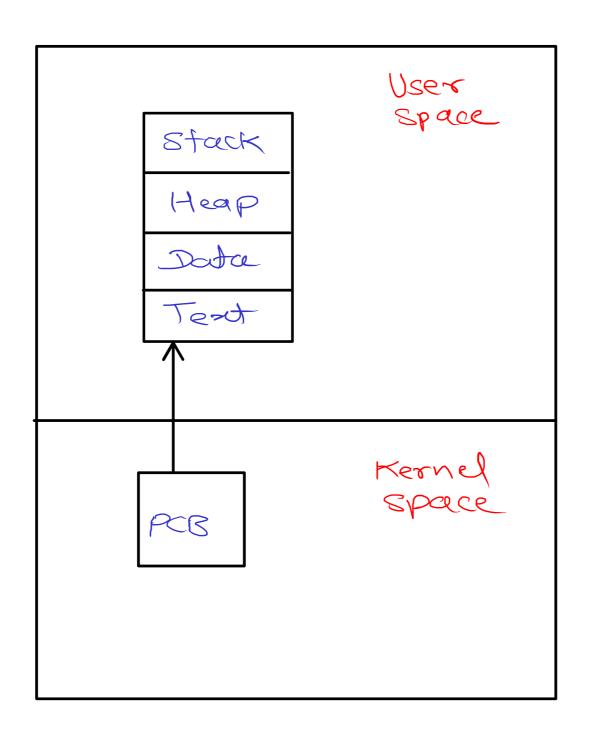
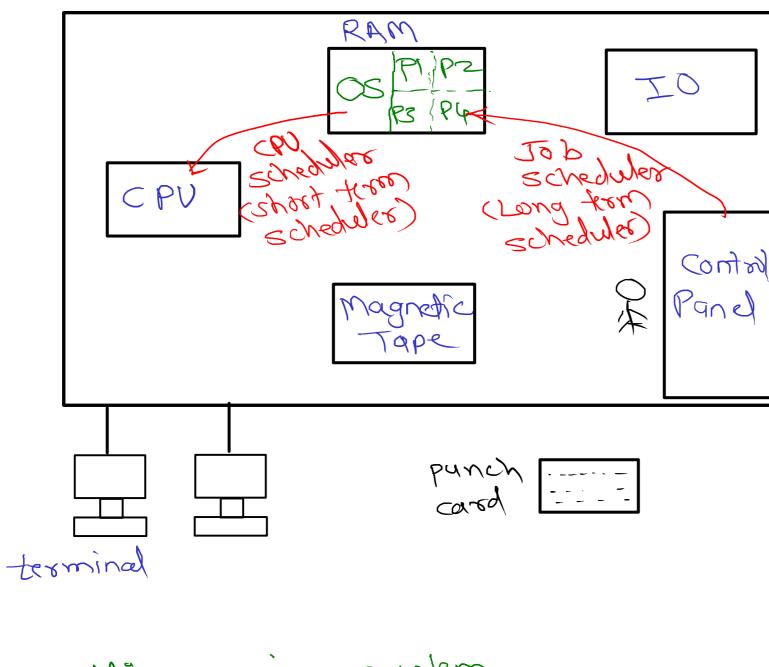
# **User space and Kernel Space**



#### **Types of Systems**



es multiprocessing system

1) Resident Monitor

2> Botch System

3) Multiprogramming System
- Degree of Multiprogramming

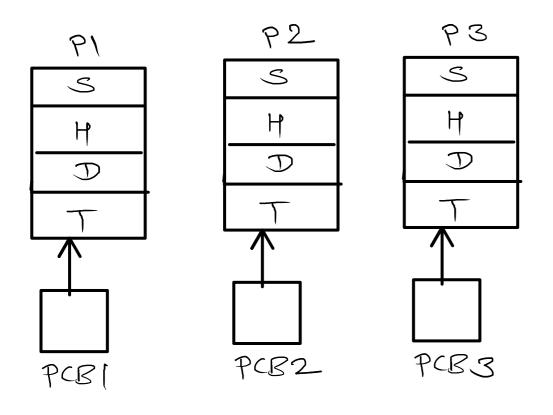
CPU time/burst - time spent on CPU IO time/burst - time spent on ID CPU burst > IO burst > CPU bound IO burst > CPU bound IO burst > CPU burst > IO bound

