

Equivalence of Hidden Markov Models with Continuous Observations

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Abstract

We consider the Sequential Probability Ratio Test applied to Hidden Markov Models. Given two Hidden Markov Models and a sequence of observations generated by one of them, the Sequential Probability Ratio Test attempts to decide which model produced the sequence. We show relationships between the execution time of such an algorithm and Lyapunov exponents of random matrix systems. Further, we give complexity results about the execution time taken by the Sequential Probability Ratio Test.

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1 Introduction

2 Preliminaries

We write \mathbb{N} for the set of non-negative integers, \mathbb{Q} for the set of rationals and \mathbb{Q}_+ for the set of positive rationals. For $d \in \mathbb{N}$ and a finite set Q we use the notation $|Q|$ for the number of elements in Q , $[d] = \{1, \dots, d\}$ and $[Q] = \{1, \dots, |Q|\}$. Vectors $\mu \in \mathbb{R}^N$ are viewed as row vectors. The norm $\|\mu\|$ is assumed to be the l_1 norm: $\|\mu\| = \sum_{i=1}^N |\mu_i|$. A matrix $M \in \mathbb{R}^{N \times N}$ is *stochastic* if M is non-negative and $\sum_{j=1}^N M_{i,j} = 1$ for all $i \in [N]$.

Throughout this paper, we use Σ to denote a finite set of *observations*. The set Σ^n is the set of words over Σ of length n and $\Sigma^* = \bigcup_{n=0}^{\infty} \Sigma^n$.

► **Definition 1.** A Hidden Markov Model (HMM) is a triple (Q, Σ, Ψ) where Q is a finite set of states, Σ is a set of observations, and the observation mass matrix $\Psi : \Sigma \rightarrow [0, \infty)^{|Q| \times |Q|}$ specifies the transitions such that $\sum_{a \in \Sigma} \Psi(a)$ is a stochastic matrix.

► **Example 2.** The HMM from the introduction is the triple $(\{q_1, q_2\}, \mathbb{R}, \Psi)$ with

$$\Psi(x) = TBC!?!? \tag{1}$$

We extend Ψ to the mapping $\Psi : \Sigma^* \rightarrow [0, \infty)^{|Q| \times |Q|}$ with $\Psi(x_1 \dots x_n) = \Psi(x_1) \times \dots \times \Psi(x_n)$ for $x_1, \dots, x_n \in \Sigma$. For the empty word ϵ we define $\Psi(\epsilon)$ to be the identity matrix. We say that $A \subseteq \Sigma^n$ is a *cylinder set* if $A = A_1 \times \dots \times A_n$ and $A_i \in \mathcal{G}$ for $i \in [n]$. For every n there is an induced measure space $(\Sigma^n, \mathcal{G}^n, \lambda^n)$ where \mathcal{G}^n is the smallest σ -algebra containing all cylinder sets in Σ^n and $\lambda^n(A_1 \times \dots \times A_n) = \prod_{i=1}^n \lambda(A_i)$ for any cylinder set $A_1 \times \dots \times A_n$. Let $A \subseteq \Sigma^n$ and write $A\Sigma^\omega$ for the set of infinite words over Σ where the first n observations fall in the set A . Given a HMM (Q, Σ, Ψ) and initial distribution π on Q viewed as vector $\pi \in \mathbb{R}^{|Q|}$, there is an induced probability space $(\Sigma^\omega, \mathcal{G}^*, \mathbb{P}_\pi)$ where Σ^ω is the set of infinite words over Σ , and \mathcal{G}^* is the smallest σ -algebra containing (for all $n \in \mathbb{N}$) all sets $A\Sigma^\omega$ where $A \subseteq \Sigma^n$ is a cylinder set and \mathbb{P}_π is the unique probability measure such that



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42 $\mathbb{P}_\pi(A\Sigma^\omega) = \sum_{w \in A} \|\pi\Psi(w)\|$ for any cylinder set $A \subseteq \Sigma^n$. Let $\mathbb{1}_F$ be the indicator random
 43 variable for the set $F \in \mathcal{G}^*$, i.e. $\mathbb{1}_F(w) = 1$ if $w \in F$ and $\mathbb{1}_F(w) = 0$ if $w \notin F$.

44 Let (Q, Σ, Ψ) be an HMM and let π_1, π_2 be two initial distributions. The *total variation*
 45 *distance* is

$$46 \quad d(\pi_1, \pi_2) = \sup_{E \in \mathcal{G}^*} |\mathbb{P}_{\pi_1}(E) - \mathbb{P}_{\pi_2}(E)|.$$

47 This supremum is actually a maximum due to Hahn's decomposition theorem; i.e., there is
 48 an event $S \in \mathcal{G}^*$ such that $d(\pi_1, \pi_2) = \mathbb{P}_{\pi_1}(S) - \mathbb{P}_{\pi_2}(S)$. We call π_1 and π_2 *distinguishable* if
 49 $d(\pi_1, \pi_2) = 1$. In [2] it is shown that distinguishability is decidable in polynomial time.

50 One could define the distinguishability of two pairs (C_1, π_1) and (C_2, π_2) where $C_i =$
 51 (Q_i, Σ, Ψ_i) are HMMs and π_i are initial distributions for $i = 1, 2$. We do not need that though,
 52 as we can define, in a natural way, a single HMM over the disjoint union of Q_1 and Q_2 and
 53 consider instead the distinguishability of π_1 and π_2 (where π_1, π_2 are appropriately padded
 54 with zeros).

55 – For the rest of the paper we assume that (Q, Σ, Ψ) is an HMM. L_n depending on π_i etc.

56 ► **Definition 3.** Let π_1 and π_2 be initial distributions. For $w \in \Sigma^\omega$ we write w_n for the
 57 length n prefix of w . For any $n \in \mathbb{N}$, the likelihood ratio L_n is a random variable on Σ^ω
 58 given by $L_n(w) = \frac{\|\pi_1\Psi(w_n)\|}{\|\pi_2\Psi(w_n)\|}$.

59 It follows from [2, Proposition 6] that L_n is a martingale and the following lemma holds due
 60 to Doob's forward convergence theorem[7].

61 ► **Lemma 4.** We have $\lim_{n \rightarrow \infty} L_n$ exists \mathbb{P}_{π_2} -almost surely.

62 The following result characterises when the limit is 0 and follows from [2].

63 ► **Theorem 5.** We have $\lim_{n \rightarrow \infty} L_n = 0$ \mathbb{P}_{π_2} -a.s. if and only if π_1 and π_2 are distinguishable.

64 ► **Example 6.** We will illustrate convergence of the likelihood ratio using an example from [4]
 65 where the authors use HMMs to model sleep cycles. They took measurements of 51 healthy and
 66 51 diseased individuals and using electrodes attached to the scalp, they read electrical signal
 67 data as part of an electroencephalography (EEG) during sleep. They split the signal into 30
 68 second intervals and mapped it onto the simplex $\Delta^3 = \{(x_1, x_2, x_3, x_4) \in \mathbb{R}_+^4 \mid \sum_{i=1}^4 x_i = 1\}$.
 69 For each individual this results in a time series of points in Δ^3 . They modelled this data
 70 using two HMMs, each with 5 states, for healthy and diseased individuals using a numerical
 71 maximum likelihood estimate. This resulted in two embedded Markov chains defined by the
 72 row-stochastic block diagonal matrices

$$73 \quad T_1 = \begin{bmatrix} 0.793 & 0.099 & 0.035 & 0.064 & 0.009 \\ 0.078 & 0.769 & 0.006 & 0.144 & 0.003 \\ 0.018 & 0.004 & 0.833 & 0.134 & 0.012 \\ 0.022 & 0.094 & 0.054 & 0.827 & 0.002 \\ 0.011 & 0.005 & 0.035 & 0.005 & 0.945 \end{bmatrix}, T_2 = \begin{bmatrix} 0.641 & 0.109 & 0.031 & 0.040 & 0.015 \\ 0.202 & 0.699 & 0.008 & 0.089 & 0.003 \\ 0.026 & 0.002 & 0.823 & 0.062 & 0.035 \\ 0.123 & 0.189 & 0.114 & 0.808 & 0.016 \\ 0.007 & 0.001 & 0.024 & 0.001 & 0.931 \end{bmatrix}.$$

74 Their HMMs are state-labelled. For each state i , they fit a Dirichlet pdf f_i describing
 75 the distribution of observations in Δ^3 emitted at state i . The pdfs of diseased and healthy
 76 individuals were so similar that they used the same pdf for both HMMs which was estimated
 77 from the whole population. Thus the two HMMs differ only in the transition probabilities.

78 The f_i are continuous and since in this paper we assume finite observation alphabets, we
 79 partition the simplex into the sets

$$80 \quad U_k = \{x \in \Delta^3 \mid f_k(x) \geq \sup_i f_i(x)\}$$

for $k = 1, \dots, 5$. The set U_k contains the points in Δ^3 most likely to be produced in state k . We assign a letter a_k for each U_k , and define a set of observations $\Sigma = \{a_1, \dots, a_5\}$. Thus, the probability of producing letter a_k from state i is given as $O_{i,k} = \int_{U_k} f_i(x)dx$. We estimated the entries of O using a numerical Monte Carlo technique. We generated 100,000 samples from all 5 Dirichlet distributions in their paper which yielded the estimate

$$O = \begin{pmatrix} 0.9172 & 0.0803 & 0 & 0.0002 & 0.0024 \\ 0.0719 & 0.8606 & 0 & 0.0665 & 0.0010 \\ 0 & 0.0007 & 0.8546 & 0.1055 & 0.0392 \\ 0.0008 & 0.0998 & 0.0663 & 0.8257 & 0.0075 \\ 0.0109 & 0.0094 & 0.1046 & 0.0334 & 0.8416 \end{pmatrix}.$$

Since we consider transition labelled HMMs, we define observation density functions Ψ_1, Ψ_2 by

$$\Psi_m(a_k)_{i,j} = (T_m)_{i,j} O_{i,k}$$

for $m = 1, 2$. Let $Q = [10]$. We construct the block diagonal HMM (Q, Σ, Ψ) where

$$\Psi(a) = \begin{pmatrix} \Psi_1(a) & 0 \\ 0 & \Psi_2(a) \end{pmatrix}$$

for each $a \in \Sigma$. Let π_1 and π_2 be the Dirac distributions on states 1 and 6 respectively. These initial distributions correspond to healthy and diseased individuals started from sleep state 1. Using the algorithm from [2] one can show that π_1 and π_2 are distinguishable.

We sampled runs of the HMM started from π_1 and π_2 and plotted the corresponding sequences of $\ln L_n$. We refer to each of these two plots as a *log-likelihood plot*; see Figure 1.

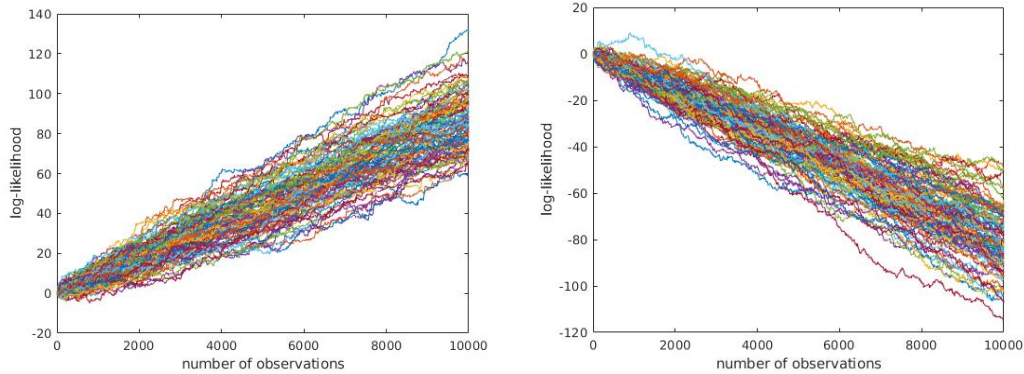


Figure 1 The two images show two log-likelihood plots of sample runs produced by π_1 and π_2 , respectively.

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By Theorem 5 it follows that $\ln L_n$ converges to ∞ and $-\infty$ in the case of words produced by π_1 and π_2 respectively. This is affirmed by Figure 1. Both log-likelihood plots also appear to follow a particular slope. This suggests that we can distinguish between words produced by π_1 and π_2 by tracking the value of $\ln L_n$ to see whether it crosses a lower or upper bound. This is the intuition behind the *Sequential Probability Ratio Test* (SPRT).

3 Sequential Probability Ratio Test

The SPRT runs as follows. We specify two error probabilities $\alpha, \beta > 0$ such that the SPRT gives the result π_2 from observations produced by π_1 with probability at most α and similarly the test gives the result π_1 from observations produced by π_2 with probability at most β . We then set $A = \ln \frac{\alpha}{1-\beta}$ and $B = \ln \frac{1-\alpha}{\beta}$. The algorithm continues to read observations and computes the value of $\ln L_n$ until $\ln L_n$ leaves the interval $[A, B]$. If $\ln L_n \leq A$ the test result is π_2 and if $\ln L_n \geq B$ the test result is π_1 . We may express SPRT as a random variable $\text{SPRT}_{\alpha, \beta} : \Sigma^\omega \rightarrow \{\pi_1, \pi_2\}$. We also define the stopping time

$$N_{\alpha, \beta} = \min\{n \in \mathbb{N} \mid \ln L_n \notin [A, B]\}.$$

We have that $N_{\alpha, \beta}$ is monotone decreasing in the sense that for $\alpha' \leq \alpha$ and $\beta' \leq \beta$ we have $N_{\alpha', \beta'} \geq N_{\alpha, \beta}$. When π_1 and π_2 are distinguishable $N_{\alpha, \beta}$ is \mathbb{P}_{π_2} -almost surely finite by Theorem 5.

