

# Discrete Math - T1

October 16, 2019

1. Determine whether each of the following sentences is a statement.
  - (a)  $y + 3$  is a positive integer
    - i. *NO*
  - (b)  $128 = 26$ 
    - i. *YES*
  - (c)  $x = 26$ 
    - i. *NO*
  - (d) Is 2 a positive number?
    - i. *NO*
2. Give the negation of the following statement.
  - (a)  $2 + 7 \leq 11$ 
    - i.  $2 + 7 > 11$
  - (b)  $2 + 1 = 3$ 
    - i.  $2 + 1 \neq 3$
  - (c) 2 is an even integer and 8 is an odd integer.
    - i. 2 is NOT an even integer OR 8 is NOT an odd number
  - (d) Today is Wednesday
    - i. Today is NOT Wednesday
3. Let  $p$ ,  $q$ ,  $r$  be the propositions
  - $p$  : You have a flu.
  - $q$  : You miss the final exam.
  - $r$  : You pass the course
  - (a)  $p \vee q \vee r$ 
    - i. You have a flu OR you miss the final exam OR you pass the course.
  - (b)  $(p \wedge q) \vee (\sim q \wedge r)$

- i. You have a flu AND you miss the final exam OR you DO NOT miss the final exam AND you pass the course.
- (c)  $\sim p \wedge \sim q \wedge r$ 
  - i. You DO NOT have a flu, AND you DO NOT miss the final exam, AND you pass the course.
- 4. Let  $h$  : “ John is healthy.”  $w$  : “John is wealthy.”  $s$  : “John is wise.” Use the indicated letters and logical connectors to represent the following compound statements.
  - (a) John is healthy and wealthy.
    - i.  $h \wedge w$
  - (b) John is healthy and not wise.
    - i.  $h \wedge \sim s$
  - (c) John is healthy and wealthy but not wise.
    - i.  $h \wedge w \wedge \sim s$
  - (d) John is not wealthy but he is healthy and wise.
    - i.  $\sim w \wedge h \wedge s$
  - (e) John is either wealthy or healthy, or both.
    - i.  $w \vee h$
  - (f) John is wealthy or he is healthy but not wealthy and healthy.
    - i.  $w \vee h$
  - (g) John is neither healthy nor wealthy.
    - i.  $\sim h \wedge \sim w$
  - (h) John is neither healthy, wealthy, nor wise.
    - i.  $\sim h \wedge \sim w \wedge \sim s$
- 5. Determine the truth or falsity of each of the following statement.
  - (a)  $2 > 3$  and 3 is a positive integer.
    - i. False AND True  $\equiv False$
  - (b)  $2 < 3$  or 3 is not a positive integer.
    - i. True OR False  $\equiv True$
  - (c) 2 is a prime but 3 is not a prime.
    - i.  $T \wedge F \equiv False$
  - (d) It is not true that 2 is not a prime or 3 is prime.
    - i.  $\sim (F \vee T) \equiv False$
  - (e) It is false that 2 is prime or multiple of 4 (Ask later:  $\sim T \vee F$ )
    - i.  $\sim (T \wedge F) \equiv True$ .

6. Find the truth value of each proposition if p and r are true and q is false.

$$p = T$$

$$r = T$$

$$q = F$$

(a)  $\sim p \wedge (q \vee r)$

$$\begin{aligned}\sim T \wedge (F \vee T) &\equiv \sim T \wedge T \\ &\equiv F\end{aligned}$$

(b)  $p \wedge (\sim (q \vee \sim r))$

$$\begin{aligned}p \wedge (\sim (q \vee \sim r)) &\equiv T \wedge (\sim (F \vee \sim T)) \\ &\equiv T \wedge (\sim F) \\ &\equiv T \wedge T \\ &\equiv T\end{aligned}$$

(c)  $(r \wedge \sim q) \vee (p \wedge r)$

$$\begin{aligned}(r \wedge \sim q) \vee (p \wedge r) &\equiv (T \wedge \sim F) \vee (T \wedge T) \\ &\equiv (T \wedge T) \vee T \\ &\equiv T \vee T \\ &\equiv T\end{aligned}$$

(d)  $(q \wedge r) \wedge (p \wedge \sim r)$

$$\begin{aligned}(q \wedge r) \wedge (p \wedge \sim r) &\equiv (F \wedge T) \wedge (T \wedge \sim T) \\ &\equiv F \wedge (T \wedge F) \\ &\equiv F \wedge F \\ &\equiv F\end{aligned}$$

