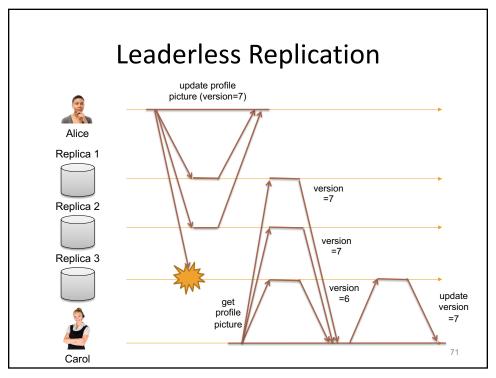


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# **Leaderless Replication**

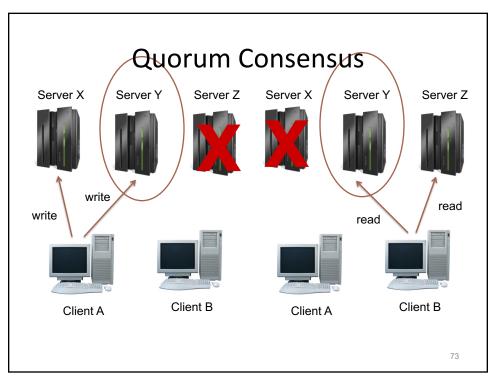
- No leader to enforce order of writes
- Amazon Dynamo
- How to prevent stale reads?

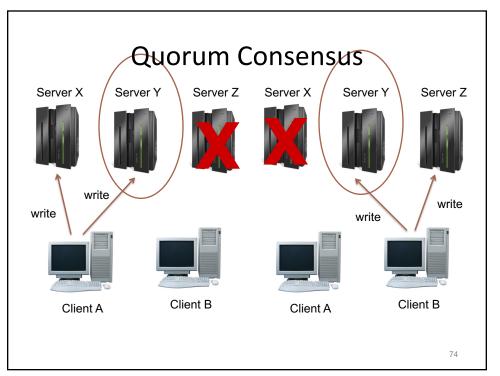


# Fixing Stale Data

- Read Repair
- Anti-Entropy

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### **Quorum Consensus**

- Each replicated object has an update and a read quorum
- Rules
  - A quorum read should "intersect" any prior quorum write at ≥ 1 processes
  - A quorum write should also intersect any other quorum write
- So, in a group of size N:
  - $-Q_r + Q_w > N$ , and
  - $-Q_w + Q_w > N$

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### Quorum example

- X is replicated at {a,b,c,d,e}
- Possible values?

```
-Q_w = 1, Q_r = 5 (violates Q_w + Q_w > 5)

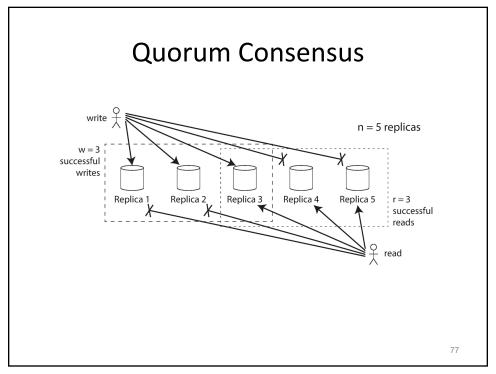
-Q_w = 2, Q_r = 4 (same issue)
```

