Séance 3

Le standard de chiffrement actuel : AES

AES: Advanced Encryption System

- ▶ 1996 : Evalutation du DES ⇒ il faut un remplaçant
- Polémique : NSA soupçonnée d'avoir introduit des trappes
- ▶ 1997 : Appel à candidature internationale et publique
 - 15 propositions, 5 finalistes
 - Rijndael (Daemen, Rijmen BE) 10/12/14 rondes Bloc: 128 bits; Clé: 128/192/256 bits
 - 2. Serpent (Anderson, Biham, Knudsen UK) 32 rondes Bloc: 128 bits; Clé: 128/192/256 bits
 - Twofish (Schneier& al US) 16 rondes Bloc: 128 bits; Clé: 128/192/256 bits
 - RC6 (Rivest US) 20 rondes
 Bloc: 128 bits; Clé: 128/192/256 bits
 - MARS (Coppersmith/IBM US) 16 rondes Bloc: 128 bits; Clé: 128→ 448 bits (128+32k bits)

And the winner is

2000 : Standard NIST : AES-Rijndael

Critères de sélection

- Sécurité
- Coût de l'implantation
- Paramètres comme vitesse, latence, complexité
- Flexibilité : Implantation sur des processeurs 8 bits, dans les cartes à puce, dans du matériel dédié.

Les inventeurs

- Joan Daemen (à Gauche)
 - ▶ 1965 : Naissance à Achel
 - 1988 : Rejoint l'UCL (Université Catholique de Louvain)
 - 1995 : Thèse sur un algorithme de chiffrement de sa conception 3 – way
 - 1997 : Concepteur de Rijndael
- Vincent Rijmen (à Droite)
 - 1970 : Naissance à Louvain
 - 1993 : Diplome d'électronique de l'UCL
 - 1997 : Thèse (Cryptanalysis and design of iterated block ciphers)



AES-Rijndael : les grandes lignes

Définition

Algorithme itératif de chiffrement par blocs appliquant sur un même bloc, 10/12/14 fois une fonction de ronde.

Fonction de ronde

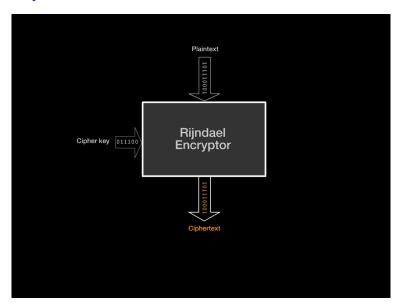
- ByteSub : non-linéarité
- ShiftRow : diffusion entre les colonnes
- MixColumn : diffusion entre les octets à l'intérieur des colonnes
- Addition de la clé de ronde : confusion, dépendance de la clé
- La clé: 128, 192 ou 256 bits
- ▶ Les blocs : 128 bits (fixe pour l'AES-Rijndael), 192, 256 bits
- Blocs découpés en matrice 4 x 4 dont chaque élément est représenté par 8bits
- Opérations sur des octets

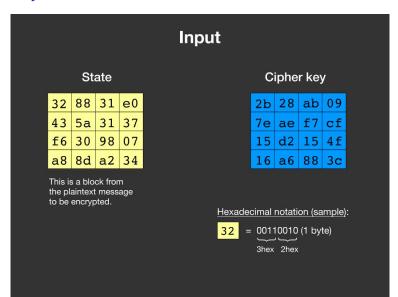
Relation nombre de rondes / Taille de la clé / Taille du bloc

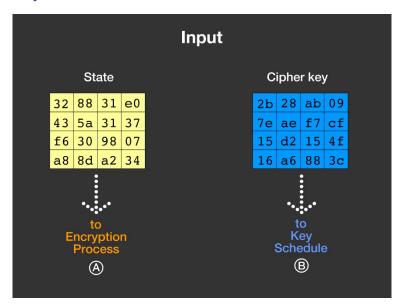
Taille	Rondes Taille de la clé				
du bloc	128 bits	192 bits	256 bits		
128 bits	10	12	14		
192 bits	12	12	14		
256 bits	14	14	14		

