
MODULE *crosslink2*

EXTENDS *TLC*, *Naturals*, *Sequences*, *utils*

CONSTANTS *BcNodes*, *BftNodes*, *CrossLink2Nodes*

CONSTANTS *Sigma*, *L*

VARIABLES *bc_chains*, *bft_chains*, *crosslink2_chains*

INSTANCE *definitions*

Init \triangleq

$\wedge bc_chains = [i \in 1 \dots BcNodes \mapsto \langle BcGenesisBlock \rangle]$
 $\wedge bft_chains = [i \in 1 \dots BftNodes \mapsto \langle BftGenesisBlock \rangle]$
 $\wedge crosslink2_chains = [i \in 1 \dots CrossLink2Nodes \mapsto CrossLink2GenesisBlock]$

Next \triangleq

$\vee \exists n \in 1 \dots BcNodes :$
 $\wedge bc_chains' = [bc_chains \text{ EXCEPT } ![n] = Append($
 $\quad bc_chains[ChooseBestBcChain], [$
 $\quad \quad context_bft \mapsto ChooseContextBft,$
 $\quad \quad hash \mapsto ChooseBestBcTip + 1])]$
 $\wedge \text{UNCHANGED } \langle bft_chains, crosslink2_chains \rangle$
 $\vee \exists m \in 1 \dots BftNodes :$
 $\wedge bft_chains' = [bft_chains \text{ EXCEPT } ![m] = Append($
 $\quad bft_chains[ChooseBestBftChain], [$
 $\quad \quad headers_bc \mapsto PruneLasts(ChooseBcView, Sigma),$
 $\quad \quad hash \mapsto ChooseBestBftTip + 1)])]$
 $\wedge \text{UNCHANGED } \langle bc_chains, crosslink2_chains \rangle$
 $\vee \exists c \in 1 \dots CrossLink2Nodes :$
 $\wedge crosslink2_chains' = [crosslink2_chains \text{ EXCEPT } ![c] = [$
 $\quad fin \mapsto PruneFirsts(bc_chains[ChooseBestBcChain], Sigma)]]$
 $\wedge \text{UNCHANGED } \langle bc_chains, bft_chains \rangle$
 $\vee \text{UNCHANGED } \langle bc_chains, bft_chains, crosslink2_chains \rangle$

Spec $\triangleq Init \wedge \Box [Next]_{\langle bc_chains, bft_chains, crosslink2_chains \rangle}$

Type checking

BcChainsTypeCheck $\triangleq bc_chains \in Seq(Seq([context_bft : Nat, hash : Nat]))$

BftChainsTypeCheck $\triangleq bft_chains \in$

$Seq(Seq([headers_bc : Seq([context_bft : Nat, hash : Nat]), hash : Nat]))$

CrossLink2ChainsTypeCheck $\triangleq crosslink2_chains \in$

$Seq([fin : Seq([context_bft : Nat, hash : Nat])])$

Lemma: Linear Prefix

If $A \preceq_{\star} C$ and $B \preceq_{\star} C$ then $A \star_{\star} B$.

$BcLinearPrefix \triangleq$
 $\forall a, b, c \in 1 \dots BcNodes :$
 LET $A \triangleq bc_chains[a]$
 $B \triangleq bc_chains[b]$
 $C \triangleq bc_chains[c]$
 IN $IsPrefix(A, C) \wedge IsPrefix(B, C) \Rightarrow$
 $IsPrefix(A, B) \vee IsPrefix(B, A)$

$BftLinearPrefix \triangleq$
 $\forall a, b, c \in 1 \dots BftNodes :$
 LET $A \triangleq bft_chains[a]$
 $B \triangleq bft_chains[b]$
 $C \triangleq bft_chains[c]$
 IN $IsPrefix(A, C) \wedge IsPrefix(B, C) \Rightarrow$
 $IsPrefix(A, B) \vee IsPrefix(B, A)$

Definition: Agreement on a view

An execution of Π has Agreement on the view $V : Node \times Time \rightarrow \star chain$ iff for all times t, u and all Π nodes i, j (potentially the same) such that i is honest at time t and j is honest at time u , we have $V_i^t \star_{\star} V_j^u$.

