# Epistemic Logic with Common and Distributed Knowledge

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# Group or Collective Knowledge

A group of agents A can "know" that  $\varphi$  in several different epistemic modes:

- $E_A\varphi$ : everybody in A knows that  $\varphi$  (or: A have mutual knowledge that  $\varphi$ )
- $C_A\varphi$ : it is a common knowledge among A that  $\varphi$
- $D_A\varphi$ : A have distributed knowledge that  $\varphi$

# Language of epistemic logic with common and distributed knowledge

### Definition (Syntax of ELCD)

$$\varphi := \rho \mid \neg \varphi \mid \varphi \wedge \psi \mid K_i \varphi \mid E_A \varphi \mid C_A \varphi \mid D_A \varphi$$

- where  $p \in \mathcal{PV}$ ,  $i \in Agt$  and  $A \subseteq Agt$
- E<sub>A</sub>φ is not included in the name of the logic because it is easily definable as Λ<sub>i∈A</sub> K<sub>i</sub>φ
- also,  $K_i \varphi$  is definable as  $D_{\{i\}} \varphi$

# Group Knowledge: Semantics

or truth definitions

In state q, everybody in A knows  $\varphi$  if and only if for every state q prime which is indistinguishable from q by a group indistinguishability relation holds  $\varphi$ 

•  $\mathcal{M}, q \models E_A \varphi$  iff  $\mathcal{M}, q' \models \varphi$  for every q' such that  $q \sim_A^E q'$ , where  $\sim_A^E = \bigcup_{i \in A} \sim_i$ 

$$E_{A}\varphi = \bigwedge_{i \in A} K_{i}\varphi$$

•  $\mathcal{M}, q \models C_A \varphi$  iff  $\mathcal{M}, q' \models \varphi$  for every q' such that  $q \sim_A^C q'$ , where  $\sim_A^C$  is the transitive closure of  $\sim_A^E$ 

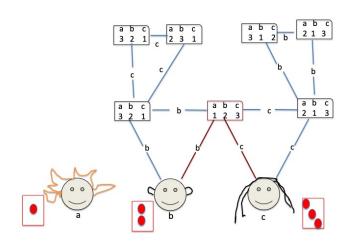
$$C_{A}\varphi = E_{A}\varphi \wedge E_{A}E_{A}\varphi \wedge E_{A}E_{A}E_{A}\varphi \wedge \dots$$

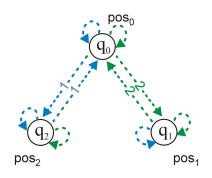
•  $\mathcal{M}, q \models D_A \varphi$  iff  $\mathcal{M}, q' \models \varphi$  for every q' such that  $q \sim_A^D q'$ , where  $\sim_A^D = \bigcap_{i \in A} \sim_i$ 

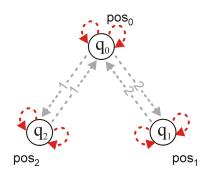
#### **Intuitions**

- $E_A\varphi$ : everyone in A knows, but not necessarily know that others know
- $D_A\varphi$ : if agents in A communicate their knowledge to each other, they will all come to know  $\varphi$  (actually: this is only true for positive knowledge and under some additional conditions)
- $C_A\varphi$ : holds for example when all agents observe the same event  $\varphi$  and see each other observing it; or they are playing a game and have common knowledge of the rules of the game.