Démonstration : Fonctionnement de l'AES

- Démo (www.cryptool.com)
- ► Corps finis, voir slide 56



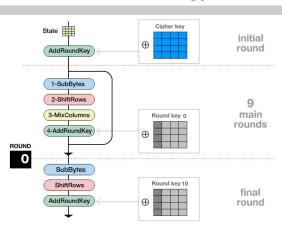




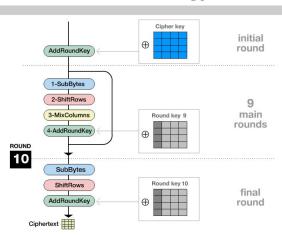


given plaintext block using 4 different transformations in the initial round, the 9 main rounds and the final round)

Encryption Process



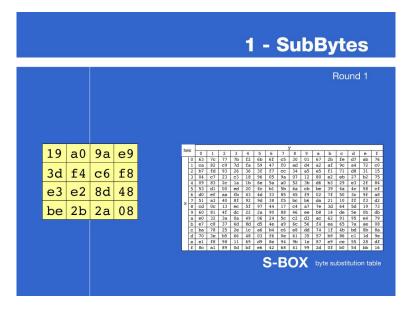
Encryption Process



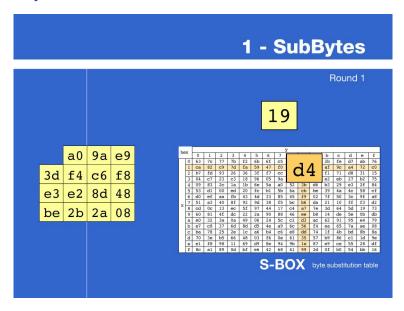
The 4 types of transformations:

1-SubBytes
2-ShiftRows
3-MixColumns

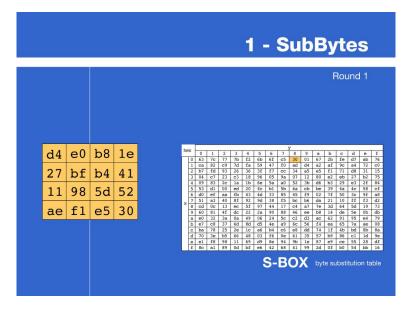
4-AddRoundKey



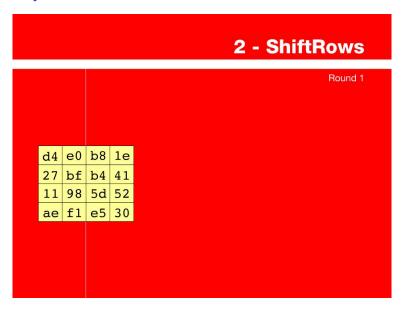
Pierre-Louis CAYREL Protection de l'information 17/56

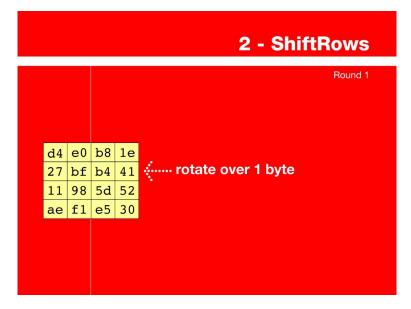


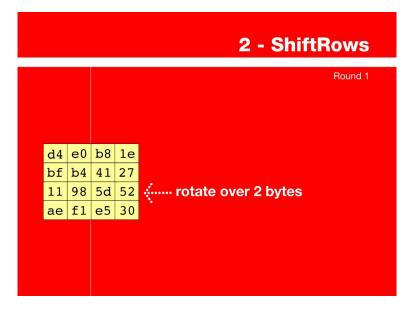
Pierre-Louis CAYREL Protection de l'information 18/56

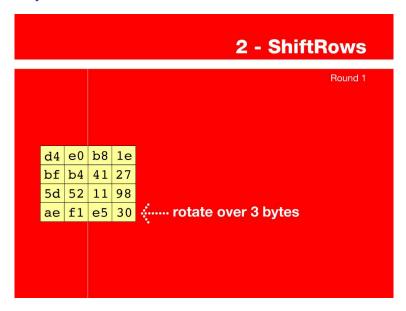


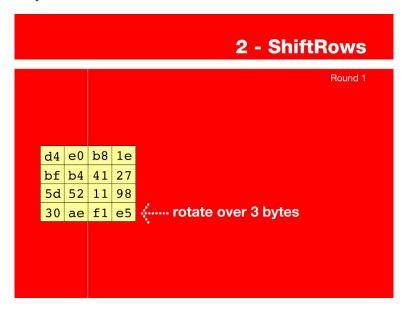
Pierre-Louis CAYREL Protection de l'information 19/56

