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EXTENDS TLC, Naturals, Sequences, utils
CONSTANTS BcNodes, BftNodes, CrossLink2Nodes
CONSTANTS ByzBft, ByzCl
CONSTANTS Sigma, L
VARIABLES bc_chains, bft_chains, crosslink2_chains
INSTANCE definitions
Init \triangleq
     \land bc\_chains = [i \in 1 ... BcNodes \mapsto \langle BcGenesisBlock \rangle]
     \land bft\_chains = [i \in 1 .. BftNodes \mapsto \langle BftGenesisBlock \rangle]
     \land crosslink2\_chains = [i \in 1 .. CrossLink2Nodes \mapsto CrossLink2GenesisBlock]
HonestBc \triangleq
    \exists n \in 1 ... BcNodes :
        LET
             base \stackrel{\triangle}{=} bc\_chains[BestBcChainIdx]
             bft \triangleq bft\_chains[BestBftChainIdx]
             tip \stackrel{\triangle}{=} base[Len(base)].hash
             next \triangleq tip + 1in
        \land bc\_chains' = [bc\_chains \ EXCEPT \ ![n] = Append(base, [
             context\_bft \mapsto bft[Len(bft)].hash,
            hash \mapsto next])]
        \land UNCHANGED \langle bft\_chains, crosslink2\_chains \rangle
HonestBft \triangleq
    \exists n \in 1 ... BftNodes:
        LET
             base \stackrel{\Delta}{=} bft\_chains[BestBftChainIdx]
                    \stackrel{\triangle}{=} \stackrel{\circ}{bc\_chains[BestBcChainIdx]}
             tip \stackrel{\triangle}{=} base[Len(base)].hash
             next \triangleq tip + 1
             hdrs \stackrel{\Delta}{=} PruneLasts(bc, Sigma)IN
        \land bft\_chains' = [bft\_chains \ EXCEPT \ ![n] = Append(base, [
                 headers\_bc \mapsto hdrs,
                 hash \mapsto next])]
        \land UNCHANGED \langle bc\_chains, crosslink2\_chains \rangle
ByzantineBft \triangleq
    \exists n \in ByzBft:
        LET
             base \triangleq bft\_chains[BestBftChainIdx]
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- MODULE crosslink2

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\stackrel{\triangle}{=} bc\_chains[BestBcChainIdx]
                   \stackrel{\scriptscriptstyle \Delta}{=} base[Len(base)].hash
               Byzantine node can create an arbitrary faulty block within a range
             byz \stackrel{\Delta}{=} tip + (CHOOSE inc \in 2...10 : TRUE)
             hdrs \stackrel{\triangle}{=} PruneLasts(bc, Sigma)IN
         \land bft\_chains' = [bft\_chains \ EXCEPT \ ![n] = Append(base, [
             headers\_bc \mapsto hdrs,
             hash
                            \mapsto byz])]
         \land UNCHANGED \langle bc\_chains, crosslink2\_chains \rangle
HonestCrosslink \stackrel{\triangle}{=}
     \exists n \in 1 .. CrossLink2Nodes :
        LET
             fin \triangleq PruneFirsts(bc\_chains[BestBcChainIdx], Sigma)
             ba \triangleq LocalBa(fin, bc\_chains[BestBcChainIdx])
         \land crosslink2\_chains' = [crosslink2\_chains \ EXCEPT \ ![n] = [
             fin \mapsto fin,
             ba \mapsto ba
         \land UNCHANGED \langle bc\_chains, bft\_chains \rangle
     \vee UNCHANGED \langle bc\_chains, bft\_chains, crosslink2\_chains \rangle
Next \triangleq
     \vee HonestBc
     \lor HonestBft
     \lor HonestCrosslink
     \lor ByzantineBft
Spec \stackrel{\Delta}{=} Init \wedge \Box [Next]_{\langle bc\_chains, \, bft\_chains, \, crosslink2\_chains \rangle}
Type checking
BcChainsTypeCheck \triangleq bc\_chains \in Seq(Seq([context\_bft : Nat, hash : Nat]))
BftChainsTypeCheck \triangleq bft\_chains \in
Seq(Seq([headers\_bc:Seq([context\_bft:Nat, hash:Nat]), hash:Nat])) \\ CrossLink2ChainsTypeCheck \triangleq crosslink2\_chains \in
     Seq([fin:Seq([context\_bft:Nat,hash:Nat]),ba:Seq([context\_bft:Nat,hash:Nat])])
Assumptions
Assume BftThresholdOK
Lemma: Linear Prefix
If A \leq_{\star} C and B \leq_{\star} C then A \underset{\star}{\underline{*}} B.
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BcLinearPrefix \triangleq
     \forall a, b, c \in 1 ... BcNodes:
         LET A \stackrel{\triangle}{=} bc\_chains[a]
                B \triangleq bc\_chains[b]
                C \stackrel{\Delta}{=} bc\_chains[c]
               IsPrefix(A, C) \wedge IsPrefix(B, C) \Rightarrow
                  IsPrefix(A, B) \vee IsPrefix(B, A)
BftLinearPrefix \triangleq
     \forall a, b, c \in 1 \dots \textit{BftNodes} :
         LET A \stackrel{\triangle}{=} bft\_chains[a]
                B \stackrel{=}{=} bft\_chains[b]
                C \stackrel{\Delta}{=} bft\_chains[c]
         IN IsPrefix(A, C) \wedge IsPrefix(B, C) \Rightarrow
                  IsPrefix(A, B) \vee IsPrefix(B, A)
Definition: Agreement on a view
```

An execution of Π has Agreement on the view $V: Node \times Time \rightarrow \star chain$ iff for all times t, uand all Π nodes i, j (potentially the same) such that i is honest at time t and j is honest at time u, we have $V_i^t \stackrel{\star}{\underline{\times}} V_i^u$.

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BcViewAgreement \triangleq
    \forall i, j \in 1 \dots BcNodes:
        \vee IsPrefix(bc\_chains[i], bc\_chains[j])
        \vee IsPrefix(bc\_chains[j], bc\_chains[i])
BftViewAgreement \triangleq
    \forall i, j \in HonestBftNodes:
        \vee IsPrefix(bft\_chains[i], bft\_chains[j])
        \vee IsPrefix(bft\_chains[j], bft\_chains[i])
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Definition: Final agreement

An execution of $\Pi_{\star bft}$ has Final Agreement iff for all bftvalid blocks C in honest view at time t and C' in honest view at time t', we have $bftlastfinal(C)

<math>
\frac{\star}{b} ft \star bftlastfinal(C')$

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BftFinalAgreement \triangleq
   \forall i, j \in HonestBftNodes:
       \vee IsPrefix(BftLastFinal(i), BftLastFinal(j))
       \vee IsPrefix(BftLastFinal(j), BftLastFinal(i))
```