7. Construct a truth table for the following compound statements.

(a)  $(p \vee q) \vee (q \vee r)$

i.

p	q	r	$(p \vee q)$	$(q \vee r)$	$(p \vee q) \vee (q \vee r)$
0	0	0	0	0	0
0	0	1	0	1	1
0	1	0	1	1	0
0	1	1	1	0	1
1	0	0	1	0	1
1	0	1	1	1	0
1	1	0	0	1	1
1	1	1	0	0	0

(b)  $(p \downarrow q) \wedge (q \downarrow r)$

$p$	$q$	$r$	$(p \downarrow q)$	$(q \downarrow r)$	$(p \downarrow q) \wedge (q \downarrow r)$
0	0	0	1	1	1
0	0	1	1	0	0
0	1	0	0	0	0
i. 0	1	1	0	0	0
1	0	0	0	1	0
1	0	1	0	0	0
1	1	0	0	0	0
1	1	1	0	0	0

(c)  $(p|q) \wedge r$

$p$	$q$	$r$	$(p q)$	$(p q) \wedge r$
0	0	0	1	0
0	0	1	1	1
0	1	0	1	0
i. 0	1	1	1	1
1	0	0	1	0
1	0	1	1	1
1	1	0	0	0
1	1	1	0	0

(d)  $(p|q) \vee (p|r)$

$p$	$q$	$r$	$(p q)$	$(p r)$	$(p q) \vee (p r)$
0	0	0	1	1	1
0	0	1	1	1	1
0	1	0	1	1	1
i. 0	1	1	1	1	1
1	0	0	1	1	1
1	0	1	1	0	1
1	1	0	0	1	1
1	1	1	0	0	0

8. Determine which of the pairs of statement forms are logically equivalent. Justify your answers using truth tables.

**Note:** The ideal solution is to place both the statements in one table. But we do not have enough space.

(a)  $\sim (p \wedge (\sim (p \wedge q)))$  and  $p \wedge (\sim q \wedge \sim p)$

$p$	$q$	$(p \wedge q)$	$\sim (p \wedge q)$	$p \wedge (\sim (p \wedge q))$	$\sim (p \wedge (\sim (p \wedge q)))$
0	0	0	1	0	<b>1</b>
i. 0	1	0	1	0	<b>1</b>
1	0	0	1	1	<b>0</b>
1	1	1	0	0	<b>1</b>

$p$	$q$	$\sim q$	$\sim p$	$\sim q \wedge \sim p$	$p \wedge (\sim q \wedge \sim p)$
0	0	1	1	1	<b>0</b>
0	1	0	1	0	<b>0</b>
1	0	1	0	0	<b>0</b>
1	1	0	0	0	<b>0</b>

iii.  $\therefore \sim (p \wedge (\sim (p \wedge q))) \not\equiv p \wedge (\sim q \wedge \sim p)$

(b)  $(p \downarrow p) \downarrow (q \downarrow q)$  and  $p \wedge q$

$p$	$q$	$(p \downarrow p)$	$(q \downarrow q)$	$(p \downarrow p) \downarrow (q \downarrow q)$	$p \wedge q$
0	0	1	1	<b>0</b>	<b>0</b>
0	1	1	0	<b>0</b>	<b>0</b>
1	0	0	1	<b>0</b>	<b>0</b>
1	1	0	0	<b>1</b>	<b>1</b>

ii.  $\therefore (p \downarrow p) \downarrow (q \downarrow q) \equiv p \wedge q$

(c)  $(p \vee q) \wedge r$  and  $(p \wedge r) \vee (q \wedge r)$

$p$	$q$	$r$	$(p \vee q)$	$(p \vee q) \wedge r$	$(p \wedge r)$	$(q \wedge r)$	$(p \wedge r) \vee (q \wedge r)$
0	0	0	0	<b>0</b>	0	0	<b>0</b>
0	0	1	0	<b>0</b>	0	0	<b>0</b>
0	1	0	1	<b>0</b>	0	0	<b>0</b>
0	1	1	1	<b>1</b>	0	1	<b>1</b>
1	0	0	1	<b>0</b>	0	0	<b>0</b>
1	0	1	1	<b>1</b>	1	0	<b>1</b>
1	1	0	0	<b>0</b>	0	0	<b>0</b>
1	1	1	0	<b>0</b>	1	1	<b>0</b>