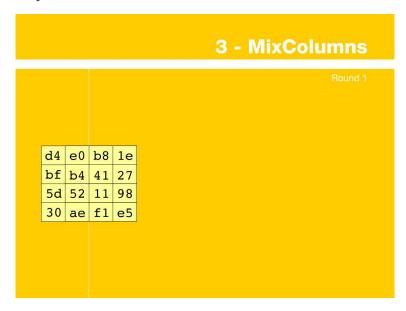


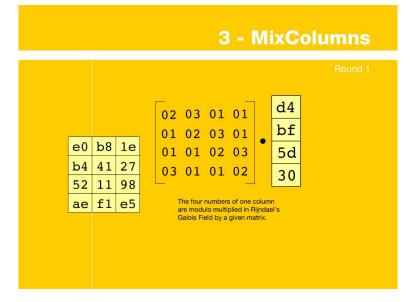




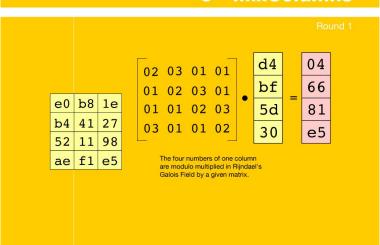


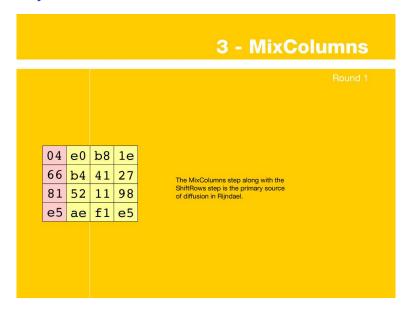




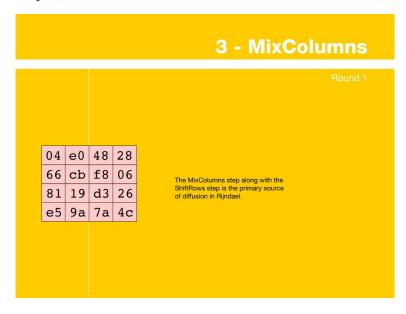






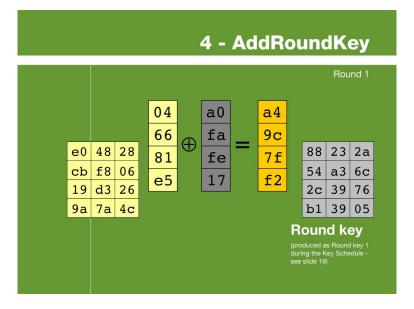


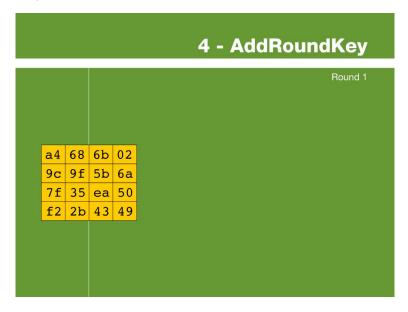
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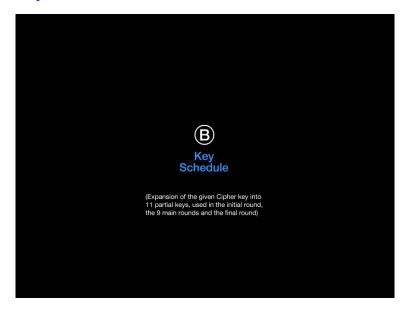


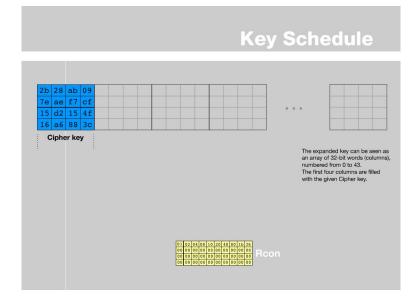


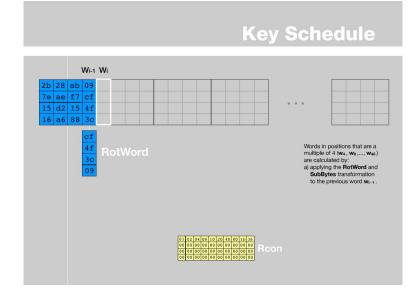
These transformations are applied to the State for 9 more rounds. The final round does not include the MixColumns transformation.

