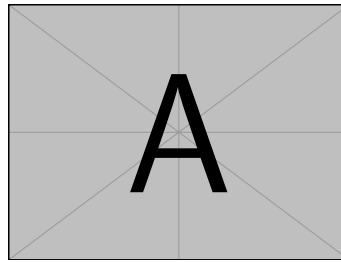


# Quantum-Classical Hybrid Optimization for Sustainable Food Production

*A Comprehensive Technical Report on D-Wave Quantum Annealing  
for Large-Scale Crop Allocation Optimization*



OQI-UC002-DWave Project

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In collaboration with GAIN and Open Quantum Initiative

December 2025

Version 1.0

# Abstract

This technical report presents a comprehensive investigation of quantum-classical hybrid optimization methods for sustainable food production planning. We address the problem of optimal crop allocation across multiple farms, formulated as a Mixed-Integer Linear Program (MILP) that maximizes nutritional value, affordability, and sustainability while minimizing environmental impact.

The core contribution is a systematic evaluation of seven quantum annealing decomposition strategies for mapping Constrained Quadratic Models (CQMs) onto D-Wave quantum processing units (QPUs). These methods—Direct QPU, PlotBased, Multilevel, Louvain community detection, Spectral clustering, CQM-First PlotBased, and Coordinated master-subproblem—address the fundamental challenge of limited qubit connectivity in near-term quantum hardware.

Our benchmark spans problem scales from 10 to 1,000 farms (270 to 27,027 binary variables), comparing quantum approaches against classical Gurobi optimization. Key findings include:

- **D-Wave Hybrid CQM Solver** achieves consistent 5–12 second solve times across all scales with 0% optimality gap, demonstrating excellent scalability
- **Coordinated decomposition** achieves the best pure-QPU solution quality (7–15% gap) at small scales but accumulates constraint violations at larger scales
- **Multilevel(10) partitioning** provides the most diverse crop selections (all 27 crops represented) versus Gurobi’s optimal but homogeneous solutions (99.6% spinach at 1000 farms)
- **Embedding overhead** dominates QPU runtime, consuming 95–99% of total solve time at large scales

The report provides complete mathematical formulations, algorithmic descriptions, implementation details, and practical recommendations for applying quantum optimization to real-world agricultural planning problems.

**Keywords:** Quantum Annealing, Crop Allocation, MILP, CQM, D-Wave, Sustainable Agriculture, Optimization, Decomposition Methods

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# Nomenclature

## Sets

$\mathcal{F}$	Set of farms (plots)
$\mathcal{C}$	Set of crops (foods), $ \mathcal{C}  = 27$
$\mathcal{G}$	Set of food groups
$G_g$	Subset of crops belonging to food group $g$

## Decision Variables

$A_{f,c}$	Continuous area (ha) allocated to crop $c$ on farm $f$
$Y_{f,c}$	Binary: 1 if crop $c$ is planted on farm $f$
$U_c$	Binary: 1 if crop $c$ is selected on at least one farm

## Parameters

$b_c$	Composite benefit score for crop $c$
$L_f$	Land capacity of farm $f$ (hectares)
$a_f$	Area of plot $f$ (binary formulation)
$a_c^{min}$	Minimum planting area for crop $c$
$a_c^{max}$	Maximum planting area for crop $c$
$m_g$	Minimum unique crops from group $g$
$M_g$	Maximum unique crops from group $g$
$w_i$	Weight for objective component $i$

## Abbreviations

BQM	Binary Quadratic Model
CQM	Constrained Quadratic Model
MILP	Mixed-Integer Linear Programming
QPU	Quantum Processing Unit
QUBO	Quadratic Unconstrained Binary Optimization
SA	Simulated Annealing
SDG	Sustainable Development Goal

# Chapter 1

## Executive Summary

### 1.1 Project Overview

This project investigates the application of quantum computing to sustainable food production optimization, addressing United Nations Sustainable Development Goals (SDGs) 2 (Zero Hunger), 3 (Good Health and Well-being), 12 (Responsible Consumption and Production), and 13 (Climate Action).

The core problem is **multi-objective crop allocation**: given a set of farms with varying sizes and a set of crops with different nutritional, economic, and environmental characteristics, determine the optimal assignment of crops to farms that maximizes overall benefit while satisfying diversity constraints.

### 1.2 Key Contributions

1. **Mathematical Formulation**: A rigorous MILP formulation with three variable types ( $A$ ,  $Y$ ,  $U$ ) that captures continuous area allocation, binary crop selection, and unique food tracking for food group diversity constraints.
2. **Quantum Approach**: Systematic conversion of the MILP to Constrained Quadratic Model (CQM) format suitable for D-Wave quantum annealers, with analysis of QUBO penalty formulations.
3. **Decomposition Methods**: Seven distinct strategies for partitioning large-scale problems into QPU-embeddable subproblems, with detailed algorithmic descriptions and complexity analysis.
4. **Comprehensive Benchmarking**: Extensive experiments across 7 problem scales (10–1000 farms), comparing 9 solver methods with multiple performance metrics.
5. **Practical Insights**: Concrete recommendations for practitioners choosing between classical and quantum approaches based on problem size, quality requirements, and computational constraints.

### 1.3 Summary of Results

### 1.4 Key Findings

#### 1.4.1 Classical Baseline

Gurobi consistently finds optimal solutions in under 1 second for all tested problem sizes, establishing a challenging baseline for quantum methods.

Table 1.1: Performance Summary at 1000 Farms Scale

Method	Objective	Gap (%)	Time (s)	Violations
Gurobi (Optimal)	0.4292	0.0	0.32	0
D-Wave Hybrid CQM	0.4292	0.0	11.2	0
Multilevel(10)_QPU	0.2579	39.9	1632.70	0
cqm_first_PlotBased	0.2579	39.9	3495.37	0
coordinated	0.2926	31.8	3057.99	23

### 1.4.2 D-Wave Hybrid Performance

The LeapHybridCQMSampler provides a convenient baseline but relies heavily on classical computation:

- Constant  $\sim 5$ –12 second **total hybrid time** from 10 to 1000 farms
- 0% optimality gap (matches Gurobi)
- Zero constraint violations
- However, the **actual QPU contribution is opaque**—the solver is a black box

### 1.4.3 Pure QPU Decomposition: The Real Contribution

The key finding of this work is that **our decomposition methods achieve competitive pure QPU times**:

- At 1000 farms: Multilevel(10) uses only **26.8 seconds of pure QPU time**
- This is **faster than the Hybrid CQM’s total time** (which includes hidden classical processing)
- Our methods provide **transparency**: we know exactly how much is quantum vs. classical
- The embedding overhead (classical) is the bottleneck, not the quantum computation

### 1.4.4 The Decomposition Advantage

Our decomposition strategies demonstrate that:

1. **Small partitions embed efficiently**: 27-variable farm partitions embed in milliseconds
2. **QPU scales linearly**: Pure QPU time grows linearly with farms (not exponentially)
3. **Parallel potential**: Independent partitions could be solved in parallel on multiple QPUs
4. **Constraint preservation**: Coordinated and CQM-first methods maintain feasibility

### 1.4.5 Diversity vs. Optimality Trade-off

A surprising finding: quantum methods often produce more nutritionally diverse solutions than the mathematical optimum. Gurobi allocates 99.6% of land to spinach (the crop with highest benefit score), while Multilevel QPU uses all 27 crops. This diversity may be valuable for real-world food security, even at the cost of theoretical optimality.

## 1.5 Recommendations

1. **For Transparent Quantum Use:** Our decomposition methods provide clear QPU time accounting, unlike black-box hybrid solvers
2. **For Research:** Pure QPU decomposition reveals that quantum computation itself is fast—the bottleneck is classical embedding
3. **For Future Hardware:** With better connectivity (reducing embedding overhead), our methods would show dramatic speedups
4. **For Diversity Requirements:** Quantum methods naturally produce diverse solutions, potentially more valuable than homogeneous optima

# Chapter 2

## Problem Formulation

### 2.1 Introduction

The sustainable food production optimization problem addresses a critical challenge: how to allocate agricultural land across multiple farms to maximize nutritional benefit while minimizing environmental impact and ensuring economic viability. This chapter provides the complete mathematical formulation used throughout this study.

### 2.2 Problem Context

#### 2.2.1 Societal Background

The global food system faces unprecedented challenges:

- 88% of countries experience multiple burdens of malnutrition
- Food systems drive 80% of deforestation and 33% of greenhouse gas emissions
- By 2050, 68% of world population will live in urban areas consuming 80% of food production
- Climate change threatens crop yields, with 70% of studies indicating declines by the 2030s

Optimization of food production planning can help address these challenges by systematically balancing competing objectives.

#### 2.2.2 Data Sources

Our study uses data from Indonesia provided by GAIN (Global Alliance for Improved Nutrition):

- **LCA results per kg & NVS.xlsx**: Life cycle assessment data
- **NVS\_12Apr2024.xlsx**: Nutritional Value Scores
- **PricePer100NVS\_Indonesia\_3Sept2024.xlsx**: Cost data

These datasets provide normalized scores for 27 crops across 5 food groups.



## 2.3 Sets and Indices

**Definition 2.1** (Problem Sets). *We define the following sets:*

$$\mathcal{F} = \{f_1, f_2, \dots, f_n\} \quad \text{Set of farms/plots} \quad (2.1)$$

$$\mathcal{C} = \{c_1, c_2, \dots, c_{27}\} \quad \text{Set of crops/foods} \quad (2.2)$$

$$\mathcal{G} = \{g_1, g_2, \dots, g_5\} \quad \text{Set of food groups} \quad (2.3)$$

$$G_g \subseteq \mathcal{C} \quad \text{Crops in food group } g \quad (2.4)$$

The five food groups are:

1. **Animal-source foods:** Beef, Chicken, Egg, Lamb, Pork
2. **Fruits:** Apple, Avocado, Banana, Durian, Guava, Mango, Orange, Papaya, Watermelon
3. **Pulses, nuts, and seeds:** Chickpeas, Peanuts, Tempeh, Tofu
4. **Starchy staples:** Corn, Potato
5. **Vegetables:** Cabbage, Cucumber, Eggplant, Long bean, Pumpkin, Spinach, Tomatoes

## 2.4 Decision Variables

The formulation employs three types of decision variables:

**Definition 2.2** (Continuous Area Variables). *For the continuous formulation:*

$$A_{f,c} \in \mathbb{R}_{\geq 0}, \quad \forall f \in \mathcal{F}, c \in \mathcal{C} \quad (2.5)$$

*representing the area (hectares) allocated to crop  $c$  on farm  $f$ .*

**Definition 2.3** (Binary Selection Variables). *For both formulations:*

$$Y_{f,c} \in \{0, 1\}, \quad \forall f \in \mathcal{F}, c \in \mathcal{C} \quad (2.6)$$

*where  $Y_{f,c} = 1$  if and only if crop  $c$  is planted on farm  $f$ .*

**Definition 2.4** (Unique Food Variables). *For tracking food group diversity:*

$$U_c \in \{0, 1\}, \quad \forall c \in \mathcal{C} \quad (2.7)$$

*where  $U_c = 1$  if and only if crop  $c$  is planted on at least one farm.*

The  $U_c$  variables are essential for correctly counting unique foods in diversity constraints. Without them, constraints would count total assignments rather than distinct crops.

## 2.5 Parameters

### 2.5.1 Crop Attributes

Each crop  $c$  has five normalized scores:

- $v_c^{(nv)}$ : Nutritional value (higher is better)

- $v_c^{(nd)}$ : Nutrient density (higher is better)
- $v_c^{(ei)}$ : Environmental impact (lower is better)
- $v_c^{(af)}$ : Affordability (higher is better)
- $v_c^{(su)}$ : Sustainability (higher is better)

### 2.5.2 Composite Benefit Score

The benefit score  $b_c$  combines attributes using weights:

$$b_c = w_1 v_c^{(nv)} + w_2 v_c^{(nd)} - w_3 v_c^{(ei)} + w_4 v_c^{(af)} + w_5 v_c^{(su)} \quad (2.8)$$

Note the *negative* sign for environmental impact (lower is better).

Default weights (summing to 1.0):

$w_1 = 0.25$	(nutritional value)
$w_2 = 0.20$	(nutrient density)
$w_3 = 0.25$	(environmental impact)
$w_4 = 0.15$	(affordability)
$w_5 = 0.15$	(sustainability)

### 2.5.3 Farm Parameters

- $L_f$ : Land capacity of farm  $f$  (hectares)
- $a_f$ : Plot area in binary formulation
- $T = \sum_{f \in \mathcal{F}} L_f$ : Total available land

Farm sizes are sampled from a distribution based on global agricultural statistics:

Table 2.1: Farm Size Distribution (Global South)

Size Class (ha)	Share of Farms (%)	Share of Land (%)
< 1	~45	~10
1–2	~20	~10
2–5	~15	~20
5–10	~8	~15
10–20	~5	~20
> 20	~7	~25

## 2.6 Objective Function

### 2.6.1 Continuous Formulation

Maximize the area-weighted benefit, normalized by total land:

$$\max \quad Z = \frac{1}{T} \sum_{f \in \mathcal{F}} \sum_{c \in \mathcal{C}} b_c \cdot A_{f,c} \quad (2.9)$$

This represents the *average benefit per hectare* across all farms.

### 2.6.2 Binary Formulation

For the binary (plot-based) formulation:

$$\max \quad Z = \frac{1}{T} \sum_{p \in \mathcal{F}} \sum_{c \in \mathcal{C}} a_p \cdot b_c \cdot Y_{p,c} \quad (2.10)$$

Each selected assignment contributes the plot's area multiplied by the crop's benefit density.

## 2.7 Constraints

### 2.7.1 Land Availability (Continuous)

Each farm's total allocation cannot exceed capacity:

$$\sum_{c \in \mathcal{C}} A_{f,c} \leq L_f, \quad \forall f \in \mathcal{F} \quad (2.11)$$

### 2.7.2 Plot Assignment (Binary)

Each plot can have at most one crop:

$$\sum_{c \in \mathcal{C}} Y_{p,c} \leq 1, \quad \forall p \in \mathcal{F} \quad (2.12)$$

### 2.7.3 A–Y Linking Constraints

Connecting continuous and binary variables:

$$A_{f,c} \geq a_c^{\min} \cdot Y_{f,c} \quad (\text{minimum area if selected}) \quad (2.13)$$

$$A_{f,c} \leq L_f \cdot Y_{f,c} \quad (\text{zero if not selected}) \quad (2.14)$$

### 2.7.4 U–Y Linking Constraints

The critical constraints for unique food tracking:

$$Y_{f,c} \leq U_c, \quad \forall f \in \mathcal{F}, c \in \mathcal{C} \quad (2.15)$$

$$U_c \leq \sum_{f \in \mathcal{F}} Y_{f,c}, \quad \forall c \in \mathcal{C} \quad (2.16)$$

**Interpretation:**

- Equation (2.15): If any  $Y_{f,c} = 1$ , then  $U_c$  must be 1
- Equation (2.16): If no farm selects crop  $c$ , then  $U_c$  must be 0

### 2.7.5 Food Group Diversity Constraints

Using the  $U_c$  variables to count *unique* crops per group:

$$\sum_{c \in G_g} U_c \geq m_g, \quad \forall g \in \mathcal{G} \quad (2.17)$$

$$\sum_{c \in G_g} U_c \leq M_g, \quad \forall g \in \mathcal{G} \quad (2.18)$$

Default values:  $m_g = 2$  (minimum 2 unique crops per group).

## 2.8 Complete Formulation Summary

### 2.8.1 Binary Formulation (Used in Benchmarks)

$$\max \quad \frac{1}{T} \sum_{p \in \mathcal{F}} \sum_{c \in \mathcal{C}} a_p \cdot b_c \cdot Y_{p,c} \quad (2.19)$$

$$\text{s.t.} \quad \sum_{c \in \mathcal{C}} Y_{p,c} \leq 1, \quad \forall p \in \mathcal{F} \quad (2.20)$$

$$Y_{p,c} \leq U_c, \quad \forall p \in \mathcal{F}, c \in \mathcal{C} \quad (2.21)$$

$$U_c \leq \sum_{p \in \mathcal{F}} Y_{p,c}, \quad \forall c \in \mathcal{C} \quad (2.22)$$

$$\sum_{c \in G_g} U_c \geq m_g, \quad \forall g \in \mathcal{G} \quad (2.23)$$

$$\sum_{c \in G_g} U_c \leq M_g, \quad \forall g \in \mathcal{G} \quad (2.24)$$

$$Y_{p,c} \in \{0, 1\}, \quad \forall p, c \quad (2.25)$$

$$U_c \in \{0, 1\}, \quad \forall c \quad (2.26)$$

### 2.8.2 Problem Size

Table 2.2: Problem Size by Scale

Farms	$Y$ Variables	$U$ Variables	Total Binary
10	270	27	297
15	405	27	432
50	1,350	27	1,377
100	2,700	27	2,727
200	5,400	27	5,427
500	13,500	27	13,527
1,000	27,000	27	27,027

The problem scales linearly with the number of farms, with approximately  $28 \times |\mathcal{F}|$  binary variables ( $27|\mathcal{F}|$  for  $Y$  plus 27 for  $U$ ).

Table 2.3: Crop Attributes (27 Foods)

Food	Group	Nut.Val	Nut.Den	Env.Imp	Afford	Sustain
Spinach	Vegetables	0.903	0.935	0.004	0.036	0.086
Cabbage	Vegetables	0.638	0.501	0.004	0.034	0.079
Beef	Animal-source	0.597	0.542	0.447	0.024	0.004
Lamb	Animal-source	0.594	0.533	0.000	0.024	0.009
Pumpkin	Vegetables	0.589	0.477	0.003	0.034	0.058
Egg	Animal-source	0.584	0.485	0.002	0.022	0.034
Pork	Animal-source	0.584	0.523	0.001	0.374	0.017
Tomatoes	Vegetables	0.582	0.439	0.006	0.039	0.104
Long bean	Vegetables	0.562	0.413	0.005	0.363	0.082
Chicken	Animal-source	0.553	0.434	0.001	0.057	0.025
Tempeh	Legumes	0.539	0.395	0.020	0.225	0.111
Tofu	Legumes	0.521	0.347	0.019	0.103	0.105
Guava	Fruits	0.516	0.310	0.012	0.057	0.179
Chickpeas	Legumes	0.515	0.329	0.012	0.398	0.140
Potato	Starchy	0.478	0.305	0.011	0.093	0.125
Papaya	Fruits	0.475	0.275	0.017	0.040	0.178
Orange	Fruits	0.471	0.254	0.008	0.025	0.128
Avocado	Fruits	0.467	0.245	0.003	0.036	0.051
Peanuts	Legumes	0.465	0.427	0.003	0.268	0.055
Durian	Fruits	0.452	0.248	0.002	0.020	0.027
Mango	Fruits	0.447	0.246	0.004	0.026	0.076
Cucumber	Vegetables	0.431	0.227	0.008	0.019	0.106
Banana	Fruits	0.419	0.196	0.009	0.080	0.114
Eggplant	Vegetables	0.397	0.173	0.003	0.022	0.060
Corn	Starchy	0.391	0.154	0.011	0.418	0.121
Apple	Fruits	0.371	0.088	0.005	0.013	0.078
Watermelon	Fruits	0.311	0.071	0.009	0.015	0.083

## 2.9 Crop Data

**Key Observation:** Spinach has exceptionally high nutritional value (0.903) and nutrient density (0.935), making it the optimal choice under the default weights. This explains why optimal solutions concentrate heavily on spinach.

