

Analizador Sintático - Gramática LL(1)

1. Gramática

$S \rightarrow \text{EXPR}$

$\text{EXPR} \rightarrow (\text{OPERAND OPERAND OPERATOR})$

$| (\text{if EXPR then EXPR else EXPR})$

$| (\text{for INT EXPR})$

$| (\text{INT RES})$

$| (\text{NUM MEM})$

$| (\text{MEM})$

$| \text{NUM}$

$\text{OPERAND} \rightarrow \text{EXPR} | \text{NUM} | \text{MEM}$

$\text{OPERATOR} \rightarrow + | - | * | / | \% | ^ | | | < | > | == | != | <= | >=$

$\text{NUM} \rightarrow \text{INT} | \text{REAL}$

$\text{INT} \rightarrow [0-9]^+$

$\text{REAL} \rightarrow [0-9]^+.[0-9]^+ | [0-9]^+e[+-]?[0-9]^+$

2. Conjuntos FIRST

$\text{FIRST}(S) = \{ (, \text{INT}, \text{REAL} \}$

$\text{FIRST}(\text{EXPR}) = \{ (, \text{INT}, \text{REAL} \}$

$\text{FIRST}(\text{OPERAND}) = \{ (, \text{INT}, \text{REAL}, \text{MEM} \}$

$\text{FIRST}(\text{OPERATOR}) = \{ +, -, *, /, \%, ^, |, <, >, ==, !=, <=, >= \}$

$\text{FIRST}(\text{NUM}) = \{ \text{INT}, \text{REAL} \}$

3. Conjuntos FOLLOW

$\text{FOLLOW}(S) = \{ \$ \}$

$\text{FOLLOW}(\text{EXPR}) = \{), \text{then}, \text{else}, \$, \text{OPERATOR} \}$

$\text{FOLLOW}(\text{OPERAND}) = \{ (, \text{INT}, \text{REAL}, \text{MEM}, \text{OPERATOR} \}$

$\text{FOLLOW}(\text{OPERATOR}) = \{) \}$

$\text{FOLLOW}(\text{NUM}) = \{), \text{MEM}, \text{RES}, \text{OPERATOR} \}$

4. Tabela de Derivação LL(1)

$M[S, ()] = \text{EXPR}$

$M[S, \text{INT}] = \text{EXPR}$

$M[S, \text{REAL}] = \text{EXPR}$

$M[\text{EXPR}, ()] = (\text{PARSE_PAREN_EXPR}$

$M[\text{EXPR}, \text{INT}] = \text{NUM}$

$M[\text{EXPR}, \text{REAL}] = \text{NUM}$

$M[\text{PARSE_PAREN_EXPR}, \text{if}] = \text{if EXPR then EXPR else EXPR)}$

$M[\text{PARSE_PAREN_EXPR}, \text{for}] = \text{for INT EXPR)}$

$M[\text{PARSE_PAREN_EXPR}, \text{INT}] = \text{CHECK_SECOND}$

$M[\text{PARSE_PAREN_EXPR}, \text{REAL}] = \text{CHECK_SECOND}$

$M[\text{PARSE_PAREN_EXPR}, \text{MEM}] = \text{CHECK_TYPE}$

$M[\text{PARSE_PAREN_EXPR}, \text{OPERAND}] = \text{OPERAND OPERAND OPERATOR)}$

$M[\text{CHECK_SECOND}, \text{RES}] = \text{INT RES)}$

$M[\text{CHECK_SECOND}, \text{MEM}] = \text{NUM MEM)}$

$M[\text{CHECK_SECOND}, \text{default}] = \text{OPERAND OPERAND OPERATOR)}$

$M[\text{CHECK_TYPE},] = \text{MEM)}$

$M[\text{CHECK_TYPE}, \text{default}] = \text{OPERAND OPERAND OPERATOR)}$

$M[\text{OPERAND}, ()] = \text{EXPR}$

$M[\text{OPERAND}, \text{INT}] = \text{NUM}$

$M[\text{OPERAND}, \text{REAL}] = \text{NUM}$

$M[\text{OPERAND}, \text{MEM}] = \text{MEM}$

$M[\text{NUM}, \text{INT}] = \text{INT}$

$M[\text{NUM}, \text{REAL}] = \text{REAL}$

5. Notas sobre a gramática

Esta gramática implementa a Notação Polonesa Reversa (RPN) onde:

1. As operações binárias seguem o formato: (OPERAND OPERAND OPERATOR)
2. Suporta expressões condicionais: (if CONDIÇÃO then EXPR_VERDADEIRO else EXPR_FALSO)
3. Suporta loops for: (for CONTADOR EXPRESSÃO)
4. Implementa comandos especiais:
 - (N RES) - Recupera o resultado de N linhas anteriores
 - (V MEM) - Armazena um valor V na memória
 - (MEM) - Recupera o valor armazenado na memória

A gramática está em conformidade com os requisitos da atividade avaliativa e suporta todas as operações necessárias:

- Operações aritméticas: +, -, *, /, %, ^, |
- Operações de comparação: <, >, ==, !=, <=, >=
- Estruturas de controle: if-then-else, for