Round 3 76 35 6a 50 42 96 87 53 87 53 42 96 96 87 53 89 £1 1a 3b 3b 89 £1 1a 80 £2 43 7a 7a 7a 7a 7a 7a 7a 7		Start of round	After SubBytes	After ShiftRows	After MixColumns	Round key
Round 1	Input	43 5a 31 37 f6 30 98 07			E	7e ae f7 cf 15 d2 15 4f
Round 2	Round 1	3d f4 c6 f8 e3 e2 8d 48	27 bf b4 41 11 98 5d 52	bf b4 41 27 5d 52 11 98	66 cb f8 06 81 19 d3 26	fa 54 a3 6c fe 2c 39 76
Round 3	Round 2	9c 9f 5b 6a 7f 35 ea 50	de db 39 02 d2 96 87 53	db 39 02 de 87 53 d2 96	4d 4b e7 6b ca 5a ca b0	c2 96 35 59 95 b9 80 f6
Round 4	Round 3	8f dd d2 32 5f e3 4a 46	73 c1 b5 23 cf 11 d6 5a	c1 b5 23 73 d6 5a cf 11	ec 0b c0 25 09 63 cf d0	80 16 23 7a 47 fe 7e 88
Round 5 92 63 b1 b8 7f 63 35 be 4f ft b 68 6c 4f 96 ae d2 fb 64 4f 96 ae d2 fb 64 4f 96 ae d2 fb 64 4f 96 4f 96 ae d2 fb 64 4f 96 4f	Round 4	6c 1d e3 5f 4e 9d b1 58	50 a4 11 cf 2f 5e c8 6a	a4 11 cf 50 c8 6a 2f 5e	d6 31 c0 b3 da 38 10 13	44 52 71 0b a5 5b 25 ad
	Round 5	92 63 b1 b8 7f 63 35 be	4f fb c8 6c d2 fb 96 ae	fb c8 6c 4f 96 ae d2 fb	d1 11 3a 4c a9 d1 33 c0	d1 83 f2 f9 c6 9d b8 15

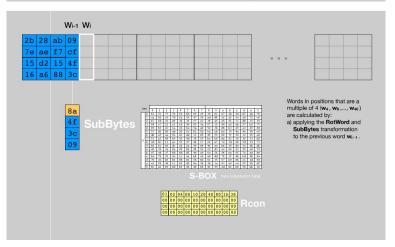
	Start of round	After SubBytes	After ShiftRows	After MixColumns	Round key
Round 6	f1 c1 7c 5d 00 92 c8 b5 6f 4c 8b d5 55 ef 32 0c	a1 78 10 4c 63 4f e8 d5 a8 29 3d 03 fc df 23 fe	a1 78 10 4c 4f e8 d5 63 3d 03 a8 29 fe fc df 23	4b 2c 33 37 86 4a 9d d2 8d 89 f4 18 6d 80 e8 d8	7a fd 41 fd
Round 7	26 3d e8 fd 0e 41 64 d2 2e b7 72 8b 17 7d a9 25	f7 27 9b 54 ab 83 43 b5 31 a9 40 3d f0 ff d3 3f	f7 27 9b 54 83 43 b5 ab 40 3d 31 a9 3f f0 ff d3	14 46 27 34 15 16 46 2a b5 15 56 d8 bf ec d7 43	0e f3 b2 4f
Round 8	5a 19 a3 7a 41 49 e0 8c 42 dc 19 04 b1 1f 65 0c	be d4 0a da 83 3b e1 64 2c 86 d4 f2 c8 c0 4d fe	be d4 0a da 3b e1 64 83 d4 f2 2c 86 fe c8 c0 4d	00 b1 54 fa 51 c8 76 1b 2f 89 6d 99 d1 ff cd ea	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$
Round 9	ea 04 65 85 83 45 5d 96 5c 33 98 b0 f0 2d ad c5	87 f2 4d 97 ec 6e 4c 90 4a c3 46 e7 8c d8 95 a6	87 f2 4d 97 6e 4c 90 ec 46 e7 4a c3 a6 8c d8 95	47 40 a3 4c 37 d4 70 9f 94 e4 3a 42 ed a5 a6 bc	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$
Round 10	eb 59 8b 1b 40 2e a1 c3 f2 38 13 42 le 84 e7 d2	e9 cb 3d af 09 31 32 2e 89 07 7d 2c 72 5f 94 b5	e9 cb 3d af 31 32 2e 09 7d 2c 89 07 b5 72 5f 94		$\begin{array}{c ccccccccccccccccccccccccccccccccccc$
Output	39 02 dc 19 25 dc 11 6a 84 09 85 0b 1d fb 97 32				
	Ciphertext				

