Multi-agent learning

Hypothesis Testing

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Basic concepts

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- 2. A simple hypothesis testing algorithm that leads to ϵ -Nash 1ϵ of the time.

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- 3. Tightening learning parameters (= simulated annealing).

Part I: Motivation



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A hypothesis \approx an opponent model formed on the basis of a large and (therefore hopefully) representative sample of stimuli.

Tim Salmon's experiment

		Player <i>B</i>				
		a	b	C	d	
Player A	a	2,0	2,0	0,2	0,2	
	b	0,2	2,0	0,2	2,0	
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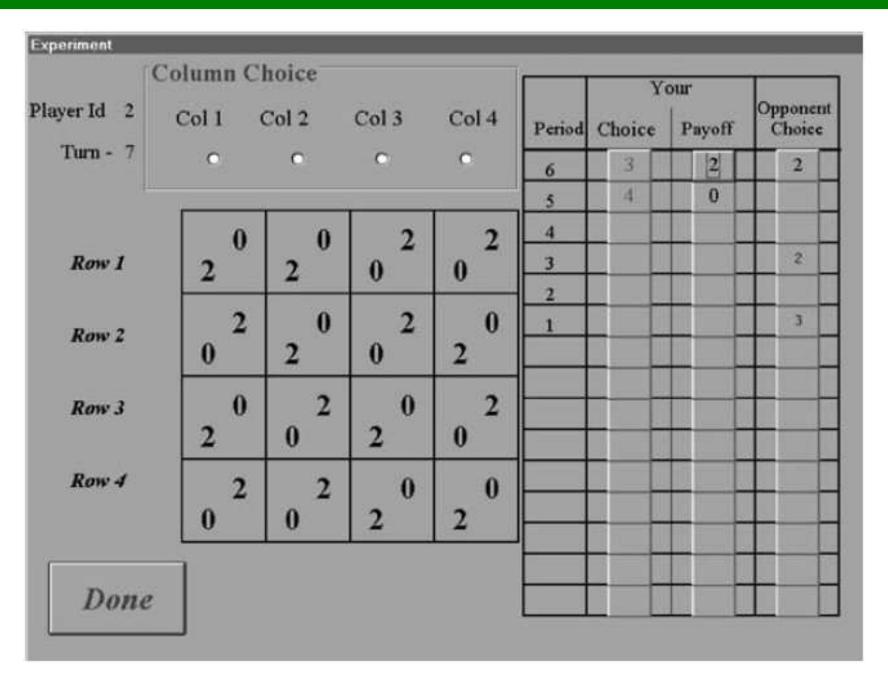
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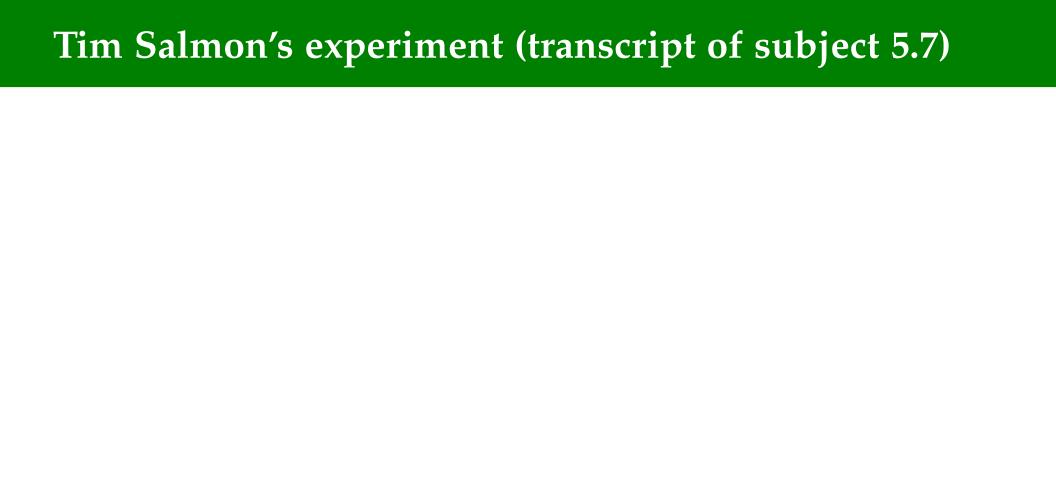
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T. Salmon (2004). "Evidence for Learning to Learn Behavior in Normal Form Games," in: *Theory and Decision*, Vol. **56**, No. 4, pp. 367-404.