► **Theorem 7.** *Suppose π_1 and π_2 are distinguishable. Let α, β be inputs to the SPRT. Then by choosing $A = \ln \frac{\alpha}{1-\beta}$ and $B = \ln \frac{1-\alpha}{\beta}$, $\mathbb{P}_{\pi_1}(\text{SPRT}_{\alpha, \beta} = \pi_2) \leq \alpha$ and $\mathbb{P}_{\pi_2}(\text{SPRT}_{\alpha, \beta} = \pi_1) \leq \beta$.*

In the following we consider the SPRT with respect to the measure \mathbb{P}_{π_2} . This is without loss of generality as there is a dual version of the SPRT, say $\overline{\text{SPRT}}$ with $\overline{L}_n = 1/L_n$ instead of L_n , such that $\overline{\text{SPRT}}_{\beta, \alpha} = \text{SPRT}_{\alpha, \beta}$.

3.1 Expectation of $N_{\alpha, \beta}$

Consider the two single state HMMs where $p_1 \neq p_2$.



(The Dirac distributions on) s_1 and s_2 are distinguishable. Further, the increments $\ln L_{n+1} - \ln L_n$ are i.i.d. and $0 > \mathbb{E}_{s_1}[\ln L_1 - \ln L_0] = p_2 \ln \frac{p_1}{p_2} + (1 - p_2) \ln \frac{1-p_1}{1-p_2} =: \ell$. Intuitively as ℓ gets more negative, the HMMs become more different.¹ Indeed, Wald [6] shows that the expected stopping time $\mathbb{E}_{s_2}[N_{\alpha, \beta}]$ and ℓ are inversely proportional:

$$\mathbb{E}_{s_2}[N_{\alpha, \beta}] = \frac{\beta \ln \frac{1-\alpha}{\beta} + (1-\beta) \ln \frac{\alpha}{1-\beta}}{\ell}. \quad (2)$$

This Wald formula cannot hold general (multi-state) HMMs. The increments $\ln L_{n+1} - \ln L_n$ need not be independent and $\mathbb{E}_{s_1}[\ln L_{n+1} - \ln L_n]$ can be different for different n . Further, $|\ln L_{n+1} - \ln L_n|$ can be unbounded –REF–.

Nevertheless, in Figure 1 we observed that $\ln L_n$ appears to decrease linearly (on the π_2 plot). Indeed we show, in Theorem 11 below, that the limit $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n$ exists \mathbb{P}_{π_2} -almost surely. Intuitively it corresponds to the average slope of the log-likelihood plot for π_2 . In the two state case, there is a simple proof of this using the law of large numbers:

$$\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n = \lim_{n \rightarrow \infty} \frac{1}{n} \sum_{i=0}^{n-1} \ln L_{i+1} - \ln L_i = \mathbb{E}_{\pi_2}[\ln L_1 - \ln L_0] = \ell \quad \mathbb{P}_{\pi_2}\text{-a.s.}$$

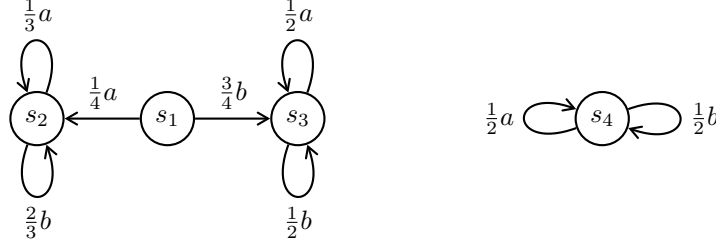
The number ℓ is called a likelihood exponent, as defined generally in the following definition.

¹ In fact, ℓ is the *KL-divergence* of the distributions f_1, f_2 where $f_i(a) = p_i$ and $f_i(b) = 1 - p_i$ for $i = 1, 2$.

137 ► **Definition 8.** A number ℓ is a likelihood exponent for initial distributions π_1 and π_2 if
 138 $\mathbb{P}_{\pi_2}(\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n = \ell) > 0$.

139 ► **Example 9.** In the case of Example 6 we show in the appendix that $\Lambda_{\pi_1, \pi_2} = \{\ell\}$. The
 140 slope of the right hand side of Figure 1 suggests that $\ell \approx -\frac{80}{10000} = -0.008$.

141 ► **Example 10.** Even for fixed π_1, π_2 there may be multiple likelihood exponents. Consider
 142 the following pair of HMMs.



143
 144 We observe two different likelihood exponents depending on the first letter produced. If
 145 the first letter is a then $L_1 = \frac{1}{2}$ and if the first letter is b then $L_1 = \frac{3}{2}$. If the first letter
 146 is a then $\ln L_{n+1} - \ln L_n$ are i.i.d. for $n > 1$ and $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n = \frac{1}{2} \ln \frac{2}{3} + \frac{1}{2} \ln \frac{4}{3} = \frac{1}{2} \ln \frac{8}{9}$
 147 like the two state example above. If the first letter is b then $L_n = \frac{3}{2}$ for all $n > 1$ and
 148 $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n = 0$.

149 We denote the set of likelihood exponents for initial distributions π_1 and π_2 by Λ_{π_1, π_2} .
 150 Further, we write $\Lambda = \bigcup_{\pi_1, \pi_2} \Lambda_{\pi_1, \pi_2}$. In addition, for $\ell \in \Lambda$ we define the event $E_\ell =$
 151 $\{\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n = \ell\}$. In Example 10, we have $\Lambda_{e_1, e_4} = \{\frac{1}{2} \ln \frac{8}{9}, 0\}$. Writing $\ell := \frac{1}{2} \ln \frac{8}{9}$ it
 152 follows that $\mathbb{P}_{s_4}(E_\ell) = \mathbb{P}_{s_4}(E_0) = \frac{1}{2}$.

153 ► **Theorem 11.** The set of likelihood exponents satisfies $\Lambda \subset [-\infty, 0]$ and $|\Lambda| \leq |Q|^2$. Further
 154 $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n$ exists (and by definition is in Λ) \mathbb{P}_{π_2} -almost surely for any π_1, π_2 .

155 We investigate how $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n$ influences the performance of the SPRT when α and
 156 β are small. Intuitively we expect a steeper slope in the likelihood plot (such as in Figure 1)
 157 to lead to faster termination. In the two state case, Wald's formula (2) becomes

$$158 \quad \mathbb{E}_{s_2}[N_{\alpha, \beta}] = \frac{\beta \ln \frac{1-\alpha}{\beta} + (1-\beta) \ln \frac{\alpha}{1-\beta}}{\ell} \sim \frac{\ln \alpha}{\ell} \quad (\text{as } \alpha, \beta \rightarrow 0), \quad (3)$$

159 where we use the asymptotic notation \sim defined in the preliminaries. In Theorem 12 below
 160 we generalise Equation (3) to arbitrary HMMs. Indeed a very similar asymptotic identity
 161 holds. In the case that $\Lambda = \{\ell\}$ and $\ell \in (-\infty, 0)$ then $\mathbb{E}_{s_2}[N_{\alpha, \beta}] \sim \frac{\ln \alpha}{\ell}$ as $\alpha, \beta \rightarrow 0$. If
 162 $|\Lambda| > 1$ then we condition our expectation on $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n$.

163 ► **Theorem 12.** [Generalised Wald Formula] Let ℓ be a likelihood exponent and let π_1 and
 164 π_2 be initial distributions. If $\ell \in (-\infty, 0)$ then

$$165 \quad \mathbb{E}_{\pi_2}[N_{\alpha, \beta} \mid E_\ell] \sim \frac{\ln \alpha}{\ell} \quad (\text{as } \alpha, \beta \rightarrow 0).$$

166 If $\ell = 0$ then there exist $\alpha, \beta > 0$ such that

$$167 \quad \mathbb{E}_{\pi_2}[N_{\alpha, \beta} \mid E_\ell] = \infty. \quad (4)$$

168 If $\ell = -\infty$ then

$$169 \quad \sup_{\alpha, \beta} \mathbb{E}_{\pi_2}[N_{\alpha, \beta} \mid E_\ell] < \infty.$$

The theorem above pertains to $\mathbb{E}_{\pi_2}[N_{\alpha,\beta}]$. In the next subsection we give additional information about the distribution of $N_{\alpha,\beta}$.

3.2 Distribution of $N_{\alpha,\beta}$

3.2.1 Likelihood Exponent 0

► **Example 13.** We continue with Example 10 to illustrate the second case in Theorem 12. By picking $\alpha = \frac{1}{4}, \beta = \frac{1}{4}$ the bounds for the SPRT are $A = \ln \frac{1}{3}$ and $B = \ln 3$. If the first letter is b then $\ln L_n = \ln \frac{3}{2}$ for all $n > 1$, thus never crosses the SPRT bounds and $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n = 0$. Hence with probability $\frac{1}{2}$ the SPRT fails to terminate and $N_{\alpha,\beta} = \infty$. It follows that $\mathbb{P}_{\pi_2}(E_0) = \frac{1}{2}$ and $E_{\pi_2}[N_{\alpha,\beta} \mid E_0] = \infty$. Also, $E_{\pi_2}[N_{\alpha,\beta}] = \infty$.

The second part of Theorem 12 says that the expectation of $N_{\alpha,\beta}$ conditioned under E_0 is infinite. The following proposition strengthens this statement. Conditioning under E_0 , the probability that $N_{\alpha,\beta}$ is infinite converges to 1 as the errors $\alpha, \beta \rightarrow 0$. Recall that $N_{\alpha,\beta}$ is monotone decreasing. It follows that $\{N_{\alpha',\beta'} = \infty\} \subset \{N_{\alpha,\beta} = \infty\}$ if $\alpha' \leq \alpha$ and $\beta' \leq \beta$.

► **Proposition 14.** The following equalities hold up to \mathbb{P}_{π_2} -null sets:

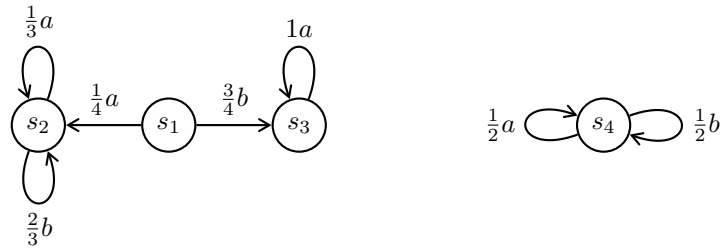
$$E_0 = \left\{ \lim_{n \rightarrow \infty} L_n > 0 \right\} = \bigcup_{\alpha,\beta} \{N_{\alpha,\beta} = \infty\}.$$

Thus, $\lim_{\alpha,\beta \rightarrow 0} \mathbb{P}_{\pi_2}(N_{\alpha,\beta} = \infty) = \mathbb{P}_{\pi_2}(E_0)$.

It follows with Theorem 5 that π_1 and π_2 are distinguishable if and only if $\mathbb{P}_{\pi_2}(E_0) = 0$ if and only if $\mathbb{P}_{\pi_2}(N_{\alpha,\beta} < \infty) = 1$ for all α, β .

3.2.2 Likelihood Exponent $-\infty$

► **Example 15.** Consider now a modification of Example 10 where state s_3 has the b loop removed.



The likelihood exponents are $-\infty$ and $\ell := \frac{1}{2} \ln \frac{8}{9}$ so that $\Lambda = \{-\infty, \ell\}$. Also, $\mathbb{P}_{s_4}(E_{-\infty}) = \mathbb{P}_{s_4}(E_\ell) = \frac{1}{2}$. Up to \mathbb{P}_{s_4} -null sets the events $E_{-\infty}$, $b\Sigma^\omega$ and $ba^*b\Sigma^\omega$ are equal. The event $ba^*b\Sigma^\omega$ represents the right chain producing an observation which the left chain cannot produce, causing the SPRT to terminate for any α, β . Therefore conditioned on $E_{-\infty}$, the random variable $N_{\alpha,\beta} - 1$ is bounded by a geometric random variable with parameter $\frac{1}{2}$. Hence $\sup_{\alpha,\beta} \mathbb{E}_{\pi_2} [N_{\alpha,\beta} \mid E_{-\infty}] \leq 1 + 2$.

We define the stopping time $N_\perp = \min\{n \in \mathbb{N} \mid L_n = 0\}$. Note that $\sup_{\alpha,\beta} N_{\alpha,\beta} \leq N_\perp$ since $\{L_n = 0\} \subseteq \{L_n \leq \frac{\alpha}{1-\beta}\}$ for all α, β . The following proposition states that the reverse inequality also holds.

► **Proposition 16.** We have $\sup_{\alpha,\beta} N_{\alpha,\beta} = N_{\perp}$. Thus, $\lim_{\alpha,\beta \rightarrow 0} \mathbb{P}_{\pi_2}(N_{\alpha,\beta} < \infty) = \mathbb{P}_{\pi_2}(E_{-\infty})$.

Applying this to Example 15, we have $\sup_{\alpha,\beta} \mathbb{E}_{\pi_2} [N_{\alpha,\beta} \mid E_{-\infty}] = 3$.

3.2.3 Likelihood Exponent in $(-\infty, 0)$

Conditioned on E_{ℓ} where $\ell \in (-\infty, 0)$, Theorem 12 shows that $N_{\alpha,\beta}$ scales with $\frac{\ln \alpha}{\ell}$ in expectation. The following result shows that this relationship also holds \mathbb{P}_{π_2} -almost surely.

