Transactions

Dominic Duggan
Stevens Institute of Technology

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Transaction and Data Managers Transactions Data (and Lock) Managers read update read update Transactions are stateful

Atomicity

- Either a transaction happens completely, or it does not happen at all
- If a transaction happens, it appears to happen as a single indivisible action

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User 1 Output Insert into emails (recipient_id, body, unread_flag) values(2, 'Hello,' true) Output Database Insert into emails (recipient_id, body, unread_flag) values(2, 'Hello,' true) Output User 1 Output User 1

Consistency

- If the system has certain invariants, then if they held before a transaction, they hold after the completion of that transaction
 - Law of conservation of money
- That doesn't mean that the invariants need to be true throughout the transaction
- Application Property

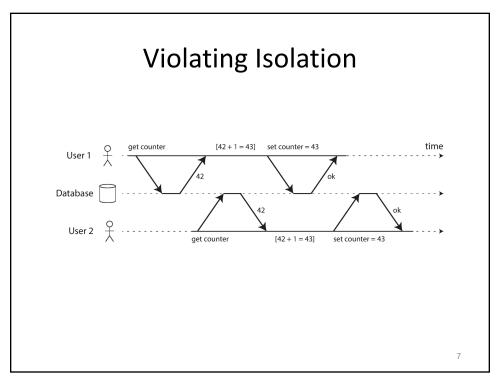
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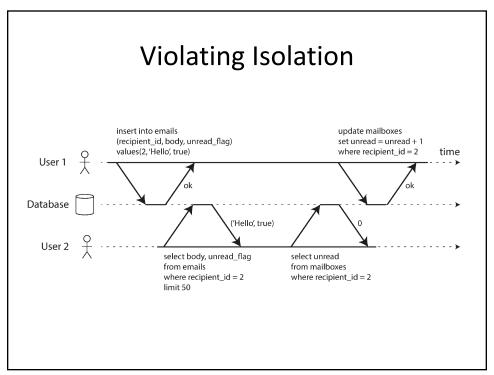
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Isolation

- Final result looks as though all transactions ran sequentially in some (system-dependent) order
 - They appear to have run serially, rather than in parallel
- Implemented by no database...

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Durability

- Once a transaction is successfully committed, regardless of what happens afterward, its results become permanent
- Specifically, no failure after the commit can undo the results or cause them to be lost
- At least written to (stable) log...

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IMPLEMENTATION

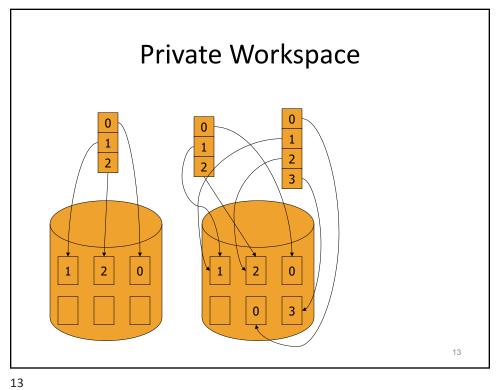
Two Implementation Methods

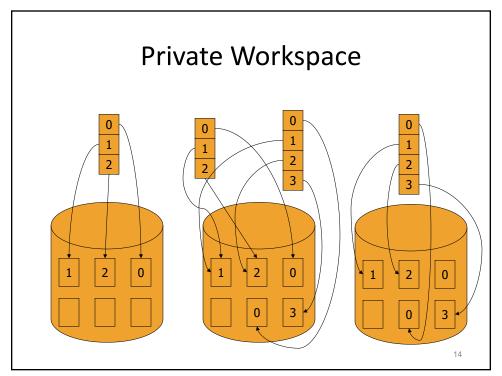
- Private Workspace
 - Shadow paging
- Writeahead Log
 - Log of updates, including original values

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Private Workspace





Writeahead Log

- DB modified in place
- Log record:
 - transaction identifier
 - file and block that are being changed
 - old and new values of the block
- DB only changed after log record has been successfully written
 - File update may happen some time after log written
 - File/database access slower than log append

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Writeahead Log Example

```
x = 0; (1) [x = 0/1]

y = 0;

BEGIN_TRANSACTION; (2) [x = 0/1]

(1) x = x + 1; [y = 0/2]

(2) y = y + 2;

(3) x = y * y; (3) [x = 0/1]

END_TRANSACTION; [y = 0/2]

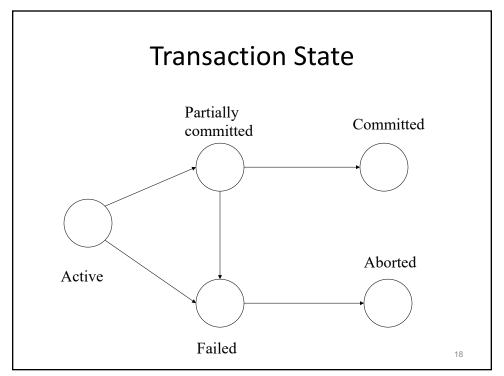
[x = 1/4]
```

Writeahead Log

- Success: commit record is written to the log
- Abort: roll back any changes

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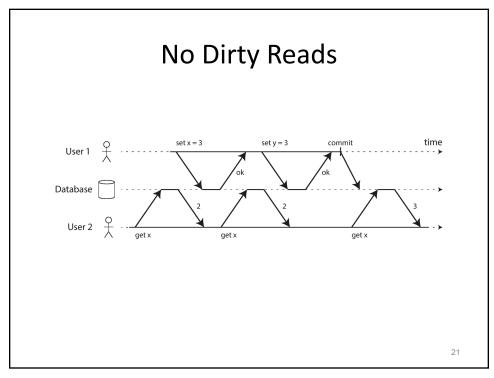


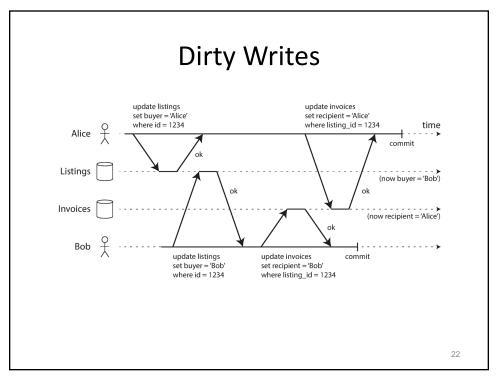


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Read Committed

- No dirty reads
 - When reading from the database, you will only see data that has been committed
- No dirty writes
 - When writing to the database, you will only overwrite data that has been committed





Implementing Read Committed

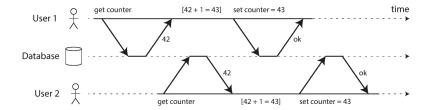
- Prevent dirty writes
 - Lock database record until commit
- Prevent dirty reads
 - Other transactions are given old version of data until commit

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Read Committed

• Does Read Committed prevent this?



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Read Skew

• Problem for e.g. backups, analytic queries