Definition: Prefix Consistency

An execution of $\Pi_{\star bc}$ has Prefix Consistency at confirmation depth σ , iff for all times $t \leq u$ and all nodes i, j (potentially the same) such that i is honest at time t and j is honest at time u, we have that $ch_i^t \upharpoonright_{\star bc}^{\sigma} \leq_{\star bc} ch_j^u$

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BcPrefixConsistency \triangleq
    \forall i, j \in 1 \dots BcNodes:
        Len(bc\_chains[i]) \leq Len(bc\_chains[j]) \Rightarrow
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IsPrefix(PruneFirsts(bc\_chains[i], Sigma), bc\_chains[j])
```

Definition: Prefix Agreement

An execution of $\Pi_{\star bc}$ has Prefix Agreement at confirmation depth σ iff it has Agreement on the view $(i,t) \mapsto ch_i^{t}|_{\star bc}^{\sigma}$.

 $BcPrefixAgreement \triangleq$

 $\forall i \in 1 \dots BcNodes$:

 $IsPrefix(PruneFirsts(bc_chains[i], Sigma), bc_chains[i])$

Definition: *-linear

A function $S: I \to \star block$ is *-linear iff for every $t, u \in I$ where $t \leq u$ we have $S(t) \leq_{\star} S(u)$

 $BcLinear(T, U) \triangleq IsPrefix(T, U)$

Definition: Local finalization linearity

Node i has Local finalization linearity up to time t iff the time series of $\star bc$ -blocks $fin_i^{r \le t}$ is $\star bc$ -linear.

 $LocalFinalizationLinearity \triangleq \Box[$

 $\forall i \in 1 .. CrossLink2Nodes:$

 $BcLinear(crosslink2_chains[i].fin, crosslink2_chains'[i].fin)]_{crosslink2_chains}$

Lemma: Local fin-depth

In any execution of Crosslink 2, for any node i that is honest at time t, there exists a time $r \leq t$ such that $fin_i \leq ch_i^r \lceil_{\star}^{\sigma}b_c$

 $LocalFinDepth \triangleq$

 $\forall i \in 1 ... CrossLink2Nodes :$

IsPrefix(crosslink2_chains[i].fin, bc_chains[BestBcChainIdx])

Definition: Assured Finality

An execution of Crosslink 2 has Assured Finality iff for all times t, u and all nodes i, j (potentially the same) such that i is honest at time t and j is honest at time u, we have $fin_i^t \stackrel{\star}{\succeq}_{bc} fin_i^u$.

 $AssuredFinality \triangleq$

 $\forall i, j \in 1 ... CrossLink2Nodes :$

 $\lor \mathit{IsPrefix}(\mathit{crosslink2_chains}[\mathit{i}].\mathit{fin}, \, \mathit{crosslink2_chains}[\mathit{j}].\mathit{fin})$

 $\vee IsPrefix(crosslink2_chains[j].fin, crosslink2_chains[i].fin)$

Theorem: Ledger prefix property

For any node i that is honest at time t, and any confirmation depth μ , $\operatorname{fin}_i^t \leq (\mathsf{ba}_\mu)_i^t$

 $LedgerPrefixProperty \triangleq$

 $\forall i \in 1 .. CrossLink2Nodes :$

 $IsPrefix(crosslink2_chains[i].fin, crosslink2_chains[i].ba)$

Lemma: Local ba-depth

In any execution of Crosslink 2, for any confirmation depth $\mu \leq \sigma$ and any node i that is honest at time t, there exists a time $r \leq t$ such that $(\mathsf{ba}_\mu)^i_i \preceq_\mathsf{bc} \mathsf{ch}^i_i \lceil^\mu_\mathsf{bc}$.