

Yesterday: how to calc. a feas. point if the starting point is not feasible.

Degeneracy: when there exists one or more zeros in b (rhs).

Some notation

Dictionary:

$$g = \sum_{j=1}^n c_j x_j$$

$$x_{n+i} = b_i - \sum a_{ij} x_j, \quad 1 \leq i \leq m$$

$$x_j \geq 0, \quad 1 \leq j \leq n+m$$

Notation.

ω is a more convenient

$$\left(\begin{array}{c} x_1 \\ \vdots \\ x_n \\ \hline w_1 \\ \vdots \\ w_m \end{array} \right) \rightarrow \left(\begin{array}{c} x_1 \\ \vdots \\ x_n \\ \hline x_{n+1} \\ \vdots \\ x_{n+m} \end{array} \right)$$

\brace{NBVs}
 \brace{BVs}

- Set of (indices of) NBVs: N
- —————, ————— BVs: B
- Initially, $N = \{1, \dots, n\}$
 $B = \{n+1, \dots, n+m\}$
- Dictionary after zero or more steps:

$$y = \bar{y} + \sum_{j \in N} \bar{c}_j x_j$$

"bar": coefficients
in this instance.
fixed data.

$$x_i = \bar{b}_i - \sum \bar{a}_{ij} x_j, i \in B$$

$$x_j \geq 0, j \in B \cup N$$

union

Degeneracy (read ch. 3)

$$\max c^T x \text{ s.t. } Ax \leq b, x \geq 0$$

Up to now: $b > 0$. Let's now look at a situation where at least one of $\{b_1, \dots, b_m\}$ is zero.

Def. A dict. is called degenerate

if $b_i = 0$ for some i from 1 to m .

Example:

$$y = 5 + x_3 - x_1$$

$$x_2 = 5 + 2x_3 - 3x_1$$

$$x_4 = 7 - 4x_1$$

$$x_5 = 0 + x_1$$

degeneracy

*
(unbounded problem)

solution is not optimal since we have a positive coefficient for x_3 , but there are no x_3 with a neg. coefficient in the constraints, so there is no solution (x_3 can be augmented (increased) as much as we want. *)

Remember, unbounded means there is no solution.

Example 2 :

$$\begin{array}{l} \underline{f = 5 + x_3 - x_1} \\ x_2 = 5 - 2x_3 - 3x_1 \\ x_4 = 7 - 4x_1 \\ x_5 = 0 + x_1 \end{array}$$

We interchange x_3 and x_2 .

$$x_3 = \frac{5}{2} - \frac{1}{2}x_2 - \frac{3}{2}x_1$$

New dict. :

$$\varphi = 5 + \left(\frac{5}{2} - \frac{1}{2}x_2 - \frac{3}{2}x_1 \right) - x_1$$

$$= \left(\frac{15}{2} \right) \dots$$

φ is improved!

Now we've seen two examples where the degeneracy (zero) has had no effect.

Example 3 :

note the degeneracy

$$\max \frac{3}{4}x_1 - 20x_2 + \frac{1}{2}x_3 - 6x_4$$

$$\text{s.t. } \frac{1}{4}x_1 - 8x_2 - x_3 + 9x_4 \leq 0$$

$$\frac{1}{2}x_1 - 12x_2 - \frac{1}{2}x_3 + 3x_4 \leq 0$$

$$x_3 < 1$$

$$x_i \geq 0, i \in \{1, 2, 3, 4\}$$

Exercise: try to make a pivot step
in pplex and see what happens.

BV	φ
w, w_2, w_3	0
x_1, w_2, w_3	0
x_1, x_2, w_3	0
x_3, x_2, w_3	0
x_3, x_4, w_3	0
w_1, x_1, w_3	0
w_1, w_2, w_3	0
	:

We will never leave! Cycling!

We're always in the same point.

How should the software handle this?

