

ATL with Imperfect Information

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Strategies under Uncertainty

- ATL includes no notion of **knowledge** (or, dually **uncertainty**)
- ...which makes reasoning in ATL rather unrealistic for MAS
- in this lecture, we show how to introduce knowledge and uncertainty into reasoning about strategic abilities

Motivating Example: Rescue Robots

Properties to express

- the robots **can** rescue person i
- the robots **can** rescue person i , and they **know that they can**
- the robots **can** rescue person i , and they **know how to do it**

Motivating Example: Rescue Robots

Different notions of knowledge (Moore)

- the robots **know that they can rescue i** : *knowledge de dicto*
- know **that** they have an action, but they may not know which action
- the robots **know how to rescue i** : *knowledge de re*
- know **which action** to perform
- the second notion is much more useful

Strategies and Knowledge

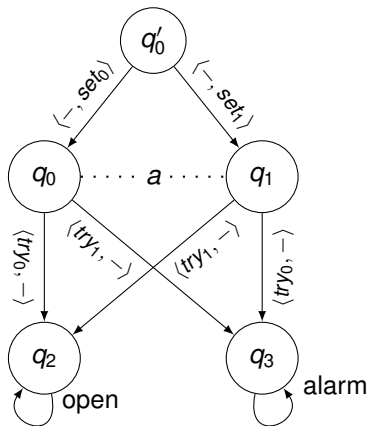
How can we reason about **multi-step games with imperfect information**?

Let's put **ATL** and **epistemic logic** together:

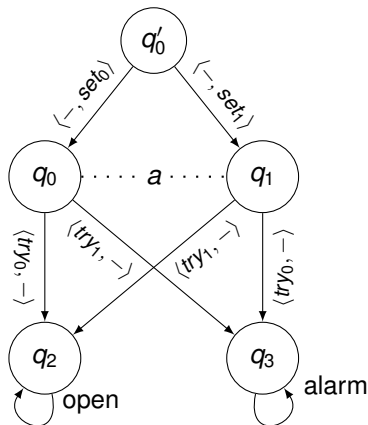
- we extend CGS with **indistinguishability relations** \sim_a , one per agent
- we add **epistemic operators** to **ATL**
- **independent combination**: the semantics is given by the union of semantic clauses of epistemic logic and ATL

\leadsto **Problems!**

Schobbens' Robber



Schobbens' Robber



Robber (a) does not
know what the code is

but:

in q'_0 ,

$\ll a \gg F_{\text{open}}$ is true!

strategy:

$q_0 \mapsto \text{try}_0$,

$q_1 \mapsto \text{try}_1$,

even worse: in q_0 and q_1 ,

$K_a \ll a \gg F_{\text{open}}$ is true

this does not make sense!

Strategies and Knowledge

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Executable strategies = **uniform strategies**

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In order to identify a strategy as successful, the agents must check its outcome paths from **indistinguishable states**

Uniform Strategies

Definition (Uniform strategy)

Strategy s_a is **uniform** iff it specifies the same choices for indistinguishable situations:

- (no recall:) if $q \sim_a q'$ then $s_a(q) = s_a(q')$
- (perfect recall:) if $h \approx_a h'$ then $s_a(h) = s_a(h')$
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A collective strategy is uniform iff it consists only of uniform individual strategies

Strategies and Knowledge

Note:

Having a successful strategy does not imply knowing that we have it!

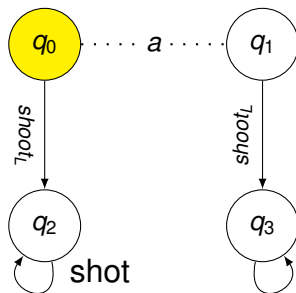
Strategies and Knowledge

Note:

Having a successful strategy does not imply knowing that we have it!

Knowing that a successful strategy exists does not imply knowing the strategy itself!

Example: Poor Duck with Fixed Gun



There is a uniform strategy (same action in q_0 and q_1), but it only *works* from q_0 , and it is not known to the agent that it works

Levels of Strategic Ability

Our cases for $\langle\langle A \rangle\rangle\gamma$ under imperfect information:

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Knowing How to Play

Case [4]: knowing how to play

- single agent case: we take into account the paths **starting from indistinguishable states** (i.e., $\bigcup_{q' \sim_a q} out(q', s_A)$)
- what about coalitions?
- question: **in what sense** should the coalition **know** the strategy?
- common knowledge (C_A), mutual knowledge (E_A), distributed knowledge (D_A)?

