

### CLOUD COMPUTING CONCEPTS with Indranil Gupta (Indy)

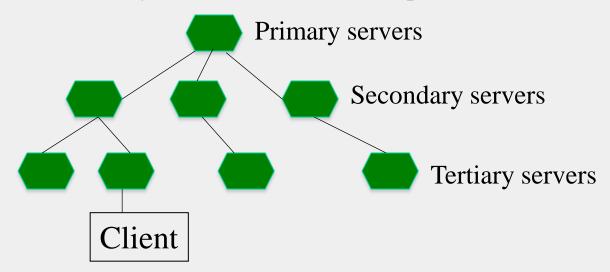
TIME AND ORDERING

Lecture C

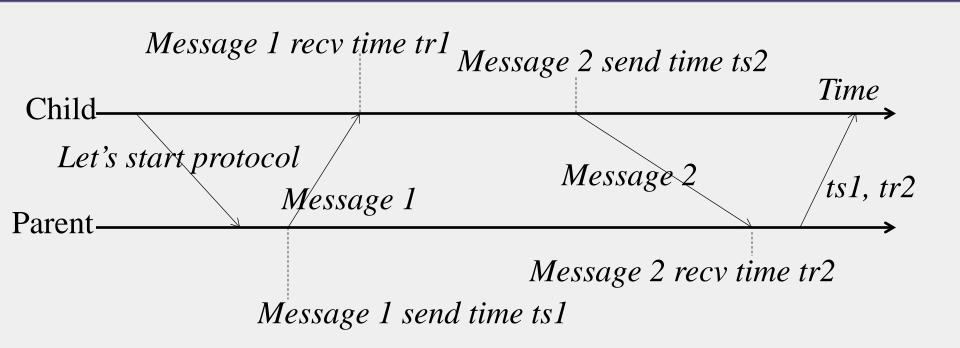
NTP

### NTP = Network Time Protocol

- NTP servers organized in a tree
- Each client = a leaf of tree
- Each node synchronizes with its tree parent



### NTP Protocol



### What the Child Does

- Child calculates *offset* between its clock and parent's clock
- Uses *ts1*, *tr1*, *ts2*, *tr2*
- Offset is calculated as

$$o = (tr1 - tr2 + ts2 - ts1)/2$$

# Why o = (trl - tr2 + ts2 - ts1)/2?

- Offset o = (tr1 tr2 + ts2 ts1)/2
- Let's calculate the error
- Suppose real offset is *oreal* 
  - Child is ahead of parent by *oreal*
  - Parent is ahead of child by -oreal
- Suppose one-way latency of Message 1 is *L1* (*L2* for Message 2)
- No one knows *L1* or *L2*!
- Then

$$tr1 = ts1 + L1 + oreal$$
  
 $tr2 = ts2 + L2 - oreal$ 

## Why o = (trl - tr2 + ts2 - tsl)/2?

#### • Then

$$tr1 = ts1 + L1 + oreal$$
  
 $tr2 = ts2 + L2 - oreal$ 

### Subtracting second equation from the first

```
oreal = (tr1 - tr2 + ts2 - ts1)/2 + (L2 - L1)/2
=> oreal = o + (L2 - L1)/2
=> |oreal - o| < |(L2 - L1)/2| < |(L2 + L1)/2|
- Thus, the error is bounded by the round-triptime
```

### And yet...

- We still have a non-zero error!
- We just can't seem to get rid of error
  - Can't, as long as message latencies are non-zero
- Can we avoid synchronizing clocks altogether and still be able to order events?