

# Tutorial 12

January 28, 2020

1. Apply the rules of Boolean algebra, verify the following.

(a)  $y \wedge (x \vee (x' \wedge (y \vee y')))) = y$

i. Start from LHS

$$\begin{aligned} y(x + (\bar{x} \cdot (y + \bar{y}))) &= y(x + (\bar{x} \cdot t)) \text{ Inverse law} \\ &= y(x + \bar{x}) \text{ Identity law} \\ &= y(t) \text{ Inverse law} \\ &= y \text{ Identity law} \\ &= RHS \end{aligned}$$

(b)  $((x \wedge y') \wedge (z \vee (x \wedge y')))' = x' \vee y$

i. Start from LHS

$$\begin{aligned} ((x\bar{y})(z + xy'))' &= \overline{(x\bar{y})} + \overline{(z + xy')} \text{ DeMorgan's Law} \\ &= \bar{x} + y + \bar{z} \cdot \overline{(xy')} \text{ DeMorgan's Law} \\ &= \bar{x} + y + \bar{z} \cdot (\bar{x} + y) \text{ DeMorgan's Law} \\ &= \bar{x} + y + \bar{z}\bar{x} + \bar{z}y \text{ Distributive law} \\ &= \bar{x} + \bar{z}\bar{x} + y + \bar{z}y \text{ Associative law} \\ &= \bar{x}(1 + \bar{z}) + y(1 + \bar{z}) \text{ Distributive law} \\ &= \bar{x} + y \text{ Null law} \\ &= x' \vee y \\ &= RHS \end{aligned}$$

2. Simplify the following Boolean functions.

(a)  $f(x, y) = (x \wedge y') \vee (x' \wedge y) \vee (x \wedge y)$

$$\begin{aligned} f(x, y) &= (x \wedge y') \vee (x' \vee x) \wedge y \\ &= (x \wedge y') \vee t \wedge y \\ &= (x \wedge y') \vee y \\ &= (x \vee y) \wedge 1 \\ f(x, y) &= (x \vee y) \end{aligned}$$

$$\begin{aligned}
\text{(b) } f(x, y, z) &= (x' \wedge y' \wedge z') \vee (x \wedge y' \wedge z') \vee (x' \wedge y \wedge z') \vee (x \wedge y \wedge z') \\
f(x, y, z) &= (x' \wedge y' \wedge z') \vee (x \wedge y' \wedge z') \vee (x' \wedge y \wedge z') \vee (x \wedge y \wedge z') \\
&= (x'y'z') + (xy'z') + (x'yz') + (xyz') \\
&= (x' + x)y'z' + (x' + x)yz' \\
&= y'z' + yz' \\
&= z'(y' + y) \\
f(x, y, z) &= z'
\end{aligned}$$

$$\begin{aligned}
\text{(c) } f(x, y, z) &= (x \wedge y) \vee [x \wedge (y \wedge z)'] \\
(x \wedge y) \vee [x \wedge (y \wedge z)'] &= (x \wedge y) \vee [x \wedge (y' \vee z')] \\
&= (x \wedge y) \vee (x \wedge y') \vee (x \wedge z') \\
&= x \wedge (y \vee y') \vee (x \wedge z') \\
&= x \wedge 1 \vee (x \wedge z') \\
&= x \vee (x \wedge z') \\
&= x \wedge (1 \vee z') \\
(x \wedge y) \vee [x \wedge (y \wedge z)'] &= x \\
f(x, y, z) &= x
\end{aligned}$$

3. Find the disjunctive normal form of the Boolean function  $f = f(x, y, z)$  with the given truth table. Then, simplify the expression by constructing the Karnaugh Map

$x$	$y$	$z$	$f(x, y, z)$
0	0	0	1
0	0	1	1
0	1	0	0
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	0

(a)

- i. Find PDNF (all ones):

$$f(x, y, z) = \bar{x}\bar{y}\bar{z} + \bar{x}\bar{y}z + x\bar{y}\bar{z} + xy\bar{z}$$

- ii. Construct karnaugh map

	$y'$	$y'$	$y$	$y$
$x'$	1	1	0	0
$x$	1	0	0	1
	$z'$	$z$	$z$	$z'$

A.

$$\bar{x}\bar{y}\bar{z} + \bar{x}\bar{y}z + x\bar{y}\bar{z} + xy\bar{z} = x'y' + xz'$$

$x$	$y$	$z$	$f(x, y, z)$
0	0	0	1
0	0	1	1
0	1	0	1
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	0
1	1	1	1

(b)

i. Find PDNF

A.  $\bar{x}\bar{y}\bar{z} + \bar{x}\bar{y}z + \bar{x}y\bar{z} + \bar{x}yz$

ii. Construct Karnaugh Map

	$y'$	$y'$	$y$	$y$
A.	$x'$	1	1	1
	$x$	0	0	1
		$z'$	$z$	$z'$

B. Conclusion:  $\bar{x}\bar{y}\bar{z} + \bar{x}\bar{y}z + \bar{x}y\bar{z} + \bar{x}yz \equiv \mathbf{x' + yz}$

4. In the following questions, Karnaugh maps of functions are given, write the simplified Boolean expression for these functions.