ii.  $(p \vee q) \wedge r \equiv (p \wedge r) \vee (q \wedge r)$

9. Use truth table to determine each of the statement forms below is a tautology, contradiction or contingency.

(a)  $((\sim p \wedge q) \wedge (q \wedge r)) \wedge \sim q$

$p$	$q$	$r$	$\sim p$	$(\sim p \wedge q)$	$(q \wedge r)$	$((\sim p \wedge q) \wedge (q \wedge r))$	$\sim q$	$((\sim p \wedge q) \wedge (q \wedge r)) \wedge \sim q$
0	0	0	1	0	0	0	1	<b>0</b>
0	0	1	1	0	0	0	1	<b>0</b>
0	1	0	1	1	0	0	0	<b>0</b>
0	1	1	1	1	1	1	0	<b>0</b>
1	0	0	0	0	0	0	1	<b>0</b>
1	0	1	0	0	0	0	1	<b>0</b>
1	1	0	0	0	0	0	0	<b>0</b>
1	1	1	0	0	1	0	0	<b>0</b>

ii.  $\therefore ((\sim p \wedge q) \wedge (q \wedge r)) \wedge \sim q \equiv c$  (contradiction)

(b)  $(\sim p \vee q) \vee (p \wedge \sim q)$

$p$	$q$	$\sim p$	$(\sim p \vee q)$	$\sim q$	$(p \wedge \sim q)$	$(\sim p \vee q) \vee (p \wedge \sim q)$
0	0	1	1	1	0	1
0	1	1	1	0	0	1
1	0	0	0	1	1	1
1	1	0	1	0	0	1

ii.  $\therefore (\sim p \vee q) \vee (p \wedge \sim q) \equiv t$  (Tautology)

10. Simplify the following statements using law of logical equivalences.

(a)  $(p \vee q) \wedge \sim (\sim p \wedge q)$

$$\begin{aligned}
 (p \vee q) \wedge \sim (\sim p \wedge q) &\equiv (p \vee q) \wedge (p \vee \sim q) \\
 &\equiv (p + q)(p + \bar{q}) \\
 &\equiv pp + pq + p\bar{q} + p\bar{q} \\
 &\equiv p + pq + p\bar{q} \\
 &\equiv p(1 + q + \bar{q}) \\
 &\equiv p(1) \\
 &\equiv p
 \end{aligned}$$

(b)  $((p \vee q) \wedge (p \vee \sim q)) \vee q$

$$\begin{aligned}
 ((p \vee q) \wedge (p \vee \sim q)) \vee q &\equiv ((p + q)(p + \bar{q})) + q \\
 &\equiv (pp + pq + p\bar{q} + q\bar{q}) + q \\
 &\equiv (p(1 + q + \bar{q}) + 1) + q \\
 &\equiv (p(1)) + q \\
 &\equiv p + q \\
 &\equiv p \vee q
 \end{aligned}$$

11. Use the law of logical equivalences to verify the following logical equivalences.

(a)  $p \vee q \vee (\sim p \wedge \sim q \wedge r) \equiv p \vee q \vee r$

i. LHS (Using a lot of OR distributive law)

$$\begin{aligned}
 p + q + \bar{p}\bar{q}r &\equiv p + \bar{p}\bar{q}r + q \\
 &\equiv p + \bar{p}(\bar{q}r) + q \\
 &\equiv (p + \bar{p})(p + \bar{q}r) + q \\
 &\equiv 1(p + \bar{q}r) + q \\
 &\equiv p + (q + \bar{q}r) \\
 &\equiv p + ((q + \bar{q}) \cdot (q + r)) \\
 &\equiv p + (q + r) \\
 &\equiv p + q + r \\
 &\equiv RHS
 \end{aligned}$$

$$(b) \quad \sim (p \downarrow q) \equiv (\sim p | \sim q)$$

i. LHS

$$\begin{aligned} \sim (p \downarrow q) &\equiv \sim (\sim (p \vee q)) \\ &\equiv p \vee q \end{aligned}$$

ii. RHS

$$\begin{aligned} (\sim p | \sim q) &\equiv \sim (\sim p \wedge \sim q) \\ &\equiv \sim (\sim p \wedge \sim q) \\ &\equiv (p \vee q) \text{ (DeMorgan's)} \\ &\equiv LHS \end{aligned}$$