Agent-Based Privacy Management for Social Media

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Privacy

- "Right to be let alone" (Warren and Brandeis, Harvard Law Review, 1890)
- "A state in which one is not observed or disturbed by other people." (Merriam-Webster Dictionary)
- "Someone's right to keep their personal matters and relationships secret" (Cambridge Dictionary)
- "The right and ability of an individual to define and live his or her life in a self-determined fashion" (substantive privacy) (Dennedy, Fox, and Finneran, "The Privacy Engineer's Manifesto", 2014)
 - By the individual
 - By others
 - By using the data about the person



Data Privacy

- What data are private?
 - Name, email
 - Financial information
 - Political opinions
 - Sexual orientation
 - Racial or ethnic origin
 - Medical conditions
- Privacy vs. Security

The Organization for Economic Cooperation and Development (OECD) Guidelines

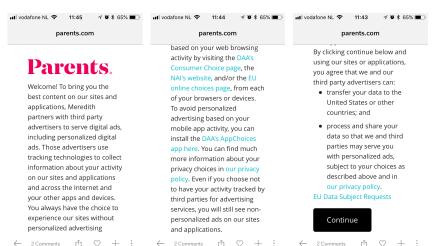
- Collection Limitation: Should be collected lawfully and fairly
- Data Quality Limitation: Should be
 - relevant: Allow appropriate content to be accessed (e.g., age is relevant for checking credit history but phone number is not)
 - accurate: Allow owners to update if necessary
- Purpose Definition Required: Specify explicitly why that information is being shared
- Use Limitation Principle: How it is going to be used (e.g., share with third parties?)
- Accountability Principle: Data sharer will be kept accountable for not abiding with rules

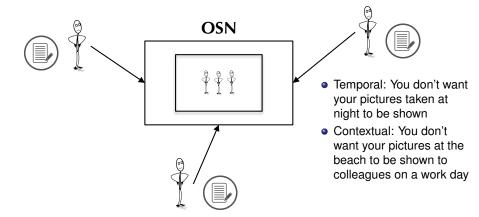
Authorization Types

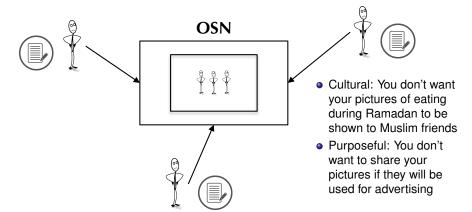
- Opt out (Default is to share)/Opt in (Default is not to share)
- Implied Consent (Your email address appearing on the instructor's list for possible future communication)
- Informed Consent (Explicitly explained how and which information will be used)
- Expressed Consent (Explicitly specified by the user by checking a box or similar)

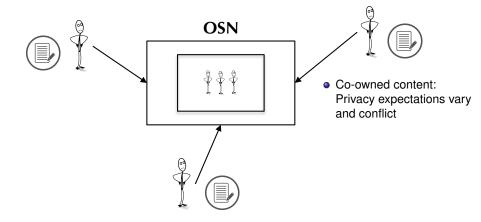
General Data Protection Regulation (GDPR)

Informed Consent: Explains what and how information is used









- Lane v. Facebook: A Class-action lawsuit
 - Sean Lane purchases a diamond ring from Overstock.com.
 - This information shows up on the newsfeed of many of his friends, including his fiancee.
 - This was result of Beacon app, with opt-out privacy options.
 - Facebook ended up paying \$9.5M
 - Moral: Information propagates
- Celebrity Stalking (from ABC News)
 - iPhones embed picture locations into the picture (known as geotags)
 - Geotags can easily be deciphered by apps, reveling the location even when not intended
 - Not only bad for celebrities (Craiglist pictures)
 - Moral: Information implies other information

8 / 58

Understanding Privacy Violations

	No inference	Inference
User	(i) OSN showing the user's media without consent or user wrongly configuring privacy constraints	(iii) Identifying user's location from a geotag in the pictures
Others	(ii) Friend tags the user and makes the picture public where the user did not want to be seen	(iv) Friend tags the user revealing friendship status even when the user had hid her friend list

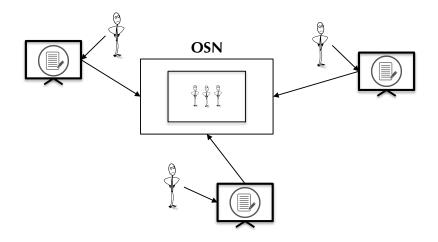
We have conducted an online survey with 330 participants. More than 96% of the participants face privacy violations that occur through inferences.



Dennis wants his friends to see his pictures but not his location.

	No inference	Inference
User	(i) Dennis checks in at a restaurant.	(iii) Dennis shares a picture without declaring his loca- tion. It turns out that his picture is geotagged.
Others	(ii) Charlie shares a picture with everyone. He tags Dennis in it as well.	(iv) Charlie checks in at a restaurant. At the same time, Den- nis shares a picture of Charlie.

Agent-Based Privacy Management



How to Manage the Privacy of Users?