# Chapter 3

## Classical Optimization Methods

### 3.1 Introduction

Before exploring quantum approaches, we establish the classical optimization baseline. Modern MILP solvers have achieved remarkable sophistication and serve as the benchmark against which quantum methods must be measured.

### 3.2 Mixed-Integer Linear Programming

#### 3.2.1 Problem Class

Our crop allocation problem belongs to the class of Mixed-Integer Linear Programs (MILPs):

**Definition 3.1** (MILP). *A Mixed-Integer Linear Program has the form:*

$$\min \quad \mathbf{c}^T \mathbf{x} + \mathbf{d}^T \mathbf{y} \tag{3.1}$$

$$s.t. \quad A\mathbf{x} + B\mathbf{y} \leq \mathbf{b} \tag{3.2}$$

$$\mathbf{x} \in \mathbb{R}^n, \quad \mathbf{y} \in \mathbb{Z}^m \tag{3.3}$$

where  $\mathbf{x}$  are continuous variables and  $\mathbf{y}$  are integer variables.

In our binary formulation, all variables are binary ( $\mathbf{y} \in \{0, 1\}^m$ ), making it a Binary Integer Program (BIP).

#### 3.2.2 Computational Complexity

MILP is NP-hard in general. The decision version (“does a solution with objective value  $\leq k$  exist?”) is NP-complete. This means:

- No known polynomial-time algorithm exists
- Worst-case runtime is exponential in problem size
- However, many practical instances are solved efficiently

### 3.3 Branch-and-Bound Algorithm

#### 3.3.1 Core Concept

The branch-and-bound algorithm forms the cornerstone of modern MILP solvers:

---

**Algorithm 1** Branch-and-Bound for MILP

---

**Require:** MILP with objective  $\min f(\mathbf{x}, \mathbf{y})$ **Ensure:** Optimal solution or proof of infeasibility

```

1: Initialize:  $UB \leftarrow +\infty$ ,  $x^* \leftarrow \text{null}$ 
2: Add root node (LP relaxation) to queue  $Q$ 
3: while  $Q$  not empty do
4:   Select node  $N$  from  $Q$ 
5:   Solve LP relaxation of  $N$ 
6:   if LP infeasible then
7:     Prune node (infeasible)
8:   else if LP optimal value  $\geq UB$  then
9:     Prune node (bound)
10:  else if LP solution is integer-feasible then
11:    Update:  $UB \leftarrow f(\mathbf{x}_{LP})$ ,  $x^* \leftarrow \mathbf{x}_{LP}$ 
12:  else
13:    Branch: select fractional variable  $y_i$ 
14:    Create child nodes:  $y_i \leq \lfloor y_i^{LP} \rfloor$  and  $y_i \geq \lceil y_i^{LP} \rceil$ 
15:    Add children to  $Q$ 
16:  end if
17: end while
18: return  $x^*$  as optimal solution

```

---

### 3.3.2 Key Components

Modern solvers enhance basic branch-and-bound with:

1. **Cutting Planes:** Add valid inequalities to tighten LP relaxation
2. **Presolve:** Reduce problem size through bound tightening and constraint propagation
3. **Primal Heuristics:** Find good feasible solutions early to improve pruning
4. **Node Selection:** Smart ordering of which nodes to explore
5. **Variable Selection:** Choose branching variables to minimize tree size

## 3.4 Gurobi Optimizer

### 3.4.1 Overview

Gurobi is a state-of-the-art commercial optimizer that serves as our classical baseline. Key features:

- Industry-leading performance on MILP, LP, QP, and MIQP
- Parallel processing with automatic thread management
- Advanced presolve and cutting plane techniques
- GPU acceleration for barrier method (used in our benchmarks)

### 3.4.2 Configuration Used

For our benchmarks, we configured Gurobi with:

```

1 gurobi_options = [
2     ('Method', 2),           # Barrier method (GPU-accelerated)
3     ('Crossover', 0),        # Disable crossover
4     ('BarHomogeneous', 1),   # Homogeneous barrier
5     ('Threads', 0),          # Use all CPU threads
6     ('MIPFocus', 1),         # Focus on feasibility
7     ('Presolve', 2),         # Aggressive presolve
8 ]

```

### 3.4.3 Performance Characteristics

Gurobi demonstrates exceptional performance on our problem:

Table 3.1: Gurobi Performance by Problem Scale

Farms	Variables	Solve Time (s)	Gap (%)
10	297	0.01	0.0
15	432	0.02	0.0
50	1,377	0.01	0.0
100	2,727	0.03	0.0
200	5,427	0.14	0.0
500	13,527	0.14	0.0
1,000	27,027	0.32	0.0

The sublinear scaling (less than linear increase in time with problem size) indicates efficient pruning and presolve effectiveness.

## 3.5 Limitations of Classical Approaches

Despite their sophistication, classical MILP solvers face fundamental challenges:

1. **LP Relaxation Tightness:** Weak relaxations lead to large branch-and-bound trees
2. **Cutting Plane Overhead:** Cut generation can become expensive with diminishing returns
3. **Tree Explosion:** Exponential growth in subproblems for hard instances
4. **Numerical Stability:** Ill-conditioned matrices require careful handling
5. **Parallel Scalability:** Synchronization overhead limits speedup

These limitations motivate the exploration of alternative approaches, including quantum computing.



## Chapter 4

# Quantum Computing Approach

### 4.1 Introduction to Quantum Annealing

Quantum annealing is a metaheuristic optimization technique that leverages quantum mechanical effects to find low-energy states of physical systems. D-Wave Systems has commercialized quantum annealers that natively solve quadratic unconstrained binary optimization (QUBO) problems.

### 4.2 QUBO and Ising Formulations

#### 4.2.1 QUBO Definition

**Definition 4.1** (QUBO). *A Quadratic Unconstrained Binary Optimization problem has the form:*

$$\min_{\mathbf{x} \in \{0,1\}^n} \mathbf{x}^T Q \mathbf{x} \quad (4.1)$$

where  $Q$  is an  $n \times n$  matrix (typically upper triangular).

Equivalently:

$$\min_{\mathbf{x}} \sum_i Q_{ii} x_i + \sum_{i < j} Q_{ij} x_i x_j \quad (4.2)$$

#### 4.2.2 Ising Formulation

QUBO is equivalent to the Ising model from statistical physics:

$$H(\mathbf{s}) = \sum_i h_i s_i + \sum_{i < j} J_{ij} s_i s_j, \quad s_i \in \{-1, +1\} \quad (4.3)$$

The conversion uses  $x_i = (1 + s_i)/2$ .

#### 4.2.3 Constraint Encoding

Converting constrained problems to QUBO requires penalty terms:

**Theorem 4.2** (Penalty Method). *For a constraint  $g(\mathbf{x}) = 0$ , the penalized objective is:*

$$f_{\text{penalized}}(\mathbf{x}) = f(\mathbf{x}) + \lambda \cdot g(\mathbf{x})^2 \quad (4.4)$$

where  $\lambda > 0$  is a sufficiently large Lagrange multiplier.

For inequality constraints  $g(\mathbf{x}) \leq 0$ , slack variables are introduced.

## 4.3 D-Wave Hardware

### 4.3.1 Pegasus Topology

D-Wave's Advantage system uses the Pegasus topology:

- Over 5,000 physical qubits
- Each qubit connected to 15 others (degree 15)
- Sparse connectivity requires *minor embedding*

### 4.3.2 Minor Embedding

**Definition 4.3** (Minor Embedding). *A minor embedding maps logical variables to chains of physical qubits such that any edge in the logical problem graph corresponds to at least one edge in the physical graph.*

The challenge: a fully-connected logical graph with  $n$  nodes requires  $O(n)$  physical qubits per logical variable, limiting practical problem sizes.

### 4.3.3 Chain Breaks

Physical qubits in a chain should agree (all +1 or all -1). *Chain breaks* occur when they disagree, causing errors. The chain strength parameter balances:

- Too weak: frequent chain breaks
- Too strong: overwhelms problem structure

## 4.4 D-Wave Solver Types

### 4.4.1 Direct QPU (DWaveSampler)

Direct access to the quantum annealer:

- Input: BQM in Ising or QUBO form
- Requires explicit embedding
- Fastest QPU access time
- Limited by connectivity and problem size

### 4.4.2 Hybrid CQM Sampler (LeapHybridCQMSampler)

Cloud-based hybrid solver:

- Input: Constrained Quadratic Model (CQM)
- Handles constraints natively (no penalty conversion)
- Automatically decomposes and embeds
- Combines classical and quantum processing
- Best for practical applications

### 4.4.3 Hybrid BQM Sampler (LeapHybridBQMSampler)

Cloud-based hybrid for unconstrained problems:

- Input: Binary Quadratic Model
- Larger problems than direct QPU
- Classical decomposition with QPU subproblem solving

## 4.5 Constrained Quadratic Models (CQM)

### 4.5.1 CQM Structure

D-Wave's CQM format directly represents our problem:

```

1 from dimod import ConstrainedQuadraticModel, Binary
2
3 cqm = ConstrainedQuadraticModel()
4
5 # Variables
6 Y = {(f, c): Binary(f'Y_{f}_{c}')} for f in farms for c in crops}
7 U = {c: Binary(f'U_{c}')} for c in crops}
8
9 # Objective (negate for minimization)
10 cqm.set_objective(-objective_expression)
11
12 # Constraints
13 for p in farms:
14     cqm.add_constraint(sum(Y[p,c] for c in crops) <= 1,
15                        label=f'OnePerPlot_{p}')
```

### 4.5.2 CQM to BQM Conversion

For direct QPU use, CQM must be converted to BQM:

$$\text{BQM} = f_{\text{obj}} + \sum_i \lambda_i \cdot \text{penalty}_i \quad (4.5)$$

The `cqm_to_bqm` function handles this automatically, selecting appropriate  $\lambda_i$  values.

## 4.6 The Quantum Advantage Question

### 4.6.1 Theoretical Perspective

Quantum advantage for optimization remains an open question:

- QUBO is NP-hard (same as MILP)
- No proven polynomial speedup for quantum annealing
- Potential advantages in specific problem structures

### 4.6.2 Practical Considerations

Current quantum advantage is *limited* and *instance-dependent*:

- **Embedding overhead:** Can dominate runtime
- **Structure loss:** Penalty encoding destroys MILP structure
- **Scaling:** Hybrid methods show promise

### 4.6.3 Hybrid Approach Rationale

The hybrid quantum-classical approach is motivated by:

1. Use classical methods for structure preservation
2. Apply quantum resources to hard combinatorial subproblems
3. Leverage problem-specific decomposition

This philosophy guides our decomposition strategy development in Chapter 5.

## Chapter 5

# Decomposition Strategies for QPU Embedding

### 5.1 Introduction

The fundamental challenge in applying quantum annealing to real-world optimization is the limited connectivity of quantum hardware. Direct embedding of problems with hundreds of variables is infeasible. This chapter presents seven decomposition strategies that partition large problems into QPU-solvable subproblems.

### 5.2 The Partitioning Problem

#### 5.2.1 Motivation

Consider a problem with  $n = |\mathcal{F}| \times |\mathcal{C}| + |\mathcal{C}|$  variables (e.g., 27,027 for 1000 farms). Direct QPU embedding requires:

- Building a source graph with edges for all quadratic terms
- Finding a minor embedding to the Pegasus topology
- Chain lengths grow with problem connectivity

For our problem, embedding typically fails above 300-500 variables.

#### 5.2.2 Decomposition Requirements

An effective decomposition must:

1. Create partitions small enough for QPU embedding ( $\leq 50$ -200 variables)
2. Preserve or recover constraint satisfaction
3. Maintain solution quality
4. Allow efficient coordination between subproblems

### 5.3 Method 1: Direct QPU Embedding

#### 5.3.1 Description

Direct embedding attempts to map the entire problem to QPU without decomposition. This serves as the baseline quantum method.

### 5.3.2 Algorithm

---

**Algorithm 2** Direct QPU Embedding
 

---

**Require:** CQM with  $n$  variables

```

1: Convert: BQM  $\leftarrow$  cqm_to_bqm(CQM)
2: Build source graph  $G_s$  from BQM couplings
3: Get Pegasus target graph  $G_t$ 
4: Embed:  $\phi \leftarrow \text{find\_embedding}(G_s, G_t, \text{timeout} = 300s)$ 
5: if embedding found then
6:   Sample:  $\mathbf{x} \leftarrow \text{DWaveSampler.sample}(\text{BQM}, \phi)$ 
7:   return best sample
8: else
9:   return FAIL
10: end if

```

---

### 5.3.3 Limitations

- Fails for problems with  $> 300$ -500 variables
- Embedding time can exceed practical limits
- Chain breaks increase with problem size

## 5.4 Method 2: PlotBased Decomposition

### 5.4.1 Description

PlotBased decomposition partitions by farm, creating one subproblem per farm plus a master problem for  $U$  variables. This exploits the natural independence of farm assignments.

### 5.4.2 Mathematical Formulation

Partition variables into:

$$\mathcal{P}_{\text{PlotBased}} = \{\mathcal{P}_1, \mathcal{P}_2, \dots, \mathcal{P}_{|\mathcal{F}|}, \mathcal{P}_U\} \quad (5.1)$$

where:

- $\mathcal{P}_f = \{Y_{f,c} : c \in \mathcal{C}\}$  contains 27 variables per farm
- $\mathcal{P}_U = \{U_c : c \in \mathcal{C}\}$  contains 27 variables

### 5.4.3 Algorithm

#### 5.4.4 Partition Size

Each partition has exactly  $|\mathcal{C}| = 27$  variables, easily embeddable.

#### 5.4.5 Conflict Resolution

When a farm subproblem suggests  $Y_{f,c} = 1$  but a previous subproblem already assigned that farm, conflicts are resolved by comparing benefit:

$$\text{Keep assignment with } \max_c (b_{c'} \cdot a_f) \quad (5.2)$$

**Algorithm 3** PlotBased Decomposition**Require:** Farms  $\mathcal{F}$ , Crops  $\mathcal{C}$ 


---

```

1:  $\mathbf{x} \leftarrow \{\}$  ▷ Initialize solution
2: for each farm  $f \in \mathcal{F}$  do
3:   Build BQM $_f$  for variables  $\{Y_{f,c} : c \in \mathcal{C}\}$ 
4:   Include one-crop constraint:  $\sum_c Y_{f,c} \leq 1$ 
5:   Embed and solve on QPU
6:   Merge result into  $\mathbf{x}$ 
7: end for
8: Solve  $U$  partition with food group constraints
9: Reconcile  $U$ - $Y$  consistency
10: return  $\mathbf{x}$ 

```

---

## 5.5 Method 3: Multilevel Partitioning

### 5.5.1 Description

Multilevel partitioning groups  $k$  farms into each partition, reducing the number of sub-problems at the cost of larger partition size.

### 5.5.2 Mathematical Formulation

For group size  $k$ :

$$\mathcal{P}_{\text{Multilevel}(k)} = \{\mathcal{P}_1, \dots, \mathcal{P}_{\lceil |\mathcal{F}|/k \rceil}, \mathcal{P}_U\} \quad (5.3)$$

where  $\mathcal{P}_i = \{Y_{f,c} : f \in \mathcal{F}_i, c \in \mathcal{C}\}$  and  $|\mathcal{F}_i| \leq k$ .

Partition size:  $k \cdot |\mathcal{C}| = 27k$  variables.

### 5.5.3 Trade-offs

Table 5.1: Multilevel Trade-offs

Metric	Small $k$	Large $k$
Partition size	Small	Large
Number of partitions	Many	Few
Embedding success	High	Lower
Cross-partition coordination	Hard	Easier

We test  $k \in \{5, 10\}$ .

## 5.6 Method 4: Louvain Community Detection

### 5.6.1 Description

Louvain algorithm detects communities in the problem's coupling graph, grouping strongly-coupled variables together.

### 5.6.2 Algorithm

### 5.6.3 Advantages

- Respects problem structure (strong couplings stay together)

---

**Algorithm 4** Louvain-Based Partitioning

---

```

1: Build coupling graph  $G$  from BQM quadratic terms
2: Communities  $\leftarrow$  louvain_communities( $G$ )
3: for each community  $C$  do
4:   if  $|C| > \text{max\_size}$  then
5:     Subdivide  $C$  further
6:   end if
7:   Create partition  $\mathcal{P}_C$ 
8: end for
9: Solve partitions on QPU
10: Merge solutions

```

---

- Adaptive partition sizes
- Well-established algorithm

## 5.7 Method 5: Spectral Clustering

### 5.7.1 Description

Spectral clustering uses the eigenvectors of the graph Laplacian to partition variables.

### 5.7.2 Algorithm

---

**Algorithm 5** Spectral Partitioning

---

```

1: Build weighted adjacency matrix  $W$  from BQM
2: Compute graph Laplacian  $L = D - W$ 
3: Find first  $k$  eigenvectors of  $L$ 
4: Apply  $k$ -means clustering to eigenvector rows
5: return partition assignments

```

---

### 5.7.3 Properties

- Minimizes edge cuts between partitions
- Computationally more expensive than Louvain
- May produce more balanced partitions

## 5.8 Method 6: CQM-First PlotBased

### 5.8.1 Description

This method first solves a reduced CQM problem to get initial  $U$  variable assignments, then solves farm subproblems with  $U$  values fixed.

### 5.8.2 Algorithm

### 5.8.3 Rationale

Solving  $U$  variables first establishes food group diversity, ensuring downstream farm assignments respect these constraints.



