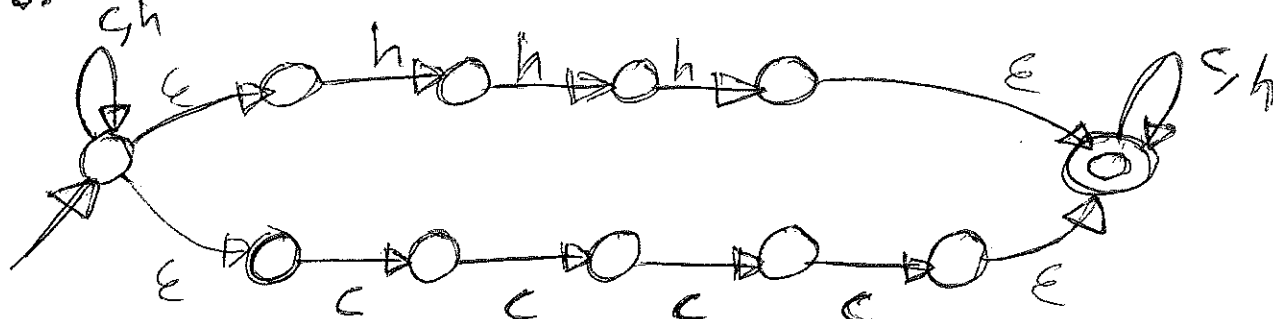


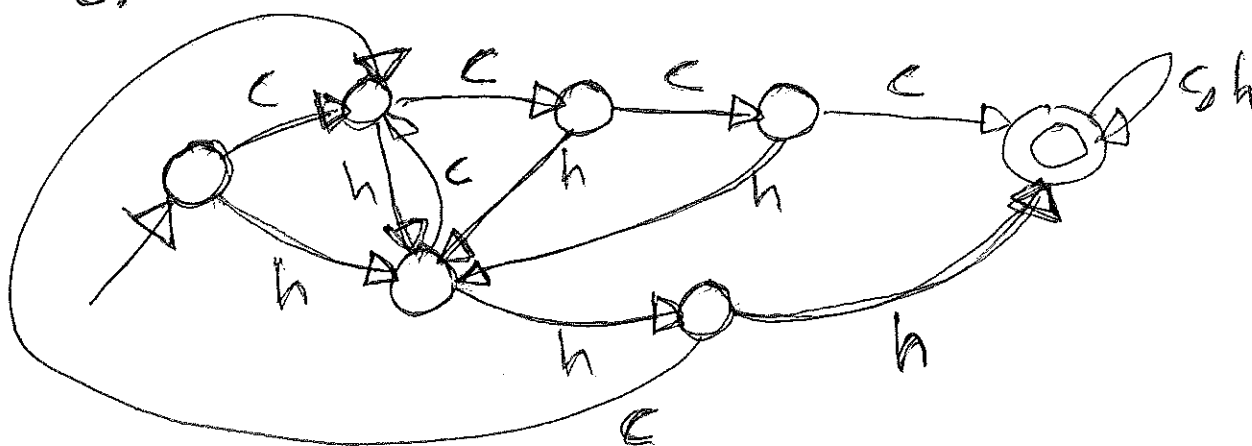
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1. a. $(c|h)^*(hhh|cccc)(c|h)^*$

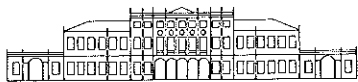
b.



c.



d. $\{a^n b^n a^n \mid n \geq 0\}$



5. a. $S \rightarrow S1 \mid AB2 \mid AB3$
 $A \rightarrow A4 \mid 5$
 $B \rightarrow B6 \mid \epsilon$

Left recursion:

$A \rightarrow A4 \quad B \rightarrow B6$

$S \rightarrow S1$

Left factoring problems:

$S \rightarrow AB2 \mid AB3$

New grammar:

$S \rightarrow ABS_2S_3$

$S_2 \rightarrow 2 \mid 3$

$S_3 \rightarrow 1S_3 \mid \epsilon$

$A \rightarrow 5A_2$

$A_2 \rightarrow 4A_2 \mid \epsilon$

$B \rightarrow 6B \mid \epsilon$

```
func S() {  
  A(); B(); S2(); S3();  
}
```

```
func S2() {  
  if (token == 2)  
    scan();  
  else if (token == 3)  
    scan();  
  else  
    error(...);  
  end if  
}
```

```
func S3() {  
  if (token == 1)  
    scan();  
    S3();  
  end if;  
}
```