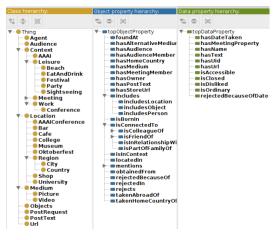
- How to represent the actual privacy preferences of users?
- How to <u>elicit</u> or <u>learn</u> the privacy preferences from users?
- How to <u>advise</u> the users to take actions that are in line with their privacy preferences?
- How to <u>detect</u> potential privacy violations on a user's side?
- How to agree on how a co-owned content will be shared?

Representations of Privacy Preferences

- Access control: Regulate who can view, edit, use resources
- Role-Based: Users take up roles and act in accordance (RBAC)
- Relation-Based: Capture relations among users
- Attribute-Based: Rules based on values of attributes
- Policy-Based: Enable rules to work in harmony

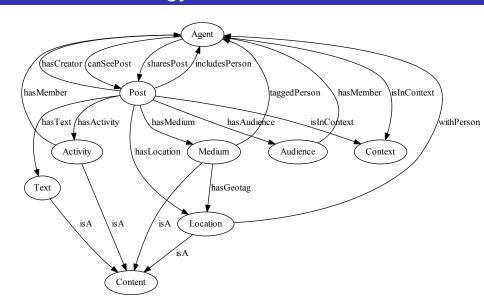
Semantic Representations

Rely on a knowledge representation, such as an ontology, for reasoning on the content.



- Concepts represent a class of individuals (e.g., wig:wig is an instance of Object).
- Object properties relate different individuals with a specific relation (e.g., includesObject relates a :Medium to a :wig).
- Data properties relate data values to individuals (e.g., isOrdinary relates : wig to either true or false).

Content Ontology



Privacy Rules

"Alice does not want her colleagues to see her leisure photos."

```
hasAudience(?postRequest, ?audience),
hasAudienceMember(?audience, ?audienceMember), Leisure(?context),
hasMedium(?postRequest, ?medium),
: isInContext(?medium, ?context),
isColleagueOf(?audienceMember, :alice) →
rejects(:alice, ?postRequest), rejectedIn(?audience, ?postRequest),
rejectedBecauseOf(?audience, ?audienceMember)
```