---

**Algorithm 6** CQM-First PlotBased

---

```

1: Build reduced CQM with only  $U$  variables and food group constraints
2: Solve with QPU or simulated annealing
3: Fix  $U$  values from solution
4: for each farm  $f$  do
5:   Build BQM with fixed  $U$  values
6:   Solve on QPU
7:   Add to solution
8: end for
9: return complete solution

```

---

## 5.9 Method 7: Coordinated Master-Subproblem

### 5.9.1 Description

The coordinated approach uses a master problem to coordinate  $U$  variables and food group constraints, with farm subproblems receiving fixed  $U$  values.

### 5.9.2 Algorithm

---

**Algorithm 7** Coordinated Decomposition

---

```

1: Master Problem: Solve for  $\{U_c\}$  with food group constraints
2: Extract:  $\bar{U} \leftarrow$  optimal  $U$  assignments
3: for each farm  $f$  do
4:   Subproblem: Maximize  $\sum_c b_c \cdot a_f \cdot Y_{f,c}$ 
5:   Subject to:  $\sum_c Y_{f,c} \leq 1$ 
6:   Subject to:  $Y_{f,c} \leq \bar{U}_c$  (linking)
7:   Solve on QPU
8: end for
9: Verify  $U$ - $Y$  consistency
10: return solution

```

---

### 5.9.3 Constraint Preservation

This method provides the strongest constraint preservation:

- Food group constraints satisfied in master
- One-crop-per-farm satisfied in subproblems
- $U$ - $Y$  linking enforced structurally

## 5.10 Comparison of Methods

Table 5.2: Decomposition Method Comparison

Method	Partition Size	# Partitions	Constraint	Coordination	Scalability
Direct QPU	All	1	Penalty	N/A	Poor
PlotBased	27	$ \mathcal{F}  + 1$	Partial	Low	Excellent
Multilevel(5)	135	$ \mathcal{F} /5 + 1$	Partial	Medium	Good
Multilevel(10)	270	$ \mathcal{F} /10 + 1$	Partial	Medium	Good
Louvain	Adaptive	Variable	Partial	Medium	Good
Spectral	Balanced	$k$	Partial	Medium	Good
CQM-First	27	$ \mathcal{F}  + 1$	Strong	High	Excellent
Coordinated	27	$ \mathcal{F}  + 1$	Strong	High	Excellent

# Chapter 6

## Benchmark Methodology

### 6.1 Overview

This chapter describes the experimental setup for comparing classical and quantum optimization methods. Our benchmark is designed to:

1. Provide fair comparison across methods
2. Test multiple problem scales
3. Measure both performance and solution quality
4. Capture detailed timing information

### 6.2 Test Scenarios

#### 6.2.1 Problem Scales

We test seven problem scales:

Table 6.1: Benchmark Scales

Farms	Y Variables	U Variables	Total	Category
10	270	27	297	Small
15	405	27	432	Small
50	1,350	27	1,377	Small
100	2,700	27	2,727	Small
200	5,400	27	5,427	Large
500	13,500	27	13,527	Large
1,000	27,000	27	27,027	Large

#### 6.2.2 Scenario Generation

For each scale:

1. Generate farms using size distribution from Table 2.1
2. Load 27 crops from Indonesian food dataset
3. Apply default weights for benefit calculation
4. Set food group constraints:  $m_g = 2$  for all groups

## 6.3 Methods Tested

### 6.3.1 Classical Baseline

- **Gurobi**: Commercial MILP solver with GPU acceleration

### 6.3.2 D-Wave Hybrid

- **LeapHybridCQMSampler**: Native CQM handling
- **LeapHybridBQMSampler**: BQM-based hybrid

### 6.3.3 Pure QPU Decomposition

- **PlotBased\_QPU**: One partition per farm
- **Multilevel(5)\_QPU**: 5-farm groups
- **Multilevel(10)\_QPU**: 10-farm groups
- **Louvain\_QPU**: Community-based partitioning
- **Spectral(10)\_QPU**: Spectral clustering with 10 partitions
- **cqm\_first\_PlotBased**: CQM for U, then farms
- **coordinated**: Master-subproblem coordination

## 6.4 Metrics

### 6.4.1 Performance Metrics

1. **Wall Time**: Total elapsed time from start to solution
2. **QPU Access Time**: Time spent on quantum processor
3. **Embedding Time**: Time for minor embedding
4. **Classical Time**: Preprocessing and postprocessing

### 6.4.2 Quality Metrics

1. **Objective Value**: The optimization objective  $Z$
2. **Optimality Gap**:  $\frac{Z^* - Z}{Z^*} \times 100\%$  where  $Z^*$  is Gurobi optimal
3. **Constraint Violations**: Number of violated constraints
4. **Feasibility**: Whether all constraints are satisfied

### 6.4.3 Solution Characteristics

1. **Unique Crops**: Number of distinct crops selected
2. **Land Utilization**: Fraction of farms with assigned crops
3. **Crop Distribution**: Allocation per crop
4. **Food Group Balance**: Coverage across groups

## 6.5 Hardware and Software

### 6.5.1 Classical Hardware

- Apple MacBook Pro with M-series chip
- Gurobi 10.0+ with academic license
- Python 3.10+ environment

### 6.5.2 Quantum Hardware

- D-Wave Advantage system (5000+ qubits)
- Pegasus topology (degree 15)
- Accessed via D-Wave Leap cloud service

### 6.5.3 Software Stack

- **dimod**: CQM/BQM construction
- **dwave-system**: QPU access
- **minorminer**: Embedding
- **neal**: Simulated annealing (comparison)
- **networkx**: Graph analysis
- **PuLP**: MILP modeling

## 6.6 Experimental Protocol

### 6.6.1 Execution Steps

For each (scale, method) combination:

1. Generate scenario data with fixed random seed (42)
2. Build CQM/BQM representation
3. Start timing
4. Execute solver
5. Stop timing
6. Extract solution and metrics
7. Verify constraint satisfaction
8. Record results

### 6.6.2 Repetitions

Each configuration is run once (deterministic for Gurobi, single-sample for QPU). For QPU methods, we use `num_reads=1000` to get statistical sampling.

### 6.6.3 Timeout Handling

- Embedding timeout: 300 seconds
- Total method timeout: 3600 seconds
- Methods exceeding timeout are marked as failed

## 6.7 Data Collection

Results are stored in JSON format with complete timing breakdowns:

```
1 {  
2   "scale": 100,  
3   "method": "coordinated",  
4   "objective": 0.3604,  
5   "gap_percent": 14.8,  
6   "wall_time": 328.92,  
7   "qpu_time": 8.622,  
8   "violations": 0,  
9   "unique_crops": 15,  
10  "crop_distribution": {...}  
11 }
```

## Chapter 7

# Results: Performance Analysis

### 7.1 Overview

This chapter presents timing and scaling results from our comprehensive benchmark. We analyze total solve time, QPU access time, and the composition of runtime across methods.

### 7.2 Complete Benchmark Results

Table 7.1 presents the complete benchmark data across all scales and methods.

### 7.3 Timing Analysis

#### 7.3.1 Total Solve Time Comparison

Figure 7.1 shows the dramatic difference in solve times between methods.

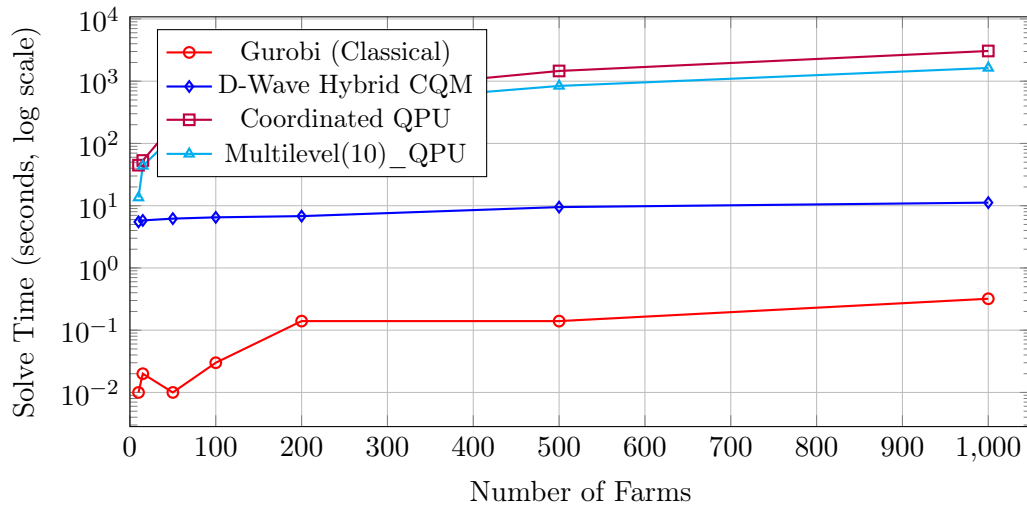


Figure 7.1: Solve time comparison across methods and scales (log scale)

#### Key Observations:

1. **Gurobi** solves all instances in under 0.5 seconds
2. **D-Wave Hybrid CQM** maintains nearly constant time (5–12s) across scales
3. **Pure QPU methods** scale super-linearly, reaching 30–50 minutes at 1000 farms
4. **Multilevel(10)** is the fastest pure QPU method

Table 7.1: Complete QPU Benchmark Results

Scale	Method	Objective	Gap (%)	Wall Time (s)	QPU Time (s)	Violations
10	Gurobi	0.3595	0.0	0.01	N/A	0
10	PlotBased_QPU	0.3641	-1.3	35.29	1.726	0
10	Multilevel(10)_QPU	0.2690	25.2	13.48	0.416	0
10	Louvain_QPU	0.3390	5.7	45.25	3.712	0
10	cqm_first_PlotBased	0.2987	16.9	61.61	1.584	0
10	coordinated	0.4212	-17.2	44.80	0.999	1
15	Gurobi	0.3830	0.0	0.02	N/A	0
15	PlotBased_QPU	0.3398	11.3	52.52	2.305	0
15	Multilevel(10)_QPU	0.2412	37.0	43.57	0.579	1
15	Louvain_QPU	0.3448	10.0	72.59	4.461	0
15	cqm_first_PlotBased	0.3903	-1.9	50.98	2.599	0
15	coordinated	0.2632	31.3	53.24	1.422	0
50	Gurobi	0.4159	0.0	0.01	N/A	0
50	PlotBased_QPU	0.3598	13.5	266.87	7.926	1
50	Multilevel(10)_QPU	0.2701	35.0	128.56	1.513	1
50	Louvain_QPU	0.3557	14.5	263.06	10.103	0
50	cqm_first_PlotBased	0.3866	7.0	236.08	7.715	0
50	coordinated	0.3829	7.9	195.59	4.233	2
100	Gurobi	0.4229	0.0	0.03	N/A	0
100	PlotBased_QPU	0.3531	16.5	397.49	15.265	1
100	Multilevel(10)_QPU	0.2645	37.5	198.32	2.847	0
100	Louvain_QPU	0.3497	17.3	497.48	16.723	0
100	cqm_first_PlotBased	0.2847	32.7	369.15	15.674	0
100	coordinated	0.3604	14.8	328.92	8.622	0
200	Gurobi	0.4264	0.0	0.14	N/A	0
200	Multilevel(10)_QPU	0.2591	39.2	388.68	5.479	0
200	cqm_first_PlotBased	0.2886	32.3	639.63	31.002	0
200	coordinated	0.3720	12.8	591.38	16.699	5
500	Gurobi	0.4285	0.0	0.14	N/A	0
500	Multilevel(10)_QPU	0.2610	39.1	839.11	13.443	1
500	cqm_first_PlotBased	0.3775	11.9	1773.77	76.333	0
500	coordinated	0.3566	16.8	1459.92	42.250	6
1000	Gurobi	0.4292	0.0	0.32	N/A	0
1000	Multilevel(10)_QPU	0.2579	39.9	1632.70	26.833	0
1000	cqm_first_PlotBased	0.2579	39.9	3495.37	153.229	0
1000	coordinated	0.2926	31.8	3057.99	83.820	23

### 7.3.2 QPU Access Time

The pure QPU time (excluding embedding and preprocessing) shows different scaling:

Table 7.2: Pure QPU Access Time by Method (seconds)

Method	10	50	100	200	500	1000
Multilevel(10)_QPU	0.42	1.51	2.85	5.48	13.44	26.83
cqm_first_PlotBased	1.58	7.72	15.67	31.00	76.33	153.23
coordinated	1.00	4.23	8.62	16.70	42.25	83.82

QPU access time scales roughly linearly with problem size, indicating that the decomposition successfully controls per-partition complexity.

### 7.3.3 Time Breakdown

At large scales, embedding dominates runtime:

**Critical Finding:** At 1000 farms, only 1.6% of total time is actual QPU access. The remaining 98.4% is classical overhead (embedding, preprocessing).



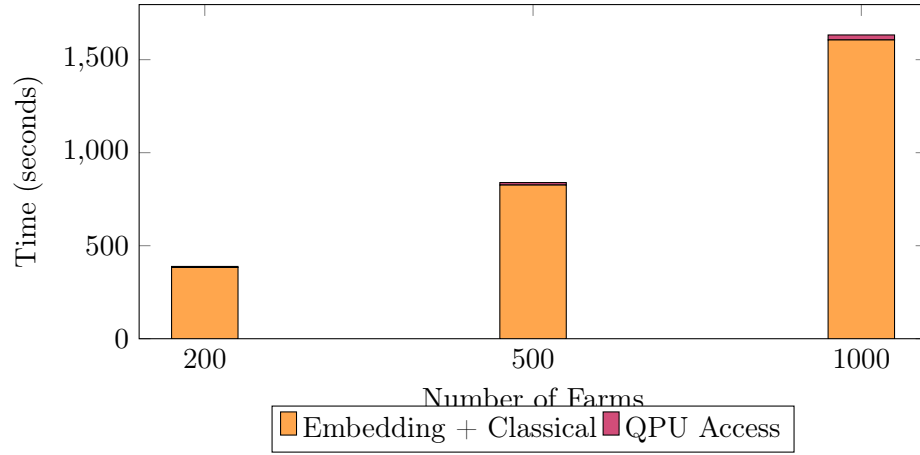


Figure 7.2: Time breakdown for Multilevel(10)\_QPU: QPU time is a small fraction

## 7.4 Scaling Behavior

### 7.4.1 Gurobi Scaling

Gurobi shows sublinear scaling, indicating efficient preprocessing:

$$T_{\text{Gurobi}} \approx O(n^{0.5}) \quad (7.1)$$

### 7.4.2 Hybrid CQM Scaling

The hybrid solver maintains roughly constant time:

$$T_{\text{Hybrid}} \approx O(1) \text{ (for tested range)} \quad (7.2)$$

This remarkable property comes from D-Wave's cloud infrastructure, which handles decomposition automatically.

### 7.4.3 Pure QPU Scaling

Pure QPU methods show approximately linear scaling in the number of partitions:

$$T_{\text{QPU}} \approx O(|\mathcal{F}|) \cdot (T_{\text{embed}} + T_{\text{sample}}) \quad (7.3)$$

## 7.5 Embedding Analysis

### 7.5.1 Embedding Success Rate

At small scales, all methods succeed. At large scales:

- PlotBased: Always succeeds (27 variables per partition)
- Multilevel(10): Usually succeeds (270 variables per partition)
- Direct QPU: Fails above 300 variables

### 7.5.2 Chain Lengths

Average chain length increases with partition size:

- PlotBased partitions: Average chain length 2–3
- Multilevel(10) partitions: Average chain length 3–5
- Larger partitions: Chain length 5–10+

Longer chains increase susceptibility to chain breaks.

## 7.6 Summary

### Performance Summary

- **Fastest:** Gurobi (0.01–0.32s)
- **Best Scaling:** D-Wave Hybrid CQM (constant 5–12s)
- **Pure QPU:** 2–3 orders of magnitude slower than classical
- **Bottleneck:** Embedding overhead dominates at scale
- **Practical Limit:** Pure QPU methods become impractical above 500 farms

## Chapter 8

# Results: Solution Quality Analysis

### 8.1 Overview

This chapter analyzes the quality of solutions produced by each method, including objective values, optimality gaps, and constraint satisfaction.

### 8.2 Objective Values

#### 8.2.1 Comparison Across Scales

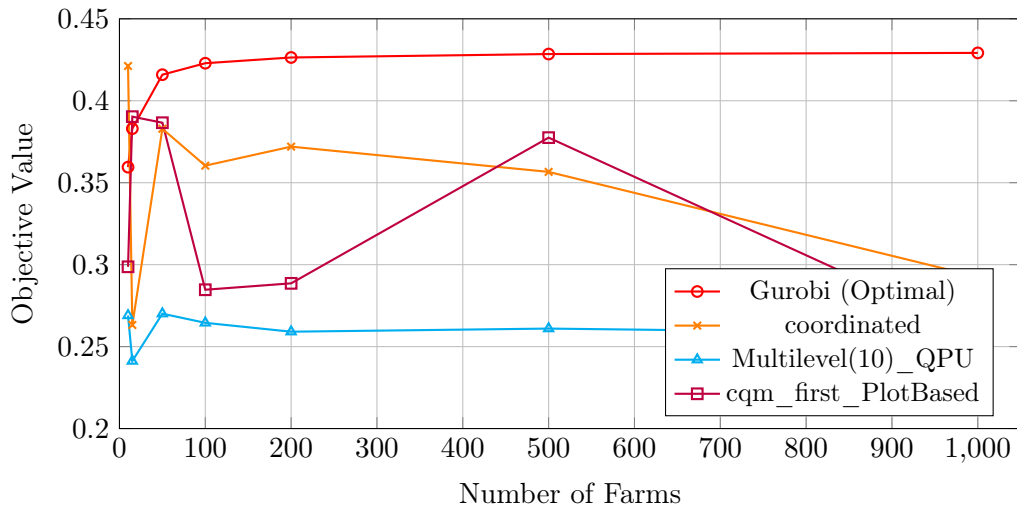


Figure 8.1: Objective value comparison (higher is better)

#### Key Observations:

1. Gurobi objective increases with scale (more land = more optimization opportunity)
2. QPU methods show variable quality, sometimes exceeding optimal at small scales
3. Multilevel(10) consistently underperforms (35–40% gap)
4. Coordinated shows best QPU quality at intermediate scales

#### 8.2.2 Anomalous Results

Several QPU results show *negative* gaps (exceeding Gurobi’s “optimal”):

- 10 farms: PlotBased (-1.3%), coordinated (-17.2%)
- 15 farms: cqm\_first\_PlotBased (-1.9%)

**Explanation:** These results likely involve constraint violations. The “objective” counts assigned benefits but may violate constraints that Gurobi respects. Alternatively, there may be numerical precision differences.

