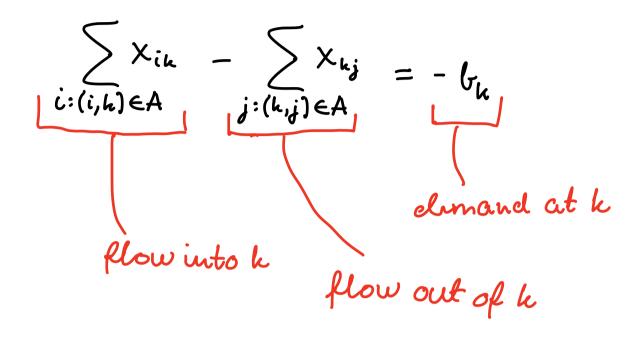


Network with supply/dimarel. Neg. numbers = dimarel.

Objective: move supply over to clemand modes.

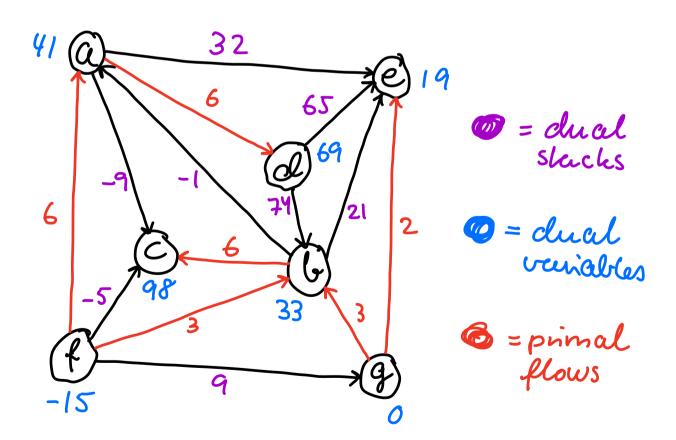


Net flow must be equal to the demand at k.

Obj. fune.: min. $\sum CijXij$ (i,j) $\in A$ min. cost of transportation

Initially, one vow in the hoch - arc incidence matrix is redundant (meaning the sum of all rows in A is -1.

A does not have rank m at this point, so we can delete one row.



7 noces and 6 arcs in the Spanning tree.

At least two nodes must have only one are connection.

Note that we can start with means there is a any node c=-6 c=-6 c=-6 c=-6 c=-6 c=-2 c=-2

c, d'and e are the leaf nocles in this case.

Step 1

stant at one of the leaves to find a tree solution.

the Xrc arc.

(F)
$$\times_{fb} + \times_{fa} = 9$$
 — f has two outgoing arcs whis needs to have sum of 9 .

We have now covered all 6 arcs (note that we did not have to look at g). It is the redundant noch in this case.

Note that any point that is not a Ceaf can be a redundant noch.

Our initial solution looks as follows:

$$X_{bc} = 6$$

 $X_{ge} = 2$
 $X_{ad} = 6$
 $X_{ea} = 6$
 $X_{bc} = 3$
 $X_{gc} = 3$

(all others are zero)

Is it optimal?

We can check the corresponding dual solution to determine this. (more on this later)

Coeff. matrix after debiting G: (notice also the permutation)

Carcs of our spannint tr.

Xad Xfa Xfb Xbc Xgb Xge

Cl 1

Q -1 1

f -1 -1

C b e 1 -1 1

The redundant noch (g) is called the root noch.

The vernoining (m-1, m-)-matrix is non-singular.

Now we look at (D) to determine if we are optimal.

olf (i,j) belongs to the spanning tree then:

· Xij is BV for (P). (red arcs)

o Zij is NBV for (D) => Zij = 0

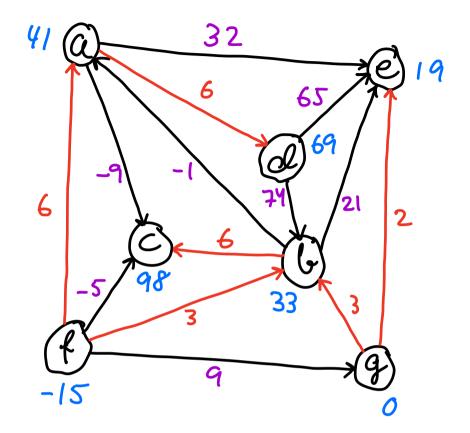
yj-yi = Cij for m nodes and M-1 arcs of the spanning tree.

