

```
54 def ave_number_of_moves(self, tries=10000):
```

1a. Correct. 2

2

```
75 class SingleDieTerrainRace(TerrainRace):
76
77     def _move(self):
78         move_length = self._random_move_parameters()
79         move_strength = move_length
80         if self.debug:
81             self._show_position(move_length, move_strength)
82         while (self._pos < len(self.terrain) and move_length and move_strength >= self.
            ↳ terrain[self._pos]):
83             move_length -= 1
84             self._pos += 1
85
86     def _random_move_parameters(self):
87         return randint(1, 6)
```

1b. This works, but is not ideal, since you were asked to only change what needs to be changed, but this has a lot of duplication from the original class. 2

2

```
61 class HardTerrainRace(TerrainRace):
62     def _move(self):
63         move_length, move_strength = self._random_move_parameters()
64         if self.debug:
65             self._show_position(move_length, move_strength)
66         while (self._pos < len(self.terrain)
67             and move_length
68             and move_strength >= self.terrain[self._pos]):
69             move_length -= 1
70             self._pos += 1
71             if move_length == 0:
72                 self._pos -= 1
```

1c. Ideally this should use `super` so you don't have to copy so much code from the original code. But also this doesn't work. See this example:

```
>>> t = HardTerrainRace([2,2,2,1,4], True)
>>> t.number_of_moves()
*2 2 2 1 4    3 6
 2 2*2 1 4    2 6
 2 2 2*1 4    5 1
 2 2 2 1*4    1 1
 2 2 2 1*4    5 3
 2 2 2 1*4    4 2
 2 2 2 1*4    4 2
 2 2 2 1*4    6 6
8
```

It backs up one step after the first two moves, but not after the later ones. 2

2

```

94  """
95
96  A.  $O(nm^2)$  - the first for loop is the array x so that's going to be variable n,
97  the next for loop is array y so that's going to be m and inside that one has another
98  for loop traversing through the same array y, this will be another m. the inner-most
99  for loop is going to have a constant of 1/2 yielding a more accurate big O as  $O(n*(m^2/2))$ ,
100  but that 1/2 constant doesn't matter, so we're left with  $n(m^2)$ 
101
102  B.  $O(nm)$  - the top for loop is going over array x (variable n). the outer for loop below is
103  going over the same, n, while the inner for loop is going over array y (variable m). the more
104  accurate big O, I suppose is  $O(n+nm)$  / factored out as  $O(n(1+m))$ , but the nm is a higher
105       $\hookrightarrow$  order so we don't care about
106  the lone n.
107
108  C.  $O(m\log n)$  - the while loop is traversing over array x (variable n) and the for loop is
109       $\hookrightarrow$  traversing
110  array y (variable m). But the while loop has  $n = n//2$  which will give it a logarithmic
111       $\hookrightarrow$  complexity,
112  making it  $\log n$ . so in total we have  $m\log n$ 
113
114  D.  $O(m^3)$  - this one has all 3 for loops traversing through the same array, array y (as
115       $\hookrightarrow$  variable m).
116  this yields  $m^3$ . the statements in the inner-most for loop with reassigning the array doesn't
117       $\hookrightarrow$  add to the
118  complexity.
119
120  """

```

**2a** 2.5  
**2b** 2.5  
**2c** 2.5  
**2d** 2.5

2.5  
 2.5  
 2.5  
 2.5

```
121 class Node:
122     def __init__(self, data):
123         self.next = None
124         self.data = data
125
126
127 class LinkedListStack:
128     def __init__(self):
129         self.head = None
130         self.count = 0
131
132     def isEmpty(self):
133         # will see if there's nothing in the first node/head
134         if self.head is None:
135             return True
136         else:
137             return False
138         """
139         Returns True if list is empty; False otherwise
140
141         :return: bool
142         """
143
144     def push(self, item):
145         # What I'm doing is under the assumption I got from
146         # the instructions which is that the newest node
147         # created will be sent to the end of the linked list.
148         # The wording threw me off a bit.
149
150         if self.head is None:
151             self.head = Node(item)
152         else:
153             new_node = Node(item)
154             new_node.next = self.head
155             self.head = new_node
156         self.count += 1
157         """
158         Pushes 'Node' item on top of stack
159
160         :parameter:item: 'Node' being pushed
161         :return: None
162         """
163
164     def pop(self):
165         # can call isEmpty method instead of rewriting same
166         # conditionals to check if the linked list is empty
167         if self.isEmpty():
168             return 'Error'
169         else:
170             delNode = self.head
171             self.head = self.head.next
```

```
172         delNode.next = None
173         self.count -= 1
174         return delNode.data
175     """
176     Pops 'Node' off the top of the stack;
177     throws error if stack is empty
178
179     :return: Node
180     """
181
182     def peek(self):
183         # do same isEmpty check first as before
184         if self.isEmpty():
185             return None
186         # just return the top node's data as it'll be the
187         # self.head here
188         else:
189             return self.head.data
190     """
191     Returns Node on top of stack without removing it;
192     Returns None if stack is empty
193
194     :return: Node
195     """
196
197     def size(self):
198         return self.count
199     """
200     Returns the number of nodes that comprise the stack
201
202     :return: int
203     """
```