► **Proposition 17.** Let $\ell \in \Lambda$ and assume $\ell \in (-\infty, 0)$. We have

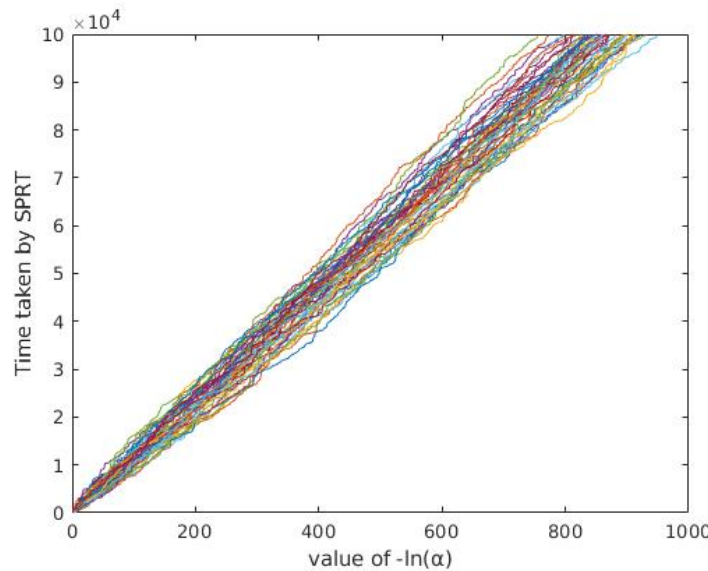
$$\mathbb{P}_{\pi_2} \left(N_{\alpha,\beta} \sim \frac{\ln \alpha}{\ell} \text{ (as } \alpha, \beta \rightarrow 0) \mid E_{\ell} \right) = 1.$$

In fact, we prove the first part of using Proposition 17. If there is a bound $M \in \mathbb{N}$ such that \mathbb{P}_{π_2} -almost surely $\frac{N_{\alpha,\beta}}{-\ln \alpha} \leq M$, the first part of Theorem 12 would follow from Proposition 17 by the Dominated convergence theorem. However this is not the case in general. Instead we show the following lemma and use Vitali's convergence theorem.

► **Lemma 18.** The set of random variables $\left\{ \frac{N_{\alpha,\beta}}{-\ln \alpha} \mid 0 < \alpha, \beta \leq \frac{1}{2} \right\}$ is uniformly integrable with respect to the measure \mathbb{P}_{π_2} ; i.e.

$$\lim_{K \rightarrow \infty} \sup_{\alpha,\beta} \mathbb{E}_{\pi_2} \left[-\frac{N_{\alpha,\beta}}{\ln \alpha} \mathbb{1}_{\frac{N_{\alpha,\beta}}{-\ln \alpha} \geq -K} \right] = 0.$$

► **Example 19.** Recall Example 6 concerning sleep cycle HMMs with $\Lambda = \{\ell\}$. Figure 2 demonstrates the asymptotic relationship in Proposition 17. Each of the 50 lines correspond to a sample run and we record the value of $N_{\alpha,\beta}$ for $0 \leq -\ln \alpha = -\ln \beta \leq 1000$. From the figure we estimate $-\frac{1}{\ell}$ as $\frac{10^5}{800} = 125$. This coincides with the estimate in Example 9.



■ **Figure 2** The time taken by the SPRT for $0 \leq -\ln \alpha = -\ln \beta \leq 1000$.

220 **Proof of Proposition 17.** Since $\Psi_{\min}^n \leq L_n \leq \Psi_{\min}^{-n}$ it follows that

$$221 \quad N_{\alpha,\beta} \geq \frac{\min\{\ln \frac{\alpha}{1-\beta}, \ln \frac{\beta}{1-\alpha}\}}{\ln \Psi_{\min}}$$

222 Hence $N_{\alpha,\beta} \rightarrow \infty$ \mathbb{P}_{π_2} -almost surely as $\alpha, \beta \rightarrow 0$. Consider the case $\ell_k \in (-\infty, 0)$. Let
 223 $U_{\alpha,\beta} = \{w \in \Sigma^\omega \mid \ln L_{N_\alpha} \leq \ln \frac{\alpha}{1-\beta}\}$. The set $\bigcap_{\alpha,\beta \in (0,1]} U_{\alpha,\beta}^c \subseteq \{L_n \text{ is unbounded}\}$. Hence,
 224 $\lim_{\alpha,\beta \rightarrow 0} \mathbb{1}_{U_{\alpha,\beta}} = 1$ \mathbb{P}_{π_2} -almost surely. Conditioned on V_k it follows that

$$225 \quad 0 \leq \mathbb{1}_{U_{\alpha,\beta}} \frac{\ln \frac{\alpha}{1-\beta} - \ln L_{N_{\alpha,\beta}}}{N_{\alpha,\beta}} \leq \mathbb{1}_{U_{\alpha,\beta}} \frac{\ln L_{N_{\alpha,\beta}-1} - \ln L_{N_{\alpha,\beta}}}{N_{\alpha,\beta}} \rightarrow 0 \text{ as } \alpha \rightarrow 0.$$

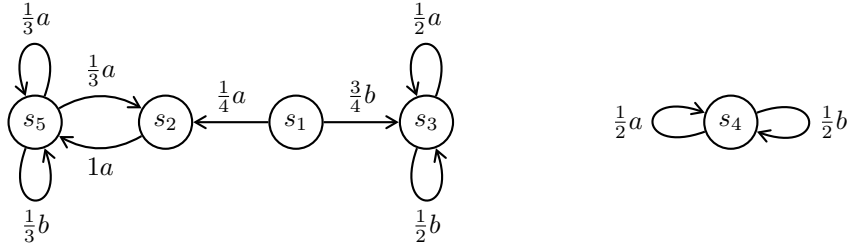
226 And so

$$227 \quad \lim_{\alpha,\beta \rightarrow 0} \frac{\ln \alpha}{N_{\alpha,\beta}} = \lim_{\alpha \rightarrow 0} \frac{\ln \frac{\alpha}{1-\beta}}{N_{\alpha,\beta}} = \lim_{\alpha \rightarrow 0} \frac{\ln L_{N_{\alpha,\beta}}}{N_{\alpha,\beta}} = \ell_k.$$

228 ◀

229 3.3 Qualitative Problems

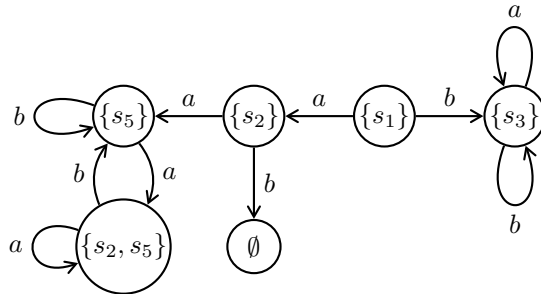
230 ► **Example 20.** Consider another adaption of Example 10.



231

232 If the first letter produced by s_4 is a b , then $L_n = \frac{3}{2}$ for all $n \in \mathbb{N}$. If the first two letters are
 233 ab then $L_1 = \frac{1}{2}$ and $L_n = 0$ for $n \geq 2$. If the first two letters are aa then $s_5 \in \text{supp } e_1 \Psi(aaw)$
 234 for all words $w \in \Sigma^*$. Therefore in this case, up to a \mathbb{P}_{s_4} -null set, $L_n > 0$ for all $n \in \mathbb{N}$ and
 235 we show in the appendix that there is $\ell \in (-\infty, 0)$ such that $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n = \ell$. Thus,
 236 $\Lambda_{s_1, s_4} = \{-\infty, \ell, 0\}$.

237 The likelihood ratio $L_n = 0$ if and only if $\text{supp } \pi_1 \Psi(w_n) = \emptyset$. In order to track the
 238 support of $\pi_1 \Psi(w)$, we consider the left part of the HMM as an NFA with s_1 as the initial
 239 state and its determinisation as shown in the DFA below.



240

241 Almost surely, s_4 produces a word that drives this DFA into a bottom SCC. This bottom
 242 SCC determines $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n$.

In general, s_4 is not a single state so the change of support of $\pi_1 \Psi(w_n)$ depends also on the state of the producing HMM started from π_2 . We construct a Markov chain that takes into account both $\text{supp } \pi_1 \Psi(w_n)$ and also the state of the producing chain.

Let $\mathcal{S} = \{(\mu, q) \mid \mu \in 2^Q, q \in Q\}$ be a set of states and let

$$P_{(\mu_1, s_1), (\mu_2, s_2)} = \sum_{a \in \Sigma} \Psi(a)_{i_1, i_2} \mathbb{1}_{\text{supp } \mu_1 \Psi(a) = \text{supp } \mu_2}$$

be a transition matrix. The pair (\mathcal{S}, P) is a Markov chain –REQUIRES PROOF?. We then define the initial distribution

$$(u_{\pi_1, \pi_2})_{(\mu, q)} = \begin{cases} (\pi_2)_q & \text{supp } \mu = \text{supp } \pi_1 \\ 0 & \text{else.} \end{cases}$$

► **Lemma 21.** Consider the set $\{(\vec{0}, q) \mid q \in Q\}$

► **Lemma 22.** Let $C \subset \mathcal{S}$ be a bottom SCC of (\mathcal{S}, P) . there exists an $\ell \in (-\infty, 0]$ such that for all π_1 and for all states $q \in Q$ such that $(\text{supp } \pi_1, q) \in C$ we have $\Lambda_{\pi_1, \delta_q} = \{\ell\}$.

► **Lemma 23.** Let $C \subset \mathcal{S}$ be a bottom SCC of (\mathcal{S}, P) . For any $(\mu, q) \in C$, there is a partition of μ written $\mathcal{U}(\mu, q)$ such that the mapping $\phi : \mathcal{U}(\mu, q) \rightarrow \mathcal{P}(\mathcal{G})$ given by

$$\phi(\rho) = \{(s, q) \mid s \in \rho\}$$

is injective and for all $\rho \in \mathcal{U}(\mu, q)$, $\phi(\rho) \subseteq C$ for some right bottom SCC of \mathcal{G} .
the set of likelihood exponents $\Lambda_{\mu/\|\mu\|, \delta_q}$ contains the single element

$$\sup \{\ell \in \Lambda_{\rho, \delta_q} \mid (p, q) \text{ in right bottom SCC of } G \text{ and } \text{supp } \rho \subseteq \text{supp } \mu\}.$$

Proof. – Prove the set is non-empty and $\ell > -\infty$

Let the set be maximised on (p, q) . Then, there is $(\rho, q) \in \mathcal{S}$ such that $p \in \text{supp } \rho$.

We have

$$\begin{aligned} -\infty &< \lim_{n \rightarrow \infty} \frac{1}{n} \ln \|\rho \Psi(w)\| \\ &= \lim_{n \rightarrow \infty} \frac{1}{n} \ln \sum_{s \in \text{supp } \rho} \|\delta_s \Psi(w)\| \\ &= \sup_{s \in \text{supp } \rho} \lim_{n \rightarrow \infty} \frac{1}{n} \ln \|\delta_s \Psi(w)\| \\ &\leq \lim_{n \rightarrow \infty} \frac{1}{n} \ln \|\rho \Psi(w)\|. \end{aligned}$$

We have $\sup_{s \in \text{supp } \rho} \lim_{n \rightarrow \infty} \frac{1}{n} \ln \|\delta_s \Psi(w)\| = \sup_{s \in \text{supp } \rho} \Lambda_{\delta_s, \delta_q}$

► **Lemma 24.** There is a mapping $\phi : \mathcal{S} \rightarrow \Lambda$ such that for any $\ell \in \Lambda$, the events $\{\text{hitting } \phi^{-1}(\ell)\}$ and $\{\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n = \ell\}$ are equal.

– each bottom SCC is associated with a likelihood exponent. We can find those bottom SCCs that are associated with $-\infty$ and the zero SCCs.

Let (p_2, p_5) be the state distributions on s_2 and s_5 respectively. We have $\|(p_2, p_5) \Psi(a)\| = \|(\frac{1}{3}p_5 + p_2, \frac{1}{3}p_5)\| \leq \|(p_2, p_5)\|$ and $\|(p_2, p_5) \Psi(b)\| = \|(\frac{1}{3}p_5, 0)\| \leq \frac{1}{3} \|(p_2, p_5)\|$.

► **Theorem 25.** Given an HMM and initial distributions π_1, π_2 ,

- 277 1. one can compute $\mathbb{P}_{\pi_2}(E_{-\infty})$ and $\mathbb{P}_{\pi_2}(E_0)$ in PSPACE;
- 278 2. one can decide whether $\mathbb{P}_{\pi_2}(E_0) = 0$ (i.e., $0 \notin \Lambda_{\pi_1, \pi_2}$) in polynomial time;
- 279 3. deciding whether $\mathbb{P}_{\pi_2}(E_0) = 1$, whether $\mathbb{P}_{\pi_2}(E_{-\infty}) = 0$, and whether $\mathbb{P}_{\pi_2}(E_{-\infty}) = 1$ are
- 280 all PSPACE-complete problems.

281 the *mortality* problem, which asks, given a finite set of states Q , a finite alphabet Σ , and
 282 a function $\Phi : \Sigma \rightarrow \{0, 1\}^{Q \times Q}$, whether there exists a word $w \in \Sigma^*$ such that $\Phi(w)$ is the
 283 zero matrix. The mortality problem can be viewed as a special case of the NFA universality
 284 problem (given an NFA, does it accept every word?). Like NFA universality, the mortality
 285 problem is PSPACE-complete [CITE SHALLIT].

286 Concerning $\mathbb{P}_{\pi_2}(E_{-\infty})$ (cf. Theorem 25.3), we actually show a stronger result, namely
 287 that any nontrivial approximation of $\mathbb{P}_{\pi_2}(E_{-\infty})$ is PSPACE-hard. The proof is also based
 288 on the mortality problem.

