



<https://hao-ai-lab.github.io/dsc291-s24/>

DSC 291: ML Systems

Spring 2024

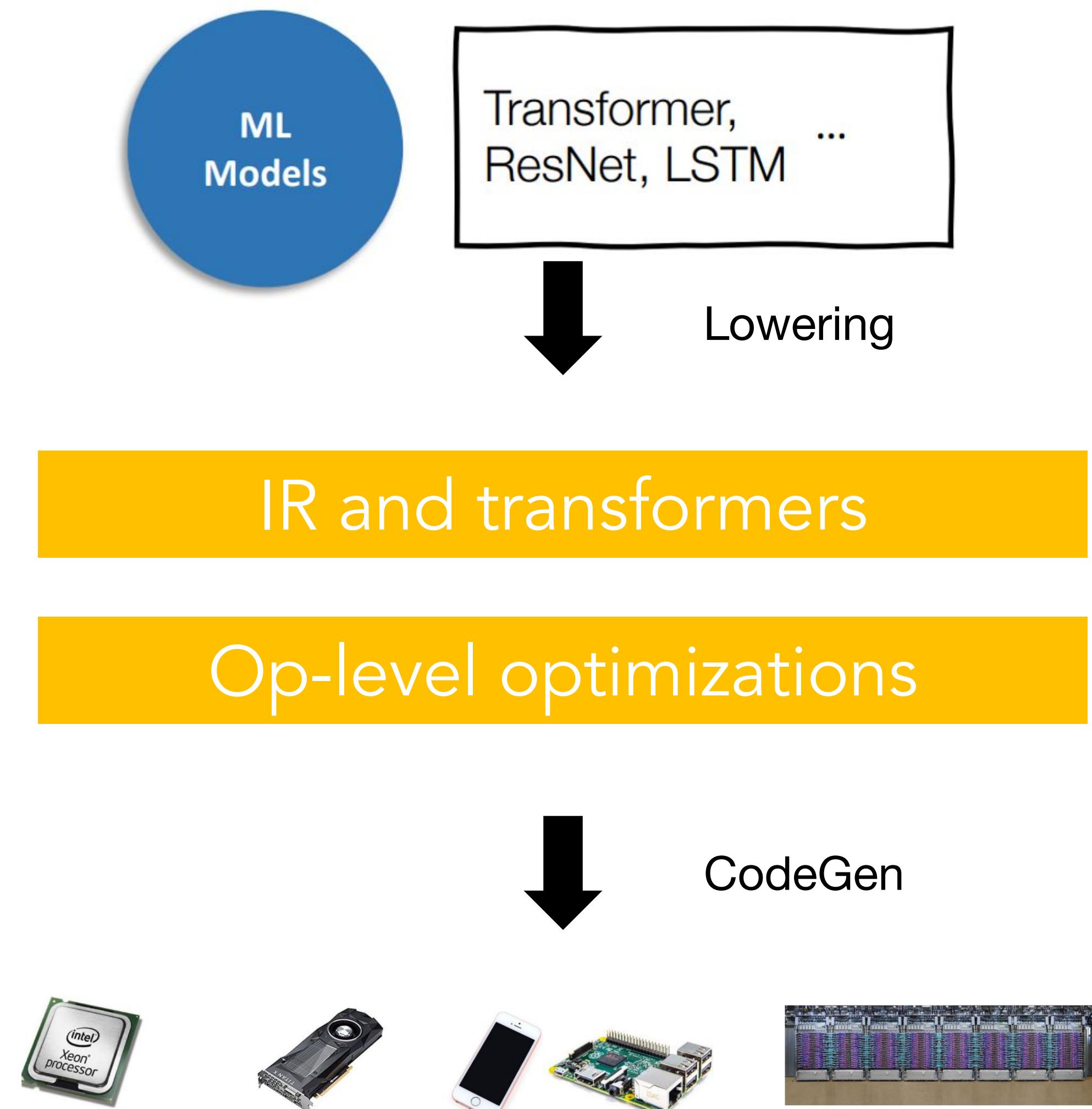
LLMs

Parallelization

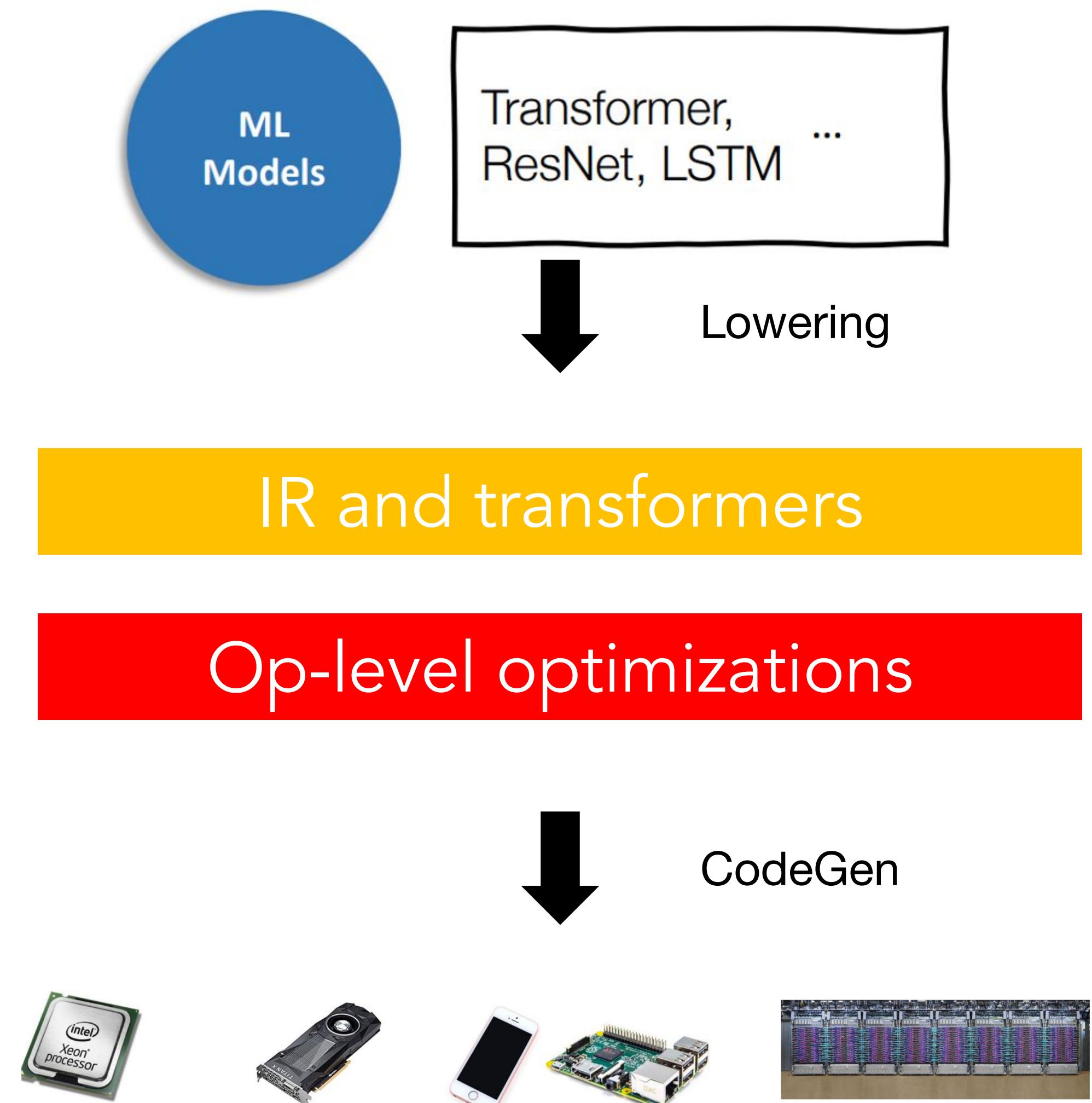
Single-device Optimization

Basics

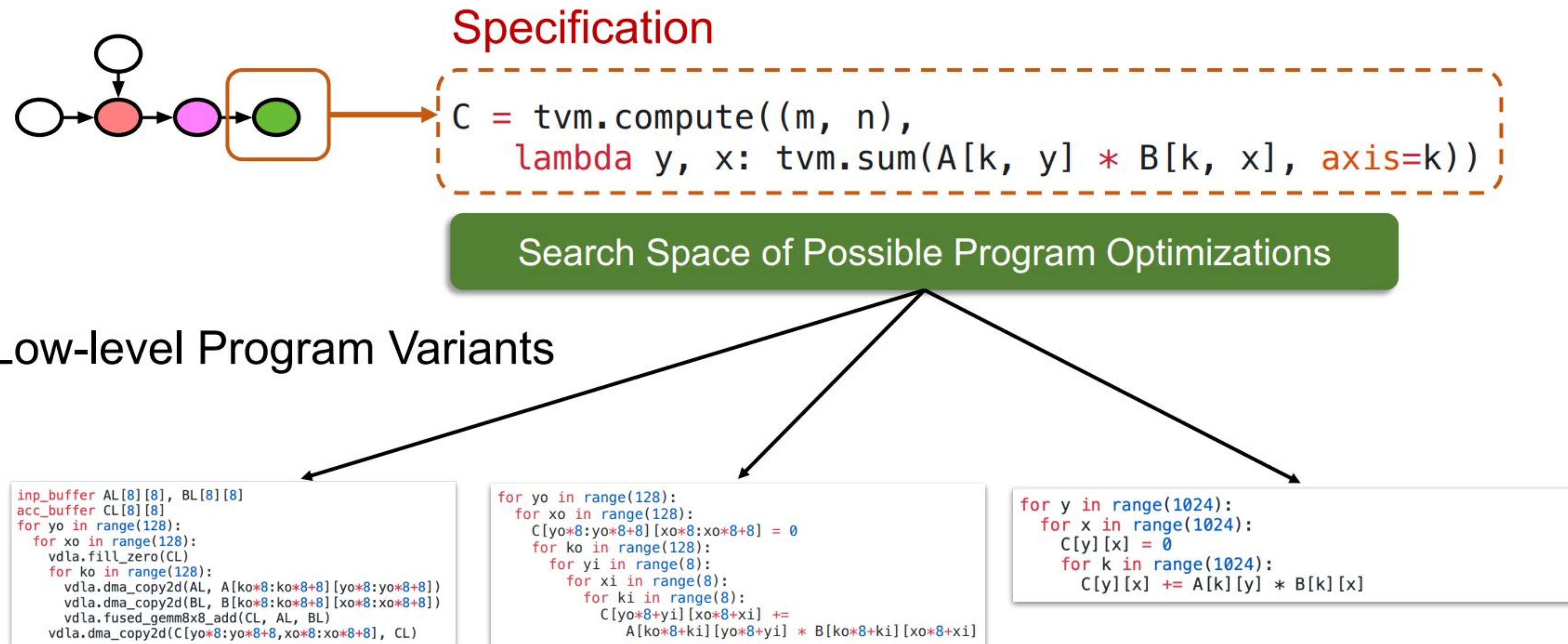
Compilation Process



Compilation Process



Lower-level code optimization



Low-level Loop Representation

```
@dot-add(x, w, b)
  for i, j in grid(16, 16):
    Y[i, j] = 0
    for i, j, k in grid(16, 16, 16):
      Y[i, j] += x[i, k] * w[k, j]
    for i, j in grid(16, 16):
      Z[i, j] = Y[i, j] + b[j]
```

Multi-dimensional buffer

Loop nests

Array computation

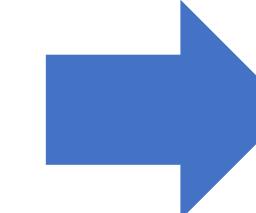
Transforming Loops: Loop Splitting

Code

```
for x in range(128):
    C[x] = A[x] + B[x]
```

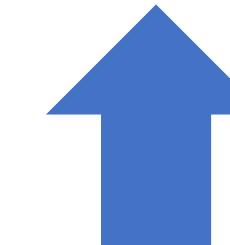


```
for xo in range(32):
    for xi in range(4):
        C[xo * 4 + xi]
        = A[xo * 4 + xi] + B[xo * 4 + xi]
```



```
def gpu_kernel():
    C[threadIdx.x * 4 + blockIdx.x] = . . .

for xi in range(4):
    for xo in range(32):
        C[xo * 4 + xi]
        = A[xo * 4 + xi] + B[xo * 4 + xi]
```



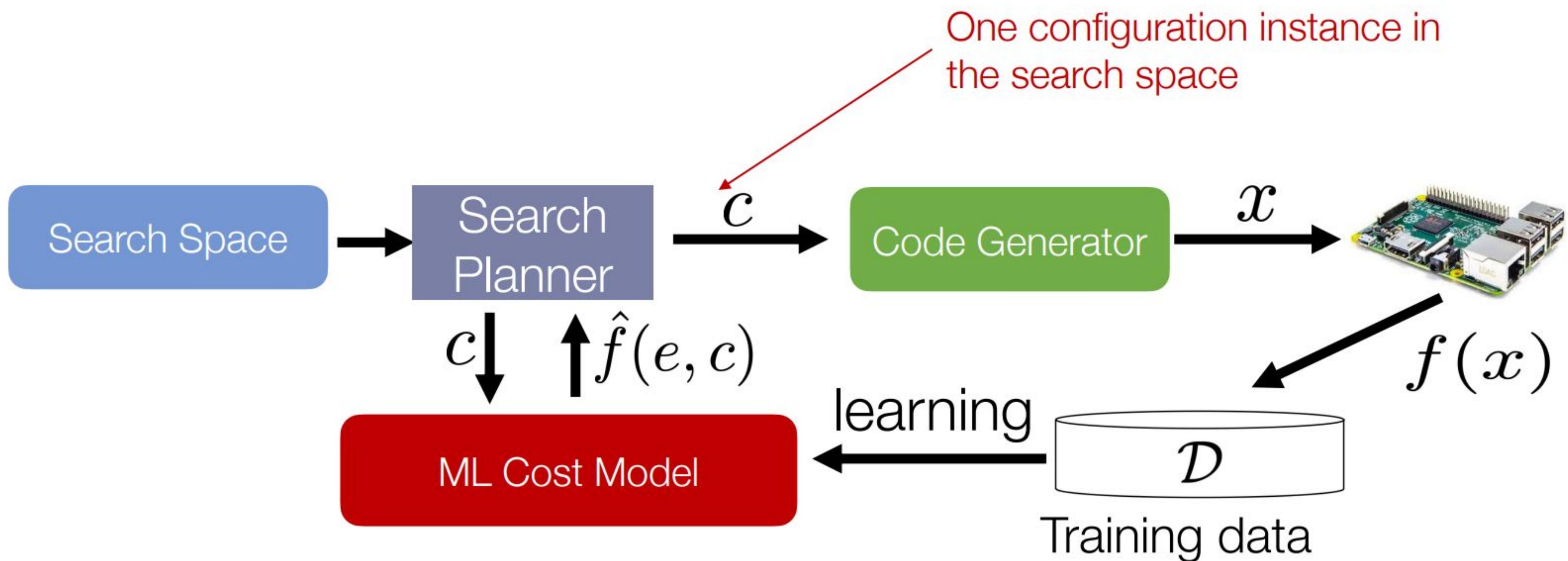
Problems

- We need to enumerate so many possibilities
- We need to fit with each device (register/cache sizes)
- We need to apply this to so many operators

Core Research Problems

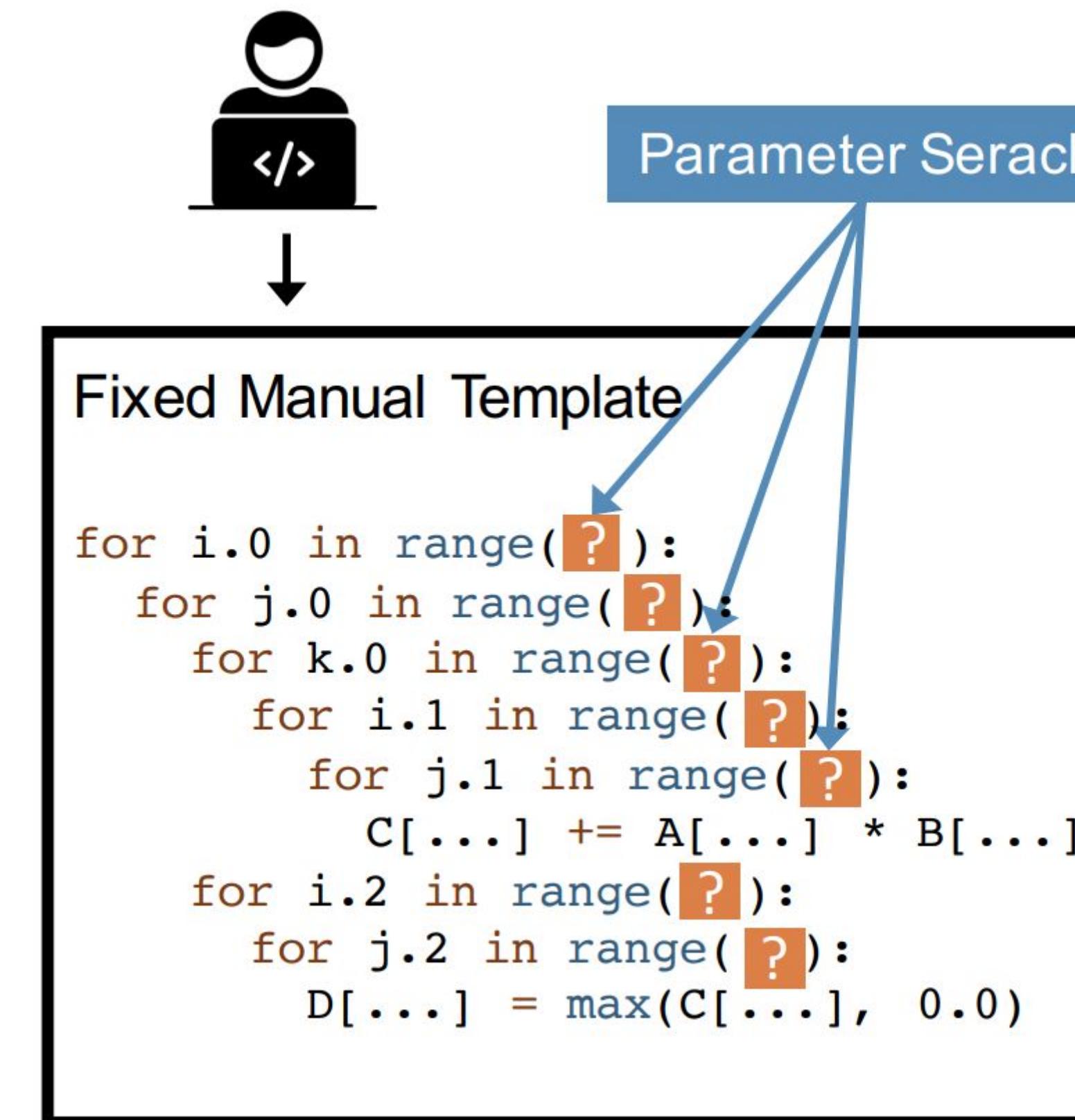
- We need to enumerate so many possibilities
 - How to represent all possibilities
 - What is the problem of missing some possibilities?
- We need to find the (close-to-)optimal values(register/cache sizes)
 - How to search?
- We need to apply this to so many operators and devices
 - How to reduce search space
 - How to generalize?