Four versions of ATL (Pierre-Yves Schobbens)

- ATL_{IR} : perfect **I**nformation and perfect **R**ecall
- ATL_{Ir} : perfect **I**nformation and imperfect **r**ecall
- ATL_{iR} : imperfect **i**nformation and perfect **R**ecall
- ATL_{ir} : imperfect **i**nformation and imperfect **r**ecall

- ATL_{ir} : Alternating-time logic with **imperfect information** and **imperfect recall** (Schobbens 2004)
- $\langle\langle a \rangle\rangle_{ir} \gamma$: agent a **knows how to play to enforce** γ from all the states she considers possible

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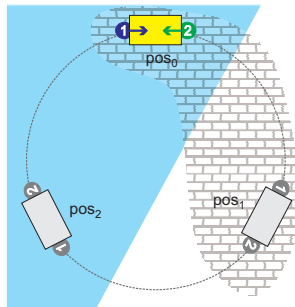
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- $\langle\langle a \rangle\rangle_{ir} \gamma$: agent a **knows how to play to enforce** γ from all the states she considers possible
- what about coalitions?
- $\langle\langle A \rangle\rangle_{ir} \gamma$: agents A know how to play in the sense of “**everybody knows**” (E_A)

Semantics of ATL_{ir}

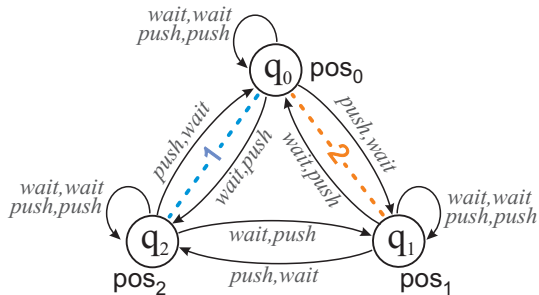
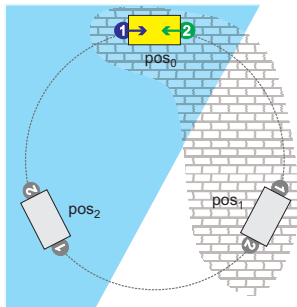
Definition (Semantics of ATL_{ir})

- $M, q \models \langle\langle A \rangle\rangle_{ir} X \varphi$ iff there is a collective **uniform** strategy s_A such that, for every path $\lambda \in \bigcup_{q' \sim_A^E q} out(q', s_A)$, we have $M, \lambda[1] \models \varphi$
- $M, q \models \langle\langle A \rangle\rangle_{ir} G \varphi$ iff there is a collective **uniform** strategy s_A such that, for every path $\lambda \in \bigcup_{q' \sim_A^E q} out(q', s_A)$, we have $M, \lambda[i] \models \varphi$ for all $i \geq 0$
- $M, q \models \langle\langle A \rangle\rangle_{ir} \varphi_1 U \varphi_2$ iff there is a collective **uniform** strategy s_A such that, for every path $\lambda \in \bigcup_{q' \sim_A^E q} out(q', s_A)$, we have $M, \lambda[i] \models \varphi_2$ for some $i \geq 0$, and $M, \lambda[j] \models \varphi_1$ for all $0 \leq j < i$;

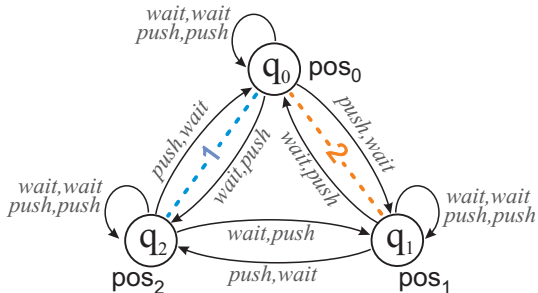
Example: Robots and Carriage



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$$\text{pos}_0 \rightarrow \neg \langle\langle 1 \rangle\rangle_{ir} G \neg \text{pos}_1$$

$$\text{pos}_0 \rightarrow \neg \langle\langle 1, 2 \rangle\rangle_{ir} G \neg \text{pos}_1$$

$$\text{pos}_0 \rightarrow \langle\langle 1, 2 \rangle\rangle_{ir} F \text{pos}_1$$

Fixpoint (Non-)Equivalences

Interesting: $\langle\langle A \rangle\rangle_{ir}$ **are not fixpoint operators** any more!

Theorem

*The following formulae are **not** valid for ATL_{ir} :*

- $\langle\langle A \rangle\rangle_{ir} G \varphi \quad \leftrightarrow \quad \varphi \wedge \langle\langle A \rangle\rangle_{ir} X \langle\langle A \rangle\rangle_{ir} G \varphi$
- $\langle\langle A \rangle\rangle_{ir} \varphi_1 U \varphi_2 \quad \leftrightarrow \quad \varphi_2 \vee \varphi_1 \wedge \langle\langle A \rangle\rangle_{ir} X \langle\langle A \rangle\rangle_{ir} \varphi_1 U \varphi_2.$

Fixpoint (Non-)Equivalences

Conjecture

Strategy for A cannot be synthesized incrementally.

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Strategy for A cannot be synthesized incrementally.

Indeed...

Theorem (Schobbens 2004; Jamroga & Dix 2006)

Model checking ATL_{ir} is Δ_2 -complete in the number of transitions in the model and the length of the formula.