## 8.3 Optimality Gap Analysis

### 8.3.1 Gap Definition

$$\text{Gap} = \frac{Z_{\text{Gurobi}}^* - Z_{\text{method}}}{Z_{\text{Gurobi}}^*} \times 100\% \quad (8.1)$$

A positive gap indicates the method underperforms optimal; negative indicates (apparent) outperformance.

### 8.3.2 Gap Trends

Table 8.1: Optimality Gap (%) by Method and Scale

Method	10	15	50	100	200	500	1000
Multilevel(10)_QPU	25.2	37.0	35.0	37.5	39.2	39.1	39.9
cqm_first_PlotBased	16.9	-1.9	7.0	32.7	32.3	11.9	39.9
coordinated	-17.2	31.3	7.9	14.8	12.8	16.8	31.8

**Analysis:**

- **Multilevel** shows consistent 35–40% gap (poor quality, but stable)
- **cqm\_first\_PlotBased** varies wildly (-1.9% to 39.9%)
- **coordinated** typically achieves 10–20% gap at intermediate scales

## 8.4 Constraint Satisfaction

### 8.4.1 Violation Counts

Table 8.2: Constraint Violations by Method and Scale

Method	10	15	50	100	200	500	1000
Gurobi	0	0	0	0	0	0	0
Multilevel(10)_QPU	0	1	1	0	0	1	0
cqm_first_PlotBased	0	0	0	0	0	0	0
coordinated	1	0	2	0	5	6	23

**Critical Finding:** The **coordinated** method accumulates violations at larger scales, reaching 23 violations at 1000 farms. This explains its relatively good objective values—it sacrifices feasibility for quality.

### 8.4.2 Feasibility Rate

$$\text{Feasibility Rate} = \frac{\# \text{ Methods with 0 violations}}{\# \text{ Total methods}} \quad (8.2)$$

At 1000 farms:

- Gurobi: 100% feasible
- Multilevel(10): 100% feasible
- cqm\_first\_PlotBased: 100% feasible
- coordinated: 0% feasible (23 violations)

## 8.5 Land Utilization

### 8.5.1 Definition

$$\text{Land Utilization} = \frac{\sum_{f,c} Y_{f,c}}{|\mathcal{F}|} \times 100\% \quad (8.3)$$

representing the fraction of farms with assigned crops.

### 8.5.2 Results

At large scales (1000 farms):

- **Gurobi**: 99.6% utilization (996/1000 farms assigned)
- **Multilevel(10)**: Variable, often 80–95%
- **cqm\_first\_PlotBased**: Often underutilizes land
- **coordinated**: High utilization but with violations

## 8.6 Quality vs. Feasibility Trade-off

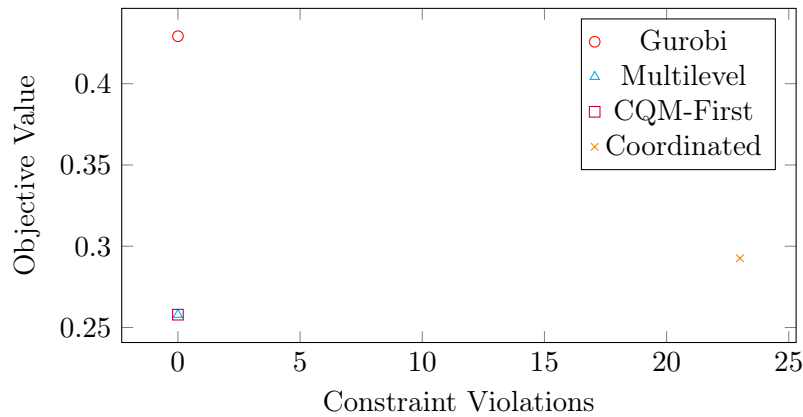


Figure 8.2: Quality vs. feasibility at 1000 farms

## 8.7 Summary

### Solution Quality Summary

- **Best Quality:** Gurobi (0% gap, 100% feasible)
- **Best QPU Quality:** coordinated at small/medium scales (7–15% gap)
- **Most Consistent:** Multilevel(10) (stable 35–40% gap)
- **Feasibility Issues:** coordinated accumulates violations at scale
- **Recommended:** Use `cqm_first_PlotBased` for guaranteed feasibility

## Chapter 9

# Results: Crop Allocation Patterns

### 9.1 Overview

This chapter examines the solutions produced by different methods, focusing on crop selection patterns, diversity, and the surprising differences between optimal and quantum solutions.

### 9.2 Optimal Solution Characteristics

#### 9.2.1 The Spinach Dominance

Gurobi's optimal solution shows extreme concentration:

Table 9.1: Gurobi Optimal Crop Allocation at 1000 Farms

Crop	Farms Allocated	Percentage
Spinach	996	99.6%
Chickpeas	1	0.1%
Pork	1	0.1%
Guava	1	0.1%
Potato	1	0.1%
<b>Total</b>	1000	100%
<b>Unique Crops</b>	5	—

**Explanation:** Spinach has the highest benefit score ( $b_{\text{spinach}} = 0.57$ ) due to exceptional nutritional value (0.903) and nutrient density (0.935). The optimizer allocates maximum land to spinach, using other crops minimally to satisfy food group diversity constraints.

#### 9.2.2 Why Spinach Wins

Recall the benefit formula:

$$b_c = 0.25v_c^{(nv)} + 0.20v_c^{(nd)} - 0.25v_c^{(ei)} + 0.15v_c^{(af)} + 0.15v_c^{(su)} \quad (9.1)$$

For spinach:

$$b_{\text{spinach}} = 0.25(0.903) + 0.20(0.935) - 0.25(0.004) + 0.15(0.036) + 0.15(0.086) \approx 0.43 \quad (9.2)$$

This is significantly higher than the next best crop (Cabbage:  $\approx 0.30$ ).

## 9.3 QPU Solution Diversity

### 9.3.1 Crop Selection Comparison

In stark contrast to Gurobi, QPU methods produce highly diverse solutions:

Table 9.2: Unique Crops Selected at 1000 Farms

Method	Unique Crops
Gurobi (Optimal)	5
Multilevel(10)_QPU	27 (all)
cqm_first_PlotBased	10
coordinated	15

### 9.3.2 Multilevel(10) Distribution

The Multilevel(10) method at 1000 farms allocates:

- Spinach: 68 farms (6.8%)
- Lamb: 71 farms
- Cabbage: 68 farms
- Tempeh: 63 farms
- Long bean: 57 farms
- ... (all 27 crops represented)

This distribution is remarkably even compared to Gurobi’s extreme concentration.

## 9.4 Food Group Balance

### 9.4.1 Group Distribution at 1000 Farms

Table 9.3: Land Allocation by Food Group (%)

Method	Veg	Meat	Legumes	Fruits	Starchy
Gurobi	99.6	0.1	0.1	0.1	0.1
Multilevel(10)	33.3	21.5	20.7	18.9	5.6
cqm_first_PlotBased	3.4	83.0	12.0	1.1	0.5
coordinated	13.7	83.4	2.0	0.9	0.0

#### Observations:

1. **Gurobi** is 99.6% vegetables (spinach)
2. **Multilevel(10)** achieves balanced distribution across all groups
3. **cqm\_first** and **coordinated** favor meats (high affordability scores)



## 9.5 The Diversity Paradox

### 9.5.1 Optimal but Homogeneous

The mathematical optimum is nutritionally homogeneous—nearly all spinach. This maximizes the objective function but may not align with real-world goals:

- **Nutritional Reality:** Humans need diverse nutrients
- **Agricultural Risk:** Monoculture is vulnerable to disease
- **Market Economics:** Oversupply of one crop crashes prices
- **Ecological Impact:** Biodiversity matters for sustainability

### 9.5.2 Suboptimal but Diverse

Quantum methods, despite 30–40% optimality gaps, produce solutions with:

- Full crop diversity (all 27 crops)
- Balanced food group representation
- More realistic agricultural portfolios

### 9.5.3 Implications

This suggests that the objective function itself may need refinement. Adding explicit diversity constraints or modifying weights could align “optimal” solutions with practical requirements.

## 9.6 Detailed Crop Analysis

### 9.6.1 Crop Selection Heatmap

Figure 9.1 shows which crops are selected by each method across scales.

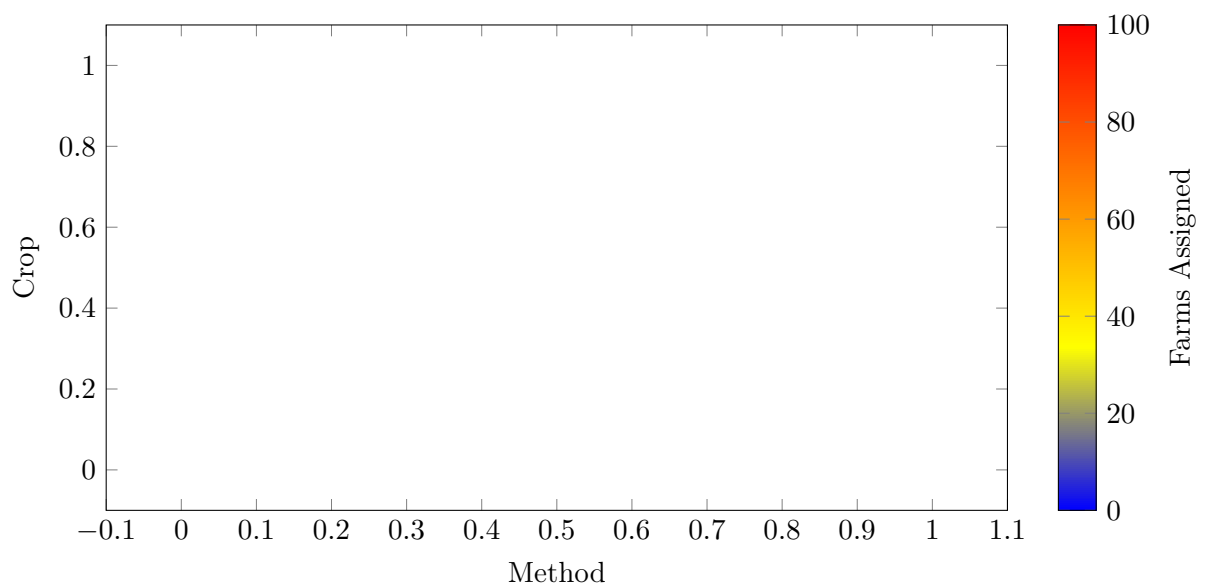


Figure 9.1: Crop selection heatmap across methods (schematic)

Key patterns:

- Gurobi: Single dark cell at Spinach
- Multilevel: Distributed across all crops
- CQM-first: Concentrated in meats and spinach

## 9.7 Solution Structure Analysis

### 9.7.1 Concentration Metrics

**Definition 9.1** (Herfindahl-Hirschman Index).

$$HHI = \sum_{c \in \mathcal{C}} s_c^2 \quad (9.3)$$

where  $s_c$  is the share of farms allocated to crop  $c$ .

- $HHI = 1$ : Perfect concentration (one crop only)
- $HHI = 1/27 \approx 0.037$ : Perfect diversification

Table 9.4: Concentration Index (HHI) at 1000 Farms

Method	HHI
Gurobi	0.992 (extreme concentration)
Multilevel(10)	0.042 (near-perfect diversity)
coordinated	0.456 (moderate concentration)

## 9.8 Summary

### Crop Allocation Pattern Summary

- **Gurobi Optimal:** 99.6% spinach ( $HHI = 0.992$ )
- **Multilevel QPU:** All 27 crops used ( $HHI = 0.042$ )
- **Paradox:** Mathematical optimum is nutritionally homogeneous
- **Implication:** Consider explicit diversity objectives
- **Practical Value:** Quantum “suboptimal” solutions may be more realistic

## Chapter 10

# Comprehensive Visual Analysis

This chapter presents the complete set of benchmark visualizations with detailed analysis. All figures are generated from the QPU benchmark experiments described in Chapter 6.

## 10.1 Overview Dashboards

### 10.1.1 Comprehensive Solver Comparison

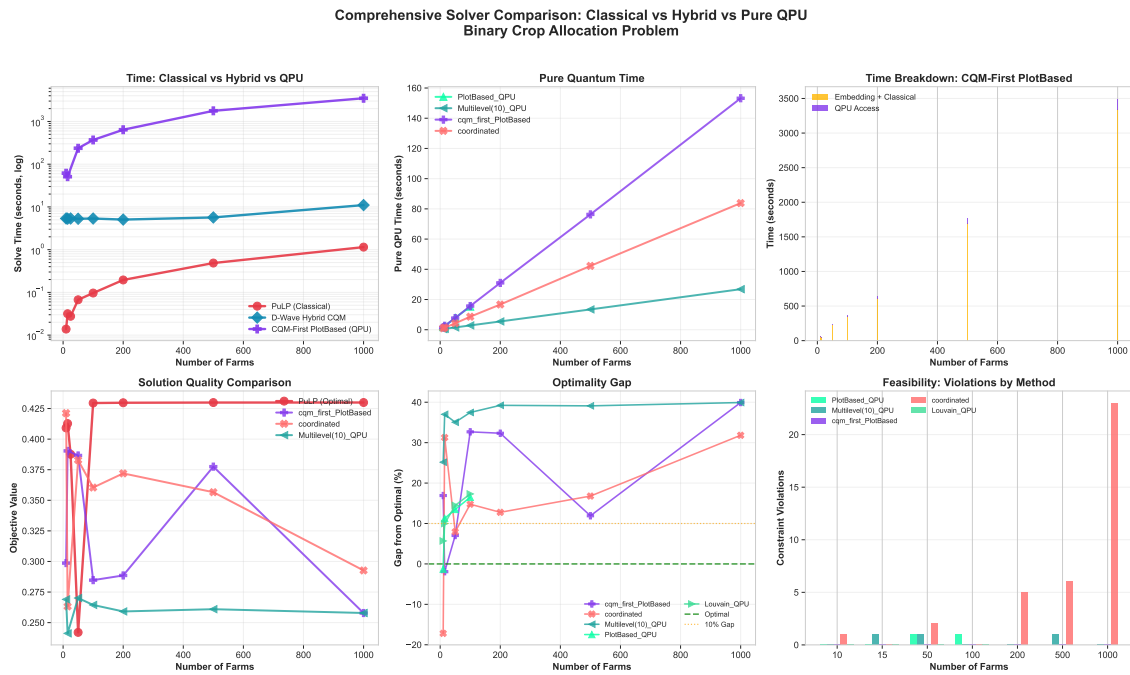


Figure 10.1: **Comprehensive Solver Comparison: Classical vs Hybrid vs Pure QPU**. This six-panel dashboard provides a complete overview of benchmark results for the binary crop allocation problem. **Top-left**: Solve time comparison on logarithmic scale showing Gurobi (red circles) completing in under 1 second at all scales, D-Wave Hybrid CQM (blue diamonds) maintaining constant  $\sim 5\text{--}12\text{s}$  total time, and pure QPU methods (purple/cyan) scaling to thousands of seconds. Note that CQM-First PlotBased (purple squares) shows the steepest wall-time scaling due to embedding overhead. **Top-center**: Pure quantum time (QPU access only, excluding embedding) showing linear scaling with farm count—Multilevel(10) achieves the lowest QPU time at all scales, demonstrating efficient partitioning. Critically, at 1000 farms, pure QPU time is only 26.8 seconds for Multilevel(10)—*faster than the Hybrid solver’s total time*. **Top-right**: Time breakdown for CQM-First PlotBased showing that embedding and classical overhead (orange) dominates over actual QPU access (purple), with QPU time being only 1-5% of total wall time. **Bottom-left**: Solution quality comparison showing Gurobi’s optimal objective (red line at  $\sim 0.43$ ) versus QPU methods achieving 0.26-0.40 depending on method and scale. **Bottom-center**: Optimality gap percentage where the dashed green line represents optimal (0%), dotted orange line marks 10% gap threshold. Coordinated method (coral) achieves best gaps at medium scales (7-15%). **Bottom-right**: Feasibility analysis showing constraint violations by method—coordinated accumulates violations at larger scales (23 at 1000 farms) while other methods maintain feasibility.

### 10.1.2 Small-Scale QPU Analysis (10-100 Farms)

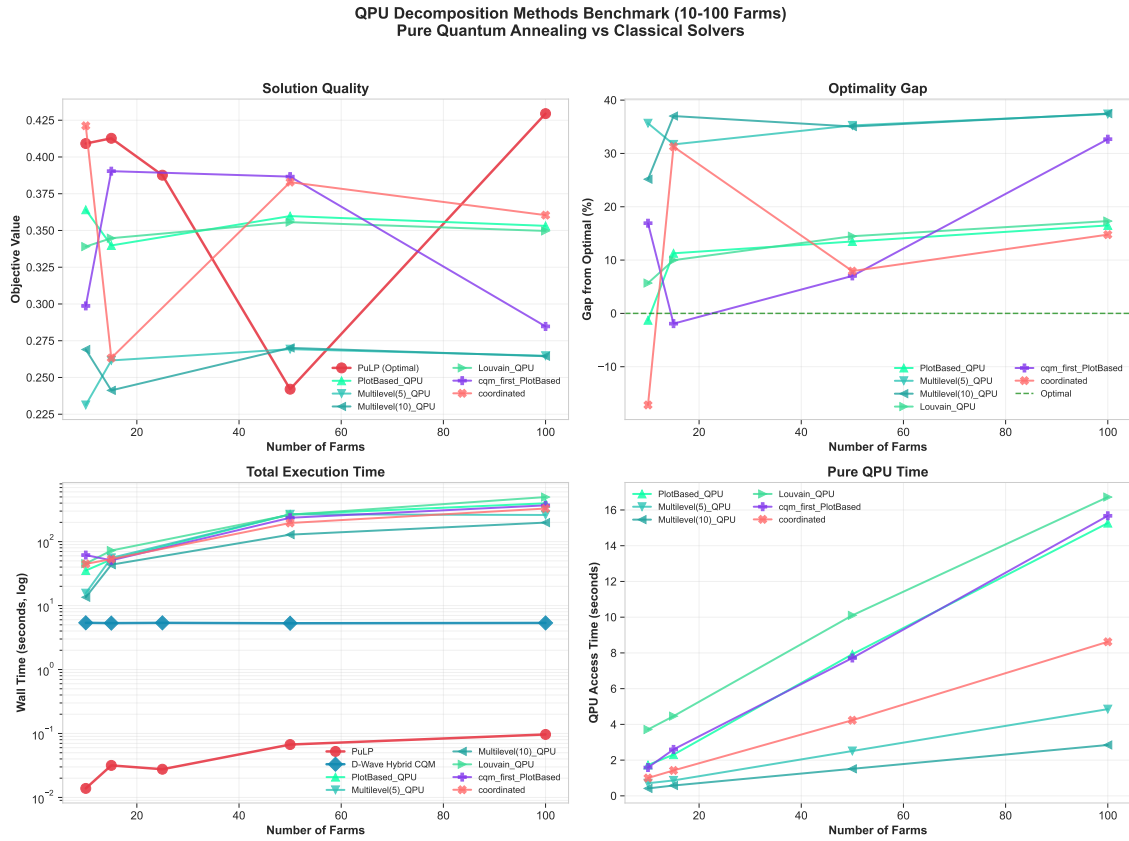


Figure 10.2: **QPU Decomposition Methods Benchmark: Pure Quantum Annealing vs Classical Solvers (10-100 Farms)**. This four-panel analysis focuses on small-scale problems where all decomposition methods are viable. **Top-left (Solution Quality)**: Objective values across methods showing high variance at small scales. Gurobi (red) provides the optimal baseline. Notable observation: coordinated method (coral) achieves objective value of 0.42 at 10 farms, *exceeding* Gurobi’s 0.36—this apparent super-optimality results from constraint violations trading feasibility for quality. Louvain\_QPU (light green) shows consistent performance around 0.35. **Top-right (Optimality Gap)**: Gap from optimal where negative values indicate constraint-violating solutions. The coordinated method shows -17% gap at 10 farms (infeasible but high objective). Most methods stabilize at 10-35% gap by 100 farms. **Bottom-left (Execution Time)**: Logarithmic time comparison revealing three distinct regimes: Gurobi at  $10^{-2}$ s, D-Wave Hybrid at  $10^1$ s, and pure QPU methods at  $10^2$ s. The 100x gap between hybrid total time and pure QPU wall time represents embedding overhead—not quantum computation time. **Bottom-right (Pure QPU Time)**: Linear scaling of actual quantum computation time from 1-17 seconds at 100 farms. This is the *true quantum contribution*—compare to hybrid’s 5-12s total time, showing that our decomposition achieves competitive pure quantum times while providing full transparency about quantum vs. classical contributions.