Search via Learned Cost Model



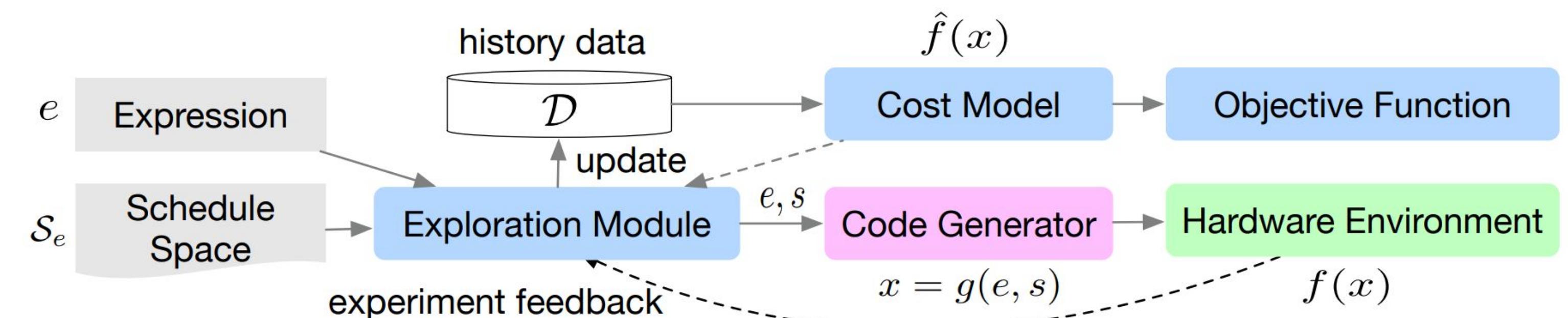
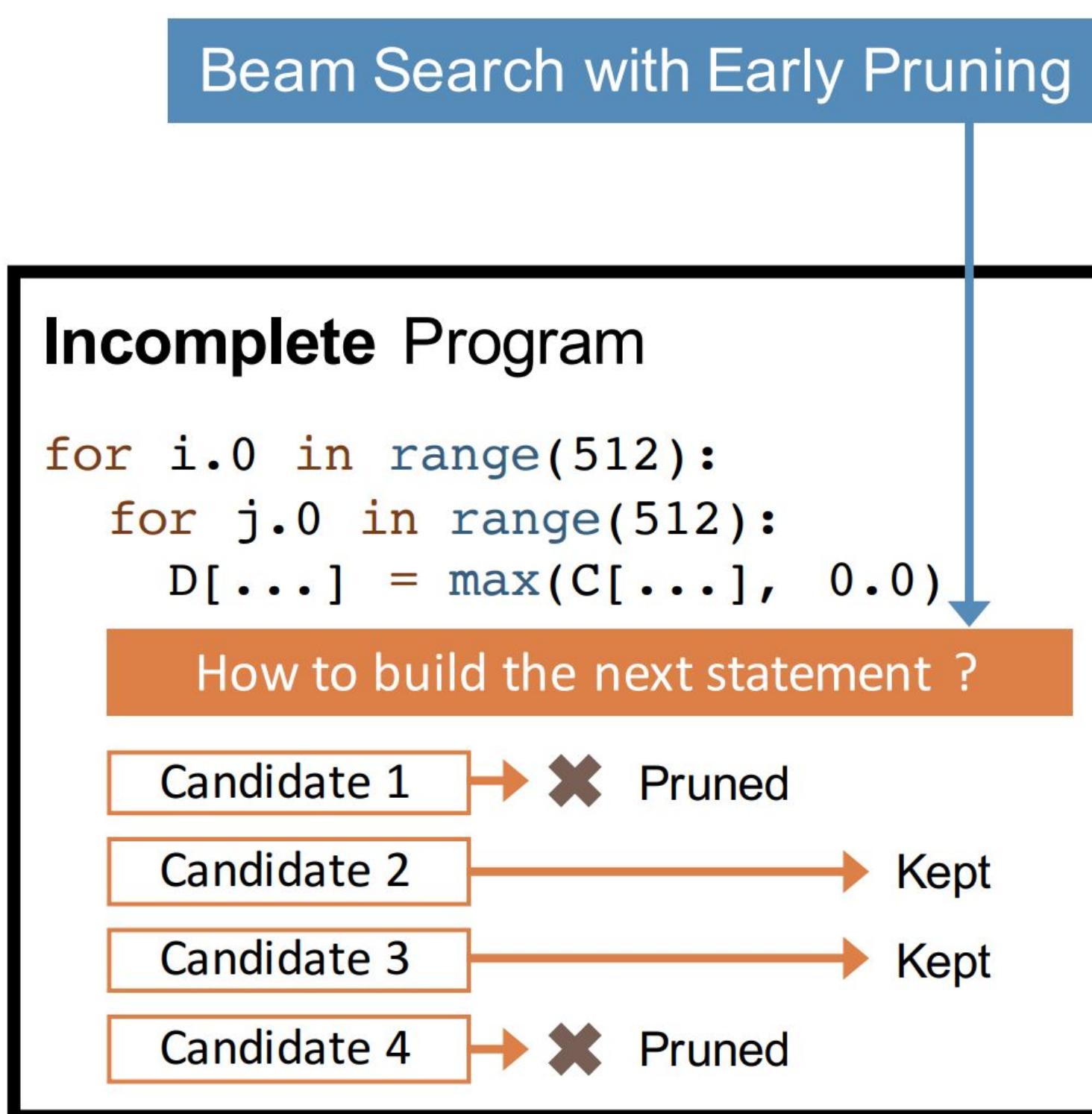
Search Space Definition e.g. Template based

- Issue: still need experts to write templates



How to Search

- Sequential Construction using Early pruning
- Cost Model



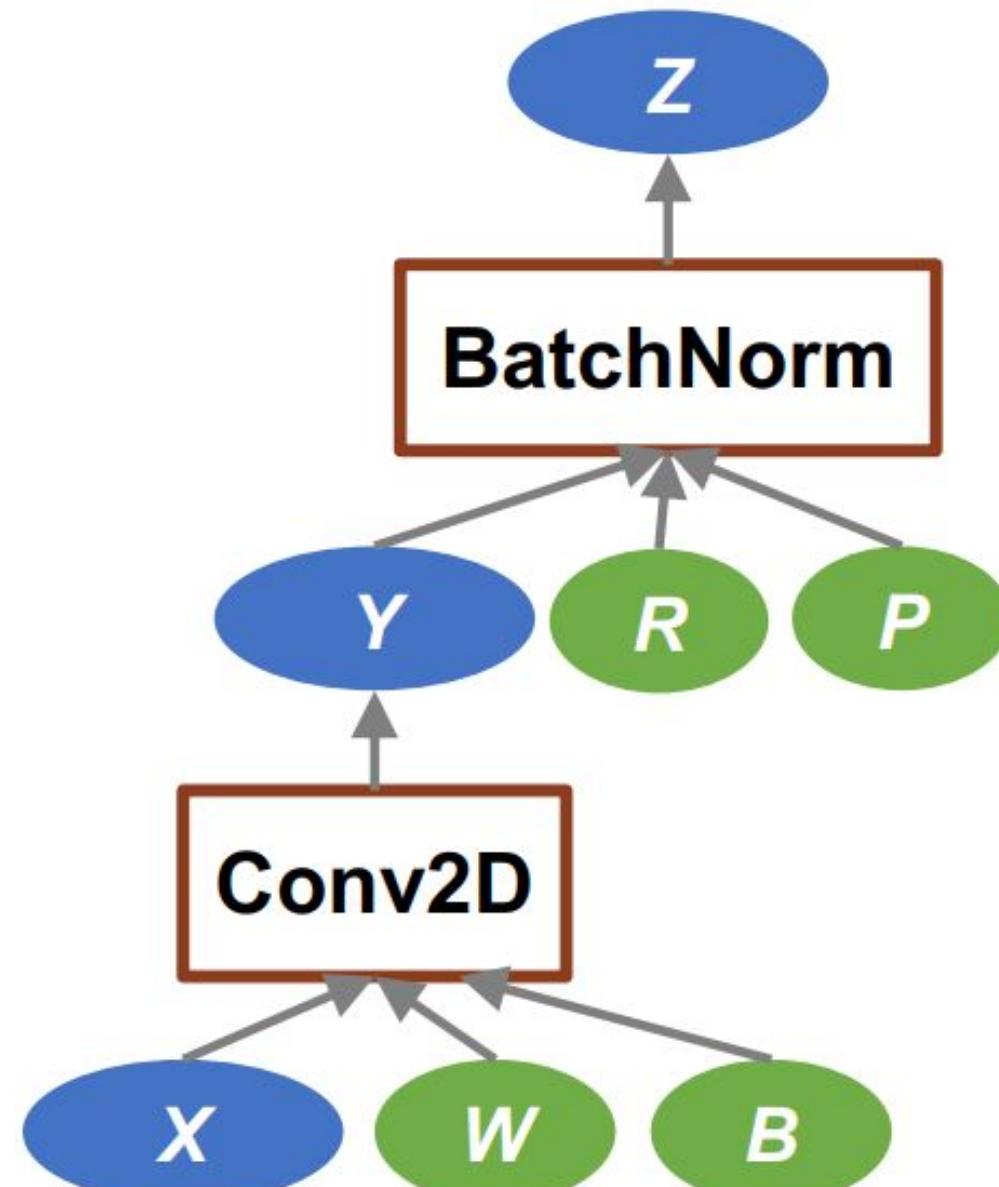
Summary: Operator Compiler

- Program abstraction
 - Represent the program/optimization of interest
- Build Search space through a set of transformations
 - Good coverage of common optimizations like tiling
- Effective Search
 - Accurate cost models
 - Transferability

Agenda on this part

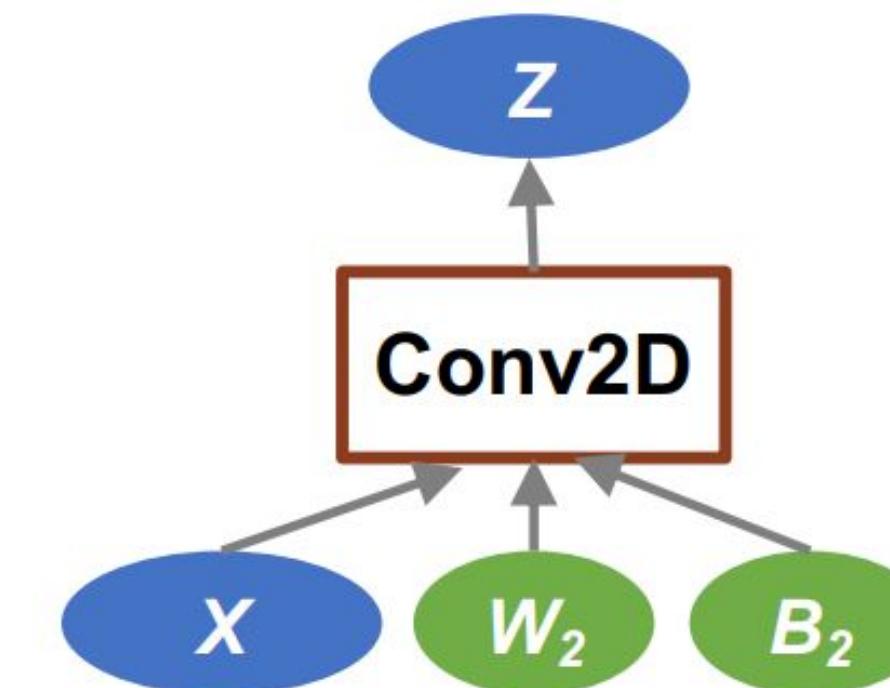
- ML Compilation Overview
 - Operator compilation
 - **Graph optimization**
- Memory Optimization
 - Activation checkpointing
 - Quantization and Mixed precision
- Two Guest Talks covering details in compilation, JIT, graph fusion, and beyond:
 - Meta PyTorch lead developer: Jason Ansel

Recall: fusing conv and bn



$$Z(n, c, h, w) = \left(\sum_{d, u, v} X(n, d, h + u, w + v) * W_2(c, d, u, v) \right) + B_2(n, c, h, w)$$

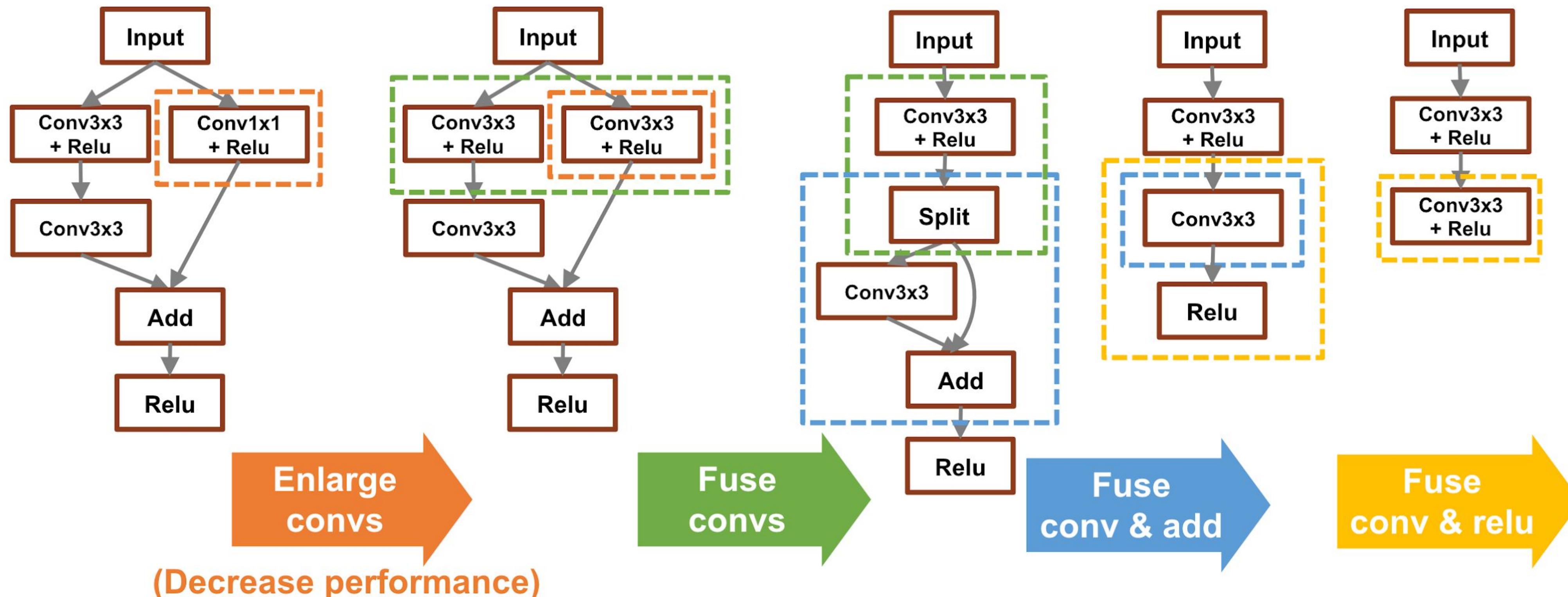
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$$W_2(n, c, h, w) = W(n, c, h, w) * R(c)$$

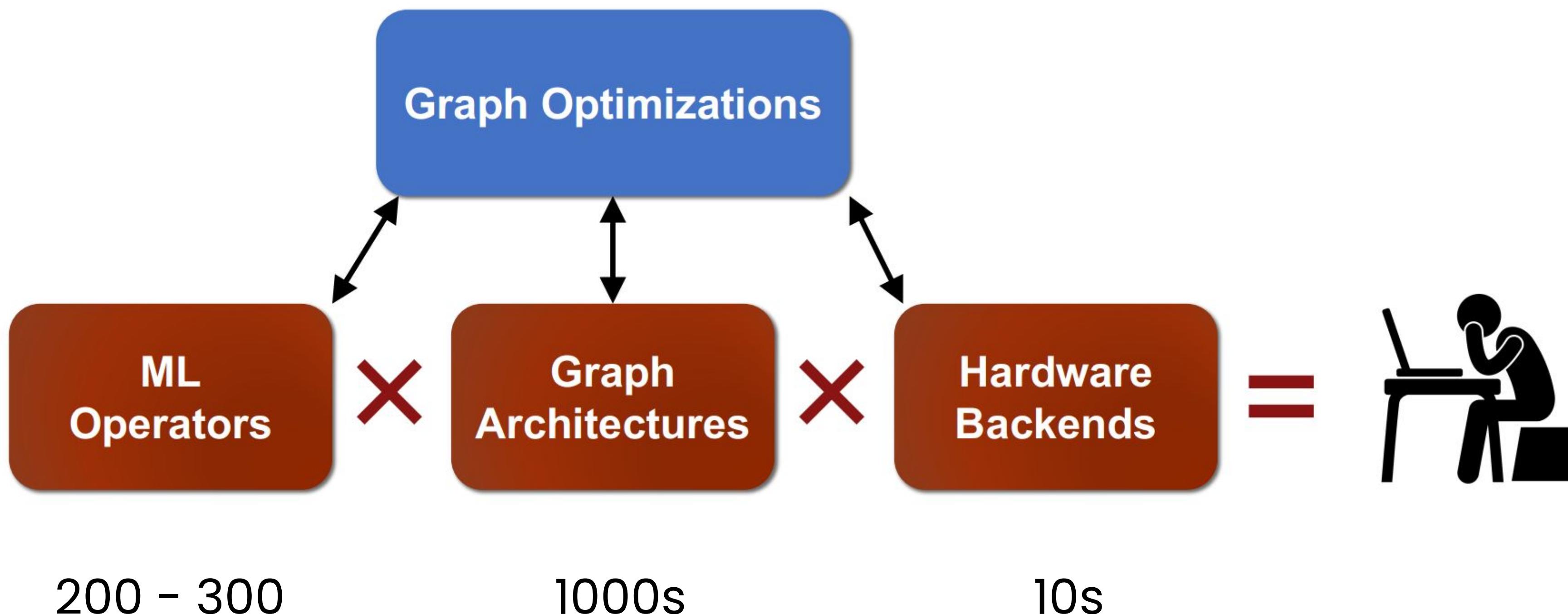
$$B_2(n, c, h, w) = B(n, c, h, w) * R(c) + P(c)$$

Recall: ResNet



- The final graph is 30% faster on V100 but 10% slower on K80.