Tim Salmon's experiment (user interface)





Round 10: Previous to the last choice I had chosen Column two three times in a row, hoping that once the other player won on that column it or she would choose a different row. However, when that player kept picking either 1 or 2, I decided to switch to Column three, where the layout was the exact opposite, hoping that my hypothesis that the other player was either staying on or switching between one and two was correct. I think it was.

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Round 41: By now I had realized that it was only playing 1 and 4 because they're was no column I could pick that would cover both of them, so it was just a matter of guessing how long it would stay on one or the other. Even though it kept losing by choosing row 1 I thought it might keep trying figuring I would eventually not want to test my luck any more.

Part II: Hypothesis testing

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- The next model is chosen randomly all of the time.



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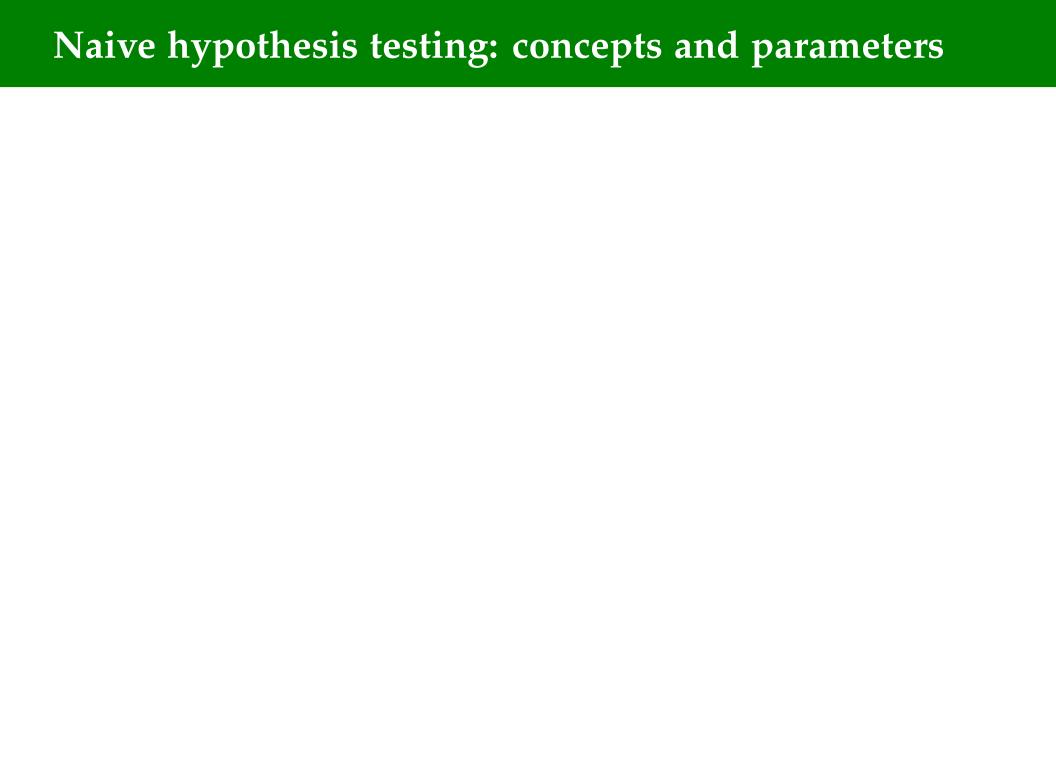
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Part III: Naive hypothesis testing



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- A randomisation factor ρ_i for determining the opponent's model.