- We remain in the same point
- φ stays the same
- Sometimes different vertices with the same value for φ are presented

Example 4 (illustration of degeneracy):

$$\begin{array}{l} \downarrow \\ \underbrace{y = 5 + x_3 - x_1} \\ x_2 = 5 + 2x_3 - 3x_1 \\ x_4 = 7 - 4x_1 \\ \rightarrow x_5 = \boxed{0} - x_3 + x_1 \end{array}$$

Problem arises if the pivot row has index i_0 and $\bar{b}_0 = 0^*$.

We interchange x_3 and x_5

$$x_3 = -x_5 + x_1$$

New dict. :

$$\begin{aligned} g &= 5 - x_5 + \cancel{x_1} - \cancel{x_1} \\ &= 5 - x_5 \end{aligned}$$

Notice that after the pivot step
we got the same obj. value.

If there is a zero in the pivot
row, the value of the obj. func.
is not improving!

- Degeneracy (or degeneration) is common (that is, $b_{i_0} = 0$ for some $i \in \{1 \dots m\}$)
- Cycling is rare. Note that degeneracy does not imply cycling
- Cycling can be avoided at some cost (we explain two possibilities.)

Theorem :

Either the simplex method terminates or there is a cycle.

Proof. Assume the contrary.

Assume it does not terminate and that it doesn't cycle. Since a dict. is uniquely determined by a partition of $n+m$ variables into m basis variables and n non-basis variables.

There are exactly $\binom{n+m}{m}$ such partitions.

upper bound for
dictionaries and
vertices!

Since SM doesn't terminate, one of the already calculated dictionaries will be obtained after at most $\binom{n+m}{m}$ steps.

Furthermore, since the value of the obj. func. $\mathcal{G} = \mathbf{c}^T \mathbf{x}$ is not decreasing at the dictionaries (or the vertices) the method is obviously cycling. **Contradiction Δ .**

$$\mathcal{G}^1 \leq \mathcal{G}^2 \leq \mathcal{G}^3 \leq \dots \mathcal{G}^{(\frac{n+m}{m})} \cdot \mathcal{G}^1$$

after at most $(\frac{n+m}{m})$ steps

↑
dictionary

$$\Rightarrow \mathcal{G}^1 = \mathcal{G}^2 = \dots \mathcal{G}^{(\frac{n+m}{m})}$$

This implies
cycling.

How to prevent cycling?

Technique I : Lexicographic SM

Introduction of m positive perturbations $\epsilon_i > 0$, $i = 1 \dots m$ of the right hand sides (b) with:

$0 < \epsilon_m < \epsilon_{m-1} < \dots < \epsilon_1 <$ all other data of the LP.

"Each ϵ_i acts on an entirely different scale from all other ϵ_j and the data of the problem. What we mean by this is that no lin. comb. of the ϵ_i w/ coeff. that may arise in the course of the SM can ever produce

a number whose size is of the same order as the data in the problem.

Similarly any ϵ_i can never escalate to a level of ϵ_j ($i \neq j$).

Trick: we do not assign any values to ϵ_i (ϵ_i are considered parameters).

Example :



$$\underline{g = 6x_1 + 4x_2}$$

$$w_1 = 0 + 9x_1 + 4x_2$$

$$\rightarrow w_2 = \textcircled{0} - 4x_1 - 2x_2$$

$$w_3 = 1 - x_2$$

Now, introduce perturbations to the zeros.

$0 < \epsilon_3 < \epsilon_2 < \epsilon_1 <$ all other data.

$$\begin{array}{l}
 \underbrace{\mathcal{L} =}_{w_1 = 0 + \boxed{\epsilon_1}} \underbrace{6x_1 + 4x_2}_{+ 9x_1 + 4x_2} \\
 w_2 = 0 + \boxed{\epsilon_2} + 4x_1 - 2x_2 \\
 w_3 = 1 + \boxed{\epsilon_3} - x_2
 \end{array}$$

ϵ_i obtain columns like variables.

- Now, determine pivot column as usual: x_1
- Pivot row is w_2
- Interchange these as usual

$$x_1 = \frac{e_2}{q} - \frac{1}{q} w_2 - \frac{1}{z} x_2$$

(continues next lecture)