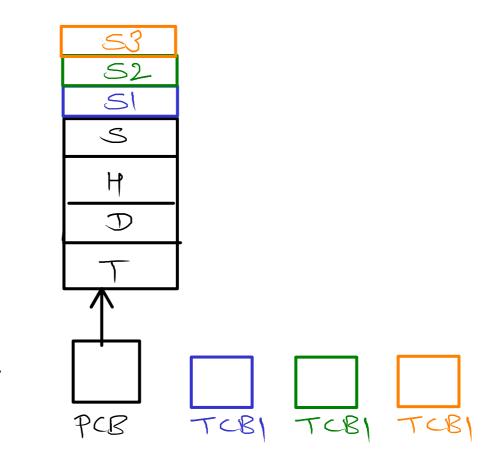
4) Time sharing system)
nultitasking system
Response time < 1 sec

2) Thread based Multitasking (multithreading)

> s) Multiuser system multiple users can control the system

#### **Process Vs Thread**

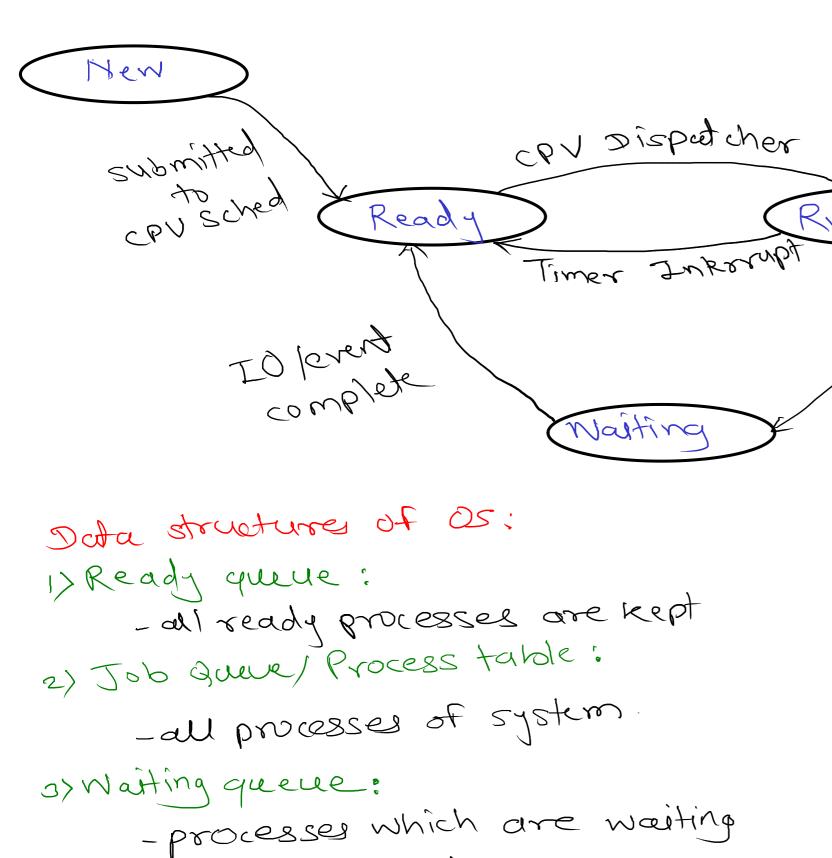




Thread - Lightweight process

- -for every process one thread is created
- process is just a container of resources.
- the threads in side that container are scheduled on CPU

# **Process Life Cycle**



- worting queues one multiple

for IO/some event.