Calculate (D).

We wooh at the constraints of the dual problem.

Start with the voot woch (in our case @).

Let y; = 0. — This is our g.



NBV for (D), (hence zero.

arc (q, e): ye-yq+zeq=Cqe ye-yq=19-0=19 $\Rightarrow ye=19$

arc (g, b): $y_{r} - y_{g} = 33 - 0 = 33$ => $y_{r} = 33$

$$avc(a,d): yd - ya = 2f$$

$$= yd = 69$$

(yi): part of the dual solution

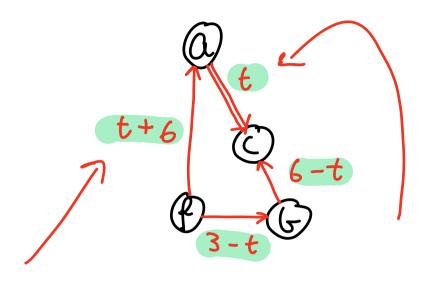
Stack variables:

curent spanning tree is not optimal since (D) is not feasible!!

Pivot step in the network SM Choose an infectible dual slach variable representing an NBV are which will become our new BV.

New BV: we choose Zac = -9 (the "most" infeasible NBV arc.

New NBV: assign a positive value t > 0 to the new BV (a,c)=> bralance countr. along the cycle (a,c)-(c,b)-(f,b)-(f,a) change accordingly.



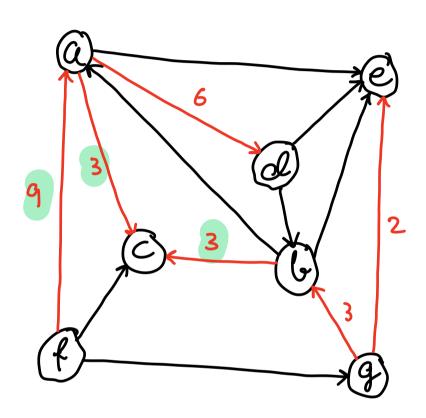
A new t goes out of @ so a new additional t must go in.

How big can t be without breaking the balance constraints?

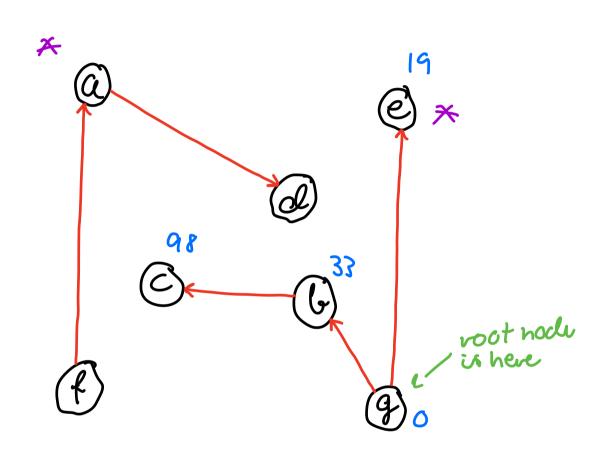
Answer: 3

 $\times \mu = 0$ for $t = 3 \Rightarrow (\mu, \mu)$ (ecues) the trusis.

New solution:



Optimal? check for fear of (D).



Delited the leaving arc (f.b).

Do not include the entering arc
(a,c). We obtain two disjoint sets.*

How to update the dual variables?

calc. starts with the root node

=> the dual solutions in the part
containing the root noch will not
change (c, 6, 9, e).

$$(f, a): ya - yf = 56$$

 $\Rightarrow yf = -6$

Previous dual solution:

$$y_{\alpha} = 41$$

 $y_{\beta} = -15$
 $y_{\beta} = 69$
+ 9 gives us the new values.

Rule:

we choose the infeasible dual slach variable Zac = -9 => (a,c) became new BV => New value of Zac after this iteration becomes zero.

After iteration: 0 = ya - yı + cac

Before iteration: -9 = ya - yı + cac

ya = ya + 9

belongs to

root nocle.

Analogously for all nocles not belonging to the root noch part (a, f, d).

Stach veriables: (for the NBV arcs)

Zae-ya+ye=Cae
ya+q not changing

Zae = Zae + 9 41 32

Optimal? No, since (a,b) in the new sol. = -10 Do two more iterations to get the optimal sp. tree.