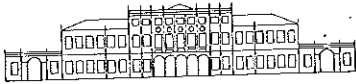
```
func A() {  
  if (token == 5)  
    scan();  
    A2();  
  else  
    error(...);  
  end if  
}
```

```
func A2() {  
  if (token == 4)  
    scan(); A2();  
  end if;  
}
```

```
func B() {  
  if (token == 6)  
    scan(); B();  
  end if  
}
```

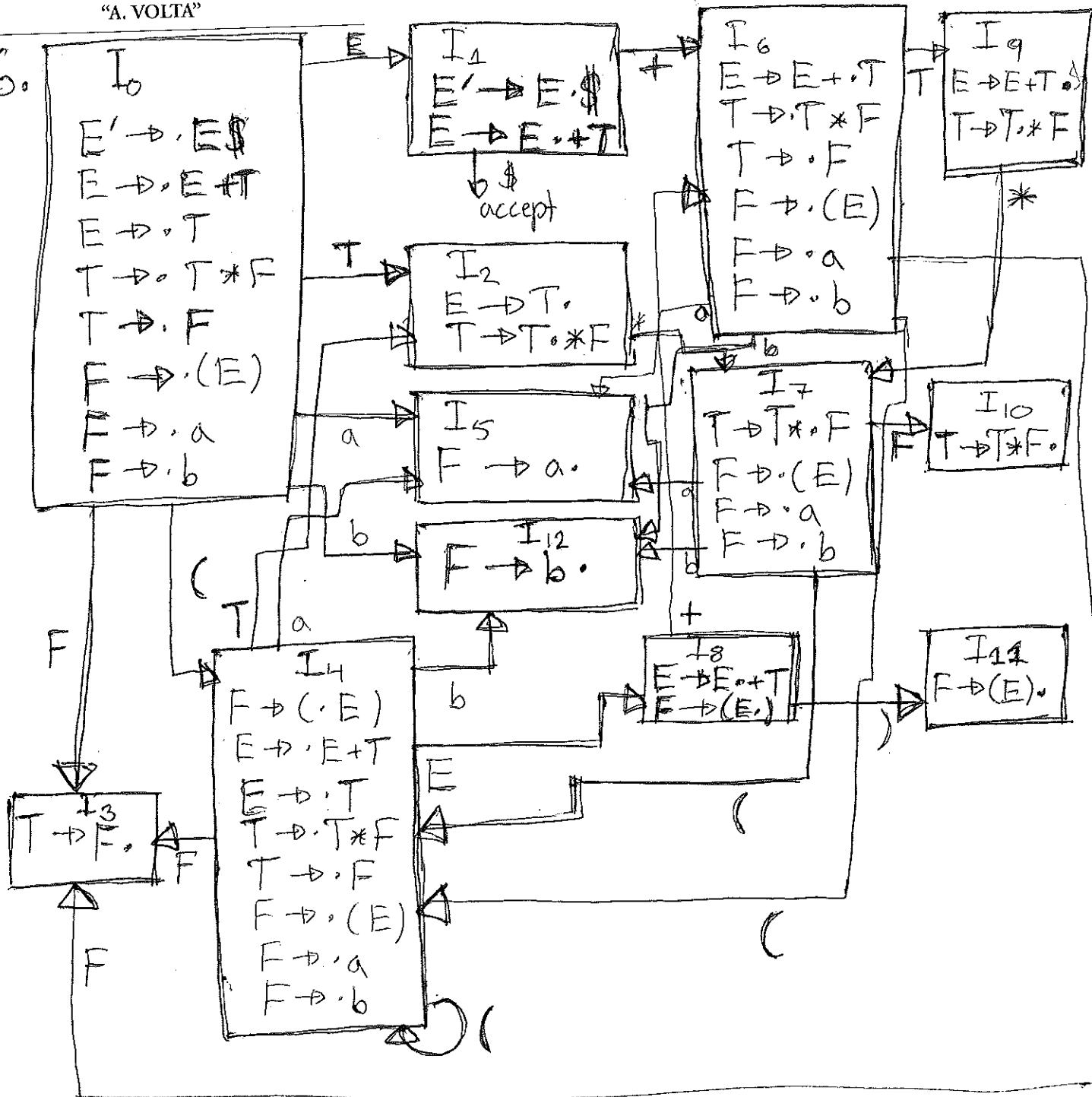
```
func main() {  
  scan();  
  S();  
  if (not eof)  
    error(...);  
  end if;  
}
```

b. The program stack that is used in the different function calls.



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6.



Shift-reduce conflicts in I_2 and I_9 .