 $D_{b,c}(a1 \wedge b2 \wedge c3)$ 





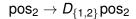






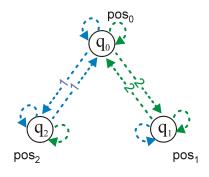


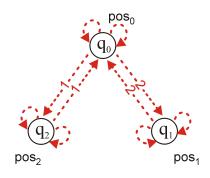


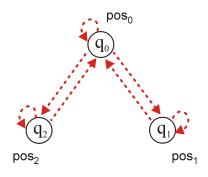


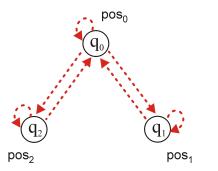




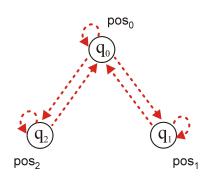






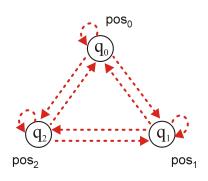


$$\mathsf{pos}_2 \to \neg E_{\{1,2\}} \mathsf{pos}_2$$



$$\mathsf{pos}_2 \to \neg E_{\{1,2\}} \mathsf{pos}_2 \\ \mathsf{pos}_2 \to E_{\{1,2\}} \neg \mathsf{pos}_1$$

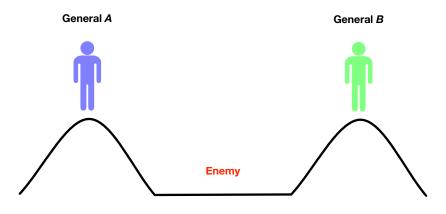
## Example: Common Knowledge



$$\mathsf{pos}_2 o 
eg C_{\{1,2\}} \mathsf{pos}_2 \ \mathsf{pos}_2 o 
eg C_{\{1,2\}} 
eg \mathsf{pos}_1$$

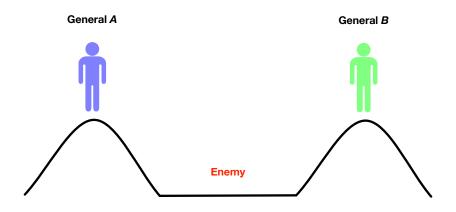
# Coordinated Attack Problem (aka Byzantine Generals)

- only simultaneous attack will succeed
- suppose General A sends General B a proposal to attack at dawn and waits for confirmation from B
- messenger may be intercepted by the enemy



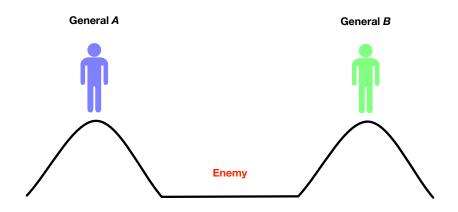
 $\leftarrow$ 

 $E_{A,B}$ attack\_at\_dawn  $\land \neg K_B K_A E_{A,B}$ attack\_at\_dawn



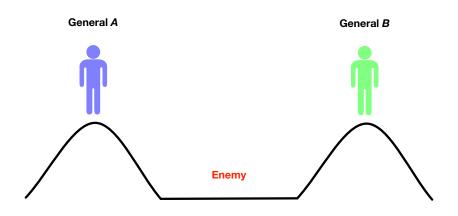
 $\Rightarrow$ 

 $E_{A,B}E_{A,B}$ attack\_at\_dawn  $\land \neg K_AK_BE_{A,B}E_{A,B}$ attack\_at\_dawn

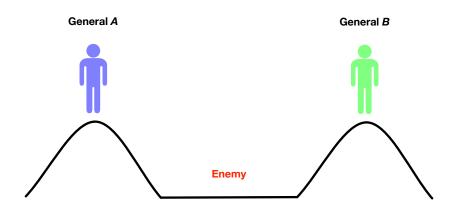


 $\Leftarrow$ 

 $E_{A,B}E_{A,B}E_{A,B}$ attack\_at\_dawn  $\land \neg K_BK_AE_{A,B}E_{A,B}E_{A,B}$ attack\_at\_dawn



 $E_{A,B}^{k}$ attack\_at\_dawn . . .  $\land \neg C_{A,B}$ attack\_at\_dawn



## Application of epistemic logic in distributed systems

- the paper using the notion of knowledge to analyse distributed systems received Gödel prize in 1997:
- see Halpern and Moses Knowledge and Common Knowledge in a Distributed Environment. Journal of the ACM 37(3): 549-587 (1990)