$BcViewAgreement \triangleq$
 $\forall i, j \in 1 \dots BcNodes :$
 $\vee IsPrefix(bc_chains[i], bc_chains[j])$
 $\vee IsPrefix(bc_chains[j], bc_chains[i])$

$BftViewAgreement \triangleq$
 $\forall i, j \in 1 \dots BftNodes :$
 $\vee IsPrefix(bft_chains[i], bft_chains[j])$
 $\vee IsPrefix(bft_chains[j], bft_chains[i])$

Definition: Final agreement

An execution of $\Pi_{\star bft}$ has Final Agreement iff for all *bftvalid* blocks C in honest view at time t and C' in honest view at time t' , we have $bftlastfinal(C) \star_{bft} \star bftlastfinal(C')$.

$BftFinalAgreement \triangleq$
 $\forall i, j \in 1 \dots BftNodes :$
 $\vee IsPrefix(BftLastFinal(i), BftLastFinal(j))$
 $\vee IsPrefix(BftLastFinal(j), BftLastFinal(i))$

Definition: Prefix Consistency

An execution of $\Pi_{\star bc}$ has Prefix Consistency at confirmation depth σ , iff for all times $t \leq u$ and all nodes i, j (potentially the same) such that i is honest at time t and j is honest at time u , we have that $ch_i^t \upharpoonright_{\star bc}^{\sigma} \preceq_{\star bc} ch_j^u$.

$$\begin{aligned}
BcPrefixConsistency &\triangleq \\
&\forall i, j \in 1 \dots BcNodes : \\
&\quad Len(bc_chains[i]) \leq Len(bc_chains[j]) \Rightarrow \\
&\quad IsPrefix(PruneFirsts(bc_chains[i], Sigma), bc_chains[j])
\end{aligned}$$

Definition: Prefix Agreement

An execution of $\Pi_{\star bc}$ has Prefix Agreement at confirmation depth σ iff it has Agreement on the view $(i, t) \mapsto ch_i^t \upharpoonright_{\star bc}^\sigma$.

$$\begin{aligned}
BcPrefixAgreement &\triangleq \\
&\forall i \in 1 \dots BcNodes : \\
&\quad IsPrefix(PruneFirsts(bc_chains[i], Sigma), bc_chains[i])
\end{aligned}$$

Definition: *-linear

A function $S : I \rightarrow \star block$ is *-linear iff for every $t, u \in I$ where $t \leq u$ we have $S(t) \preceq_\star S(u)$

$$BcLinear(T, U) \triangleq IsPrefix(T, U)$$

Definition: Local finalization linearity

Node i has Local finalization linearity up to time t iff the time series of $\star bc$ -blocks $fin_i^{r \leq t}$ is $\star bc$ -linear.

$$\begin{aligned}
LocalFinalizationLinearity &\triangleq \square[\\
&\forall i \in 1 \dots CrossLink2Nodes : \\
&\quad BcLinear(crosslink2_chains[i].fin, crosslink2_chains'[i].fin)]_{crosslink2_chains}
\end{aligned}$$

Lemma: Local fin-depth

In any execution of Crosslink 2, for any node i that is honest at time t , there exists a time $r \leq t$ such that $fin_i \preceq ch_i^r \upharpoonright_{\star bc}^\sigma$

$$\begin{aligned}
LocalFinDepth &\triangleq \\
&\forall i \in 1 \dots CrossLink2Nodes : \\
&\quad IsPrefix(crosslink2_chains[i].fin, bc_chains[ChooseBestBcChain])
\end{aligned}$$

Definition: Assured Finality

An execution of Crosslink 2 has Assured Finality iff for all times t, u and all nodes i, j (potentially the same) such that i is honest at time t and j is honest at time u , we have $fin_i^t \not\prec_{bc} fin_j^u$.

$$\begin{aligned}
AssuredFinality &\triangleq \\
&\forall i, j \in 1 \dots CrossLink2Nodes : \\
&\quad \vee IsPrefix(crosslink2_chains[i].fin, crosslink2_chains[j].fin) \\
&\quad \vee IsPrefix(crosslink2_chains[j].fin, crosslink2_chains[i].fin)
\end{aligned}$$