289 ► **Proposition 26.** *There is a polynomial-time computable function that maps any instance*
 290 *of the mortality problem to an HMM and initial distributions π_1, π_2 so that if the instance is*
 291 *positive then $\mathbb{P}_{\pi_2}(E_{-\infty}) = 1$ and if the instance is negative then $\mathbb{P}_{\pi_2}(E_{-\infty}) = 0$. Thus, any*
 292 *nontrivial approximation of $\mathbb{P}_{\pi_2}(E_{-\infty})$ is PSPACE-hard.*

293 **Proof sketch.** Given an instance Φ of the mortality problem, construct an HMM (Q, Σ, Ψ)
 294 so that Φ and Ψ have the same graph [DEFINED?] and $\sum_{a \in \Sigma} \Psi(a)$ is stochastic. Define
 295 π_1 as a uniform distribution on Q . Define π_2 as a Dirac distribution on a fresh state that
 296 emits letters from Σ uniformly at random. Thus, if Φ is a positive instance of the mortality
 297 problem then $\mathbb{P}_{\pi_2}(E_{-\infty}) = 1$, and if Φ is a negative instance then $\mathbb{P}_{\pi_2}(E_{-\infty}) = 0$. ◀

298 4 Stefan

299 4.1 Representing Likelihood Exponents

300 In the following we show that, although it is hard to *compute* likelihood exponents, one
 301 can efficiently represent them in terms of *Lyapunov exponents*. The definition of Lyapunov
 302 exponents is based on the following definition.

303 ► **Definition 27.** *A matrix system is a triple $\mathcal{M} = (Q, \Sigma, \Psi)$ where Q is a finite set of*
 304 *states, Σ is a finite set of observations, and $\Psi : \Sigma \rightarrow \mathbb{R}_{\geq 0}^{Q \times Q}$ specifies the transitions.*
 305 *(Note that an HMM is a matrix system.) A Lyapunov system is a pair $\mathcal{S} = (\mathcal{M}, \rho)$ where*
 306 *$\mathcal{M} = (Q, \Sigma, \Psi)$ is a matrix system and $\rho : \Sigma \rightarrow (0, 1]$ is a probability distribution on Σ*
 307 *with full support [PROBABLY INTRODUCE A NOTATION FOR A PROBABILITY*
 308 *DISTRIBUTION IN THE PRELIMINARIES.], such that the directed graph (Q, E) with*
 309 *$E = \{(q, r) \mid \sum_{a \in \Sigma} \Psi_{q,r}(a) > 0\}$ is strongly connected.*

310 We can identify the probability distribution ρ from this definition with the single-state
 311 HMM $(\{s\}, \Sigma, \Psi_\rho)$ where $\Psi_\rho(a)_{s,s} = \rho(a)$ for all $a \in \Sigma$. In this way, ρ produces a random
 312 infinite word from Σ^ω . For $w \in \Sigma^\omega$ we write w_n for the length- n prefix of w . The following
 313 lemma is known from [CITE].

314 ► **Lemma 28** ([CITE]). *Let $((Q, \Sigma, \Psi), \rho)$ be a Lyapunov system. Then there is $\ell \in \mathbb{R}$ such*
 315 *that, for all $q \in Q$, \mathbb{P}_ρ -a.s., either $e_q \Psi(w_n) = \vec{0}$ [NOTATION DEFINED?] for some $n \in \mathbb{N}$*
 316 *or the limit $\lim_{n \rightarrow \infty} \frac{1}{n} \ln \|e_q \Psi(w_n)\|$ exists and equals ℓ .*

317 For a Lyapunov system \mathcal{S} we call $\ell(\mathcal{S}) = \ell$ from the lemma the *Lyapunov exponent* defined
 318 by \mathcal{S} . We prove the following theorem, which implies Theorem 11. [CHECK $|Q|^2$ vs. $|Q|^2 + 1$]

► **Theorem 29.** *Given an HMM (Q, Σ, Ψ) we can compute in polynomial time $2K \leq 2|Q|^2$ Lyapunov systems $\mathcal{S}_1^1, \mathcal{S}_1^2, \mathcal{S}_2^1, \mathcal{S}_2^2, \dots, \mathcal{S}_K^1, \mathcal{S}_K^2$ such that for any initial distributions π_1, π_2 the limit $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n$ exists \mathbb{P}_{π_2} -a.s. and lies in*

$$\Lambda \subseteq \{-\infty\} \cup \{\ell(\mathcal{S}_1^1) - \ell(\mathcal{S}_1^2), \dots, \ell(\mathcal{S}_K^1) - \ell(\mathcal{S}_K^2)\}.$$

In particular, the HMM has at most $|Q|^2 + 1$ likelihood exponents.

In the rest of the section we provide more details on the construction underlying Theorem 29. As an intermediate concept (between the given HMM and the Lyapunov systems from Theorem 29) we define *generalized Lyapunov systems*.

First, for two matrix systems $\mathcal{M}_1 = (Q_1, \Sigma, \Psi_1)$ and $\mathcal{M}_2 = (Q_2, \Sigma, \Psi_2)$ with finite Q_1, Q_2, Σ and transitions $\Psi_1, \Psi_2 : \Sigma \rightarrow \mathbb{R}_{\geq 0}^{Q_1 \times Q_2}$ we define the directed graph $G_{\mathcal{M}_1, \mathcal{M}_2} = (Q_1 \times Q_2, E)$ such that there is an edge from (q_1, q_2) to (r_1, r_2) if there is $a \in \Sigma$ with $\Psi_1(a)_{q_1, r_1} > 0$ and $\Psi_2(a)_{q_2, r_2} > 0$.

A *generalized Lyapunov system* is a triple $\mathcal{S} = (\mathcal{M}, \mathcal{H}, C)$ where $\mathcal{M} = (Q_1, \Sigma, \Psi_1)$ is matrix system and $\mathcal{H} = (Q_2, \Sigma, \Psi_2)$ is a strongly connected [DEFINED?] HMM and $C \subseteq Q_1 \times Q_2$ is a bottom SCC of $G_{\mathcal{M}, \mathcal{H}}$. Given a generalized Lyapunov system, one can efficiently compute an “equivalent” Lyapunov system:

► **Lemma 30.** *Let $\mathcal{S} = ((Q_1, \Sigma, \Psi_1), (Q_2, \Sigma, \Psi_2), C)$ be a generalized Lyapunov system.*

1. *There is $\ell(\mathcal{S}) \in \mathbb{R}$ such that, for all $\pi_1 \in \mathbb{R}_{\geq 0}^{Q_1}$ and all probability distributions [NOTATION?] $\pi_2 \in [0, 1]^{Q_2}$ with $\{q \mid \pi_1(q) > 0\} \times \text{supp}(\pi_2) \subseteq C$, we have \mathbb{P}_{π_2} -a.s. that either $\pi_1 \Psi_1(w_n) = \vec{0}$ [NOTATION?] for some $n \in \mathbb{N}$ or the limit $\lim_{n \rightarrow \infty} \frac{1}{n} \ln \|\pi_1 \Psi_1(w_n)\|$ exists and equals $\ell(\mathcal{S})$.*

2. *One can compute in polynomial time a Lyapunov system \mathcal{S}' such that $\ell(\mathcal{S}) = \ell(\mathcal{S}')$.*

Let $\mathcal{H} = (Q, \Sigma, \Psi)$ be an HMM. Let $R \subseteq Q \times Q$ be a (not necessarily bottom) SCC of the graph $G_{\mathcal{H}, \mathcal{H}}$ such that $Q_2 := \{q_2 \in Q \mid \exists q_1 \in Q : (q_1, q_2) \in R\}$ is a bottom SCC of the graph of $\sum_{a \in \Sigma} \Psi(a)$. We call such R a *right-bottom* SCC. Clearly there are at most $|Q|^2$ right-bottom SCCs. Towards Theorem 29 we want to define, for each right-bottom SCC R , two generalized Lyapunov systems $\mathcal{S}_R^1, \mathcal{S}_R^2$.

For a function of the form $\Phi : \Sigma \rightarrow \mathbb{R}^{Q \times Q}$ and $P \subseteq Q$ we write $\Phi|_P : \Sigma \rightarrow \mathbb{R}^{P \times P}$ for the function with $\Phi|_P(a)(q, r) = \Phi(a)(q, r)$ for all $a \in \Sigma$ and $q, r \in P$; i.e., $\Phi|_P(a)$ denotes the principal submatrix obtained from $\Phi(a)$ by restricting it to the rows and columns indexed by P .

Define $\Psi'(a, r)_{q, r} := \Psi(a)_{q, r}$ for all $a \in \Sigma$ and $q, r \in Q$. Then $(Q, \Sigma \times Q, \Psi')$ is an HMM, which is similar to \mathcal{H} , but which emits, in addition to an observation from Σ , also the next state. Since Q_2 is a bottom SCC of the graph of $\sum_{a \in \Sigma} \Psi(a)$, the HMM $\mathcal{H}_2 := (Q_2, \Sigma \times Q_2, \Psi'|_{Q_2})$ is strongly connected. This HMM \mathcal{H}_2 will be used both in \mathcal{S}_R^1 and in \mathcal{S}_R^2 .

Next, define $\bar{\Psi} : (\Sigma \times Q) \rightarrow [0, 1]^{(Q \times Q) \times (Q \times Q)}$ by

$$\bar{\Psi}(a, r_2)_{(q_1, q_2), (r_1, r_2)} := \Psi(a)_{q_1, r_1} \quad \text{for all } a \in \Sigma \text{ and } q_1, q_2, r_1, r_2 \in Q.$$

Now define $\mathcal{S}_R^1 := (\mathcal{M}^1, \mathcal{H}_2, C^1)$, where $\mathcal{M}^1 := (R, \Sigma \times Q_2, \bar{\Psi}|_R)$ and $C^1 := \{((q_1, q_2), q_2) \mid (q_1, q_2) \in R\}$. Finally, denoting by $R' \subseteq Q_2 \times Q_2$ the SCC of the graph $G_{\mathcal{H}, \mathcal{H}}$ that contains the “diagonal” vertices $(q, q) \in Q_2 \times Q_2$, define $\mathcal{S}_R^2 := (\mathcal{M}^2, \mathcal{H}_2, C^2)$, where $\mathcal{M}^2 := (R', \Sigma \times Q_2, \bar{\Psi}|_{R'})$ and $C^2 := \{((q_1, q_2), q_2) \mid (q_1, q_2) \in R'\}$.

For sets $U, V \subseteq Q \times Q$ write $U \xrightarrow{*}_{G_{\mathcal{H}, \mathcal{H}}} V$ denote that there are $u \in U$ and $v \in V$ such that v is reachable from u in $G_{\mathcal{H}, \mathcal{H}}$. We are ready to state the following key technical lemma:

► **Lemma 31.** *Given an HMM (Q, Σ, Ψ) , let $\mathcal{R} \subseteq 2^{Q \times Q}$ be the set of its right-bottom SCCs, and, for $R \in \mathcal{R}$, let $\mathcal{S}_R^1, \mathcal{S}_R^2$ be the generalized Lyapunov systems defined above. Then,*

$$\Lambda \subseteq \{-\infty\} \cup \{\ell(\mathcal{S}_R^1) - \ell(\mathcal{S}_R^2) \mid R \in \mathcal{R}\}.$$

NEW AND STRONGER VERSION:

► **Lemma 32.** *Given an HMM (Q, Σ, Ψ) , let $\mathcal{R} \subseteq 2^{Q \times Q}$ be the set of its right-bottom SCCs, and, for $R \in \mathcal{R}$, let $\mathcal{S}_R^1, \mathcal{S}_R^2$ be the generalized Lyapunov systems defined above. Then, for any initial distributions π_1, π_2 , the limit $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n$ exists \mathbb{P}_{π_2} -a.s. and lies in*

$$\{-\infty\} \cup \{\ell(\mathcal{S}_R^1) - \ell(\mathcal{S}_R^2) \mid R \in \mathcal{R}, \text{supp}(\pi_1) \times \text{supp}(\pi_2) \xrightarrow{*}_{G_{\mathcal{H}, \mathcal{H}}} R\}.$$

Thus, $\Lambda_{\pi_1, \pi_2} \subseteq \{-\infty\} \cup \{\ell(\mathcal{S}_R^1) - \ell(\mathcal{S}_R^2) \mid R \in \mathcal{R}, \text{supp}(\pi_1) \times \text{supp}(\pi_2) \xrightarrow{*}_{G_{\mathcal{H}, \mathcal{H}}} R\}.$

Proof sketch. Let π_1, π_2 be initial distributions. Very loosely speaking, we show in the appendix that on \mathbb{P}_{π_2} -almost every run w there is a right-bottom SCC R which “traps” “most” of the mass of $\pi_1 \Psi(w_n)$ and $\pi_2 \Psi(w_n)$. This can be made meaningful and formal using (the cross-product systems) $\mathcal{S}_R^1, \mathcal{S}_R^2$. We then show that on \mathbb{P}_{π_2} -almost every such run w , for both $i = 1, 2$, the limit $\lim_{n \rightarrow \infty} \frac{1}{n} \ln \|\pi_i \Psi(w_n)\|$ exists and equals $\ell(\mathcal{S}_R^i)$ (or $\pi_i \Psi(w_n) = \vec{0}$ for some n). It follows that

$$\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n = \lim_{n \rightarrow \infty} \frac{1}{n} \ln \frac{\|\pi_1 \Psi(w_n)\|}{\|\pi_2 \Psi(w_n)\|} = \ell(\mathcal{S}_R^1) - \ell(\mathcal{S}_R^2). \quad \blacktriangleleft$$

With Lemma 32 at hand, the proof of Theorem 29 is easy:

Proof of Theorem 29. As argued before, the set \mathcal{R} of right-bottom SCCs of the given HMM has at most $|Q|^2$ elements. These right-bottom SCCs R and the associated generalized Lyapunov systems $\mathcal{S}_R^1, \mathcal{S}_R^2$ can be computed in polynomial time. By Lemma 32 we have $\Lambda = \bigcup_{\pi_1, \pi_2} \Lambda_{\pi_1, \pi_2} \subseteq \{-\infty\} \cup \{\ell(\mathcal{S}_R^1) - \ell(\mathcal{S}_R^2) \mid R \in \mathcal{R}\}$. By Lemma 30.2, for each $R \in \mathcal{R}$ one can compute in polynomial time an equivalent Lyapunov system. \blacktriangleleft

5 Computational Aspects

– We may define a mapping on this Markov chain onto a set of at most $|Q|^2$ HMMs such that

- In the case of deterministic HMM, this exponential markov chain has at most $|Q|^2$ reachable states.

