

EXTENDS *Integers*

CONSTANTS *R, G, B*

ASSUME $\wedge R \in 1 \dots 100$

$\wedge G \in 1 \dots 100$

$\wedge B \in 1 \dots 100$

$\wedge R + G + B > 0$

--fair algorithm *beans*{

variable *r* = *R*, *g* = *G*, *b* = *B*;

 { *S*: **while** (TRUE)

 { **either**

A1r: { **await** (*r* > 1); same color and red
 r := *r* - 2; *g* := *g* + 1; *b* := *b* + 1;
 } ;

or

A1g: { **await** (*g* > 1); same color and green
 g := *g* - 2; *r* := *r* + 1; *b* := *b* + 1;
 } ;

or

A1b: { **await** (*b* > 1); same color and blue
 b := *b* - 2; *r* := *r* + 1; *g* := *g* + 1;
 } ;

or

A2rg: { **await** (*r* > 0 \wedge *g* > 0); different color red and green
 r := *r* - 1; *g* := *g* - 1; *b* := *b* + 1;
 } ;

or

A2bg: { **await** (*b* > 0 \wedge *g* > 0); different color blue and green
 b := *b* - 1; *g* := *g* - 1; *r* := *r* + 1;
 } ;

or

A2rb: { **await** (*r* > 0 \wedge *b* > 0); different color red and blue
 r := *r* - 1; *b* := *b* - 1; *g* := *g* + 1;
 } ;

 } end while

 } end alg

} * end alg

BEGIN TRANSLATION

VARIABLES *r, g, b, pc*

vars \triangleq $\langle r, g, b, pc \rangle$

Init \triangleq Global variables

$\wedge r = R$

$\wedge g = G$

$$\begin{aligned} &\wedge b = B \\ &\wedge pc = \text{"S"} \end{aligned}$$

$$\begin{aligned} S \triangleq & \wedge pc = \text{"S"} \\ &\wedge \vee \wedge pc' = \text{"A1r"} \\ &\vee \wedge pc' = \text{"A1g"} \\ &\vee \wedge pc' = \text{"A1b"} \\ &\vee \wedge pc' = \text{"A2rg"} \\ &\vee \wedge pc' = \text{"A2bg"} \\ &\vee \wedge pc' = \text{"A2rb"} \\ &\wedge \text{UNCHANGED } \langle r, g, b \rangle \end{aligned}$$

$$\begin{aligned} A1r \triangleq & \wedge pc = \text{"A1r"} \\ &\wedge (r > 1) \\ &\wedge r' = r - 2 \\ &\wedge g' = g + 1 \\ &\wedge b' = b + 1 \\ &\wedge pc' = \text{"S"} \end{aligned}$$

$$\begin{aligned} A1g \triangleq & \wedge pc = \text{"A1g"} \\ &\wedge (g > 1) \\ &\wedge g' = g - 2 \\ &\wedge r' = r + 1 \\ &\wedge b' = b + 1 \\ &\wedge pc' = \text{"S"} \end{aligned}$$

$$\begin{aligned} A1b \triangleq & \wedge pc = \text{"A1b"} \\ &\wedge (b > 1) \\ &\wedge b' = b - 2 \\ &\wedge r' = r + 1 \\ &\wedge g' = g + 1 \\ &\wedge pc' = \text{"S"} \end{aligned}$$

$$\begin{aligned} A2rg \triangleq & \wedge pc = \text{"A2rg"} \\ &\wedge (r > 0 \wedge g > 0) \\ &\wedge r' = r - 1 \\ &\wedge g' = g - 1 \\ &\wedge b' = b + 1 \\ &\wedge pc' = \text{"S"} \end{aligned}$$

$$\begin{aligned} A2bg \triangleq & \wedge pc = \text{"A2bg"} \\ &\wedge (b > 0 \wedge g > 0) \\ &\wedge b' = b - 1 \\ &\wedge g' = g - 1 \\ &\wedge r' = r + 1 \\ &\wedge pc' = \text{"S"} \end{aligned}$$

$$\begin{aligned}
A2rb &\triangleq \wedge pc = \text{"A2rb"} \\
&\wedge (r > 0 \wedge b > 0) \\
&\wedge r' = r - 1 \\
&\wedge b' = b - 1 \\
&\wedge g' = g + 1 \\
&\wedge pc' = \text{"S"}
\end{aligned}$$

$$Next \triangleq S \vee A1r \vee A1g \vee A1b \vee A2rg \vee A2bg \vee A2rb$$

$$\begin{aligned}
Spec &\triangleq \wedge Init \wedge \Box [Next]_{vars} \\
&\wedge WF_{vars}(Next)
\end{aligned}$$

END TRANSLATION

$$\begin{aligned}
TypeOK &\triangleq \wedge r \in 0 \dots 200 \\
&\wedge g \in 0 \dots 200 \\
&\wedge b \in 0 \dots 200
\end{aligned}$$

$$Termination \triangleq \Diamond((r + g + b) < 2)$$

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\ * Modification History
\ * Last modified Tue Sep 30 22:18:41 EDT 2014 by Siddharth
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The invariant functions are given as:

1. $(r + g + b) > 0$
2. $r \geq 0$
3. $g \geq 0$
4. $b \geq 0$

The above are verified my model checking. The first invariant is however given as $(r + g + b) > 1$ in the model to avoid stuttering. When the invariant is violated we know that the program has reached a fixed point.

The metric/variant function can be given as the total number of beans:

$No.B = (r + g + b)$ is the metric function.

We observe from the program that the $No.B$ either stays the same or decreases for each action and hence selected as the metric funtion.

Fixed point:

To find the FP , we take the guards and negate it to get:

$$\begin{aligned}
&r < 2 \\
&\hat{g} < 2 \\
&\hat{b} < 2 \\
&\hat{(r < 1)} \vee (g < 1) \\
&\hat{(b < 1)} \vee (g < 1) \\
&\hat{(r < 1)} \vee (b < 1)
\end{aligned}$$

But we know that atleast one of them should be '1' from the invariant $(r + g + b) > 0$.

Therefore the fixed points states are given as $(\langle r \rangle, \langle g \rangle, \langle b \rangle)$:

1. $\langle 1,0,0 \rangle$
2. $\langle 0,1,0 \rangle$
3. $\langle 0,0,1 \rangle$

We also know that the program terminates when $(r + g + b) < 2$, or more specifically when $r + b + g = 1$. (*This can be verified using TLA + model checking*). This is also given by the metric function $No.B(r + g + b)$ the value of which either decreases or stays the same after every action.

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