Problems of High-level Graph Optimizations



Problem: Infeasible to manually design graph optimizations for all cases

Summary of Limitations

Robustness

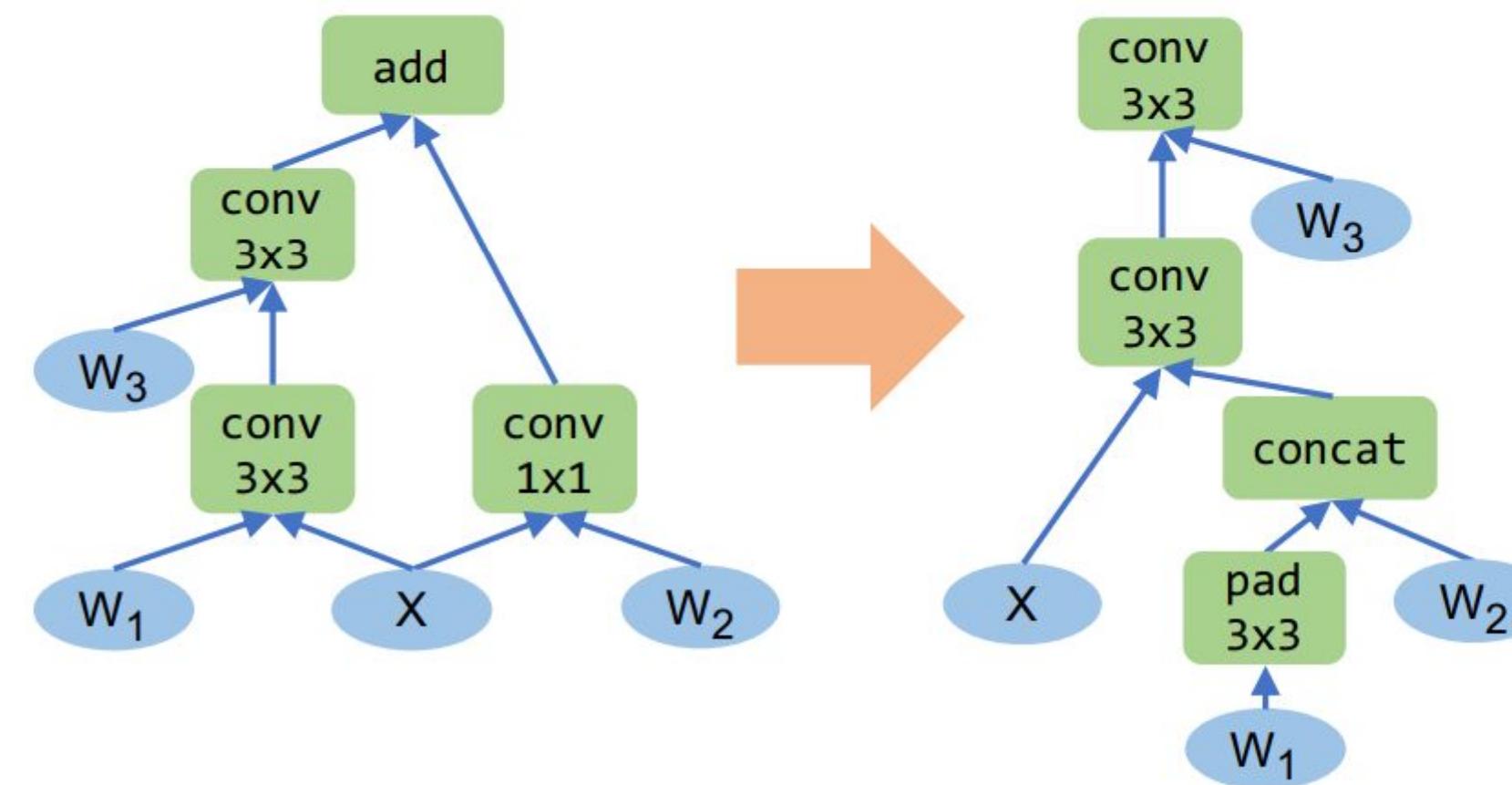
Experts' heuristics do not apply to all DNNs/hardware

Scalability

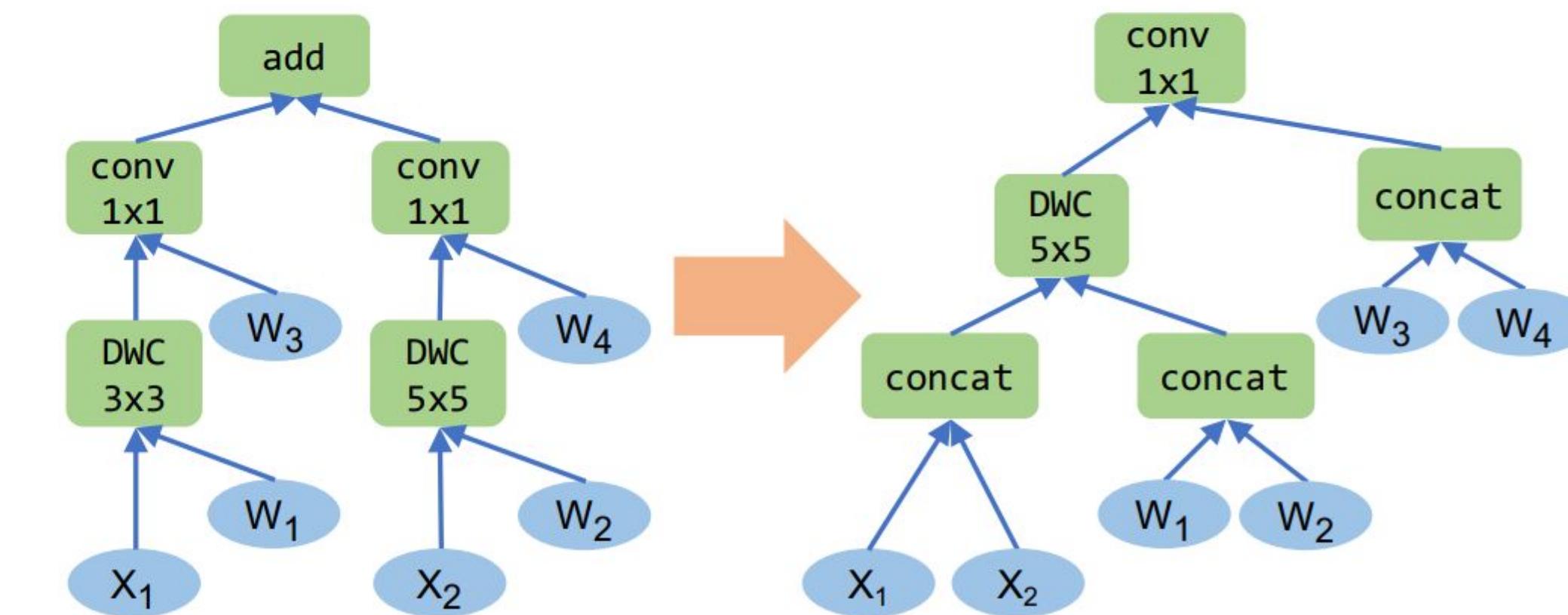
New operators and graph structures require more rules

Performance

Miss subtle optimizations for specific DNNs/hardware



Only apply to **specific hardware**



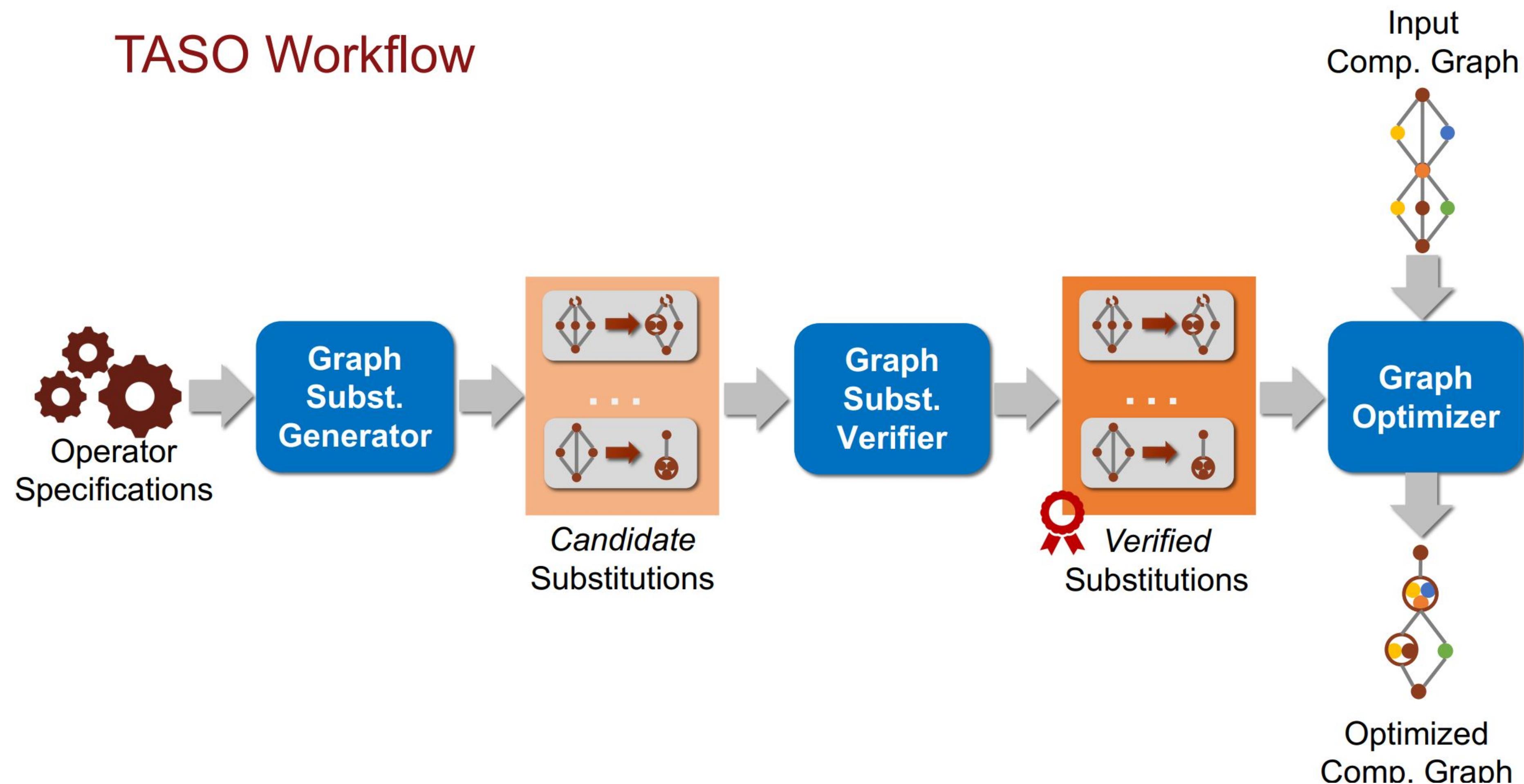
Only apply to **specialized graph structures**

Automate Graph Transformation Big Ideas

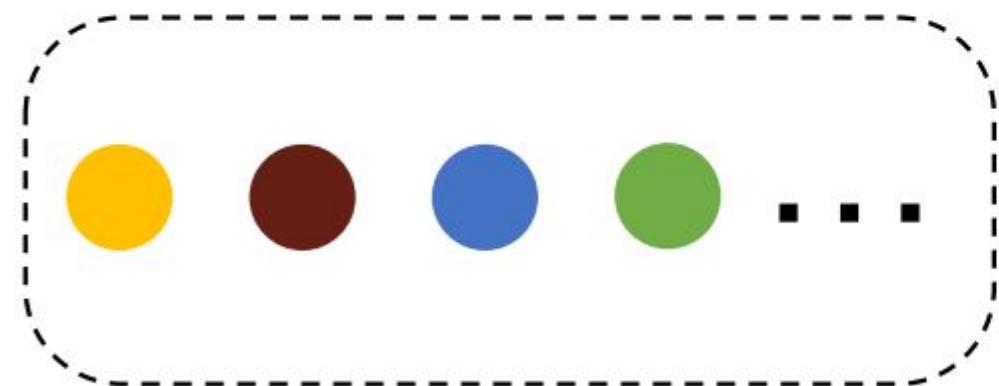
Key idea: replace manually-designed graph optimizations with automated generation and verification of graph substitutions for tensor algebra

- Less engineering effort: 53,000 LOC for manual graph optimizations in TensorFlow → 1,400 LOC
- Better performance: outperform existing optimizers by up to 3x
- Correctness: formally verified

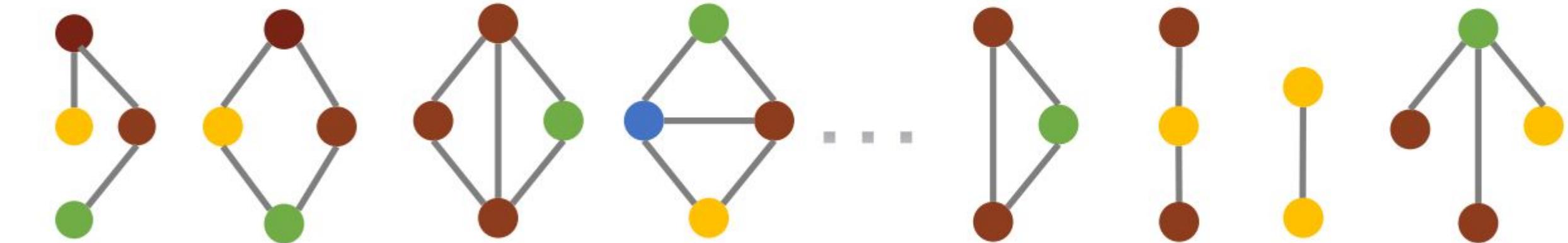
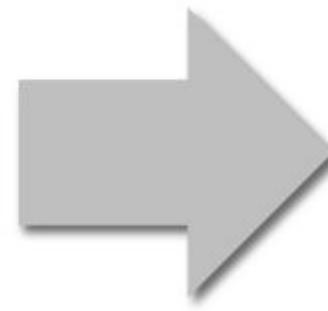
TASO: Enumerate and Verify ALL possible graph



