

Rectangular maximum volume

Scoltech

5 GOSLINGS

Our Plan

- [https://github.com/ArkadiyAliev/
Rectangular-maximum-volume](https://github.com/ArkadiyAliev/Rectangular-maximum-volume)
- Problem statement
- Known algorithms
- Speed comparison
- Quality comparison
- Application in low-rank approximation
- Contribution of each team member

Problem statement

Volume of a matrix

Let $A \in \mathbb{C}^{m \times n}$, $m \geq n$. The volume of A is defined as

$$\text{vol}(A) = \sqrt{\det(A^* A)}$$

Problem statement

Let $A \in \mathbb{C}^{m \times n}$, $m \geq n, r \geq n$. Our goal is to find a submatrix $\hat{A} \in \mathbb{C}^{r \times n}$ with largest volume.

Dominant submatrix

Let $A \in \mathbb{C}^{m \times n}$, $m \geq n, r \geq n$ be of a full column rank. $\hat{A} \in \mathbb{C}^{r \times n}$ is a submatrix of A . \hat{A} is called dominant submatrix of a matrix A , if the volume of \hat{A} does not increase by swapping one of its rows with another row from A .

In general, finding the maximum volume submatrix is a NP-complete problem.

maxvol

- ① Input: $A \in \mathbb{R}^{M \times r}, M > r$
- ② Output: $r \times r$ matrix — approximate solution of square maxvol problem.
- ③ Complexity: $O(Mr^2)$

rect-maxvol

- ① Input: $A \in \mathbb{R}^{M \times r}, n, M \geq n > r$. r indices of rows of approximate solution of square maxvol problem.
- ② Output: $n \times r$ matrix, — approximate solution of rectangular maxvol problem.
- ③ Complexity: $O(M(n^2 - r^2))$

pre-maxvol

- ① Input: $A \in \mathbb{R}^{M \times n}$, the required rank r .
- ② Output: set of column indices $\mathcal{I}, |\mathcal{I}| = r$, corresponding to the submatrix, whose volume differs from the maximum by no more than $r!$ times.
- ③ Complexity: $O(Mnr)$

DominantC

- ① Input: $A \in \mathbb{R}^{M \times r}$, set of indices $\mathcal{I}, |\mathcal{I}| = n$ to update.
- ② Output: set of row indices \mathcal{I}_1 , which corresponds to a dominant submatrix $n \times r$.
- ③ Complexity: $O(Mn \cdot \text{iter})$ ($O(Mnr \cdot \ln(r))$ after pre-maxvol).
Here iter is number of switches.

maxvol2

- ① Input: $A \in \mathbb{R}^{M \times r}$, set of indices $\mathcal{I}, |\mathcal{I}| = r$ to update by adding rows, $n \geq r \in \mathbb{N}$.
- ② Output: set of row indices $\mathcal{I}_1, |\mathcal{I}_1| = n$, which corresponds to a submatrix $n \times r$ with indices chosen greedily to maximize the volume.
- ③ Complexity: $O(M(n^2 - r^2))$

maxvol2-householder

- ① Input: $A \in \mathbb{R}^{M \times r}$, set of indices $\mathcal{I}, |\mathcal{I}| = r$ to update by adding rows, $n \geq r \in \mathbb{N}$.
- ② Output: set of row indices $\mathcal{I}_1, |\mathcal{I}_1| = n$, which corresponds to a submatrix $n \times r$ with indices chosen greedily to maximize the volume.
- ③ Complexity: $O(Mr(n - r))$

Wrong derivation of formulas in (Osinsky A.I., 2019)