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Naive hypothesis testing: algorithm

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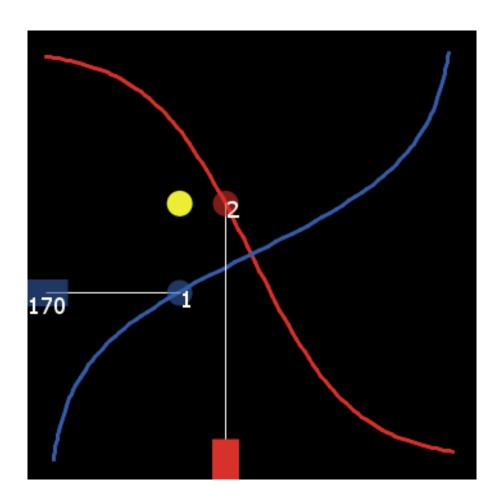
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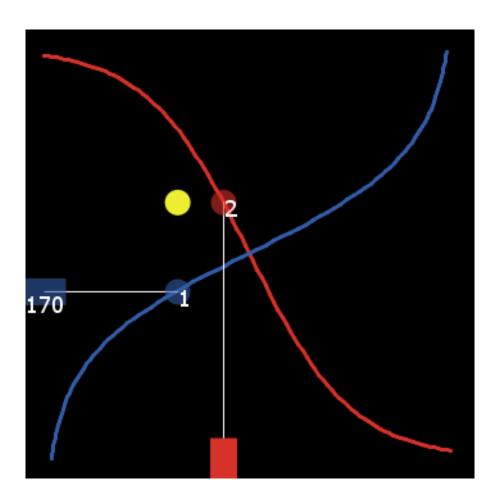
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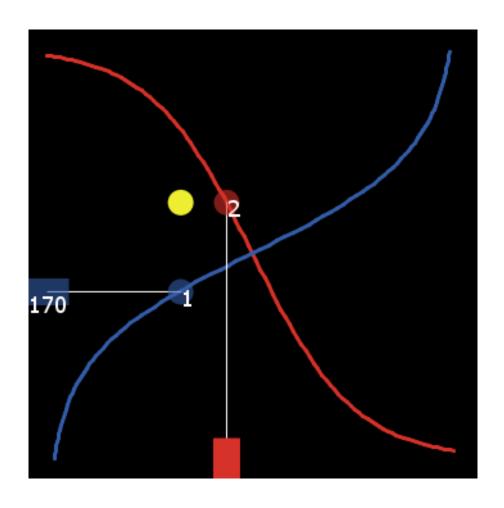
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Else, the hypothesis is rejected, and a new hypothesis is generated. With probability $1 - \rho$ the new hypothesis is taken equal to ϕ^{-it} . Else, it is a random element of Δ_{-i} .



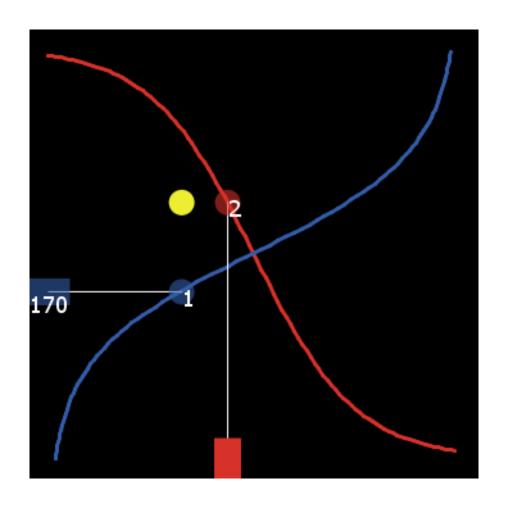


■ Player 1: blue, Player 2: red.