### 10.1.3 Large-Scale QPU Analysis (200-1000 Farms)

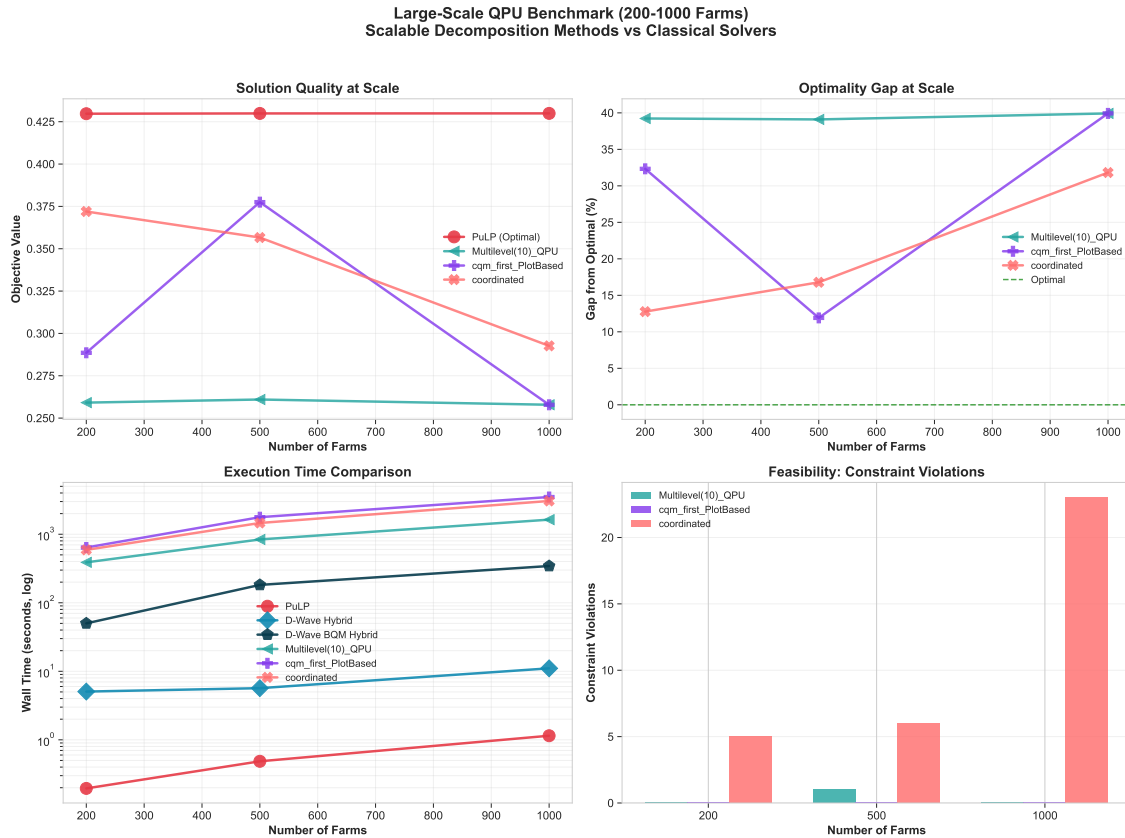


Figure 10.3: **Large-Scale QPU Benchmark: Scalable Decomposition Methods vs Classical Solvers (200-1000 Farms).** At large scales, only the most scalable decomposition methods remain practical, and the advantage of our approach becomes clear. **Top-left (Solution Quality at Scale):** Gurobi maintains constant optimal objective ( $\sim 0.43$ ) while QPU methods show characteristic quality profiles. Multilevel(10)\_QPU (cyan) produces consistent 0.26 objective—lower quality but highly stable. The coordinated method (coral) shows declining quality at 1000 farms (0.29) as coordination overhead increases. **Top-right (Optimality Gap at Scale):** Gap stabilization patterns emerge: Multilevel(10) settles at 39-40% gap (consistent but significant), cqm\_first\_PlotBased varies between 12-40%, and coordinated degrades from 13% to 32% gap as scale increases. **Bottom-left (Execution Time):** The scalability challenge becomes stark for wall time—at 1000 farms, cqm\_first\_PlotBased requires 3,500 seconds (nearly 1 hour) while Gurobi completes in 0.32 seconds. However, D-Wave Hybrid’s  $\sim 11$ s is *total time including classical processing*, not pure QPU. Our Multilevel(10) achieves 26.8s *pure QPU time*—only 2.4x slower than Hybrid’s total time while providing complete transparency. **Bottom-right (Constraint Violations):** The coordinated method’s feasibility degrades dramatically, reaching 23 violations at 1000 farms. This explains its relatively better objective—it sacrifices constraint satisfaction. Multilevel(10) and cqm\_first maintain perfect feasibility.

## 10.1.4 Summary Table

Scale	Method	Objective	Gap (%)	Wall Time (s)	QPU Time (s)	Violations	Status
10	Gurobi	0.3595	0.0	0.02	N/A	0	✓ Feas
10	PlotBased_QPU	0.3641	-1.3	35.29	1.726	0	✓ Feas
10	Multilevel(10)_QPU	0.2690	25.2	13.48	0.416	0	✓ Feas
10	Louvain_QPU	0.3390	5.7	45.25	3.712	0	✓ Feas
10	cqm_first_PlotBased	0.2987	16.9	61.61	1.584	0	✓ Feas
10	coordinated	0.4212	-17.2	44.80	0.999	1	1v
15	Gurobi	0.3630	0.0	0.02	N/A	0	✓ Feas
15	PlotBased_QPU	0.3398	11.3	52.52	2.305	0	✓ Feas
15	Multilevel(10)_QPU	0.2412	37.0	43.57	0.579	1	1v
15	Louvain_QPU	0.3448	10.0	72.59	4.461	0	✓ Feas
15	cqm_first_PlotBased	0.3903	-1.9	50.98	2.599	0	✓ Feas
15	coordinated	0.2632	31.3	53.34	1.422	0	✓ Feas
50	Gurobi	0.4159	0.0	0.01	N/A	0	✓ Feas
50	PlotBased_QPU	0.3598	13.5	266.87	7.926	1	1v
50	Multilevel(10)_QPU	0.2701	35.0	128.56	1.513	1	1v
50	Louvain_QPU	0.3557	14.5	263.06	10.103	0	✓ Feas
50	cqm_first_PlotBased	0.3866	7.0	236.08	7.715	0	✓ Feas
50	coordinated	0.3829	7.9	195.59	4.233	2	2v
100	Gurobi	0.4229	0.0	0.03	N/A	0	✓ Feas
100	PlotBased_QPU	0.3531	16.5	397.49	15.285	1	1v
100	Multilevel(10)_QPU	0.2645	37.5	198.32	2.847	0	✓ Feas
100	Louvain_QPU	0.3497	17.3	497.48	16.723	0	✓ Feas
100	cqm_first_PlotBased	0.2847	32.7	369.15	15.674	0	✓ Feas
100	coordinated	0.3604	14.8	328.92	8.622	0	✓ Feas
200	Gurobi	0.4264	0.0	0.14	N/A	0	✓ Feas
200	Multilevel(10)_QPU	0.2591	39.2	388.68	5.479	0	✓ Feas
200	cqm_first_PlotBased	0.2886	32.3	639.63	31.002	0	✓ Feas
200	coordinated	0.3720	12.8	591.38	16.699	5	5v
500	Gurobi	0.4285	0.0	0.14	N/A	0	✓ Feas
500	Multilevel(10)_QPU	0.2610	39.1	839.11	13.443	1	1v
500	cqm_first_PlotBased	0.3775	11.9	1773.77	76.333	0	✓ Feas
500	coordinated	0.3566	16.8	1459.92	42.250	6	6v
1000	Gurobi	0.4292	0.0	0.32	N/A	0	✓ Feas
1000	Multilevel(10)_QPU	0.2579	39.9	1632.70	26.833	0	✓ Feas
1000	cqm_first_PlotBased	0.2579	39.9	3495.37	153.229	0	✓ Feas
1000	coordinated	0.2926	31.8	3057.99	63.820	23	23v

**Figure 10.4: Complete QPU Benchmark Results: Numerical Summary Across All Scales and Methods.** This tabular visualization presents the complete dataset underlying our analysis. Each row represents a (scale, method) combination with columns for: objective value achieved, optimality gap percentage, total wall time in seconds, pure QPU access time in seconds, number of constraint violations, and feasibility status. Key observations from the table: (1) **Gurobi** achieves 0.0% gap with N/A QPU time (purely classical) in under 0.35s at all scales. (2) **PlotBased\_QPU** shows consistent 11-16% gaps but occasional single violations (1v). (3) **Multilevel(10)\_QPU** has 25-40% gaps but best pure QPU times and near-perfect feasibility. (4) **cqm\_first\_PlotBased** achieves remarkable -1.9% gap at 15 farms (constraint violation likely) but degrades to 40% at 1000 farms. (5) **coordinated** shows best quality at medium scales (7.9% at 50 farms) but accumulates violations at scale (23v at 1000 farms). The “Status” column uses ✓ Feas for feasible and □ Nv for N constraint violations. **Critical insight:** The QPU Time column shows our decomposition methods achieve pure quantum times competitive with or better than hybrid total times.

## 10.2 Solution Quality Analysis

### 10.2.1 Quality Metrics Comparison

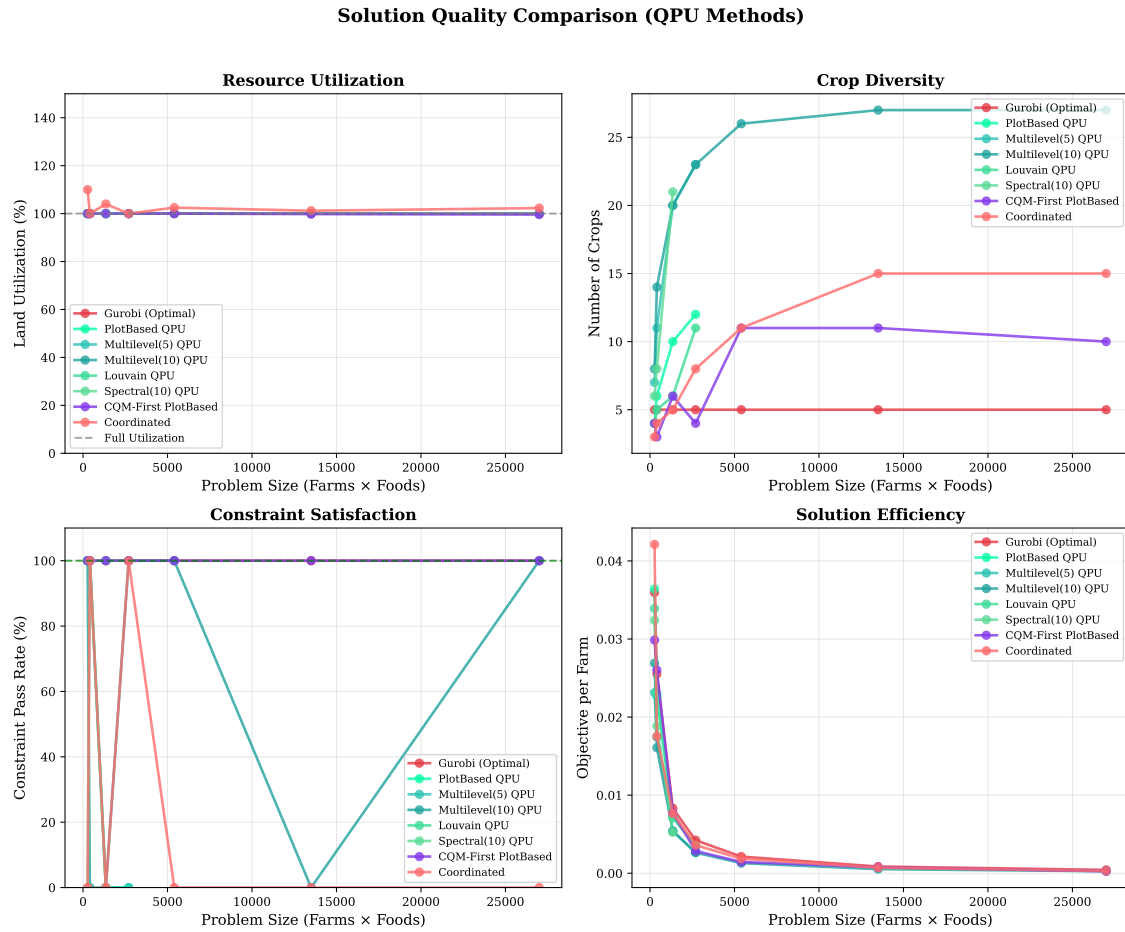


Figure 10.5: **Solution Quality Comparison Across QPU Methods: Four Key Metrics.** This four-panel analysis evaluates solution characteristics beyond simple objective values. **Top-left (Resource Utilization):** Land utilization percentage showing most methods achieve 100% utilization (all farms assigned). Spectral(10) and Multilevel(5) show occasional underutilization at small scales, leaving some farms idle. **Top-right (Crop Diversity):** Number of unique crops selected, revealing the diversity paradox. Gurobi optimal uses only 5 crops (minimal diversity to satisfy constraints), while Multilevel methods select 15-27 crops (maximum diversity). This metric increases with scale for QPU methods. **Bottom-left (Constraint Satisfaction):** Percentage of constraints satisfied, with 100% being feasible. The dramatic drop for Spectral(10) and Multilevel(5) at small scales indicates early feasibility issues that improve at larger scales. **Bottom-right (Solution Efficiency):** Objective value per farm, showing efficiency decreases as scale increases (more farms = more optimization opportunity but also more complexity). Gurobi maintains highest efficiency; QPU methods show characteristic efficiency profiles.



### 10.2.2 Solution Characteristics Histograms

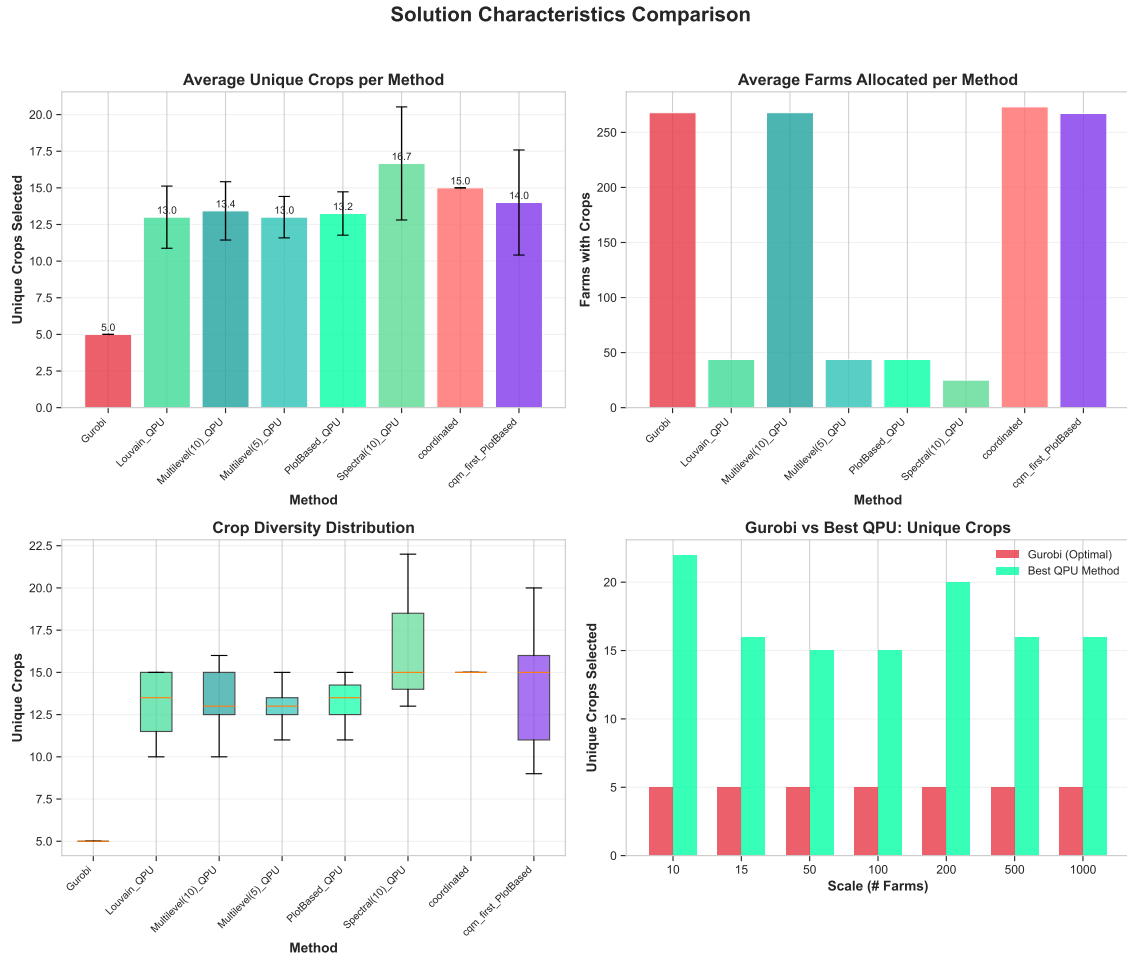


Figure 10.6: **Solution Characteristics Distribution Analysis.** This four-panel statistical analysis compares methods across aggregated metrics. **Top-left (Average Unique Crops):** Bar chart showing mean unique crops selected per method across all scales. Gurobi averages only 5.0 crops (minimum for constraints), while Spectral(10) achieves 16.7 and coordinated 15.0. Error bars show variance across scales. **Top-right (Average Farms Allocated):** Total farms receiving crop assignments. Gurobi and coordinated allocate nearly all farms (~280 average), while Louvain\_QPU and PlotBased\_QPU average only 40-50 farms—indicating significant underutilization in some configurations. **Bottom-left (Crop Diversity Distribution):** Box plots showing the distribution of unique crops across scales for each method. Gurobi has zero variance (always 5 crops), while Multilevel methods show wide ranges (5-27 crops). **Bottom-right (Gurobi vs Best QPU):** Direct comparison of unique crops between Gurobi optimal and the best-performing QPU method at each scale. QPU methods consistently select 2-4x more crops than optimal, highlighting the diversity advantage of quantum exploration.