Our task is to calculate C and C' on the basis of C_0 . The expression (9) shows that it is sufficient to find $(\hat{A}^* \hat{A})^{-1}$. Since

$$\hat{A}^* \hat{A} = \hat{A}_0^* \hat{A}_0 + aa^* = (\hat{A}_0^* \hat{A}_0) \left(I + (\hat{A}_0^* \hat{A}_0)^{-1} aa^* \right),$$

then for the inverse

$$\begin{aligned} (\hat{A}^* \hat{A})^{-1} &= (\hat{A}_0^* \hat{A}_0)^{-1} \left(I + (\hat{A}_0^* \hat{A}_0)^{-1} aa^* \right)^{-1} \\ &= (\hat{A}_0^* \hat{A}_0)^{-1} \left(I - \frac{(\hat{A}_0^* \hat{A}_0)^{-1} aa^*}{1 + a^* (\hat{A}_0^* \hat{A}_0)^{-1} a} \right). \end{aligned} \quad (10)$$

The expression in the brackets can be simplified by introducing the notation $c^* = C_{i,:}$. Indeed,

$$\begin{aligned} c^* &= C_{0i,:} = A_{i,:} \hat{A}_0^+ = a^* (\hat{A}_0^* \hat{A}_0)^{-1} \hat{A}_0^*, \\ c^* c &= a^* (\hat{A}_0^* \hat{A}_0)^{-1} \hat{A}_0^* \hat{A}_0 (\hat{A}_0^* \hat{A}_0)^{-1} a = a^* (\hat{A}_0^* \hat{A}_0)^{-1} a. \end{aligned}$$

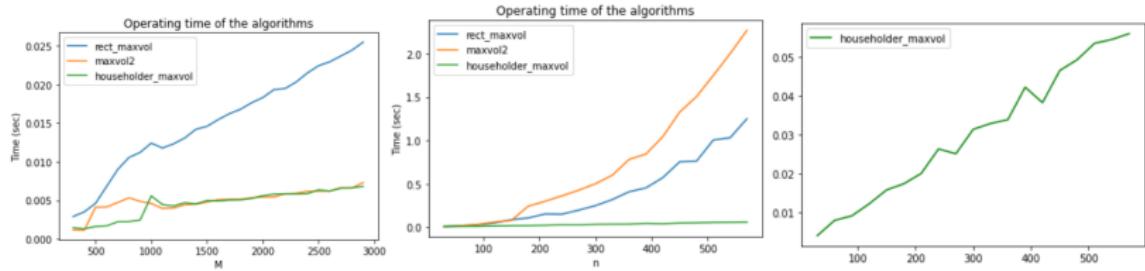
Also,

$$(\hat{A}_0^* \hat{A}_0)^{-1} aa^* = aa^* (\hat{A}_0^* \hat{A}_0)^{-1}.$$

Speed comparison

Input: $A \in \mathbb{R}^{M \times r}$, set of indices $\mathcal{I}, |\mathcal{I}| = r$ to update by adding rows, $n \geq r \in \mathbb{N}$.

Output: submatrix $n \times r$.



(a) Dependence on M ,
 $r = 10, n = 20$

(b) Dependence on n ,
 $r = 20, M = 600$

(c) Dependence on n ,
 $r = 20, M = 600$

Figure: Speed comparison for rect-maxvol, maxvol2, maxvol2-householder

Quality comparison

Input: $A \in \mathbb{R}^{M \times r}, M > r$

Output: $r \times r$ matrix — approximate solution of square maxvol problem.