3 Excellent. 10

10

```

209  """
210
211  Write your answer here.
212
213  I would use 1 outer dictionary as my main access point to the rest of the data. I believe
    ↳ this will work with the caveat
214  that the telecommunications company has either a finite number of rankings, or if they come
    ↳ up with a new ranking then
215  update the dictionary with it, as you'll see later on.
216  The length of the dictionary, meaning the number of keys minus 1, would be predetermined by
    ↳ the number of rankings the
217  telecommunications company has. So if they have rankings 1,2,3,4, and 5 then the the keys of
    ↳ the dictionary would be the
218  same, (1,2,3,4,5). Next the values of these keys would be structured with singly linked lists
    ↳ , with the self.head in
219  each singly linked list constantly being updated wth the customer next in line. The linked
    ↳ list will be filled out in
220  order of first come first serve FOR SPECIFICALLY THAT RANKING, not just for ALL customers
    ↳ with varying rankings. And
221  just as the example presented in the instructions portrayed, if there are two customers in
    ↳ rank 3 (which is the highest
222  ranking defined by the telecommunications company), one customer with rank 2, and one
    ↳ customer with rank 1 then the
223  operators will be give both customers in rank 3 priority over the customers with rank 2 and
    ↳ 1, and the customers with
224  rank 3 will be dealt with in order of which called in first, as customers are added to the
    ↳ tails of linked lists as they
225  come in.
226  Operations of storing data and printing are as follows:
227  customers with ranking are fed into the dictionary and their ranking is matched against the
    ↳ dictionary's keys, then the
228  customer is added to the tail-end of the linked list corresponding to that specific key. Now,
    ↳ once a customer is taken
229  care of, in any linked list, they are removed as the self.head of that linked list and the
    ↳ self.head.next is assigned as
230  the new self.head, continuously updating as callers are dealt with or if they drop from the
    ↳ call. All customers in the
231  highest priority ranking (in this case rank 3) must be dealt with before allocating
    ↳ representatives to lesser rankings.
232  This can be accomplished with a simple isEmpty() method, commonly made in linked list classes
    ↳ (or even from the previous
233  question in this exam), for the highest ranking key's value's linked list.
234  For time complexities:
235  Insertion and deletion of nodes in the singly linked list both have average and worst case
    ↳ time complexities of  $O(1)$ 
236  which is nice for this example. The dictionary's time complexity on average is  $O(1)$  as well
    ↳ as it functions as a hashmap
237  in python, requiring the inner parameters to be has functions. This is advantageous as other
    ↳ times collisions can occur
238  and it'd make the worst time complexity be  $O(n)$ . Next, the operation of the previously
    ↳ mentioned isEmpty() method to

```

239 check if the higher ranking's linked list has customer waiting is a constant. All in all, I  
240     ↪ think this is a  
241 semi-efficient solution with a minimized time complexity.

242 Another thing I thought of which might be cool, but not mentioned by the question, is if the  
243     ↪ self.head customer is  
244 talking to a call representative and their call is deemed a different ranking as the customer  
245     ↪ and representative talk.  
246 In such a case, the customer put back into the stream of data inputted to the dictionary and  
247     ↪ put at the tail end of the  
248 newly designated, appropriate ranking.