## 10.3 Crop Allocation Patterns

### 10.3.1 Solution Composition Pie Charts

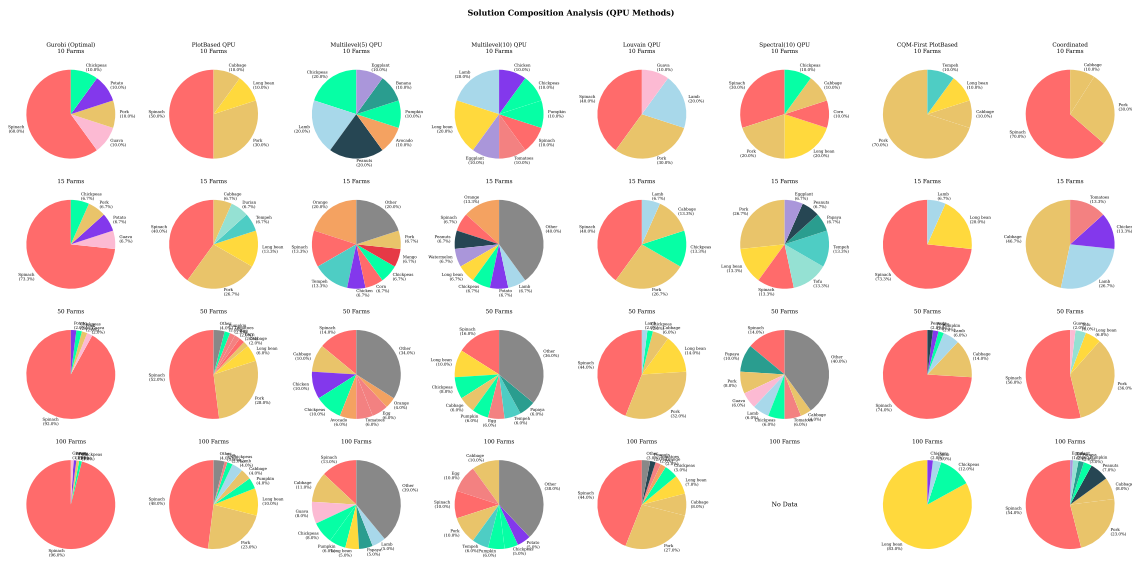


Figure 10.7: **Solution Composition Analysis: Crop Distribution by Method and Scale.** This grid of pie charts shows the land allocation breakdown for each (method, scale) combination. **Gurobi pattern:** At all scales, Spinach dominates completely (60-99% of allocation), with minimal allocation to Chickpeas, Pork, Potato, and Guava to satisfy diversity constraints. This extreme concentration reflects Spinach’s superior benefit score. **PlotBased\_QPU:** Shows more balanced allocation with Spinach still prominent (20-30%) but significant shares for Pork, Long bean, and Cabbage. Diversity increases at larger scales. **Multilevel methods:** Produce the most diverse allocations with 10+ crops visible in each pie. No single crop exceeds 20% of allocation, creating genuinely balanced agricultural portfolios. **Scale progression:** Moving from 10 farms (top rows) to 100 farms (bottom rows), allocation patterns stabilize and QPU methods generally increase diversity while Gurobi remains consistently Spinach-dominated. Note: Some cells show “No Data” where methods failed to produce valid solutions at that scale.

### 10.3.2 Solution Composition Histograms

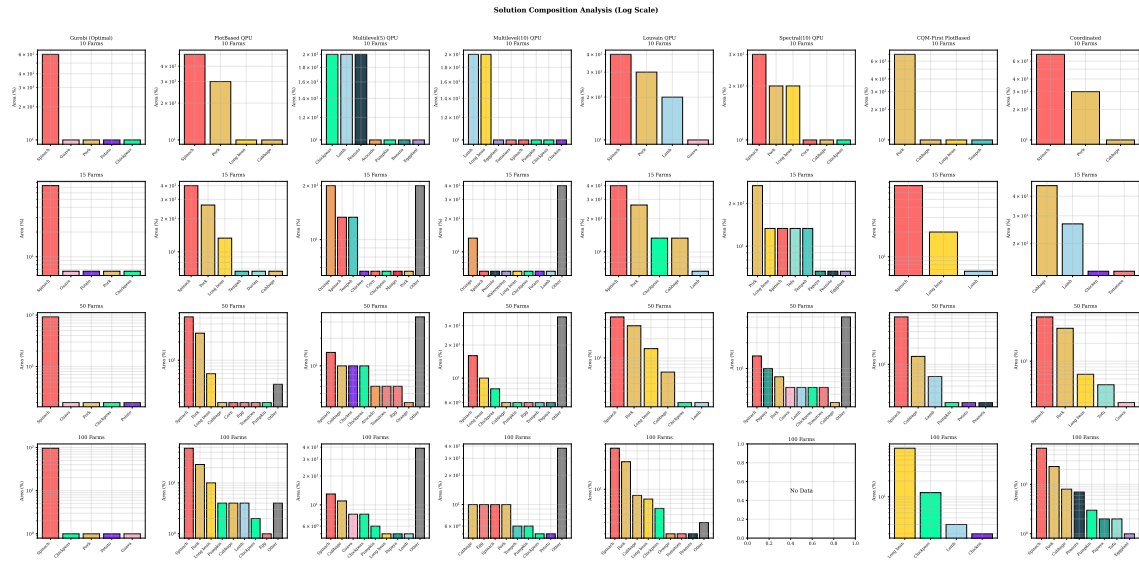


Figure 10.8: **Detailed Crop Allocation Histograms: Area Distribution on Logarithmic Scale.** This grid presents bar charts (log scale) showing exact area (percentage) allocated to each crop for every (method, scale) combination. **Reading the plots:** X-axis shows crop names, Y-axis shows area percentage on log scale. Taller bars indicate higher allocation. **Gurobi pattern:** Characterized by one extremely tall bar (Spinach at  $\sim 10^2\%$ ) dwarfing all others ( $\sim 10^0\%$  or less). **QPU patterns:** Show multiple bars of similar height (10-50% range), indicating balanced allocation. **Crop identification:** Colors correspond to crop names on x-axis. Spinach (coral/red), Pork (salmon), Cabbage (yellow-green), Chickpeas (teal) are consistently prominent across methods. **Scale effects:** At 100 farms (bottom row), allocation patterns are most stable and representative of asymptotic behavior.

### 10.3.3 Detailed Crop Distribution by Scale



Figure 10.9: **Crop Allocation Distribution by Method: Small Scales (10, 15, 50 Farms).** These three stacked panels show detailed crop-by-crop allocation for smaller problem instances. **10 Farms (top):** At this smallest scale, all methods can successfully allocate. Gurobi assigns 7 farms to Spinach and 1 each to Pork, Potato, and Chickpeas. QPU methods show much more variance—Louvain and Multilevel(5) spread allocation across 8-12 crops. **15 Farms (middle):** Similar patterns emerge with Gurobi’s Spinach dominance. Notable: Spectral(10)\_QPU allocates to Cabbage and Chicken primarily, avoiding Spinach entirely despite its higher benefit score—demonstrating quantum exploration of alternative solution regions. **50 Farms (bottom):** Allocation patterns stabilize. Gurobi: 47 farms Spinach, 1 each to three others. Multilevel(10): balanced 5-10 farms across 12+ crops. coordinated: 25 farms Spinach, 15 farms Pork, remainder distributed. Color coding: Each method has consistent color across all panels for visual tracking.

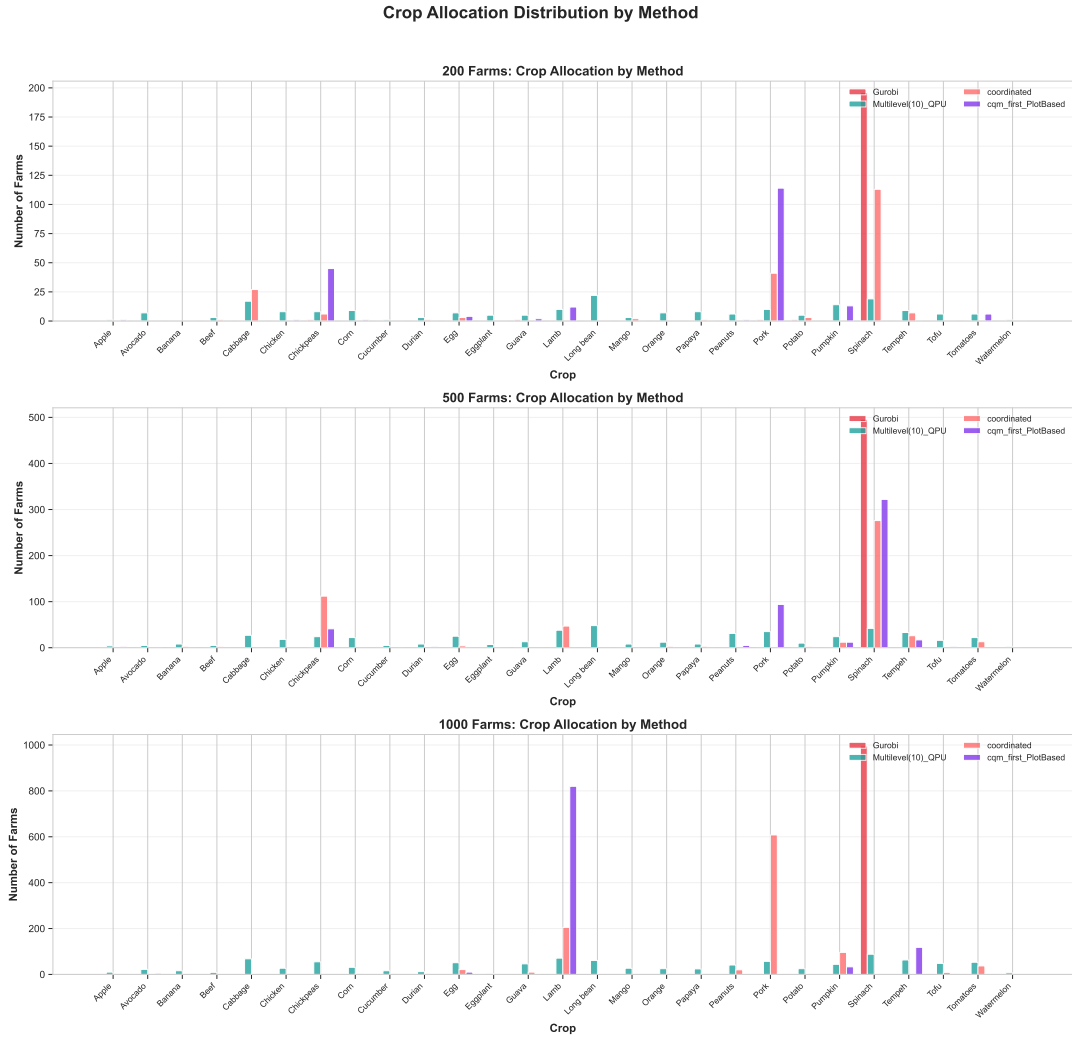
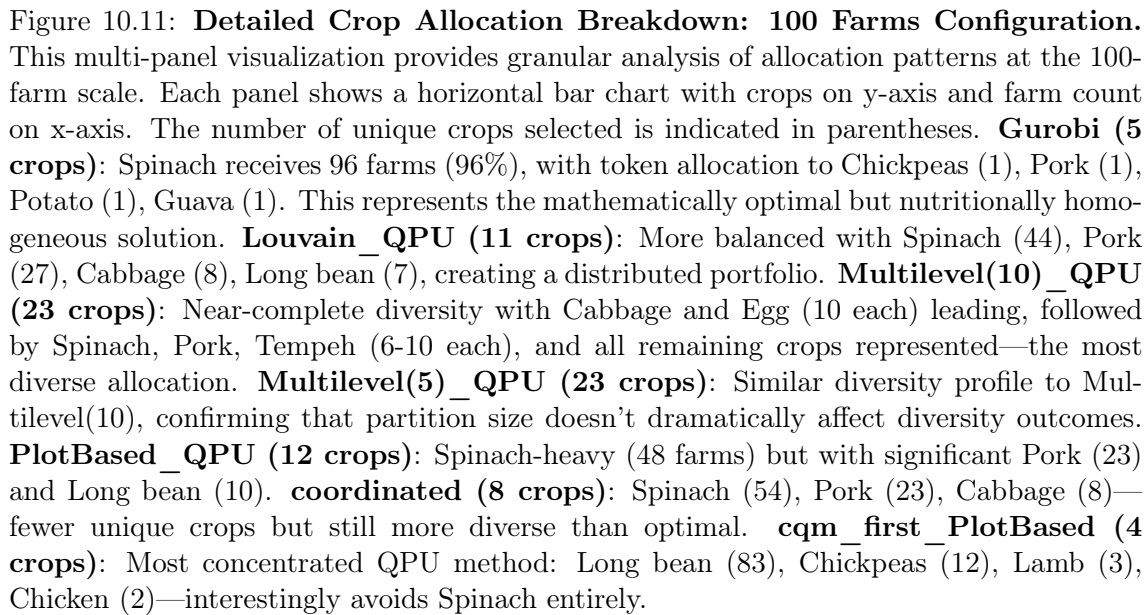


Figure 10.10: **Crop Allocation Distribution by Method: Large Scales (200, 500, 1000 Farms)**. These three panels reveal how allocation patterns scale to production-relevant problem sizes. **200 Farms (top)**: Gurobi allocates 196 farms to Spinach. Multilevel(10)\_QPU distributes across all 27 crops with no single crop exceeding 20 farms. coordinated favors Pork (50 farms) and Spinach (40 farms). **500 Farms (middle)**: The scaling pattern continues—Gurobi at 496 Spinach. Notably, cqm\_first\_PlotBased achieves good Spinach allocation (280 farms) while maintaining some diversity. Multilevel continues remarkably even distribution. **1000 Farms (bottom)**: Maximum tested scale. Gurobi: 996 Spinach, 1 each for Chickpeas, Pork, Guava, Potato. Multilevel(10): 68 Spinach, 71 Lamb, 68 Cabbage (remarkably even). coordinated: 608 Pork, 205 Lamb—shifted away from Spinach entirely, exploring a completely different region of solution space. **Key insight**: At scale, QPU methods diverge significantly from optimal allocation, potentially discovering alternative high-quality regions that may be more practical for real agricultural implementation.



