

Q.5. STRASSENS's Matrix multiplication

Practical implementation of Strassen's matrix multiplication algorithm usually switch to the brute force method after matrix sizes become smaller than some crossover point. Run an experiment to determine such crossover point on your computer system.

NORMAL MULTIPLICATION

```
function [ ct ] = multiply( a,b,n )
%UNTITLED4 Summary of this function goes here
% Detailed explanation goes here
ct=0;
c=zeros(n,n);
for i=1:n
    for j=1:n
        for k=1:n
            c(i,j)=c(i,j)+a(i,k)*b(k,j);
            ct=ct+1;
        end
    end
end
end
%disp(ct);
end
```

STRASSEN MULTIPLICATION

```
function [ C ct ] = strassen( A, B, ct )
n = size(A, 1);
%disp(n);
if (n == 1)
    C = A * B;
    ct = ct + 1;
else
    A11 = A(1: n/2 , 1: n/2);
    A12 = A(1: n/2 , n/2+1: n);
    A21 = A(n/2+1: n , 1: n/2);
    A22 = A(n/2+1: n , n/2+1: n);
    B11 = B(1: n/2 , 1: n/2);
    B12 = B(1: n/2 , n/2+1: n);
    B21 = B(n/2+1: n , 1: n/2);
    B22 = B(n/2+1: n , n/2+1: n);

    [M1 ct] = strassen(A11 + A22, B11 + B22, ct);
    [M2 ct] = strassen(A21 + A22, B11, ct);
    [M3 ct] = strassen(A11, B12 - B22, ct);
    [M4 ct] = strassen(A22 , B21 - B11, ct);
    [M5 ct] = strassen(A11 + A12, B22, ct);
    [M6 ct] = strassen(A21 - A11, B11 + B12, ct);
    [M7 ct] = strassen(A12 - A22, B21 + B22, ct);

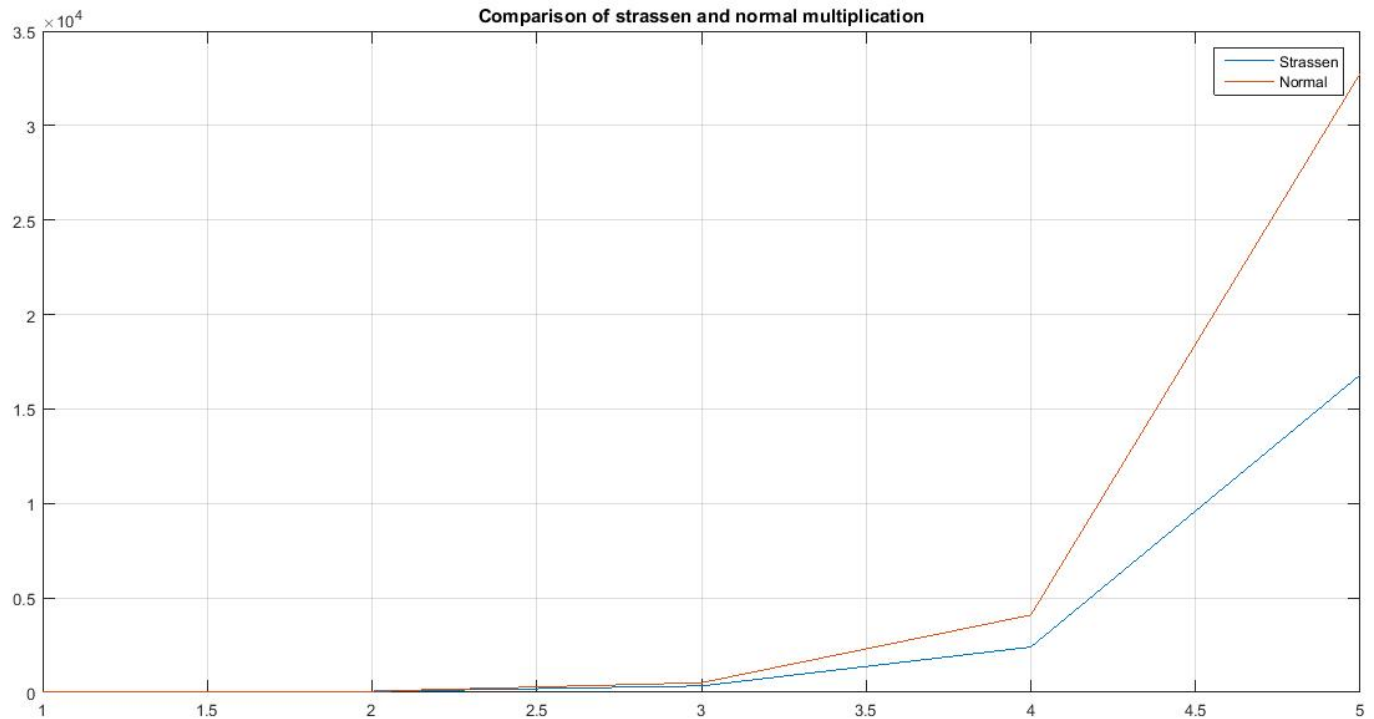
    C(1: n/2 , 1: n/2) = M1 + M4 - M5 + M7;
    C(1: n/2 , n/2+1: n) = M3 + M5;
```

```
C(n/2+1: n , 1: n/2) = M2 + M4;
C(n/2+1: n , n/2+1: n) = M1 + M3 - M2 + M6;
```

```
end
end
```