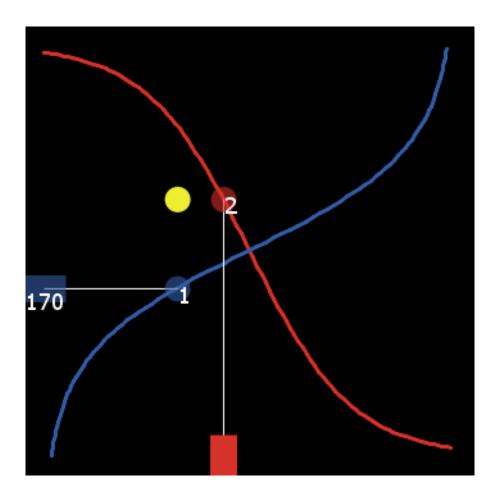
actions = empirical frequency of Player 1.

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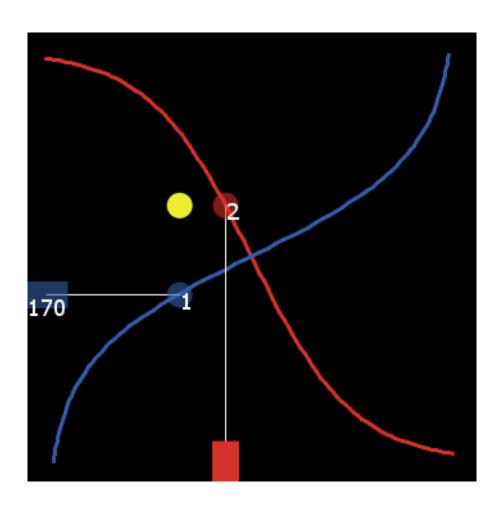
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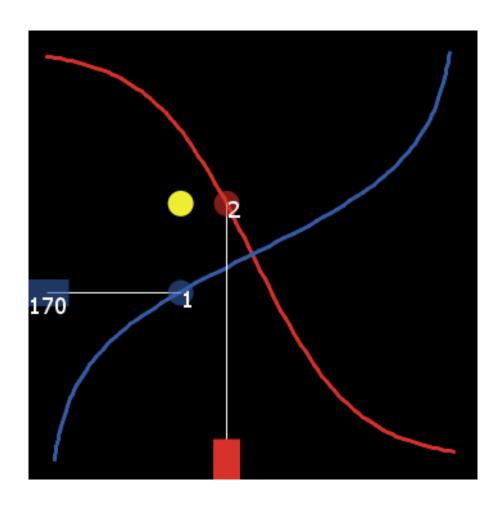
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Vertical axis rectangle:Player 1's model of Player 2's actions = empirical frequency of Player 2.



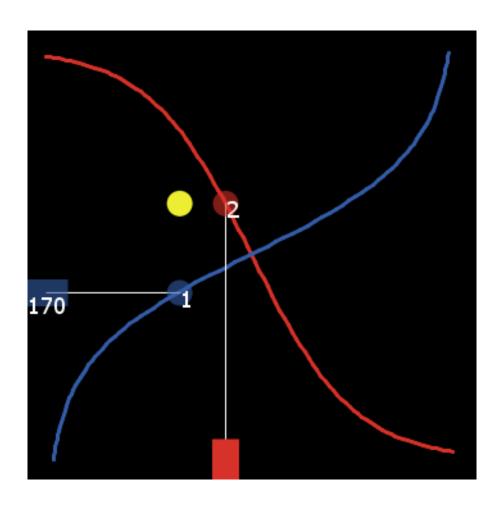
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- — Red **ball**: Player 2's smoothed reply to what it observes.