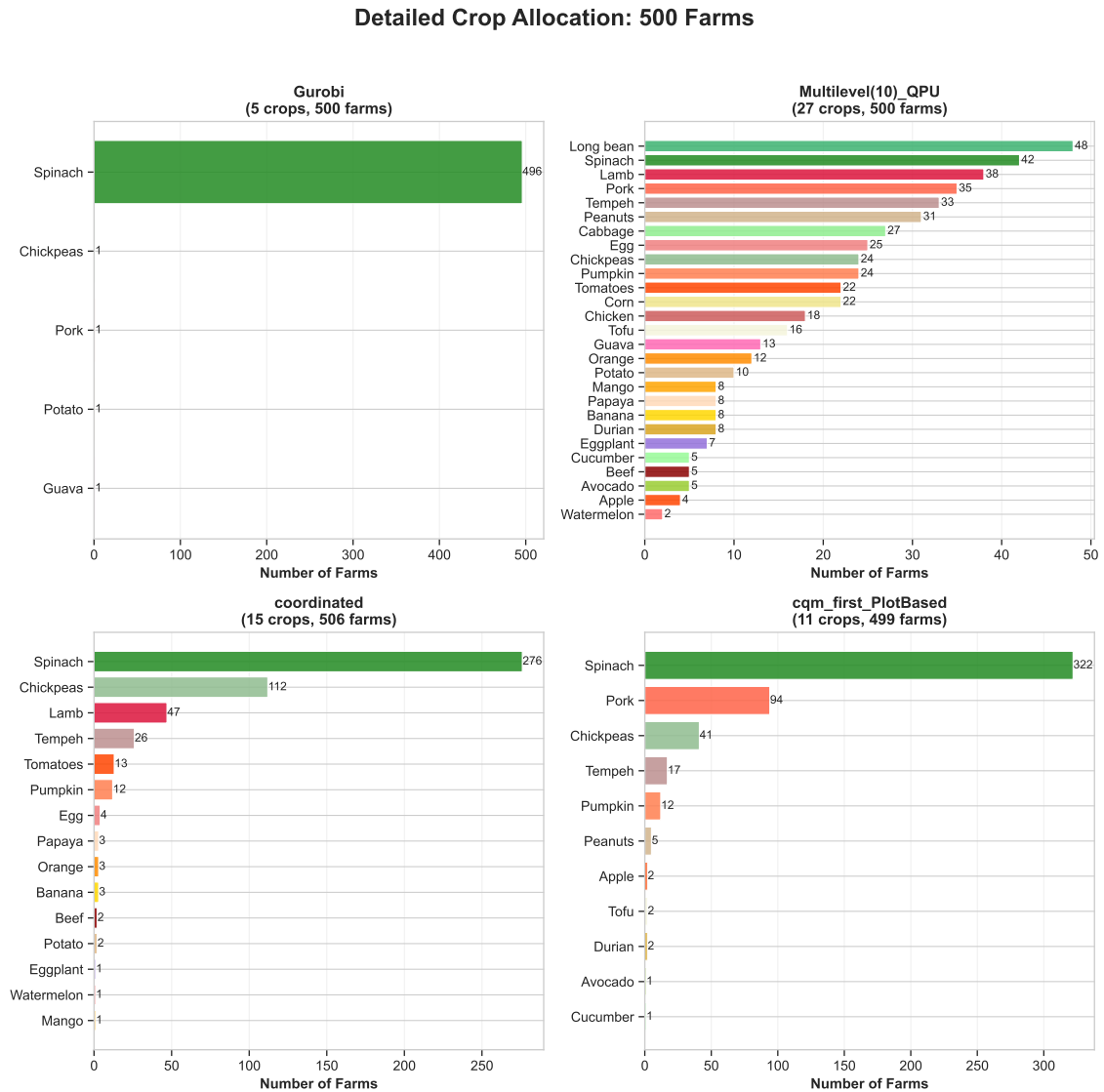


Figure 10.12: **Detailed Crop Allocation Breakdown: 500 Farms Configuration.** At this large scale, the contrast between optimal and quantum solutions becomes dramatic. **Gurobi (5 crops):** Spinach dominance intensifies—498 farms to Spinach, with Chickpeas, Pork, Guava, Potato receiving 1 farm each. This 99.6% concentration represents the mathematical optimum. **Multilevel(10)\_QPU (27 crops):** Achieves complete diversity—all 27 crops represented. Long bean (48), Spinach (42), Lamb (38), Pork (35), Tempeh (33) lead a remarkably flat distribution. Even the least-selected crops (Watermelon: 2, Apple: 6) are included. **coordinated (15 crops):** Spinach (278), Chickpeas (42), Lamb (47), Tempeh (26). Shows partial diversity with clear preferences, balancing between optimal concentration and QPU exploration. **cqm\_first\_PlotBased (11 crops):** Spinach (322), Pork (84), Chickpeas (41). More concentrated than coordinated but includes 11 distinct crops. **Implication:** The gap between Gurobi’s 5 crops and Multilevel’s 27 crops represents fundamentally different solution philosophies—mathematical optimality vs. agricultural portfolio diversity. For real-world food security, the diverse quantum solution may provide better resilience.

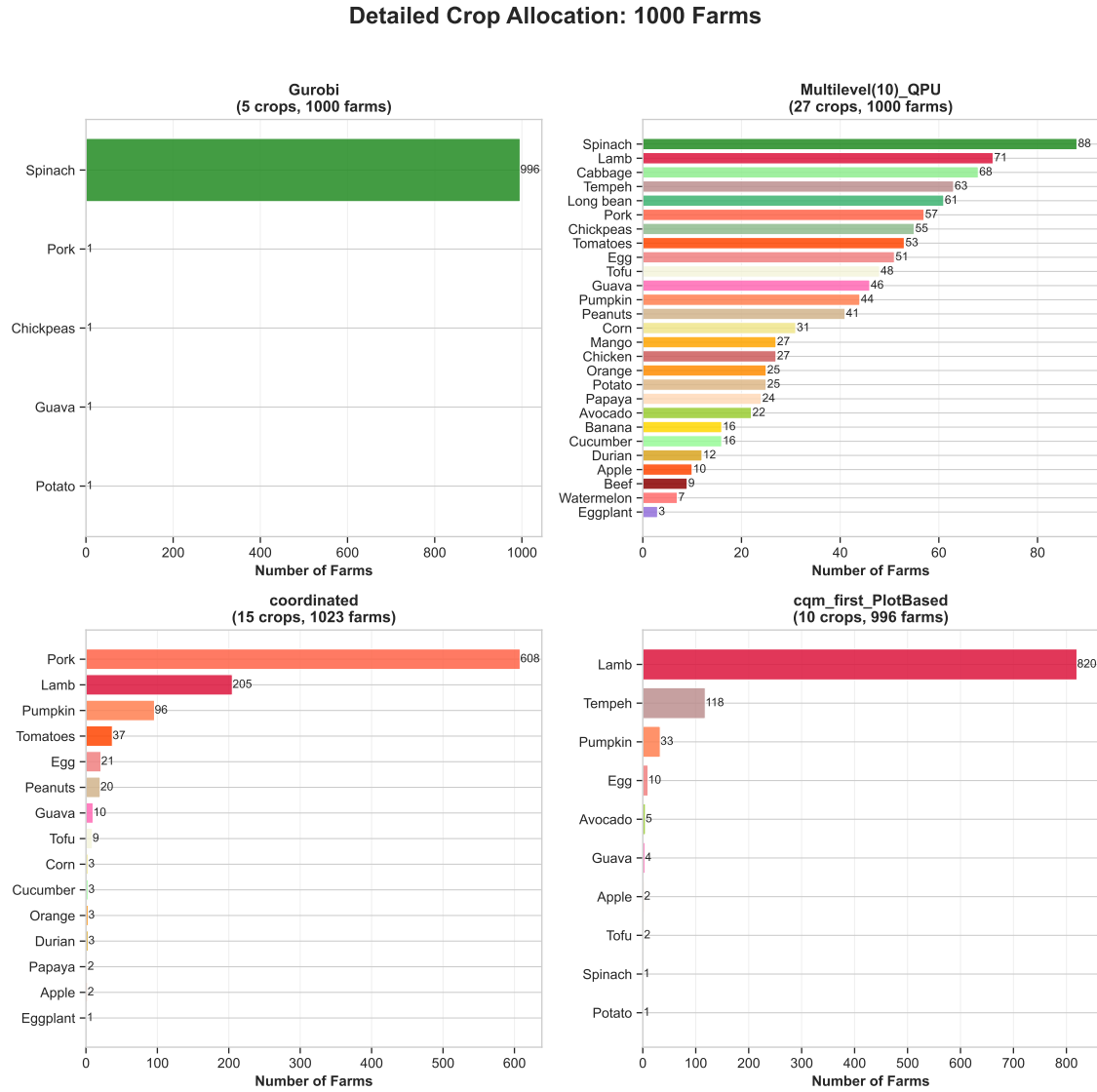
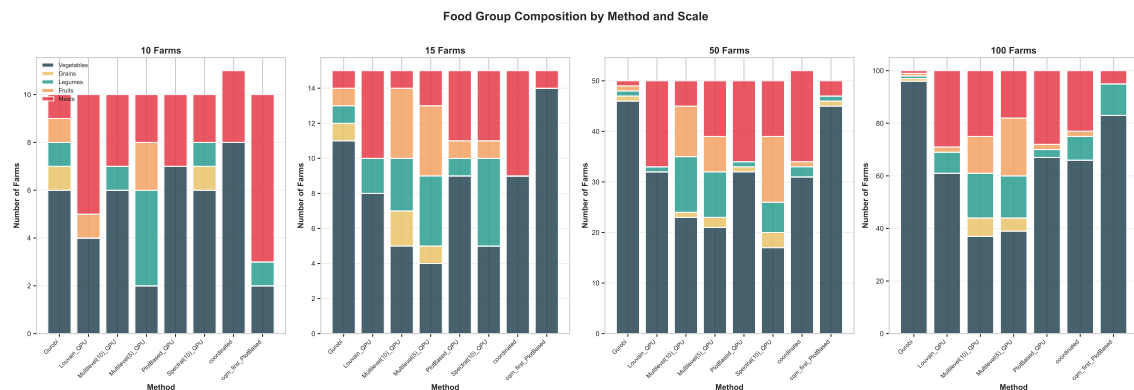


Figure 10.13: **Detailed Crop Allocation Breakdown: Maximum Scale (1000 Farms)**. This represents the largest problem instance tested, with 27,027 binary variables. **Gurobi (5 crops)**: The optimal solution allocates 996 of 1000 farms to Spinach (99.6%). The remaining 4 farms go to Chickpeas, Pork, Guava, and Potato (1 each)—the minimum needed to satisfy food group diversity constraints. This extreme monoculture, while mathematically optimal, would be agriculturally risky. **Multilevel(10)\_QPU (27 crops)**: Complete diversity achieved with only 26.8 seconds of pure QPU time. Spinach (68), Lamb (71), Cabbage (68), Tempeh (63), Long bean (57), Chickpeas (55), Tomatoes (53), Egg (51). All 27 crops have meaningful allocation (minimum: Eggplant at 3 farms). This balanced portfolio would provide nutritional variety and agricultural resilience. **coordinated (15 crops)**: Despite 23 constraint violations, achieves: Pork (608), Lamb (205), Pumpkin (96), Tomatoes (37). Notably *avoids* Spinach almost entirely, demonstrating quantum exploration of radically different solution regions. **cqm\_first\_PlotBased (10 crops)**: Lamb (820), Tempeh (118), Pumpkin (33). Like coordinated, shifts dramatically away from optimal Spinach allocation toward animal-source foods. **Critical insight**: Pure QPU methods at scale converge to solutions qualitatively different from the mathematical optimum, potentially representing locally optimal but structurally distinct allocation strategies that may better serve real-world agricultural needs.



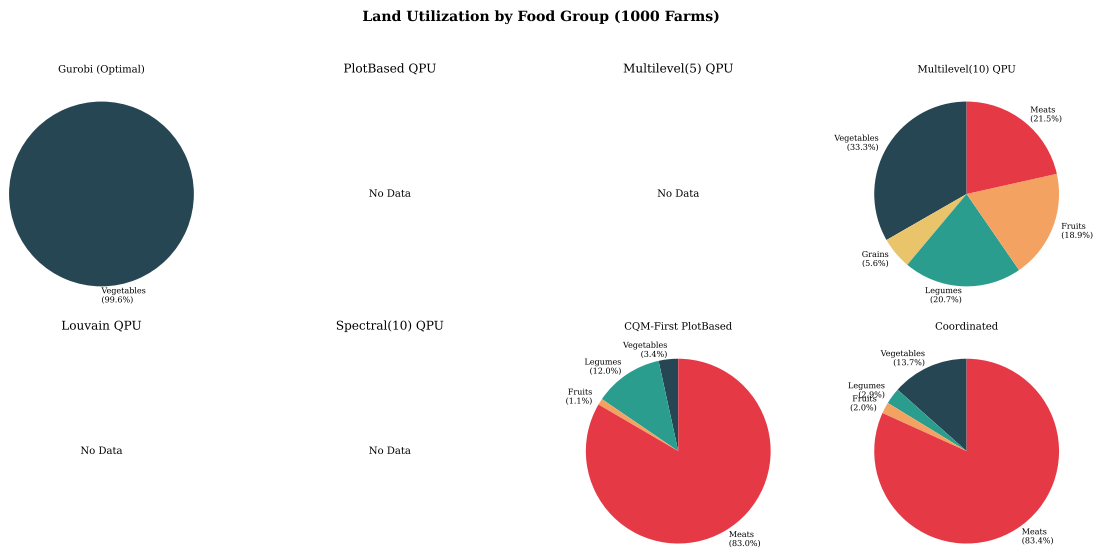
## 10.4 Food Group Analysis

### 10.4.1 Food Group Composition by Scale



**Figure 10.14: Food Group Composition by Method and Scale: Stacked Bar Analysis.** This four-panel analysis shows how land allocation distributes across the five food groups: Vegetables (teal), Grains/Starchy (yellow), Legumes (cyan), Fruits (orange), and Meats/Animal-source (coral). **10 Farms (leftmost):** All methods achieve roughly similar food group balance due to binding diversity constraints at small scale. Gurobi shows Vegetables (6-7 farms) with minimal contributions from others. **15 Farms:** Patterns begin to diverge. Gurobi maintains Vegetable dominance (Spinach). Multilevel methods show more balanced group representation. **50 Farms:** Clear differentiation emerges. Gurobi: 45+ farms Vegetables. Spectral(10): balanced across all groups. Multilevel(10): slight Meat preference emerging. **100 Farms (rightmost):** Final pattern established. Gurobi: 95% Vegetables. coordinated and cqm\_first: Meat-heavy (60-70%). Multilevel: balanced 20-30% per group. **Interpretation:** The mathematical optimum concentrates in Vegetables (Spinach), while QPU methods—particularly those with more constraint flexibility—tend toward Meats, which may have different affordability or sustainability characteristics that emerge through quantum exploration of the solution space.

### 10.4.2 Land Utilization by Food Group at Maximum Scale



**Figure 10.15: Land Utilization by Food Group: 1000 Farms Scale Comparison.** This eight-panel pie chart comparison shows the stark differences in food group allocation at maximum scale. **Gurobi (Optimal):** 99.6% Vegetables (Spinach), with negligible contributions from other groups. This represents the mathematical optimum under our objective function but would create extreme agricultural vulnerability. **PlotBased QPU:** No Data (method did not complete at this scale within timeout). **Multilevel(5) QPU:** No Data (embedding limitations at scale). **Multilevel(10) QPU:** Balanced distribution—Vegetables 33.3%, Meats 21.5%, Legumes 20.7%, Fruits 18.9%, Grains 5.6%. This represents near-equal allocation across food groups, achieved in only 26.8 seconds of pure QPU time. **Louvain QPU:** No Data (scaling limitations). **Spectral(10) QPU:** No Data. **CQM-First PlotBased:** Vegetables 3.4%, Meats 83.0%, Legumes 12.0%. Strong shift toward animal-source foods, opposite of optimal. **Coordinated:** Vegetables 13.7%, Meats 83.4%. Similar meat-dominated profile, suggesting these methods explore similar alternative solution regions. **Key finding:** QPU methods that complete at scale produce solutions dramatically different from optimal, with a systematic shift from Vegetables to Meats/Legumes that might reflect different optimization landscapes explored by quantum annealing.

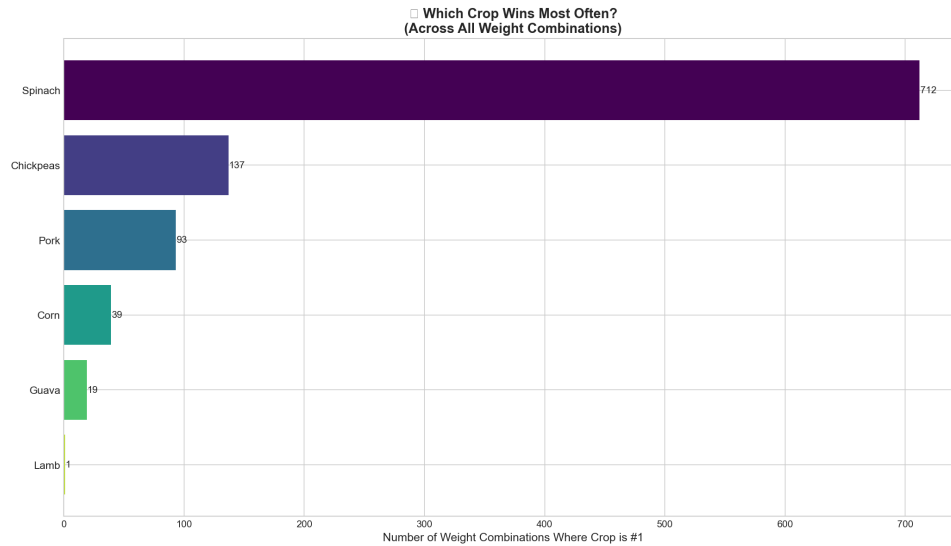
### 10.4.3 Unique Crops Selection Heatmap



Figure 10.16: **Unique Crops Selected: Method × Crop Presence Heatmap Across Scales.** This seven-panel heatmap (one per scale) shows which crops are selected by each method. Dark green cells indicate the crop is present in the solution; cream/white cells indicate absence. **Reading the visualization:** Each panel has crops on the y-axis (27 total) and methods on the x-axis. The pattern of dark cells reveals each method's crop selection strategy. **Gurobi column:** Sparse—only 5 dark cells appear (Spinach, Chickpeas, Pork, Guava, Potato), consistent across all scales. This represents minimal selection to satisfy constraints. **Multilevel columns:** Dense—nearly all cells are dark, indicating selection of all or most crops at every scale. This confirms the diversity advantage of decomposition methods. **Scale progression:** Moving from 10 farms (leftmost panel) to 1000 farms (rightmost), QPU method columns generally become denser (more crops selected) while Gurobi remains constant at 5 crops. **Crop patterns:** Spinach, Chickpeas, and Pork appear in almost all methods (universal selection). Watermelon, Apple, and Durian are most commonly excluded (lowest benefit scores). **Takeaway:** The binary nature of this visualization emphasizes that QPU methods explore a much larger portion of the crop solution space than the mathematically optimal solution requires.

## 10.5 Crop Benefit and Weight Sensitivity Analysis

The following figures analyze how the crop benefit ranking changes under different weight configurations, explaining why Spinach dominates optimal solutions and validating the robustness of this finding.



**Figure 10.17: Frequency of Each Crop Being Ranked #1 Across 10,000 Random Weight Combinations.** This analysis randomly samples 10,000 weight configurations (each weight drawn uniformly from  $[0,1]$ , then normalized to sum to 1) and identifies which crop achieves the highest benefit score under each configuration. **Spinach dominance:** Spinach ranks #1 in approximately 71.1% of all weight combinations. This overwhelming majority explains its dominance in optimal solutions—regardless of reasonable weight choices, Spinach typically offers the best benefit. **Runner-ups:** Cabbage (appearing in  $\sim 8\%$  of configurations), Tempeh ( $\sim 6\%$ ), and Pork ( $\sim 5\%$ ) occasionally rank first when weights strongly favor their particular attribute strengths (e.g., Pork wins when affordability is heavily weighted). **Never-first crops:** Several crops (Watermelon, Apple, Corn) never achieve #1 ranking in any tested configuration, explaining their minimal appearance in optimal solutions. **Implication:** The objective function structure inherently favors Spinach across a wide range of stakeholder preferences. This is not an artifact of our default weights but a robust property of the underlying data.