( per device one)

Running / TO levent request CPV scheduling Contenias; 1) CPU utilization: (max) 2) Thoughput: (max) -amount of work done in unit time 3) Maiting time: (min) - total time spent by process into ready queue. 4) Response time: (min) -time form arraival into reddy queue appo First time getting e se ecutival. 5) Furn Arround time: (min) -total time spent into memory.

Terminated

## **Types of Scheduling**

1) Running > Terminated of Non-Preemptive Scheduling
2) Running > Neuting
3) Running > Ready
4) Waiting > Ready
5) Preemptive Scheduling
6) Waiting > Ready

# **CPU Scheduling Algorithms**

) FCFS

2) SJF

3> Prioriy

4) RR

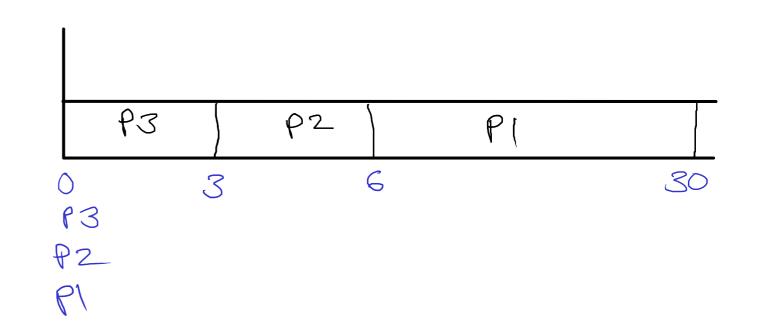
5> Fair Share

# FCFS (First Come First Serve) (Non-Preemptive scheduling)

	Process	Arrival	CPU Burst	WT	RT	TAT
1	P1	0	24	0	$\circ$	24
	P2	0	3	24	24	27
1	P3	0	3	27	27	30

	Process	Arrival	CPU Burst	WT	RT	TAT
	P3	0	3	0	$\bigcirc$	S
	P2	0	3	3	3	6
$\downarrow$	P1	0	24	6	6	30

P	\	P2		P3 [
Ŏ	2	L9	27	30
Pl				
P2 P3				
P3				



### **Convoy's Effect**

- due to arrival of longer process early into ready queue, all other processes has to wait for longer time
  - we do not have any control on sequenece of processes

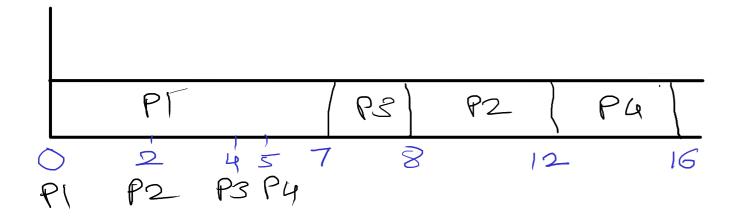
# SJF (Shortest Job First)

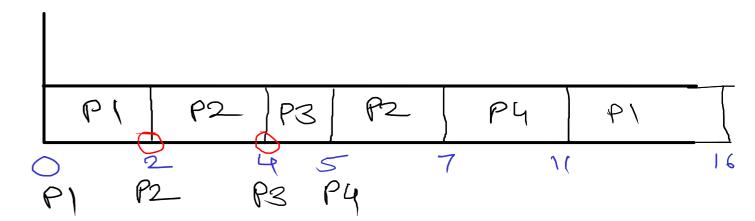
(Non Preemplive Scheduling)

(Preemptive Scheduling) (Shortest Remaining time First)

Process	Arrival	CPU Burst	WT	RT	TAT
P1	0	7	0	$\bigcirc$	7
P2	2	4	6	6	$\bigcirc$
Р3	4	1	3	3	4
P4	5	4	フ	7	\ \

				. 1—	$\bigcirc$ $\leftarrow$	T 01
Process	Arrival	CPU Burst		WT	₩	1 + 1
P1	0	7	5	9	$\bigcirc$	16
P2	2	4	2,0	·	$\bigcirc$	5
P3	4	1	D	0	$\mathcal{O}$	1
P4	5	4	4,0	2	2	Ç