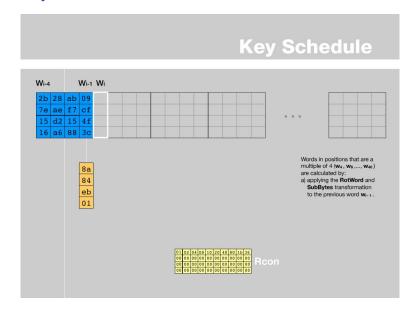


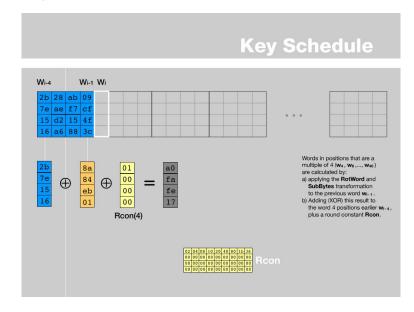


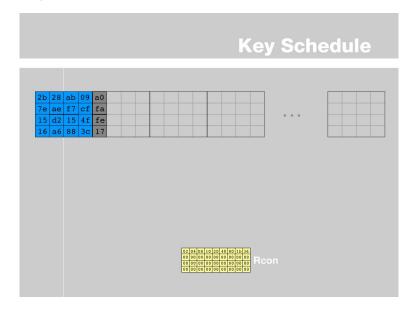


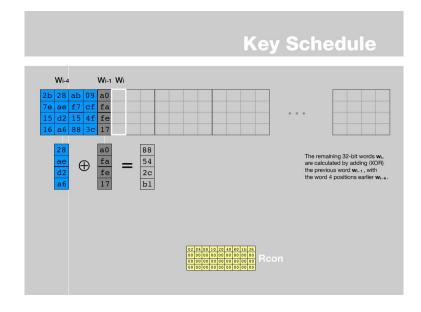


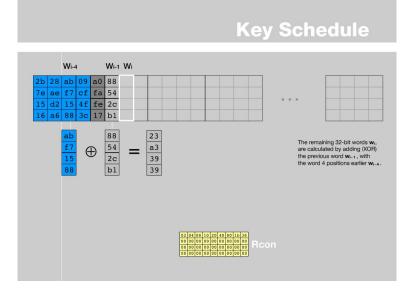


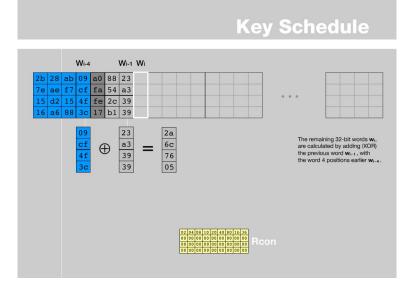


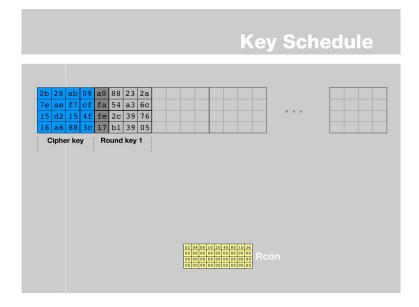


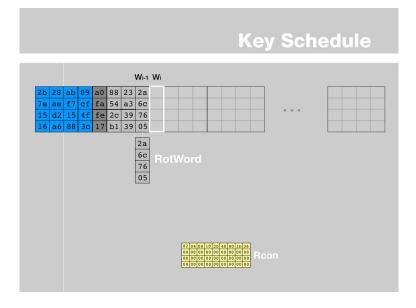


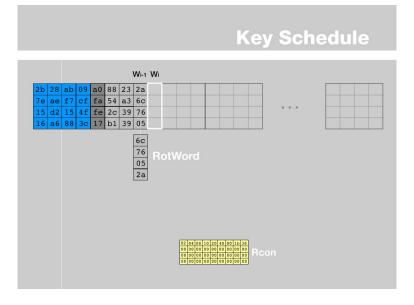


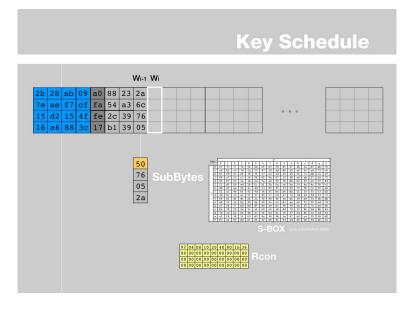


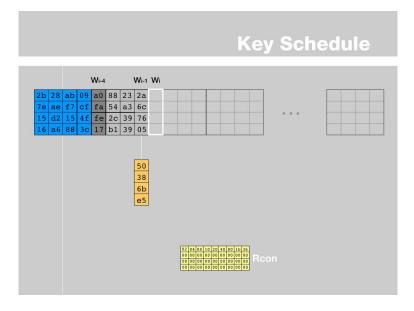


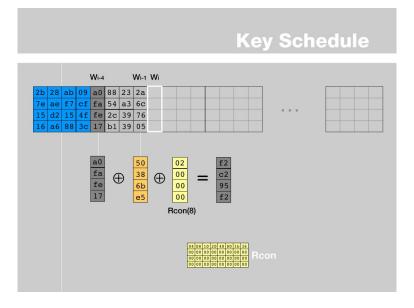


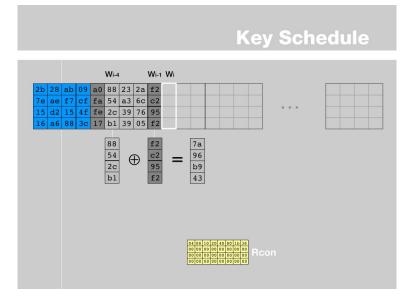


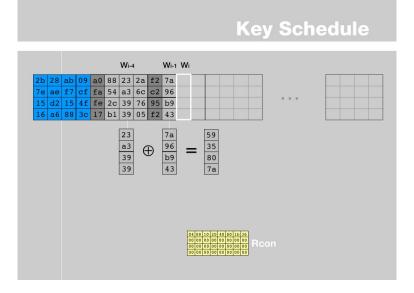








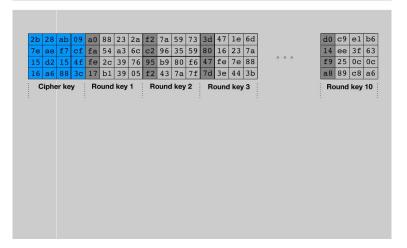






							Wi-4		,	Wi-1	Wi						
2b	28	ab	09	a0	88	23	2a	f2	7a	59							
7e	ae	f7	cf	fa	54	a3	6c	c2	96	35							
15	d2	15	4f	fe	2c	39	76	95	b9	80							
16	a 6	88	3с	17	b1	39	05	f2	43	7a							
							05		ļ	7a			7 f				
										0	4 08 10 0 00 00 0 00 00 0 00 00	20 40 00 00 00 00	80 1b 00 00 00 00 00 00	36 00 00 00	Rc		

Key Schedule



Corps finis quelques mots

Opération MixColumn :

$$\begin{pmatrix} 02 & 03 & 01 & 01 \\ 01 & 02 & 03 & 01 \\ 01 & 01 & 02 & 03 \\ 03 & 01 & 01 & 02 \end{pmatrix} \cdot \begin{pmatrix} d4 \\ bf \\ 5d \\ 30 \end{pmatrix} = \begin{pmatrix} 04 \\ 66 \\ 81 \\ e5 \end{pmatrix}$$

Opération de multiplication matrice vecteur.

$$r_{i,j} = \sum_{k=0}^{n-1} a_{i,k} b_k$$

où fait on cette opération?

▶ Dans GF(2^m) (Corps de Galois à 2^m éléments)