MAIN FUNCTION

```
r = [1:1:5];
s = zeros(length(r), 1);
n = zeros(length(r), 1);
for i = 1:length(r)
    x(i)=2^r(i);
    a = randi(100, x(i));
    b = randi(100, x(i));
    [dummy,s(i)] = strassen(a,b, 0);
    n(i) = multiply(a,b,x(i));
end
plot(r, s,r, n);
grid on;
title('Comparison of strassen and normal multiplication');
legend('Strassen', 'Normal');
```



Q.6. QUICK SELECT

Use the **QUICK SELECT** algorithm to find **3rd largest element** in an array of n integers. Analyze the performance of **QUICK SELECT** algorithm for the different instance of size 50 to 500 element. Record your observation with the *number of comparison made vs. instance*.

//PARTITION

```
function [ arr, index, n ] = partition_quickSelect( arr, left, right, n)

    pivot = arr(1);
    i = left;
    j = right;

    while(i < j)
        while(pivot >= arr(i) && i < right)
            i = i + 1;
            n = n + 1;
        end

        while(pivot < arr(j) && j > left )
            j = j - 1;
            n = n + 1;
        end

        if(i < j)
            temp = arr(i);
            arr(i) = arr(j);
            arr(j) = temp;
        end
    end

    arr(1) = arr(j);
    arr(j) = pivot;
    index = j;
end
```