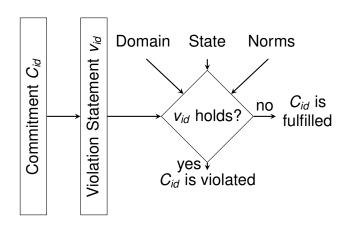
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rejects(:alice, ?postRequest), rejectedIn(?audience, ?postRequest),
rejectedBecauseOf(?audience, ?audienceMember)
```

We can build software agents that can reason on users' privacy preferences.

Detection Privacy Violations with PriGuard¹



¹Nadin Kökciyan and Pınar Yolum. "PriGuard: A Semantic Approach to Detect Privacy Violations in Online Social Networks". In: *IEEE Transactions on Knowledge and Data Engineering* 28.10 (2016), pp. 2724–2737.

Representation of Privacy Requirements

- Commitments are a powerful representation for modeling multiagent interactions.
- Here used to represent the privacy agreement between a user and the OSN.
- A commitment is denoted as a four-place relation:
 C(debtor; creditor; antecedent; consequent)

 $C_1(:osn; :dennis; isFriendOf(:dennis,X), sharesPost(:dennis,P), MediumPost(P); canSeePost(X,P))$ Friends of Dennis are allowed to see medium posts of Dennis

C2(:osn;:dennis; isFriendOf(:dennis,X), sharesPost(:dennis,P), LocationPost(P); not(canSeePost(X,P)))
Friends of Dennis are not allowed to see location posts of Dennis

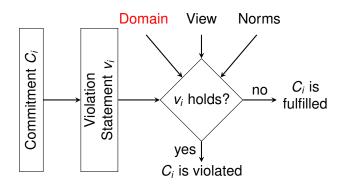
Violation Statements

- A violation occurs when the debtor fails to bring about the condition of a commitment.
- We identify violation statements according to the commitments.
- In a commitment, the condition is true if the antecedent is true that can be represented as the rule: precondition → condition.
- A violation statement is modeled as the negation of this rule:
 - violation: *precondition*,not(*condition*)

 $C_1(:\texttt{osn}; :\texttt{dennis}; is \textit{FriendOf}(:\texttt{dennis}, X), \textit{sharesPost}(:\texttt{dennis}, P), \texttt{MediumPost}(P); \textit{canSeePost}(X, P)) \\ v_1: is \textit{FriendOf}(:\texttt{dennis}, X), \textit{sharesPost}(:\texttt{dennis}, P), \texttt{MediumPost}(P), \texttt{not}(\textit{canSeePost}(X, P)) \\$



The Social Network Domain



The Social Network Domain: Axioms

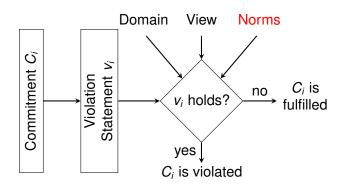
Agent, Post, Audience, Context, Content ☐ T	Leisure, Meeting, Work ⊆ Context
Beach, EatAndDrink, Party, Sightseeing ⊑ Leisure	Bar, Cafe, College, Museum, University ⊑ Location
Picture, Video ⊑ Medium	Medium, Text, Location ⊆ Content
Post $\sqcap \exists sharesPost^-$. Agent $\equiv \exists R_sharedPost$. Self	$LocationPost \equiv \exists R_locationPost.Self$
LocationPost \equiv Post $\sqcap \exists hasLocation$.Location	$MediumPost \equiv Post \sqcap \exists hasMedium.Medium$
TaggedPost \equiv Post \sqcap $\exists isAbout$.Agent	TextPost \equiv Post \sqcap \exists hasText. Text

The Social Network Domain: Axioms

Role Inclusions Role Restrictions

$canSeePost \sqsubseteq U_a$	$\exists canSeePost$. $\top \sqsubseteq Agent$, $\top \sqsubseteq \forall canSeePost$. Post
$hasAudience \sqsubseteq U_a$	\exists has Audience. $\top \sqsubseteq Post$, $\top \sqsubseteq \forall$ has Audience. \exists has Audience. $\top \sqsubseteq \leq 1$ has Audience. \top
$hasGeotag \sqsubseteq U_a$	$\exists hasGeotag. \top \sqsubseteq Medium, \top \sqsubseteq \forall hasGeotag. Location, \top \sqsubseteq \leq 1 hasGeotag. \top$
$hasLocation \sqsubseteq U_a$	$\exists hasLocation. \top \sqsubseteq Post, \top \sqsubseteq \forall hasLocation. Location, \top \sqsubseteq \leq 1 hasLocation. \top$
$hasMedium \sqsubseteq U_a$	$\exists hasMedium$. $\top \sqsubseteq Post$, $\top \sqsubseteq \forall hasMedium$. Medium
hasMember ⊑ U _a	$\exists hasMember$. $\top \sqsubseteq Audience$, $\top \sqsubseteq \forall hasMember$. $Agent$
$isAbout \sqsubseteq U_a$	∃isAbout. T ⊑ Post, T ⊑ ∀isAbout. Agent
$isConnectedTo \sqsubseteq U_a$	$\exists isConnectedTo. \top \sqsubseteq Agent, \top \sqsubseteq \forall isConnectedTo. Agent, isConnectedTo \equiv isConnectedTo$
isFriendOf ⊑ isConnectedTo	
$taggedPerson \sqsubseteq U_a$	$\exists taggedPerson$. $\top \sqsubseteq Medium$, $\top \sqsubseteq \forall taggedPerson$.Agent

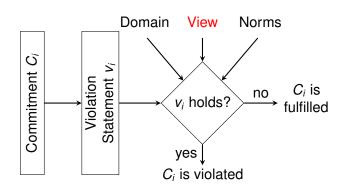
Norms



Norms

```
N_1: sharesPost(X,P) \rightarrow canSeePost(X,P) [Agent can see the posts that it shares.] N_2: sharesPost(X,P) \wedge hasAudience(P,A) \wedge hasMember(A,M) \rightarrow canSeePost(M,P) [Audience of a post can see the post.] N_3: hasMedium(P,M) \wedge taggedPerson(M,X) \rightarrow isAbout(P,X) [Post is about agents tagged in a medium.] N_4: Post(P) \wedge hasMedium(P,M) \wedge hasGeotag(M,T) \rightarrow \texttt{LocationPost}(P) [Geotagged medium gives away the location.]
```

View



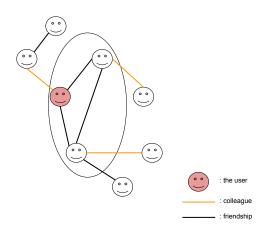
View

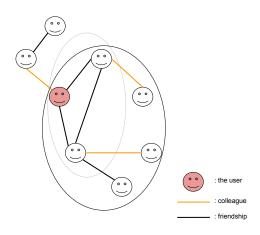
ABSN view captures a given state of the network.

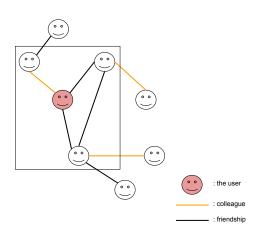
Table: Charlie shares a post :pc1

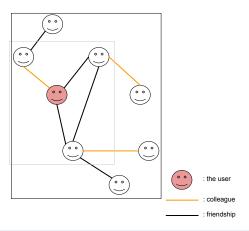
ClassAssertion(Agent :alice)	ClassAssertion(Agent :bob)
ClassAssertion(Agent :charlie)	ClassAssertion(Agent :dennis)
ClassAssertion(Agent : eve)	ClassAssertion(Audience : audience)
ClassAssertion(Post :pcl)	ClassAssertion(Picture:pictureConcert)
ObjectPropertyAssertion(isFriendOf :alice :bob)	ObjectPropertyAssertion(isFriendOf :alice :charlie)
ObjectPropertyAssertion(isFriendOf :bob :charlie)	ObjectPropertyAssertion(isFriendOf :charlie :dennis)
ObjectPropertyAssertion(isFriendOf :dennis :eve)	
ObjectPropertyAssertion(sharesPost : charlie :pcl)	ObjectPropertyAssertion(hasAudience :pcl :audience)
ObjectPropertyAssertion(hasMedium : pcl :pictureConcert)	ObjectPropertyAssertion(taggedPerson:pictureConcert:alice)
ObjectPropertyAssertion(hasMember : audience : alice)	ObjectPropertyAssertion(hasMember : audience :dennis)
ObjectPropertyAssertion(hasMember : audience : eve)	ObjectPropertyAssertion(hasMember : audience : bob)











Used to extend the current view. At the final extension, we have the *global view*.



Detection Algorithm

Algorithm 1: DepthLimitedDetection (C, m=MAX)