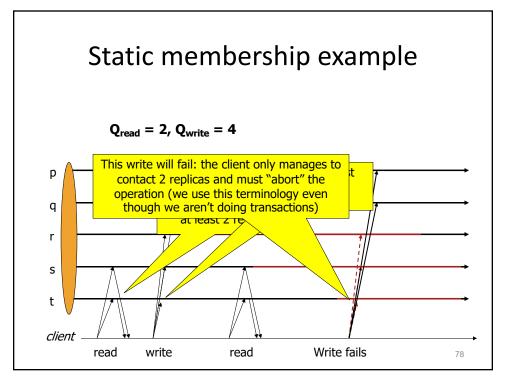
$$-Q_{w} = 3, Q_{r} = 3$$

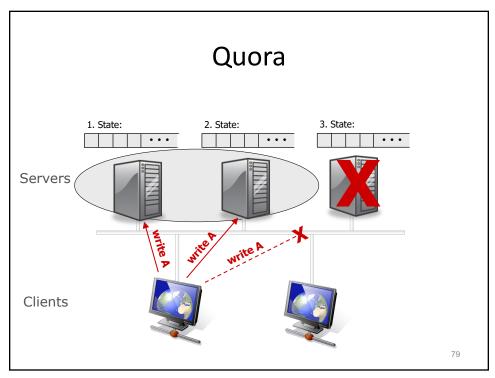
$$-Q_{w} = 4, Q_{r} = 2$$

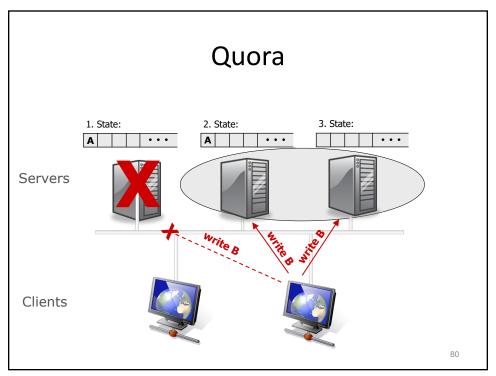
 $-Q_w = 5$ ,  $Q_r = 1$  (violates availability)

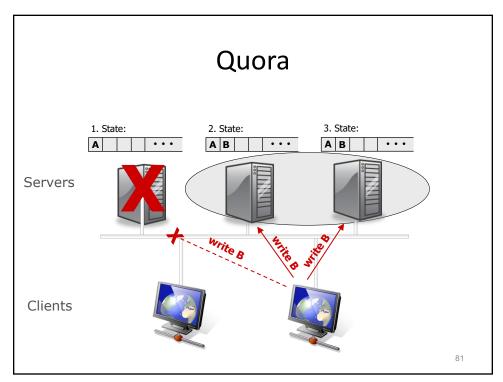
• Probably prefer Q<sub>w</sub>=4, Q<sub>r</sub>=2

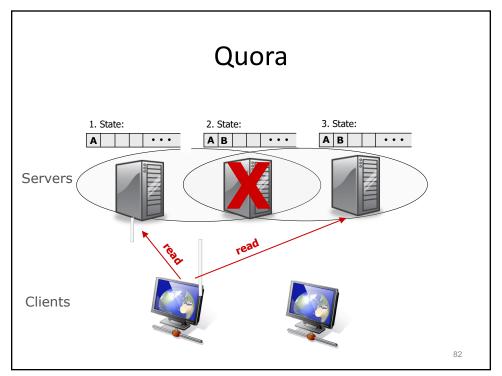








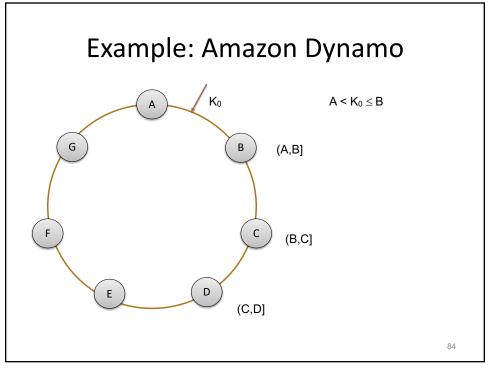


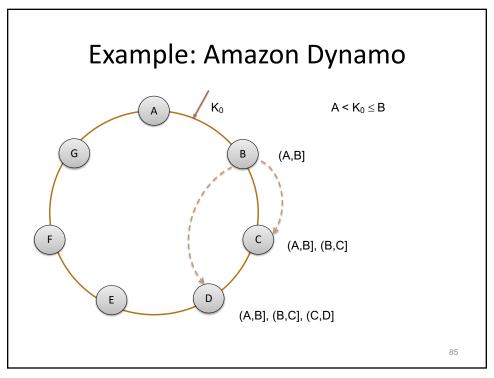


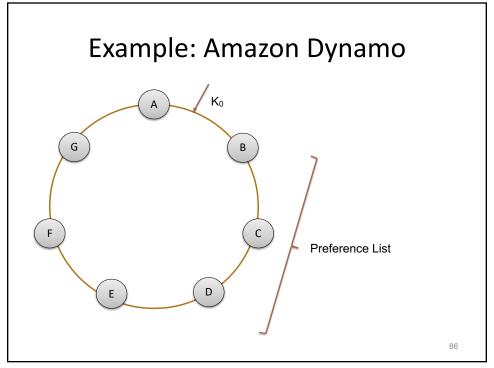
## **Leaderless Replication**

- Client contacts available replicas
- Read quorum, write quorum
- Read operations repair missed updates

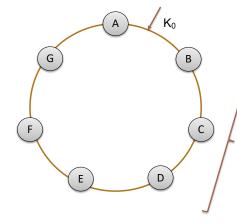
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### Example: Amazon Dynamo

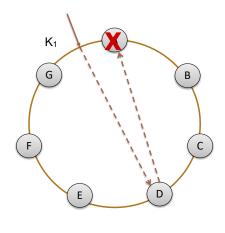


- Consistent hashing
  - "Virtual Nodes"
    - Nodes back up other nodes
    - Load balancing
- "Preference list"
  - e.g. B, C, D for key  $K_0$

