O(1)$ ) access to processes using direct indexing.

#### Optimizations:

- Uses dynamic memory allocation (`new PCB*[]`) for flexible storage.
- Ensures proper memory deallocation in the destructor (`~PCBTable()`).
- Avoids fragmentation with a structured array-based implementation.

### Ready Queue (`ReadyQueue`)

The `ReadyQueue` implements a priority queue using a max heap, ensuring that higher-priority processes are scheduled before lower-priority ones.

#### Optimizations:

- Uses heapify operations (`heapifyUp()` and `heapifyDown()`) to maintain heap order efficiently.
- Implements dynamic allocation of the heap (`heap = new PCB*[capacity]`).
- Copy constructor and assignment operator prevent memory leaks.

## Logarithmic Runtime Analysis

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### Heap Operations (Priority Queue Complexity)

The max heap guarantees logarithmic runtime for key operations:

Operation	Time Complexity
Insertion ( <code>addPCB()</code> )	$O(\log n)$
Deletion ( <code>removePCB()</code> )	$O(\log n)$
Access highest-priority PCB ( <code>heap[0]</code> )	$O(1)$

Insertion (`heapifyUp()`) moves newly added PCBs logarithmically ( $O(\log n)$ ) up the heap, while removal (`heapifyDown()`) extracts the highest-priority PCB, ensuring efficient scheduling.

### PCB Table Operations Complexity

Operation	Time Complexity
Retrieve PCB ( <code>getPCB()</code> )	$O(1)$
Add PCB ( <code>addPCB()</code> )	$O(1)$

Array indexing enables fast lookups ( $O(1)$ ), making PCB retrieval efficient.

## Efficiency and Performance Considerations

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### Space Complexity

Data Structure	Space Complexity
PCBTable ( <code>PCBTable</code> )	$O(n)$
ReadyQueue (Heap-based priority queue)	$O(n)$

Heap scales dynamically, efficiently utilizing memory, while the PCBTable preallocates storage ( $O(n)$ ), reducing flexibility but ensuring fast access.

### Scalability Considerations

- The heap-based ReadyQueue scales logarithmically ( $O(\log n)$ ), making it efficient for large workloads.
- The PCBTable provides instant lookups ( $O(1)$ ) but is limited by static array sizing.

### Potential Bottlenecks & Further Optimizations

1. **Heap Resizing:** The fixed-size heap limits scalability. **Optimization:** Implement dynamic resizing (`std::vector<PCB*>`) for adaptive growth.

2. **Memory Management in PCBTable:** Each `PCB` is allocated separately, leading to fragmentation. **Optimization:** Use memory pools or preallocated buffers to improve efficiency.

## Conclusion

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This system provides efficient process scheduling using:

- Array-based `PCB` storage ( $O(1)$  lookups).
- Heap-based priority queue ( $O(\log n)$  insertions/removals).