//QUICKSELECT

```
function [ m, n ] = quickSelect(a, k, n)

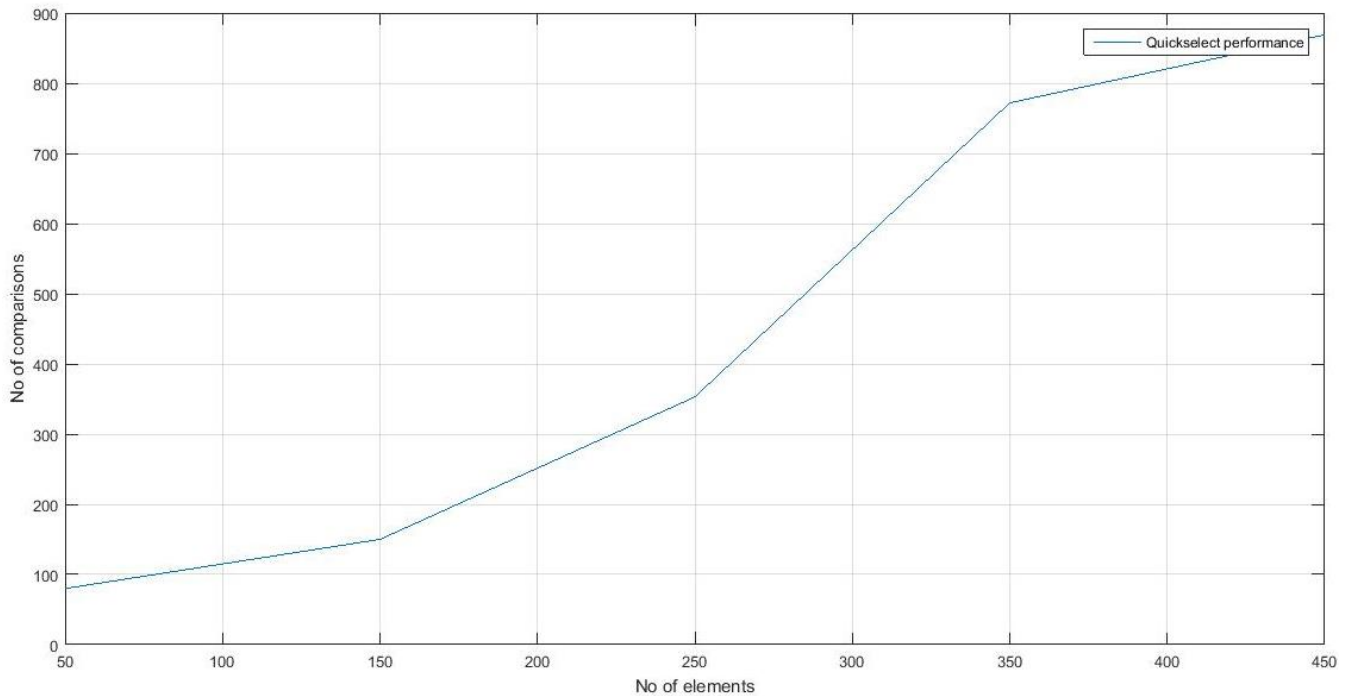
    l = 1;
    s = length(a);
    [a, pivot, n] = partition_quickSelect(a, l, s, n);
    if(pivot == k)
        m = a(pivot);
    elseif (pivot > k)
        s = pivot-1;
        a = a(1:s);
        [m, n] = quickSelect(a, k, n);
    end
```

```

else
    l = pivot + 1;
    k = k - pivot;
    a = a(l:s);
    [m, n] = quickSelect(a, k, n);
end
end

//MAIN FUNCTION
x = [50:25:500];
y = zeros(1,length(x));
fprintf('\nSize of array\t3rd largest element\n');
for i = 1:length(x)
    a = round(rand(x(i))*1000);
    %disp(a);
    [m, y(i)] = quickSelect(a, 3, 0);
    fprintf('\t%d\t\t\t\t\t%d\n',x(i),m);
end
plot(x,y);
xlabel('No of elements');
ylabel('No of comparisons');
legend('Quickselect performance');
grid on;

```



Size of array	3rd largest element
50	50
75	136
100	7
125	26
150	13
175	11
200	8
225	2
250	9
275	17
300	9
325	6
350	11
375	3
400	10
425	9
450	8
475	4
500	4

Q.7. Matrix Chain Multiplication Comparision

Given a matrix chain $A_1 \dots A_n$ with the dimension of each of the matrices given by the vector $\mathbf{p} = \langle 12, 21, 65, 18, 24, 93, 121, 16, 41, 31, 47, 5, 47, 29, 76, 18, 72, 15 \rangle$. ($n=17$) Write and run both the dynamic programming and memorized versions of this algorithm to determine the minimum number of multiplications that are needed (use type **longint**) and the factorization that produces this best case number of multiplications. Run each of the two programs over an appropriately large number of times (put each in a loop to run repeatedly x times) and obtain the times at the beginning and end of the run. Use these times to determine the comparative runtimes of the two algorithms.

//MATRIX CHAIN DYNAMIC PROGRAMMING

```
function [ c ] = MatrixChainOrder_DP( p,n )
%UNTITLED6 Summary of this function goes here
% Detailed explanation goes here
    m = zeros(n,n);
    for L = 2:n
        for i = 1:n-L+1
```

```

        j = i+L-1;
        m(i,j) = intmax();
        for k = i:j-1
            q = m(i,k) + m(k+1,j) + p(i)*p(k+1)*p(j+1);
            if q < m(i,j)
                m(i,j) = q;
            end
        end
    end
end
c = m(1,n);
%display(m);

end

```

//MAIN FUNCTION

```

p = [12,21,65,18,24,93,121,16,41,31,47,5,47,29,76,18,72,15];

%count1 = MatrixChainOrder_Memorised(p,2,18);
count2 = MatrixChainOrder_DP(p,17);
%display(count1);
fprintf('No. of operations required are : %d\n',count2);
%display(count2);