$FOLLOW(E') = \{ \$ \}$

$FOLLOW(E) = \{ \$,), + \}$

$FOLLOW(T) = \{ \$,), *, + \}$

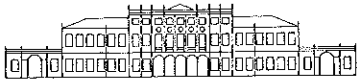
$FOLLOW(F) = \{ \$,), +, * \}$

But $*$ not in $FOLLOW(E)$

Thus $SLR(1)$.


$$\begin{aligned} (5) & F \rightarrow (E) \\ (6) & F \vdash a \\ (7) & F \vdash b \end{aligned}$$

0
 0a5
 0F3
 0T2
 0T2*7
 0T2*7b12
 0T2*7F10
 0T2
 0E1
 0E1+6
 0E1+6b12
 0E1+6F3
 0E1+6T9
 0E1+6T9*7
 0E1+6T9*7(4
 0E1+6T9*7(4a5
 0E1+6T9*7(4F3
 0E1+6T9*7(4T2
 0E1+6T9*7(4E8
 0E1+6T9*7(4E8+6
 0E1+6T9*7(4E8+6a5
 0E1+6T9*7(4E8+6F3
 0E1+6T9*7(4E8+6T9
 0E1+6T9*7(4E8
 0E1+6T9*7(4E8)11
 0E1+6T9*7F10
 0E1+6T9
 0E1
 ACCEPT!



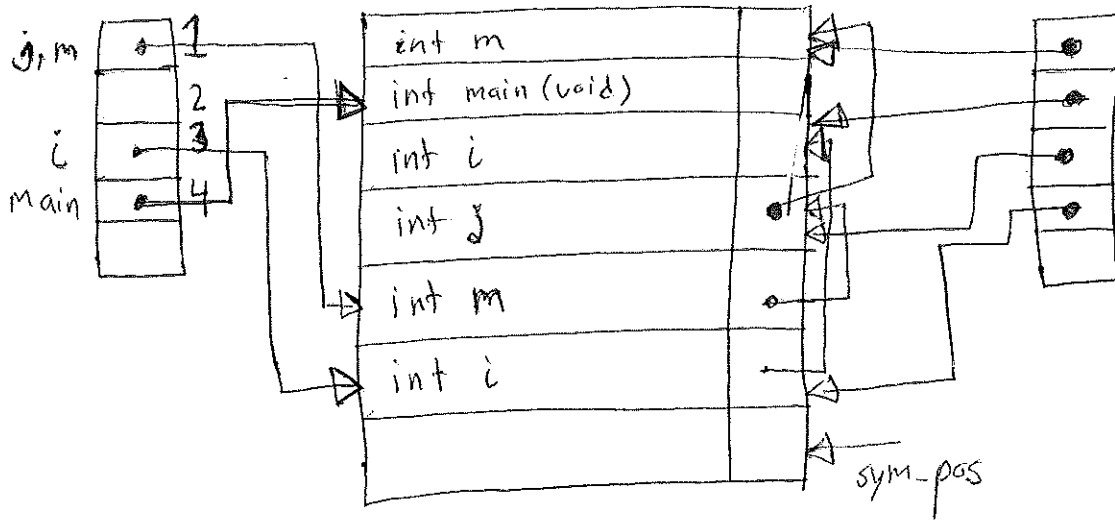
Symbol table

7.

Hash table

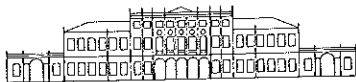
Hashlink

Block table



For looking up m in $i = m * 2$ we simply hash m , get an entry to the hash table, follow the link and we find the m we were looking for. If we hadn't found correct m we would have had to follow hash/back links.

Use the block table to get to the beginning of the current block and then enumerate the variables from there.



8. Rewrite the grammar to:

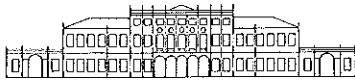
$\langle \text{loop} \rangle ::= \text{for } \langle \text{id-range} \rangle \text{ loop } \langle \text{statement-list} \rangle$
 $\text{end for};$

$\langle \text{id-range} \rangle ::= \langle \text{id} \rangle \text{ in } \langle \text{start-expr} \rangle : \langle \text{step-expr} \rangle : \langle \text{end-expr} \rangle;$

$\langle \text{loop} \rangle ::= \text{for } \langle \text{id-range} \rangle \text{ loop } \langle \text{statement-list} \rangle \text{ end for};$
{
 GEN(JMP, $\langle \text{id-range} \rangle.\text{addr}$, 0, 0);
 QUADRUPLE[$\langle \text{id-range} \rangle.\text{quad}$][1] = current-quad + 1; }
}