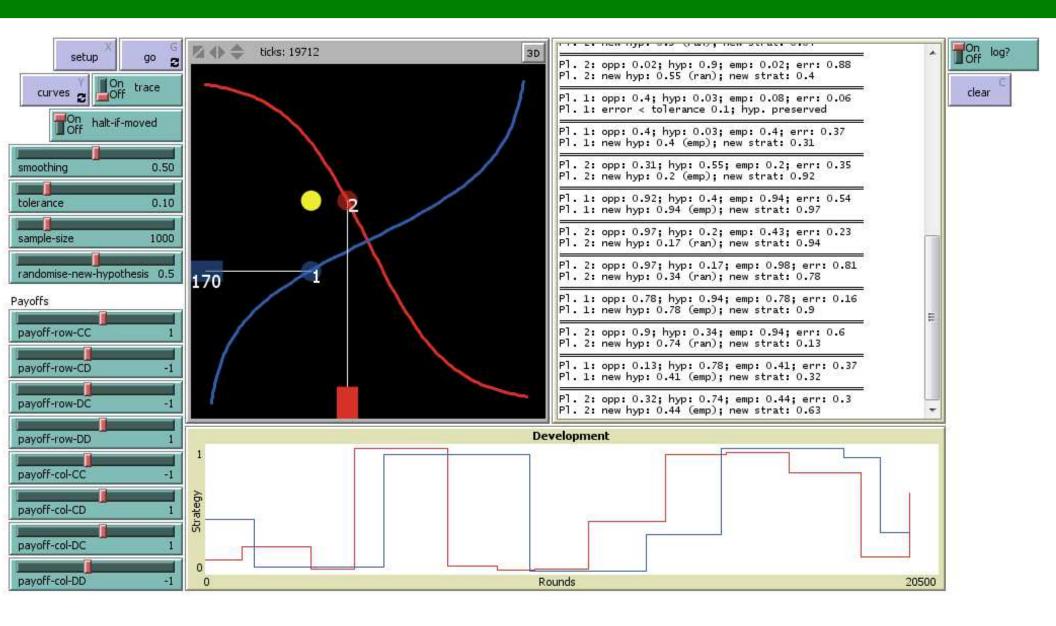
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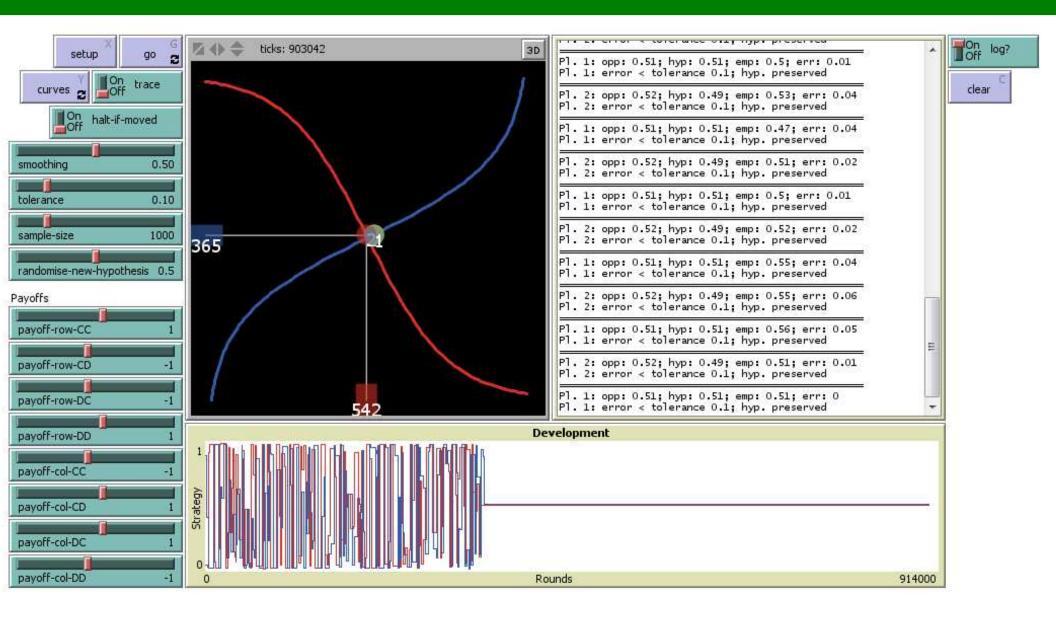


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- Yellow ball: strategy profile.