Graph Substitution Generator



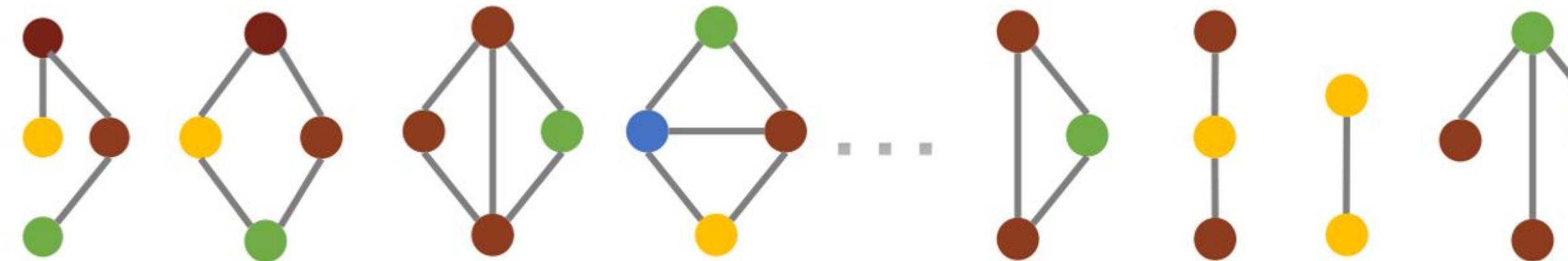
Operators supported by
hardware backend



Enumerate **all possible** graphs up to a
fixed size using available operators

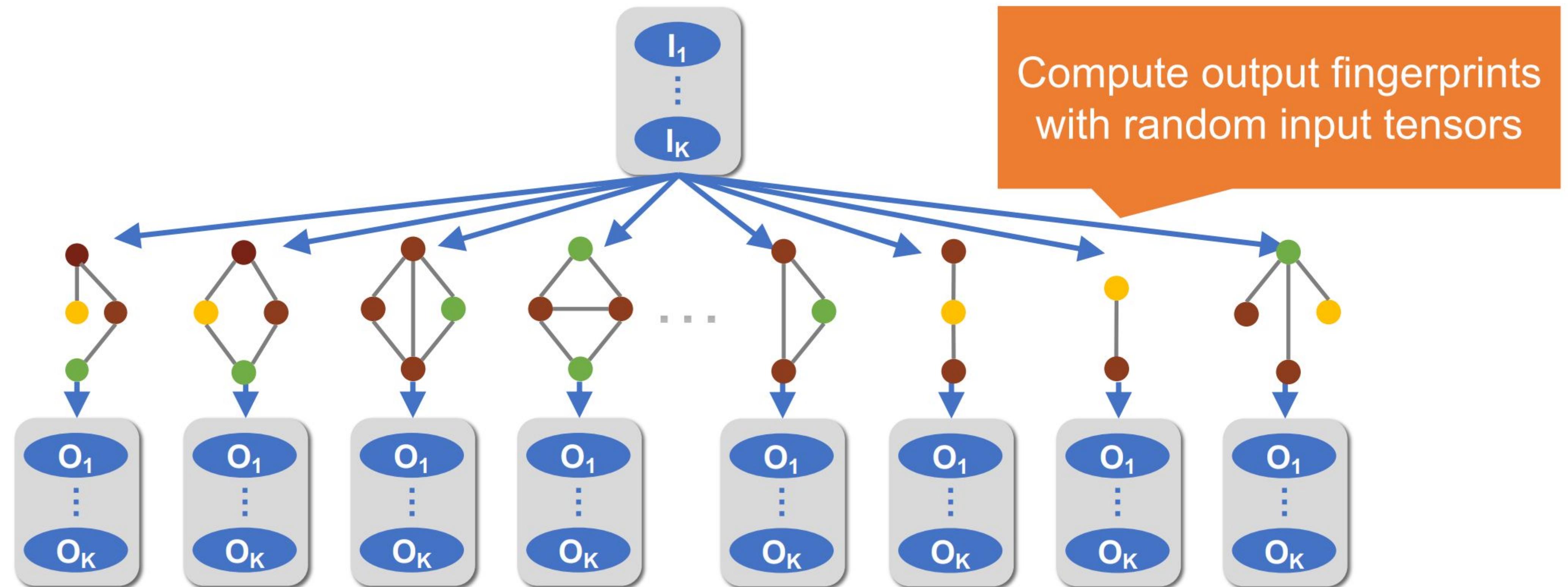
There are many subgraphs even only given 4 Ops

66M graphs with up to 4 operators



A substitution = a pair of equivalent graphs

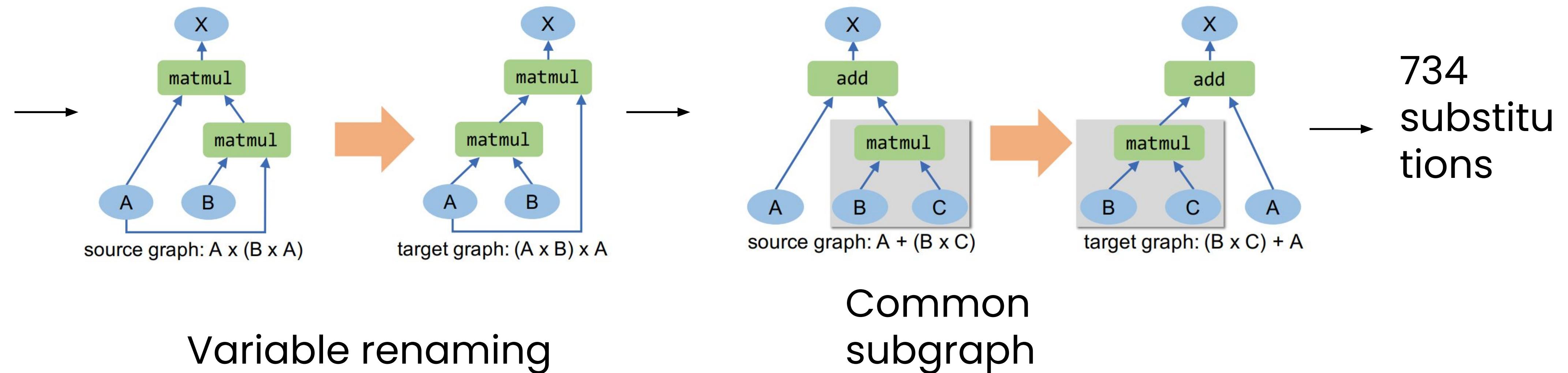
Graph Substitution Generator



TASO generates 28744 substitutions by enumerating graphs with up to 4 ops

Pruning repeated graphs

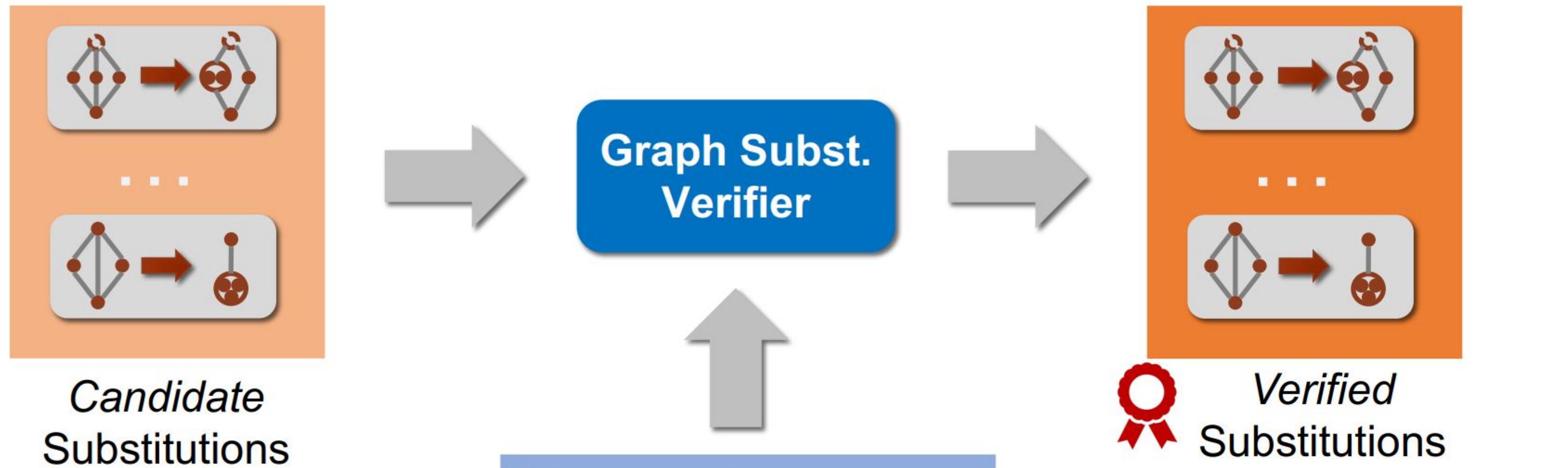
28744
substitution
s



Can we trust graph substitutions?

- We have $f(a) = g(b)$, $f(b) = g(b)$
 - But can we say: $f(x) = g(x)$ for $\forall x$
 - We need to verify formally.

Substitution Verifier



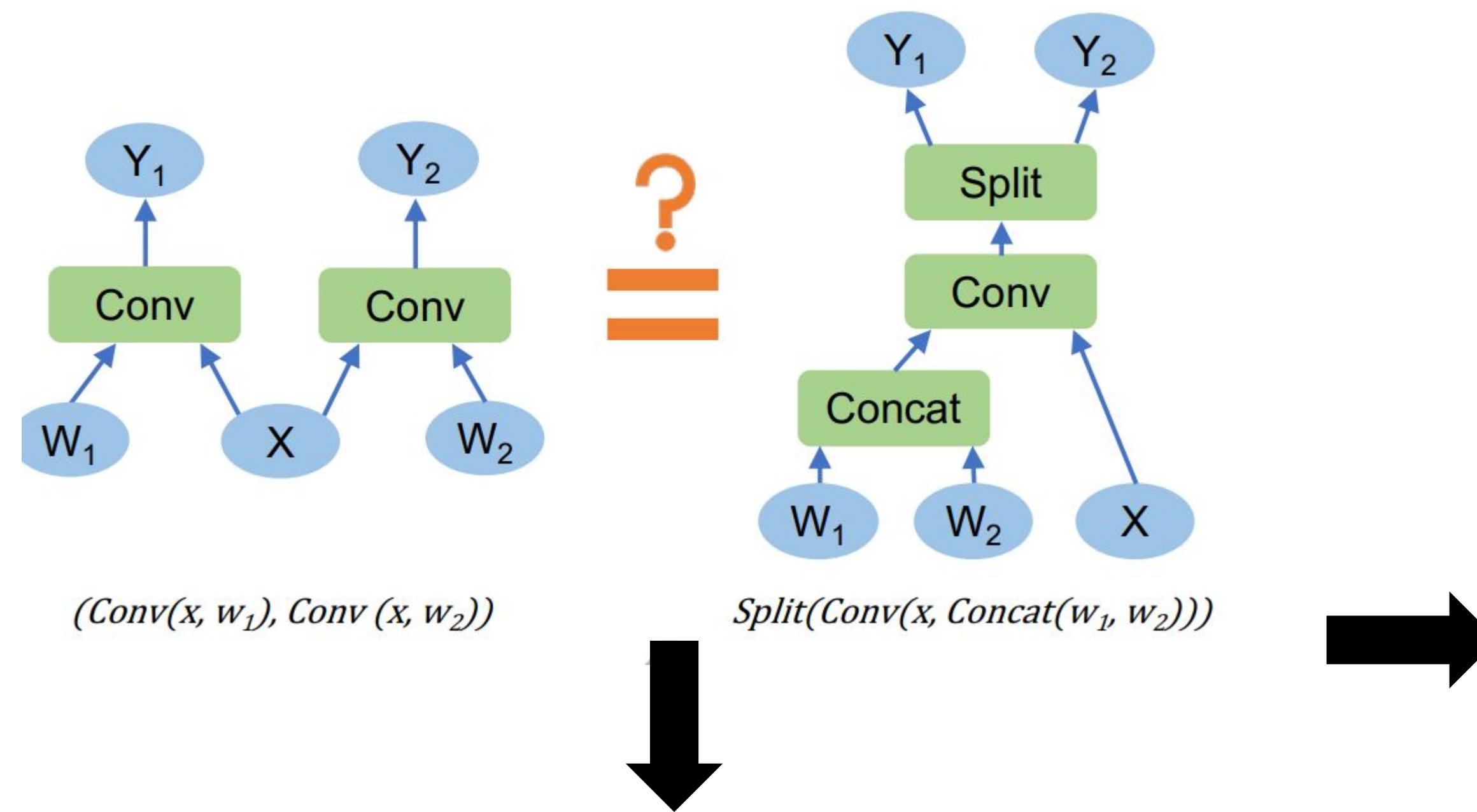
P1. conv is distributive
over concatenation
P2. conv is bilinear
...
Pn.

$$\forall x, w_1, w_2 . \quad Conv(x, Concat(w_1, w_2)) = \\ Concat(Conv(x, w_1), Conv(x, w_2))$$

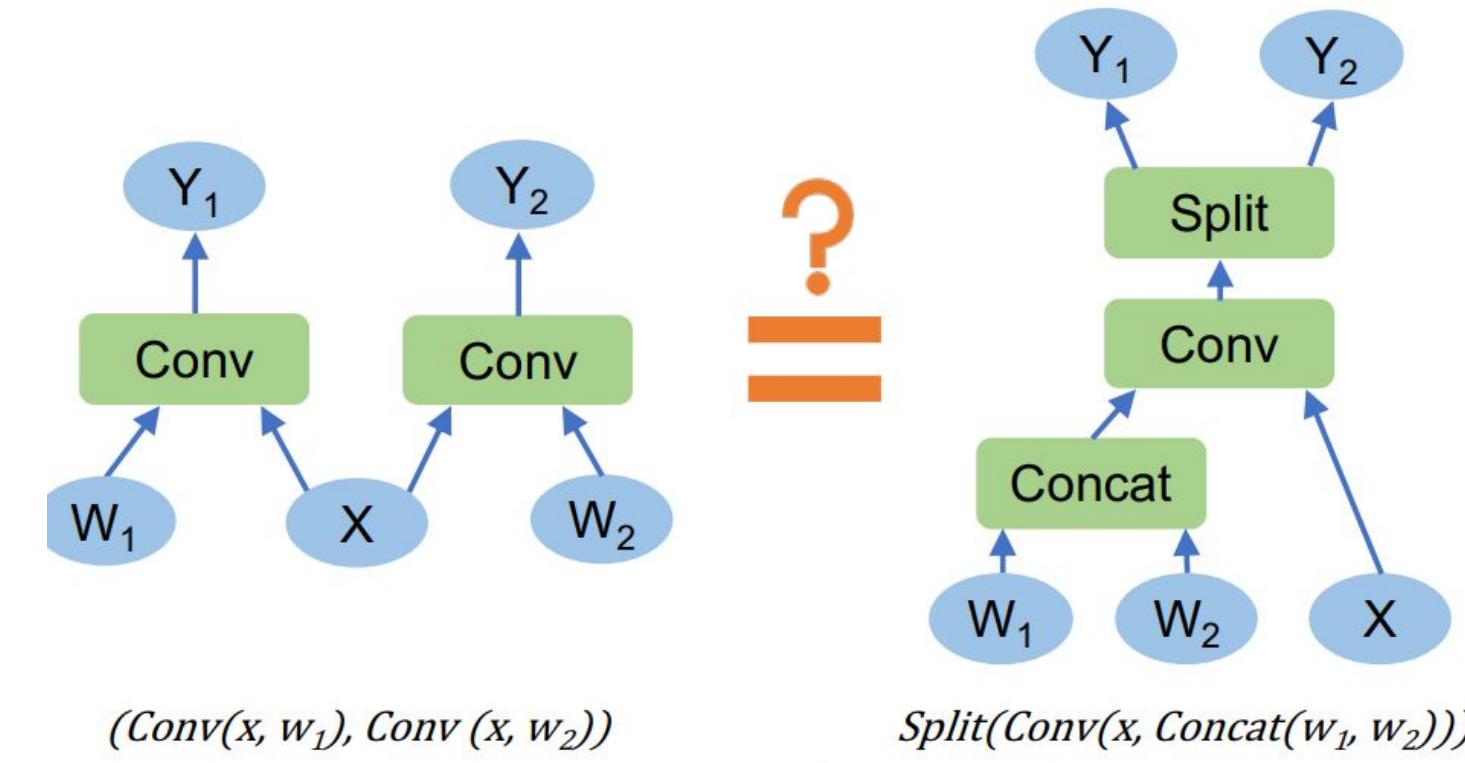


Idea: writing specifications are easier than
actually, conducting the optimizations

How to Verify



Automated
theorem
prover



$\forall x, w_1, w_2 .$
 $(Conv(x, w_1), Conv(x, w_2))$
 $= Split(Conv(x, Concat(w_1, w_2)))$

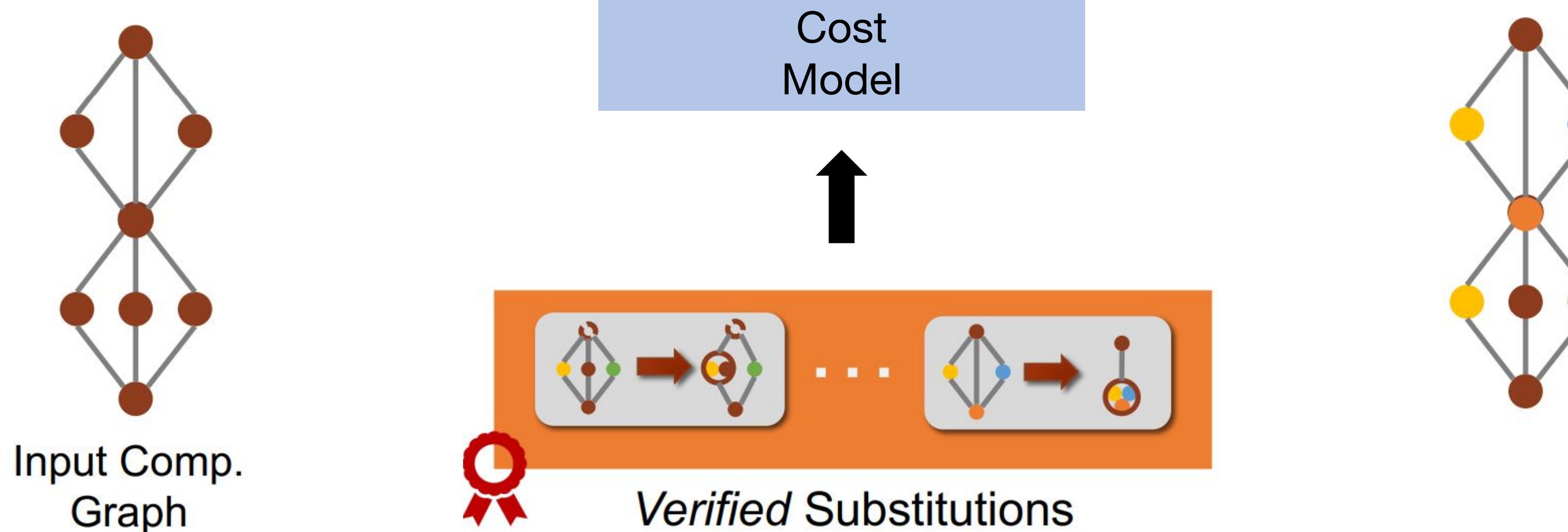
P1. $\forall x, w_1, w_2 .$
 $Conv(x, Concat(w_1, w_2)) =$
 $Concat(Conv(x, w_1), Conv(x, w_2))$

P2. ...