```
Input: C, the commitment to be checked
   Input: m, the maximum number of iterations
   Output: V, the set of privacy violations
   Data: KB, the knowledge base (domain + norms)
 1 S \leftarrow \text{initView}(C.creditor);
2 V \leftarrow \{\}, iterno \leftarrow 0;
3 \ vstatement \leftarrow C.antecedent, not(C.consequent);
 4 while iterno < m do
       KB \leftarrow \mathsf{updateKB}(KB, S);
       V \leftarrow V \cup \mathsf{checkViolations}(KB, vstatement);
       iterno \leftarrow iterno + 1;
       if V = \{\} then
           S \leftarrow \mathsf{extendView}(S);
10
       else
           return V;
11
12 return V:
```

Theoretical Results

Theorem (Soundness)

Given an ABSN that is correctly represented with a KB, and a commitment C that represents a privacy requirement $PR_{a,i}^t$, if DEPTHLIMITEDDETECTION returns a violation, then is Violated $(PR_{a,i}^t, ABSN)$ holds.

Theoretical Results

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Proof: Assume that DEPTHLIMITEDDETECTION detects a violation, which is not true. This may occur only if one of the following holds:

- S contains incorrect information.
- KB does not contain the necessary information.
- vstatement is computed incorrectly so that it does not reflect a privacy violation.



Completeness

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Given a commitment C, DEPTHLIMITEDDETECTION always returns a privacy violation, if one exists.

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Lemma

Given a violation statement of a commitment v_i and a knowledge base KB, if there is a privacy violation in KB, checkViolations returns it.

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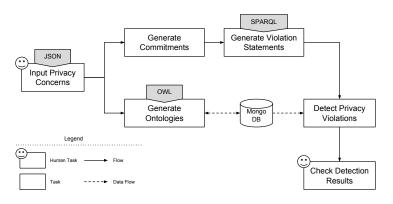
Lemma

Given a violation statement of a commitment v_i and a knowledge base KB, if there is a privacy violation in KB, checkViolations returns it.

Lemma

extendView can eventually create the global view.

A Facebook Application: PriGuardTool²



²Nadin Kökciyan and Pınar Yolum. "PriGuardTool: A Web-Based Tool to Detect Privacy Violations Semantically". In: *Engineering Multi-Agent Systems: 4th International Workshop, EMAS 2016, Singapore, Singapore, May 9-10, 2016, Revised, Selected, and Invited Papers.* Ed. by Matteo Baldoni et al. Springer International Publishing, 2016, pp. 81–98.

27 / 58

Running Example