#### **Starvation**

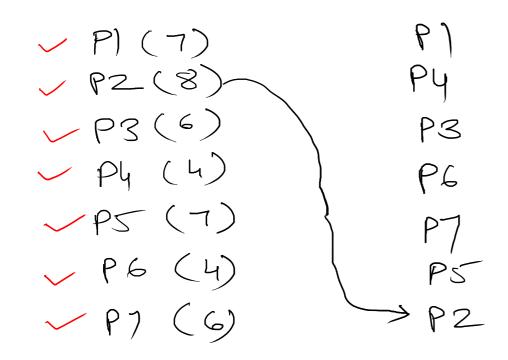
- due to longer CPU time process will get last chance to execute
  - there is no solution for starvation in SJF

#### **Priority**

# ( Non Preemptive Scheduling)

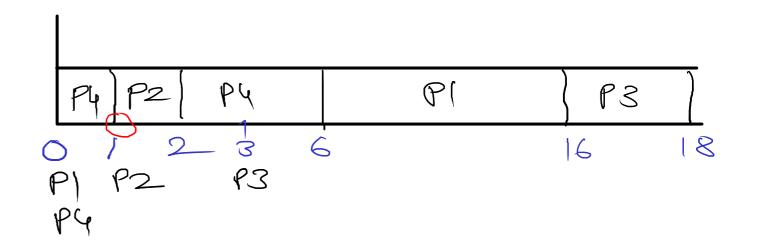
				,	$\wedge$ $-$	ナハイ
<b>Process</b>	Arrival	CPU Burst	Priority	NT	RT	( )
P1	0	10	3	\ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \	6	(6
P2	0	1	1(4)		Ô	<b>\</b>
P3	0	2	4 (L)	16	16	(8
P4	0	5	2	1	ſ	6

P2	PP	PI	P3	
0 1	E		16	8
P) P2		( (-0.0	dt's char	-+)
P3		( dan	01130160	
PY				



# (Preemptive Scheduling)

				-		<b>—</b>	
Process	Arrival	<b>CPU Burst</b>	Priority				TAT
P1	0	10	3	0,61	$\varsigma$	6	
P2	1	1	1	0	0	0	
Р3	3	2	4	2-10	13	13	
P4	0	5	2	40	\	$\bigcirc$	



#### **Starvation**

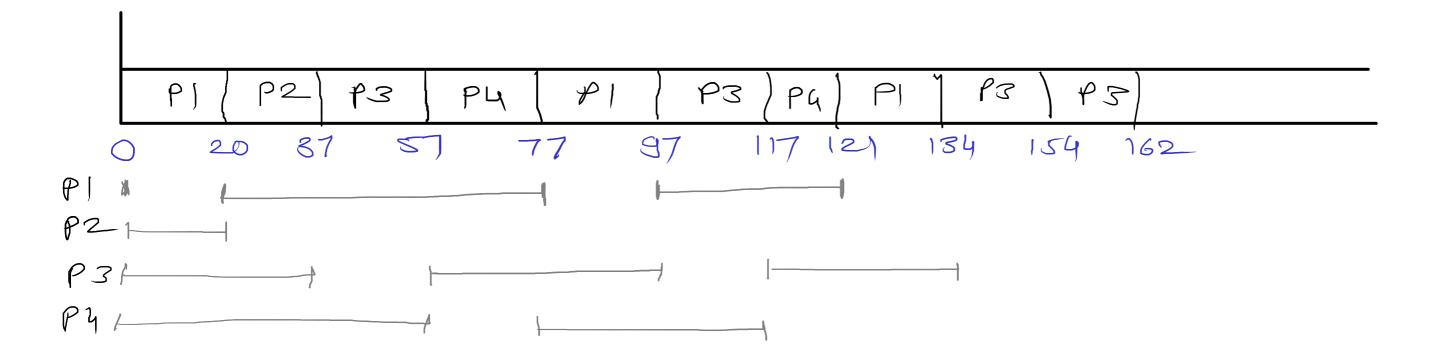
due to low priority, process will not get enough CPU time for execution