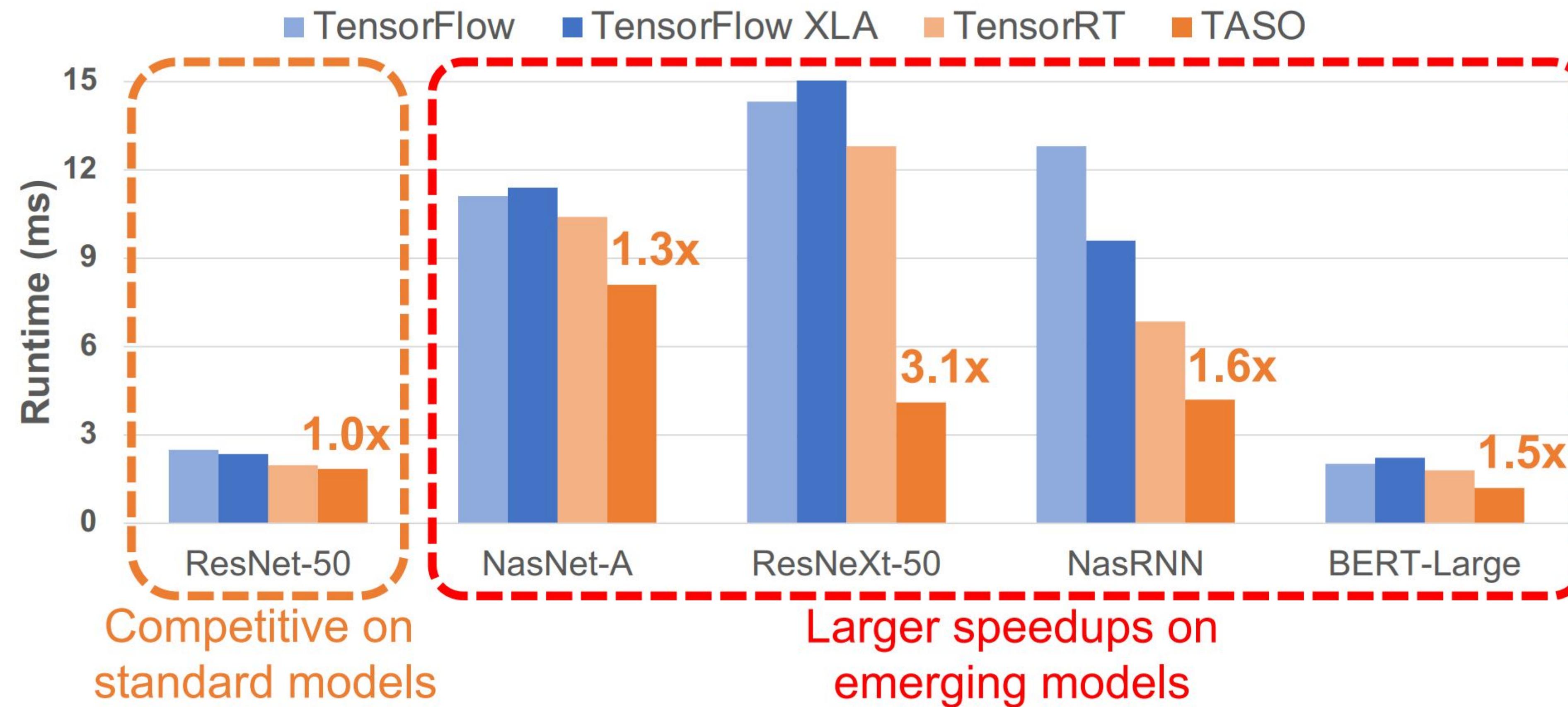
- Generating 743 substitutions = 5 mins
- Verify against 43 op specs = 10 mins
- Supporting a new op requires experts to write specs = 1400 LoC
- vs. 53K LoC of manual optimization in TF

Incorporating substitutions

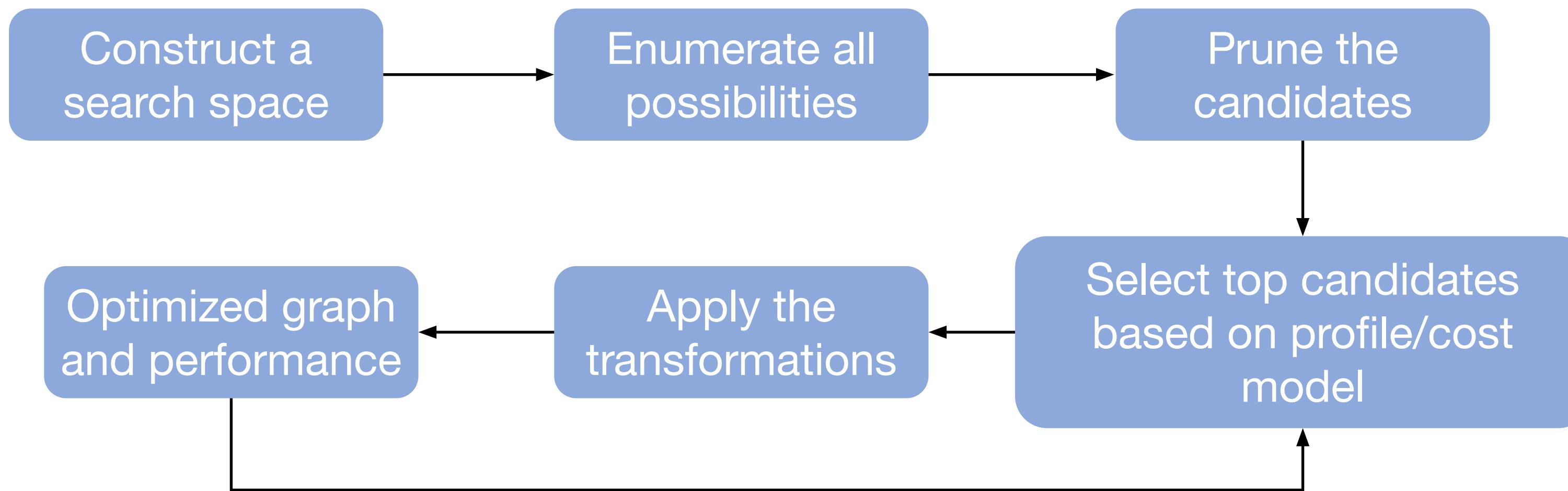
- Goal: apply verified substitutions to obtain an optimized graph
- Cost Model
 - Based on the sum of individual operator's cost
 - Profile each operator's cost on the target hardware
- Traverse the graph, apply substitutions, calculate cost, use backtracking



Performance (as of 2019)



Summary of Graph Optimization

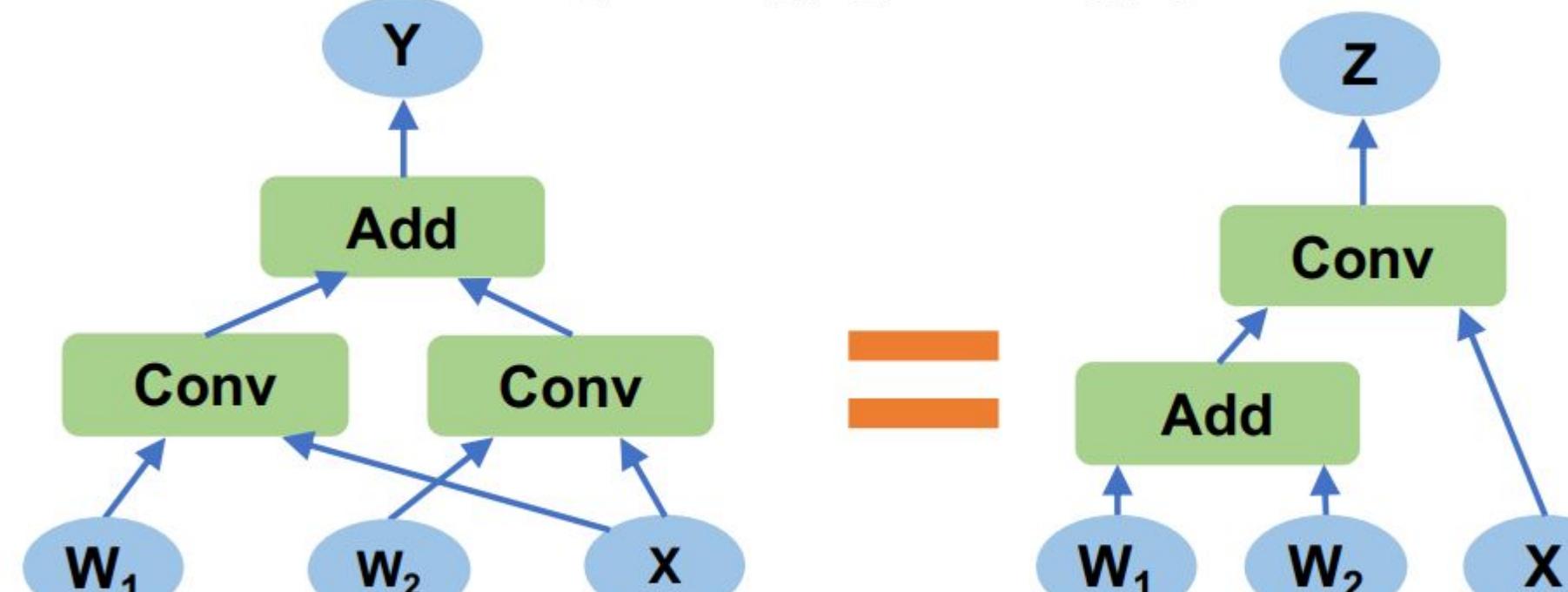


Limitations

- The best optimization is not covered by search space
- Search is too slow
- Evaluation of the resulting graph is too expensive
 - Limits your trial-and-error times

A Failure Example

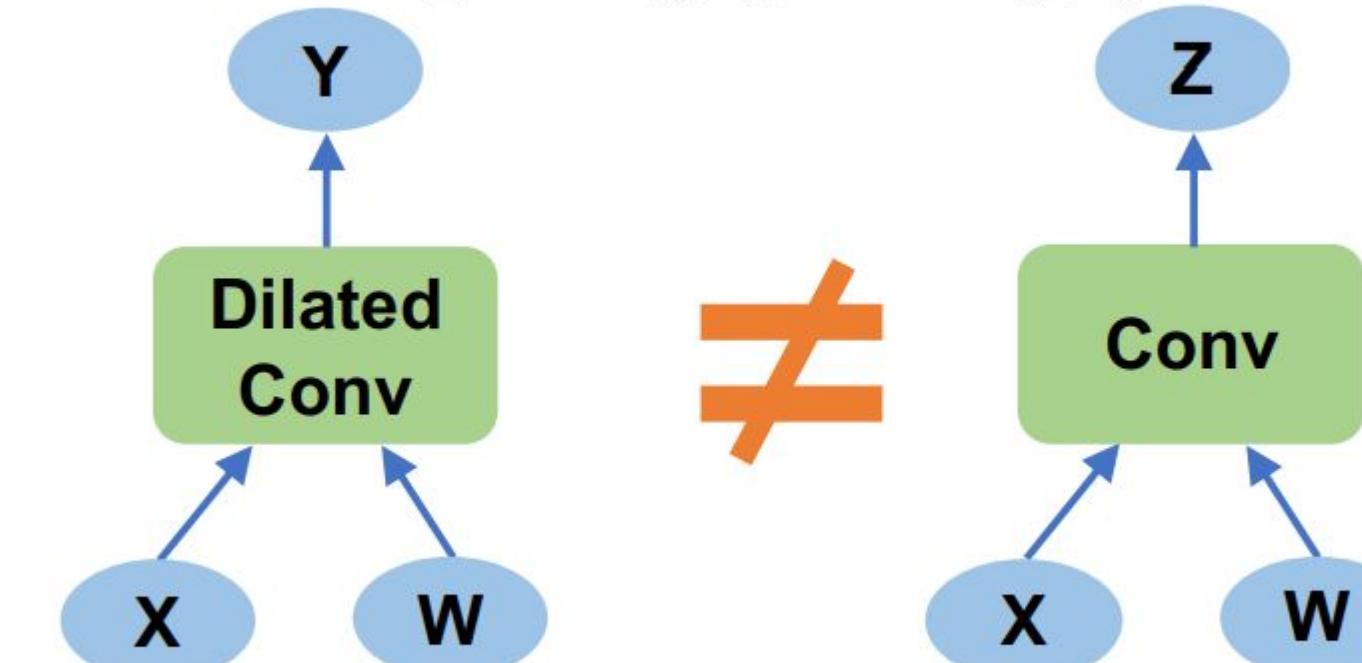
$$\forall p. Y[p] = Z[p]$$



Fully Equivalent Transformations

- Math-equivalent
- Missing some optimization opportunities
- Better performance
- Not fully equivalent -> accuracy loss

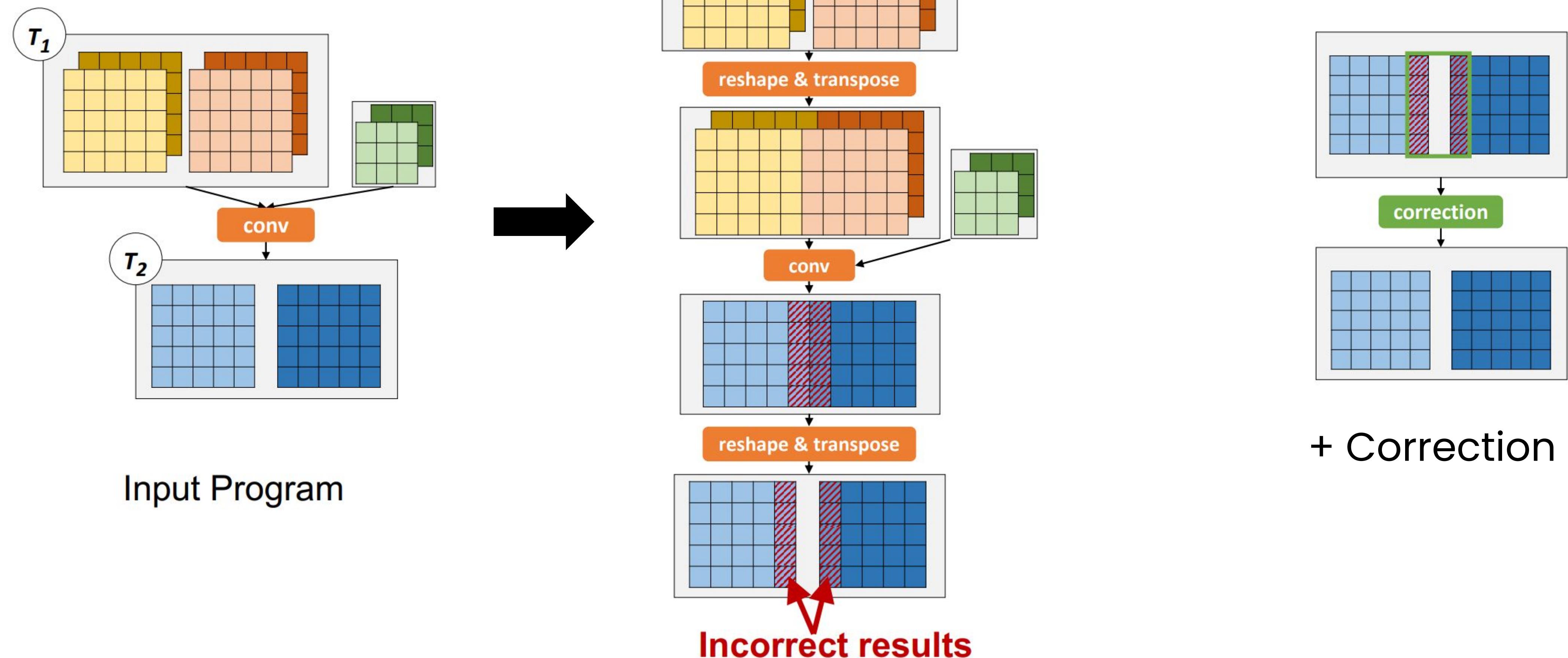
$$\exists p. Y[p] \neq Z[p]$$



Partially Equivalent Transformations

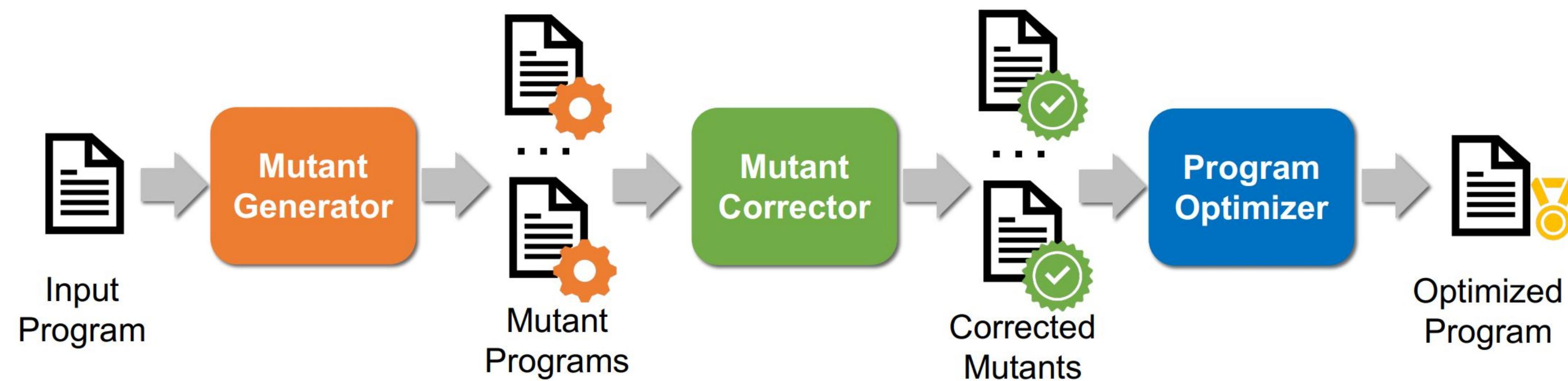
How about: exploit the larger space partially equivalent transformations for performance while still preserve correctness?