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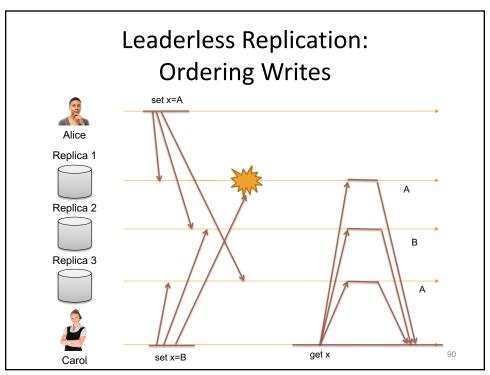
### Hinted handoff



- Assume # replicas for K<sub>1</sub>
   N = 3
- A down ⇒ send replica to e.g. D
- D is hinted that the replica belongs to A
- D will deliver to A when A is recovered
- "Always writeable"

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# Sloppy Quorum R/W read and write quora R + W > N Latency dictated by slowest of R (or W) replicas ...but which N? Pref List (N=3) ⇒ consistent Sloppy Quorum: any N (hinted handoff) R + W < N</li> better latency



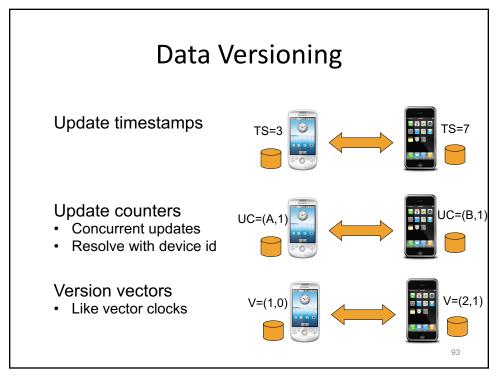
# Leaderless Replication: Ordering Writes

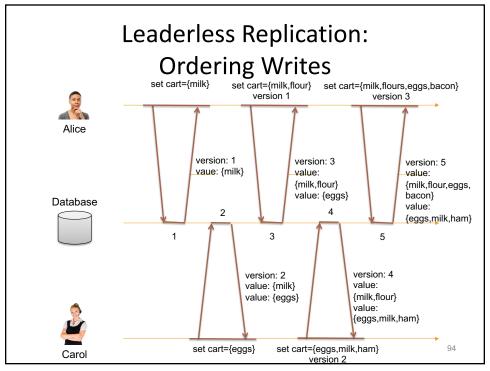
- Last Write Wins (LWW)
  - writes may be lost (not durable)
  - Cassandra: keep version history
  - Concurrent writes?
- Detecting conflicts
  - "Happened-before"
  - Version numbers
  - Update must read and merge current values (conflict resolution)

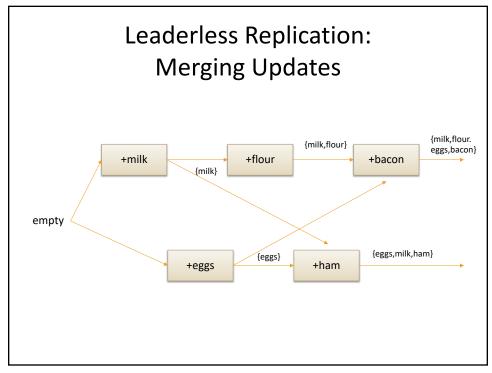
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### **Data Versioning**

- Write: return before update applied at all replicas
- Read: return many versions of the same object
- Challenge: object has distinct version subhistories
- Solution: vector clocks for reconciliation

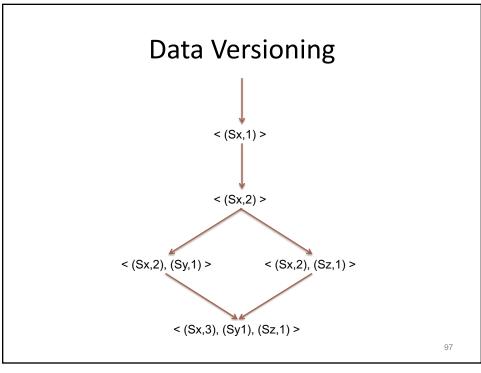






### **Merging Writes**

- Server stores version number with each record
- Client read: All values not overwritten
- Client write (after read): merge together all versions in last read
- Server write: overwrite all earlier versions



### **Limitations of Quorum Consistency**

- Sloppy Quorum (Dynamo)
  - Qw and Qr on different nodes from "home" nodes
- · Concurrent writes must be merged
- Concurrent write and read, read returns new or old value
- Failed write not rolled back, may be returned from read
- Recovering node may get stale value from another node
- · Problems with linearizability
- Measuring staleness: Ordering of writes?