



# A TLA<sup>+</sup> and PLUSCAL Overview

Some koans and exercises

---

Rafael Luque

September 6, 2019

OSOCO

# Table of contents

1. Introduction
2. **TLA<sup>+</sup>** and **PLUSCAL** Semantics
3. **PLUSCAL** Koans
4. Test the Water
5. Workshop of Spec Writing

# Introduction

---

# TLA<sup>+</sup> and PLUSCAL Semantics

---

# About PLUSCAL

- Developed by Lamport in 2009 to make **TLA<sup>+</sup>** more accessible to programmers.
- **PLUSCAL** provides a pseudocode-like structure on top of **TLA<sup>+</sup>**.
- Adds additional syntax: **assignments** and **labels**.
- **PLUSCAL** is language that compiles down to **TLA<sup>+</sup>**.

# Spec layout

```
---- ❶ MODULE example ❷ ----  
EXTENDS Integers ❸  
  
(* --algorithm wire ❹  
  
    variables ❺  
        people = { "bob", "alice" },  
        acc = [ alice |-> 5, bob |-> 5 ];  
  
    begin ❻  
        skip;  
    end algorithm; ❹ *)  
===== ❶
```

# Spec layout

- ❶ **TLA<sup>+</sup>** specs must start with at least four - at each side of **MODULE** and four = at the end.
- ❷ The module name must be the same as filename.
- ❸ **EXTENDS** is the import keyword.
- ❹ **(\*...\*)** is the block comment. **PLUSCAL** spec starts with `--algorithm <name>` and ends with `end algorithm`.
- ❺ Initialization of variables.
- ❻ Where the algorithm is implemented.

Four basic values in **TLA<sup>+</sup>**:

**String** Must be written en double quotes.

**Integer** Floats are not supported.

**Boolean** Written as TRUE and FALSE.

**Model value** A kind of symbol value.



# Standard operators

Operator	Meaning	Example
<code>x = y</code>	Equals	<code>&gt;&gt; 1 = 2</code> <code>FALSE</code>
<code>x /= y</code>	Not Equals	<code>&gt;&gt; 1 /= 2</code> <code>TRUE</code>
<code>x /\ y</code>	And	<code>&gt;&gt; TRUE /\ FALSE</code> <code>FALSE</code>
<code>x \/ y</code>	Or	<code>&gt;&gt; TRUE \/ FALSE</code> <code>TRUE</code>
<code>~x</code>	Not	<code>&gt;&gt; ~TRUE</code> <code>FALSE</code>
<code>x := y</code>	Assignment	PlusCal only

# Arithmetic Operators

- If you `EXTENDS Integers` get the arithmetic operators: `+`, `-`, `%` and `*`.
- Decimal division is not supported, only the integer division: `\div`.
- You also get the range operator `..` where `a..b` is the set `{a, a+1, ..., b-1, b}`.

TLA<sup>+</sup> has four complex types:

- Sets.
- Tuples or sequences.
- Structures.
- Functions.

## Sets

Unordered collections of elements of the same type.

`{1, 2, 42}`

`{{TRUE}, {FALSE, TRUE}, {}}`

# Set Operators

Operator	Meaning	Example
<code>x \in set</code>	Is member of	<pre>&gt;&gt; 1 \in 1..3 TRUE</pre>
<code>x \notin set</code>	Is not member of	<pre>&gt;&gt; 1 \notin 1..2 FALSE</pre>
<code>set1 \subseq set2</code>	Is subset of	<pre>&gt;&gt; {1, 2} \subseq {1, 2, 3} TRUE</pre>
<code>set1 \union set2</code>	Union	<pre>&gt;&gt; {1, 2} \union (2..3) {1, 2, 3}</pre>
<code>set1 \intersect set2</code>	Intersection	<pre>&gt;&gt; {1, 2} \intersect (2..3) {2}</pre>
<code>set1 \ set2</code>	Difference	<pre>&gt;&gt; {1, 2} \ (2..3) {1}</pre>
<code>Cardinality(set)</code>	Cardinality (requires EXTENDS FiniteSets)	<pre>&gt;&gt; Cardinality({1, 2}) 2</pre>

# Set Transformations

## Filter sets

```
{x \in set: conditional}
```

```
>> {x \in 1..3: x > 2}  
{3}
```

## Map sets

```
{expression: x \in set}
```

```
>> {x * 3: x \in 1..3}  
{3, 6, 9}
```

# Tuples or Sequences

## Tuples or Sequences

Ordered collections of elements with the index starting at 1.

```
tuple := <<1, TRUE, {1, 2}>>;
```

```
>> tuple[1]
```

```
1
```

```
>> tuple[3]
```

```
{1, 2}
```

# Sequence Operators

If you **EXTENDS Sequences** you get the following additional operators:

Operator	Meaning	Example
Head(seq)	Head	>> Head(<<1, 2>> 1
Tail(seq)	Tail	>> Tail(<<1, 2, 3>> <<2, 3>>
Append(seq, x)	Append	>> Append(<<1, 2>>, 3 <<1, 2, 3>>
seq1 \o seq2	Combine	>> <<1, 2>> \o <<3>> <<1, 2, 3>>
Len(seq)	Length	>> Len(<<1, 2>> 2



# Structures or Structs

## Structures or Structs

A **map** of strings to values.

```
>> [a |-> 1, b |-> <<1, {2, 3}>>].b  
    <<1, {2, 3}>>
```

# Assignments

Assign an **existing** variable to a value with `:=`.