Motivating Example



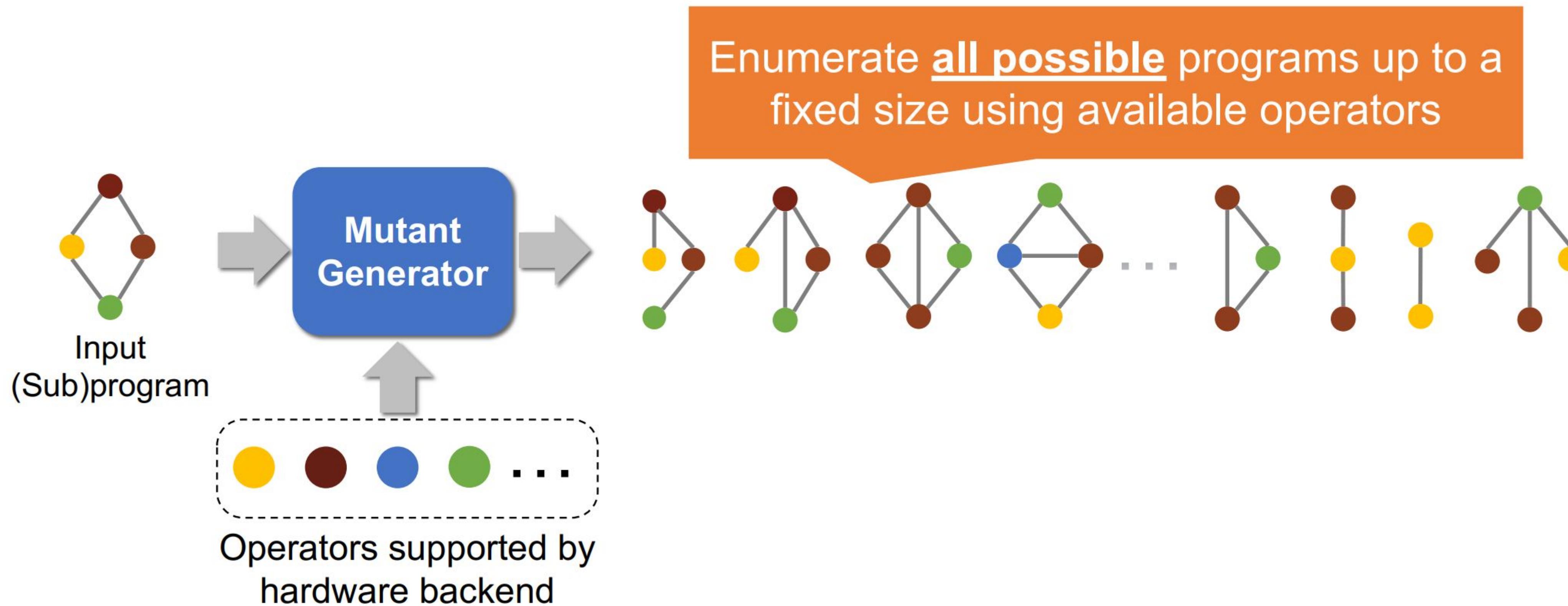
- Partial equivalent transformations + correction yield 1.2x speedup
- Which would otherwise be impossible in fully equivalent transformations

Partially Equivalent Transformations

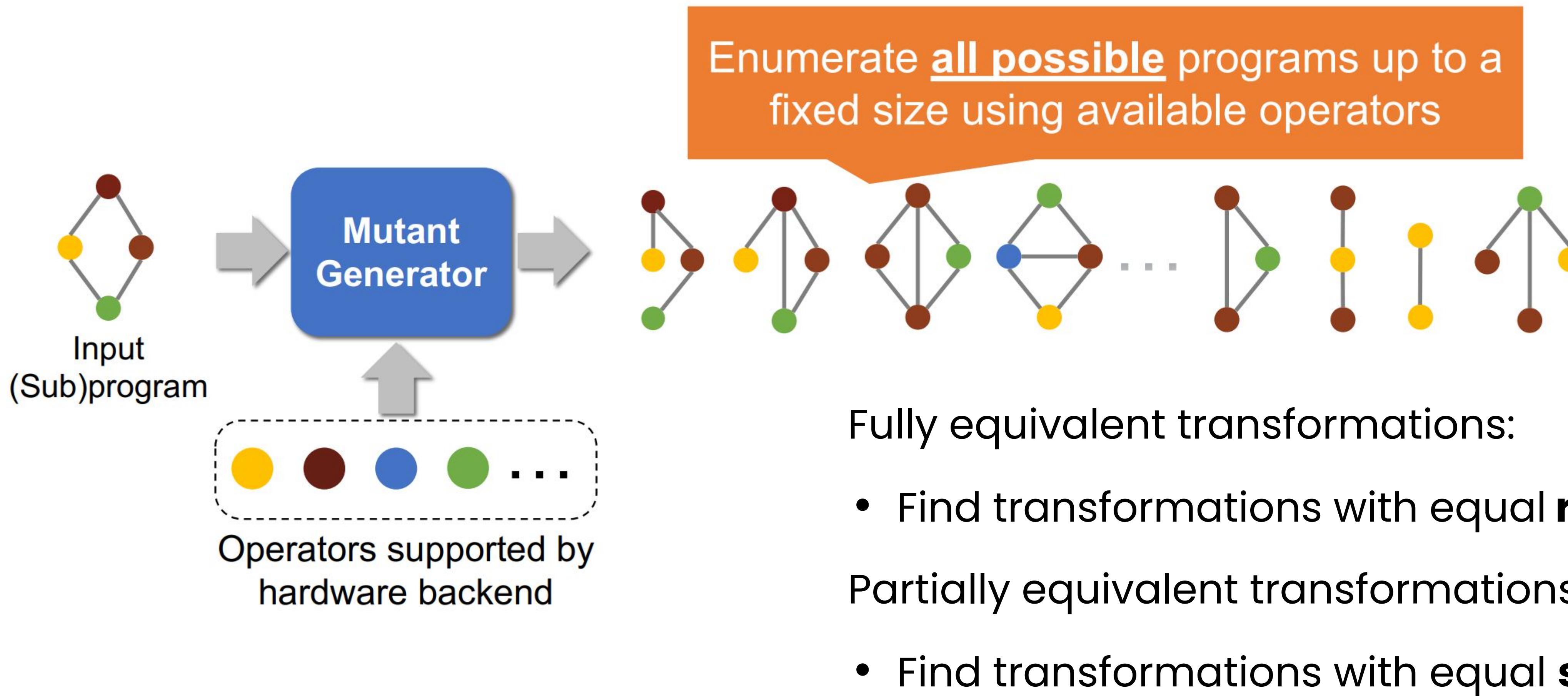


- How to mutate?
- How to correct?

Mutant Generator: Step 1

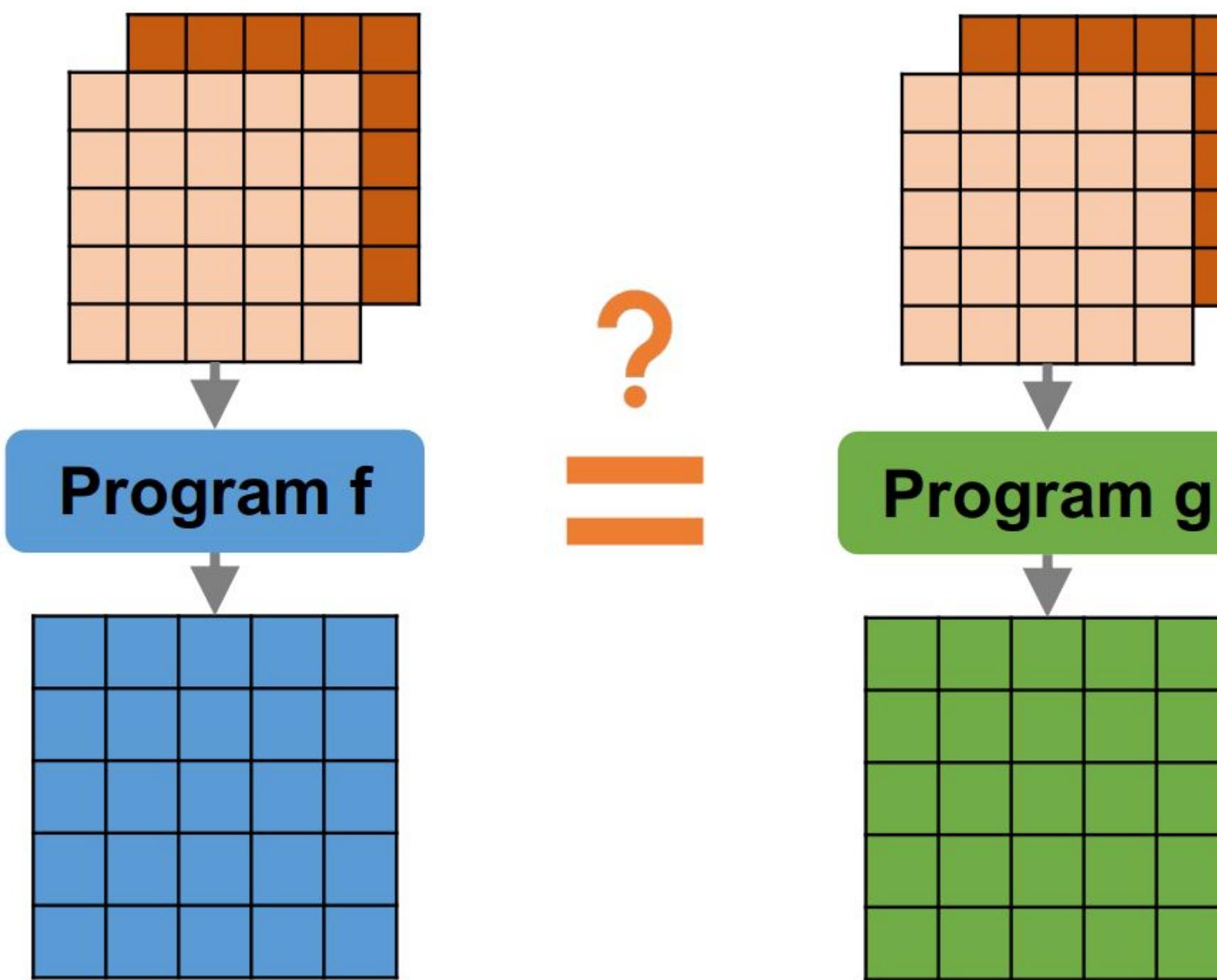


Mutant Generator: Step 2



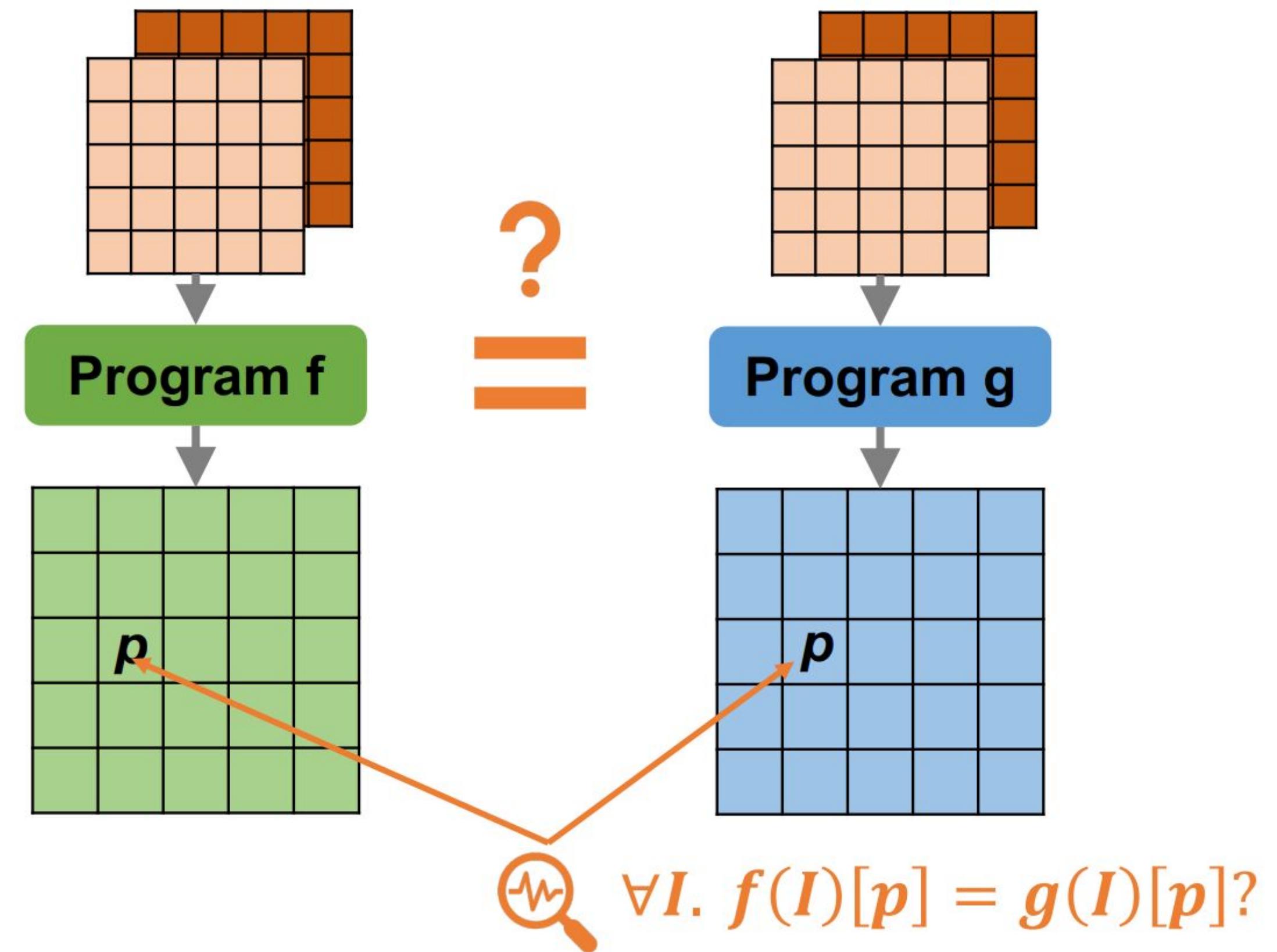
How to Detect and Correct?

- Which part of the computation is not equivalent?
- How to correct the results?



By Enumeration

- For each possible input I
 - For each position p
 - Check if $f(I)[p] == g(I)[p]$
- Complexity $O(m \times n)$:
 - m : possible inputs
 - n : output shape
- How to reduce enumeration effort?

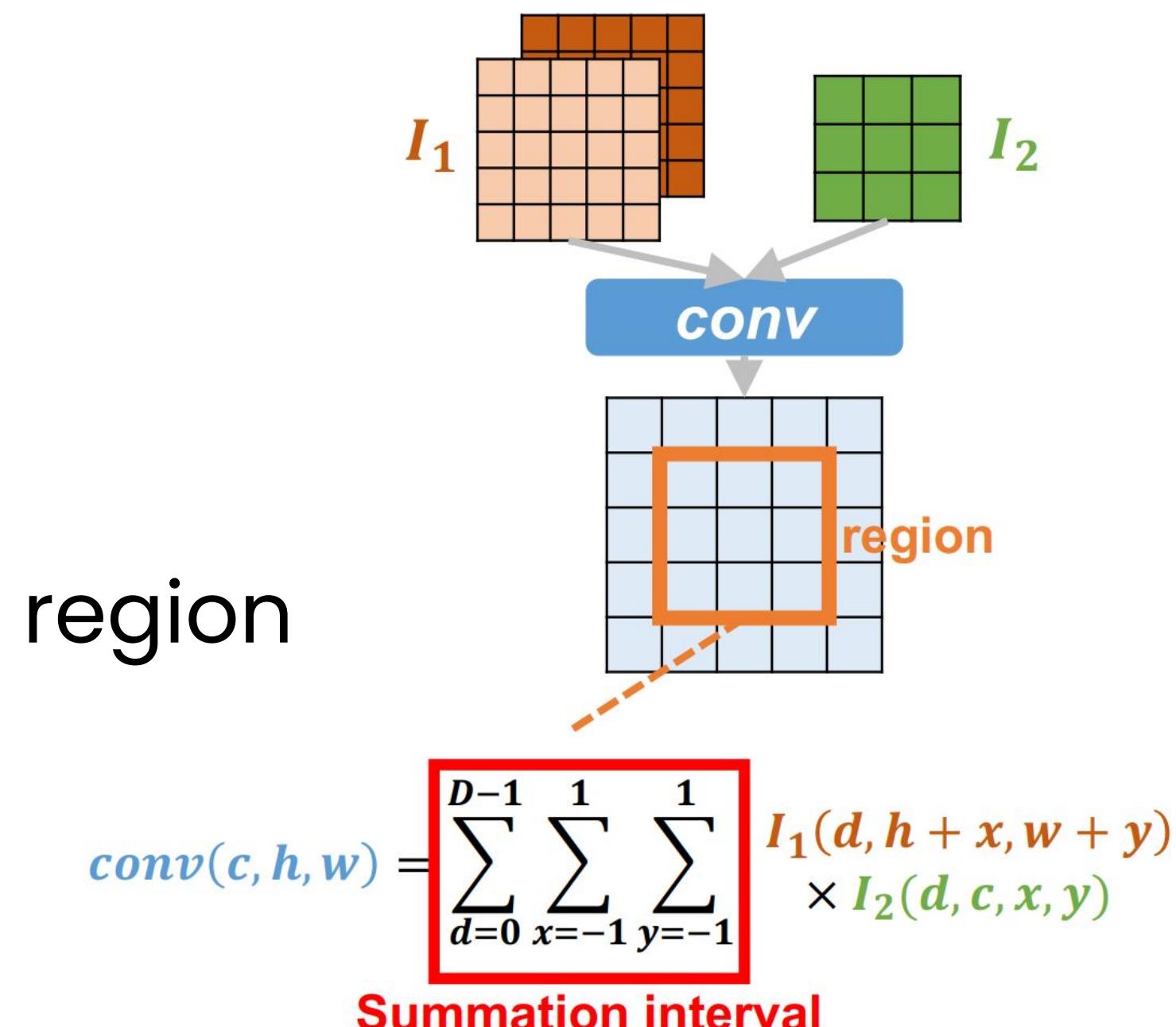


How to reduce n?