```
Dennis wants his friends to see his pictures but not his location. He posts a picture without declaring
his location. However, it turns out that his picture is geotagged.
C1(:osn, :dennis, isFriendOf(:dennis, X), isAbout(P, :dennis), LocationPost(P), not(canSeePost(X,P)))
V1 - :osn, :dennis, isFriendOf(:dennis, X), isAbout(P, :dennis), LocationPost(P), canSeePost(X,P))
                      SELECT ?x ?p WHERE {
                       ?x osn:isFriendOf osn:dennis .
                      ?p osn:isAbout osn:dennis .
                      ?p rdf:type osn:LocationPost.
                      FILTER EXISTS (?x osn:canSeePost ?p) }
```

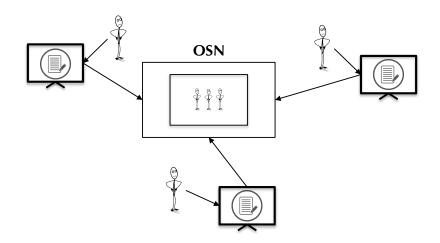
PRIGUARD: Performance Results

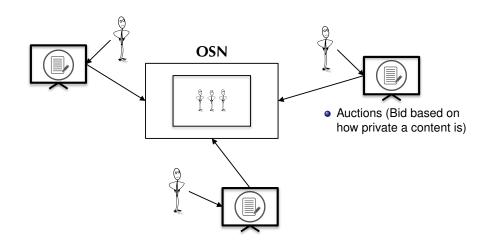
ABSN	depth=0	depth=1	depth=2	G
				•
$(\#\mathcal{A},\!\#\mathcal{R})$	(1,0)	(39,412)	(535,5347)	(535,5347)
G_1 : #Axioms	2175	4267	29959	29959
Time	3ms	4.74ms 30.19ms		29.79ms
$(\#\mathcal{A},\!\#\mathcal{R})$	(1,0)	(51,579)	(1035,27783)	(1035,27783)
G ₂ : #Axioms	2175	5079	125703	125703
Time	2.96ms	5.49ms	123.95ms	122.46ms
$(\#\mathcal{A},\!\#\mathcal{R})$	(1,0)	(123,4199)	(1046,27795)	(4039,88234)
G_3 : #Axioms	2175	20423	125883	403555
Time	3.09ms	18.01ms	121.15ms	530.01ms
$(\#\mathcal{A},\!\#\mathcal{R})$	(1,0)	(37,235)	(848,8543)	(60001,728596)
G ₄ : #Axioms	2175	3535	46463	3636547
Time	3.07ms	4.13ms	47.09ms	18397.26ms
$(\#\mathcal{A},\!\#\mathcal{R})$	(1,0)	(157,2669)	(2787,74217)	(65328,1435168)
G ₅ : #Axioms	2175	14711	332463	6526759
Time	3.11ms	19.03ms	406.91ms	25890.27ms

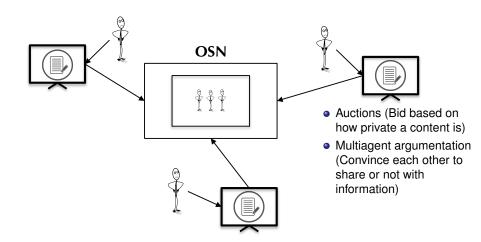
PROTOSS³

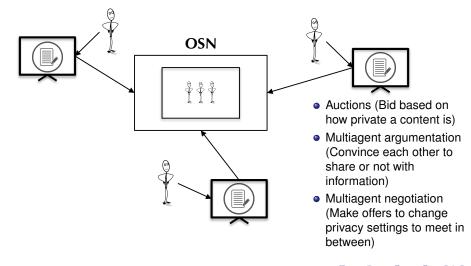
- Rather than checking a single state check the model
- Automatic verification of properties based on the model of a system
- The system is represented as a state transition graph
- The property is represented as a logic formula
- Model checker verifies whether the property holds for the system
- Example:
 - System: the social network (with entire details of relations, agreements, and so on)
 - Property: location of user X is accessible by user Y
 - Model checking will say whether this property holds

³Özgür Kafalı, Akın Günay, and Pınar Yolum. "Detecting and predicting privacy violations in online social networks". In: *Distributed* and *Parallel Databases* 32.1 (2014), pp. 161–190.









Should consider the privacy of relevant users.

- Should consider the privacy of relevant users.
- Should enable relevant users express opinions on a post before it is revealed.

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- Should consider the privacy of relevant users.
- Should enable relevant users express opinions on a post before it is revealed.
- Should be automatic.
- Should enable customized privacy constraints.
- Should protect against violations that occur with inference.

PANO: Privacy Auctioning⁴

- Clarke-Tax mechanism provides an auction mechanism where participants bid for different possible actions in the environment.
- Participants whose bids are decisive on the final action are taxed according to the value they put on it.
- Extentions:
 - Group-wise spending: Earned currencies can only be used in new contents with same co-owners to overcome abuse.
 - Boundaries: Limitations to minimum and maximum bids in order to prevent richer users dominating the decisions.

⁴Onuralp Ulusoy and Pinar Yolum. "PANO: Privacy Auctioning for Online Social Networks". In: AAMAS. 2018, pp. 2103–2105.

PANO Example

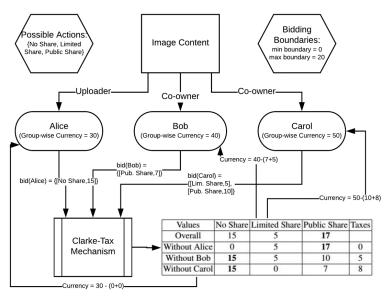
Table 1: Four User Bids for Sharing an Image

Users	No Share	Limited Share	Public Share
Alice	3	5	0
Bob	15	2	0
Carol	5	8	5
Dave	2	6	18

Table 2: Clarke-Tax Mechanism Example - Decision and Taxes

Values	No	Limited	Public	Taxes
Overall	25	21	23	
Without Alice	22	16	23	1
Without Bob	10	19	23	13
Without Carol	20	13	18	0
Without Dave	23	15	5	0

How does PANO work?

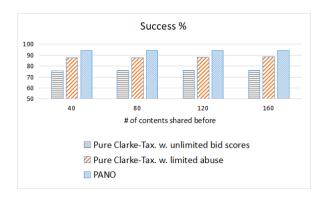


Challenges

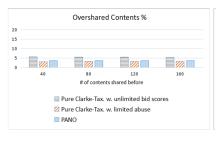
- How to generate a bid given the privacy requirements of the user and importance of a content?
- How to model others to see what bids they give?
- Is it necessary to go into an auction every single time?

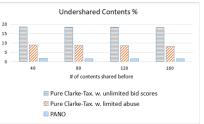
How Well is the Privacy Preserved?

Success is defined as the percentage of users who view the content as specified in the privacy policies of the agents



Evaluation - Oversharing & Undersharing





Convince Others to Keep Private⁵

- Protecting privacy collaboratively
- Users discuss on a post before it is shared.
- Discussion is conducted automatically.
 - Each agent is equipped with an ontology and the semantic rules.
 - Agents discuss on a post by providing each other with arguments using a distributed algorithm.
 - At the end of the discussion, we find the justified arguments.

⁵Nadin Kökciyan, Nefise Yaqlikci, and Pınar Yolum. "An Argumentation Approach for Resolving Privacy Disputes in Online Social Networks". In: ACM Transactions on Internet Technology (2017).

Assumption-Based Argumentation (ABA)

- ABA framework is a four-tuple $\langle \mathcal{L}, \mathcal{R}, \mathcal{A}, \mathcal{C} \rangle$ (Dung *et al.*, 2009).
- Each rule in \mathcal{R} consists of a body $\sigma_1,...,\sigma_m$ and a head σ_0 where $\sigma_1,...,\sigma_m \to \sigma_0$ ($m \ge 0$, $\sigma_i \in \mathcal{L}$).
 - Facts are rules with an empty body
 (e.g., {→ includesObject(:medium,:wig)}).
- Assumption set A includes the weak points of arguments.
- \bullet Contrary mapping ${\cal C}$ includes the contraries of the assumptions.



How to manage the privacy of co-owned data?

Alice would like to share a picture taken with Bob. Bob does not like to share party pictures online.



Derivation of Arguments

An argument has the form $S \vdash^R \sigma$ where $S \subseteq A$, $R \subseteq R$, $\sigma \in \mathcal{L}$.

Derivation of Arguments

An object that can found in a shop is an ordinary object. I_{A} : foundAt(?object, ?shop) \rightarrow isOrdinary(?object, true)

An argument has the form $S \vdash^R \sigma$ where $S \subseteq A$, $R \subseteq R$, $\sigma \in \mathcal{L}$.

Table: SWRL Rules