► **Lemma 33.** *Let (Q, Σ, Ψ) be a strongly connected HMM with initial distribution π_1 and let $(\{q\}, \Sigma, \rho)$ be an HMM with a single state with initial distribution e_q . Then the likelihood exponent $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n$ exists and \mathbb{P}_{e_q} -almost surely.*

$$\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n = \lim_{n \rightarrow \infty} \frac{1}{n} \ln \frac{\|\pi_1 \times \pi_2 \bar{\Psi}\|}{\|\pi_2 \times \pi_2 \bar{\Psi}\|} = \lim_{n \rightarrow \infty} \frac{1}{n} \ln \frac{\|\pi_1 \times \pi_2 \bar{\Psi}\|}{\|e_q \rho\|} - \lim_{n \rightarrow \infty} \frac{1}{n} \ln \frac{\|\pi_2 \times \pi_2 \bar{\Psi}\|}{\|e_q \rho\|} \quad (5)$$

For the general HMMs we simulate the joint state distribution using a new HMM $(\bar{Q}, \bar{\Sigma}, \bar{\Psi})$. This HMM has state space $\bar{Q} = (Q \cup \{\perp\}) \times Q$ where the states (\perp, q) are fresh states for each $q \in Q$. The set of observations is $\bar{\Sigma} = \Sigma$. Let $a \in \Sigma$ and let $l_1, l_2 \in Q \cup \{\perp\}$ and $r_1, r_2 \in Q$. We define

$$\bar{\Psi}(a, r)_{(l_1, r_1), (l_2, r_2)} = \begin{cases} \Psi(a)_{l_1, l_2} & l_1, l_2 \in Q, r = r_2, \Psi(a)_{r_1, r_2} > 0 \\ 1 & l_1 \in Q, l_2 = \perp, r = r_2, \sum_{s \in Q} \Psi(a)_{l_1, s} = 0 \\ 1 & l_1 = l_2 = \perp, r = r_2 \\ 0 & \text{else.} \end{cases}$$

398 ► **Lemma 34.** *The triple $(\overline{Q}, \overline{\Sigma}, \overline{\Psi})$ is an HMM and $\|\pi_1 \times \pi_2 \overline{\Psi}((a_1, r_1), \dots, (a_n, r_n))\| =$
 399 $\|\pi_1 \Psi(a_1 \dots a_n)\|$*

400 **Proof.** In the case that $l_1 \in Q$ and $\sum_{s \in Q} \Psi(a)_{l_1, s} > 0$ we have

$$\begin{aligned}
 401 \quad \sum_{a \in \Sigma} \sum_{r \in Q} \sum_{l_2 \in Q \cup \{\perp\}} \sum_{r_2 \in Q} \overline{\Psi}(a, r)_{(l_1, r_1), (l_2, r_2)} &= \sum_{a \in \Sigma} \sum_{r \in Q} \sum_{l_2 \in Q} \sum_{r_2 \in Q} \delta_{r=r_1} \Psi(a)_{l_1, l_2} \sum_{b \in \Sigma} \Psi(b)_{r_1, r_2} \\
 402 &= \sum_{a \in \Sigma} \sum_{l_2 \in Q} \Psi(a)_{l_1, l_2} \sum_{b \in \Sigma} \sum_{r_2 \in Q} \Psi(b)_{r_1, r_2} \\
 403 &= 1.
 \end{aligned}$$

404
 405
 406 In the case that $l_1 \in Q$ and $\sum_{s \in Q} \Psi(a)_{l_1, s} = 0$ or in the case that $l_1 = \perp$,

$$\begin{aligned}
 407 \quad \sum_{a \in \Sigma} \sum_{r \in Q} \sum_{l_2 \in Q \cup \{\perp\}} \sum_{r_2 \in Q} \overline{\Psi}(a, r)_{(l_1, r_1), (l_2, r_2)} &= \sum_{a \in \Sigma} \sum_{r \in Q} \sum_{r_2 \in Q} \delta_{r=r_2} \Psi(a)_{r_1, r_2} \\
 408 &= \sum_{a \in \Sigma} \sum_{r_2 \in Q} \Psi(a)_{r_1, r_2} \\
 409 &= 1.
 \end{aligned}$$

411 Now consider a word $(a_1, r_2), (a_2, r_3), \dots, (a_n, r_{n+1}) \in (\Sigma \times Q)^n$. We have ◀

412 ► **Lemma 35.**

413 Write $O : (Q \cup \{\perp\}) \times Q \rightarrow Q$ for the function $O(l, r) = r$. Let \mathcal{R} be the set of strongly con-
 414 nected components of $(\overline{Q}, \overline{\Sigma}, \overline{\Psi})$. The image $O(\mathcal{R})$ is a set of strongly connected components of
 415 (Q, Σ, Ψ) . We define $\mathcal{R}' = \{R \in \mathcal{R} \mid O(R) \text{ is a bottom connected component of } (Q, \Sigma, \Psi)\}$.

416 For $R \in \mathcal{R}'$ we may define an associated HMM (R, Σ', Ψ') where $\Sigma' = \Sigma \times Q \cup \{\$ \}$ and

$$417 \quad \Psi'(a)_{(l_1, r_1), (l_2, r_2)} = \{ \tag{6}$$

418 ► **Proposition 36.** *Whether $0 \in \Lambda$ can be decided in polynomial time. If $0 \in \Lambda$ then*
 419 *$\sup_{\alpha, \beta} \mathbb{P}_{\pi_2}(N_{\alpha, \beta} = \infty) = \mathbb{P}_{\pi_2}(E_0)$ can be computed in PSPACE. It is PSPACE-complete to*
 420 *decide whether $\mathbb{P}_{\pi_2}(E_0) = 1$.*

$$421 \quad \pi_1 \times \pi_2 = \pi_1^T \pi_2 \text{ we have } \|\pi_1 \times \pi_2 \overline{\Psi}(a, r_2)\| = \|\pi_1 \Psi(a)\|$$

422 ► **Lemma 37.** *It is PSPACE-hard to decide whether $\mathbb{P}_{\pi_2}(E_0) = 1$.*

423 **Proof.** We reduce from the universality decision problem in NFAs. Let $\mathcal{N} = (Q, \Sigma, \delta, q_0, F)$
 424 be an NFA. Define the HMM (Q', Σ', Ψ) where $Q' = Q \cup \{q_\$, q_1\}$ and $\Sigma' = \Sigma \cup \{\$ \}$. Then,
 425 we may define Ψ as

$$426 \quad \Psi(a)_{i, j} = \begin{cases} \delta(a)_{i, j} / \sum_{k=1}^{|Q|} \delta(a)_{i, k} & a \in \Sigma, i \in Q/F, j \in Q \\ \delta(a)_{i, j} / (1 + \sum_{k=1}^{|Q|} \delta(a)_{i, k}) & a \in \Sigma', i \in F, j \in Q \\ 1 / (1 + \sum_{k=1}^{|Q|} \delta(a)_{i, k}) & a \in \Sigma', i \in F, j = s_\$ \\ 1 / |\Sigma'| & a \in \Sigma, i = j = q_1 \\ 1 / |\Sigma'| & a = \$, i = q_1, j = q_\$ \\ 1 & i = j = q_\$ \\ 0 & \text{else.} \end{cases}$$

Next, let $\pi_1 = e_{q_0}$ and $\pi_2 = e_{q_1}$ be initial distributions. If every word is accepted by \mathcal{N} then for every $w \in (\Sigma)^*$, the event $w\$ \Sigma^\omega = E_0$ since $\lim_{n \rightarrow \infty} L_n > 0$. But up to a \mathbb{P}_{π_2} -null set, $\Sigma^*(\Sigma')^\omega = \Sigma^*(\Sigma')^\omega$ which is the set of words that can be generated starting the HMM at π_2 . it follows that $\mathbb{P}_{\pi_2}(E_0) = 1$. If there is a word $w \in (\Sigma)^*$ which is not final, then the word $w\$$ can be generated with non-zero probability by π_2 but is a mortal word for π_1 . Hence $\mathbb{P}_{\pi_2}(E_0) < 1$. \blacktriangleleft

► **Proposition 38.** *Deciding whether $-\infty \in \Lambda$ is PSPACE-complete. If $-\infty \in \Lambda$ then $\mathbb{E}_{\pi_2}[N_\perp | E_{-\infty}] < \infty$ and both $\mathbb{P}_{\pi_2}(E_{-\infty})$ and $\mathbb{E}_{\pi_2}[N_\perp | E_{-\infty}]$ are computable in PSPACE. It is PSPACE-complete to decide whether $\mathbb{P}_{\pi_2}(E_{-\infty}) = 1$.*

Proof. It follows that $\mathbb{P}_{\pi_2}(N = n) = u_{\pi_1, \pi_2} P^n \chi_K$ since

$$\mathbb{P}_{\pi_2}(N = n) = \mathbb{P}_{\pi_2}(L_n(w) = 0) \iff \mathbb{P}_{\pi_2}(\|\pi_1 \Psi(w)\| = 0) \iff \mathbb{P}_{\pi_2}(\text{supp}(\pi_1 \Psi(w)) = \mathbf{0}).$$

Therefore $\mathbb{E}_{\pi_2}[N]$ is equal to the expected hitting time of K in the Markov chain (\mathcal{S}, P) which has $|Q| \times 2^{|Q|}$ states.

– system can be computed by a PSPACE transducer. – Kemeny, Snell book citation \blacktriangleleft

► **Lemma 39.** *Deciding whether $-\infty \in \Lambda$ is PSPACE-hard.*

Proof. Let the matrices $M_1, \dots, M_D \in \{0, 1\}^{N \times N}$ be an instance of the mortality problem. If the i th row of $\sum_{d=1}^D M_d$ is a zero row, we may define M'_1, \dots, M'_D as the set of matrices M_1, \dots, M_D but with the i th row and column removed. M'_1, \dots, M'_D is mortal if and only if M_1, \dots, M_D is mortal.

– if there is remaining state after the product M' then add an additional letter.

Therefore, without loss of generality, we may assume $\sum_{d=1}^D M_d$ contains no zero rows. The matrices $\overline{M}_1, \dots, \overline{M}_D$ defined by $\overline{M}_{i,j} = M_{i,j} / \sum_{k=1}^N M_{i,k}$ are mortal if and only if M_1, \dots, M_D are mortal. Finally, define the HMM (Q, Σ, Ψ) where $Q = [N + 1]$, $\Sigma = \{a_1, \dots, a_d\}$ and Ψ is defined in blocks

$$\Psi(a_i) = \begin{pmatrix} \overline{M}_i & 0 \\ 0 & 1/|\Sigma| \end{pmatrix}.$$

Let $\pi_1 = (1/N, \dots, 1/N, 0)$ and $\pi_2 = (0, \dots, 0, 1)$. These initial distributions lead to a set of likelihood exponents Λ and It follows that $-\infty \in \Lambda$ if and only if M_1, \dots, M_d are mortal. \blacktriangleleft

Define an directed graph as follows. The vertex set is $V = (Q \cup \{\perp\}) \times Q$ where \perp is a fresh state. There is an edge from $(q_1, q_2) \in V$ to $(r_1, r_2) \in V$ if any of the following three conditions hold.

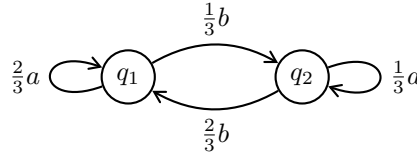
We say an HMM (Q, Σ, Ψ) is strongly connected if the Markov chain $\sum_{a \in \Sigma} \Psi(a)$ is strongly connected.

► **Theorem 40.** *We may compute in polynomial time a set of at most $|Q|^2$ strongly connected HMMs (Q_k, Σ, Ψ_k) and pairs of initial states (q_k, r_k) with corresponding likelihood ratios L_n^k such that for all $\ell \in \Lambda$ there is L_n^k such that $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n^k = \ell$.*

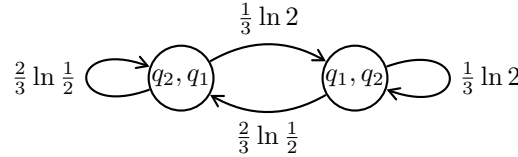
Markov chain has states labelled (q_i, q_j) for $i, j = 1, \dots, |Q|$ and also a special absorbing state \perp . The transition probabilities are taken from the outgoing transitions from q_j . It may be that a letter produced from q_j cannot be produced from q_i in which case the Markov chain transitions to \perp . The full definition is

5.1 Deterministic Chains

A HMM (Q, Σ, Ψ) is deterministic if for all $a \in \Sigma$, all rows of $\Psi(a)$ contain exactly one non-zero entry. If q is a starting state, for any word in $w \in \Sigma$ the vector $\delta_q \Psi(w)$ contains at most one non-zero entry. A simple example of a deterministic chain is the following example.



Suppose two copies of the chain are started from states q_1 and q_2 respectively, if the chains emit the same word, they will be in opposite states. Let the chain started from q_1 emit word $aab \in \Sigma^3$. The chain will now be in state q_2 . Suppose the next letter emitted is b . The value of $L_4 = L_3 \times \frac{2}{3}/\frac{1}{3} = 2L_3$ and $\ln L_4 = \ln L_3 + \ln 2$. In general deterministic chains have the property that $\ln L_n = R_n + \ln L_{n-1}$ where R_n is a random variable depending only the transition chosen. In the example above, we can build a Markov chain to characterise the joint distribution of states. We are interested in $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n$ with words produced from starting state q_2 so the transition probabilities are lifted from the second chain. In the diagram below we also include the value of R_n on each transition.