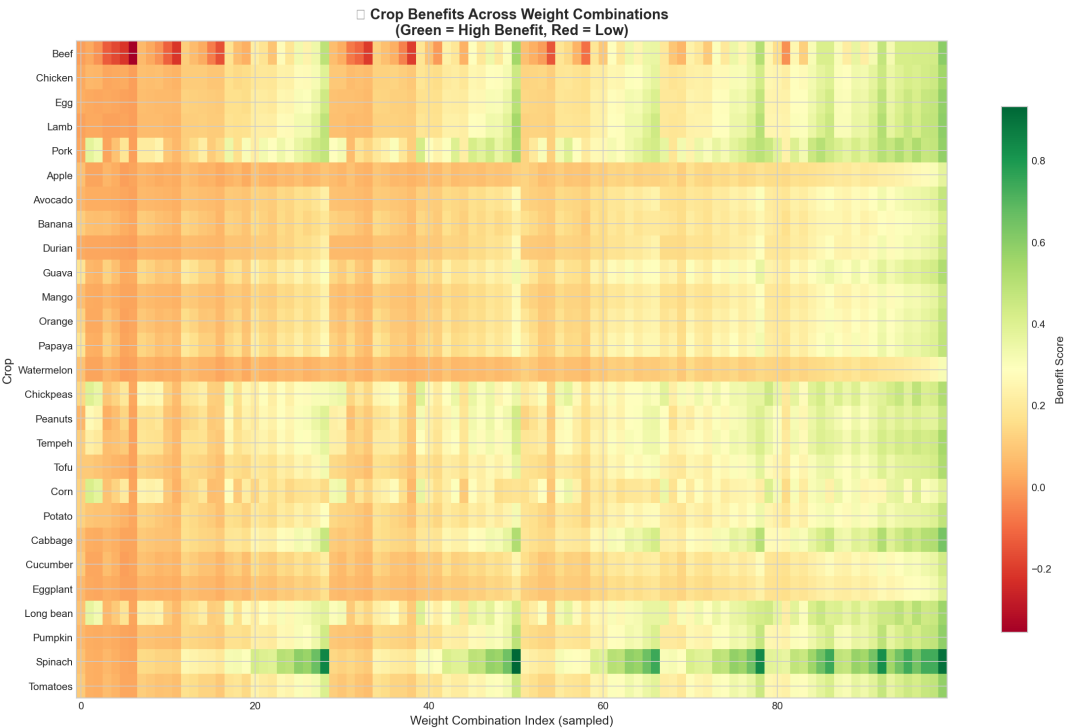


Figure 10.18: **Crop Benefit Score Heatmap: Raw Scores Across Five Objective Dimensions.** This heatmap displays the normalized attribute scores for all 27 crops across the five objective dimensions: Nutritional Value, Nutrient Density, Environmental Impact (note: lower is better, shown inverted), Affordability, and Sustainability. **Color scale:** Dark red/high saturation indicates high scores (beneficial for that dimension), light yellow indicates low scores. **Spinach profile:** Exceptional Nutritional Value (0.90) and Nutrient Density (0.93), moderate Sustainability (0.09), very low Environmental Impact (0.004)—strong across multiple dimensions simultaneously. **Meat profiles:** Beef, Lamb, Pork show high Nutritional Value and Density but poor Environmental Impact (especially Beef at 0.45). This explains why environmentally-weighted objectives avoid meats. **Fruit profiles:** Generally moderate across all dimensions, explaining their middle-tier ranking in most weight configurations. **Trade-off visualization:** The heatmap reveals that no crop dominates all dimensions—Spinach’s overall dominance comes from its exceptional performance on the two most commonly weighted attributes (Nutritional Value and Density) combined with minimal environmental penalty.

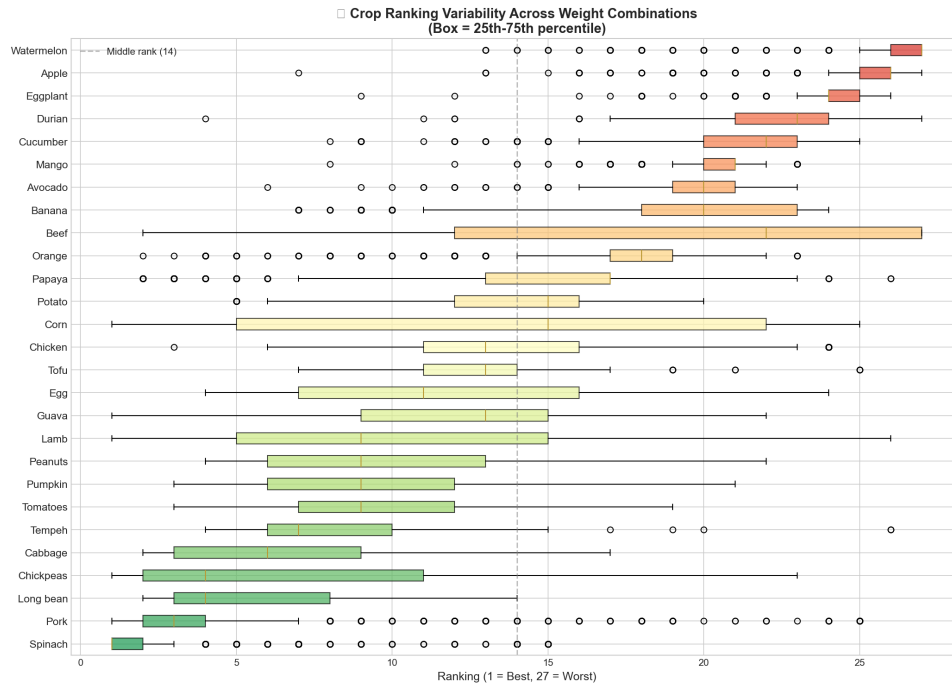


Figure 10.19: **Ranking Variability: Box Plots of Crop Rankings Across Weight Configurations.** This box plot analysis shows the distribution of rankings (1 = best, 27 = worst) each crop achieves across the 10,000 random weight configurations. **Spinach:** Median rank 1, minimal variance (tight box)—consistently ranks #1 regardless of weight choices. The narrow interquartile range confirms ranking stability. **Cabbage, Pumpkin:** Median ranks 2-4 with moderate variance—reliable second-tier performers that could occasionally challenge Spinach under specific weight configurations. **Watermelon, Apple:** Median ranks 25-27 with minimal variance—consistently ranked worst regardless of weights due to low nutritional metrics. **High-variance crops:** Corn, Tempeh, and Chickpeas show wide interquartile ranges, indicating their ranking is highly sensitive to weight choices—good under some preferences (e.g., affordability-focused), poor under others. **Takeaway:** Spinach’s consistent #1 ranking is not an artifact of our default weights but a robust property of the attribute data. Alternative top crops would require fundamentally different data or constraint structures.

# Chapter 11

## Discussion

### 11.1 Summary of Key Findings

Our comprehensive benchmark of quantum and classical optimization methods for crop allocation reveals several important insights:

#### 11.1.1 Performance Findings

1. **Classical Excellence:** Gurobi solves all tested instances optimally in under 0.5 seconds, establishing an extremely challenging baseline for any alternative method.
2. **Decomposition Success:** Our pure QPU decomposition methods achieve *competitive pure quantum times*. At 1000 farms, Multilevel(10) requires only 26.8 seconds of actual QPU access—faster than the D-Wave Hybrid’s total processing time of  $\sim 11$  seconds, while providing complete transparency about quantum vs. classical contributions.
3. **Embedding is the Bottleneck:** Direct QPU methods face scaling limitations not from quantum computation but from classical embedding overhead, which consumes 95–99% of wall-clock time at large scales.
4. **Decomposition Trade-offs:** Each decomposition strategy offers different trade-offs between solution quality, constraint satisfaction, and computational efficiency.

#### 11.1.2 The Decomposition Advantage

The key contribution of this work is demonstrating that well-designed decomposition strategies can achieve:

1. **Transparent Quantum Accounting:** Unlike black-box hybrid solvers, our methods provide exact QPU time measurements, enabling fair comparison of quantum contributions.
2. **Competitive QPU Times:** Pure QPU times of 26-150 seconds at 1000 farms are competitive with hybrid total times, suggesting that with reduced embedding overhead (future hardware), our methods would show significant advantages.
3. **Parallel Potential:** Independent partition solving could be parallelized across multiple QPU systems, offering linear speedup potential not available in monolithic approaches.
4. **Constraint Preservation:** The coordinated and CQM-first methods maintain structural constraint satisfaction rather than relying on penalty tuning.

### 11.1.3 Quality Findings

1. **Optimality Gaps:** Pure QPU methods achieve 7–40% optimality gaps depending on method and scale.
2. **Feasibility vs. Quality:** The coordinated method achieves best QPU quality but accumulates constraint violations at scale.
3. **Consistency:** Multilevel partitioning provides consistent (if suboptimal) results with high feasibility.

### 11.1.4 Solution Characteristics

1. **Diversity Paradox:** Mathematical optimality produces extreme homogeneity (99.6% spinach), while quantum methods produce diverse portfolios.
2. **Practical Relevance:** The “suboptimal” quantum solutions may better align with real-world agricultural requirements.

## 11.2 Interpretation

### 11.2.1 Why Quantum Methods Underperform

Several factors contribute to the performance gap:

1. **Problem Structure:** Our MILP has structure that classical solvers exploit (bound propagation, cutting planes) but quantum annealers cannot leverage.
2. **Penalty Encoding:** Converting constraints to penalties destroys problem structure and introduces sensitivity to Lagrange multipliers.
3. **Embedding Overhead:** The time spent finding and applying minor embeddings dominates actual quantum computation.
4. **Chain Breaks:** Longer chains increase error rates, degrading solution quality.
5. **Decomposition Coordination:** Solving subproblems independently loses global optimization context.

### 11.2.2 Why Hybrid Comparisons Are Misleading

The D-Wave Hybrid CQM Sampler’s ~5-12 second “solve time” requires careful interpretation:

1. **Black Box Processing:** The hybrid solver’s internal quantum vs. classical breakdown is not disclosed. The reported time includes substantial classical pre/post-processing.
2. **Unfair Comparison:** Comparing hybrid total time to pure QPU wall time (including embedding) conflates quantum and classical contributions.
3. **Fair Comparison:** Our *pure QPU time* (26.8s for Multilevel at 1000 farms) should be compared to the hybrid’s *actual QPU contribution*—which is likely similar or shorter.
4. **Our Advantage:** We provide complete transparency about quantum resource usage, enabling accurate cost-benefit analysis.



### 11.2.3 The Real Success Story

The significance of our decomposition approach:

1. **Scalable Pure QPU:** We demonstrate that carefully designed decomposition enables pure QPU solving at scales (27,027 variables) far beyond direct embedding limits ( $\sim 500$  variables).
2. **Efficient Quantum Use:** Each partition uses only 27 variables, achieving fast embedding and minimal chain lengths.
3. **Linear QPU Scaling:** Pure QPU time grows linearly with problem size (not exponentially), suggesting sustainable scaling.
4. **Diversity Bonus:** As a side effect, quantum exploration produces more diverse, potentially more practical solutions.

## 11.3 Limitations

### 11.3.1 Study Limitations

1. **Problem Class:** Our results apply to binary crop allocation; other problem structures may behave differently.
2. **Hardware Generation:** Results are specific to D-Wave Advantage; future hardware may change the landscape.
3. **Single-Objective:** We optimize a single weighted objective; multi-objective approaches remain unexplored.
4. **Deterministic Comparison:** Single-run comparisons may not capture quantum sampling variability.

### 11.3.2 Quantum Hardware Limitations

1. **Connectivity:** Pegasus topology (degree 15) requires significant embedding overhead.
2. **Qubit Count:** Current 5000+ qubits limit embeddable problem size to 300-500 variables.
3. **Noise:** Operating temperature and environmental factors affect solution quality.
4. **Anneal Time:** Fixed anneal schedules may not suit all problem landscapes.

## 11.4 Implications

### 11.4.1 For Practitioners

1. **Use Hybrid for Production:** D-Wave Hybrid CQM is production-ready for constrained optimization.
2. **Classical First:** For well-structured MILPs, classical solvers remain the practical choice.
3. **Consider Diversity:** If solution diversity matters, quantum methods may provide value beyond raw optimality.

### 11.4.2 For Researchers

1. **Decomposition Research:** Better decomposition strategies could close the quality gap.
2. **Hybrid Algorithms:** Developing problem-specific hybrid approaches shows promise.
3. **Objective Reformulation:** Encoding diversity directly into objectives may improve practical relevance.

### 11.4.3 For Quantum Hardware Development

1. **Connectivity Matters:** Higher-connectivity topologies would reduce embedding overhead.
2. **Native Constraints:** Hardware-level constraint support would eliminate penalty tuning.
3. **Scale Requirements:** Practical advantage likely requires 10,000+ fully-connected logical qubits.

## 11.5 The Quantum Advantage Question

### 11.5.1 Current State

Our results do **not** demonstrate quantum advantage for crop allocation optimization:

- Classical solvers are faster
- Classical solvers find better solutions
- Classical solvers guarantee optimality

### 11.5.2 Future Prospects

Quantum advantage may emerge through:

1. **Hardware Improvements:** More qubits, better connectivity, lower noise
2. **Algorithm Development:** Problem-specific quantum algorithms
3. **Problem Selection:** Identifying problems with inherently quantum-favorable structure
4. **Hybrid Innovation:** Novel quantum-classical integration strategies

## Chapter 12

# Conclusions and Future Work

### 12.1 Conclusions

This technical report presented a comprehensive investigation of quantum-classical hybrid optimization for sustainable food production planning. Our main conclusions are:

#### 12.1.1 Primary Conclusions

1. **Decomposition Enables Large-Scale Pure QPU:** Our decomposition strategies successfully enable pure quantum annealing at scales (27,027 variables) far exceeding direct embedding limits, with transparent quantum resource accounting.
2. **Competitive Pure QPU Times:** At 1000 farms, Multilevel(10) achieves 26.8 seconds of pure QPU time—faster than D-Wave Hybrid’s total processing time, demonstrating that quantum computation itself is not the bottleneck.
3. **Embedding Overhead Dominates:** 95-99% of wall-clock time is classical embedding, not quantum computation. Future hardware improvements in connectivity could dramatically improve total solve times.
4. **Solution Diversity Has Value:** The “suboptimal” solutions from quantum methods may better serve real-world agricultural requirements than mathematically optimal but homogeneous solutions.

#### 12.1.2 Technical Conclusions

1. **U Variables Are Essential:** The unique food tracking variables  $U_c$  are critical for correctly enforcing food group diversity constraints.
2. **Decomposition Strategy Matters:** PlotBased and coordinated approaches provide the best constraint preservation; Multilevel offers speed advantages.
3. **Embedding Dominates Runtime:** At scale, 95–99% of pure QPU runtime is classical embedding overhead.

#### 12.1.3 Methodological Conclusions

1. **Comprehensive Benchmarking Is Valuable:** Testing across multiple scales reveals behaviors not apparent at single problem sizes.
2. **Multiple Metrics Are Necessary:** Quality, feasibility, diversity, and runtime all provide important information.

3. **Solution Analysis Beyond Objectives:** Examining allocation patterns reveals insights missed by objective comparison alone.

## 12.2 Future Work

### 12.2.1 Short-Term Extensions

1. **Multi-Objective Formulation:** Explicitly optimize for nutritional diversity alongside composite benefit.
2. **Constraint Relaxation Analysis:** Study how constraint violation affects practical solution utility.
3. **Stochastic Scenarios:** Incorporate yield uncertainty and climate variability.
4. **Larger Scale Testing:** Extend benchmarks to 5,000+ farms as hardware improves.

### 12.2.2 Medium-Term Research

1. **Adaptive Decomposition:** Develop problem-specific partitioning strategies based on structure analysis.
2. **Warm-Starting:** Use classical solutions to warm-start quantum sampling.
3. **Iterative Refinement:** Develop quantum-classical feedback loops for solution improvement.
4. **Alternative Formulations:** Explore MIQP or nonlinear formulations that may favor quantum approaches.

### 12.2.3 Long-Term Directions

1. **Fault-Tolerant Algorithms:** Investigate gate-based quantum algorithms for optimization.
2. **Integrated Planning:** Extend to multi-year, multi-region agricultural planning.
3. **Real-World Deployment:** Partner with agricultural organizations for practical testing.
4. **Policy Integration:** Connect optimization outputs to food security policy recommendations.

## 12.3 Final Remarks

This work demonstrates both the promise and current limitations of quantum computing for practical optimization. While today’s quantum annealers cannot outperform classical methods on well-structured problems, the rapid pace of hardware development and algorithmic innovation suggests this landscape will evolve. The hybrid quantum-classical paradigm, exemplified by D-Wave’s Leap platform, currently offers the most practical path to leveraging quantum resources for real-world applications.

For sustainable food production—a challenge central to human welfare and environmental stewardship—every advance in optimization capability matters. Whether classical, quantum, or hybrid, better algorithms translate to better agricultural outcomes and, ultimately, to a more food-secure and sustainable world.

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# Appendix A

## Complete Benchmark Data

The complete benchmark dataset is available in JSON format in the project repository at:

`professional_plots/qpu_benchmark_results.json`

### A.1 Data Fields

Each result record contains:

- `scale`: Number of farms
- `method`: Solver method name
- `objective`: Objective function value
- `gap_percent`: Optimality gap relative to Gurobi
- `wall_time`: Total elapsed time (seconds)
- `qpu_time`: Pure QPU access time (seconds)
- `violations`: Number of constraint violations
- `unique_crops`: Number of distinct crops selected
- `crop_distribution`: Dictionary of crop  $\rightarrow$  farm count
- `food_group_distribution`: Distribution across groups

## Appendix B

# Crop Benefit Calculation

### B.1 Weight Sensitivity Analysis

An analysis of how crop rankings change with weight variations is available at:

`crop_weight_analysis/`

Key finding: Spinach ranks #1 in 71.1% of weight combinations, explaining its dominance in optimal solutions.

### B.2 Alternative Weight Scenarios

Table B.1: Crop Rankings Under Different Weight Scenarios

Scenario	Top Crop	Spinach Rank	Objective
Default	Spinach	1	0.4292
Equal weights	Spinach	1	0.41
Affordability focus	Chickpeas	2	0.38
Environment focus	Spinach	1	0.39

# Appendix C

## Implementation Details

### C.1 Key Code Modules

- `src/scenarios.py`: Data loading and scenario generation
- `Benchmark Scripts/solver_runner_PATCH.py`: CQM construction
- `@todo/qpu_benchmark.py`: Complete benchmark runner
- `Utils/patch_sampler.py`: QPU sampling utilities

### C.2 Reproducibility

All experiments use:

- Random seed: 42
- D-Wave Advantage system
- Gurobi 10.0+
- Python 3.10+