- Can we just check out a few (or even just one) position at $f(l)[p]$ and assert the (in-)correctness?
- Answer: Yes for 80% of the computation
- Reason: Neural nets computation are mostly Multi-Linear
- Define Multi-linear: f is multi-linear if the output is linear to all inputs
 - $f(I_1, \dots, X, \dots, I_n) + f(I_1, \dots, Y, \dots, I_n) = f(I_1, \dots, X + Y, \dots, I_n)$
 - $\alpha f(I_1, \dots, X, \dots, I_n) = f(I_1, \dots, \alpha X, \dots, I_n)$

How to reduce n

- Theorem 1: For two Multi-linear functions f and g , if $f=g$ for $O(1)$ positions in a region, then $f=g$ for all positions in the region
- Implications: only need to examine $O(1)$ positions for each region
 - Reduce $O(mn) \rightarrow O(m)$
Group all output positions with an identical summation interval into a region

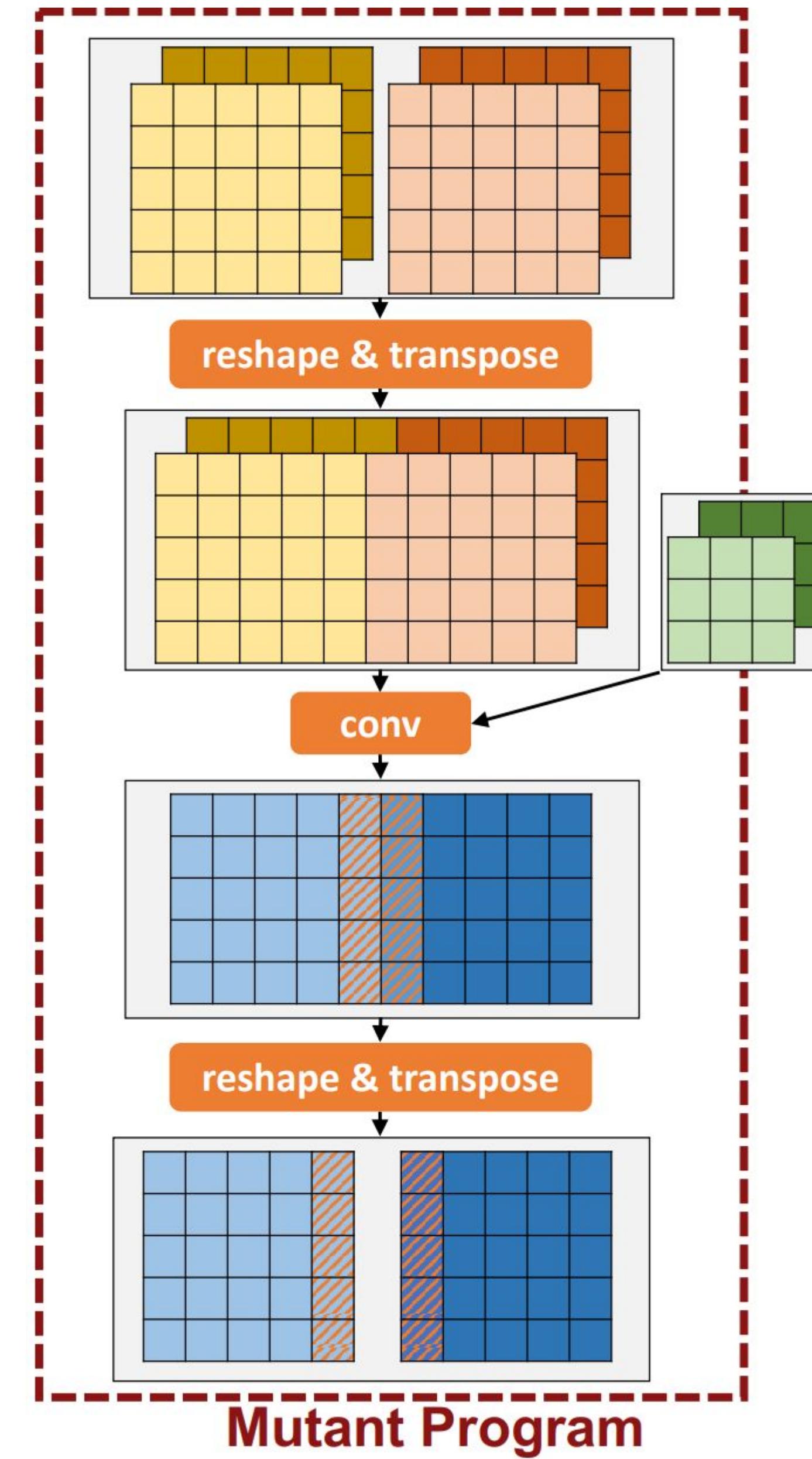


How to reduce m?

- Theorem 2: if $\exists I, f(I)[p] \neq g(I)[p]$, then the probability that f and g give identical results on t random inputs is $\left(\frac{1}{2^{31}}\right)^t$
- Implications: Run t random tests with random input, and if all t passed, it is very unlikely f and g are inequivalent
- $O(mn) \rightarrow O(m) \rightarrow O(t)$ ($t \ll m$)

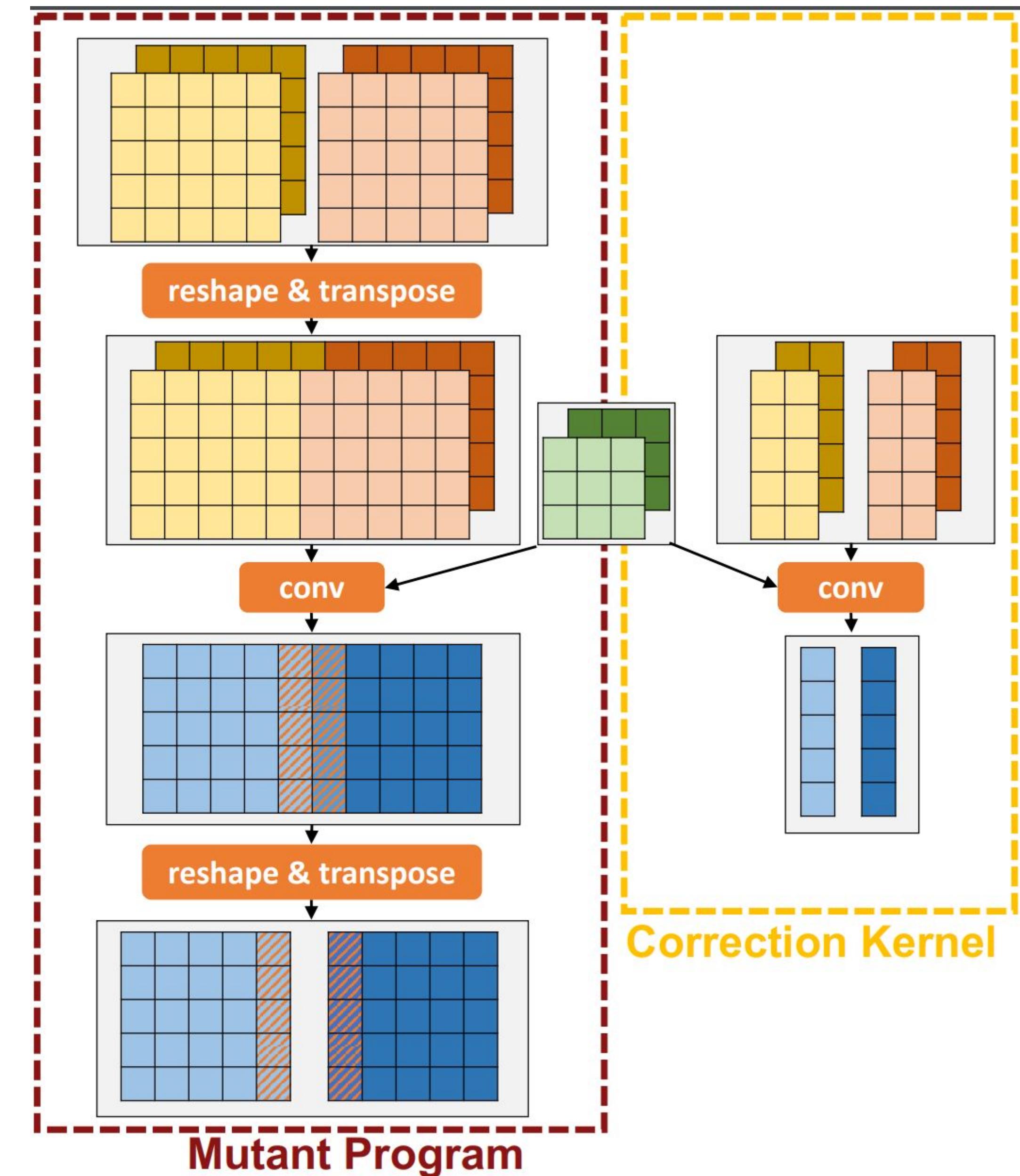
Correct the Mutant

- Goal: quickly and efficiently correcting the outputs of a mutant program



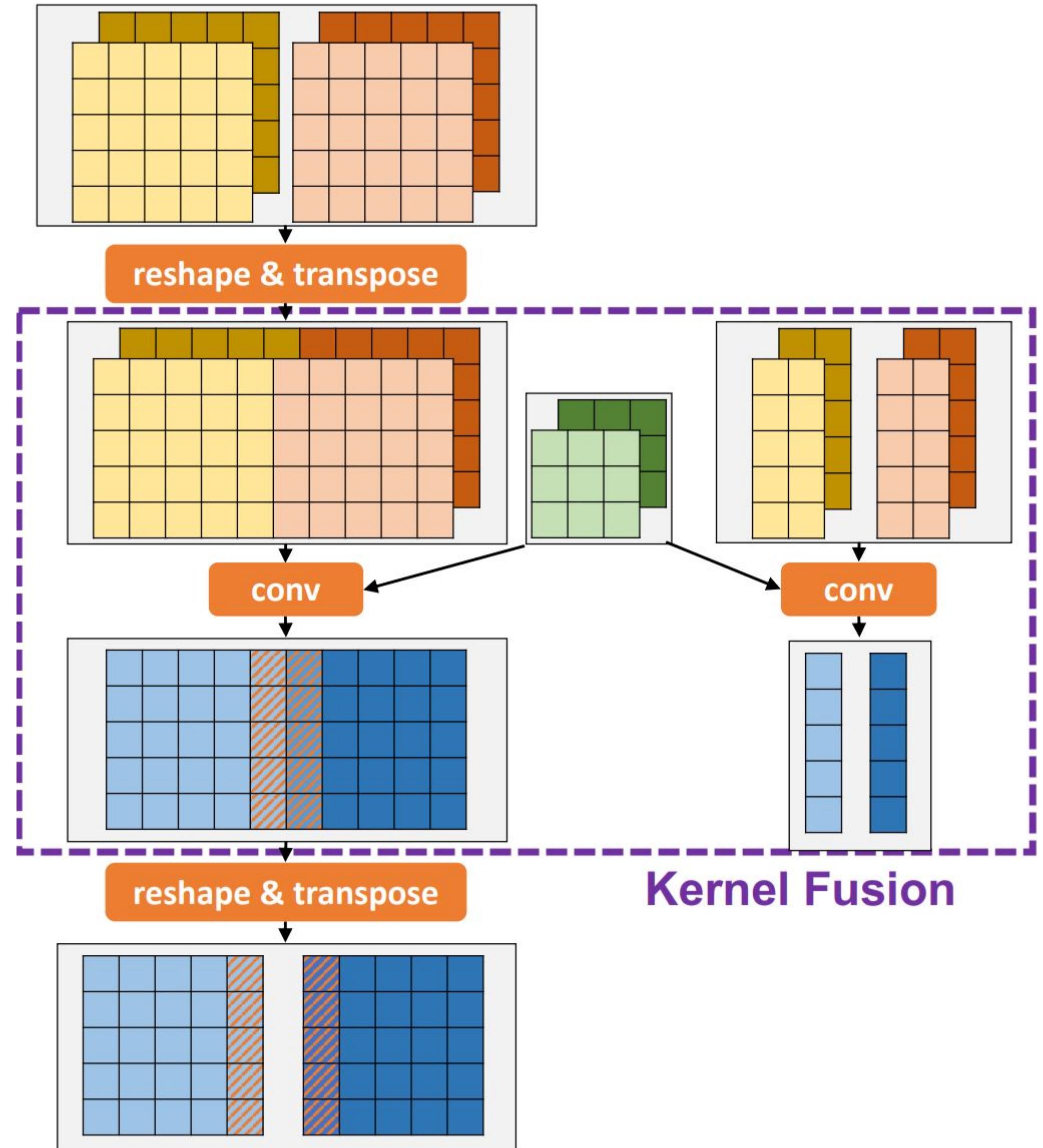
Correct the Mutant

- Goal: quickly and efficiently correcting the outputs of a mutant program
- Step 1: recompute the incorrect outputs using the original program

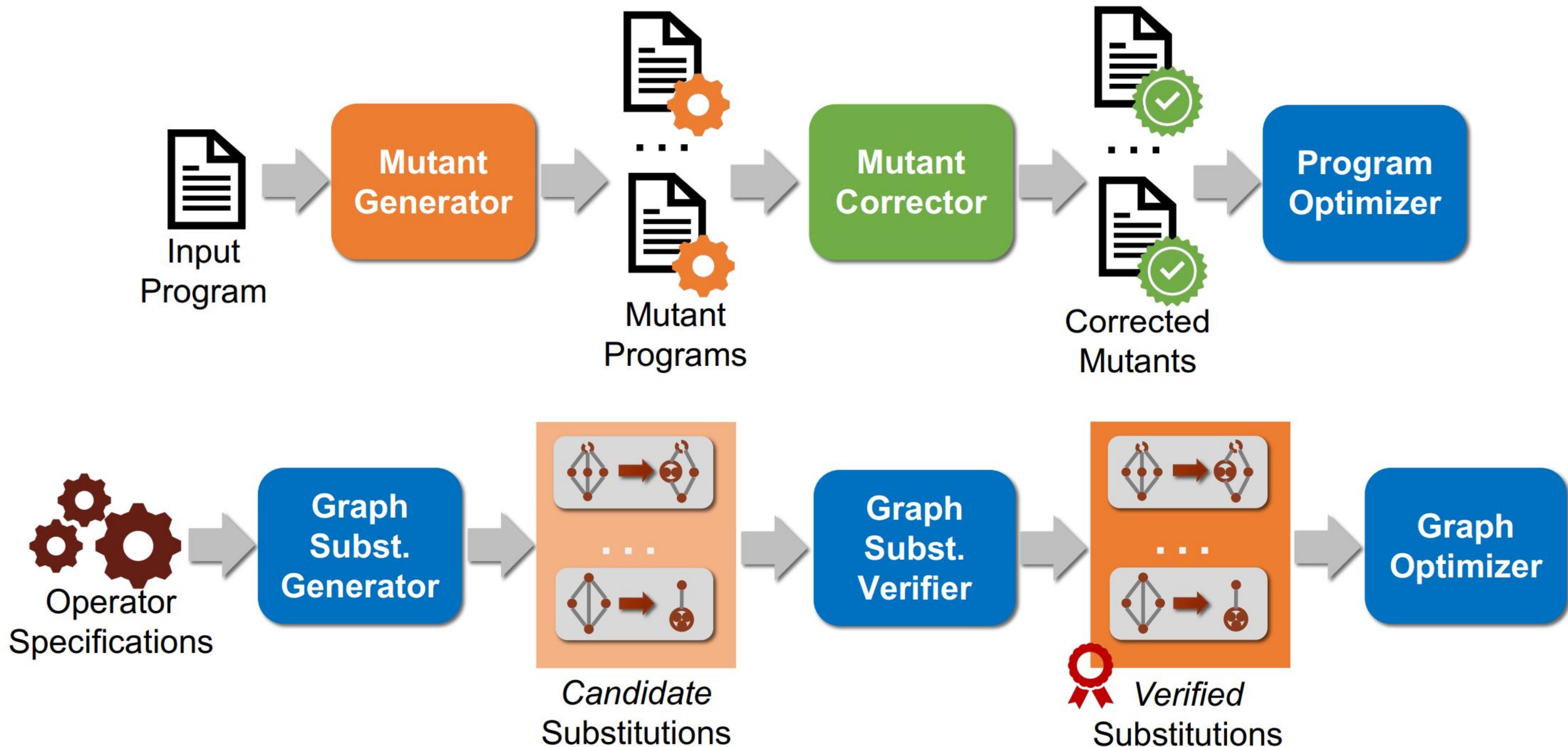


Correct the Mutant

- Goal: quickly and efficiently correcting the outputs of a mutant program
- Step 1: recompute the incorrect outputs using the original program
- Step 2: opportunistically fuse correction kernels with other operators