The Markov chain has stationary distribution $(\frac{2}{3}, \frac{1}{3})$. By the strong ergodic theorem for Markov chains, $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n = \frac{1}{3} \ln 2 + \frac{2}{3} \ln \frac{1}{2} = -\frac{1}{3} \ln 2$.

In general there may be a finite number of possible limits of $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n$. Moreover, these limits may not be rational numbers but expressible as linear combinations of natural logarithms of rational numbers. We say a possible limit of $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n$ is ln-symbolically computable if we can compute a set of $2M$ rational numbers where $M \leq |Q|^2$, $\alpha_1, \dots, \alpha_M$ and β_1, \dots, β_M such that

$$\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n = \sum_{l=1}^M \alpha_l \ln \beta_l.$$

For the general deterministic HMM we simulate the joint state distribution using a Markov chain M . This Markov chain has states labelled (q_i, q_j) for $i, j = 1, \dots, |Q|$ and also a special absorbing state \perp . The transition probabilities are taken from the outgoing transitions from q_j . It may be that a letter produced from q_j cannot be produced from q_i in which case the Markov chain transitions to \perp . The full definition is

$$\begin{aligned}
M_{(i,j),(i^*,j^*)} &= \sum_{a \in \Sigma} \Psi(a)_{j,j^*} \delta_{\Psi(a)_{i,i^*} > 0} \\
M_{(i,j),\perp} &= 1 - \sum_{(i^*,j^*) \in [Q \times Q]} M_{(i,j),(i^*,j^*)} \\
M_{\perp,\perp} &= 1 \\
M_{\perp,(i^*,j^*)} &= 0.
\end{aligned}$$

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500 We write T for the set of transitions of M . Let $C_0 = \{\perp\}$ and let C_1, \dots, C_K be the
501 remaining bottom connected components of M . It follows that the sets C_0, \dots, C_K are
502 disjoint.

503 ► **Theorem 41.** *Let (Q, Σ, Ψ) be a deterministic HMM and let i_0, j_0 be the indices of starting
504 states then we may compute a set of symbolic numbers $\{\lambda_1, \dots, \lambda_K\} \subseteq [-\infty, 0]$ and a
505 probability distribution over these numbers given as p_1, \dots, p_K in $O(|Q|^8)$ time such that,*

$$\mathbb{P}_{\delta_{j_0}} \left(\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n = \lambda_k \right) = p_k.$$

506 for all $k = 1, \dots, K \leq |Q|^2$.

508 **Proof.** First assume that the pair (i_0, j_0) is in the bottom connected component C_k with
509 $k \geq 1$ of M . Fix $n \in \mathbb{N}$ then for any $a_1 \dots a_n \in \Sigma^n$, since Ψ is deterministic, there are unique
510 sets of states indexed by i_1, \dots, i_n and j_1, \dots, j_n such that

$$\frac{1}{n} \ln L_n = \frac{1}{n} \ln \frac{\Psi(a_1)_{i_0,i_1} \Psi(a_2)_{i_1,i_2} \dots \Psi(a_n)_{i_{n-1},i_n}}{\Psi(a_1)_{j_0,j_1} \Psi(a_2)_{j_1,j_2} \dots \Psi(a_n)_{j_{n-1},j_n}} = \frac{1}{n} \sum_{m=1}^n \ln \frac{\Psi(a_m)_{(i_{m-1},i_m)}}{\Psi(a_m)_{(j_{m-1},j_m)}}.$$

512 The state pair (i_m, j_m) has the same distribution under $\mathbb{P}_{\delta_{j_0}}$ as the state distribution of the
513 sub-Markov chain M started from (i_0, j_0) after m transitions. For each transition in M there
514 is a mapping $l : T \rightarrow (-\infty, 0]$ given by

$$l((i, j) \rightarrow (i^*, j^*)) = \ln \frac{\Psi(a)_{(i,i^*)}}{\Psi(a)_{(j,j^*)}}.$$

516 l is well defined because Ψ is deterministic and so a is unique to each transition. Let μ^k be
517 the unique stationary distribution of the BCC C_k . Then by the strong Ergodic theorem,
518 $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n$ converges to

$$\sum_{(i,j) \rightarrow (i^*,j^*) \in T} \mu_{(i,j)}^k M_{(i,j),(i^*,j^*)} l((i,j) \rightarrow (i^*,j^*)).$$

520 If the starting state (i_0, j_0) is in C_0 , then for all possible transitions from j_0 in Ψ , there
521 are no transitions from i_0 and so it follows that $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n = -\infty$. Computing the
522 bottom connected components can be done using Tarjan's algorithm in $O(|Q|^2)$ time. Then,
523 computing the stationary distributions μ^k for $k = 1, \dots, K$ takes $O(|Q|^6)$ time. Therefore,
524 computing the possible $\lambda_1, \dots, \lambda_K$ takes $O(|Q|^6)$ time.

525 We may write $-\infty = \lambda_0, \dots, \lambda_K$ for the possible limits of $\frac{1}{n} \ln L_n$ starting in C_0, \dots, C_K
526 respectively. We write $C^{-1}(\lambda_k)$ for the union of bottom connected components associated
527 with the likelihood exponent λ_k . In general, for starting states (i, j) , $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n = \lambda_k$
528 when the Markov chain M hits any of the states in $C^{-1}(\lambda_k)$. Therefore computing p_1, \dots, p_K
529 is equivalent to computing a hitting probability which can be done in $O(|Q|^6)$ because M
530 has dimension at most $|Q|^2$. This makes the whole algorithm $O(|Q|^8)$ since $K \leq |Q|^2$. ◀

5.2 Non-Deterministic Chains

- show that it's atleast PSPACE hard to approximate and computing them is as hard as computing lyapunov exponents (known to be a hard problem).

By introducing non-determinism, the likelihood ratio can no longer be split into a simple product of random variables depending only on the transitions in a Markov chain. It follows that we may not use the Strong ergodic theorem for Markov chains to show convergence of $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n$. We instead prove convergence using work by Protasov. The following theorem is proven in the appendix.

There is no known algorithm for computing each limit $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n^k$. Moreover, even if we know the values of $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n^k$, we show that computing the distribution of $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n$ over the possible limits $\{\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n^k \mid k = 1, \dots, K\}$ is PSPACE-complete. We therefore consider a series of problems leading up to the final result.

The starting point is language inclusion for *non-deterministic finite automata* (NFA). An NFA is a quintuple $(Q, \Sigma, \Delta, q_0, F)$ such that The starting point is language inclusion for *non-deterministic finite automata* (NFA). An NFA is a quintuple $(Q, \Sigma, \Delta, q_0, F)$ such that

- a finite set of states Q .
- a finite set of input symbols Σ .
- a transition function $\Delta : Q \times \Sigma \rightarrow \mathcal{P}(Q)$.
- an initial state q_0
- a set of final states $F \subseteq Q$.

The language accepted by M , $L(M)$ is the set of $a_1 \cdots a_n \in \Sigma^n$ such that there exists a sequence of states q_1, \dots, q_n such that $q_i \in \Delta(q_{i-1}, a_i)$ for $i = 1, \dots, n$ and $q_n \in F$.

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► **Theorem 5.** We have $\lim_{n \rightarrow \infty} L_n = 0$ \mathbb{P}_{π_2} -a.s. if and only if π_1 and π_2 are distinguishable.

Proof. By proposition 6 in [2] the limit L_n exists \mathbb{P}_{π_2} -a.s. and the following equalities hold.

$$\begin{aligned}
 1 - d(\pi_1, \pi_2) &= \lim_{n \rightarrow \infty} \sum_{w \in \Sigma^n} \|\pi_1 \Psi(w)\| \wedge \|\pi_2 \Psi(w)\| && \text{Theorem 7 of [2]} \\
 &= \lim_{n \rightarrow \infty} \sum_{w \in \Sigma^n} (L_n(w) \wedge 1) \|\pi_2 \Psi(w)\| \\
 &= \lim_{n \rightarrow \infty} \mathbb{E}_{\pi_2} [L_n \wedge 1] \\
 &= \mathbb{E}_{\pi_2} \left[\lim_{n \rightarrow \infty} L_n \wedge 1 \right] && \text{because } 0 \leq L_n(w) \wedge 1 \leq 1.
 \end{aligned}$$

Then, $\lim_{n \rightarrow \infty} L_n \wedge 1 = 0 \iff \lim_{n \rightarrow \infty} L_n = 0$. ◀

A

 Proofs from Section 3

► **Theorem 7.** Suppose π_1 and π_2 are distinguishable. Let α, β be inputs to the SPRT. Then by choosing $A = \ln \frac{\alpha}{1-\beta}$ and $B = \ln \frac{1-\alpha}{\beta}$, $\mathbb{P}_{\pi_1}(\text{SPRT}_{\alpha, \beta} = \pi_2) \leq \alpha$ and $\mathbb{P}_{\pi_2}(\text{SPRT}_{\alpha, \beta} = \pi_1) \leq \beta$.

Proof. We wish to control the probabilities $\mathbb{P}_{\pi_2}(L_N > B)$ and $\mathbb{P}_{\pi_1}(L_N < A)$ by choosing suitable values of A and B . Let $W_n^1 = \{w \in \Sigma^\omega \mid A \leq L_m(w) \leq B \ \forall m < n, L_n < A\}$ then

$$\begin{aligned}
 \mathbb{P}_{\pi_1}(L_N < A) &= \sum_{n=1}^{\infty} \mathbb{P}_{\pi_1}(W_n^1) = \sum_{n=1}^{\infty} \sum_{w \in W_n^1} \pi_1 \Psi(w) \mathbb{1}^T = \sum_{n=1}^{\infty} \sum_{w \in W_n^1} L_n(w) \pi_2 \Psi(w) \mathbb{1}^T \\
 &\leq A \sum_{n=1}^{\infty} \sum_{w \in W_n^1} \pi_2 \Psi(w) \mathbb{1}^T = A \sum_{n=1}^{\infty} \mathbb{P}_{\pi_2}(W_n^1) = A \mathbb{P}_{\pi_2}(L_N < A).
 \end{aligned}$$

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Similarly, we may derive $\mathbb{P}_{\pi_2}(L_N > b) \geq \frac{1}{b} \mathbb{P}_{\pi_1}(L_N > b)$ so it follows that

$$\begin{aligned}
 A &\geq \frac{\mathbb{P}_{\pi_1}(L_N < A)}{\mathbb{P}_{\pi_2}(L_N < A)} = \frac{\mathbb{P}_{\pi_1}(L_N < A)}{1 - \mathbb{P}_{\pi_2}(L_N > B)} \\
 B &\leq \frac{\mathbb{P}_{\pi_1}(L_N > B)}{\mathbb{P}_{\pi_2}(L_N > B)} = \frac{1 - \mathbb{P}_{\pi_1}(L_N < A)}{\mathbb{P}_{\pi_2}(L_N > B)}
 \end{aligned}$$

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to guarantee the error bounds $\alpha = \mathbb{P}_{\pi_1}(L_{n^*} < A)$ and $\beta = \mathbb{P}_{\pi_2}(L_{n^*} > B)$. ◀

A.1 Proof of Theorem 40

► **Theorem 11.** The set of likelihood exponents satisfies $\Lambda \subset [-\infty, 0]$ and $|\Lambda| \leq |Q|^2$. Further $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n$ exists (and by definition is in Λ) \mathbb{P}_{π_2} -almost surely for any π_1, π_2 .

The proof of Theorem 11 relies on related work in a subset of Ergodic Theory called *Lyapunov exponents*. The main papers are by Protasov [5] and Osedelets [?]. We first must define a similar object to a observation density matrix. A *random matrix product* is a triple (Q, Σ, Φ) where Q is a set of states, Σ is a set of letters and $\Phi : \Sigma \rightarrow [0, 1]^{|Q| \times |Q|}$ is a non-negative matrix valued function. We may extend Φ to Σ^* in the same way as an observation density matrix. we use the shorthand $\Phi_n : \Sigma^n \rightarrow [0, 1]^{|Q| \times |Q|}$ for Φ restricted

to Σ^n . (Q, Σ, Φ) is strongly connected if for all $i, j \in [Q]$ there exists a $w \in \Sigma^*$ such that $\Phi(w)_{i,j} > 0$. (Q, Σ, Φ) is mortal if there exists a $w \in \Sigma^*$ such that $\Phi(w) = 0$. The main theorem by Protasov is stated below.

► **Theorem 42. Protasov's Theorem** *Let (Q, Σ, Φ) be a strongly connected random matrix product, let \mathbb{P}_{ind} be an i.i.d probability measure on Σ^ω and let π be an initial distribution. If (Q, Σ, Φ) is mortal, then $\lim_{n \rightarrow \infty} \frac{1}{n} \ln \pi \Phi_n \mathbb{1}^T = -\infty$ \mathbb{P}_{ind} -a.s. otherwise there is a $\lambda \in (-\infty, 0]$ such that $\lim_{n \rightarrow \infty} \frac{1}{n} \ln \pi \Phi_n \mathbb{1}^T = \lambda$ \mathbb{P}_{ind} -a.s.*

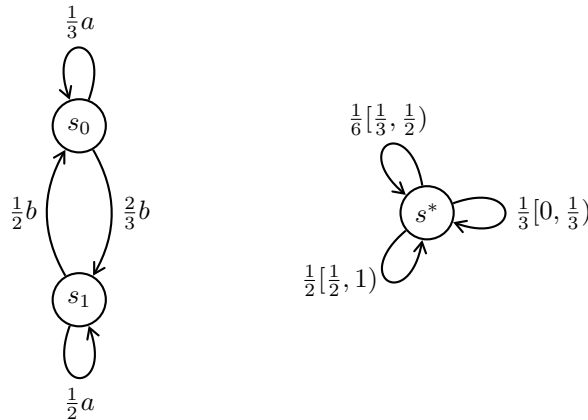
The technique we use to prove Theorem 40 involves producing an i.i.d measure \mathbb{P}_{ind} and random matrix product $(Q \times Q, \Sigma, \Phi)$ such that for all $n \in \mathbb{N}$ and initial distributions π ,

$$\mathbb{P}_{\text{ind}}\left(\frac{1}{n} \ln \pi \Phi_n \mathbb{1}^T \in A\right) = \mathbb{P}_{\pi_2}\left(\frac{1}{n} \ln \pi \Psi_n \mathbb{1}^T \in A\right).$$

Since $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n = \lim_{n \rightarrow \infty} \frac{1}{n} \ln \pi_1 \Psi_n \mathbb{1}^T - \lim_{n \rightarrow \infty} \frac{1}{n} \ln \pi_2 \Psi_n \mathbb{1}^T$ such a random matrix product would imply convergence of $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n$ due to Protasov's theorem.