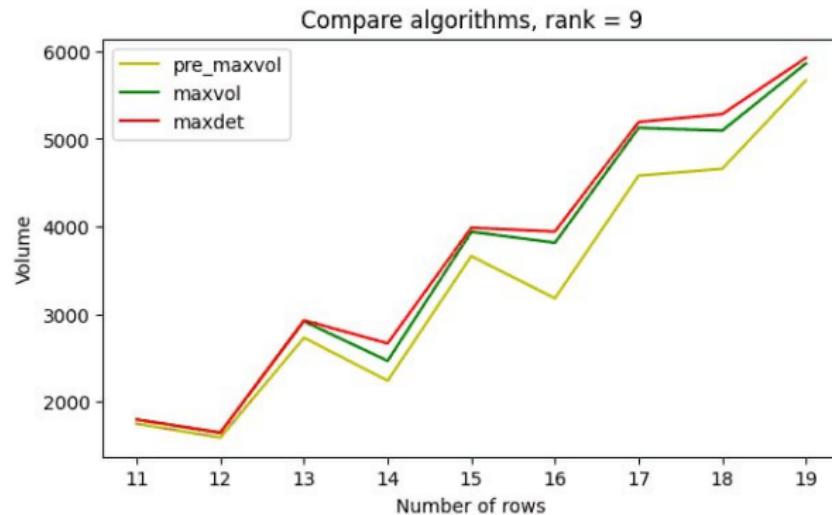


Figure: Quality comparison for premaxvol, maxvol and optimal algorithm ($r = 9$)

Quality comparison

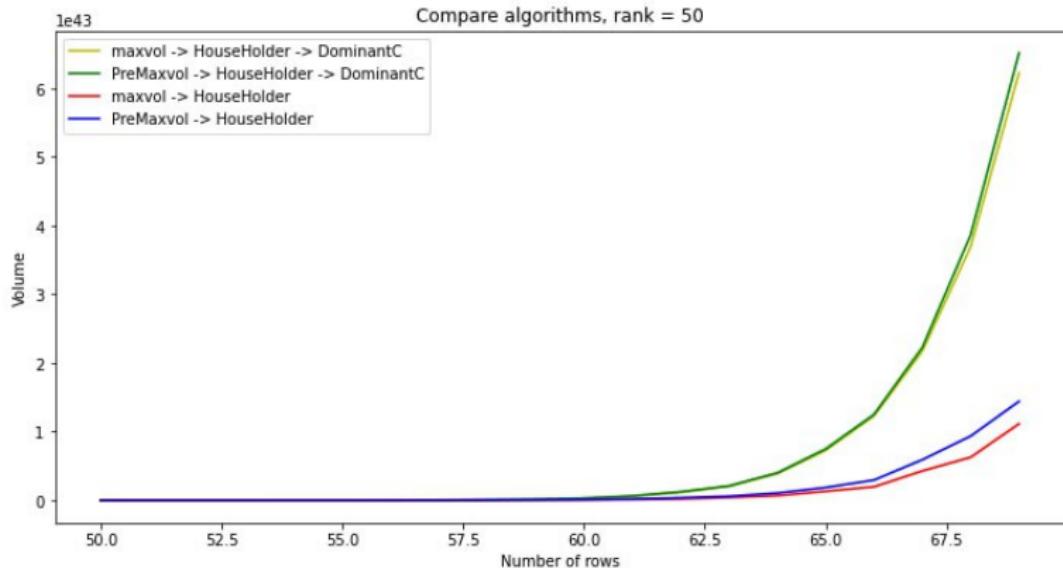


Figure: Input: Matrix $A \in \mathbb{R}^{200 \times 50}$ and r - desired number of rows.

We see that adding DominantC improves the volume of output submatrix dramatically.

Quality comparison

$A \in \mathbb{R}^{N \times 50}$, desired number of rows 100.

Output: 100×50 matrix — approximate solution of maxvol problem.

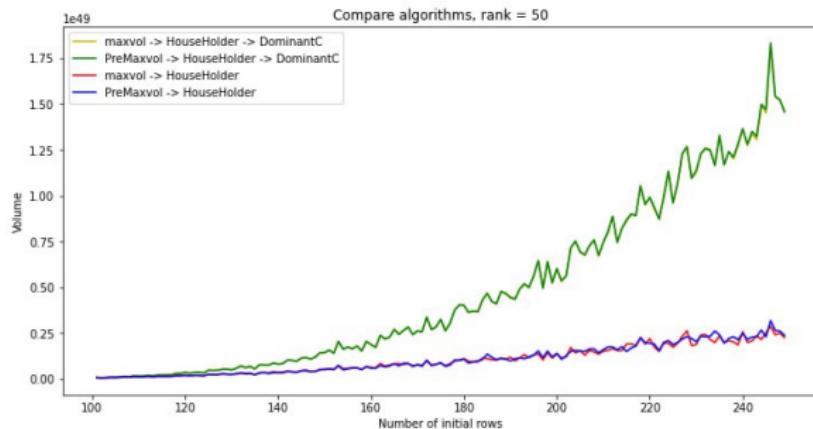


Figure: Quality comparison for different numbers of initial rows.