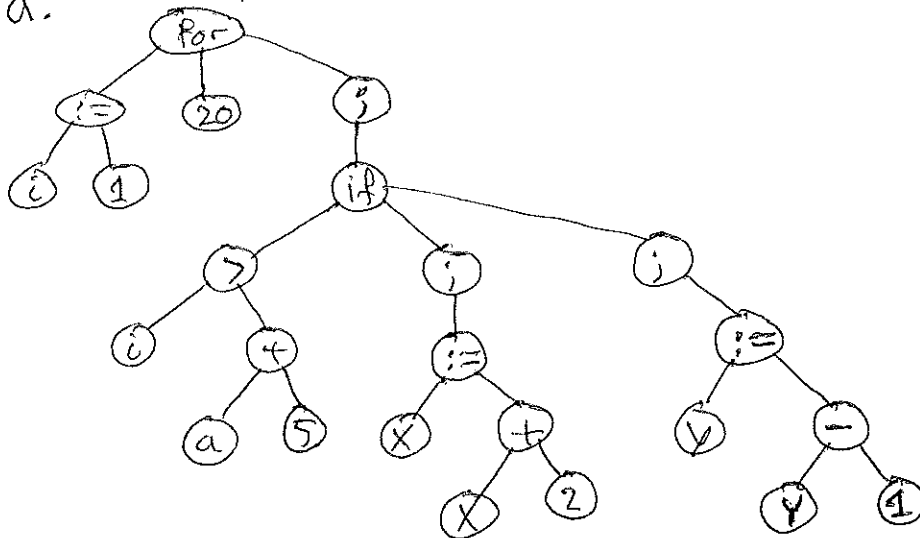
$\langle \text{id-range} \rangle ::= \langle \text{id} \rangle \text{ in } \langle \text{start-expr} \rangle : \langle \text{step-expr} \rangle : \langle \text{end-expr} \rangle$
{
 GEN(ASSIGN, $\langle \text{start-expr} \rangle.\text{addr}$, 0, $\langle \text{id} \rangle.\text{addr}$);
 int temp = current-quad + 1;
 GEN(JMP, 0, 0, 0);
 $\langle \text{id-range} \rangle.\text{addr} = \text{current-quad} + 1;$ // We can make use of the addr field as well
 int temp-var = gentemp();
 GEN(ADD, $\langle \text{id} \rangle.\text{addr}$, $\langle \text{step-expr} \rangle.\text{addr}$, temp-var);
 GEN(ASSIGN, temp-var, 0, $\langle \text{id} \rangle.\text{addr}$);
 QUADRUPLE[temp][1] = current-quad + 1;
 temp-var = gentemp();
 GEN(GTEQ, $\langle \text{end-expr} \rangle.\text{addr}$, $\langle \text{id} \rangle.\text{addr}$, temp-var);
 $\langle \text{id-range} \rangle.\text{quad} = \text{current-quad} + 1;$
 GEN(JMPF, 0, temp-var, 0);
}

$\langle \text{id-range} \rangle$ will be reduced first, then $\langle \text{loop} \rangle$
(in a bottom-up parser).



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9. a. Abstract Syntax Tree



Quadruples

1.	qassign	1	-	i
2.	qlabel	1	-	-
3.	qgteq	20	i	t ₁
4.	qjmp	4	t ₁	-
5.	qadd	a	5	t ₂
6.	qgt	i	t ₂	t ₃
7.	qjmp	2	t ₃	-
8.	qadd	x	2	t ₄
9.	qassign	t ₄	-	x
10.	qjmp	3	-	-
11.	qlabel	2	-	-
12.	qsub	y	1	t ₅
13.	qassign	t ₅	-	y
14.	qlabel	3	-	-
15.	qadd	i	1	t ₆
16.	qassign	t ₆	-	i
17.	qjmp	1	-	-
18.	qlabel	4	-	-

Postfix code

i 1 :=
 L1 20 i >= L4 JEQZ
 i a 5 + > L2 JEQZ
 x x 2 + := L3 JMP
 L2 y y 1 - :=
 L3 i i 1 + := L1 JMP
 L4

1 goto L2
 2 L1: x := x + 1
 3 L2: x := x + 1
 4 L3: x := x + 1
 5 if x = 1 then goto L1
 6 if x = 2 then goto L3
 7 if x = 3 then goto L5
 8 L4: x := x + 1
 9 L5: x := x + 1
 10. if x = 4 then goto L4.