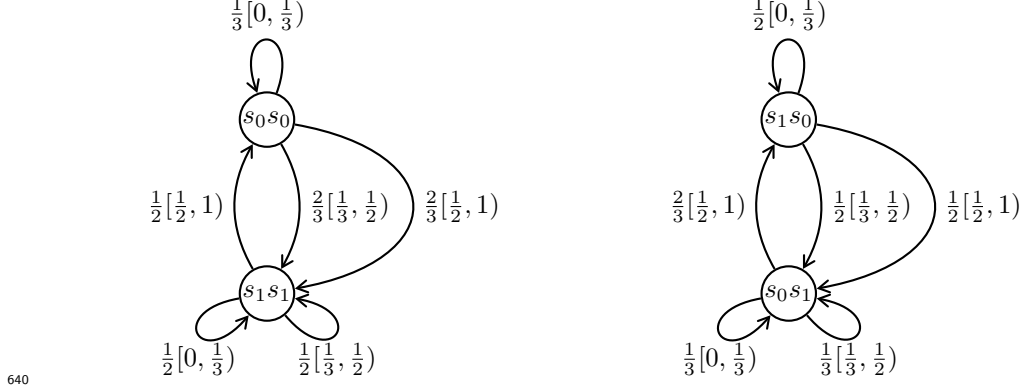
When the probability space on infinite words is derived from a general HMM, the probability of a particular letter being produced at a specific position in the infinite word depends on the state the producing HMM is in. In the constructed random matrix product, the current state of the producing HMM as started from π_2 is incorporated into the state space of the HMM started from π_1 .

To accomplish this, we simulate transitions in the producing HMM (the chain started from π_2) by sampling a uniform random number in the interval $[0, 1)$. At each state, we may partition $[0, 1)$ so that each sub-interval corresponds to specific transition and the size of the sub-interval corresponds to the probability of said transition. The union over all states of these partitions has a minimal finite σ -algebra. The atoms of this σ -algebra are also a partition of $[0, 1)$ and so we may sample them independently at random with probabilities according to their size. The transformation on a simple example is demonstrated in the diagram below. On the left is the original HMM. On the right is a single state HMM representing \mathbb{P}_{ind} . For example, with probability $\frac{1}{6}$, it produces the label $[\frac{1}{3}, \frac{1}{2})$ which corresponds to the b transition in s_0 or the a transition in s_1 .



We now construct the random matrix product $(Q \times Q, \Sigma, \Phi)$ which can also be represented by a state transition system (but without any stochastic properties). In our example, $(Q \times Q, \Sigma, \Phi)$ consists of two strongly connected components. Each state is labelled with a left and right component state corresponding to the current states of the HMM started from

π_1 and π_2 respectively. The transition weights are taken from the transitions in π_1 and the transition letters are taken from the transitions in π_2 because this random matrix product simulates $\frac{1}{n} \ln \pi_1 \Psi_n \mathbb{1}^T$ under the probability measure \mathbb{P}_{π_2} .



A.1.0.1 Random Matrix Product Construction

Consider the HMM (Q, Σ, Ψ) with finite alphabet Σ . Since for each $i \in [Q]$, $\sum_{a \in \Sigma} \sum_{j=1}^{|Q|} \Psi(a)_{i,j} = 1$ it follows that we may define a function $\rho_i : [0, 1) \rightarrow Q \times \Sigma$ such that for all $i \in [Q]$, $\mathcal{M}_{Leb}(\rho_i^{-1}\{(j, a)\}) = \Psi(a)_{i,j}$. Consider the minimal σ -algebra $\sigma\{\rho_i^{-1}\{(j, a)\} \mid i, j \in [Q], a \in \Sigma\}$ which is finite and has a set of atomic elements P of at most $|Q|^2|\Sigma|$ elements. P is also a partition of $[0, 1)$. Let $p \in P$ then $\rho_i(x)$ is constant for all $x \in p$ so we may overload the notation and consider the function $\rho_i : P \rightarrow Q \times \Sigma$. We then extend ρ_i to $\rho_i : P^n \rightarrow (Q \times \Sigma)^n$ by iteratively defining $\rho_i(ua) = \rho_i(u)\rho_{l(\rho_i(u))}(a)$ for $u \in P^n$ and $a \in P$.

We will describe the one-state chain $(\{1\}, P, \Psi_P)$ as the *singleton generator* for Ψ where $\Psi_P(p) = \mathcal{M}_{Leb}(p)$. Given an initial distribution π for (Q, Σ, Ψ) , a word generated by its singleton generator uniquely defines a path of states and letters. Let \mathbb{P}_{ind} be the measure on the set of infinite words P^ω generated by this HMM.

In order to incorporate the state space of the producing HMM into the state space of the chain started from π_1 , we define two functions $l : Q \times \Sigma \rightarrow Q$ and $r : Q \times \Sigma \rightarrow \Sigma$ where $l(q, a) = q$ and $r(q, a) = a$. We then define the matrix valued function $\Phi : P \rightarrow [0, 1]^{(Q \times Q) \times (Q \times Q)}$ as

$$\Phi(p)_{(i_1, j_1), (i_2, j_2)} = \begin{cases} \Psi(r \circ \rho_{j_1}(p))_{i_1, i_2} & l \circ \rho_{j_1}(p) = j_2 \\ 0 & \text{else.} \end{cases}$$

Further for initial distributions $\pi_1, \pi_2 \in [0, 1]^{|Q|}$ we define $\pi_1 \times \pi_2 = \pi_1^T \pi_2 \in [0, 1]^{|Q| \times |Q|}$.

► Proposition 43. Consider the HMM (Q, Σ, Ψ) with initial distributions π_1 and π_2 then for any measurable set $A \in [-\infty, 0]$ and for all $n \in \mathbb{N}$

$$\mathbb{P}_{\pi_2} \left(\frac{1}{n} \ln L_n \in A \right) = \mathbb{P}_{\text{ind}} \left(\frac{1}{n} \ln \frac{\pi_1 \times \pi_2 \Phi_n \mathbb{1}}{\pi_2 \times \pi_2 \Phi_n \mathbb{1}} \in A \right).$$

Proof. Let $i_0, j_0 \in [Q]$ be initial states, It follows that for a word $u_1 \dots u_n \in \{\rho_{j_0}(u_1 \dots u_n) = (j_1, a_1), \dots, (j_n, a_n)\}$,

$$\begin{aligned}
665 \quad \|e_{i_0}^T e_{j_0} \Phi(u)\| &= \sum_{(i_1, \dots, i_n) \in [Q]^n} \Phi(u_1)_{(i_0, j_0), (i_1, j_1)} \dots \Phi(u_n)_{(i_{n-1}, j_{n-1}), (i_n, j_n)} \\
666 \quad &= \sum_{(i_1, \dots, i_n) \in [Q]^n} \Psi(a_1)_{i_0, i_1} \dots \Psi(a_n)_{i_{n-1}, i_n} \\
667 \quad &= \|e_{i_0} \Psi(a_1 \dots a_n)\|. \\
668
\end{aligned}$$

669 It follows that for any initial distributions π_1, π_2 , $\|\pi_1^T \pi_2 \Psi^*(u)\| = \|\pi_1 \Psi(a_1 \dots a_n)\|$. Then
 670 considering the word a_1, \dots, a_n ,

$$\begin{aligned}
671 \quad \mathbb{P}_{\pi_2}(a_1 \dots a_n) &= \|\pi_2 \Psi(a_1 \dots a_n)\| \\
672 \quad &= \sum_{j_1, \dots, j_n \in [Q]^n} \|\pi_2 \Psi(a_1)_{\pi_2, j_1} \dots \Psi(a_n)_{j_{n-1}, j_n}\| \\
673 \quad &= \sum_{u_1, \dots, u_n \in \{\rho_{\pi_2}(u_1, \dots, u_n) = (j_1, a_1), \dots, (j_n, a_n)\}} \Psi_P(u_1 \dots u_n) \\
674 \quad &= \mathbb{P}_{\text{ind}}(l \circ \rho_{\pi_2}(u_1 \dots u_n) = a_1 \dots a_n) \\
675
\end{aligned}$$

676

677 Consider the bottom connected components $C_1, \dots, C_k \subseteq Q \times Q$ of the Markov chain
 678 defined by $\sum_{p \in P} \Phi(p)$ and let $d : Q \times Q \rightarrow Q \times Q$ be defined as $d(s_1, s_2) = (s_2, s_2)$. The
 679 following lemmas give the $|Q|^2$ bound in Theorem 40.

680 ► **Lemma 44.** *Let s be a starting state for an HMM (Q, Σ, Ψ) . Then $\lim_{n \rightarrow \infty} \frac{1}{n} \ln \|(s, s) \Psi_n^*\|$
 681 takes at most $|Q|$ values.*

682 **Proof.** Let P_1, \dots, P_K be the irreducible components of Φ . Consider states $(s_1, s_2), (r_1, r_2) \in$
 683 $Q \times Q$. If there is a path from (s_1, s_2) to (r_1, r_2) then there is also a path from (s_2, s_2) to
 684 (r_2, r_2) . Therefore for any end component P_i it follows that the image $d(P_i) \subseteq P_j$ for some
 685 end component P_j and so we may define a function $\rho : \{P_1, \dots, P_K\} \rightarrow \{P_1, \dots, P_K\}$ such
 686 that $\rho(P_i) = P_j$. Suppose P_i, P_j have Lyapunov exponents λ_i and λ_j respectively. Let π_1
 687 and π_2 be initial distributions such that the support of π_1 is in P_i and the support of π_2
 688 is in P_j then the likelihood ratio $L_n = \frac{\|\pi_1 \Psi_n\|}{\|\pi_2 \Psi_n\|}$ converges to a limit in the set $[0, \infty)$ with
 689 respect to the measure \mathbb{P}_{π_2} , the same limit as $\frac{\|(\pi_1, \pi_2) \Psi_n^*\|}{\|(\pi_2, \pi_2) \Psi_n^*\|}$ with respect to \mathbb{P}_{ind} . Since both
 690 $\frac{1}{n} \ln \|(\pi_1, \pi_2) \Psi_n^*\|$ and $\frac{1}{n} \ln \|(\pi_2, \pi_2) \Psi_n^*\|$ converge almost surely in the set $[-\infty, 0]$ to λ_i and
 691 λ_j respectively, $\frac{1}{n} \ln \frac{\|\pi_1 \Psi_n\|}{\|\pi_2 \Psi_n\|}$ converges in $[-\infty, 0]$ with respect to \mathbb{P}_{π_2} . Therefore $\lambda_i \leq \lambda_j$.
 692 Now consider $\lim_{n \rightarrow \infty} \frac{1}{n} \ln \|(s, s) \Psi_n^*\|$ whose possible limits is bounded by $|Q|^2$. Suppose for
 693 some word $w \in P^n$ the support of $(s, s) \Psi_n^*$ intersects an irreducible component P_i . Then
 694 it must also intersect P_j . Since $\lambda_i \leq \lambda_j$ it follows that $\lim_{n \rightarrow \infty} \frac{1}{n} \ln \|(s, s) \Psi_n^*\| \geq \lambda_j$. Since
 695 the image $\rho\{P_1, \dots, P_K\} \leq |Q|$ it follows that $\lim_{n \rightarrow \infty} \frac{1}{n} \ln \|(s, s) \Psi_n^*\|$ takes at most $|Q|$
 696 values. ◀

697 ► **Lemma 45.** *Consider an HMM (Q, Σ, Ψ) . Let $\mathcal{E} : [0, 1]^{|Q|} \times [0, 1]^{|Q|} \rightarrow [\infty, 0]$ be a
 698 parametrised random variable on Σ^ω defined by $\mathcal{E}(\pi_1, \pi_2) = \lim_{n \rightarrow \infty} \frac{1}{n} \ln \frac{\|\pi_1 \Psi(w)\|}{\|\pi_2 \Psi(w)\|}$. Then
 699 $|\{\mathcal{E}(\pi_1, \pi_2) \mid \pi_1, \pi_2 \in [0, 1]^{|Q|}\}| \leq |Q|^2$ with respect to the measure \mathbb{P}_{π_2} .*

Proof. First consider the case of dirac distributions $\pi_1 = \delta_r$ and $\pi_2 = \delta_s$. We may instead consider a bound on

$$\lim_{n \rightarrow \infty} \frac{1}{n} \ln \frac{\|(\delta_r, \delta_s)\Psi_n^*\|}{\|(\delta_s, \delta_s)\Psi_n^*\|} = \lim_{n \rightarrow \infty} \frac{1}{n} \ln \|(\delta_r, \delta_s)\Psi_n^*\| - \lim_{n \rightarrow \infty} \frac{1}{n} \ln \|(\delta_s, \delta_s)\Psi_n^*\|$$

with respect to the \mathbb{P}_{ind} measure. Let C_1, \dots, C_K be the irreducible lethal components of Φ . For $L \leq K$ without loss of generality suppose C_1, \dots, C_L be the irreducible components that are also end components containing diagonal entries. Let $R_1, \dots, R_L \subseteq Q$ be disjoint and have the property that for all $q \in R_i$, $(q, q) \in C_i$.