249 Solution 2:  
250 A doubly linked list can be used instead of a singly linked list as the pointer to the  
251     ↪ previous node can be utilized  
252 effectively in this example. Also, doubly linked lists can be used to implement binary trees  
253     ↪ unlike singly linked lists.  
254 I feel as thought this would be advantageous for something like a call center, or a  
255     ↪ telecommunications company in our  
256 case, to account for further features than proposed than in this question, such as calls  
257     ↪ dropped or elevation of  
258 ranking. The customers are inputted with their ranking in the conventional .data parameter. A  
259     ↪ isEmpty() will be used  
260 again here for the same purpose. This is constant time so we don't mind that in the grand  
261     ↪ scheme of things. Now, once  
262 the doubly linked list is not empty, the next newly created Node will be appended and it's .  
263     ↪ data parameter (its  
264 ranking) will be compared against the previous node's data and if it's greater (higher  
265     ↪ priority as defined by the  
266 company) then its .prev parameter will point to the .prev.prev's node and in turn the .prev.  
267     ↪ prev's .next parameter will  
268 point to the the node at hand. And the node at hand's .next will point to the former .prev (  
269     ↪ the node that began the  
270 comparison. This comparison will happen until the node at hand meets attains self.head status  
271     ↪ , or reaches a node with  
272 the same ranking priority as its data parameter. Additionally, this also extends to if the  
273     ↪ node at hand has a lower  
274 rank priority it will stay at the tail end of the doubly linked list, pointing to None, until  
275     ↪ more new Nodes are  
276 created for further comparison, etc. And similarly to solution 1, the customer/node at the  
277     ↪ self.head position is deleted  
278 once dealt with and prompts the self.head.next Node to become the first node/self.head .  
279 For time complexities:  
280 Initially inserting the customers with their rankings into the doubly linked list will have  
281     ↪ at worst  $O(1)$  (this is the  
282 same at best, too). Deleting the customer after they're dealt with also has a worst time  
283     ↪ complexity of  $O(1)$ . However,  
284 searching through the doubly linked list for a node to compare its ranking data will come  
285     ↪ with it a time complexity of  
286  $O(n)$  as it must traverse through the nodes one at a time (or 'n' at a time).

```
270 The same idea of reassigning customer ranking applies to this solution, but instead of the
    ↳ data reentering the input
271 stream seen in solution 1, it will be either elevated or demoted to a position closer or
    ↳ further away from the self.head
272 iff (if and only if) it's ranking is reassigned by the representative. This would cause  $O(n)$ 
    ↳ in the worst case.
273
274 ""
```

4 Great discussion overall; good join. A small comment: I think what you're describing in the first method, in so many words, is a queue no? **10**

**10**

```

281 def read_and_sort():
282     from spacy.lang.en import stop_words, English
283
284     nlp = English()
285
286     # trout file
287     with open('trout.txt') as trout_file:
288         raw_text = trout_file.read().lower()
289     # add stopwords to nlp's defined stop word vocabulary
290     with open('stopwords.txt') as stopwords_file:
291         lines = stopwords_file.read().split()
292         # print(lines)
293         for word in lines:
294             term = nlp.vocab[word]
295             term.is_stop = True
296
297     tokenizer = nlp.tokenizer
298     doc = nlp(raw_text)
299     tokenized_words = [tok.text for tok in doc if tok.is_stop == False if tok.is_alpha]
300     # tokenized_words
301
302     from collections import defaultdict
303     r = defaultdict(int)
304     for token in tokenized_words:
305         r[token] += 1
306
307     values = r.values()
308     # Return values of a dictionary
309     total = sum(values)
310     print("Token\tfreq\tprob")
311
312     with open('trout.txt.probs', 'w') as file:
313         file.write("Token\tfreq\tprob\n")
314         for key, value in r.items():
315             prob = value / total
316             file.write(f"{key}\t{value}\t{prob}\n")

```

5 Code looks great. Two minor things you didn't do, however: 1. sort the output by descending order of probability and 2. add the input file as an argument to the function. 6.5

6.5



```

322 #####
323
324 import numpy as np
325 import pandas as pd
326 import matplotlib.pyplot as plt
327 import seaborn as sns
328
329 A = np.array([[22, 28, 52],
330              [87, 12, 76],
331              [44, 61, 81],
332              [97, 4, 67],
333              [52, 14, 24],
334              [34, 82, 7]])
335
336 B = np.array([[88, 41, 22, 1, 14, 70, 7, 48, 64],
337              [60, 10, 17, 34, 5, 57, 16, 98, 36]])
338
339 # Question 6.1
340
341 # Write your code here.
342 arr = np.max(np.power(A,2))
343 print(arr)
344 # Question 6.2
345
346 # Write your code here.
347 summation = np.sum(A.argmax(axis=0))
348 print(summation)
349 # Question 6.3
350
351 # Write your code here.
352 B = np.reshape(B, (6,3))
353 # print(B)
354 #np.concatenate(np.reshape(A,B),axis=1)
355 X = np.concatenate((A,B), axis=1)
356 print(np.mean(X, axis = 0))
357 # Question 6.4
358
359 # Write your code here.
360 C = np.concatenate((A,B))
361 Y = np.random.randint(1,99, size=(6,3))
362 C = np.concatenate((C,Y))
363 #[random.uniform(low,high) for _ in range(size)]
364 #np.random.Generator.uniform(low=0.0, high=1.0, size=None)
365 # Question 6.5
366
367 # Write your code here.
368 import pandas as pd
369 df = pd.DataFrame(data=C, columns = ['X', 'Y', 'Z'])
370 print(df)
371 import seaborn as sns
372 import matplotlib.pyplot as plt

```

```
373 sns.scatterplot(data=df, x='X', y='Y', size = 'Z')
374 plt.title("A Very Informative Plot")
375 plt.show()
```

6.1 2

6.2 2

6.3 2

6.4 2

6.5 2

2

2

2

2

2

2

Total: 52.5 / 60 = 88 %, VG