```
If a post request has a medium including an unordinary object given at ChristmasParty, then it is in Party context. isInContext(?postRequest, ?context), hasMedium(?postRequest, ?medium), includesObject(?medium, ?object), ChristmasParty(?location), obtainedFrom(?object, ?location), isOrdinary(?object, false) \rightarrow Party(?context)

Bob rejects all the post requests in Party context.

P<sub>B</sub>: Party(?context), isInContext(?postRequest, ?context) \rightarrow rejects(:bob, ?postRequest)

rejects(:bob,:pr)

isOrdinary(:wig,true) Party(:context) isInContext(:pr,:context) Party(:context)

foundAt(:wig,:shopx) isOrdinary(:wig,false)
```

Figure: Deduction Trees

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Table: SWRL Rules

```
An object that can found in a shop is an ordinary object. I_a: foundAt(?object, ?shop) \rightarrow isOrdinary(?object, true)
If a post request has a medium including an unordinary object given at ChristmasParty, then it is in Party context. I_a: isInContext(?postRequest, ?context), hasMedium(?postRequest, ?medium), includesObject(?medium, ?object), I_a: ChristmasParty(?location), obtainedFrom(?object, ?location), isOrdinary(?object, false) \rightarrow Party(?context)
Bob rejects all the post requests in Party context. P_B: Party(?context), isInContext(?postRequest, ?context) \rightarrow rejects(:bob, ?postRequest)
rejects(:bob,:pr)
isOrdinary(:wig,true)
Party(:context) isInContext(:pr,:context)
Party(:context)
rejects(:bob,:pr)
isOrdinary(:wig,true)
Party(:wig,false)
risOrdinary(:wig,false)
Figure: Deduction Trees
```

 a_3 : {foundAt(:wig,:Gifty)} $\vdash^{I_{A_1}}$ isOrdinary(:wig,true) b_2 : {isOrdinary(:wig,false)} $\vdash^{I_{B_1} \cup_{i=1}^5 I_i}$ Party(:context) b_3 : {isOrdinary(:wig,false)} $\vdash^{I_{B_1} \cup_{i=1}^5 I_i}$ rejects(:bob.:pr)

Attacks between Arguments

An argument $S_1 \vdash \sigma_1$ can attack another argument $S_2 \vdash \sigma_2$ if and only if σ_1 is the contrary of one of the assumptions in S_2

```
a_3: \{foundAt(:wig,:Gifty)\} \vdash^{I_{A_1}} isOrdinary(:wig,true) 

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c_3 = (isOrdinary(:wig,false),isOrdinary(:wig,true))
```





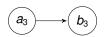
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Argumentation in Action

Table: ABA Specification