```

```

>> Ques_7

m =

Columns 1 through 10

    0    16380    29106    34290    61074    196110    219342    227214    242466    259950
    0         0    24570    33642    80514    313038    253290    267066    284042    312730
    0         0         0    28080    148986    384300    241392    282450    288206    333152
    0         0         0         0    40176    242730    222672    234480    251936    278162
    0         0         0         0         0    270072    215760    231504    248000    277456
    0         0         0         0         0         0    180048    241056    246512    293632
    0         0         0         0         0         0         0    79376    80352    134640
    0         0         0         0         0         0         0         0    20336    43648
    0         0         0         0         0         0         0         0         0    59737
    0         0         0         0         0         0         0         0         0         0
    0         0         0         0         0         0         0         0         0         0
    0         0         0         0         0         0         0         0         0         0
    0         0         0         0         0         0         0         0         0         0
    0         0         0         0         0         0         0         0         0         0
    0         0         0         0         0         0         0         0         0         0
    0         0         0         0         0         0         0         0         0         0
    0         0         0         0         0         0         0         0         0         0
    0         0         0         0         0         0         0         0         0         0

```

Columns 11 through 17

110120	112940	118675	132515	135875	145595	147575
108860	113795	118720	134675	135425	147575	146990
102035	117310	118275	144570	132560	156590	143465
96185	100415	105610	120860	122480	133820	134090
94025	99665	104320	120980	120860	133820	132380
82865	104720	103165	136040	115910	147500	126395
26600	55035	50960	90415	62165	101315	72230
16920	20680	26055	40835	43035	53835	54675
13640	23275	26400	47055	42005	59555	53270
7285	14570	18595	36900	34750	49600	46165
0	11045	13630	35695	28905	48075	40080
0	0	6815	17835	24675	31155	36555
0	0	0	103588	64206	125118	87387
0	0	0	0	39672	77256	66942
0	0	0	0	0	98496	39960
0	0	0	0	0	0	19440
0	0	0	0	0	0	0

No. of operations required are : 147575

Q.8. BINOMIAL COEFFICIENTS

Computing the **binomial coefficients** $C(n, k)$ defined by the following recursive formula:

$$C(n, k) = \begin{cases} 1, & \text{if } k = 0 \text{ or } k = n; \\ C(n-1, k) + C(n-1, k-1), & \text{if } 0 < k < n; \\ 0, & \text{otherwise.} \end{cases}$$

Write the program with three different algorithm to compute **binomial coefficients** $C(n, k)$ and compare them?

//FUNCTION

```
function [ ans ] = Binomial( n,k )
%UNTITLED5 Summary of this function goes here
% Detailed explanation goes here
c=zeros(n+1,k+1);
for i=1:n+1
    for j=1:min(i+1,k+1)
        %co=co+1;
        if j==1 || j==i+1
            c(i,j)=1;
        else
            c(i,j)=c(i-1,j-1)+c(i-1,j);
        end
    end
end
```

```

end
ans=c(n,k+1);
disp(c);
end

//MAIN FUNCTION
n=10;
k=5;
disp('Binomial Coefficients Table');
a=Binomial(n,k);
fprintf('Value of C(%d, %d) is %d \n', n, k, a );