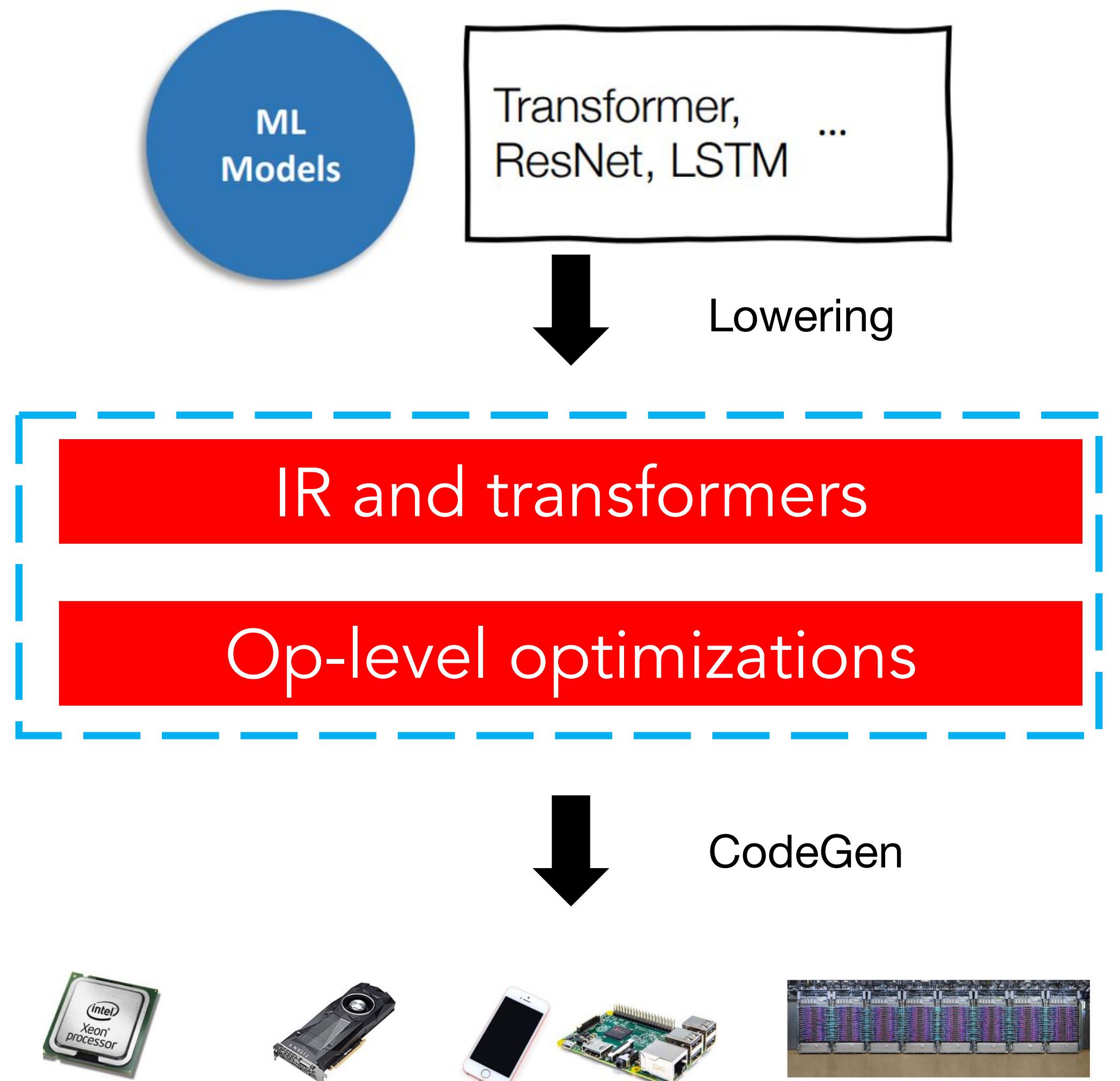
Recap



Summary & Questions to discuss

- Fully equivalent transformations vs. Partial
 - How to define search space
 - How to prune search space
 - How to verify & correct
 - How to apply to the ML graph optimization

Compilation Process

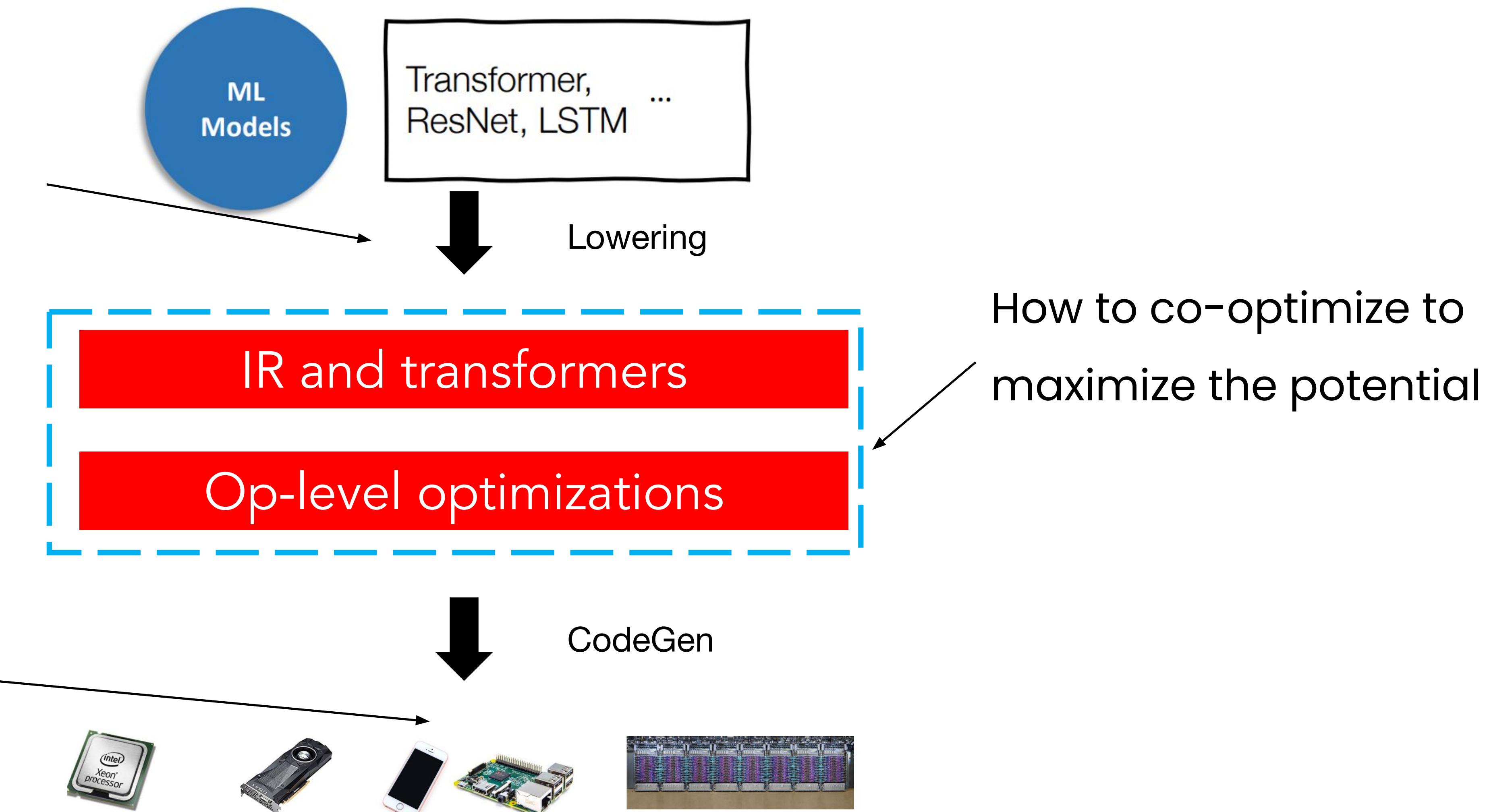


Ideally, they
should be
co-optimized?
Guest lectures

Topics will be covered later by Guests

How to lower user
program to IRs?

- The program is in PY
- Control flows
- Dynamism



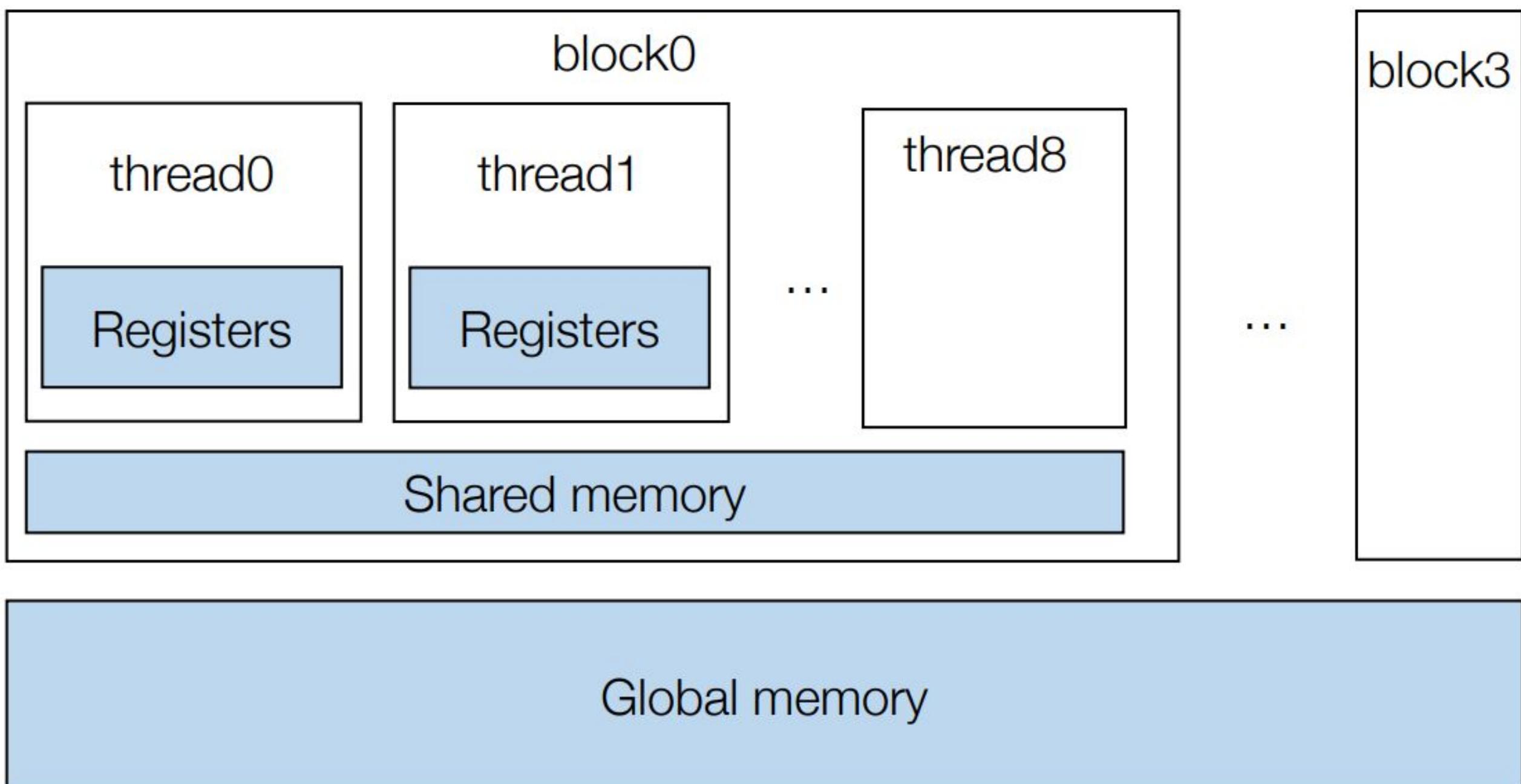
Agenda on this part

- ML Compilation Overview
 - Operator compilation
 - Graph optimization
- **Memory Optimization**
 - Activation checkpointing
 - Quantization and Mixed precision
- Two Guest Talks covering details in compilation, JIT, graph fusion, and beyond:
 - Meta PyTorch lead developer: Jason Ansel

Memory Optimization

- Checkpointing and rematerialization
- CPU Swapping
- Quantization and Mixed precision

Recap: Memory Hierarchy



Shared memory: 64 KB per core

GPU memory(Global memory):

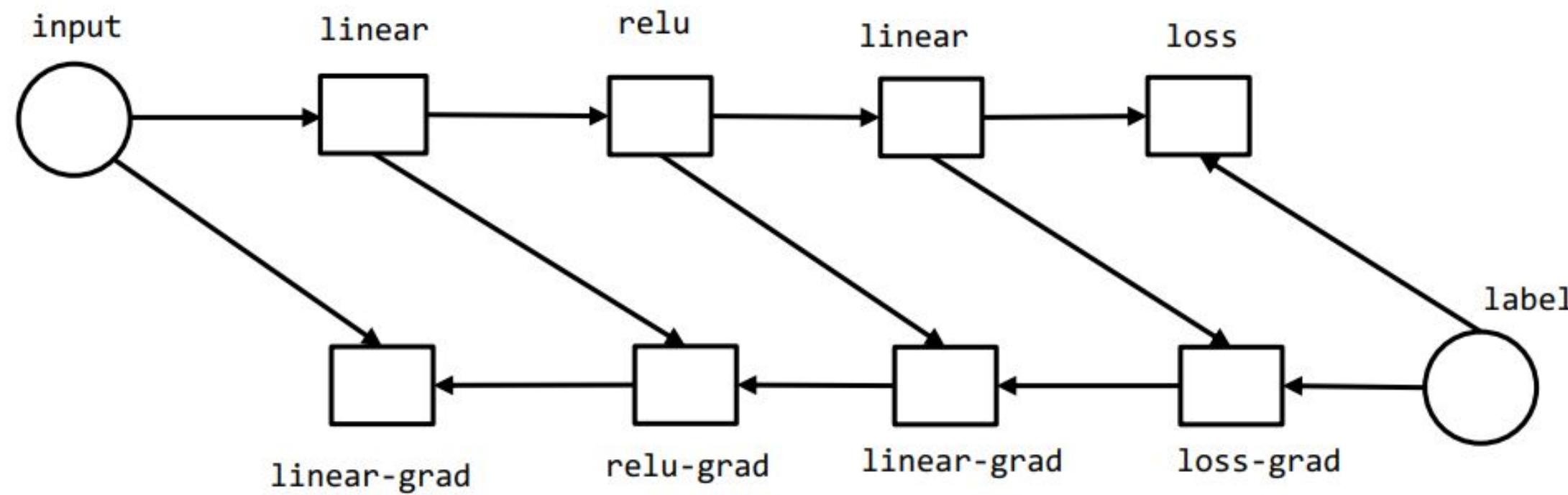
RTX3080 10GB

RTX3090 24GB

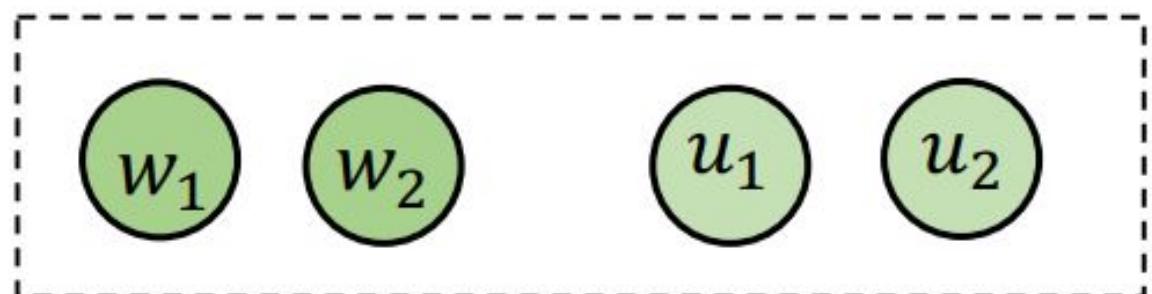
A100 40/80 GB

Source of Memory Consumption

A simplified view of a typical computational graph for training,
weights are omitted and implied in the grad steps.



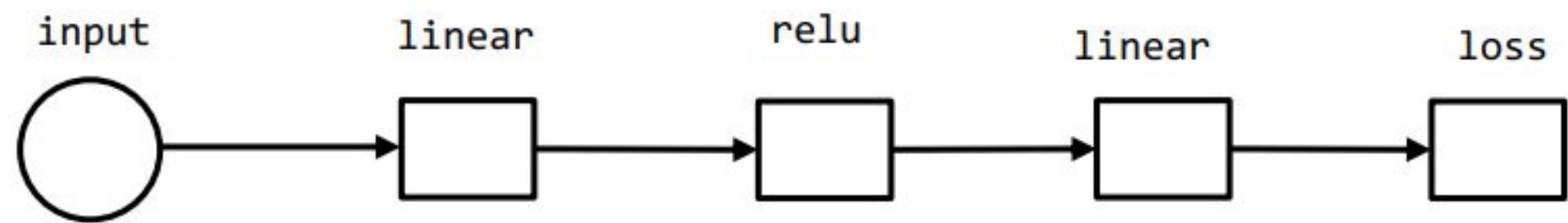
Optimizer states



Sources of memory consumption

- Model weights
- Optimizer states
- Intermediate activation values

At Inference



We only need $O(1)$ memory for computing the final output of a N layer deep network by cycling through two buffers