	$y'$	$y$
(a)	$x'$	1
	$x$	0
		1

i.  $f(x, y) = (x' \wedge y') \vee (x \wedge y)$

	$y'$	$y$
(b)	$x'$	1
	$x$	1
		0

i.  $f(x, y) = x \wedge y$

	$y'$	$y'$	$y$	$y$
(c)	$x'$	1	1	1
	$x$	1	0	0
		$z'$	$z$	$z'$

i.  $x' + z'$

	$y'$	$y'$	$y$	$y$
(d)	$x'$	1	1	0
	$x$	0	1	0
		$z'$	$z$	$z'$

i.  $x'y' + yz + yz'$

	$y'$	$y'$	$y$	$y$
(e)	$x'$	1	1	1
	$x$	0	0	1
		$z'$	$z$	$z'$

i.  $x' + yz$

(f)

	$y'$	$y'$	$y$	$y$
$x'$	0	1	0	1
$x$	1	1	0	1
	$z'$	$z$	$z$	$z'$

i.  $xy' + yz' + y'z$

5. Draw a Karnaugh map for the Boolean expression whose disjunctive normal forms are as follow. Hence find a simplified version of the expression.

(a)  $f(x, y, z) = (x' \wedge y' \wedge z) \vee (x' \wedge y \wedge z) \vee (x \wedge y \wedge z') \vee (x \wedge y \wedge z)$

i.

	$y'$	$y'$	$y$	$y$
$x'$	0	1	1	0
$x$	0	0	1	1
	$z'$	$z$	$z$	$z'$

A.  $x'z + xy$

(b)  $f(x, y, z, w) = (x \wedge y \wedge z \wedge w) \vee (x \wedge y \wedge z \wedge w') \vee (x' \wedge y \wedge z \wedge w) \vee (x' \wedge y \wedge z \wedge w') \vee (x \wedge y' \wedge z' \wedge w') \vee (x' \wedge y' \wedge z' \wedge w')$

i.

	$z'$	$z'$	$z$	$z$	
$x'$	1	0	0	0	$y'$
$x'$	0	0	1	1	$y$
$x$	0	0	1	1	$y$
$x$	1	0	0	0	$y'$
	$w'$	$w$	$w$	$w'$	

A.  $y'z'w' + yz$

6. Find the disjunctive normal form of the Boolean function  $f(x, y, z)$  with the following truth table and then draw a Karnaugh map to find a simplified version of  $f(x, y, z)$ .

(a)

$x$	$y$	$z$	$f(x, y, z)$
0	0	0	1
0	0	1	0
0	1	0	1
0	1	1	0
1	0	0	0
1	0	1	1
1	1	0	0
1	1	1	1

i. Find PDNF

A.  $\bar{x}\bar{y}\bar{z} + \bar{x}y\bar{z} + x\bar{y}z + xyz$

ii.

	$y'$	$y'$	$y$	$y$
$x'$	1	0	0	1
$x$	0	1	1	0
	$z'$	$z$	$z$	$z'$

iii.  $x'z' + xz$

7. Construct a truth table for the Boolean expression  $(x \wedge (y' \vee z)) \vee (x' \wedge (y \vee z'))$  and hence determine its disjunctive normal form. Draw a Karnaugh map and hence find a simplified version of  $f(x, y, z)$ .

(a)

$x$	$y$	$z$	$y' \vee z$	$y \vee z'$	$x \wedge (y' \vee z)$	$x' \wedge (y \vee z')$	$x \wedge (y' \vee z) \vee x' \wedge (y \vee z')$
0	0	0	1	1	0	1	1
0	0	1	1	0	0	0	0
0	1	0	0	1	0	1	1
0	1	1	1	1	0	1	1
1	0	0	1	1	1	0	1
1	0	1	1	0	0	0	0
1	1	0	0	1	1	0	1
1	1	1	1	1	1	0	1

i.  $x'y'z' + x'y'z + x'yz + xy'z' + xyz' + xyz$

(b)

	$y'$	$y'$	$y$	$y$
$x'$	1	0	1	1
$x$	1	1	1	0
	$z'$	$z$	$z$	$z'$

i.  $y'z' + xz + x'y$