```

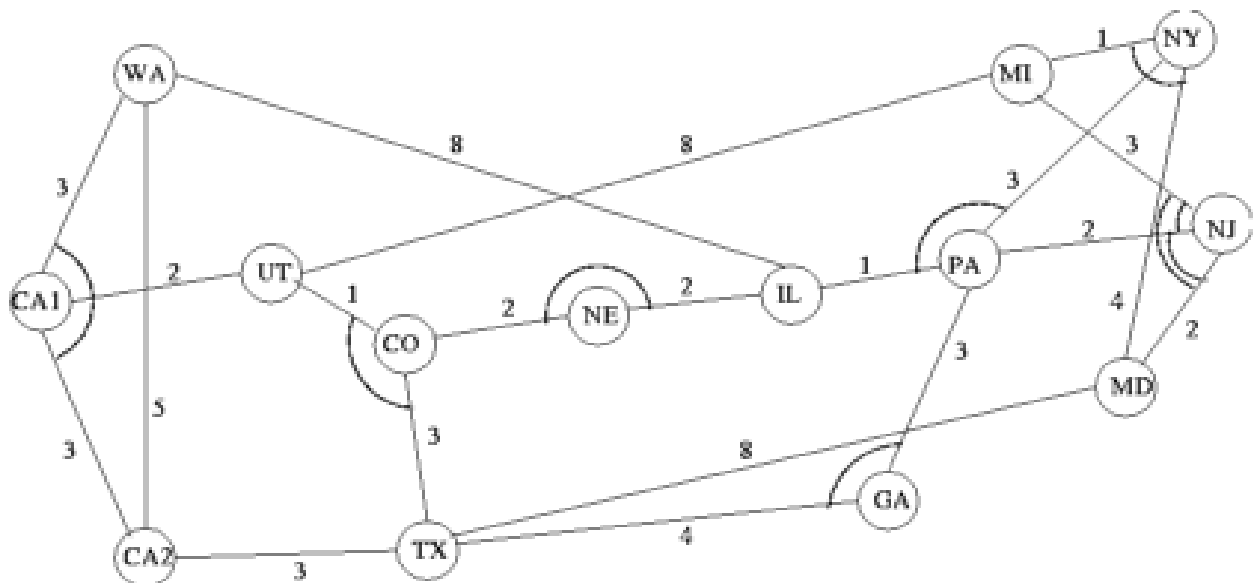
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>> Ques_8
Binomial Coefficients Table
  1   1   0   0   0   0
  1   2   1   0   0   0
  1   3   3   1   0   0
  1   4   6   4   1   0
  1   5  10  10   5   1
  1   6  15  20  15   6
  1   7  21  35  35  21
  1   8  28  56  70  56
  1   9  36  84 126 126
  1  10  45 120 210 252

Value of C(10, 5) is 252

```

Q.9. SPANNING TREE



Write a program obtain minimum cost spanning tree the above NSF network using Prim's algorithm, Kruskal's algorithm and Boruvka's algorithm..

Q.10. 0-1 KNAPSACK PROBLEM

Write a program that computes optimal solution to the 0–1 Knapsack Problem using dynamic programming? You may test your program with the following example:

There are $n = 5$ objects with integer weights $w[1..5] = \{1,2,5,6,7\}$, and values $v[1..5] = \{1,6,18,22,28\}$. Assuming a knapsack capacity of 11).

//FUNCTION

```
function [Max ,Seq] = knapsack_algo(Wts, V, Capacity)
    A = zeros(length(Wts)+1,Capacity+1);
    for j = 1:length(Wts)
        for X = 1:Capacity
            if Wts(j) > X
                A(j+1,X+1) = A(j,X+1);
            else
                A(j+1,X+1) = max( A(j,X+1), V(j) + A(j,X-Wts(j)+1));
            end
        end
    end
    Max = A(end,end);

    %Finding the optimal sequence
    Seq = zeros(length(Wts),1);
    a = Max;
    j = length(Wts);
    X = Capacity;
    while a > 0
        while A(j+1,X+1) == a
            j = j - 1;
        end
        j = j + 1;
        Seq(j) = 1;
        X = X - Wts(j);
        j = j - 1;
        a = A(j+1,X+1);
    end
end
```

//MAIN FUNCTION

```
n = 5;
Wts = [1, 2, 5, 6, 7]
V = [1, 6, 18, 22, 28]
Capacity = 11;
```

```
[max, Seq] = knapsack_algo(Wts, V, Capacity);  
fprintf('Max values is %d\n',max);  
items = find(Seq);  
fprintf('Items are \n');  
disp(items);
```

