

Section 3.1

Exercise 53

- a) Total: 2 quarters, 1 penny (51c)
- b) Total: 2 quarters, 1 dime, 1 nickel, 4 pennies (69c)
- c) Total: 3 quarters, 1 penny (76c)
- d) Total: 2 quarters, 1 dime (60c)

Exercise 55

- a) Total: 2 quarters, 1 penny (51c)
- b) Total: 2 quarters, 1 dime, 9 pennies (69c)
- c) Total: 3 quarters, 1 penny (76c)
- d) Total: 2 quarters, 1 dime (60c)

The fewest coins are given when $(n \bmod 25) \bmod 10 \geq 5$

Exercise 56

Proof by counterexample:

15c using the greedy algorithm would draw one 12c coin and 3 pennies, which is four coins, when one dime and one nickel would suffice.

Section 3.2

Exercise 27

- a) $n \log(n^2 + 1) = \mathcal{O}(n 2 \log n) = \mathcal{O}(n \log n)$
 $\mathcal{O}(n^2 \log n + n \log n) = \mathcal{O}(n^2 \log n)$
- b) $(n \log n + 1)^2 = \mathcal{O}(n^2 (\log n)^2) (\log n + 1)(n^2 + 1) = \mathcal{O}(n^2 \log n) \mathcal{O}(n^2 \log n) + \mathcal{O}(n^2 (\log n)^2) = \mathcal{O}(n^2 (\log n)^2)$

Exercise 30

- c) $\lfloor |x + 1/2| \rfloor \leq C|x| \forall x > k$ is true for (among others) $C = 2$ and $k = 2$
 $C \lfloor |x + 1/2| \rfloor \geq |x| \forall x > k$ is true for (among others) $C = 2$ and $k = 2$
Therefore both expressions is $\Theta(x)$
- e) $\log_2(x) = \frac{\log_{10}(x)}{\log_{10}(2)} \implies C \log_{10}(x) \propto \log_2(x)$ which means they are both $\Theta(\log x)$

Exercise 34

- a)

$$\begin{aligned} f(x) &= 3x^2 + x + 1 \\ g(x) &= 3x^2 \\ C_1 |3x^2| &\leq |3x^2 + x + 1| \leq C_2 |3x^2|, \forall x > k \end{aligned}$$

Since k can be chosen freely we choose $k > 0$ to get rid of absolute values

$$3C_1 x^2 \leq 3x^2 + x + 1 \leq 3C_2 x^2, \forall x > k$$

We simplify and get

$$3C_1 \leq 3 + \frac{1}{x} + \frac{1}{x^2} \leq 3C_2, \forall x > k$$

The constants $C_1 = \frac{1}{2}, C_2 = 2, k = 1$ is a solution

b) `loglog` in matlab is a nice way to plot this

```
x = logspace(-1,2);
g = 3*x.^2;
f1 = 1.5*x.^2 + 0.5*x + 0.5;
f2 = 6*x.^2 + 2*x + 2;
loglog(x,f1, x, g, x, f2);
legend('C_1 f(x)', 'g(x)', 'C_2 f(x)');
```

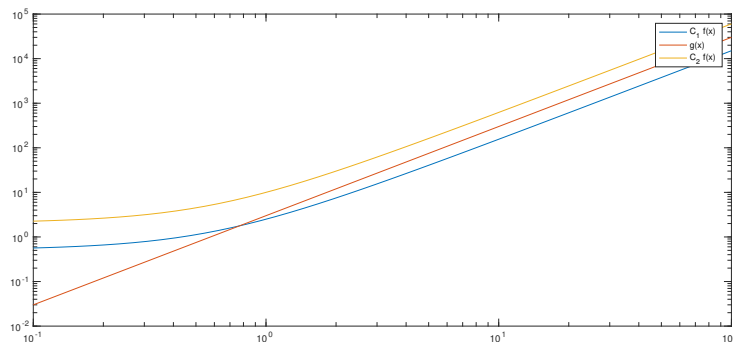


Figure 1: $f(x) = \Theta(3x^2)$ as we see here for $x > 10^0 = 1$

Exercise 42

From $f(x) = \mathcal{O}(g(x))$ we have that

$$\exists C, k \text{ such that } C|g(x)| \geq |f(x)|, \quad \forall x > k$$

If $2^{f(x)} = \mathcal{O}(2^{g(x)})$ then the following must be true

$$C|2^{g(x)}| \geq |2^{f(x)}|$$

$2^a > 0, \forall a$ so we can ignore absolute values

$$\begin{aligned} \log_2 C + \log_2 2^{g(x)} &\geq \log_2 2^{f(x)} \\ \log_2 C + g(x) &\geq f(x) \end{aligned}$$

For large values of x , $\log_2 C$ becomes insignificant, which means that $g(x) \geq f(x)$ must hold for large x if $2^{f(x)} = \mathcal{O}(2^{g(x)})$

Since $f(x) = \mathcal{O}(g(x))$ does not imply $g(x) \geq f(x), \forall x > k$ then it does not follow from $f(x) = \mathcal{O}(g(x))$ that $2^{f(x)} = \mathcal{O}(2^{g(x)})$

Section 4.1**Exercise 11**

- a) $11 + 80 \bmod 12 = 7$
The clock reads 7:00
- b) $12 + 40 \bmod 12 = 4$
The clock reads 4:00
- c) $6 + 100 \bmod 12 = 10$
The clock reads 10:00