```
\mathcal{R} = I_{A_i} \cup I_{B_i} \cup I_{B_i} \cup P_{B_i} \cup_{i=1}^7 r_i
r_1 = \{ \rightarrow isInContext(:pr,:context) \}
r_2 = \{ \rightarrow hasMedium(:pr,:medium) \}
r_3 = \{ \rightarrow includesObject(:medium,:wig) \}
r_4 = \{ \rightarrow \text{ChristmasParty(:location)} \}
r_5 = \{ \rightarrow obtainedFrom(:wig,:location) \}
r_6 = \{ \rightarrow taggedPerson(:medium,:bob) \}
r_7 = \{ \rightarrow hasUrl(:Gifty,:url) \}
A = \{as_1, as_2, as_3, as_4\}
as<sub>1</sub> = foundAt(:wig,:Gifty)
as<sub>2</sub> = not(rejects(:alice,:pr))
as<sub>3</sub> = isOrdinary(:wig,false)
as<sub>4</sub> = isAccessible(:url.false)
C = \{c_1, c_2, c_3, c_4\}
c1 = (foundAt(:wig,:Gifty)=isClosed(:Gifty,true))
c<sub>2</sub> = (not(rejects(:alice.:pr))=rejects(:bob.:pr))
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Table: Arguments

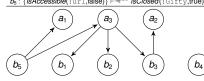
```
f: {} \| \rightarrow \| islnContext(:pr,:context)
b: {} \| \rightarrow \| hasMedium(:pr,:medium)
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b: {} \| \rightarrow \| hasMedium(:pr,:medium)
b: {} \| \rightarrow \| hasDunderson(:medium,:wig)
b: {} \| \rightarrow \| hasDunderson(:medium,:bob)
f; {} \| \rightarrow \| hasDund(:Gifty,:url)
a: {\lind(tagedPerson(:medium,:bob) \rightarrow \| houndAl(:wig,:Gifty) \rightarrow \| houndAl(:wig,:Gifty,:Indise) \rightarrow \| houndAl(:wig,:Indise) \rightarro
```

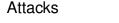
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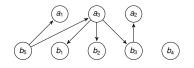
Table: Arguments





- Finds justified argument sets.
- We use credulously admissible argument sets.
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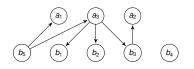