```
>> Ques_10  
  
Wts =  
  
     1     2     5     6     7  
  
V =  
  
     1     6    18    22    28  
  
Max values is 40  
Items are  
     3  
     4
```

Q.11. STRING MATCHING ALGORITHM

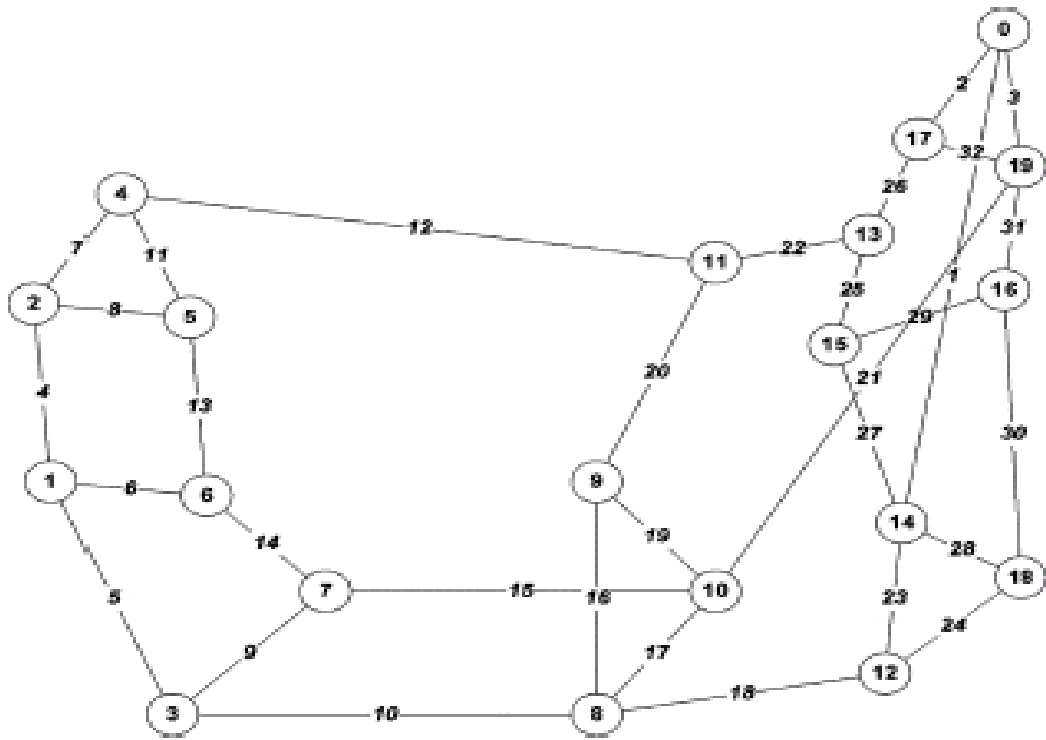
Given two strings P and T over the same alphabet set Σ , determine whether P occurs as a substring in T (or find in which position(s) P occurs as a substring in T). The strings P and T are called pattern and text respectively. Compare the efficiency of three string matching algorithms (Brute-Force