## Rule of thumb

If it's the first time you're using the variable, `=` is initialization.  
Every other time, `=` is equality and `:=` is assignment.

## assert

`assert TRUE` does nothing, `assert FALSE` raises an error.

In order to use assertions you need to add `EXTENDS TLC` to the spec.

- A **no-op**.
- To fill parts of the spec that we haven't filled out yet or conditionals that don't update anything.

```
if condition1 then
  body1
elsif condition2 then
  body2
else
  body3
end if;
```

# while loop

```
while condition do  
  body  
end while;
```

# macros

To avoid duplications in your specs you can add macros before the begin of the algorithm:

```
macro name(arg1, arg2) begin
  \* macro's body
end macro;

begin
  name(x, y);
end algorithm;
```

## Macros limitations

You can place assignments, assertions, and if statements in macros, but not while loops. You also cannot assign to any variable more than once.

We need a way to specify not just one setup, but an entire space of setups to check our specifications.



## Multiple Starting States ii

We initialize variables with `=`, but we can also initialize them with `\in`. TLC will try running the algorithm with any possible element in the set:

```
(* --algorithm in
variables x \in 1..3};

begin
  assert x < 3;
end algorithm; *)
```

### BOOLEAN set

TLA<sup>+</sup> defines **BOOLEAN** as the set {TRUE, FALSE}. This can be useful if you have a variable `isReady` \in **BOOLEAN**.

### SUBSET

SUBSET is the power set, or the set of all subsets.

```
>> SUBSET 1..2  
    {{}}, {1}, {2}, {1, 2}}
```

## Multiple Starting States v

`\X`

`set1 \X set2` is the set of all tuples where the first element is in `set1` and the second element in `set2`.

```
>> (1..2) \X BOOLEAN  
    {<<1, TRUE>>, <<1, FALSE>>, <<2, TRUE>>, <<2, FALSE>>}
```

## Multiple Starting States vi

[key: set]

If  $x \in [\text{key: set}]$ , then  $x$  is an structure where the value of  $\text{key}$  is an element in  $\text{set}$ .

```
>> [a: (1..2), b: BOOLEAN]
{[a |-> 1, b |-> TRUE], [a |-> 2, b |-> TRUE],
 [a |-> 1, b |-> FALSE], [a |-> 2, b |-> FALSE]}
```

## Beware of state explosion i

```
variables
  capacity = [trash |-> 10, recycle |-> 10],
  items = <<
    [type |-> "trash", size |-> 5],
    [type |-> "recycle", size |-> 3],
    [type |-> "recycle", size |-> 4],
    [type |-> "trash", size |-> 2]
  >>;
```

## Beware of state explosion ii

```
variables
  capacity \in [trash: 1..10, recycle: 1..10],
  item \in [type: {"trash", "recycle"}, size: 1..6],
  items \in item \X item \X item \X item;
```

1 state vs.  $10 \times 10 \times (2 \times 6)^4 = 2,073,600$  states

For single process algorithms **PLUSCAL** provides two constructs to simulate nondeterminism:



```
either
  /* branch 1
or
  /* branch 2
/* ...
or
  /* branch n
end either;
```

- TLC will check all branches.
- No way to make one branch more likely than others.
- If all branches are *macro-valid*, we can place an either inside a macro.

# Simulating Nondeterminism iii

There are two ways to use the **with** statement:

```
with var = value do
  /* body
end with;
```

```
with var \in set do
  /* body
end with;
```

- The former creates a temporary variable, the second is nondeterministic.
- TLC will check what happens to all possible assignments of var to elements of set.
- with follows **macro rules**: no double assignments and no while loops.
- You can place with statements inside macros.

## PLUSCAL Koans

---

# Test the Water

---

# Workshop of Spec Writing

---

Questions?

# References i



P. Erdős.

**A selection of problems and results in combinatorics.**

In *Recent trends in combinatorics* (Matrahaza, 1995), pages 1–6.  
Cambridge Univ. Press, Cambridge, 1995.



R.L. Graham, D.E. Knuth, and O. Patashnik.

***Concrete mathematics.***

Addison-Wesley, Reading, MA, 1989.



George D. Greenwade.

**The Comprehensive Tex Archive Network (CTAN).**

*TUGBoat*, 14(3):342–351, 1993.



D.E. Knuth.

**Two notes on notation.**

*Amer. Math. Monthly*, 99:403–422, 1992.



H. Simpson.

**Proof of the Riemann Hypothesis.**

preprint (2003), available at

<http://www.math.drofnats.edu/riemann.ps>, 2003.