Application in low-rank approximation

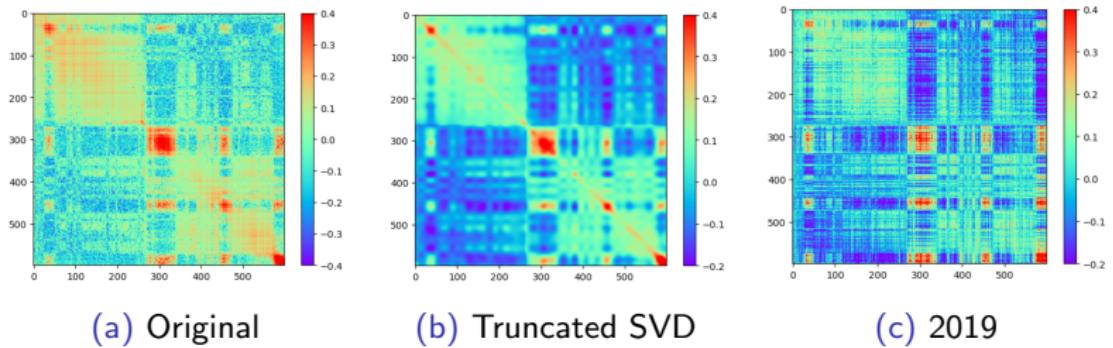
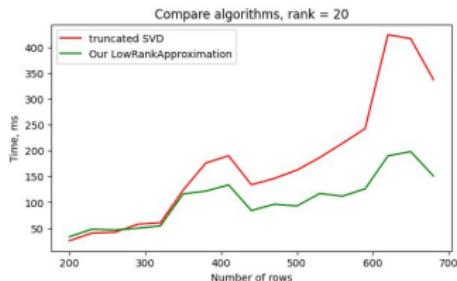


Figure: Truncated SVD vs (Osinsky A.I., 2019)

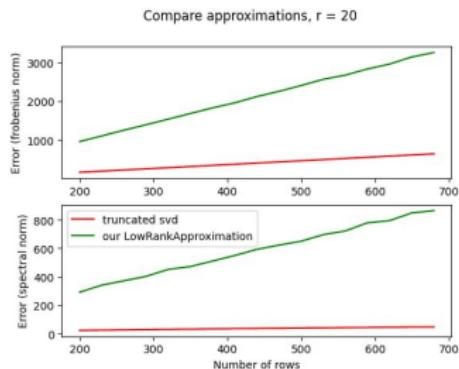
Application in low-rank approximation

Input: square matrix of size $N \times N$.

Output: low-rank approximation of this matrix by our algorithm.



(a) Speed comparison



(b) Error comparison

Contribution of each team member

- Arkadiy: LowRankApproximation, pre-maxvol, HouseHolder, DominantC implementation.
- Vsevolod: Correction of derivation of formulas, maxvol2, HouseHolder, DominantC implementation.
- Roman: maxvol, rect-maxvol implementation; numerical experiments for maxvol-rect vs maxvol2 vs HouseHolder
- Yuriy: numerical experiments for maxvol → HouseHolder → DominantC vs PreMaxvol → HouseHolder → DominantC vs maxvol→HouseHolder vs PreMaxvol → HouseHolder
- Mikhail: suggested project idea, numerical experiments for LowRankApproximation vs Truncated SVD, premaxvol vs maxvol vs random r rows vs overkill

¿PREGUNTAS?