Algorithm, Knuth-Morris-Pratt and Boyer-Moore Algorithm) by varying pattern length [1-15] for n=5000.

Q.12. MATRIX CHAIN MULTIPLICATION

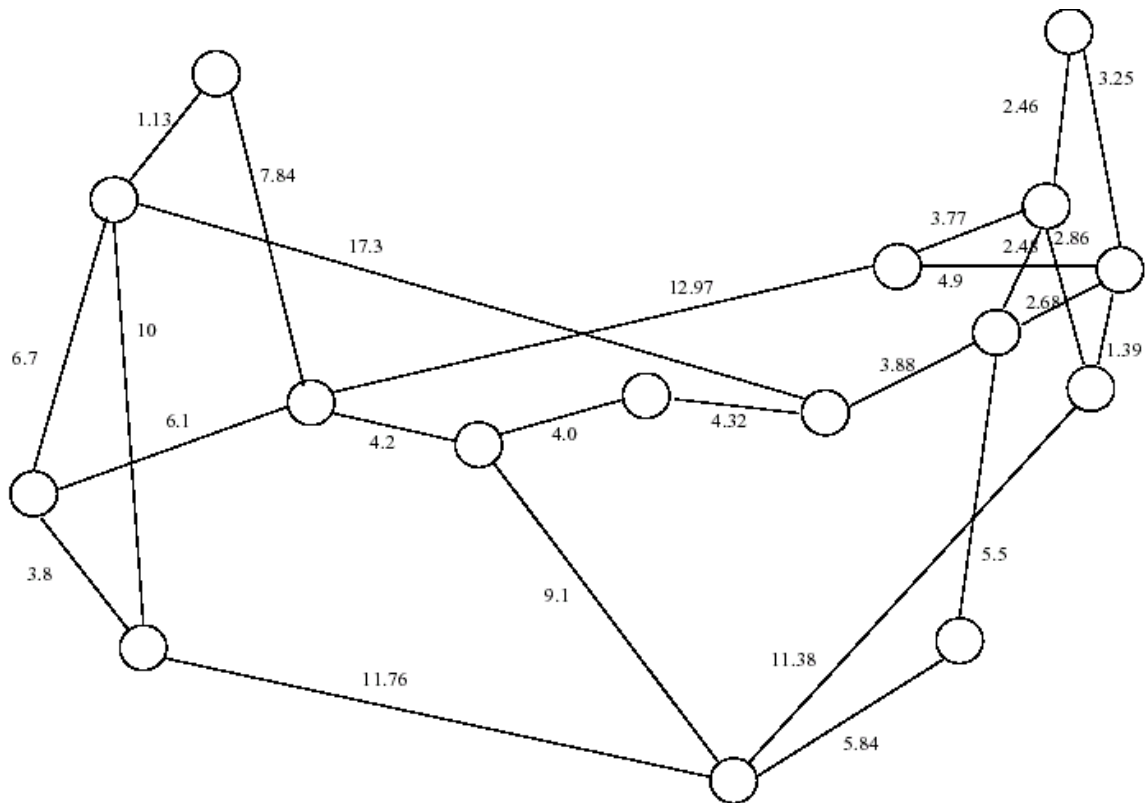
Write a program to compute the best ordering of matrix multiplication. Include the routine to print the actual ordering.

Q.13. PATHING ALGORITHM

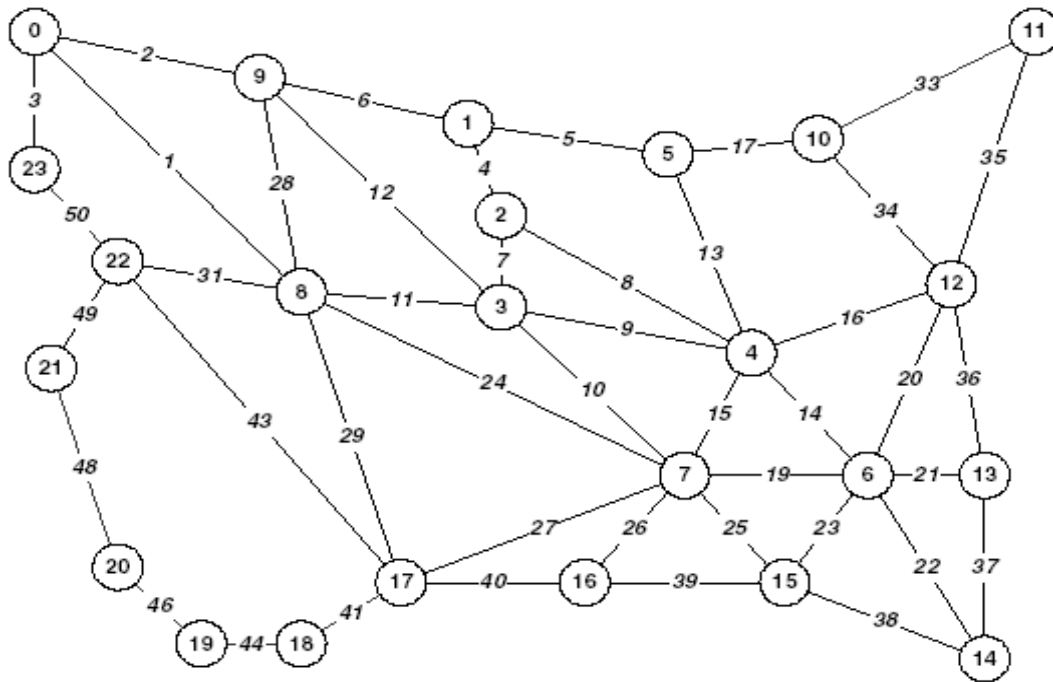
Use Floyd–Warshall algorithm (also known as Floyd's algorithm) to compute all pair shortest path for any one of the following standard network



ARPA Network



NSF Network



NATIONAL NETWORK