## Aging

increase the priority of process gradually, so that it will get scheduled

# RR (Round Robin) ( Pre emptive scheduling)

Process	CPU Burst		WT	RT	TAT	Time guantura	
P1		33,13,0	0+57+24	$\bigcirc$	134	- CPV time is	divided
P2	17	0	20	20	37	IND CHUM DUATI.	
Р3	68	48,28,8	37 +40+17	37	162		
P4	24	4,0	57+40	57	12)	T9 = 20	



First waiting time is Response time Tg=100 Lbehave like FCFS Tg=4 LCPU overhead will increase.

#### **Fair Share**

- CPU time is divided into time slices (epoch)
- some share of each epoch is given to the processes which are in ready queue.
- share is given to the process on the basis of their priority
- priority of every process is decided by its nice value
- nice values range ---> -20 to +19 (40 values)
  - \* -20 highest priority

\* +19 - lowest priority

Process	Nice Value
P1	10
P2	10
P3	10
P4	10

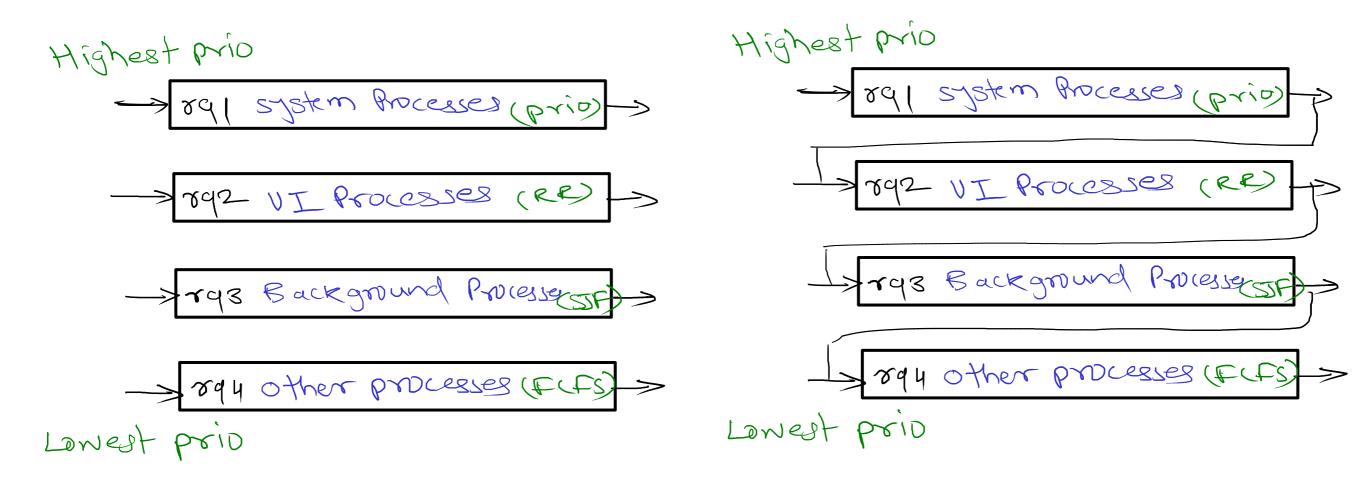
**Epoch - 100** 

Process	Nice Value
P1	5
P2	5
P3	10
P4	10

P1 P2 P3 P4 P1 P2	PS	<b>P</b> 4
-------------------	----	------------

## Multi Level Ready Queue

## Multi Level Feedback Ready Queue



#### **Linux Scheduling**

- There are two scheduling policies
  - 1. Non real time policy
    - i. SCHED OTHER
    - ii. SCHED BATCH
    - iii. SCHED IDLE
  - 2. Real time policy
    - i. SCHED\_FIFO
    - ii. SCHED\_RR