Demo: the matching pennies game (ϵ -equilibrium)





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- At round 745, your opponent starts sampling.

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- Meanwhile your are still sampling.
 - Notice different opponent behaviour in [1,1746] and [1746, 1812].

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Part IV: Towards ϵ -Nash equilibria

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Part V: Generalisation



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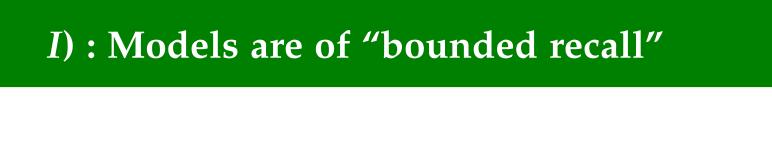
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¹Such families of tests would discriminate very sharply between models that are near the truth (nearer than τ) and those that are far from the truth (further than τ).



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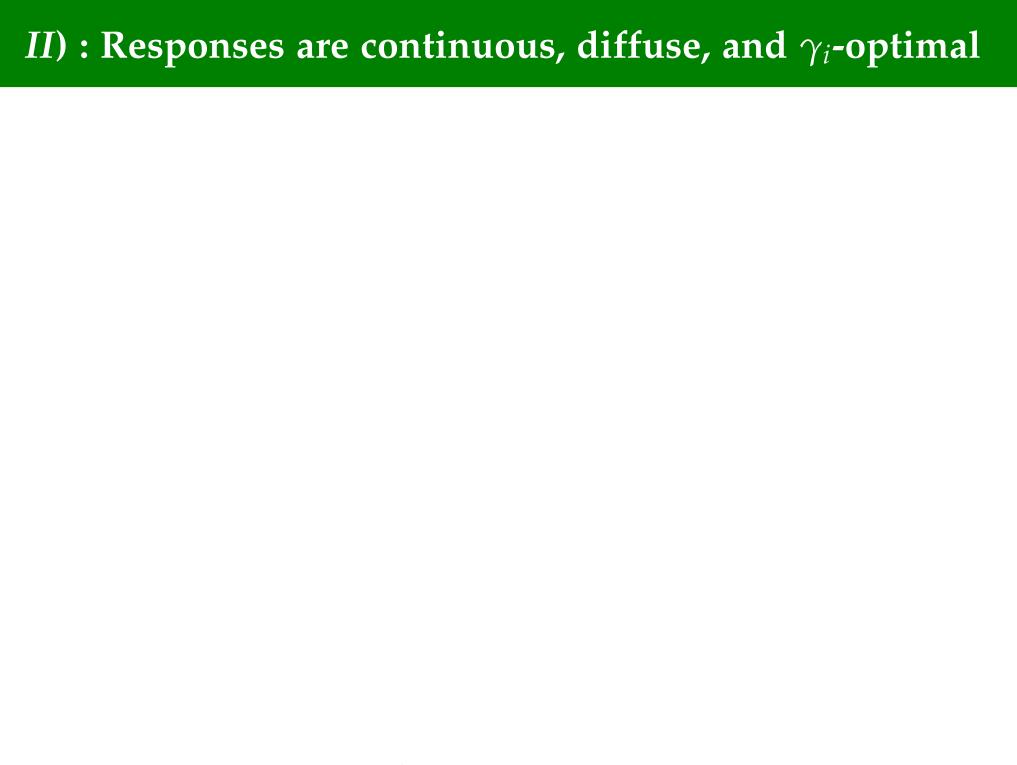
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Such a model revision process is said to be flexible.

Generalised theorem



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Informally: more or less any reasonable hypothesis testing strategy 'works' in the sense that learned behaviours are eventually close to Nash equilibrium most of the time.



