CMPT 295 Assignment 5 (2%)

- 1. [7 marks] Floating-Point Integers Let S be the set of all 32-bit IEEE floating-point values that are integers. Answer the following questions about S.
- (a) [1 mark] How many elements are in S?

Answer: There are 3 elements in S.

(b) [1 mark] What's the largest odd value of S?

Answer: It's 16,777,215.

(c) [1 mark] What are the largest two values of S? Write your answers as exact expressions.

Answer: They're $(2-2^{-23})*2^{127}$ and $(2-2^{-22})*2^{127}$.

(d) [1 mark] How many consecutive integers are in S? You shouldn't count both 0 and -0 in your total.

Answer: It's 225+1.

(e) [1 mark] The number 264 is in S. What are its nearest neighbours? In other words, what two values in S are closest to 264?

Answer: They're 263 and 265.

(f) [2 marks] Determine the number of values in S which are representable using 32-bit unsigned.

Answer: It's 1+2*224*(127-24+1).

- 2. [2 marks] Floating-Point Addition Add the following pairs of 32-bit IEEE floating-point numbers. Show all your work.
- 0x43938000 + 0xc2280000.

Answer:

 $X_e = 10000111_2 - 127_{10} = 135 - 127 = 8$

 $Y_e = 10000100_2 - 127_{10} = 132 - 127 = 5$

- \Rightarrow X=1.00100111₂*2⁸₁₀
- \Rightarrow Y=-1.0101₂*2⁵=-0.0010101₂*2⁸₁₀
- $\Rightarrow X+Y=2^{8}_{10}*(1.00100111_{2}-0.0010101_{2})=$ $2^{8}_{10}*0.111111101_{2}=2^{7}_{10}*1.11111101_{2}$
- 0x3f2aaaaa + 0x3eaaaaab.

Answer:

 $X_e = 011111110_2 - 127_{10} = 126 - 127 = -1$

 $Y_e = 011111101_2 - 127_{10} = 125 - 127 = -2$

- \Rightarrow X=1.010101010101010101010101₂*2⁻¹₁₀
- \Rightarrow Y=1.0101010101010101010101011₂ *2⁻¹
 ²=0.101010101010101010101010101011₂*2⁻¹₁₀
- \Rightarrow X+Y=2⁻¹₁₀*(1.0101010101010101010101₂-0.1010101010101010101010101₂)
- $\Rightarrow =2^{-1}_{10}*0.1010101010101010101010101_2 =2^{-1}_{10}*1.010101010101010101010101_2$
- ⇒ =0 01111101 01010101010101010101001=0x3eaaaaa9

3. [11 marks] Adding Positive Floating-Point Numbers #if only Q is empty, go to dequeue F el: (c) [7 marks] movq %rbx, %rcx sum_float.s: \$8, %rcx subq .globl sum float cmpg %rcx, %rdx sum_float: endloop je %rbp #if F is empty and Q has only 1 element, go to push end the loop leag (%rdi, %rsi, 4), %rbp compare1g jmp # endptr <- F + n #if Q has more than 1 element, dequeue Q xorps %xmm1, %xmm1 compare1: xorps %xmm2, %xmm2 %xmm0, %xmm0 xorps #clear xmm1 and xmm2. xmm1 is x. xmm2 is y. xorps %xmm3, %xmm3 -8(%rsp), %rdx leag movss (%rdi), %xmm0 #create a queue and let %rdx be the tail pointer of the queue movss -8(%rbx), %xmm3 movq %rdx, %rbx ucomiss %xmm0, %xmm3 #let %rbx be the head pointer of the queue #use xmm0 and xmm3 to temparily store the value of heads of two queues for comparison #now, %rdi and %rbp is the head and tail pointers of queue F to sum compare1 jae #%rbx and %rdx is the head and tail pointers of compare1g jmp queue Q to store sums #dequeue the smaller one jmp gool compare1g: #start the loop subq \$8, %rbx loop: movss (%rbx), %xmm1 cmpq %rdi, %rbp jmp second ile el #Assign head of Q to x and dequeue Q and go to #when F has all been dequeued, go to el for do the second dequeue advanced check compare11: cmpg %rbx, %rdx addss (%rdi), %xmm1 #if F is not empty, check whether Q is empty add \$4, %rdi ine compare1 second jmp #when two queues are not empty, go to #Assign head of F to x and dequeue F and go to compare the head of two queues do the second dequeue jmp compare1

second:

cmpq %rdi, %rbp

jle compare2g

#if F is empty, dequeue Q

cmpq %rbx, %rdx

jne compare2

#when two queues are not empty, go to compare the head of two queues

jmp compare2l

#if only Q is empty, go to dequeue F

compare2:

xorps %xmm0, %xmm0

xorps %xmm3, %xmm3

movss (%rdi), %xmm0

movss -8(%rbx), %xmm3

#use xmm0 and xmm3 to temparily store the value of heads of two queues for comparison

ucomiss %xmm0, %xmm3

jae compare2l

jmp compare2g

#dequeue the smaller one

compare2g:

subq \$8, %rbx

movss (%rbx), %xmm2

imp adding

#Assign head of Q to y and dequeue Q and go to sum x and y

compare2l:

addss (%rdi), %xmm2

add \$4, %rdi

jmp adding

#Assign head of F to y and dequeue F and go to sum x and y

adding:

addss %xmm1, %xmm2

subq \$8, %rdx

movss %xmm2, (%rdx)

#adding x and y and enqueue to Q

xorps %xmm1, %xmm1

xorps %xmm2, %xmm2

imp loop

#clear x and y and go to loop

endloop:

movss -8(%rbx), %xmm0

#let %xmm0 to store the final result at the head

of Q

pop %rbp

ret

#return the answer and finish the function

(d) [2 BONUS marks] Hardcopy: Produce a [mathematical] justification/proof as to why the elements of Q must certainly be in increasing order.

Answer:

In floating point addition like in question 2, we find out that we need to convert the two values to add in the same exponent base then shift the fraction elements. The fraction part has 23 bits in IEEE 32-bit floating point system which means the smallest fraction is 2⁻²³. Therefore, if the exponent bases of two values are the same, the fraction part after shifting and addition less than 2⁻²³ will be cut off. If a number is less than 1/2⁻²³ of the number to be added, addition has no effects on the number to be added even if we take infinity times of this addition. However, if we take adding in the increasing order, we can sum the little numbers together first to let their result be greater than 1/2-23 of the number to be added. Then the addition of this result and the big number will have effect on the final result.