Justified Argument Sets

```
\{\},\{b_5\},\{b_4\},\{b_4,b_5\},\{b_3,b_5\},\{\overline{b_3,b_4,b_5}\},
 \{b_2, b_5\}, \{b_2, b_3, b_5\}, \{b_2, b_4, b_5\},
 b_2, b_3, b_4, b_5, \{b_1, b_5\}, \{b_1, b_4, b_5\},
 \{b_1, b_3, b_5\}, \{b_1, b_3, b_4, b_5\}, \{b_1, b_2, b_5\},
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```

Credulous semantics allow for alternative argument sets

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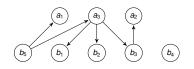
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$$b_3$$
: { $isOrdinary(:wig,false)$ } $\vdash^{I_{B_1} \cup P_{B_1} \cup \frac{5}{i=1}r_i}$ $rejects(:bob,:pr)$ is justified!

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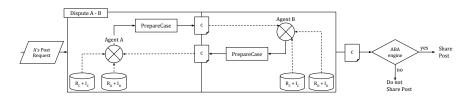
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Distributed argumentation to create ABA specification in a turntaking fashion.

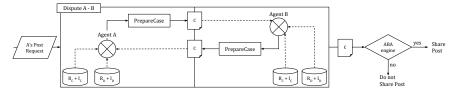
Distributed Privacy Argumentation Framework



- A case is a tuple $\langle \mathcal{R}, \mathcal{A}, \mathcal{F}, \mathcal{C}, status \rangle$
- R is a set of rules, A is a set of assumptions, F is a set of facts, C is the assumption contrary mapping and status is either ongoing or stop.
- A case includes an ABA specification, which is updated in each iteration.



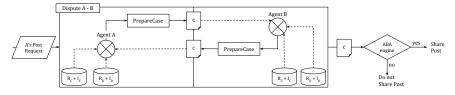
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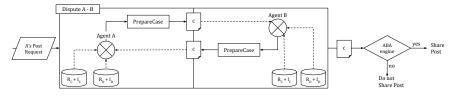
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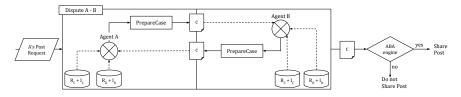
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Decentralized Rules R_D , Centralized Instances I_C I_{B_2} : $hasUrl(?shop,?url),isAccessible(?url,false) <math>\rightarrow isClosed(?shop,true)$ hasUrl(:Gifty,:url)

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Evaluation

- Lack of data: Difficult to collect, impossible to share
- User study
 - Online survey and personal interviews to gather privacy requirements and outcome expectations
 - Participants evaluate the scenarios as neutral observers or by impersonation
 - Example scenarios are shown in stages
 - User expectations are compared with the algorithms outcomes
- Multiagent simulations

Evaluation

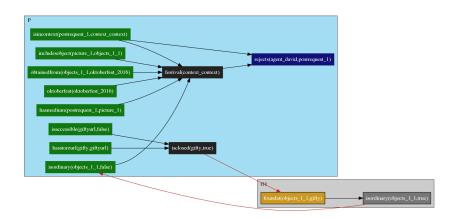
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Table: Personal Interviews and Online Survey Results

Stage	Personal Interviews (36 participants)		Online Survey (68 participants)		PriArg
	Share	Not Share	Share	Not Share	
1	5.55%	94.44%	7.35%	92.65%	Not Share
2	52.77%	47.22%	20.59%	79.41%	Share
3	2.77%	97.22%	7.35%	92.65%	Not Share



Explanation



Negotiation

- Negotiation is mostly used in e-commerce.
- Agents try to reach a mutually acceptable agreement.
- Negotiation technique consists of various components:
 - A *protocol* is a set of rules allowing agents to interact.
 - A strategy (mostly private) is used by agents to make offers and counter-offers.
 - An agreement rule determines when an agreement has been reached.

How to use negotiation technique in privacy context?

- Given a protocol, an agent starts a negotiation with other agents to publish a post.
- Each agent evaluates this post according to its own strategy.
 - It gives a response (accept or deny). The negotiator agent analyzes responses and take an action.
 - It proposes a counter-offer (e.g., a new post), which should be agreed on by agents involved in the counter-offer.

In privacy context, what is ...

An agreement? A protocol? A strategy? An offer? A counter-offer? An agreement rule?



Creating a Post Request

- The content owner puts together the content she wants to publish with the potential audience
- Her agent decides with whom the post is related
 - Sends the post request to those agents
 - Asks for feedback
 - Feedback: I don't want to see Bob in the audience; I don't want a picture on this date to be shown, etc.
 - Feedback calculated based on the Privacy Rules
 - Collects the reasons and revises the post request

Revising a Post Request

- Rejection reasons cannot conflict with each other.
- When a post request is rejected by at least one agent, the negotiator agent:
 - honors every rejection reason,
 - checks whether the resulting post request is reasonable.
- Alternatives: lots of possibilities (using priorities, past experience)

An Example Execution

Iter.	Content	Audience	Asked Agents	Evaluate	Response
1	May 1 picture	Bob, Carol, Errol, Filipo	:carol	$:$ carol $ ightarrow P_{\mathcal{C}_2}$:carol → -date
2	May 28 picture ₁	Bob, Carol, Errol, Filipo	:carol,:bob	:carol $ ightarrow$ N/A, :bob $ ightarrow$ P_{B_2}	$:$ carol \rightarrow 3, $:$ bob \rightarrow -self
3	May 28 picture ₂	Bob, Carol, Errol, Filipo	:carol,:bob	$\texttt{:carol} \to N/A, \texttt{:bob} \to N/A$	$:$ carol \rightarrow 4, $:$ bob \rightarrow 4

Preserving Privacy as Social Responsibility⁶

- Exploit reciprocity as a heuristic (e.g., this time you help me, next time I help you)
- Agents negotiate with each other on their users' preferences
- Negotiation strategies to concede on their preferences
- Given incentives through gamification

⁶Dilara Kekulluoglu, Nadin Kökciyan, and Pınar Yolum. "Preserving Privacy as Social Responsibility in Online Social Networks". In: ACM Transactions on Internet Technology (2018).

Important Criteria

- Concealment of privacy constraints (not being have to explain everything)
- Protection before exposure (checking privacy constraints prior to posting)
- Automating privacy protection (using software agents)
- Fairness (partial improvements instead of all-or-nothing approach)

Research Directions

- Deciphering user's privacy preferences⁷
 - Privacy rules can be identified based on previously shared content using machine learning algorithms
 - Asking other trusted users for privacy recommendation
- Instructing users about privacy preferences⁸
 - User studies show many users do not know what their privacy expectations or even implications
 - Making suggestions based on other trusted users for privacy recommendation or already shared content

Estimation of Images from Tags". In: AAMAS. 2018, pp.4989–1991.

⁷Berkant Kepez and Pınar Yolum. "Learning Privacy Rules Cooperatively in Online Social Networks". In: *PrAISe@ECAI*. 2016. ⁸Abdurrahman Can Kurtan and Pinar Yolum. "PELTE: Privacy

Research Directions

- Managing privacy in IoT
 - Context-Based as opposed to Policy-Based⁹
 - Common-sense reasoning as opposed to personalization
 - Scaling up methods for detection and prediction
- Privacy vs. Utility
 - Agents choosing to violate privacy for a better outcome
 - Metrics to evaluate benefit and cost for privacy
 - Agents learning their evaluations over time

⁹Nadin Kökciyan and Pinar Yolum. "Context-Based Reasoning on Privacy in Internet of Things." In: *IJCAI*. 2017.

Summary: Agents for Privacy

- Represent Privacy Preferences: Semantic representation of policies as those in knowledge representation
- Elicit Privacy Preferences: Machine learning to understand user behavior over time or gamification for understanding users
- Agent-Based Modeling: Agents act on behalf of users to detect privacy violations or avoid them in the first place
- Multiagent Agreement Technologies: Negotiation or argumentation among software agents to reach an agreement for sharing settings