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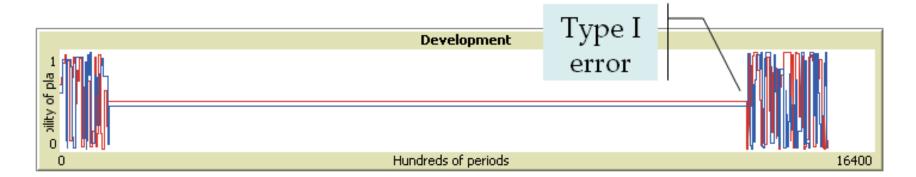
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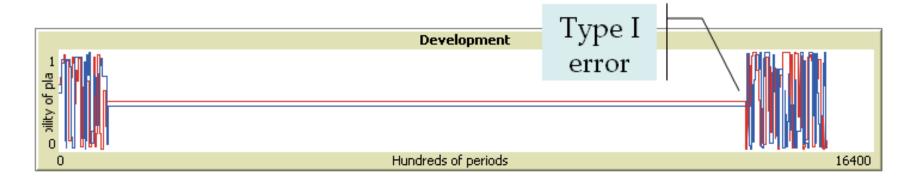
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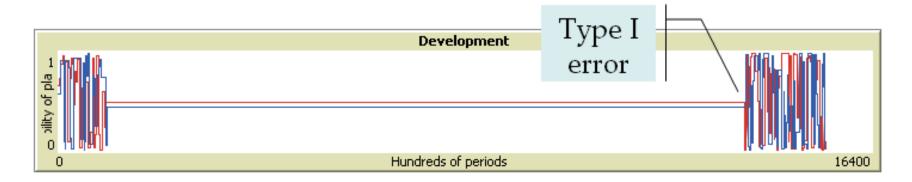
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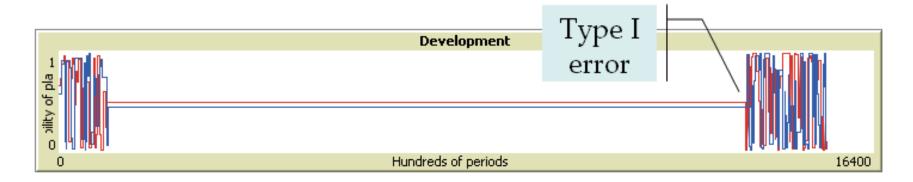


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H. Peyton Young, Dean P. Foster (2003). "Learning, Hypothesis Testing, and Nash Equilibrium," in *Games and Economic Behavior*, Vol. **45**, pp. 73-96.

Part VI: Observations

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Vs. Bayesian play. Theorem (Kalai and Lehrer). Let μ represent the true distribution of play, and let μ_i represent i's view on the distribution of play. If $\mu \ll \mu_i$, then μ_i tends to play like μ . $\mu \ll \mu_i$ is bold assumption.

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- **Young's notion of a "Hunch".** A continuous, diffuse and γ_i -optimal response is determined by two factors.
 - History of play.
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- Annealing. Tighten learning parameters slowly \Rightarrow convergence in probability to the set of Nash equilibria.

Part VII: Extensions



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Theorem (Foster and Young, 2003). Let *G* be a finite n-person game.

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Theorem (Foster and Young, 2003). Let G be a finite n-person game. If the hypothesis testing parameters are tightened sufficiently slowly, the players' responses converge set-wise^a in probability^b to the set of Nash equilibria \mathcal{N} of the repeated game.

^aNot necessarily pointwise.

^bFor all $\epsilon > 0$, $\lim_{n \to \infty} \Pr\{|X_n - X| > \epsilon\} = 0$.



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Theorem (Foster and Young, 2003). For every $\epsilon > 0$ the learning parameters may be chosen so that all players for whom prediction matters by at least ϵ , are ϵ -good predictors. Further, the learning parameters can be tightened at a rate such that all players are good predictors.