Given a state $q_i \in Q$ any letter in P defines a unique subsequent state q_j and a unique letter produced a . Therefore, projecting $(\delta_p, \delta_q)\Phi_n$ onto its right component, yields a point distribution on some state. Therefore the function $\zeta : \{(\delta_p, \delta_q)\Phi(w) \mid w \in P^*\} \rightarrow Q$ defined by $\zeta((\delta_p, \delta_q)\Phi(w)) = r$ if and only if $\text{supp} \sum_{i=1}^{|Q|} ((\delta_p)_i, \delta_q)\Phi(w)_{i,j} = (\delta_r)_j$ for all j is well defined.

We may partition Σ^ω into W_1, \dots, W_L such that $\zeta((\delta_r, \delta_s)\Phi_n(w))$ hits all states in R_k infinitely often for $w \in W_k$. It follows that $(\delta_s, \delta_s)\Phi_n(w)$ intersects the end component C_k and hits no other end components with diagonal entries. let $q \in C_k$ then $\frac{1}{n} \ln (\delta_s, \delta_s)\Phi_n(w)$ must converge almost surely on W_k to the Lyapunov exponent given by $\lim_{n \rightarrow \infty} \frac{1}{n} \ln (\delta_q, \delta_q)\Phi_n(w)$.

Similarly $\zeta((\delta_r, \delta_s)\Phi_n(w)) \in R_k$ and so $(\delta_r, \delta_s)\Phi_n(w)$ is contained in the set $Q \times R_k$ for $w \in W_k$. Since $\zeta((\delta_r, \delta_s)\Phi_n(w))$ hits all states in R_k infinitely often each irreducible component P_i such that $P_i \leq (\delta_r, \delta_s)\Phi_n(w)$ must have the property that $|P_i| \geq |R_k|$. Therefore the total number of irreducible components hit by $(\delta_r, \delta_s)\Phi_n(w)$ where $w \in W_k$ is at most $|Q|$. Since $L \leq |Q|$ the total number of possible limits for

$$\lim_{n \rightarrow \infty} \frac{1}{n} \ln \frac{\|(\delta_r, \delta_s)\Psi_n^*\|}{\|(\delta_s, \delta_s)\Psi_n^*\|}$$

is $|Q|^2$. It remains to show that $\{\mathcal{E}(\pi_1, \pi_2) \mid \pi_1, \pi_2 \in [0, 1]^{|Q|}\} \subseteq \{\mathcal{E}(\delta_r, \delta_s) \mid r, s \in Q\}$. Fix π_1 and π_2 and let us consider the possible values of

$$\lim_{n \rightarrow \infty} \frac{1}{n} \ln \|(\pi_1, \delta_s)\Psi_n^*\| - \lim_{n \rightarrow \infty} \frac{1}{n} \ln \|(\pi_2, \delta_s)\Psi_n^*\|$$

where $s \in \text{supp } \pi_2$. A consider again the partition of $\Sigma^\omega = \bigcup_{k=1}^L W_k$. For $w \in W_k$, the only end component in C_1, \dots, C_L hit by $(\pi_2, \delta_s)\Phi_n(w)$ is C_k . It follows that $\lim_{n \rightarrow \infty} \frac{1}{n} \ln \|(\pi_2, \delta_s)\Psi_n^*\| = \lim_{n \rightarrow \infty} \frac{1}{n} \ln \|(\delta_s, \delta_s)\Psi_n^*\|$. Any strongly connected component hit by $(\pi_1, \delta_s)\Psi_n^*(w)$, is also hit by $(\delta_r, \delta_s)\Psi_n^*(w)$ for some $r \in \text{supp } \pi_1$. It follows that we may partition Σ^ω so that on each part of the partition there is some $s, r \in Q$ such that

$$\lim_{n \rightarrow \infty} \frac{1}{n} \ln \frac{\|(\pi_1, \delta_s)\Psi_n^*\|}{\|(\pi_2, \delta_s)\Psi_n^*\|} = \lim_{n \rightarrow \infty} \frac{1}{n} \ln \frac{\|(\delta_r, \delta_s)\Psi_n^*\|}{\|(\delta_s, \delta_s)\Psi_n^*\|}.$$

731

We may now prove Theorem 40.

Proof of Theorem 40. We may compute Φ is $O(|Q|^4|\Sigma|)$ time and a set of strongly connected components using Tarjan's algorithm $P_1, \dots, P_K \subset Q \times Q$. On each connected component, any initial distribution converges to a constant or $-\infty$ by Theorem 42 and Proposition 43. The maximum of $|Q|^2$ possible limits is a result of and Lemma 45.

for each $k \in [K]$, $(r(P_k) \cup l(P_k), \Sigma, \Psi_{|l(P_k)} \oplus \Psi_{|r(P_k)})$ is a lossy HMM. Let $(q_k, r_k) \in P_k$ then δ_{q_k} and δ_{r_k} are initial distributions such that the corresponding likelihood ratios $\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n^k$ satisfy the requirements of the second half of the theorem. 739

A.2 Proof of the Asymptotic Wald Formula

► **Theorem 12.** [Generalised Wald Formula] Let ℓ be a likelihood exponent and let π_1 and π_2 be initial distributions. If $\ell \in (-\infty, 0)$ then

$$\mathbb{E}_{\pi_2} [N_{\alpha, \beta} \mid E_\ell] \sim \frac{\ln \alpha}{\ell} \quad (\text{as } \alpha, \beta \rightarrow 0).$$

If $\ell = 0$ then there exist $\alpha, \beta > 0$ such that

$$\mathbb{E}_{\pi_2} [N_{\alpha, \beta} \mid E_\ell] = \infty. \quad (4)$$

If $\ell = -\infty$ then

$$\sup_{\alpha, \beta} \mathbb{E}_{\pi_2} [N_{\alpha, \beta} \mid E_\ell] < \infty.$$

► **Lemma 46.** There is a number $c > 0$, computable in polynomial time, such that

$$\mathbb{P}_{\pi_2} (L_{2|Q|n} \geq \exp(-\frac{c^2}{36}n)) \leq 4 \exp(-\frac{c^2}{36}n).$$

Proof. By Proposition 17, $\lim_{\alpha, \beta \rightarrow 0} \frac{N_{\alpha, \beta}}{\ln \alpha}$ exists $\mathbb{P}_{\pi_2}(\cdot \mid V_k)$ -almost surely. Hence, the convergence is also in $\mathbb{P}_{\pi_2}(\cdot \mid V_k)$ -measure. Therefore, by the Vitali convergence theorem[1] it is sufficient to show that the set of random variables $\{\frac{N_{\alpha, \beta}}{\ln \alpha} \mid \alpha, \beta \in (0, \frac{1}{2})\}$ is uniformly integrable conditioned on V_k . In fact, because

$$\lim_{K \rightarrow \infty} \sup_{\alpha, \beta} \mathbb{E}_{\pi_2} [\frac{N_{\alpha, \beta}}{-\ln \alpha} \mathbb{1}_{\frac{N_{\alpha, \beta}}{-\ln \alpha} \geq -K}] \geq \mathbb{P}_{\pi_2}(V_k) \lim_{M \rightarrow \infty} \sup_{\alpha, \beta} \frac{1}{-\ln \alpha} \mathbb{E}_{\pi_2} [\frac{N_{\alpha, \beta}}{-\ln \alpha} \mathbb{1}_{\frac{N_{\alpha, \beta}}{-\ln \alpha} \geq -K} \mid V_k], \quad (7)$$

It is sufficient to check the uniform integrability condition without conditioning on V_k .

For fixed $M \geq \frac{144|Q|}{c^2}$, write $m_\alpha = \lfloor \frac{-M \ln \alpha}{2|Q|} \rfloor$. It follows that $\frac{2|Q|m_\alpha}{\ln \alpha} \leq M$ and $\alpha \geq \exp(-\frac{c^2}{36}m_\alpha)$. Further, $m_\alpha \geq \frac{M \ln 2}{2|Q|} - 1$. The following series of equalities hold

$$\begin{aligned} \mathbb{E}_{\pi_2} [\frac{N_{\alpha, \beta}}{-\ln \alpha} \mathbb{1}_{\frac{N_{\alpha, \beta}}{-\ln \alpha} \geq M}] &= \frac{1}{-\ln \alpha} \sum_{n=0}^{\infty} \mathbb{P}_{\pi_2}(N_{\alpha, \beta} \mathbb{1}_{N_{\alpha, \beta} \geq 2|Q|m_\alpha}) && \leq \frac{2|Q|}{-\ln \alpha} (m_\alpha + 1) \\ &\leq M \mathbb{P}_{\pi_2}(L_{2|Q|m_\alpha} \geq \alpha) + \frac{2|Q|}{-\ln \alpha} \sum_{n=m_\alpha}^{\infty} \mathbb{P}_{\pi_2}(L_{2|Q|n} \geq \alpha) \\ &\leq M \mathbb{P}_{\pi_2}(L_{2|Q|m_\alpha} \geq \exp(-\frac{c^2}{36}m_\alpha)) + \frac{2|Q|}{-\ln \alpha} \sum_{n=m_\alpha}^{\infty} \mathbb{P}_{\pi_2}(L_{2|Q|n} \geq \exp(-\frac{c^2}{36}n)) \\ &\leq 4M \exp(-\frac{c^2}{36}m_\alpha) + \frac{8|Q|}{-\ln \alpha} \sum_{n=m_\alpha}^{\infty} \exp(-\frac{c^2}{36}n) \\ &\leq 4M \exp(-\frac{c^2}{36}m_\alpha) + \frac{8|Q| \exp(-\frac{c^2}{36}m_\alpha)}{-\ln \alpha} \frac{1}{1 - \exp c^2/36} \\ &\leq 4M \exp(-\frac{c^2}{36}(\frac{M \ln 2}{2|Q|} - 1)) + \frac{8|Q| \exp(-\frac{c^2}{36}(\frac{M \ln 2}{2|Q|} - 1))}{\ln 2} \frac{1}{1 - \exp c^2/36} \\ &\rightarrow 0 \end{aligned}$$

as $K \rightarrow \infty$ where the fourth inequality follows by Lemma 46. Hence, Equation (7) must hold. ◀

► **Lemma 47.** *The events $\{L_n \rightarrow 0\}$ and $\{\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n < 0\}$ are equal up to a \mathbb{P}_{π_2} -null set.*

Proof. Fix $A < \alpha < 0$ and define the event $W_n = \{1 > L_n \geq e^{n\alpha}\}$. Then

$$\begin{aligned}
 \mathbb{P}_{\pi_2}(\lim_{n \rightarrow \infty} \frac{1}{n} \ln L_n > \alpha) &\leq \mathbb{P}_{\pi_2}(\liminf_n \{\frac{1}{n} \ln L_n \geq \alpha\}) \\
 &\leq \liminf_n \mathbb{P}_{\pi_2}(\frac{1}{n} \ln L_n \geq \alpha) \\
 &\leq \liminf_n \mathbb{P}_{\pi_2}(L_n \geq e^{n\alpha}) \\
 &= \liminf_n [\mathbb{P}_{\pi_2}(1 > L_n \geq e^{n\alpha}) + \mathbb{P}_{\pi_2}(L_n \geq 1)] \\
 &\leq \liminf_n \left[\sum_{w \in W_n} \pi_2 \Psi(w) \mathbb{1}^T + e^{An+B} \right] \\
 &\leq \liminf_n \left[e^{-n\alpha} \sum_{w \in W_n} \pi_1 \Psi(w) \mathbb{1}^T \right] \\
 &\leq \liminf_n \left[e^{-n\alpha} \mathbb{P}_{\pi_1}(L_n < 1) \right] \\
 &\leq \liminf_n \left[e^{-n\alpha} e^{An+B} \right] \\
 &= 0.
 \end{aligned}$$

781

Now suppose π_1 and π_2 are distinguishable. By Theorem 5 of [3] one may compute a $c \in (-\infty, 0)$ such that

$$\begin{aligned}
 \mathbb{P}_{\pi_1}(L_n < 1) - \mathbb{P}_{\pi_2}(L_n < 1) &\geq 1 - 2e^{cn} \\
 &\geq 1 - 2e^{n \max \Lambda}
 \end{aligned}$$

For a given HMM (Q, Σ, Ψ) , we may define a *monitor* as a function $M_n : \Sigma^n \rightarrow \{1, 2\}$. A well designed monitor reads an input word from an HMM started with either π_1 or π_2 and aims to return 1 or 2 respectively with high probability. However the following series of inequalities hold

$$\begin{aligned}
 \mathbb{P}_{\pi_2}(M_n(w) = 2) - \mathbb{P}_{\pi_1}(M_n(w) = 2) &= \mathbb{P}_{\pi_2}(M_n(w) = 1) + \mathbb{P}_{\pi_1}(M_n(w) = 2) \\
 &= \sum_{w \in \Sigma^n} \pi_1 \Psi(w) \mathbb{1}^T \delta_{M_n(w)=1} + \pi_2 \Psi(w) \mathbb{1}^T \delta_{M_n(w)=2} \\
 &\geq \sum_{w \in \Sigma^n} \pi_1 \Psi(w) \mathbb{1}^T \wedge \pi_2 \Psi(w) \mathbb{1}^T \\
 &= \mathbb{P}_{\pi_2}(L_n \leq 1) - \mathbb{P}_{\pi_1}(L_n \leq 1)
 \end{aligned}$$

which means for any monitor, to guarantee an error probability bound of at most ϵ , we require atleast $\frac{\ln(\epsilon) - \ln 2}{\max \Lambda}$ observations. This bound motivates us to investigate computability properties of Λ .